

## Natural numbers

In the beginning there were the *natural numbers*

$\mathbb{N} : 0, 1, \dots, n, n+1, \dots$

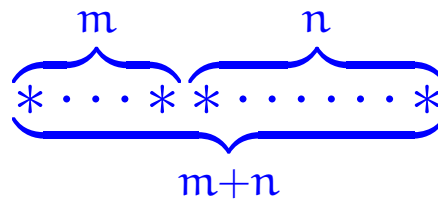
generated from *zero* by successive increment; that is, put in ML:

```
datatype
```

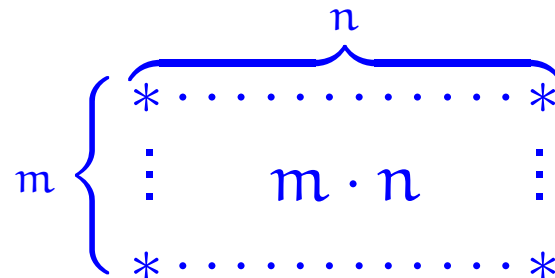
```
  N = zero | succ of N
```

The basic operations of this number system are:

► Addition



► Multiplication



The additive structure  $(\mathbb{N}, 0, +)$  of natural numbers with zero and addition satisfies the following:

► Monoid laws

$$0 + n = n = n + 0 \quad , \quad (l + m) + n = l + (m + n)$$

► Commutativity law

$$m + n = n + m$$

and as such is what in the mathematical jargon is referred to as a commutative monoid.

Also the *multiplicative structure*  $(\mathbb{N}, 1, \cdot)$  of natural numbers with one and multiplication is a commutative monoid:

► Monoid laws

$$1 \cdot n = n = n \cdot 1 \quad , \quad (l \cdot m) \cdot n = l \cdot (m \cdot n)$$

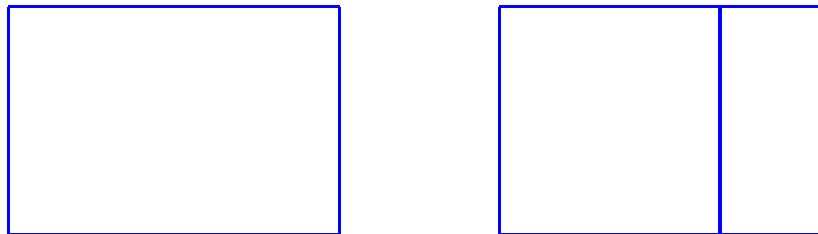
► Commutativity law

$$m \cdot n = n \cdot m$$

The additive and multiplicative structures interact nicely in that they satisfy the

► Distributive law

$$l \cdot (m + n) = l \cdot m + l \cdot n$$



and make the overall structure  $(\mathbb{N}, 0, +, 1, \cdot)$  into what in the mathematical jargon is referred to as a *commutative semiring*.

# Cancellation

The additive and multiplicative structures of natural numbers further satisfy the following laws.

▶ Additive cancellation

For all natural numbers  $k, m, n$ ,

$$k + m = k + n \implies m = n \quad .$$

▶ Multiplicative cancellation

For all natural numbers  $k, m, n$ ,

$$\text{if } k \neq 0 \text{ then } k \cdot m = k \cdot n \implies m = n \quad .$$

# Inverses

## Definition 41

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1. A number  $x$  is said to admit an additive inverse whenever there exists a number  $y$  such that  $x + y = 0$ .
2. A number  $x$  is said to admit a multiplicative inverse whenever there exists a number  $y$  such that  $x \cdot y = 1$ .



Extending the system of natural numbers to: (i) admit all additive inverses and then (ii) also admit all multiplicative inverses for non-zero numbers yields two very interesting results:

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(i) the integers

$$\mathbb{Z} : \dots -n, \dots, -1, 0, 1, \dots, n, \dots$$

which then form what in the mathematical jargon is referred to as a commutative ring, and

(ii) the rational  $\mathbb{Q}$  which then form what in the mathematical jargon is referred to as a field.

## The division theorem and algorithm

**Theorem 42 (Division Theorem)** *For every natural number  $m$  and positive natural number  $n$ , there exists a unique pair of integers  $q$  and  $r$  such that  $q \geq 0$ ,  $0 \leq r < n$ , and  $m = q \cdot n + r$ .*

## The division theorem and algorithm

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**Definition 43** *The natural numbers  $q$  and  $r$  associated to a given pair of a natural number  $m$  and a positive integer  $n$  determined by the Division Theorem are respectively denoted  $\text{quo}(m, n)$  and  $\text{rem}(m, n)$ .*

## The Division Algorithm in ML:

```
fun divalg( m , n )
  = let
    fun diviter( q , r )
      = if r < n then ( q , r )
        else diviter( q+1 , r-n )
    in
      diviter( 0 , m )
    end

fun quo( m , n ) = #1( divalg( m , n ) )

fun rem( m , n ) = #2( divalg( m , n ) )
```

**Theorem 44** *For every natural number  $m$  and positive natural number  $n$ , the evaluation of  $\text{divalg}(m, n)$  terminates, outputting a pair of natural numbers  $(q_0, r_0)$  such that  $r_0 < n$  and  $m = q_0 \cdot n + r_0$ .*

PROOF: