

# 5.2 Fibonacci Heaps

Frank Stajano

Thomas Sauerwald





Operation	Linked list	Binary heap	Binomial heap
MAKE-HEAP	O(1)	O(1)	<i>O</i> (1)
INSERT	O(1)	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$
Мінімим	$\mathcal{O}(n)$	O(1)	$\mathcal{O}(\log n)$
EXTRACT-MIN	$\mathcal{O}(n)$	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$
MERGE	$\mathcal{O}(n)$	$\mathcal{O}(n)$	$\mathcal{O}(\log n)$
DECREASE-KEY	O(1)	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$
DELETE	O(1)	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$



Operation	Linked list	Binary heap	Binomial heap	Fibon. heap
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INSERT	O(1)	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$	O(1)
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Merge	$\mathcal{O}(n)$	$\mathcal{O}(n)$	$\mathcal{O}(\log n)$	O(1)
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Operation	Binomial heap	Fibonacci heap
	actual cost	amortized cost
MAKE-HEAP	<i>O</i> (1)	<i>O</i> (1)
INSERT	$\mathcal{O}(\log n)$	<i>O</i> (1)
Мінімим	$\mathcal{O}(\log n)$	<i>O</i> (1)
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Binomial Heap: k/2 DECREASE-KEY

+ k/2 INSERT



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### Binomial Heap: k/2 DECREASE-KEY

- + k/2 INSERT
  - $c_1 = c_2 = \cdots = c_k = \mathcal{O}(\log n)$



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### Binomial Heap: k/2 DECREASE-KEY

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$$c_1 = c_2 = \cdots = c_k = \mathcal{O}(\log n)$$

$$\Rightarrow \sum_{i=1}^{k} c_i = \mathcal{O}(k \log n)$$



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Binomial Heap: k/2 DECREASE-KEY + k/2 INSERT

Fibonacci Heap: k/2DECREASE-KEY + k/2 INSERT

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Fibonacci Heap: k/2

DECREASE-KEY + k/2 INSERT

• 
$$\widetilde{c}_1 = \widetilde{c}_2 = \cdots = \widetilde{c}_k = \mathcal{O}(1)$$



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Binomial Heap: k/2 DECREASE-KEY

+ k/2 INSERT

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$$c_1 = c_2 = \cdots = c_k = \mathcal{O}(\log n)$$

$$\Rightarrow \sum_{i=1}^{k} c_i = \mathcal{O}(k \log n)$$

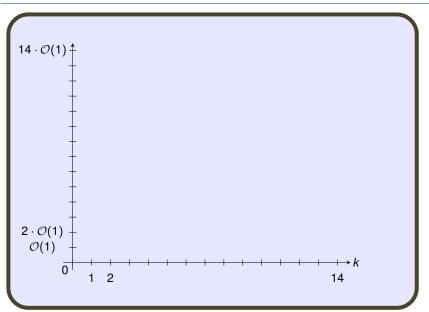
# Fibonacci Heap: k/2

DECREASE-KEY + k/2 INSERT

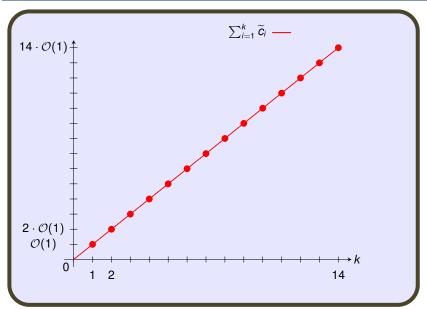
• 
$$\widetilde{c_1} = \widetilde{c_2} = \cdots = \widetilde{c_k} = \mathcal{O}(1)$$

$$\Rightarrow \sum_{i=1}^k c_i \leq \sum_{i=1}^k \widetilde{c}_i = \mathcal{O}(k)$$



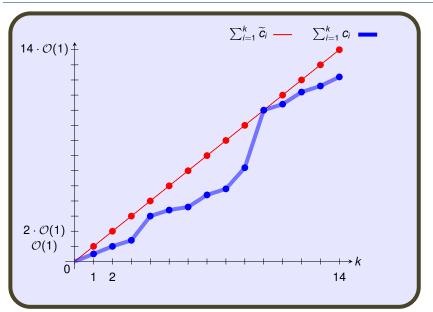




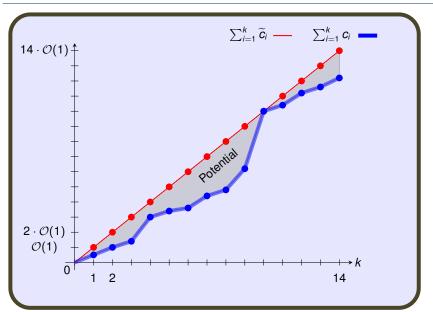




5.2: Fibonacci Heaps

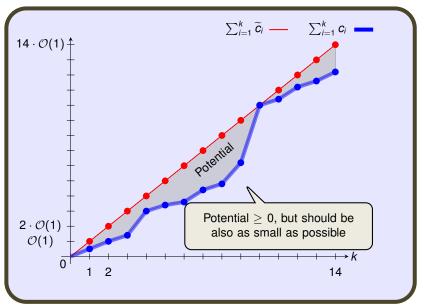








5.2: Fibonacci Heaps





### **Outline**

Structure

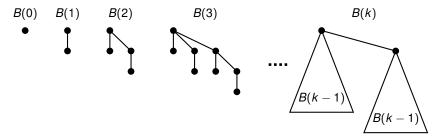
Operations

Glimpse at the Analysis



### **Reminder: Binomial Heaps**

#### **Binomial Trees**

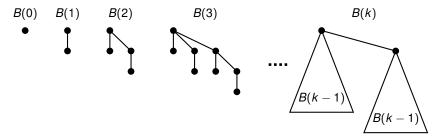


Binomial Heaps -

 Binomial Heap is a collection of binomial trees of different orders, each of which obeys the heap property

### **Reminder: Binomial Heaps**

#### **Binomial Trees**



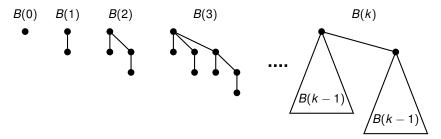
Binomial Heaps -

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- Operations:



### **Reminder: Binomial Heaps**

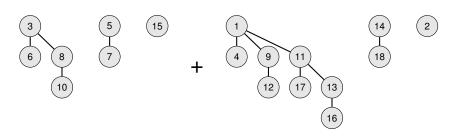
#### **Binomial Trees**



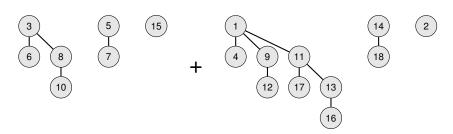
Binomial Heaps

- Binomial Heap is a collection of binomial trees of different orders, each of which obeys the heap property
- Operations:
  - MERGE: Merge two binomial heaps using Binary Addition Procedure
  - INSERT: Add B(0) and perform a MERGE
  - EXTRACT-MIN: Find tree with minimum key, cut it and perform a MERGE
  - DECREASE-KEY: The same as in a binary heap

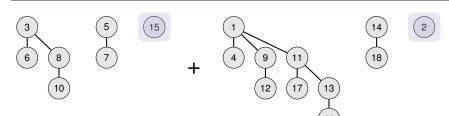




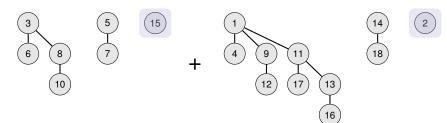






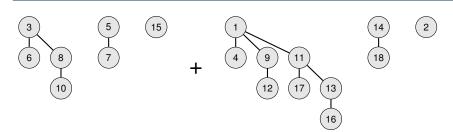






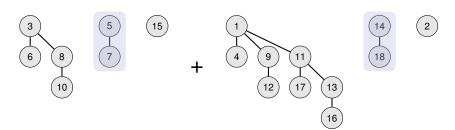






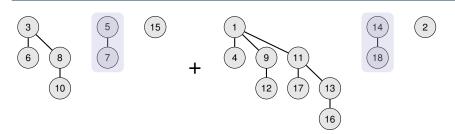


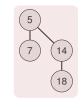






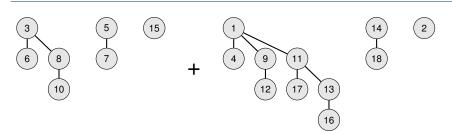








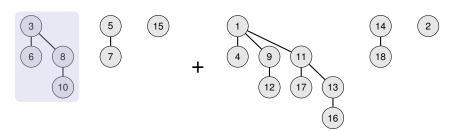


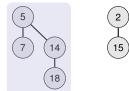




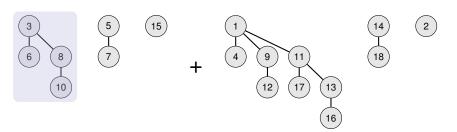


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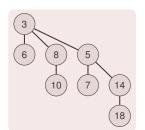




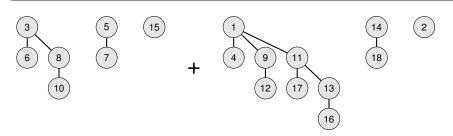




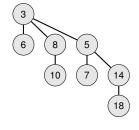




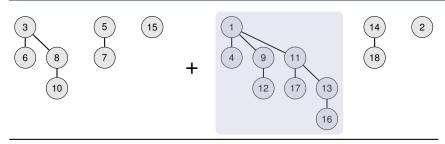




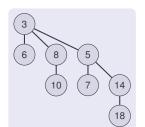




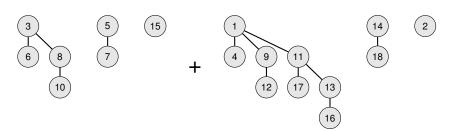


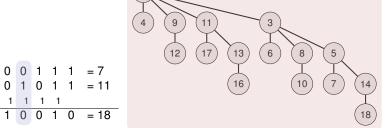






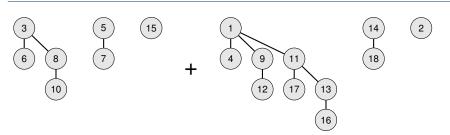


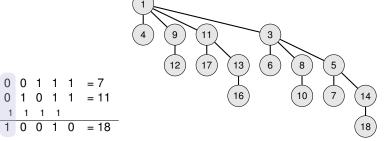




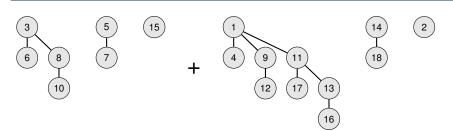


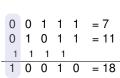
5.2: Fibonacci Heaps

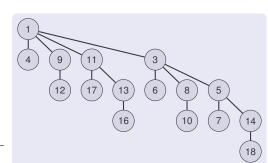






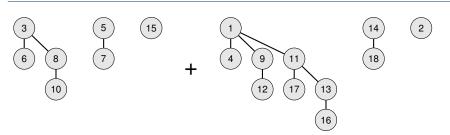


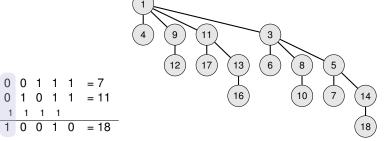






# **Merging two Binomial Heaps**





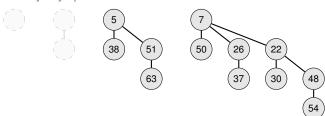


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### Binomial Heap vs. Fibonacci Heap: Structure

#### Binomial Heap:

- consists of binomial trees, and every order appears at most once
- immediately tidy up after INSERT or MERGE

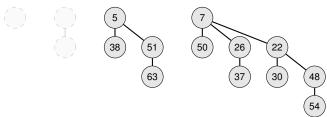




### Binomial Heap vs. Fibonacci Heap: Structure

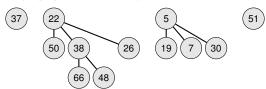
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#### Fibonacci Heap:

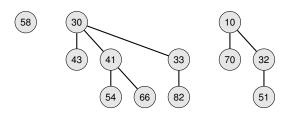
- forest of MIN-HEAPs
- lazily defer tidying up; do it on-the-fly when search for the MIN





Fibonacci Heap ————

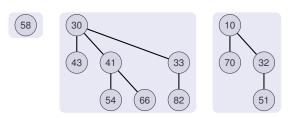
Forest of MIN-HEAPs





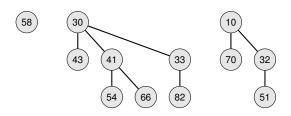
Fibonacci Heap ———

Forest of MIN-HEAPs



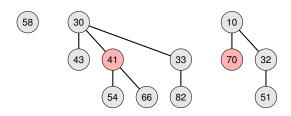


- Forest of MIN-HEAPs
- Nodes can be marked (roots are always unmarked)



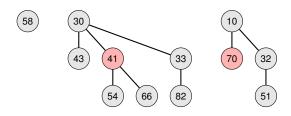


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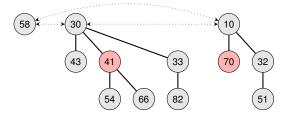


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- Nodes can be marked (roots are always unmarked)
- Tree roots are stored in a circular, doubly-linked list



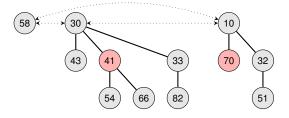


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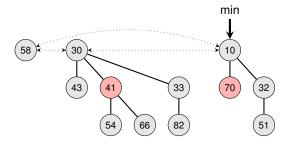


- Forest of MIN-HEAPs
- Nodes can be marked (roots are always unmarked)
- Tree roots are stored in a circular, doubly-linked list
- Min-Pointer pointing to the smallest element



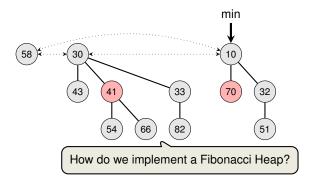


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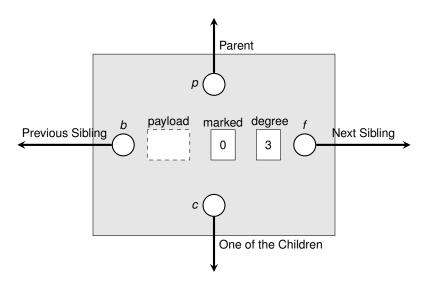




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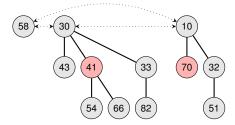






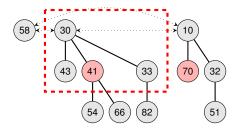
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# **Magnifying a Four-Node Portion**



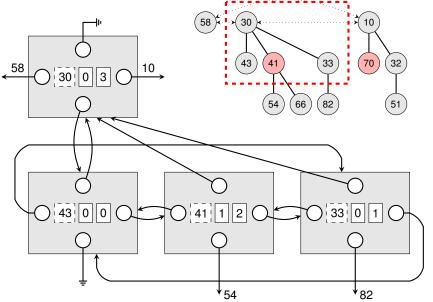


# **Magnifying a Four-Node Portion**





# **Magnifying a Four-Node Portion**





5.2: Fibonacci Heaps

#### **Outline**

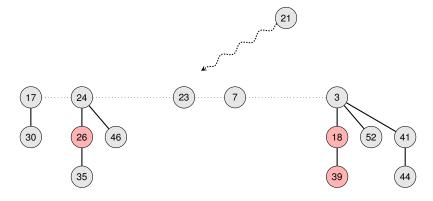
Structure

### Operations

Glimpse at the Analysis



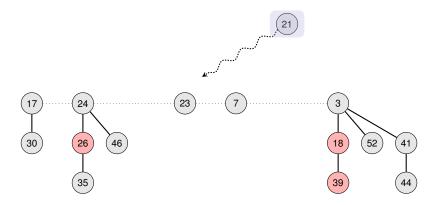






INSERT

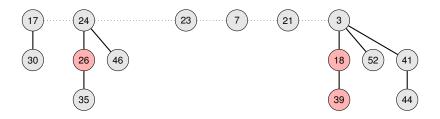
• Create a singleton tree





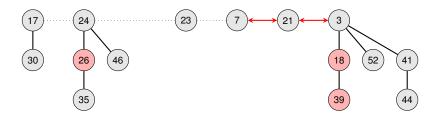
- Create a singleton tree
- Add to root list





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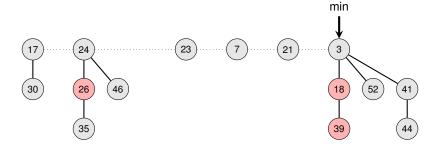






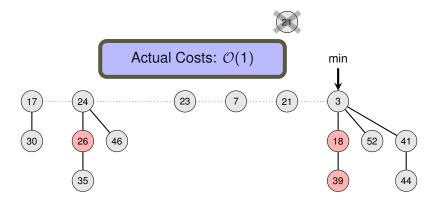
- Create a singleton tree
- Add to root list and update min-pointer (if necessary)





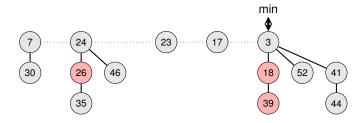


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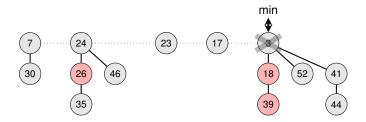






— EXTRACT-MIN ————

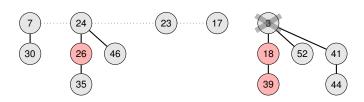
Delete min





— Extract-Min ———

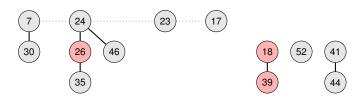
■ Delete min √





— EXTRACT-MIN ———

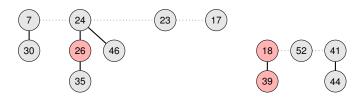
- Delete min √
- Meld childen into root list and unmark them





— EXTRACT-MIN ————

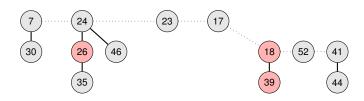
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- Meld childen into root list and unmark them





— EXTRACT-MIN ————

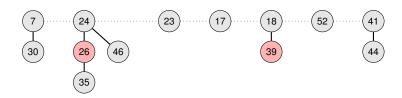
- Delete min √
- Meld childen into root list and unmark them





— EXTRACT-MIN ———

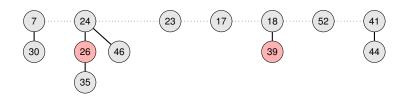
- Delete min √
- Meld childen into root list and unmark them





- EXTRACT-MIN -

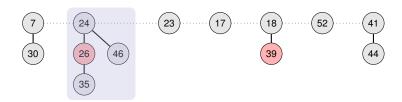
- Delete min √
- Meld childen into root list and unmark them
- Consolidate so that no roots have the same degree





- EXTRACT-MIN -

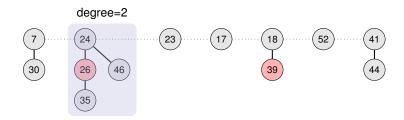
- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)





#### — EXTRACT-MIN –

- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)





- EXTRACT-MIN -

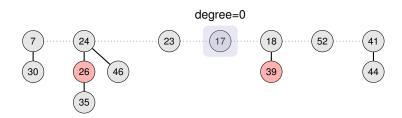
- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)





#### — EXTRACT-MIN ———

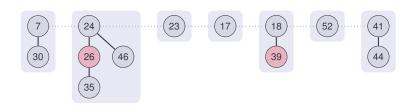
- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)





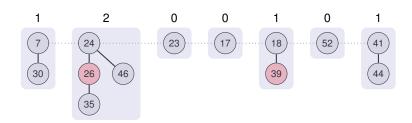
— EXTRACT-MIN ———

- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)





- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)

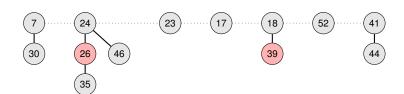




#### - EXTRACT-MIN -

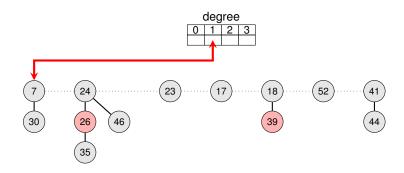
- Delete min √
- Meld childen into root list and unmark them
- Consolidate so that no roots have the same degree (# children)





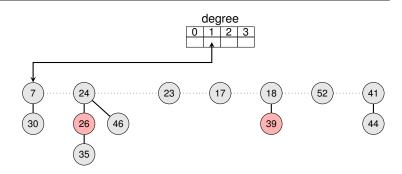


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



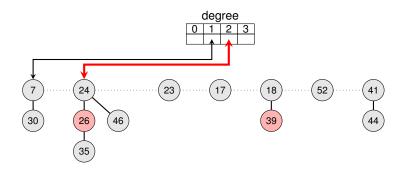


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



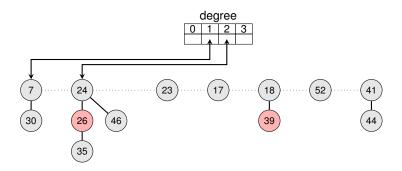


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



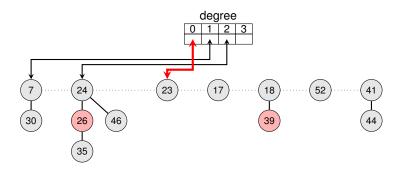


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



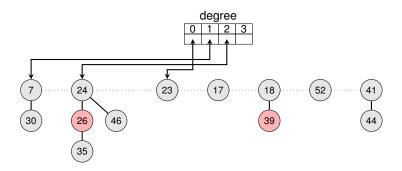


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



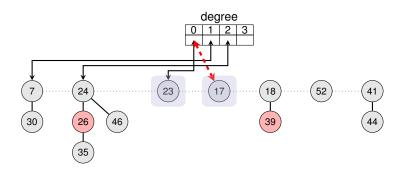


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



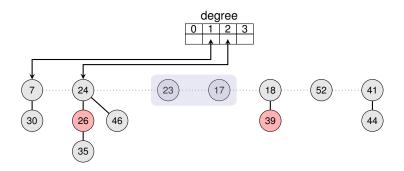


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



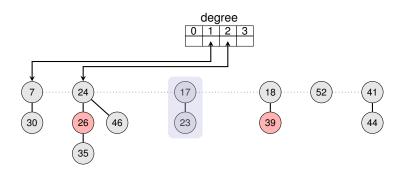


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



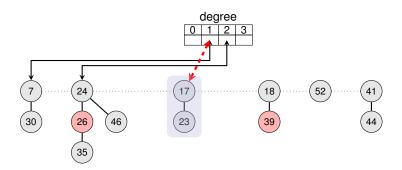


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



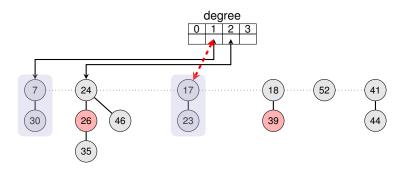


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



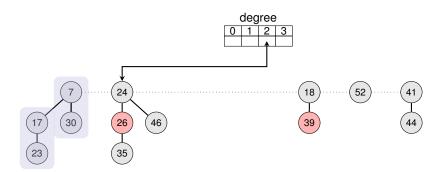


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



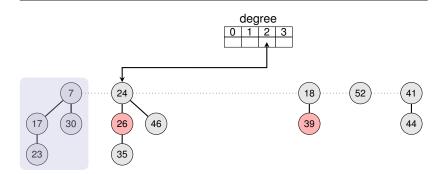


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



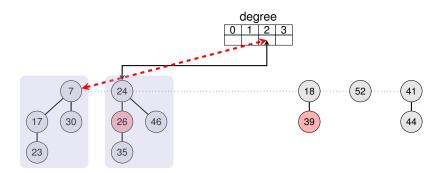


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)





- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)





#### — EXTRACT-MIN ———

- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)

# degree 0 1 2 3





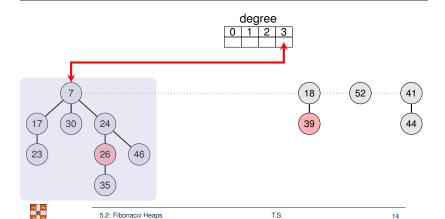
#### — EXTRACT-MIN ———

- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)

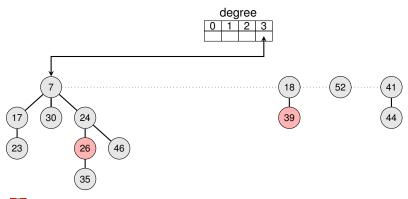
# degree 0 1 2 3



- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)

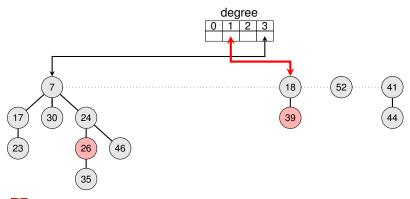


- Delete min √
- Meld childen into root list and unmark them
- Consolidate so that no roots have the same degree (# children)



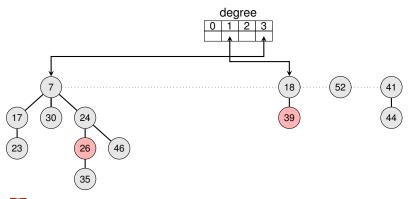
— Extract-Min ———

- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



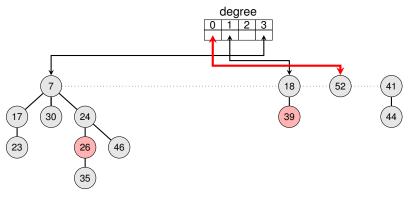


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



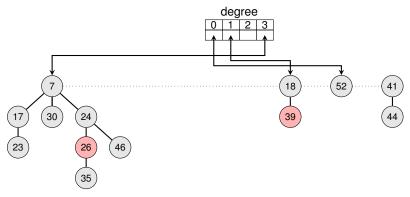


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



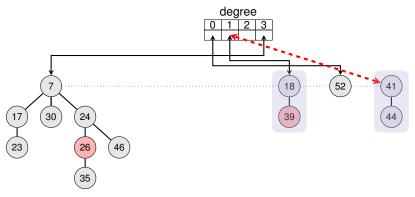


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



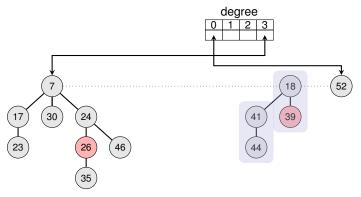


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



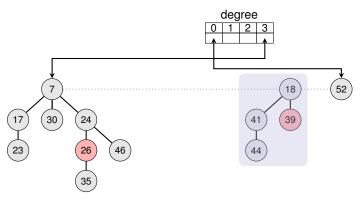


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



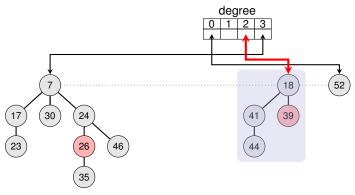


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)



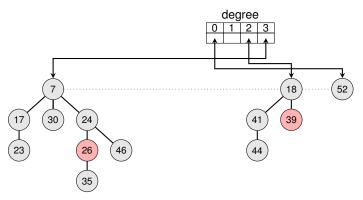


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)





- Delete min √
- Meld childen into root list and unmark them
- Consolidate so that no roots have the same degree (# children) ✓



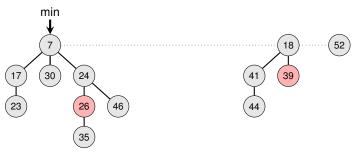


- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children) √
- Update minimum



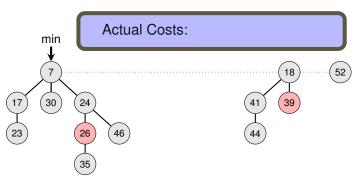


- Delete min √
- Meld childen into root list and unmark them
- Consolidate so that no roots have the same degree (# children) √
- Update minimum √





- Delete min √
- Meld childen into root list and unmark them
- Consolidate so that no roots have the same degree (# children) √
- Update minimum √





EXTRACT-MIN -

- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children) √
- Update minimum √

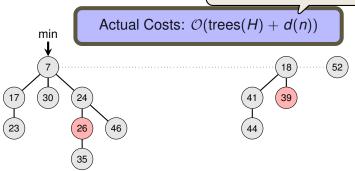
Every root becomes child of another root at most once! **Actual Costs:** min 30 39 17 23 46



- EXTRACT-MIN

- Delete min ✓
- Meld childen into root list and unmark them
- Consolidate so that no roots have the same degree (# children) √
- Update minimum √

d(n) is the maximum degree of a root in any Fibonacci heap of size n

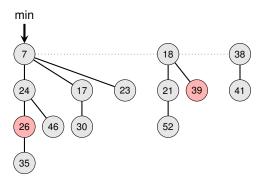




## Fibonacci Heap: Decrease-Key (First Attempt)

DECREASE-KEY of node x —

Decrease the key of x (given by a pointer)

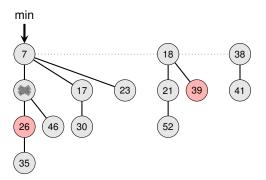




#### Fibonacci Heap: Decrease-Key (First Attempt)

DECREASE-KEY of node x —

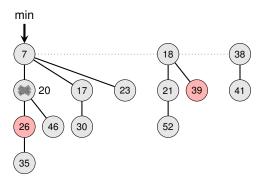
Decrease the key of x (given by a pointer)





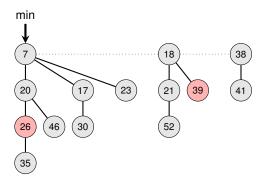
DECREASE-KEY of node x —

Decrease the key of x (given by a pointer)



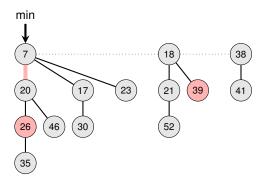


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated



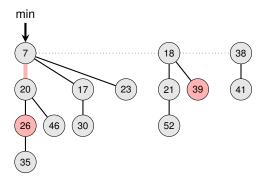


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated



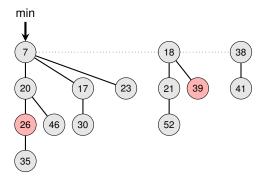


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not



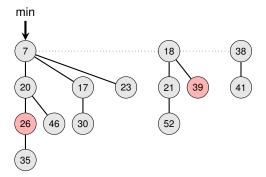


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.



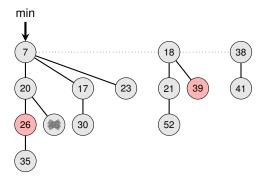


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise,



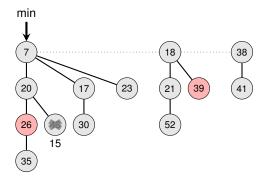


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise,



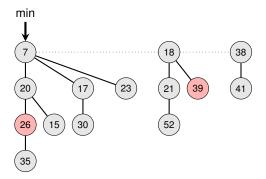


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise,



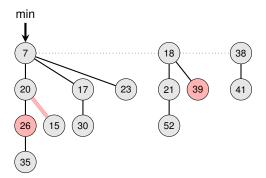


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise,



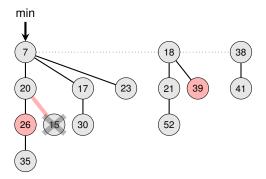


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise,



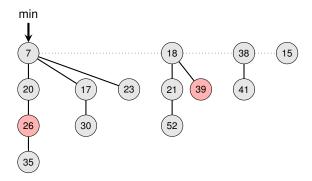


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



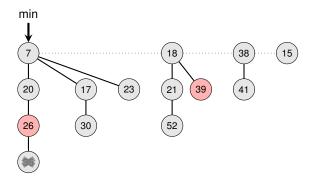


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
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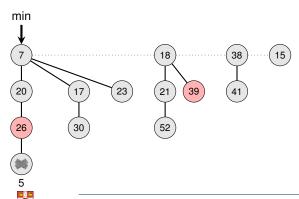
- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).





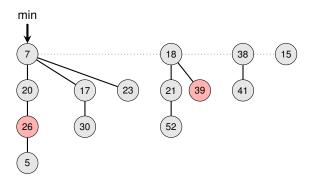
DECREASE-KEY of node x -

- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



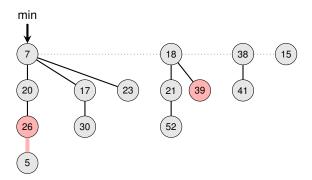
5.2: Fibonacci Heaps T.S. 15

- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



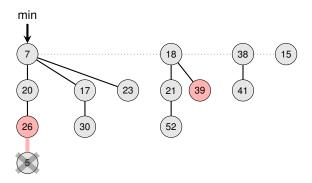


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



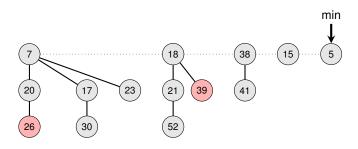


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



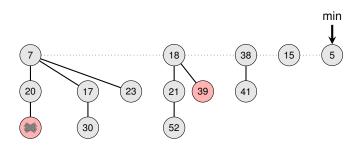


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



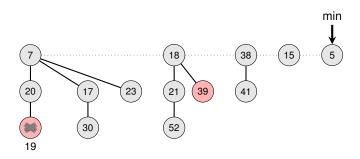


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at *x* and meld into root list (update min).



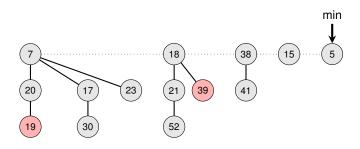


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



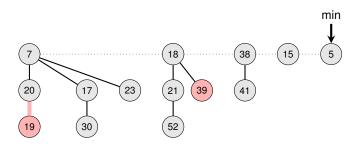


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



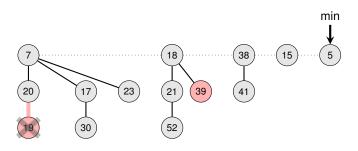


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



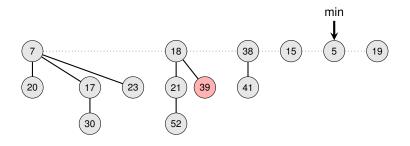


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at *x* and meld into root list (update min).



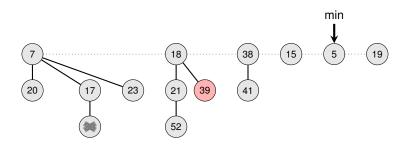


- Decrease the key of x (given by a pointer)
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  - If not, then done.
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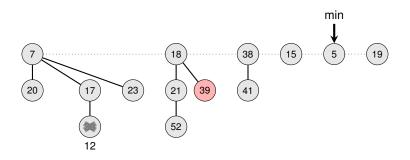


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



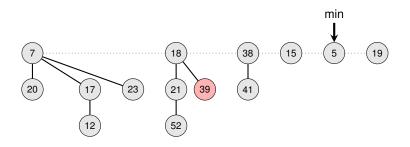


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
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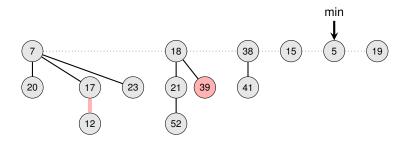


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



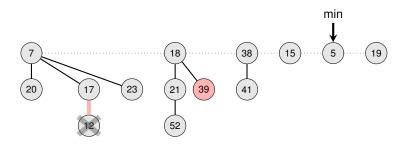


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



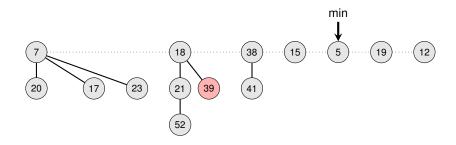


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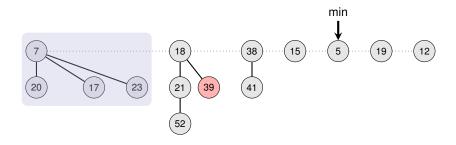


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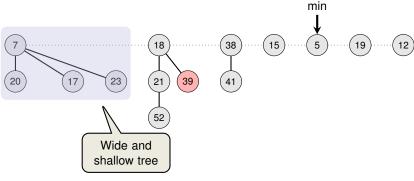


- Decrease the key of x (given by a pointer)
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  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).



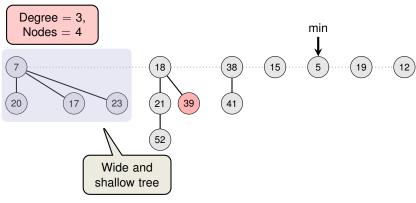


- Decrease the key of x (given by a pointer)
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  - If not, then done.
  - Otherwise, cut tree rooted at *x* and meld into root list (update min).





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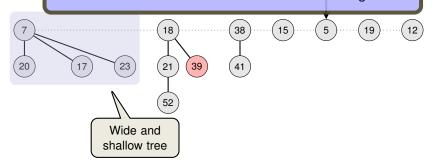




DECREASE-KEY of node x =

- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
  - If not, then done.
  - Otherwise, cut tree rooted at x and meld into root list (update min).

**Peculiar Constraint**: Make sure that each non-root node loses at most one child before becoming root

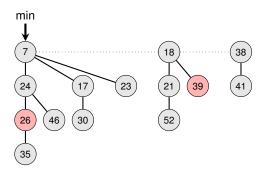




# Fibonacci Heap: DECREASE-KEY

DECREASE-KEY of node x —

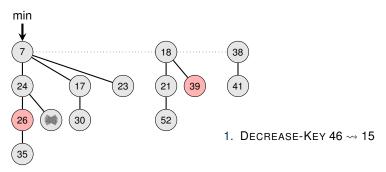
Decrease the key of x (given by a pointer)





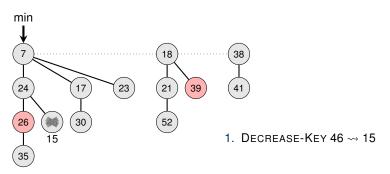
## Fibonacci Heap: DECREASE-KEY

- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)



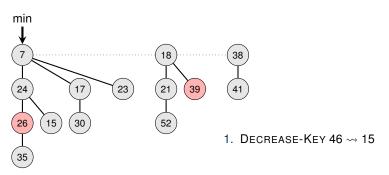


- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list



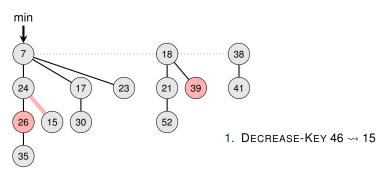


- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
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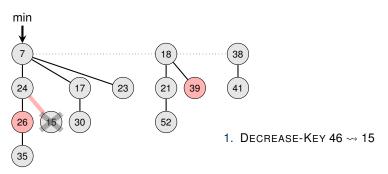
- Decrease the key of x (given by a pointer)
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DECREASE-KEY of node x -

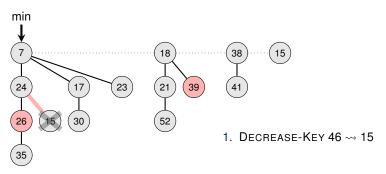
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list





16

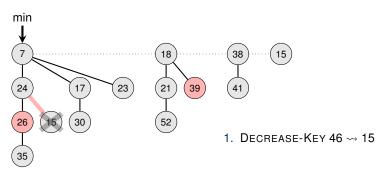
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list and:





DECREASE-KEY of node x =

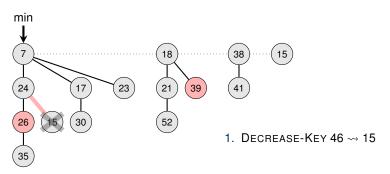
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list and:
  - Check if parent node is marked





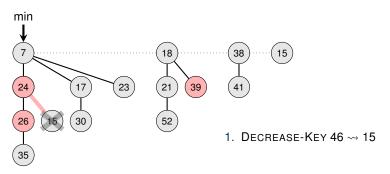
5.2: Fibonacci Heaps T.S. 16

- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list and:
  - Check if parent node is marked
    - If unmarked, mark it (unless it is a root)



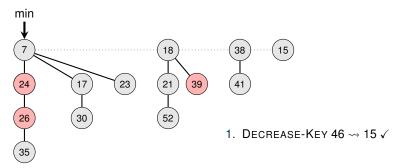


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- (Here we consider only cases where heap-order is violated)
- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list and:
  - Check if parent node is marked
    - If unmarked, mark it (unless it is a root)



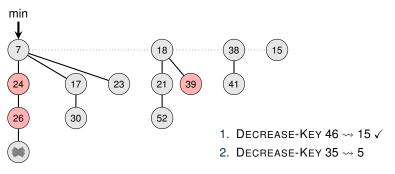


- Decrease the key of x (given by a pointer)
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- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list and:
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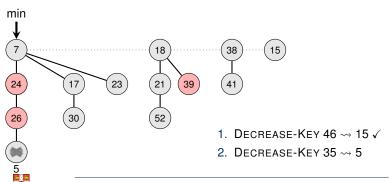
- Decrease the key of x (given by a pointer)
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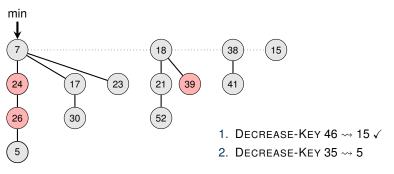
DECREASE-KEY of node x =

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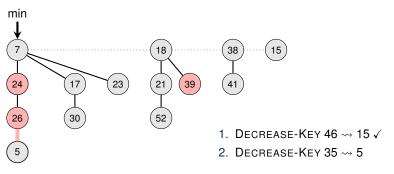
5.2: Fibonacci Heaps T.S. 16

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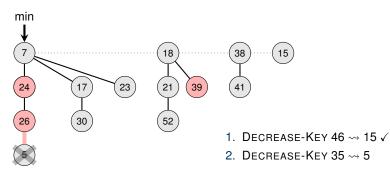
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DECREASE-KEY of node x =

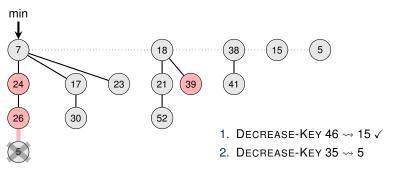
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5.2: Fibonacci Heaps T.S. 16

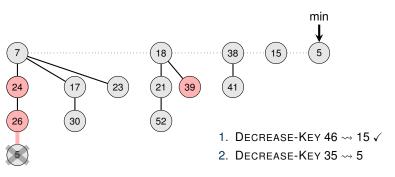
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DECREASE-KEY of node x =

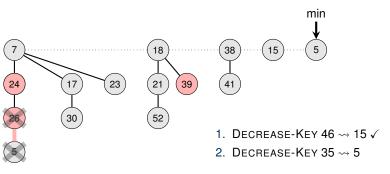
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- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list and:
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    - If unmarked, mark it (unless it is a root)
    - If marked,





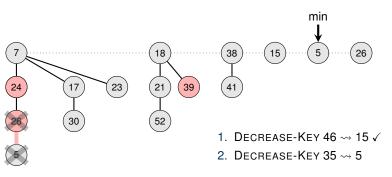
5.2: Fibonacci Heaps

- Decrease the key of x (given by a pointer)
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- $\Rightarrow$  Cut tree rooted at x, unmark x, meld into root list and:
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    - If unmarked, mark it (unless it is a root)
    - If marked, unmark and meld it into root list and recurse (Cascading Cut)



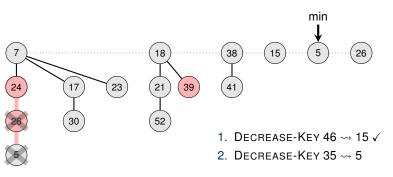


- Decrease the key of x (given by a pointer)
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  - Check if parent node is marked
    - If unmarked, mark it (unless it is a root)
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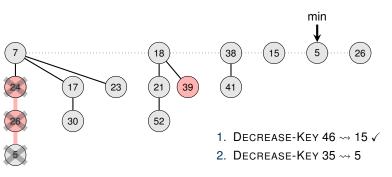


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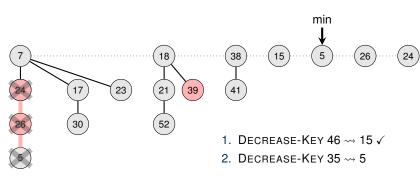


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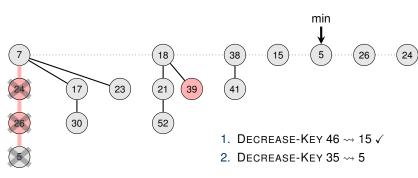




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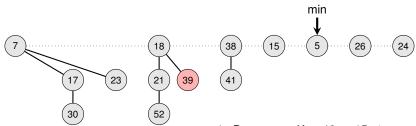


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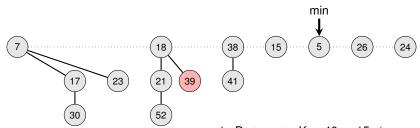
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- 1. Decrease-Key 46  $\leadsto$  15  $\checkmark$
- 2. Decrease-Key 35 → 5



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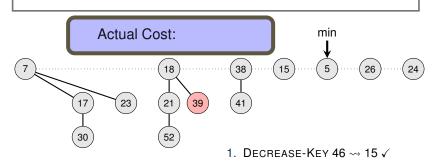


- 1. Decrease-Key 46  $\leadsto$  15  $\checkmark$
- 2. Decrease-Key 35 → 5 ✓



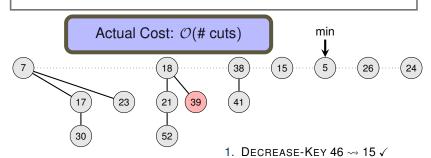
DECREASE-KEY of node x =

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- 2. Decrease-Key 35 ↔ 5 ✓



#### **Outline**

Structure

Operations

Glimpse at the Analysis



• INSERT: actual  $\mathcal{O}(1)$ 

■ EXTRACT-MIN: actual O(trees(H) + d(n))

• INSERT: actual  $\mathcal{O}(1)$ 

■ EXTRACT-MIN: actual  $\mathcal{O}(\text{trees}(H) + d(n))$ 

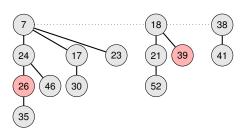
$$\Phi(H) = \mathsf{trees}(H) + 2 \cdot \mathsf{marks}(H)$$



• INSERT:  $actual \mathcal{O}(1)$ 

■ EXTRACT-MIN: actual O(trees(H) + d(n))

$$\Phi(H) = \mathsf{trees}(H) + 2 \cdot \mathsf{marks}(H)$$

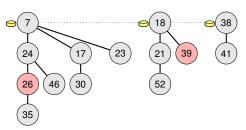




• INSERT: actual  $\mathcal{O}(1)$ 

■ EXTRACT-MIN: actual O(trees(H) + d(n))

$$\Phi(H) = \mathsf{trees}(H) + 2 \cdot \mathsf{marks}(H)$$

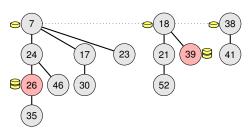




• INSERT: actual  $\mathcal{O}(1)$ 

■ EXTRACT-MIN: actual O(trees(H) + d(n))

$$\Phi(H) = \mathsf{trees}(H) + 2 \cdot \mathsf{marks}(H)$$





• INSERT:  $\operatorname{actual} \mathcal{O}(1)$  amortized  $\mathcal{O}(1)$ 

EXTRACT-MIN: actual  $\mathcal{O}(\operatorname{trees}(H) + d(n))$  amortized  $\mathcal{O}(d(n))$ 

■ DECREASE-Key: actual  $\mathcal{O}(\text{\# cuts}) \leq \mathcal{O}(\text{marks}(H))$  amortized  $\mathcal{O}(\text{1})$ 

$$\Phi(H) = \mathsf{trees}(H) + 2 \cdot \mathsf{marks}(H)$$

#### Lifecycle of a node

