

For $A, B \in \mathcal{P}(U)$, $A \cup B = \{x \in U \mid x \in A \vee x \in B\}$
 $A \cap B = \{x \in U \mid x \in A \wedge x \in B\}$

Sets and logic



$\mathcal{P}(U)$	$\{ \text{false}, \text{true} \}$
\emptyset	false
U	true
\cup	\vee
\cap	$\&$
$(\cdot)^c$	$\neg(\cdot)$

false true
 / /

$\rightarrow \mathcal{P}(\{*\}) = \{ \emptyset, \{*\} \}$

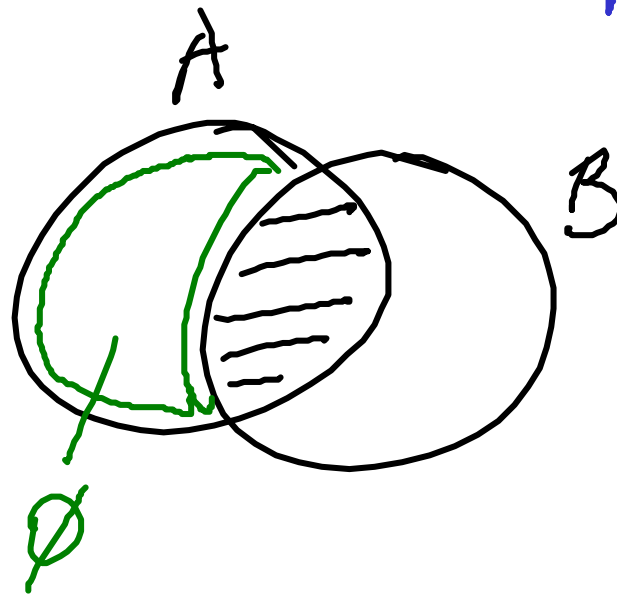
$$A \cup B = A$$

$$\Leftrightarrow B \subseteq A$$



$$A \cap B = A$$

$$\Leftrightarrow A \subseteq B$$



$$X \cup (X \cap A) = X$$

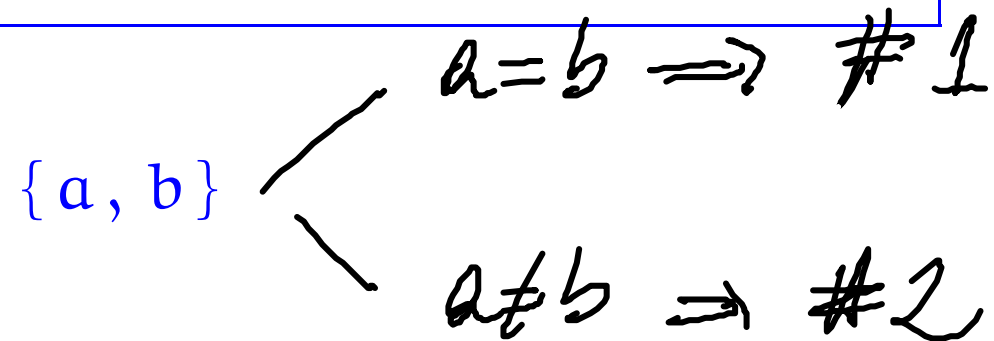
corresponds $X \cap A \subseteq X$

$$X \cap (X \cup A) = X$$

corresponds $X \subseteq X \cup A$

Unordered
Pairing axiom

For every a and b , there is a set with a and b as its only elements.



defined by

$$\forall x. x \in \{a, b\} \iff (x = a \vee x = b)$$

$$\{a, b\} = \{b, a\}$$

NB The set $\{a, a\}$ is abbreviated as $\{a\}$, and referred to as a singleton.

$$\left. \begin{array}{l} \text{for } a \neq b \\ \langle a, b \rangle \neq \langle b, a \rangle \end{array} \right\}$$

Ordered pairing

For every pair a and b , the set

$$\{\{a\}, \{a, b\}\}$$

is abbreviated as

$$\langle a, b \rangle$$

and referred to as an ordered pair.

Exercise Find other encodings for ordered pairing.

Proposition 71 (Fundamental property of ordered pairing)

For all a, b, x, y ,

$$\langle a, b \rangle = \langle x, y \rangle \iff (a = x \ \& \ b = y) .$$

PROOF: (\Leftarrow) ✓

$$(\Rightarrow) \{ \{a\}, \{a, b\} \} = \{ \{x\}, \{x, y\} \}$$

Case 1: $a = b \Rightarrow \{ \{a\}, \{a, b\} \} = \{ \{a\} \} = \{ \{x\}, \{x, y\} \}$

$$\{a\} = \{x\} \Rightarrow a = x$$

$$\{x\} = \{x, y\}$$

$$\{ \{a\} \} = \{ \{x\} \} = \{ \{x\}, \{x, y\} \} \iff x = y$$

Case 2 $a \neq b$

EXERCISE

In ML, the * type constructor.

Products

The product $A \times B$ of two sets A and B is the set

$$A \times B = \{ x \mid \exists a \in A, b \in B. x = (a, b) \}$$

where

$$\forall a_1, a_2 \in A, b_1, b_2 \in B.$$

$$(a_1, b_1) = (a_2, b_2) \iff (a_1 = a_2 \ \& \ b_1 = b_2) .$$

Thus,

$$\forall x \in A \times B. \exists! a \in A. \exists! b \in B. x = (a, b) .$$

in ordered
pairing
construction

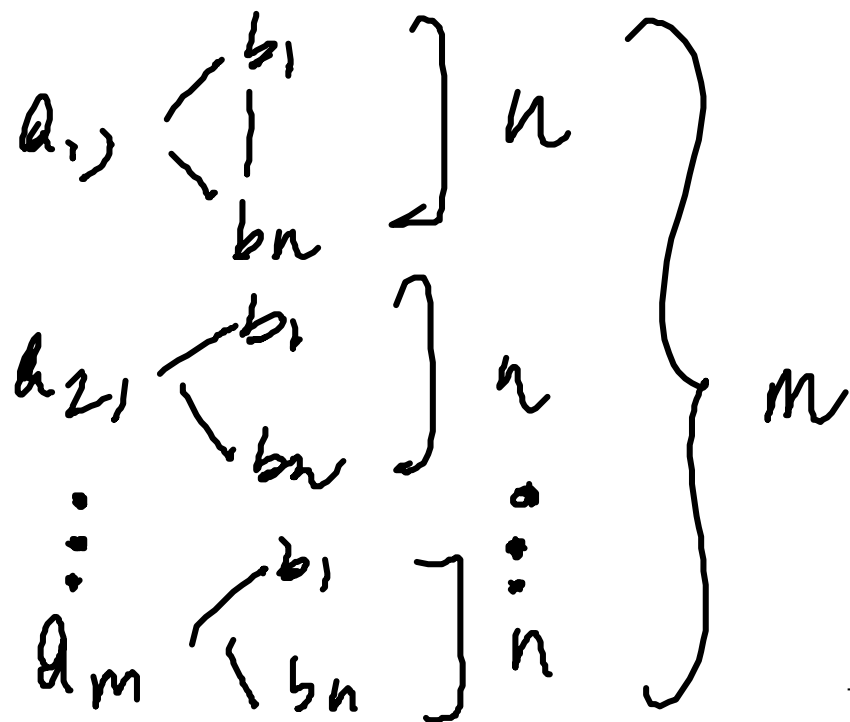
Proposition 73 For all finite sets A and B ,

$$\#(A \times B) = \#A \cdot \#B$$

PROOF IDEA: $A = \{a_1, \dots, a_m\}$ ($\#A = m$)

$B = \{b_1, \dots, b_n\}$ ($\#B = n$)

$$A \times B = \{(a_i, b_j) \mid i=1 \dots m, j=1 \dots n\}$$



total $n + \dots + n = m \cdot n$



Big unions and intersections

Given A and B we considered $A \cup B$

Given A_1, A_2, \dots, A_n we have $A_1 \cup A_2 \cup \dots \cup A_n$

Consider $\mathcal{F} \subseteq \mathcal{P}(U)$
{ a set of subsets of U

and define $\bigcup \mathcal{F} = \{x \in U \mid \exists X \in \mathcal{F}. x \in X\}$

notation

$$\bigcup_{X \in \mathcal{F}} X$$

example $\mathcal{F} = \{A_1, A_2, \dots, A_n\}$

$$\bigcup \mathcal{F} = A_1 \cup A_2 \cup \dots \cup A_n$$

Intuition: In ML

flatten: $\alpha \text{ list list} \rightarrow \alpha \text{ list}$

$[L_1, \dots, L_n] \mapsto L_1 @ \dots @ L_n$

Given $\mathcal{F} \subseteq \mathcal{P}(U)$

Consider

$$\bigcap \mathcal{F} = \{x \in U \mid \forall X \in \mathcal{F}, x \in X\}$$

$$\bigcap \{A_1, \dots, A_n\} = A_1 \cap \dots \cap A_n$$

a closure property

Theorem 74 Let

$$\mathcal{F} = \left\{ S \subseteq \mathbb{R} \mid (0 \in S) \ \& \ (\forall x \in \mathbb{R}. x \in S \implies (x+1) \in S) \right\} .$$

Then, (i) $\mathbb{N} \in \mathcal{F}$ and (ii) $\mathbb{N} \subseteq \bigcap \mathcal{F}$. Hence, $\bigcap \mathcal{F} = \mathbb{N}$.

PROOF:

$$(ii) \forall n \in \mathbb{N}. n \in \bigcap \mathcal{F}$$

$$\iff \forall n \in \mathbb{N}.$$

$$\boxed{\forall S \in \mathcal{F}. n \in S}$$

$$\updownarrow P(n)$$

Prove $\forall n \in \mathbb{N}. P(n)$
by induction.

Exercise

Example

$$\mathbb{R} \in \mathcal{F}$$

$$\mathbb{Q} \in \mathcal{F}$$

$$\mathbb{Z} \in \mathcal{F}$$

$$\{0\} \notin \mathcal{F}$$

$$\emptyset \notin \mathcal{F}$$

\mathbb{N} is the least set of numbers containing 0 and closed under successor.

For $\mathcal{F} \subseteq \mathcal{P}(\mathcal{P}(U))$,

$$\bigcup (\bigcup \mathcal{F}) = \bigcup \{ \bigcup \mathcal{A} \mid \mathcal{A} \in \mathcal{F} \}$$

or, in indexed notation,

$$\bigcup_{x \in \left(\bigcup_{\mathcal{A} \in \mathcal{F}} \mathcal{A} \right)} x = \bigcup_{\mathcal{A} \in \mathcal{F}} \left(\bigcup_{x \in \mathcal{A}} x \right)$$

IDEA:

$$F = \{ \sim, A, \neg \} \quad \mathcal{A} = \{ \sim, X, \neg \}$$

$$\begin{aligned} \cup F &= \sim \cup A \cup \neg \\ &= \{ \sim, X, \neg \} \end{aligned}$$

$$\cup(\cup F) = \sim \cup X \cup \neg$$

$$\cup \mathcal{A} = (\sim \cup X \cup \neg)$$

$$\cup \{ \cup \mathcal{A} \mid \mathcal{A} \in F \}$$

$$= \sim \cup (\sim \cup X \cup \neg) \cup \neg$$

$$\text{PROOF: } z \in \bigcup_{X \in \left(\bigcup_{\mathcal{A} \in \mathcal{F}} \mathcal{A} \right)} X$$

$$\Leftrightarrow \exists x. x \in \left(\bigcup_{\mathcal{A} \in \mathcal{F}} \mathcal{A} \right) \ \& \ z \in x$$

$$\Leftrightarrow \exists x. \exists \mathcal{A}. \mathcal{A} \in \mathcal{F} \ \& \ x \in \mathcal{A} \ \& \ z \in x$$

$$\Leftrightarrow \exists \mathcal{A}. \mathcal{A} \in \mathcal{F} \ \& \ \exists x. x \in \mathcal{A} \ \& \ z \in x$$

$$\Leftrightarrow \exists \mathcal{A}. \mathcal{A} \in \mathcal{F} \ \& \ z \in \bigcup_{x \in \mathcal{A}} x$$

$$\Leftrightarrow z \in \bigcup_{\mathcal{A} \in \mathcal{F}} \left(\bigcup_{x \in \mathcal{A}} x \right)$$

Compare the previous identity with the following one for lists.

For $L: \alpha \text{ list list list}$,

$$\text{flatten}(\text{flatten } L) = \text{flatten}(\text{map flatten } L)$$

But, this is one of the laws of a mathematical structure called a MONAD, which has become a fundamental notion in functional programming.

notation: $\cup \{A, B\} = A \cup B$

Union axiom

Every collection of sets has a union.

$$\cup \mathcal{F}$$

$$x \in \cup \mathcal{F} \iff \exists X \in \mathcal{F}. x \in X$$

For non-empty \mathcal{F} we also have

$$\bigcap \mathcal{F}$$

defined by

$$\forall x. x \in \bigcap \mathcal{F} \iff (\forall X \in \mathcal{F}. x \in X) .$$