

# **Compiler Construction**

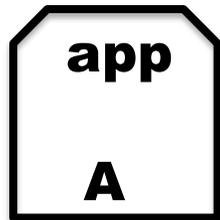
## **Lent Term 2014**

### **Lecture 9 (of 16)**

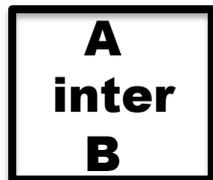
- **Assorted topics**
  - **bootstrapping**
  - **exceptions**

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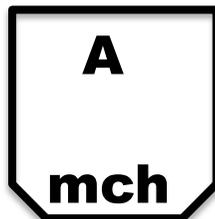
# Bootstrapping. We need some notation . . .



An application called **app** written in language **A**

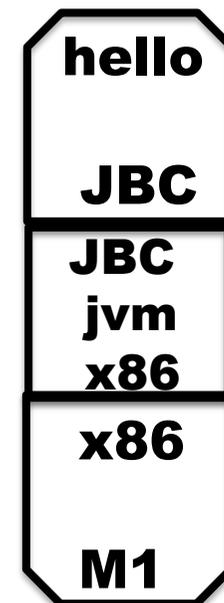
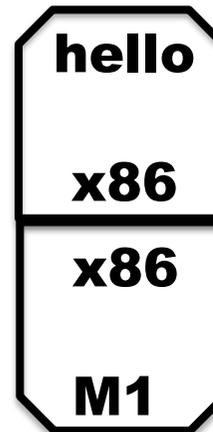


An interpreter or VM for language **A**  
Written in language **B**

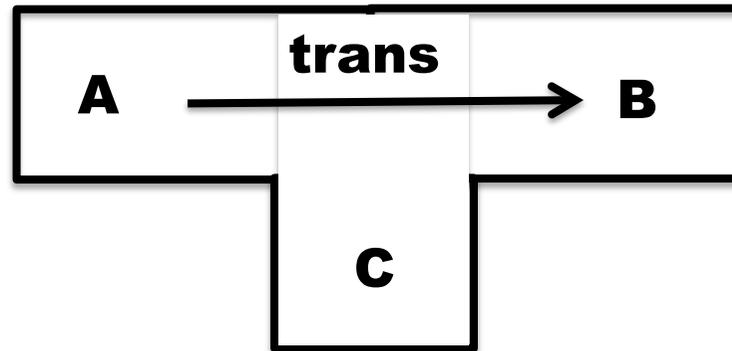


A machine called **mch** running language **A** natively.

## Simple Examples

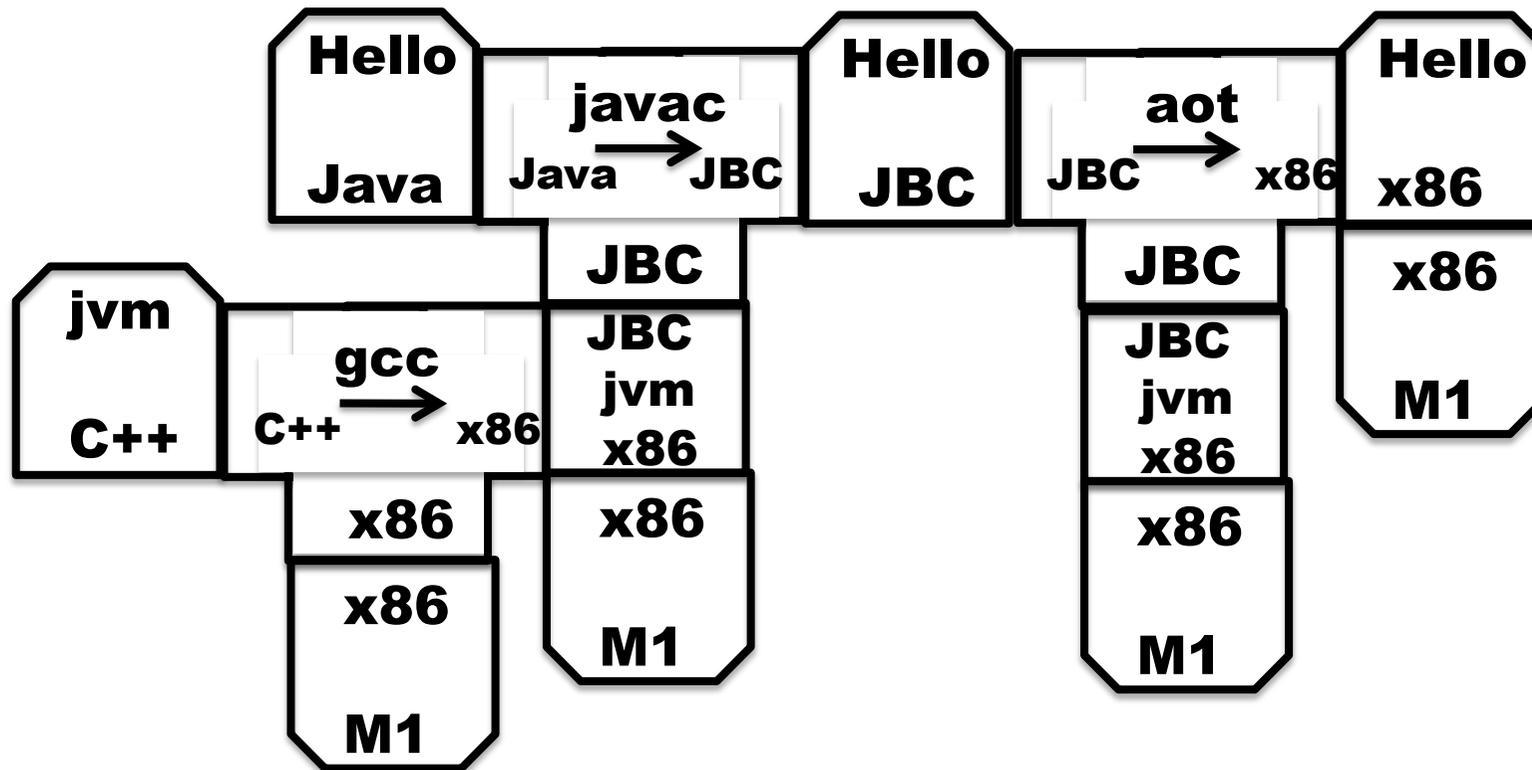


# Tombstones



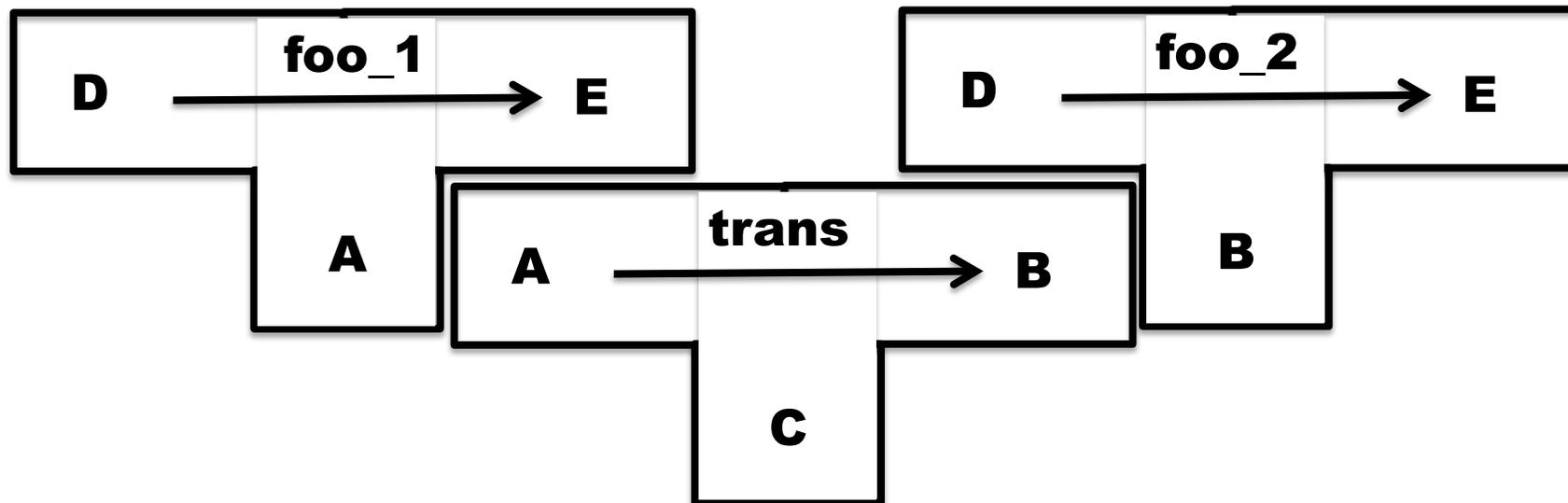
This is an application called **trans** that translates programs in language **A** into programs in language **B**, and it is written in language **C**.

# Ahead-of-time compilation



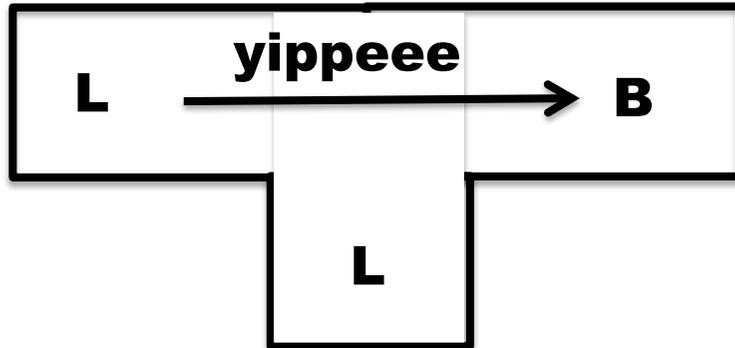
Thanks to David Greaves  
for the example.

## Of course translators can be translated

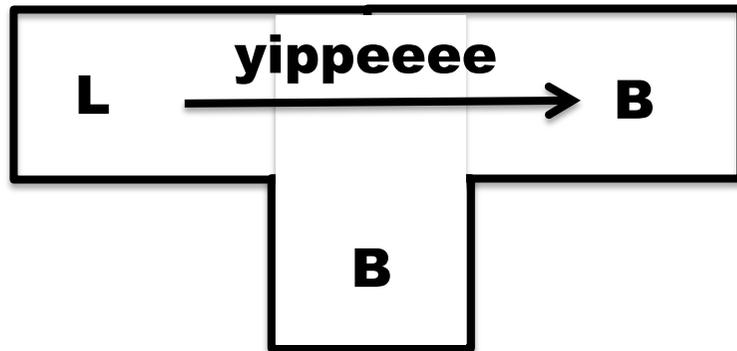


Translator **foo\_2** is produced as output from **trans** when given **foo\_1** as input.

## Our seemingly impossible task



We have just invented a really great new language **L** (in fact we claim that “**L** is far superior to C++”). To prove how great **L** is we write a compiler for **L** in **L** (of course!). This compiler produces machine code **B** for a widely used instruction set (say **B** = x86).



Furthermore, we want to compile our compiler so that it can run on a machine running **B**.

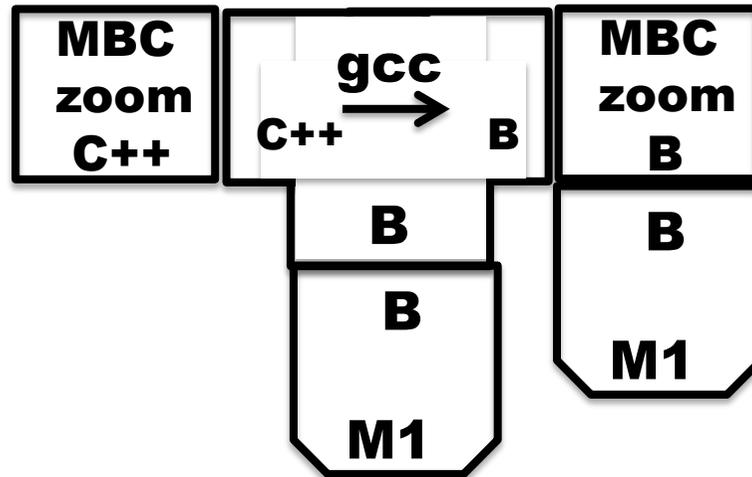
**How can we compile our compiler?**

There are many many ways we could go about this task. The following slides simply sketch out one plausible route to fame and fortune.

# Step 1

## Write a small interpreter (VM) for a small language of byte codes

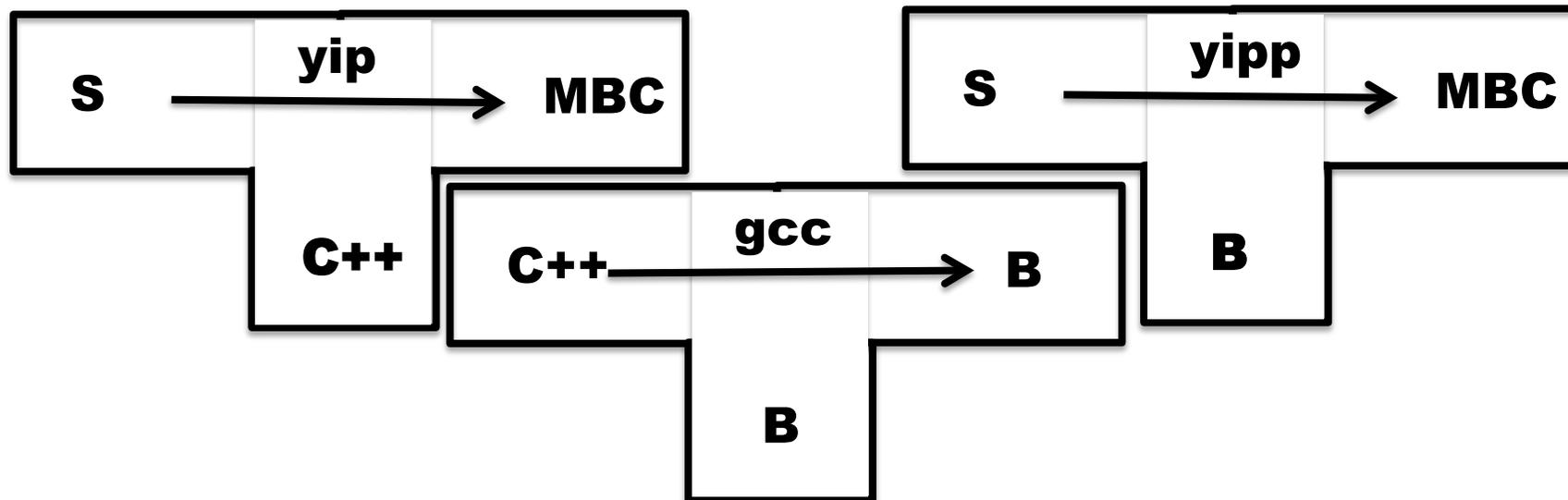
**MBC** = My Byte Codes



The **zoom** machine!

## Step 2

Pick a small subset **S** of **L** and  
write a translator from **S** to **MBC**

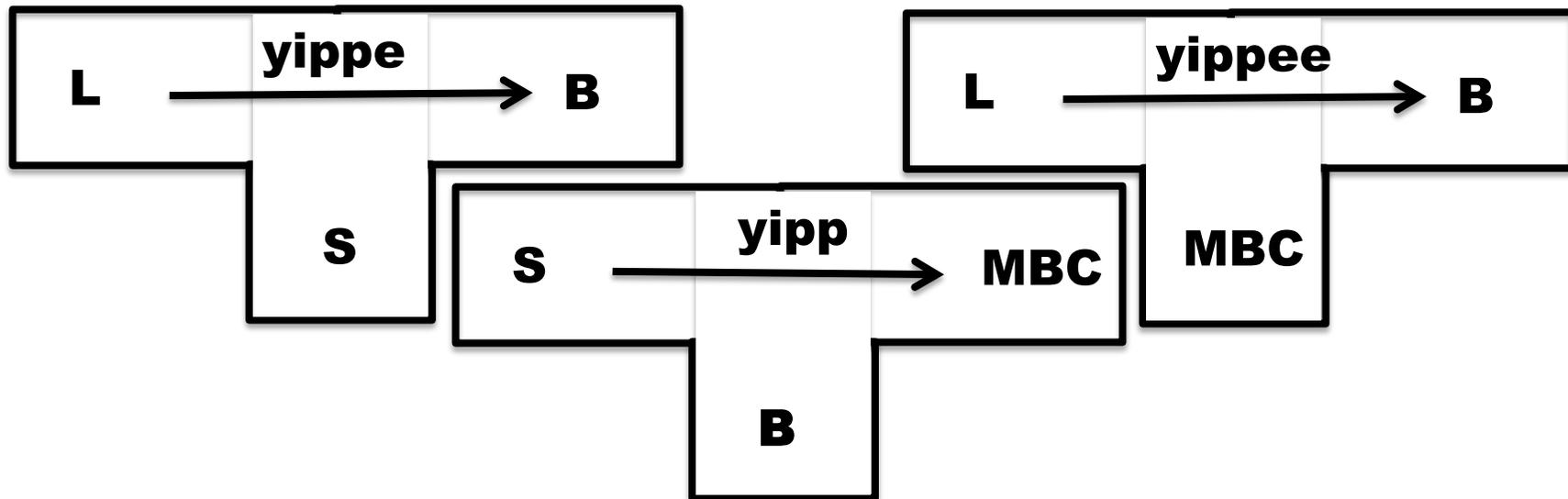


Write **yip** by hand. (It sure would be nice if we could hide the fact that this is written in C++.)

Translator **yipp** is produced as output from **gcc** when **yip** is given as input.

## Step 3

### Write a compiler for L in S

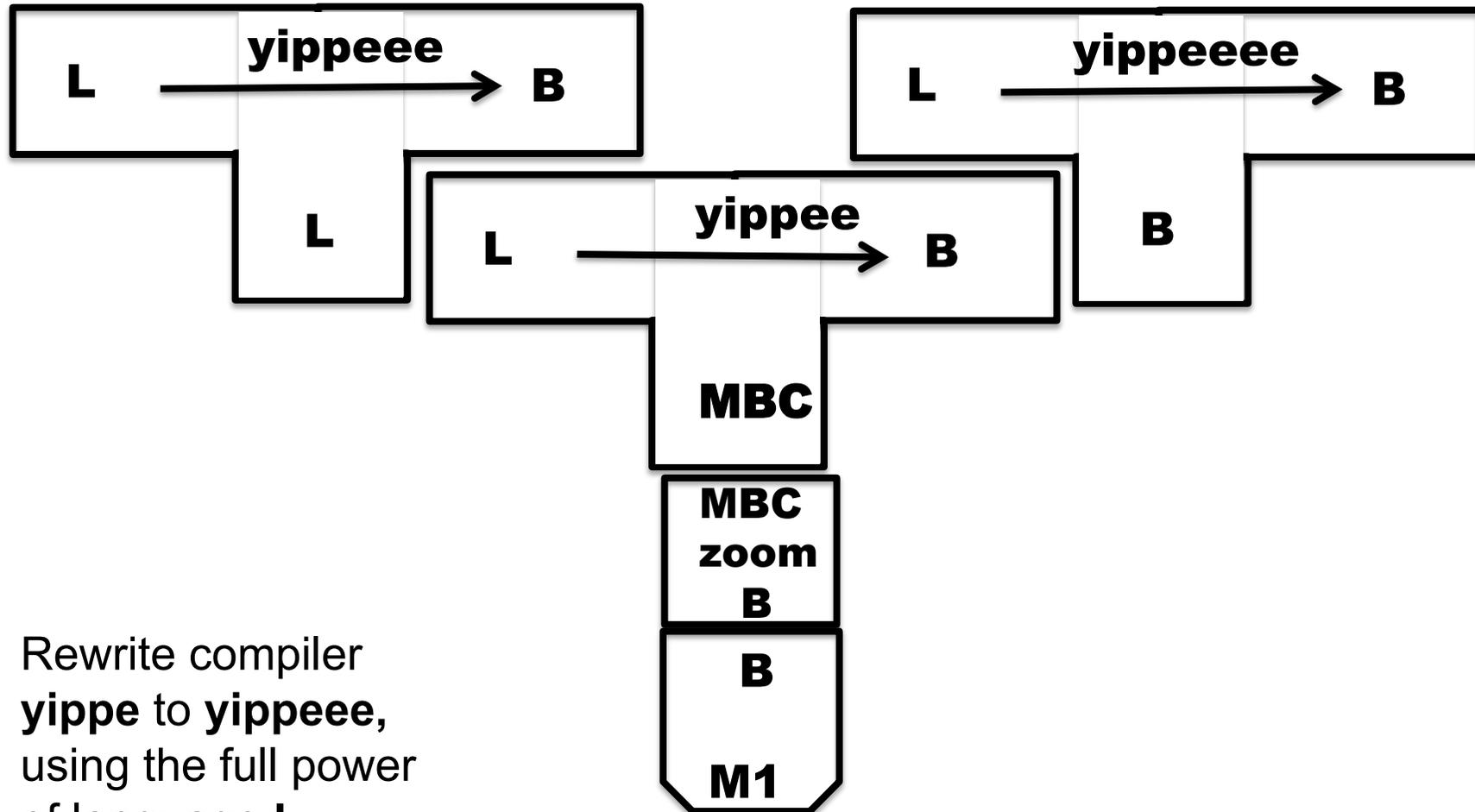


Write a compiler **yippe** for the full language **L**, but written only in the sub-language **S**.

Compile **yippe** using **yipp** to produce **yippee**

## Step 4

### Write a compiler for L in L



Rewrite compiler **yippe** to **yippeee**, using the full power of language **L**.

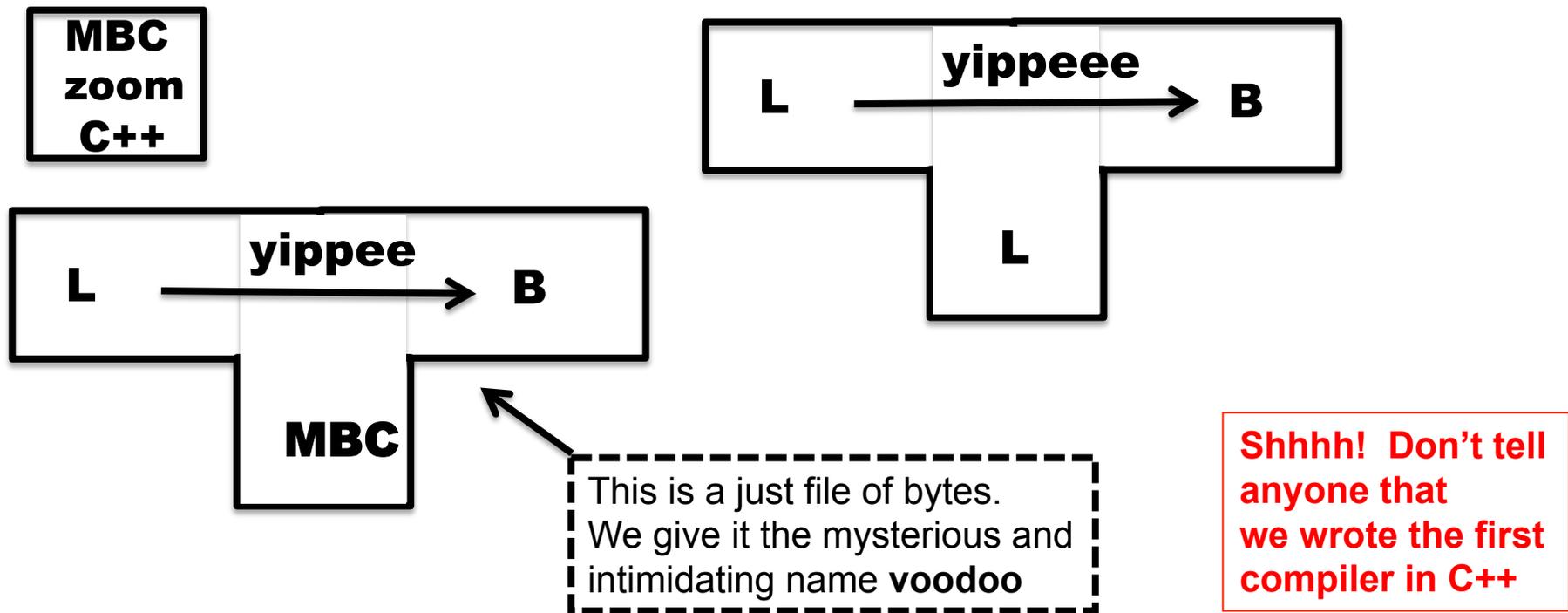
Now compile this using **yippee** to obtain our goal!



## Step 5

# Cover our tracks and leave the world mystified and amazed

Our L compiler download site contains only three components:



Our instructions:

1. Use **gcc** to compile the **zoom** interpreter
2. Use **zoom** to run **voodoo** with input **yippee** to produce output the compiler **yippee**

## New Topic : Exceptions (informal description)

`e handle f`

If expression `e` evaluates “normally” to value `v`, then `v` is the result of the entire expression.

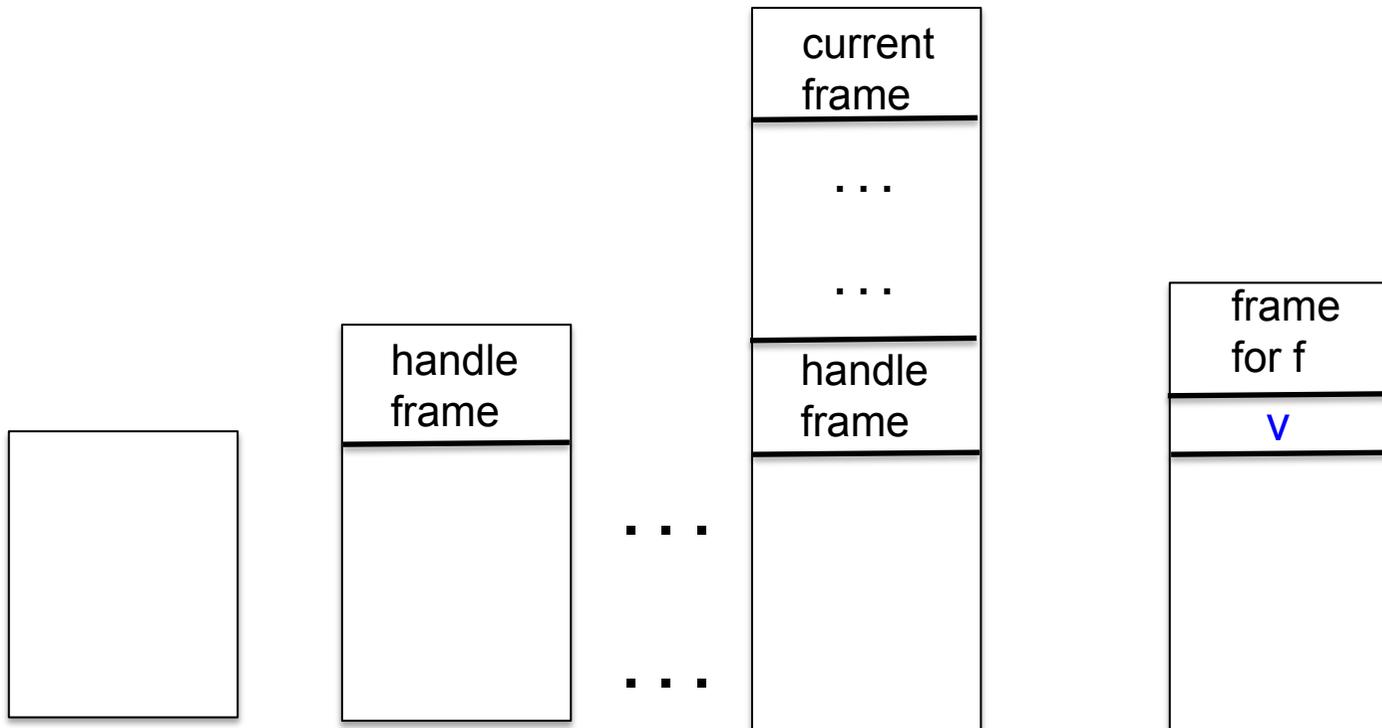
Otherwise, an exceptional value `v'` is “raised” in the evaluation of `e`, then result is `(f v')`

`raise e`

Evaluate expression `e` to value `v`, and then raise `v` as an exceptional value, which can only be “handled”.

Implementation of exceptions may require a lot of language-specific consideration and care. Exceptions can interact in powerful and unexpected ways with other language features. Think of C++ and class destructors, for example.

# Viewed from the call stack



Call stack just before evaluating code for

`e handle f`

Push a special frame for the handle

“`raise v`” is encountered while evaluating a function body associated with top-most frame

“Unwind” call stack. Depending on language, this may involve some “clean up” to free resources.

## Possible pseudo-code implementation

e handle f

```
let fun _h27 () =  
  build special "handle frame"  
  save address of f in frame;  
  ... code for e ...  
  return value of e  
in _h27 () end
```

raise e

```
... code for e ...  
save v, the value of e;  
unwind stack until first  
fp found pointing at a handle frame;  
Replace handle frame with frame  
for call to (extracted) f using  
v as argument.
```