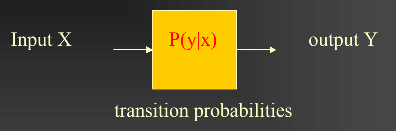
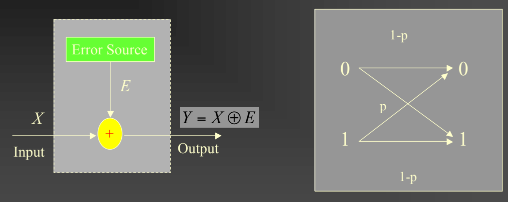
some channel models



memoryless:

- output at time i depends only on input at time i
- input and output alphabet finite

Example: binary symmetric channel (BSC)

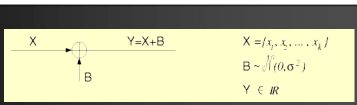


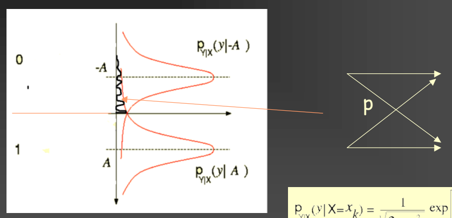
E is the binary error sequence s.t. P(1) = 1-P(0) = p

X is the binary information sequence

Y is the binary output sequence

from AWGN to BSC

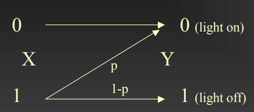




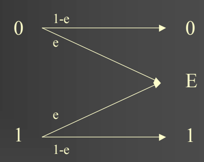
$$p_{Y|X}(y|X=x_k) = \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left[-\frac{(y-x_k)^2}{2\sigma^2}\right]$$

Homework: calculate the capacity as a function of A and σ^2

Other models



P(X=0) = P_0

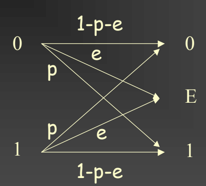


P(X=0) = P_0

Z-channel (optical) (MAC)

Erasure channel

Erasure with errors



burst error model (Gilbert-Elliot)

Random error channel; outputs independent

Error Source → P(0) = 1- P(1);

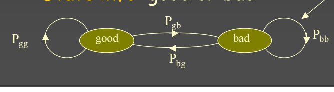
Burst error channel; outputs dependent

Error Source → P(0 | state = bad) = P(1|state = bad) = 1/2;

P(0 | state = good) = 1 - P(1|state = good) = 0.999

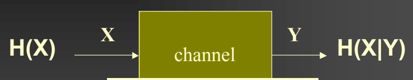
State info: good or bad

transition probability



channel capacity:

I(X;Y) = H(X) - H(X|Y) = H(Y) – H(Y|X) (Shannon 1948)



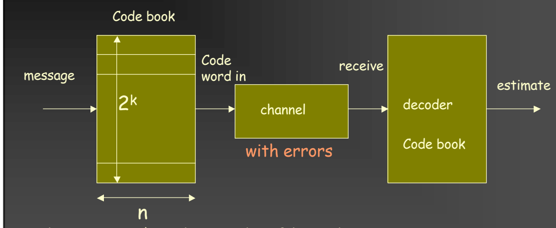
$$\max I(X;Y) = capacity$$

$$H = Entropy$$

$$E(p) = \sum p(i) * I(i)$$

notes:
capacity depends on input probabilities
because the transition probabilites are fixed

Practical communication system design



There are 2^k code words of length n

k is the number of information bits transmitted in n channel uses

Channel capacity

Definition:

The rate R of a code is the ratio k/n, where

k is the number of information bits transmitted in n channel uses

Shannon showed that: :

for R ≤ C

encoding methods exist

with decoding error probability ⇨ 0

Encoding and decoding according to Shannon

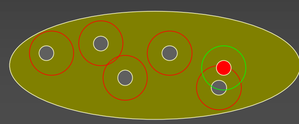
Code: 2^k binary codewords where p(0) = P(1) = ½

Channel errors: P(0 →1) = P(1 → 0) = p

i.e. # error sequences ≈ $2^{nh(p)}$

Decoder: search around received sequence for codeword

with ≈ np differences



space of 2^n binary sequences

decoding error probability

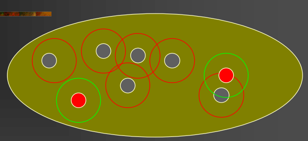
1. For t errors: |t/n-p|> Є

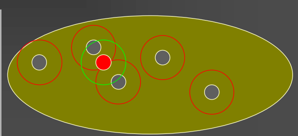
→ 0 for n → ∞

(law of large numbers)

2. > 1 code word in region

(codewords random)





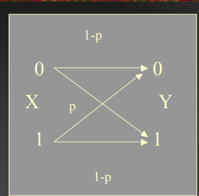
$$P(>1) \approx (2^k - 1)\frac{2^{nh(p)}}{2^n}$$

$$\to 2^{-n(1-h(p)-R)} = 2^{-n(C_{BSC}-R)} \to 0$$

$$\text{for} \quad R = \frac{k}{n} < 1 - h(p)$$

$$\text{and} \quad n \to \infty$$

channel capacity: the BSC



I(X;Y) = H(Y) – H(Y|X)

the maximum of H(Y) = 1

since Y is binary

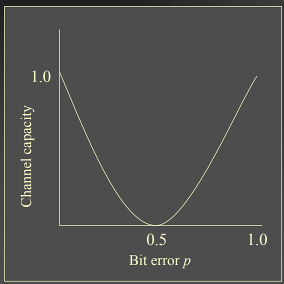
H(Y|X) = h(p)

= P(X=0)h(p) + P(X=1)h(p)

Conclusion: the capacity for the BSC C_{BSC} = 1- h(p)

Homework: draw C_{BSC} , what happens for p > ½

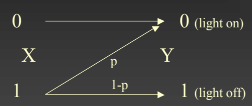
channel capacity: the BSC



Explain the behaviour!

channel capacity: the Z-channel

Application in optical communications



$P(X=0) = P_0$

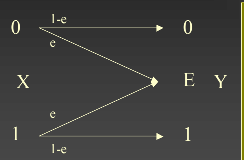
$H(Y) = h(P_0 + p(1 - P_0))$

$H(Y|X) = (1 - P_0)\, h(p)$

For capacity,
maximize $I(X;Y)$ over P_0

channel capacity: the erasure channel

Application: cdma detection



$P(X=0) = P_0$

$I(X;Y) = H(Y) - H(Y|X)$

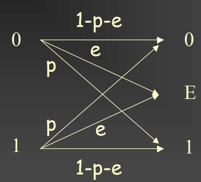
$H(Y) = h(P_0)$

$H(Y|X) = e\, h(P_0)$

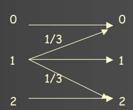
Thus $C_{erasure} = 1 - e$

(check!, draw and compare with BSC and Z)

Erasure with errors: calculate the capacity!



example



- Consider the following example

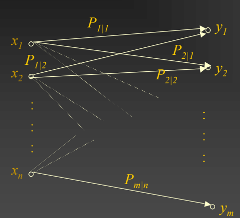
- For P(0) = P(2) = p, P(1) = 1-2p

  $H(Y) = h(1/3 - 2p/3) + (2/3 + 2p/3)$; $H(Y|X) = (1-2p)\log_2 3$

  Q: maximize H(Y) – H(Y|X) as a function of p
  Q: is this the capacity?

  hint use the following: $\log_2 x = \ln x / \ln 2$; $d \ln x / dx = 1/x$

channel models: general diagram



Input alphabet $X = \{x_1, x_2, \ldots, x_n\}$
Output alphabet $Y = \{y_1, y_2, \ldots, y_m\}$
$P_{j|i} = P_{Y|X}(y_j|x_i)$

In general:
calculating capacity needs more theory

The statistical behavior of the channel is completely defined by the channel transition probabilities $P_{j|i} = P_{Y|X}(y_j|x_i)$

* clue:

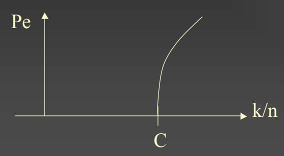
$I(X;Y)$

is convex ∩ in the input probabilities

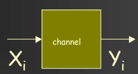
i.e. finding a maximum is simple

Channel capacity: converse

For $R > C$ the decoding error probability > 0

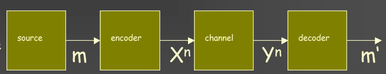


Converse: For a discrete memory less channel



$$I(X^n;Y^n) = H(Y^n) - \sum_{i=1}^{n} H(Y_i \mid X_i) \le \sum_{i=1}^{n} H(Y_i) - \sum_{i=1}^{n} H(Y_i \mid X_i) = \sum_{i=1}^{n} I(X_i;Y_i) \le nC$$

Source generates one out of 2^k equiprobable messages



Let Pe = probability that $m' \neq m$

converse $R := k/n$

$k = H(M) = I(M;Y^n) + H(M|Y^n)$

$\le \underbrace{I(X^n;Y^n)}_{X^n \text{ is a function of } M} + \underbrace{1 + k\,Pe}_{\text{Fano}}$

$\le nC + 1 + k\,Pe$

$1 - C\,n/k - 1/k \le Pe$

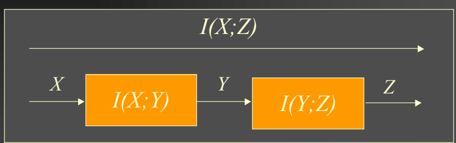
$Pe \ge 1 - C/R - 1/nR$

Hence: for large n, and $R > C$,

the probability of error $Pe > 0$

We used the data processing theorem

Cascading of Channels



The overall transmission rate $I(X;Z)$ for the cascade can not be larger than $I(Y;Z)$, that is:

$$I(X;Z) \le I(Y;Z)$$