Basic set theory and logic
Prerequisites for Language 1
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The lectures in Language 1 assume some very basic knowledge of set theory and logic. Students need to be familiar with the following concepts:

1. The basic idea of a set, set intersection, set union etc. Venn diagrams.
2. Propositional logic, interpretation with respect to a model (expressed in terms of sets).
4. First order predicate calculus, universal and existential quantifiers.
5. Basic logical inference (e.g., modus ponens, double-negation elimination).

This document contains some exercises which should serve to indicate the background that is assumed. Students who have difficulty with the exercises should work through an introductory logic book: suggestions for reading are given at the end of the document.

1 Notation

Unfortunately, logical notation is not standardized. The following notation is assumed in the examples and in the course in general:

**Set theory**

\{a, b, c\} the set containing a, b and c
\(\in\) set membership
\(\emptyset\) the empty set
\(\subset\) proper subset (i.e., not =)
\(\subseteq\) subset
\(=\) equality
\(\cap\) intersection
\(\cup\) union
\(-\) set difference

**Logic**

\(\neg\) negation
\(\wedge\) conjunction
\(\vee\) disjunction
\(\Rightarrow\) implication
\(\Leftrightarrow\) mutual implication
\(\forall\) universal quantifier
\(\exists\) existential quantifier
2 Exercises

2.1 Set theory and logic

1. Draw shaded set diagrams (Venn diagrams) corresponding to:
   (a) \( A \cap (B \cup C) \)
   (b) \( A - B \) (i.e., set difference)

2. Draw set diagrams corresponding to the following logical expressions:
   (a) \( \neg P \lor Q \)
   (b) \( \neg(P \land Q) \land R \)

3. The sentence every cat does not sleep has two possible interpretations:
   (a) No cat sleeps
   (b) It is not the case that every cat sleeps

Draw the Venn diagrams for these two interpretations, and give a specific model (i.e., statement about what is true in some world) such that b) is true but a) isn’t. Is it possible for a) to be true without b) being true?

2.2 Propositional logic and quantifier-free predicate logic

1. Draw the truth table for \( P \Rightarrow Q \)

2. Which of the following expressions are tautologies or contradictions? (Show the truth tables)
   (a) \( P \Rightarrow \neg P \)
   (b) \( \neg(P \Rightarrow (P \lor Q)) \)
   (c) \( \neg(P \lor Q) \Rightarrow P \)

3. Assume that:
   Rover barks and it-is-not-the-case-that Kitty chases Rover

   corresponds to the logical expression:
   \( \text{bark}'(r) \land \neg \text{chase}'(k, r) \)

   where \( k \) and \( r \) are constants corresponding to Kitty and Rover. We use the prime symbol (e.g., bark') to indicate we are talking about the predicate and not the word. A logical expression capturing the meaning of a sentence (approximately) is called the logical form for the sentence. I have used the artificial word ‘it-is-not-the-case-that’ to make the meaning of the sentence clearer for this example. We will also assume for the sake of the exercise that the logical connectives correspond to their English counterparts.

Show a logical form or forms for each of the following sentences:
   (a) it-is-not-the-case-that Kitty chases Rover or Rover sleeps
   (b) if Lynx sleeps then Rover barks or Rover sleeps

   The sentences may be ambiguous, in which case you should give more than one logical form.

4. The truth of a logical form can be evaluated with respect to a model (i.e., statement of what is true in some world). For instance, in the question above:
   \( \text{bark}'(r) \land \neg \text{chase}'(k, r) \)
could be evaluated with respect to the following model:

- chase': \{\langle l, k \rangle, \langle l, l \rangle, \langle l, r \rangle\}
- sleep': \{\}
- bark': \{r\}

In this case, bark'(r) is true, and chase'(k, r) is not true, so the expression as a whole is true.

Evaluate the truth of each logical form that you gave in the answer to the previous question with respect to this model.

### 2.3 First order predicate calculus (FOPC)

1. For each of the following expressions, underline the scope of each quantifier, and mark any free variables:
   
   - (a) \( \forall x[P(x) \Rightarrow [Q(x) \land P(y)]] \)
   - (b) \( \exists z[\forall x[3y[P(x) \land Q(y) \Rightarrow \neg S(x, y)] \land P(x)] \land \neg[R(z) \Rightarrow Q(y)]] \)

2. Consider the following model (where s is Sandy, k is Kim and l is Lee):

- like': \{\langle s, k \rangle, \langle l, k \rangle, \langle l, l \rangle, \langle s, s \rangle\}
- student': \{s, k\}
- beer-drinker': \{s, l\}

In this model, the following sentences are true (among others):

   - (a) like'(l, k)
     Lee likes Kim
   - (b) student'(s)
     Sandy is a student
   - (c) \( \neg\forall x[\text{student}'(x) \Rightarrow \text{beer-drinker}'(x)] \)
     It is not the case that every student is a beer drinker.

For each of the following, say whether it is true or false in the model, and give an English sentence that corresponds to the logical expression:

   - (a) \( \forall y[\forall x[\text{beer-drinker}'(x) \land \text{beer-drinker}'(y) \Rightarrow \text{like}'(x, y)]] \)
   - (b) \( \exists z[\text{student}'(z) \land \forall y[\text{beer-drinker}'(y) \Rightarrow \neg\text{like}'(z, y)]] \)

### 2.4 Translation into FOPC

English sentences can be (roughly) translated into FOPC (we will go over some ways this may be done automatically in the lectures, but for now we are just concerned with intuitively valid translations). For instance:

1. Every student drinks or smokes
   \( \forall x[\text{student}'(x) \Rightarrow (\text{drink}'(x) \lor \text{smoke}'(x))] \)

2. Kim likes every student
   \( \forall y[\text{student}'(y) \Rightarrow \text{like}'(k, y)] \)

3. Kim likes some students
   \( \exists z[\text{student}'(z) \land \text{like}'(k, z)] \)

Give FOPC equivalents for the following sentences, making sure to give all possible equivalents in the case of ambiguity:
1. Not every student drinks
2. Every student does not drink
3. Every student likes Kim and some student does not drink

2.5 Inference

Application of inference rules can be written with the premises above a horizontal line and the conclusions below it.

Give an English paraphrase of the following inferences:

\[ \forall x [\text{cat}'(x) \Rightarrow \text{animal}'(x)] \]

\[ \text{cat}'(k) \Rightarrow \text{animal}'(k) \]

\[ \text{cat}'(k) \Rightarrow \text{animal}'(k) \quad \text{cat}'(k) \]

\[ \text{animal}'(k) \]

3 Reading


A very good and quite gentle introduction geared towards the sort of logical concepts needed in computational linguistics.


This is a book with an associated software package — it’s not worth getting if you don’t intend to use the software (Windows and Macs, not Linux). It is a very good way of learning logic really thoroughly, though it is not really oriented towards linguistics. The first two-thirds of the book is the part that is relevant for the course.

The following two books go deeper and are recommended for the lectures, although there is no recommended book that covers the computational aspects of the course and no single book that covers all the formal semantics that will be discussed.

Ronnie Cann, Formal semantics, an introduction, Cambridge University Press, 1993

Recommended for coverage of logic, compositional semantics and lambda calculus. Covers a lot more formalism than we’ll need.

Kate Kearns, Semantics, Macmillan (Modern linguistics series), 2000

This covers formal aspects of semantics in a relatively gentle way, though not semantic composition. It contains much more about how semantics is used in linguistics than Cann’s book does, much of which is relevant to the course, at least as background.