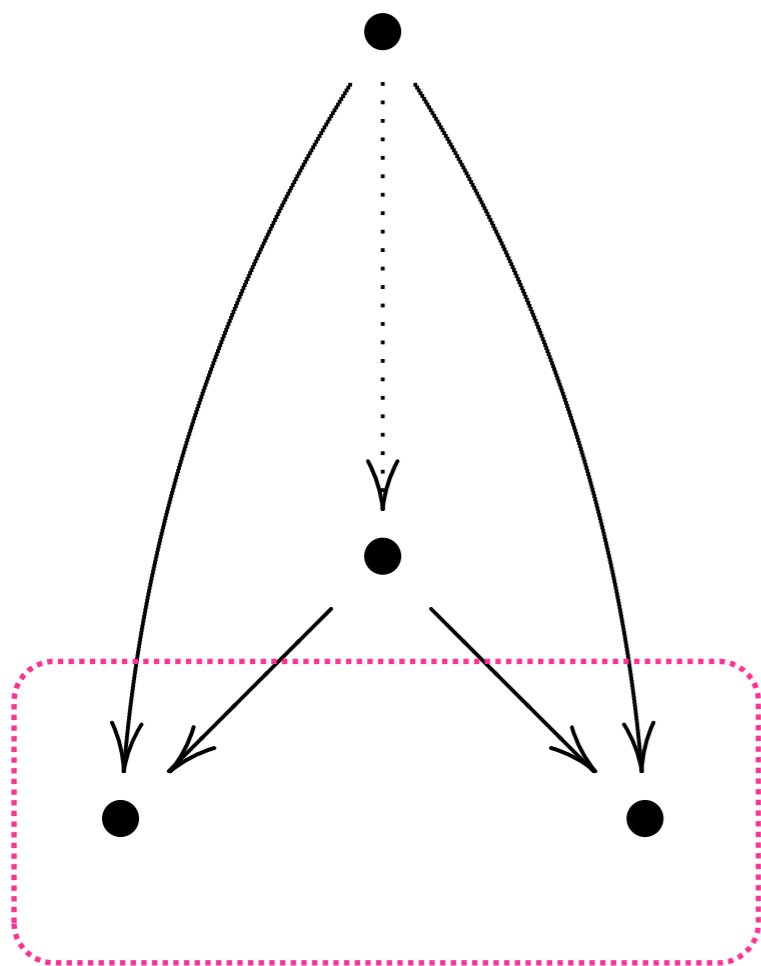
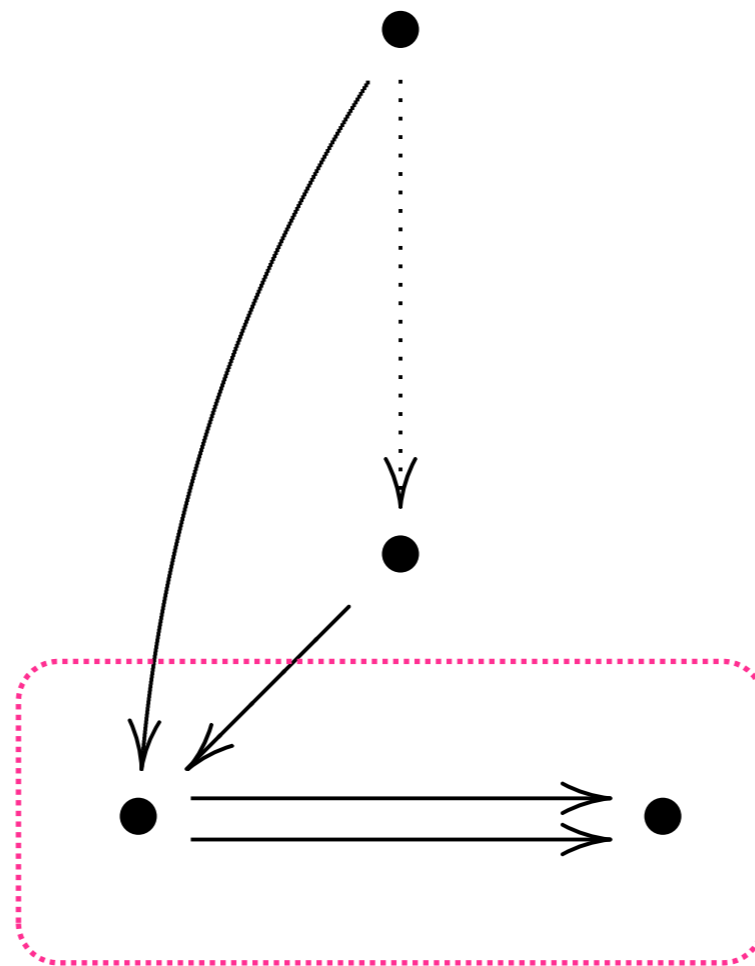


# A common theme

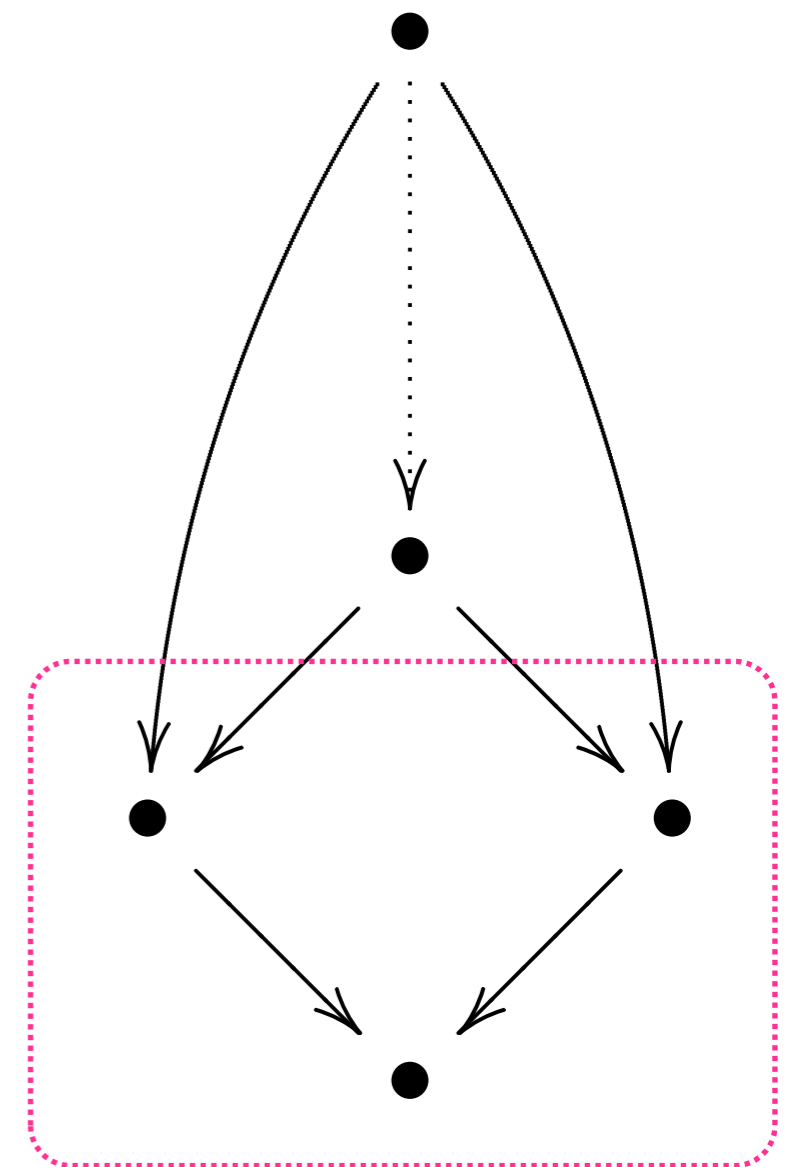
Product:



Equalizer:



Pullback:



# Diagrams

**Defn.** A **diagram**  $D$  in a category  $\mathbf{C}$  consists of:

- a graph  $G(D)$ ,
- an object  $D_v$  for each vertex  $v$  in  $G(D)$ ,
- an arrow  $D_e$  for each edge  $e$  in  $G(D)$ ,

s.t. if  $e : v \rightarrow w$  is an edge then

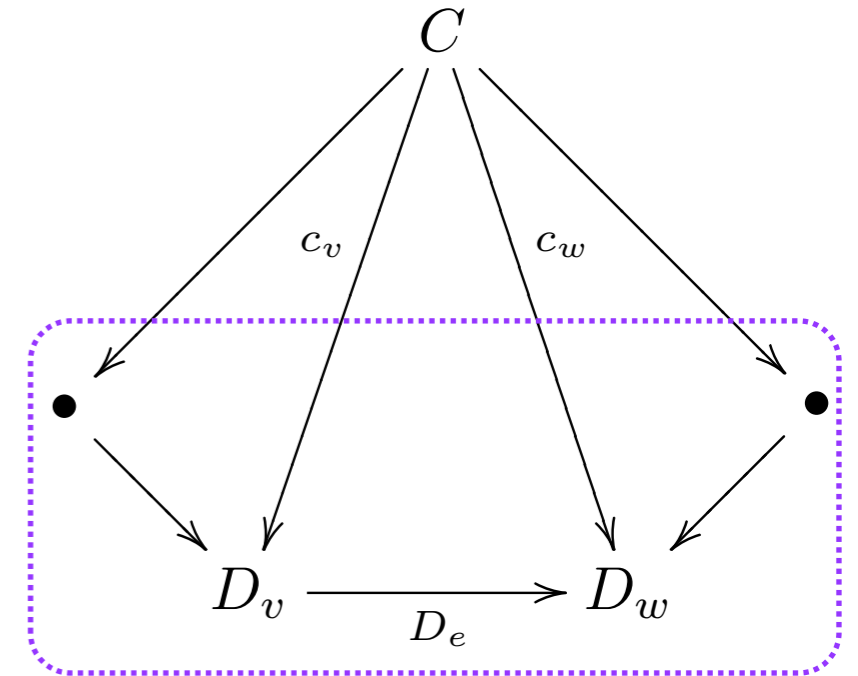
$$\text{dom}(D_e) = D_v \text{ and } \text{cod}(D_e) = D_w.$$

A diagram  $D$  is **commutative** (or **commutes**), if for any two paths in  $G(D)$  with common source and target, the compositions of the corresponding sequences of arrows are equal.

# Cones and cocones

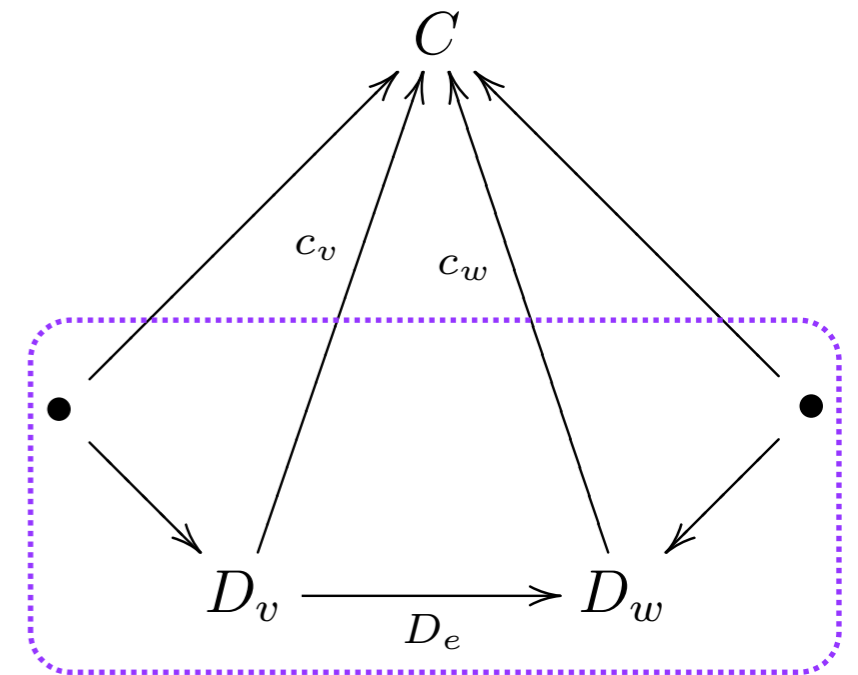
**Defn.** A **cone** for a diagram  $D$  in  $\mathbf{C}$  is:

- an object  $C$  in  $\mathbf{C}$ ,
  - an arrow  $c_v : C \rightarrow D_v$  for each vertex  $v$ ,
- s.t. for each edge  $e : v \rightarrow w$ ,  $D_e \circ c_v = c_w$ .



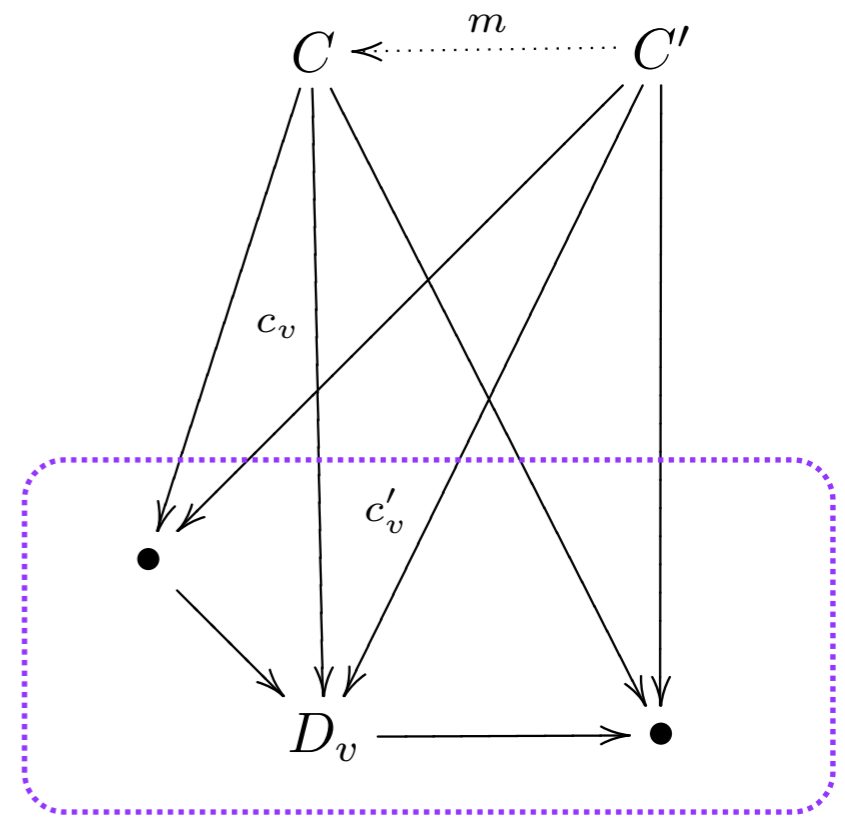
**Defn.** A **cocone** for a diagram  $D$  in  $\mathbf{C}$  is:

- an object  $C$  in  $\mathbf{C}$ ,
  - an arrow  $c_v : D_v \rightarrow C$  for each vertex  $v$ ,
- s.t. for each edge  $e : v \rightarrow w$ ,  $c_w \circ D_e = c_v$ .

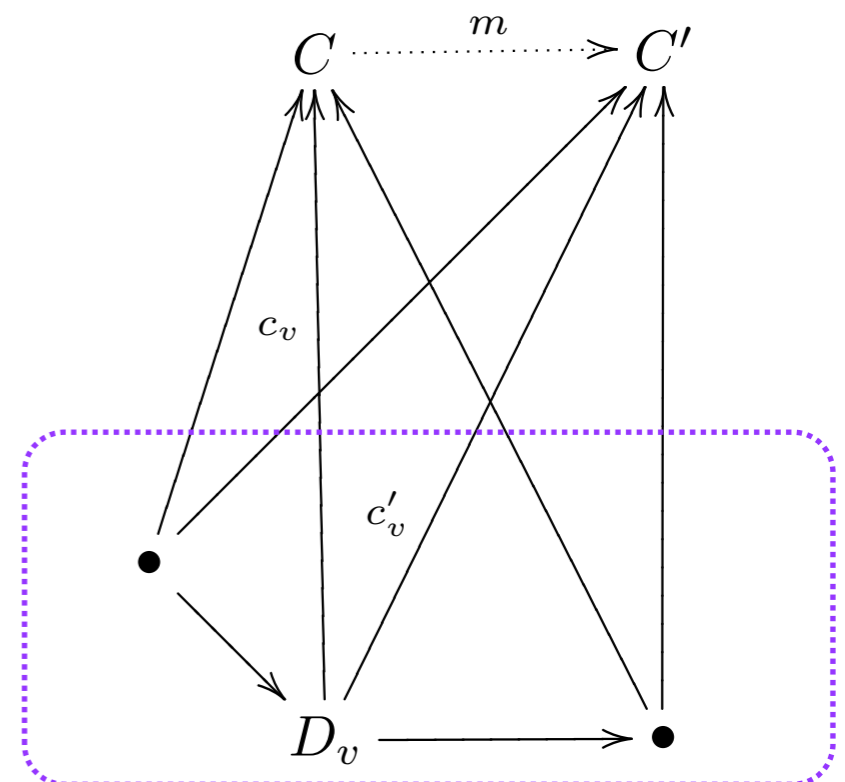


# Limits and colimits

**Defn.** A cone  $C$  is a **limit** for  $D$  if for any cone  $C'$  for  $D$  there is a *unique* arrow  $m : C' \rightarrow C$  s.t.  $c_v \circ m = c'_v$  for each vertex  $v$ .



**Defn.** A cocone  $C$  is a **colimit** for  $D$  if for any cocone  $C'$  for  $D$  there is a *unique* arrow  $m : C \rightarrow C'$  s.t.  $m \circ c_v = c'_v$  for each vertex  $v$ .



# Examples

Diagram shape	Limit	Colimit
$\emptyset$	Final object	Initial object
• •	Product	Coproduct
• $\begin{array}{c} \longrightarrow \\ \longrightarrow \end{array}$ •	Equalizer	Coequalizer
• $\longrightarrow$ • $\longleftarrow$ •	Pullback	
• $\longleftarrow$ • $\longrightarrow$ •		Pushout

# Some facts about (co)limits

**Fact.** (Co)Limits, if they exist, are unique up to isomorphism.

**Defn.** A category is (**finitely**) (**co**)complete if it has all (finite) (co)limits.

**Fact.** If a category has products of arbitrary families of objects, and equalizers, then it is complete.

**Fact.** If a category has a final object, all binary products and equalizers, then it is finitely complete.

**Note:** We only consider (co)limits of *small* diagrams.

**Fact.** A category with limits of all (large) diagrams is necessarily a preorder.

**Also:** A small complete category is necessarily a preorder.

# Limits and colimits in Sets.

**Fact.** Sets is complete and cocomplete.

**Limits:** For a diagram  $D$ , take the set

$$\{ \langle x_v, x_w, \dots \rangle \in D_v \times D_w \times \dots \mid \forall e : v \rightarrow w. D_e(x_v) = x_w \}$$

with obvious projections as the limiting cone.

**Colimits:** For a diagram  $D$ , take the set

$$(D_v + D_w + \dots) / \equiv$$

where  $\equiv$  is the least equivalence s.t.

$$D_e(x_v) \equiv x_v \text{ for all } e : v \rightarrow w, x_v \in D_v,$$

with obvious injections as the colimiting cocone.

limit = subset of product, colimit = quotient of sum