Complexity Theory
Easter 2010
Suggested Exercises 3

1. Show that a language L is in $\mathsf{co\text{-}NP}$ if, and only if, there is a nondeterministic Turing machine M and a polynomial p such that M halts in time $p(n)$ for all inputs of length x, and L is exactly the set of strings x such that *all* computations of M on input x end in an accepting state.

2. Define a *strong* nondeterministic Turing machine as one where each computation has three possible outcomes: accept, reject or maybe. If M is such a machine, we say that it accepts L, if for every $x \in L$, every computation path of M on x ends in either accept or maybe, with at least one accept *and* for $x \notin L$, every computation path of M on x ends in reject or maybe, with at least one reject.

   Show that if L is decided by a strong nondeterministic Turing machine running in polynomial time, then $L \in \mathsf{NP} \cap \mathsf{co\text{-}NP}$.

3. Consider the algorithm presented in the lecture which establishes that $\mathsf{Reachability}$ is in $\mathsf{SPACE}((\log n)^2)$. What is the time complexity of this algorithm?

   Can you generalise the time bound to the entire complexity class? That is, give a class of functions F, such that

   $$\mathsf{SPACE}((\log n)^2) \subseteq \bigcup_{f \in F} \mathsf{TIME}(f)$$

4. Show that, for every nondeterministic machine M which uses $O(\log n)$ work space, there is a machine R with three tapes (`input`, `work` and `output`) which works as follows. On input x, R produces on its output tape a description of the configuration graph for M, x, and R uses $O(\log |x|)$ space on its work tape.

   Explain why this means that if $\mathsf{Reachability}$ is in L, then $\mathsf{L} = \mathsf{NL}$.

5. Consider the language L in the alphabet $\{a, b\}$ given by $L = \{a^n b^n \mid n \in \mathbb{N}\}$. $L \notin \mathsf{SPACE}(c)$ for any constant c. Why?

6. On page 39 of the notes, a number of functions are listed as being constructible. Show that this is the case by giving, for each one, a description of an appropriate Turing machine.

   Prove that if f and g are constructible functions and $f(n) \geq n$, then so are $f(g)$, $f + g$, $f \cdot g$ and 2^f.

7. For any constructible function f, and any language $L \in \mathsf{NTIME}(f(n))$, there is a nondeterministic machine M that accepts L and all of whose computations terminate in time $O(f(n))$ for all inputs of length n. Give a detailed argument for this statement, describing how M might be obtained from a machine accepting L in time $f(n)$.

8. In the lecture, a proof of the Time Hierarchy Theorem was sketched. Give a similar argument for the following Space Hierarchy Theorem:

   Space Hierarchy. For every constructible function f, there is a language in $\mathsf{S}PACE(f(n) \cdot \log f(n))$ that is not in $\mathsf{S}PACE(f(n))$.

9. Show that, if $\mathsf{SPACE}((\log n)^2) \subseteq \mathsf{P}$, then $\mathsf{L} \neq \mathsf{P}$. (Hint: use the Space Hierarchy Theorem from Exercise 3 above.)

10. POLYLOGSPACE is the complexity class

$$\bigcup_k \mathsf{SPACE}((\log n)^k).$$

   (a) Show that, for any k, if $A \in \mathsf{SPACE}((\log n)^k)$ and $B \leq_L A$, then $B \in \mathsf{SPACE}((\log n)^k)$.

   (b) Show that there are no POLYLOGSPACE-complete problems with respect to $\leq_L$. (Hint: use (a) and the space hierarchy theorem).

   (c) Which of the following might be true: $\mathsf{P} \subseteq \mathsf{POLYLOGSPACE}$, $\mathsf{P} \supseteq \mathsf{POLYLOGSPACE}$, $\mathsf{P} = \mathsf{POLYLOGSPACE}$?