

Complex Analysis

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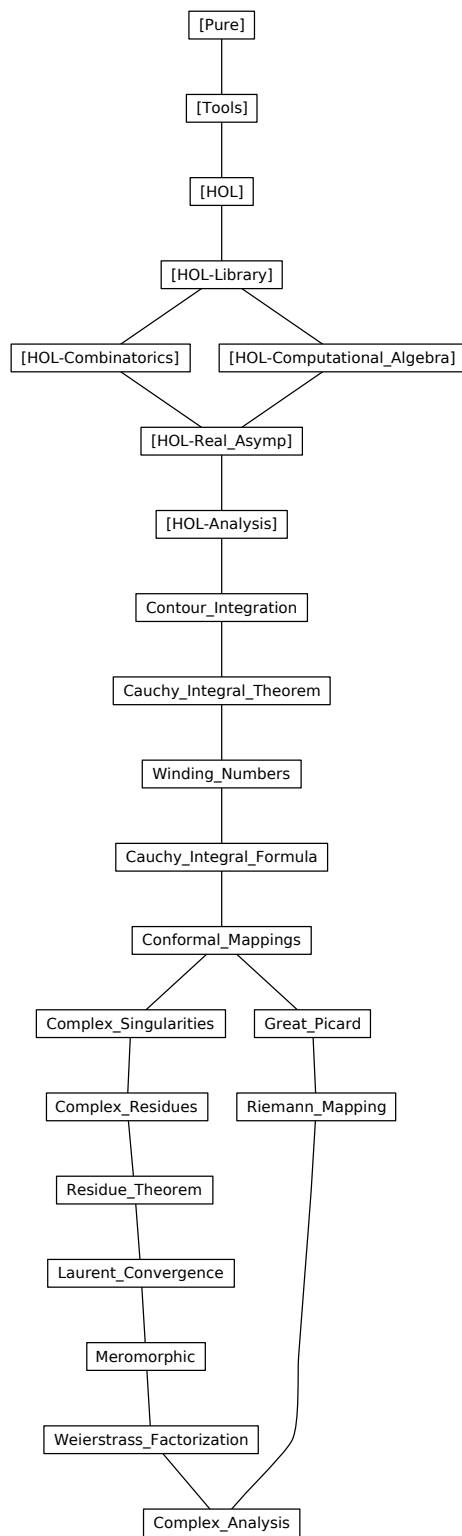
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1 Contour integration

theory *Contour_Integration*

imports *HOL-Analysis.Analysis*

begin

lemma *lhopital_complex_simple*:

assumes $(f \text{ has_field_derivative } f') \text{ (at } z)$

assumes $(g \text{ has_field_derivative } g') \text{ (at } z)$

assumes $f z = 0 \ g z = 0 \ g' \neq 0 \ f' / g' = c$

shows $((\lambda w. f w / g w) \longrightarrow c) \text{ (at } z)$

proof –

have *eventually* $(\lambda w. w \neq z) \text{ (at } z)$

by *(auto simp: eventually_at_filter)*

hence *eventually* $(\lambda w. ((f w - f z) / (w - z)) / ((g w - g z) / (w - z)) = f w / g w) \text{ (at } z)$

by *eventually_elim (simp add: assms field_split_simps)*

moreover **have** $((\lambda w. ((f w - f z) / (w - z)) / ((g w - g z) / (w - z))) \longrightarrow f' / g') \text{ (at } z)$

by *(intro tendsto_divide has_field_derivativeD assms)*

ultimately **have** $((\lambda w. f w / g w) \longrightarrow f' / g') \text{ (at } z)$

by *(blast intro: Lim_transform_eventually)*

with *assms* **show** *?thesis* **by** *simp*

qed

1.1 Definition

This definition is for complex numbers only, and does not generalise to line integrals in a vector field

definition *has_contour_integral* :: $(\text{complex} \Rightarrow \text{complex}) \Rightarrow \text{complex} \Rightarrow (\text{real} \Rightarrow \text{complex}) \Rightarrow \text{bool}$

(infixr $\langle \text{has}'_contour_integral \rangle$ 50)

where $(f \text{ has_contour_integral } i) \ g \equiv$

$((\lambda x. f(g x) * \text{vector_derivative } g \text{ (at } x \text{ within } \{0..1\}))$

$\text{has_integral } i) \ \{0..1\}$

definition *contour_integrable_on*

(infixr $\langle \text{contour}'_integrable_on \rangle$ 50)

where $f \text{ contour_integrable_on } g \equiv \exists i. (f \text{ has_contour_integral } i) \ g$

definition *contour_integral*

where $\text{contour_integral } g \ f \equiv \text{SOME } i. (f \text{ has_contour_integral } i) \ g \ \vee \ \neg f \text{ contour_integrable_on } g \ \wedge \ i=0$

lemma *not_integrable_contour_integral*: $\neg f \text{ contour_integrable_on } g \implies \text{contour_integral } g \ f = 0$

unfolding *contour_integrable_on_def* *contour_integral_def* **by** *blast*

lemma *contour_integral_unique*: $(f \text{ has_contour_integral } i) \ g \implies \text{contour_integral } g \ f$

$g f = i$
unfolding `contour_integral_def has_contour_integral_def contour_integrable_on_def`
using `has_integral_unique` **by** `blast`

lemma `has_contour_integral_cong`:
assumes $\bigwedge z. z \in \text{path_image } g \implies f z = f' z \text{ } g = g' \text{ } c = c'$
shows $(f \text{ has_contour_integral } c) \text{ } g \longleftrightarrow (f' \text{ has_contour_integral } c') \text{ } g'$
unfolding `has_contour_integral_def` `assms(2,3)`
by `(intro has_integral_cong) (auto simp: assms path_image_def intro!: assms(1))`

lemma `has_contour_integral_eqpath`:
 $\llbracket (f \text{ has_contour_integral } y) \text{ } p; f \text{ contour_integrable_on } \gamma;$
 $\text{contour_integral } p \text{ } f = \text{contour_integral } \gamma \text{ } f \rrbracket$
 $\implies (f \text{ has_contour_integral } y) \text{ } \gamma$
using `contour_integrable_on_def contour_integral_unique` **by** `auto`

lemma `has_contour_integral_integral`:
 $f \text{ contour_integrable_on } i \implies (f \text{ has_contour_integral } (\text{contour_integral } i \text{ } f))$
 i
by `(metis contour_integral_unique contour_integrable_on_def)`

lemma `has_contour_integral_unique`:
 $(f \text{ has_contour_integral } i) \text{ } g \implies (f \text{ has_contour_integral } j) \text{ } g \implies i = j$
using `has_integral_unique`
by `(auto simp: has_contour_integral_def)`

lemma `has_contour_integral_translate`:
 $(f \text{ has_contour_integral } I) ((+) z \circ g) \longleftrightarrow ((\lambda x. f (x + z)) \text{ has_contour_integral } I) \text{ } g$
by `(simp add: has_contour_integral_def add_ac)`

lemma `contour_integrable_translate`:
 $f \text{ contour_integrable_on } ((+) z \circ g) \longleftrightarrow (\lambda x. f (x + z)) \text{ contour_integrable_on } g$
by `(simp add: contour_integrable_on_def has_contour_integral_translate)`

lemma `contour_integral_translate`:
 $\text{contour_integral } ((+) z \circ g) \text{ } f = \text{contour_integral } g \text{ } (\lambda x. f (x + z))$
by `(simp add: contour_integral_def contour_integrable_translate has_contour_integral_translate)`

lemma `has_contour_integral_integrable`: $(f \text{ has_contour_integral } i) \text{ } g \implies f \text{ contour_integrable_on } g$
using `contour_integrable_on_def` **by** `blast`

Show that we can forget about the localized derivative.

lemma `has_integral_localized_vector_derivative`:
 $((\lambda x. f (g x) * \text{vector_derivative } p \text{ (at } x \text{ within } \{a..b\})) \text{ has_integral } i) \{a..b\}$
 \longleftrightarrow
 $((\lambda x. f (g x) * \text{vector_derivative } p \text{ (at } x)) \text{ has_integral } i) \{a..b\}$

proof –

have *: $\{a..b\} - \{a,b\} = \text{interior } \{a..b\}$
by (simp add: atLeastAtMost_diff_ends)
show ?thesis
by (rule has_integral_spike_eq [of $\{a,b\}$]) (auto simp: at_within_interior [of
 $\{a..b\}$])
qed

lemma integrable_on_localized_vector_derivative:

$(\lambda x. f (g x) * \text{vector_derivative } p \text{ (at } x \text{ within } \{a..b\})) \text{ integrable_on } \{a..b\}$
 \longleftrightarrow
 $(\lambda x. f (g x) * \text{vector_derivative } p \text{ (at } x)) \text{ integrable_on } \{a..b\}$
by (simp add: integrable_on_def has_integral_localized_vector_derivative)

lemma has_contour_integral:

$(f \text{ has_contour_integral } i) g \longleftrightarrow$
 $(\lambda x. f (g x) * \text{vector_derivative } g \text{ (at } x)) \text{ has_integral } i \{0..1\}$
by (simp add: has_integral_localized_vector_derivative has_contour_integral_def)

lemma contour_integrable_on:

$f \text{ contour_integrable_on } g \longleftrightarrow$
 $(\lambda t. f (g t) * \text{vector_derivative } g \text{ (at } t)) \text{ integrable_on } \{0..1\}$
by (simp add: has_contour_integral_integrable_on_def contour_integrable_on_def)

lemma has_contour_integral_mirror_iff:

assumes valid_path g
shows $(f \text{ has_contour_integral } I) (-g) \longleftrightarrow ((\lambda x. -f (-x)) \text{ has_contour_integral } I) g$

proof –

from assms **have** $g \text{ piecewise_differentiable_on } \{0..1\}$
by (auto simp: valid_path_def piecewise_C1_imp_differentiable)
then obtain S **where** finite S **and** $S: \bigwedge x. x \in \{0..1\} - S \implies g \text{ differentiable at } x \text{ within } \{0..1\}$

unfolding piecewise_differentiable_on_def **by** blast

have S' : $g \text{ differentiable at } x \text{ if } x \in \{0..1\} - (\{0, 1\} \cup S)$ **for** x

proof –

from that **have** $x \in \text{interior } \{0..1\}$ **by** auto

with $S[\text{of } x]$ that **show** ?thesis **by** (auto simp: at_within_interior[$\text{of } \{0..1\}$])

qed

have $(f \text{ has_contour_integral } I) (-g) \longleftrightarrow$
 $((\lambda x. f (-g x) * \text{vector_derivative } (-g) \text{ (at } x)) \text{ has_integral } I) \{0..1\}$
by (simp add: has_contour_integral)
also have ... $\longleftrightarrow ((\lambda x. -f (-g x) * \text{vector_derivative } g \text{ (at } x)) \text{ has_integral } I) \{0..1\}$

by (intro has_integral_spike_finite_eq[$\text{of } S \cup \{0, 1\}$])

(insert $\langle \text{finite } S \rangle S'$, auto simp: o_def fun_Compl_def)

also have ... $\longleftrightarrow ((\lambda x. -f (-x)) \text{ has_contour_integral } I) g$

by (simp add: has_contour_integral)

finally show ?thesis .

qed

lemma *contour_integral_on_mirror_iff*:

assumes *valid_path g*

shows $f \text{ contour_integrable_on } (-g) \longleftrightarrow (\lambda x. -f (-x)) \text{ contour_integrable_on } g$

by (*auto simp: contour_integrable_on_def has_contour_integral_mirror_iff assms*)

lemma *contour_integral_mirror*:

assumes *valid_path g*

shows $\text{contour_integral } (-g) f = \text{contour_integral } g (\lambda x. -f (-x))$

proof (*cases f contour_integrable_on (-g)*)

case *True with contour_integral_unique assms show ?thesis*

by (*auto simp: contour_integrable_on_def has_contour_integral_mirror_iff*)

next

case *False then show ?thesis*

by (*simp add: assms contour_integral_on_mirror_iff not_integrable_contour_integral*)

qed

1.2 Reversing a path

lemma *has_contour_integral_reversepath*:

assumes *valid_path g and f: (f has_contour_integral i) g*

shows $(f \text{ has_contour_integral } (-i)) (\text{reversepath } g)$

proof –

{ **fix** *S x*

assume *xs: g C1_differentiable_on ({0..1} - S) x \notin (-) 1 ' S 0 ≤ x x ≤ 1*

have $\text{vector_derivative } (\lambda x. g (1 - x)) (\text{at } x \text{ within } \{0..1\}) =$

$-\text{vector_derivative } g (\text{at } (1 - x) \text{ within } \{0..1\})$

proof –

obtain *f'* **where** $f': (g \text{ has_vector_derivative } f') (\text{at } (1 - x))$

using *xs*

by (*force simp: has_vector_derivative_def C1_differentiable_on_def*)

have $(g \circ (\lambda x. 1 - x)) \text{ has_vector_derivative } -1 *_{\mathbb{R}} f' (\text{at } x)$

by (*intro vector_diff_chain_within has_vector_derivative_at_within [OF f'] derivative_eq_intros | simp*)+

then have $mf': ((\lambda x. g (1 - x)) \text{ has_vector_derivative } -f') (\text{at } x)$

by (*simp add: o_def*)

show *?thesis*

using *xs*

by (*auto simp: vector_derivative_at_within_ivl [OF mf'] vector_derivative_at_within_ivl [OF f']*)

qed

} **note** $* = \text{this}$

obtain *S* **where** *S: continuous_on {0..1} g finite S g C1_differentiable_on {0..1} - S*

using *assms* **by** (*auto simp: valid_path_def piecewise_C1_differentiable_on_def*)

have $((\lambda x. - (f (g (1 - x)) * \text{vector_derivative } g (\text{at } (1 - x) \text{ within } \{0..1\}))) \text{ has_integral } -i)$

```

    {0..1}
    using has_integral_affinity01 [where m= -1 and c=1, OF f [unfolded
has_contour_integral_def]]
    by (simp add: has_integral_neg)
    then show ?thesis
    using S
    unfolding reversepath_def has_contour_integral_def
    by (rule_tac S = ( $\lambda x. 1 - x$ ) ' S in has_integral_spike_finite) (auto simp: *)
qed

```

```

lemma contour_integrable_reversepath:
  valid_path g  $\implies$  f contour_integrable_on g  $\implies$  f contour_integrable_on
(reversepath g)
  using has_contour_integral_reversepath contour_integrable_on_def by blast

```

```

lemma contour_integrable_reversepath_eq:
  valid_path g  $\implies$  (f contour_integrable_on (reversepath g)  $\longleftrightarrow$  f contour_integrable_on
g)
  using contour_integrable_reversepath valid_path_reversepath by fastforce

```

```

lemma contour_integral_reversepath:
  assumes valid_path g
  shows contour_integral (reversepath g) f = - (contour_integral g f)
proof (cases f contour_integrable_on g)
  case True then show ?thesis
    by (simp add: assms contour_integral_unique has_contour_integral_integral
has_contour_integral_reversepath)
  next
  case False then have  $\neg$  f contour_integrable_on (reversepath g)
    by (simp add: assms contour_integrable_reversepath_eq)
  with False show ?thesis by (simp add: not_integrable_contour_integral)
qed

```

1.3 Joining two paths together

```

lemma has_contour_integral_join:
  assumes (f has_contour_integral i1) g1 (f has_contour_integral i2) g2
  valid_path g1 valid_path g2
  shows (f has_contour_integral (i1 + i2)) (g1 +++ g2)
proof -
  obtain s1 s2
  where s1: finite s1  $\forall x \in \{0..1\} - s1. g1$  differentiable at x
  and s2: finite s2  $\forall x \in \{0..1\} - s2. g2$  differentiable at x
  using assms
  by (auto simp: valid_path_def piecewise_C1_differentiable_on_def C1_differentiable_on_eq)
  have 1: (( $\lambda x. f (g1 x) * \text{vector\_derivative } g1 \text{ (at } x)$ ) has_integral i1) {0..1}
  and 2: (( $\lambda x. f (g2 x) * \text{vector\_derivative } g2 \text{ (at } x)$ ) has_integral i2) {0..1}
  using assms
  by (auto simp: has_contour_integral)

```

```

have i1: (( $\lambda x. (2 * f (g1 (2 * x))) * \text{vector\_derivative } g1 \text{ (at (2 * x))}$ ) has_integral
i1) {0..1/2}
  and i2: (( $\lambda x. (2 * f (g2 (2 * x - 1))) * \text{vector\_derivative } g2 \text{ (at (2 * x - 1))}$ )
has_integral i2) {1/2..1}
  using has_integral_affinity01 [OF 1, where  $m=2$  and  $c=0$ , THEN has_integral_cmul
[where  $c=2$ ]]
    has_integral_affinity01 [OF 2, where  $m=2$  and  $c=-1$ , THEN has_integral_cmul
[where  $c=2$ ]]
  by (simp_all only: image_affinity_atLeastAtMost_div_diff, simp_all add:
scaleR_conv_of_real mult_ac)
  have g1: vector_derivative ( $\lambda x. \text{if } x^2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)$ )
(at z) =
  2 *R vector_derivative g1 (at (z*2))
  if  $0 \leq z \wedge z^2 < 1 \wedge z^2 \notin s1$  for z
proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
  show  $0 < |z - 1/2|$ 
  using that by auto
  have ((*) 2 has_vector_derivative 2) (at z)
  by (simp add: has_vector_derivative_def has_derivative_def bounded_linear_mult_left)
  moreover have (g1 has_vector_derivative vector_derivative g1 (at (z * 2)))
(at (2 * z))
  using s1 that by (auto simp: algebra_simps vector_derivative_works)
  ultimately
  show (( $\lambda x. g1 (2 * x)$ ) has_vector_derivative 2 *R vector_derivative g1 (at
(z * 2))) (at z)
  by (intro vector_diff_chain_at [simplified o_def])
  qed (use that in <simp_all add: dist_real_def abs_if_split: if_split_asm>)

  have g2: vector_derivative ( $\lambda x. \text{if } x^2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)$ )
(at z) =
  2 *R vector_derivative g2 (at (z*2 - 1))
  if  $1 < z^2 \wedge z \leq 1 \wedge z^2 - 1 \notin s2$  for z
proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
  show  $0 < |z - 1/2|$ 
  using that by auto
  have (( $\lambda x. 2 * x - 1$ ) has_vector_derivative 2) (at z)
  by (simp add: has_vector_derivative_def has_derivative_def bounded_linear_mult_left)
  moreover have (g2 has_vector_derivative vector_derivative g2 (at (z * 2 -
1))) (at (2 * z - 1))
  using s2 that by (auto simp: algebra_simps vector_derivative_works)
  ultimately
  show (( $\lambda x. g2 (2 * x - 1)$ ) has_vector_derivative 2 *R vector_derivative g2
(at (z * 2 - 1))) (at z)
  by (intro vector_diff_chain_at [simplified o_def])
  qed (use that in <simp_all add: dist_real_def abs_if_split: if_split_asm>)

have (( $\lambda x. f ((g1 +++ g2) x) * \text{vector\_derivative } (g1 +++ g2) \text{ (at } x)$ ) has_integral
i1) {0..1/2}
proof (rule has_integral_spike_finite [OF __ i1])

```

```

  show finite (insert (1/2) ((*) 2 - ' s1))
    using s1 by (force intro: finite_vimageI [where h = (*)2] inj_onI)
  qed (auto simp add: joinpaths_def scaleR_conv_of_real mult_ac g1)
  moreover have (( $\lambda x. f ((g1 +++ g2) x) * vector\_derivative (g1 +++ g2) (at x)$ ) has_integral i2) {1/2..1}
  proof (rule has_integral_spike_finite [OF _ _ i2])
    show finite (insert (1/2) (( $\lambda x. 2 * x - 1$ ) - ' s2))
      using s2 by (force intro: finite_vimageI [where h =  $\lambda x. 2*x-1$ ] inj_onI)
    qed (auto simp add: joinpaths_def scaleR_conv_of_real mult_ac g2)
  ultimately
  show ?thesis
    by (simp add: has_contour_integral has_integral_combine [where c = 1/2])
  qed

```

lemma contour_integrable_joinI:

```

  assumes f contour_integrable_on g1 f contour_integrable_on g2
    valid_path g1 valid_path g2
  shows f contour_integrable_on (g1 +++ g2)
  using assms
  by (meson has_contour_integral_join contour_integrable_on_def)

```

lemma contour_integrable_joinD1:

```

  assumes f contour_integrable_on (g1 +++ g2) valid_path g1
  shows f contour_integrable_on g1
  proof -
    obtain s1
      where s1: finite s1  $\forall x \in \{0..1\} - s1. g1$  differentiable at x
    using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def
      C1_differentiable_on_eq)
    have ( $\lambda x. f ((g1 +++ g2) (x/2)) * vector\_derivative (g1 +++ g2) (at (x/2))$ )
      integrable_on {0..1}
      using assms integrable_affinity [of _ 0 1/2 1/2 0] integrable_on_subbox
      [where a=0 and b=1/2]
    then have *: ( $\lambda x. (f ((g1 +++ g2) (x/2))/2) * vector\_derivative (g1 +++ g2) (at (x/2))$ )
      integrable_on {0..1}
    by (auto dest: integrable_cmul [where c=1/2] simp: scaleR_conv_of_real)
    have g1: vector_derivative ( $\lambda x. if x*2 \le 1$  then  $g1 (2*x)$  else  $g2 (2*x - 1)$ )
      (at (z/2)) =
      2 *_R vector_derivative g1 (at z)
    if 0 < z z < 1 z  $\notin$  s1 for z
  proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
    show 0 < |(z - 1)/2|
      using that by auto
    have  $\S$ : (( $\lambda x. x * 2$ ) has_vector_derivative 2) (at (z/2))
      using s1 by (auto simp: has_vector_derivative_def has_derivative_def
        bounded_linear_mult_left)
    have (g1 has_vector_derivative vector_derivative g1 (at z)) (at z)
      using s1 that by (auto simp: vector_derivative_works)

```

```

then show (( $\lambda x. g1 (2 * x)$ ) has_vector_derivative 2 *R vector_derivative g1
(at z)) (at (z/2))
using vector_diff_chain_at [OF §] by (auto simp: field_simps o_def)
qed (use that in <simp_all add: field_simps dist_real_def abs_if_split: if_split_asm>)
have fin01: finite ({0, 1}  $\cup$  s1)
by (simp add: s1)
show ?thesis
unfolding contour_integrable_on
by (intro integrable_spike_finite [OF fin01 _ *]) (auto simp: joinpaths_def
scaleR_conv_of_real g1)
qed

```

lemma *contour_integrable_joinD2*:

```

assumes f contour_integrable_on (g1 +++ g2) valid_path g2
shows f contour_integrable_on g2
proof -
obtain s2
where s2: finite s2  $\forall x \in \{0..1\} - s2. g2$  differentiable at x
using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def
C1_differentiable_on_eq)
have ( $\lambda x. f ((g1 +++ g2) (x/2 + 1/2)) * \text{vector\_derivative} (g1 +++ g2) (at$ 
( $x/2 + 1/2$ ))) integrable_on {0..1}
using assms integrable_affinity [of _ 1/2::real 1 1/2 1/2]
integrable_on_subbox [where a=1/2 and b=1]
by (fastforce simp: contour_integrable_on image_affinity_atLeastAtMost_diff)
then have *: ( $\lambda x. (f ((g1 +++ g2) (x/2 + 1/2))/2) * \text{vector\_derivative} (g1$ 
+++ g2) (at ( $x/2 + 1/2$ )))
integrable_on {0..1}
by (auto dest: integrable_cmul [where c=1/2] simp: scaleR_conv_of_real)
have g2: vector_derivative ( $\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)$ )
(at ( $z/2 + 1/2$ )) =
2 *R vector_derivative g2 (at z)
if  $0 < z < 1$   $z \notin s2$  for z
proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
show  $0 < |z/2|$ 
using that by auto
have §: (( $\lambda x. x * 2 - 1$ ) has_vector_derivative 2) (at ((1 + z)/2))
using s2 by (auto simp: has_vector_derivative_def has_derivative_def
bounded_linear_mult_left)
have (g2 has_vector_derivative vector_derivative g2 (at z)) (at z)
using s2 that by (auto simp: vector_derivative_works)
then show (( $\lambda x. g2 (2 * x - 1)$ ) has_vector_derivative 2 *R vector_derivative
g2 (at z)) (at ( $z/2 + 1/2$ ))
using vector_diff_chain_at [OF §] by (auto simp: field_simps o_def)
qed (use that in <simp_all add: field_simps dist_real_def abs_if_split: if_split_asm>)
have fin01: finite ({0, 1}  $\cup$  s2)
by (simp add: s2)
show ?thesis
unfolding contour_integrable_on

```

by (intro integrable_spike_finite [OF fin01 _ *]) (auto simp: joinpaths_def scaleR_conv_of_real g2)
qed

lemma *contour_integrable_join* [simp]:
 $\llbracket \text{valid_path } g1; \text{valid_path } g2 \rrbracket$
 $\implies f \text{ contour_integrable_on } (g1 \text{ +++ } g2) \longleftrightarrow f \text{ contour_integrable_on } g1$
 $\wedge f \text{ contour_integrable_on } g2$
using *contour_integrable_joinD1 contour_integrable_joinD2 contour_integrable_joinI*
by *blast*

lemma *contour_integral_join* [simp]:
 $\llbracket f \text{ contour_integrable_on } g1; f \text{ contour_integrable_on } g2; \text{valid_path } g1; \text{valid_path } g2 \rrbracket$
 $\implies \text{contour_integral } (g1 \text{ +++ } g2) f = \text{contour_integral } g1 f + \text{contour_integral } g2 f$
by (*simp add: has_contour_integral_integral has_contour_integral_join contour_integral_unique*)

1.4 Shifting the starting point of a (closed) path

lemma *has_contour_integral_shiftpath*:
assumes *f: (f has_contour_integral i) g valid_path g*
and *a: a ∈ {0..1}*
shows *(f has_contour_integral i) (shiftpath a g)*
proof –
obtain *S*
where *S: finite S and g: $\forall x \in \{0..1\} - S. g$ differentiable at x*
using *assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def C1_differentiable_on_eq)*
have ***: $((\lambda x. f (g x) * \text{vector_derivative } g \text{ (at } x)) \text{ has_integral } i) \{0..1\}$
using *assms by (auto simp: has_contour_integral)*
then have *i*: $i = \text{integral } \{a..1\} (\lambda x. f (g x) * \text{vector_derivative } g \text{ (at } x)) + \text{integral } \{0..a\} (\lambda x. f (g x) * \text{vector_derivative } g \text{ (at } x))$
by (*smt (verit, ccfv_threshold) Henstock_Kurzweil_Integration.integral_combine a add.commute atLeastAtMost_iff has_integral_iff*)
have *vd1*: $\text{vector_derivative } (\text{shiftpath } a g) \text{ (at } x) = \text{vector_derivative } g \text{ (at } (x + a))$
if $0 \leq x + a < 1 \ x \notin (\lambda x. x - a) ' S$ **for** *x*
unfolding *shiftpath_def*
proof (*rule vector_derivative_at [OF has_vector_derivative_transform_within]*)
have $((\lambda x. g (x + a)) \text{ has_vector_derivative } \text{vector_derivative } g \text{ (at } (a + x))) \text{ (at } x)$
proof (*rule vector_diff_chain_at [of _ 1, simplified o_def scaleR_one]*)
show $((\lambda x. x + a) \text{ has_vector_derivative } 1) \text{ (at } x)$
by (*rule derivative_eq_intros | simp*)
have *g* differentiable at $(x + a)$
using *g a that by force*
then show $(g \text{ has_vector_derivative } \text{vector_derivative } g \text{ (at } (a + x))) \text{ (at } (x$

```

+ a))
  by (metis add.commute vector_derivative_works)
qed
then show (( $\lambda x. g (a + x)$ ) has_vector_derivative vector_derivative g (at (x
+ a))) (at x)
  by (auto simp: field_simps)
show 0 < dist (1 - a) x
  using that by auto
qed (use that in <auto simp: dist_real_def>)

have vd2: vector_derivative (shiftpath a g) (at x) = vector_derivative g (at (x
+ a - 1))
  if  $x \leq 1$   $1 < x + a$   $x \notin (\lambda x. x - a + 1) \text{ ' } S$  for x
  unfolding shiftpath_def
proof (rule vector_derivative_at [OF has_vector_derivative_transform_within])
  have (( $\lambda x. g (x + a - 1)$ ) has_vector_derivative vector_derivative g (at
(a+x-1))) (at x)
  proof (rule vector_diff_chain_at [of _ 1, simplified o_def scaleR_one])
    show (( $\lambda x. x + a - 1$ ) has_vector_derivative 1) (at x)
      by (rule derivative_eq_intros | simp)+
    have g differentiable at (x+a-1)
      using g a that by force
    then show (g has_vector_derivative vector_derivative g (at (a+x-1))) (at
(x + a - 1))
      by (metis add.commute vector_derivative_works)
  qed
  then show (( $\lambda x. g (a + x - 1)$ ) has_vector_derivative vector_derivative g (at
(x + a - 1))) (at x)
    by (auto simp: field_simps)
  show 0 < dist (1 - a) x
    using that by auto
  qed (use that in <auto simp: dist_real_def>)

have va1: ( $\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)$ ) integrable_on ({a..1})
  using * a by (fastforce intro: integrable_subinterval_real)
have v0a: ( $\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)$ ) integrable_on ({0..a})
  using * a by (force intro: integrable_subinterval_real)
have finite ({1 - a}  $\cup (\lambda x. x - a) \text{ ' } S$ )
  using S by blast
then have (( $\lambda x. f (\text{shiftpath } a \text{ } g \text{ } x) * \text{vector\_derivative } (\text{shiftpath } a \text{ } g) \text{ (at } x)$ )
has_integral integral {a..1} ( $\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x)$ )) {0..1
- a}
  apply (rule has_integral_spike_finite
[where f =  $\lambda x. f (g(a+x)) * \text{vector\_derivative } g \text{ (at}(a+x))$ ])
  subgoal
  using a by (simp add: vd1) (force simp: shiftpath_def add.commute)
  subgoal
  using has_integral_affinity [where m=1 and c=a] integrable_integral [OF
va1]

```

```

    by (force simp add: add.commute)
  done
  moreover
  have finite ({1 - a}  $\cup$  ( $\lambda x. x - a + 1$ ) ' S)
    using S by blast
  then have (( $\lambda x. f$  (shiftpath a g x) * vector_derivative (shiftpath a g) (at x))
    has_integral integral {0..a} ( $\lambda x. f$  (g x) * vector_derivative g (at x)))
    {1 - a..1}
  apply (rule has_integral_spike_finite
    [where f =  $\lambda x. f$ (g(a+x-1)) * vector_derivative g (at(a+x-1))])
  subgoal
  using a by (simp add: vd2) (force simp: shiftpath_def add.commute)
  subgoal
  using has_integral_affinity [where m=1 and c=a-1, simplified, OF integrable_integral [OF v0a]]
  by (force simp add: algebra_simps)
  done
  ultimately show ?thesis
  using a
  by (auto simp: i has_contour_integral intro: has_integral_combine [where c = 1-a])
qed

```

lemma has_contour_integral_shiftpath_D:

```

  assumes (f has_contour_integral i) (shiftpath a g)
    valid_path g pathfinish g = pathstart g a  $\in$  {0..1}
  shows (f has_contour_integral i) g
proof -
  obtain S
  where S: finite S and g:  $\forall x \in \{0..1\} - S. g$  differentiable at x
  using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def
    C1_differentiable_on_eq)
  { fix x
  assume x: 0 < x < 1 x  $\notin$  S
  then have gx: g differentiable at x
    using g by auto
  have  $\S$ : shiftpath (1 - a) (shiftpath a g) differentiable at x
    using assms x
    by (intro differentiable_transform_within [OF gx, of min x (1-x)])
      (auto simp: dist_real_def shiftpath_shiftpath abs_if_split: if_split_asm)
  have vector_derivative g (at x within {0..1}) =
    vector_derivative (shiftpath (1 - a) (shiftpath a g)) (at x within {0..1})
  apply (rule vector_derivative_at_within_ivl
    [OF has_vector_derivative_transform_within_open
      [where f = (shiftpath (1 - a) (shiftpath a g)) and S =
{0<.. $\dots$ <1}-S]])
  using S assms x  $\S$ 
  apply (auto simp: finite_imp_closed open_Diff shiftpath_shiftpath
    at_within_interior [of _ {0..1}]) vector_derivative_works

```



```

[symmetric]
  done
} note vd = this
have fi: (f has_contour_integral i) (shiftpath (1 - a) (shiftpath a g))
  using assms by (auto intro!: has_contour_integral_shiftpath)
show ?thesis
  unfolding has_contour_integral_def
  proof (rule has_integral_spike_finite [of {0,1} ∪ S, OF _ _ fi [unfolded
has_contour_integral_def]])
    show finite ( $\{0, 1\} \cup S$ )
    by (simp add: S)
  qed (use S assms vd in <auto simp: shiftpath_shiftpath>)
qed

```

```

lemma has_contour_integral_shiftpath_eq:
  assumes valid_path g pathfinish g = pathstart g a ∈ {0..1}
  shows (f has_contour_integral i) (shiftpath a g)  $\longleftrightarrow$  (f has_contour_integral
i) g
  using assms has_contour_integral_shiftpath has_contour_integral_shiftpath_D
by blast

```

```

lemma contour_integrable_on_shiftpath_eq:
  assumes valid_path g pathfinish g = pathstart g a ∈ {0..1}
  shows f contour_integrable_on (shiftpath a g)  $\longleftrightarrow$  f contour_integrable_on g
  using assms contour_integrable_on_def has_contour_integral_shiftpath_eq by
auto

```

```

lemma contour_integral_shiftpath:
  assumes valid_path g pathfinish g = pathstart g a ∈ {0..1}
  shows contour_integral (shiftpath a g) f = contour_integral g f
  using assms
by (simp add: contour_integral_def contour_integrable_on_def has_contour_integral_shiftpath_eq)

```

1.5 More about straight-line paths

```

lemma has_contour_integral_linepath:
  shows (f has_contour_integral i) (linepath a b)  $\longleftrightarrow$ 
  (( $\lambda x. f(\text{linepath } a \ b \ x) * (b - a)$ ) has_integral i)  $\{0..1\}$ 
by (simp add: has_contour_integral)

```

```

lemma has_contour_integral_trivial [iff]: (f has_contour_integral 0) (linepath a
a)
by (simp add: has_contour_integral_linepath)

```

```

lemma has_contour_integral_trivial_iff [simp]: (f has_contour_integral i) (linepath
a a)  $\longleftrightarrow$  i=0
using has_contour_integral_unique by blast

```

```

lemma contour_integral_trivial [simp]: contour_integral (linepath a a) f = 0

```

using *has_contour_integral_trivial* *contour_integral_unique* by *blast*

1.6 Relation to subpath construction

lemma *has_contour_integral_subpath_refl* [*iff*]: (*f* *has_contour_integral* 0) (*subpath* *u* *u* *g*)
 by (*simp* *add*: *has_contour_integral_subpath_def*)

lemma *contour_integrable_subpath_refl* [*iff*]: *f* *contour_integrable_on* (*subpath* *u* *u* *g*)
 using *has_contour_integral_subpath_refl* *contour_integrable_on_def* by *blast*

lemma *contour_integral_subpath_refl* [*simp*]: *contour_integral* (*subpath* *u* *u* *g*) *f* = 0
 by (*simp* *add*: *contour_integral_unique*)

lemma *has_contour_integral_subpath*:
assumes *f*: *f* *contour_integrable_on* *g* **and** *g*: *valid_path* *g*
and *uv*: *u* ∈ {0..1} *v* ∈ {0..1} *u* ≤ *v*
shows (*f* *has_contour_integral* *integral* {*u..v*} (λ*x*. *f*(*g* *x*) * *vector_derivative* *g* (at *x*)))
 (*subpath* *u* *v* *g*)
proof (*cases* *v=u*)
case *True*
then show ?*thesis*
 using *f* by (*simp* *add*: *contour_integrable_on_def* *subpath_def* *has_contour_integral*)
next
case *False*
obtain *S* **where** *S*: ∧*x*. *x* ∈ {0..1} - *S* ⇒ *g* *differentiable* at *x* **and** *fs*: *finite* *S*
 using *g* **unfolding** *piecewise_C1_differentiable_on_def* *C1_differentiable_on_eq* *valid_path_def* by *blast*
have §: (λ*t*. *f* (*g* *t*) * *vector_derivative* *g* (at *t*)) *integrable_on* {*u..v*}
 using *contour_integrable_on* *f* *integrable_on_subinterval* *uv* by *fastforce*
then have *: ((λ*x*. *f* (*g* ((*v* - *u*) * *x* + *u*)) * *vector_derivative* *g* (at ((*v* - *u*) * *x* + *u*)))
 (*has_integral* (1 / (*v* - *u*)) * *integral* {*u..v*} (λ*t*. *f* (*g* *t*) * *vector_derivative* *g* (at *t*)))
 {0..1}
 using *uv* *False* **unfolding** *has_integral_integral*
apply *simp*
apply (*drule* *has_integral_affinity* [**where** *m=v-u* **and** *c=u*, *simplified*])
apply (*simp_all* *add*: *image_affinity_atLeastAtMost_div_diff* *scaleR_conv_of_real*)
apply (*simp* *add*: *divide_simps*)
done

have *vd*: *vector_derivative* (λ*x*. *g* ((*v*-*u*) * *x* + *u*)) (at *x*) = (*v*-*u*) *_R *vector_derivative* *g* (at ((*v*-*u*) * *x* + *u*))
if *x* ∈ {0..1} *x* ∉ (λ*t*. (*v*-*u*) *_R *t* + *u*) - ' *S* **for** *x*

```

proof (rule vector_derivative_at [OF vector_diff_chain_at [simplified o_def]])
  show (( $\lambda x. (v - u) * x + u$ ) has_vector_derivative  $v - u$ ) (at  $x$ )
    by (intro derivative_eq_intros | simp)+
qed (use  $S$  uv mult_left_le [of  $x$   $v-u$ ] that in ‹auto simp: vector_derivative_works›)

have fin: finite (( $\lambda t. (v - u) *_{\mathbb{R}} t + u$ ) - ‘ $S$ )
  using fs by (auto simp: inj_on_def False finite_vimageI)
show ?thesis
  unfolding subpath_def has_contour_integral
  apply (rule has_integral_spike_finite [OF fin])
  using has_integral_cmul [OF *, where  $c = v-u$ ] fs assms
  by (auto simp: False vd scaleR_conv_of_real)
qed

```

lemma contour_integrable_subpath:

```

assumes f contour_integrable_on g valid_path g  $u \in \{0..1\}$   $v \in \{0..1\}$ 
  shows f contour_integrable_on (subpath  $u$   $v$  g)
by (smt (verit, ccfv_threshold) assms contour_integrable_on_def contour_integrable_reversepath_eq
  has_contour_integral_subpath reversepath_subpath valid_path_subpath)

```

lemma has_integral_contour_integral_subpath:

```

assumes f contour_integrable_on g valid_path g  $u \in \{0..1\}$   $v \in \{0..1\}$   $u \leq v$ 
  shows (( $\lambda x. f(g\ x) * \text{vector\_derivative } g \text{ (at } x)$ )
    has_integral_contour_integral (subpath  $u$   $v$  g)  $f$ ) { $u..v$ }
    (is (?fg has_integral _))

```

proof –

```

  have (?fg has_integral integral { $u..v$ } ?fg) { $u..v$ }
    using assms contour_integrable_on integrable_on_subinterval by fastforce
  then show ?thesis
    by (metis (full_types) assms contour_integral_unique has_contour_integral_subpath)
qed

```

lemma contour_integral_subcontour_integral:

```

assumes f contour_integrable_on g valid_path g  $u \in \{0..1\}$   $v \in \{0..1\}$   $u \leq v$ 
  shows contour_integral (subpath  $u$   $v$  g)  $f =$ 
    integral { $u..v$ } ( $\lambda x. f(g\ x) * \text{vector\_derivative } g \text{ (at } x)$ )
  using assms has_contour_integral_subpath contour_integral_unique by blast

```

lemma contour_integral_subpath_combine_less:

```

assumes f contour_integrable_on g valid_path g  $u \in \{0..1\}$   $v \in \{0..1\}$   $w \in \{0..1\}$ 
   $u < v$   $v < w$ 
  shows contour_integral (subpath  $u$   $v$  g)  $f + \text{contour\_integral (subpath } v \text{ } w \text{ g)}$ 
 $f =$ 
    contour_integral (subpath  $u$   $w$  g)  $f$ 
by (smt (verit) Henstock_Kurzweil_Integration.integral_combine assms
  has_integral_contour_integral_subpath has_integral_iff)

```

lemma contour_integral_subpath_combine:

assumes f *contour_integrable_on* g *valid_path* g $u \in \{0..1\}$ $v \in \{0..1\}$ $w \in \{0..1\}$

shows $\text{contour_integral (subpath } u \ v \ g) f + \text{contour_integral (subpath } v \ w \ g) f =$

$\text{contour_integral (subpath } u \ w \ g) f$

proof (*cases* $u \neq v \wedge v \neq w \wedge u \neq w$)

case *True*

have $*$: $\text{subpath } v \ u \ g = \text{reversepath (subpath } u \ v \ g) \wedge$
 $\text{subpath } w \ u \ g = \text{reversepath (subpath } u \ w \ g) \wedge$
 $\text{subpath } w \ v \ g = \text{reversepath (subpath } v \ w \ g)$

by (*auto simp: reversepath_subpath*)

have $u < v \wedge v < w \vee$

$u < w \wedge w < v \vee$

$v < u \wedge u < w \vee$

$v < w \wedge w < u \vee$

$w < u \wedge u < v \vee$

$w < v \wedge v < u$

using *True* *assms* **by** *linarith*

with *assms* **show** *?thesis*

using *contour_integral_subpath_combine_less* [*of* $f \ g \ u \ v \ w$]

contour_integral_subpath_combine_less [*of* $f \ g \ u \ w \ v$]

contour_integral_subpath_combine_less [*of* $f \ g \ v \ u \ w$]

contour_integral_subpath_combine_less [*of* $f \ g \ v \ w \ u$]

contour_integral_subpath_combine_less [*of* $f \ g \ w \ u \ v$]

contour_integral_subpath_combine_less [*of* $f \ g \ w \ v \ u$]

by (*elim disjE*) (*auto simp: * contour_integral_reversepath contour_integrable_subpath valid_path_subpath algebra_simps*)

next

case *False*

with *assms* **show** *?thesis*

by (*metis add.right_neutral contour_integral_reversepath contour_integral_subpath_refl diff_0_eq_diff_eq add_0 reversepath_subpath valid_path_subpath*)

qed

lemma *contour_integral_integral*:

$\text{contour_integral } g \ f = \text{integral } \{0..1\} (\lambda x. f (g \ x) * \text{vector_derivative } g \ (\text{at } x))$

by (*simp add: contour_integral_def integral_def has_contour_integral contour_integrable_on*)

lemma *contour_integral_cong*:

assumes $g = g' \wedge x. x \in \text{path_image } g \implies f \ x = f' \ x$

shows $\text{contour_integral } g \ f = \text{contour_integral } g' \ f'$

unfolding *contour_integral_integral* **using** *assms*

by (*intro integral_cong*) (*auto simp: path_image_def*)

lemma *contour_integral_spike_finite_simple_path*:

assumes *finite* A *simple_path* $g \ g = g' \wedge x. x \in \text{path_image } g - A \implies f \ x = f' \ x$

shows $\text{contour_integral } g \ f = \text{contour_integral } g' \ f'$

```

unfolding contour_integral_integral
proof (rule integral_spike)
  have finite (g - ` A ∩ {0<..<1}) using ⟨simple_path g⟩ ⟨finite A⟩
    by (intro finite_vimage_IntI simple_path_inj_on) auto
  hence finite ({0, 1} ∪ g - ` A ∩ {0<..<1}) by auto
  thus negligible ({0, 1} ∪ g - ` A ∩ {0<..<1}) by (rule negligible_finite)
next
  fix x assume x ∈ {0..1} - ({0, 1} ∪ g - ` A ∩ {0<..<1})
  hence g x ∈ path_image g - A by (auto simp: path_image_def)
  with assms show f' (g' x) * vector_derivative g' (at x) = f (g x) * vector_derivative g (at x)
    by simp
qed

```

Contour integral along a segment on the real axis

```

lemma has_contour_integral_linepath_Reals_iff:
  fixes a b :: complex and f :: complex ⇒ complex
  assumes a ∈ Reals b ∈ Reals Re a < Re b
  shows (f has_contour_integral I) (linepath a b) ↔
    ((λx. f (of_real x)) has_integral I) {Re a..Re b}
proof -
  have [simp]: of_real (Re a) = a of_real (Re b) = b and a ≠ b
    using assms by (simp_all add: complex_eq_iff)
  have ((λx. f (of_real x)) has_integral I) (cbox (Re a) (Re b)) ↔
    ((λx. f (a + b * of_real x - a * of_real x)) has_integral I /R (Re b - Re a)) {0..1}
    by (subst has_integral_affinity_iff [of Re b - Re a _ Re a, symmetric])
      (insert assms, simp_all add: field_simps scaleR_conv_of_real)
  also have (λx. f (a + b * of_real x - a * of_real x)) =
    (λx. (f (a + b * of_real x - a * of_real x) * (b - a)) /R (Re b - Re a))
    using ⟨a ≠ b⟩ by (auto simp: field_simps fun_eq_iff scaleR_conv_of_real)
  also have (... has_integral I /R (Re b - Re a)) {0..1} ↔
    ((λx. f (linepath a b x) * (b - a)) has_integral I) {0..1} using assms
    by (subst has_integral_cmul_iff) (auto simp: linepath_def scaleR_conv_of_real algebra_simps)
  also have ... ↔ (f has_contour_integral I) (linepath a b)
    unfolding has_contour_integral_def
    using has_contour_integral_def has_contour_integral_linepath by presburger
  finally show ?thesis by simp
qed

```

```

lemma contour_integrable_linepath_Reals_iff:
  fixes a b :: complex and f :: complex ⇒ complex
  assumes a ∈ Reals b ∈ Reals Re a < Re b
  shows (f contour_integrable_on linepath a b) ↔
    (λx. f (of_real x)) integrable_on {Re a..Re b}
  using has_contour_integral_linepath_Reals_iff[OF assms, of f]
  by (auto simp: contour_integrable_on_def integrable_on_def)

```

```

lemma contour_integral_linepath_Reals_eq:
  fixes a b :: complex and f :: complex  $\Rightarrow$  complex
  assumes a  $\in$  Reals b  $\in$  Reals Re a < Re b
  shows contour_integral (linepath a b) f = integral {Re a..Re b} ( $\lambda x$ . f (of_real x))
proof (cases f contour_integrable_on linepath a b)
  case True
  thus ?thesis
    by (metis assms has_contour_integral_integral
        has_contour_integral_linepath_Reals_iff integral_unique)
  next
  case False
  thus ?thesis
    by (simp add: assms contour_integrable_linepath_Reals_iff
        not_integrable_contour_integral not_integrable_integral)
qed

```

1.7 Cauchy's theorem where there's a primitive

```

lemma contour_integral_primitive_lemma:
  fixes f :: complex  $\Rightarrow$  complex and g :: real  $\Rightarrow$  complex
  assumes a  $\leq$  b
    and  $\bigwedge x. x \in S \implies$  (f has_field_derivative f' x) (at x within S)
    and g piecewise_differentiable_on {a..b}  $\bigwedge x. x \in \{a..b\} \implies$  g x  $\in$  S
  shows (( $\lambda x$ . f'(g x) * vector_derivative g (at x within {a..b}))
        has_integral (f(g b) - f(g a))) {a..b}
proof -
  obtain K where finite K and K:  $\forall x \in \{a..b\} - K. g$  differentiable (at x within {a..b})
  and cg: continuous_on {a..b} g
  using assms by (auto simp: piecewise_differentiable_on_def)
  have continuous_on (g ' {a..b}) f
    using assms by (metis DERIV_continuous_on continuous_on_subset image_subsetI)
  then have cfg: continuous_on {a..b} ( $\lambda x$ . f (g x))
    by (rule continuous_on_compose [OF cg, unfolded o_def])
  { fix x::real
    assume a: a < x and b: x < b and xk: x  $\notin$  K
    then have g differentiable at x within {a..b}
      using K by (simp add: differentiable_at_withinI)
    then have (g has_vector_derivative vector_derivative g (at x within {a..b}))
      (at x within {a..b})
      by (simp add: vector_derivative_works has_field_derivative_def scaleR_conv_of_real)
    then have gdiff: (g has_derivative ( $\lambda u. u * vector_derivative g (at x within \{a..b\})$ ))
      (at x within {a..b})
      by (simp add: has_vector_derivative_def scaleR_conv_of_real)
    have (f has_field_derivative (f' (g x))) (at (g x) within g ' {a..b})
      using assms by (metis a atLeastAtMost_iff b DERIV_subset_image_subset_iff less_eq_real_def)
  }

```

```

then have fdiff: (f has_derivative (*) (f' (g x))) (at (g x) within g ` {a..b})
  by (simp add: has_field_derivative_def)
have (( $\lambda x. f (g x)$ ) has_vector_derivative f' (g x) * vector_derivative g (at x
within {a..b})) (at x within {a..b})
  using diff_chain_within [OF gdiff fdiff]
  by (simp add: has_vector_derivative_def scaleR_conv_of_real o_def mult_ac)
} then show ?thesis
  using assms cfg
  by (force simp: at_within_Icc_at intro: fundamental_theorem_of_calculus_interior_strong
[OF <finite K>])
qed

```

lemma *contour_integral_primitive*:

```

assumes  $\bigwedge x. x \in S \implies$  (f has_field_derivative f' x) (at x within S)
  and valid_path g path_image g  $\subseteq$  S
  shows (f' has_contour_integral (f(pathfinish g) - f(pathstart g))) g
  using assms
  apply (simp add: valid_path_def path_image_def pathfinish_def pathstart_def
has_contour_integral_def)
  apply (auto intro!: piecewise_C1_imp_differentiable contour_integral_primitive_lemma
[of 0 1 S])
done

```

corollary *Cauchy_theorem_primitive*:

```

assumes  $\bigwedge x. x \in S \implies$  (f has_field_derivative f' x) (at x within S)
  and valid_path g path_image g  $\subseteq$  S pathfinish g = pathstart g
  shows (f' has_contour_integral 0) g
  using assms by (metis diff_self contour_integral_primitive)

```

lemma *contour_integrable_continuous_linepath*:

```

assumes continuous_on (closed_segment a b) f
  shows f contour_integrable_on (linepath a b)
proof -
  have continuous_on (closed_segment a b) ( $\lambda x. f x * (b - a)$ )
    by (rule continuous_intros | simp add: assms)+
  then have continuous_on {0..1} ( $\lambda x. f (linepath a b x) * (b - a)$ )
    by (metis (no_types, lifting) continuous_on_compose continuous_on_cong
continuous_on_linepath linepath_image_01 o_apply)
  then have ( $\lambda x. f (linepath a b x)$ 
    * vector_derivative (linepath a b) (at x within {0..1}))
    integrable_on {0..1}
    by (metis (no_types, lifting) continuous_on_cong integrable_continuous_real
vector_derivative_linepath_within)
  then show ?thesis
    by (simp add: contour_integrable_on_def has_contour_integral_def integrable_on_def
[symmetric])
qed

```

lemma *has_field_der_id*: $((\lambda x. x^2/2) \text{ has_field_derivative } x) \text{ (at } x)$
by (*rule has_derivative_imp_has_field_derivative*)
(rule derivative_intros | simp)+

lemma *contour_integral_id* [*simp*]: $\text{contour_integral (linepath a b) } (\lambda y. y) = (b^2 - a^2)/2$
using *contour_integral_primitive* [*of UNIV* $\lambda x. x^2/2$ $\lambda x. x$ *linepath a b*] *contour_integral_unique*
by (*simp add: has_field_der_id*)

lemma *contour_integrable_on_const* [*iff*]: $(\lambda x. c) \text{ contour_integrable_on (linepath a b)}$
by (*simp add: contour_integrable_continuous_linepath*)

lemma *contour_integrable_on_id* [*iff*]: $(\lambda x. x) \text{ contour_integrable_on (linepath a b)}$
by (*simp add: contour_integrable_continuous_linepath*)

1.8 Arithmetical combining theorems

lemma *has_contour_integral_neg*:
 $(f \text{ has_contour_integral } i) \text{ } g \implies ((\lambda x. -(f x)) \text{ has_contour_integral } (-i)) \text{ } g$
by (*simp add: has_integral_neg has_contour_integral_def*)

lemma *has_contour_integral_add*:
 $\llbracket (f1 \text{ has_contour_integral } i1) \text{ } g; (f2 \text{ has_contour_integral } i2) \text{ } g \rrbracket$
 $\implies ((\lambda x. f1 x + f2 x) \text{ has_contour_integral } (i1 + i2)) \text{ } g$
by (*simp add: has_integral_add has_contour_integral_def algebra_simps*)

lemma *has_contour_integral_diff*:
 $\llbracket (f1 \text{ has_contour_integral } i1) \text{ } g; (f2 \text{ has_contour_integral } i2) \text{ } g \rrbracket$
 $\implies ((\lambda x. f1 x - f2 x) \text{ has_contour_integral } (i1 - i2)) \text{ } g$
by (*simp add: has_integral_diff has_contour_integral_def algebra_simps*)

lemma *has_contour_integral_lmul*:
 $(f \text{ has_contour_integral } i) \text{ } g \implies ((\lambda x. c * (f x)) \text{ has_contour_integral } (c*i)) \text{ } g$
by (*simp add: has_contour_integral_def algebra_simps has_integral_mult_right*)

lemma *has_contour_integral_rmul*:
 $(f \text{ has_contour_integral } i) \text{ } g \implies ((\lambda x. (f x) * c) \text{ has_contour_integral } (i*c)) \text{ } g$
by (*simp add: mult.commute has_contour_integral_lmul*)

lemma *has_contour_integral_div*:
 $(f \text{ has_contour_integral } i) \text{ } g \implies ((\lambda x. f x/c) \text{ has_contour_integral } (i/c)) \text{ } g$
by (*simp add: field_class.field_divide_inverse*) (*metis has_contour_integral_rmul*)

lemma *has_contour_integral_eq*:
 $\llbracket (f \text{ has_contour_integral } y) \text{ } p; \bigwedge x. x \in \text{path_image } p \implies f x = g x \rrbracket \implies (g$

has_contour_integral y) p

by (*metis* (*mono_tags*, *lifting*) *has_contour_integral_def* *has_integral_eq* *image_eqI* *path_image_def*)

lemma *has_contour_integral_bound_linepath*:

assumes (f *has_contour_integral* i) (*linepath* a b)

$0 \leq B$ **and** $B: \bigwedge x. x \in \text{closed_segment } a \ b \implies \text{norm}(f \ x) \leq B$

shows $\text{norm } i \leq B * \text{norm}(b - a)$

proof –

have $\text{norm } i \leq (B * \text{norm}(b - a)) * \text{measure } \text{lborel } (\text{cbox } 0 \ (1::\text{real}))$

proof (*rule* *has_integral_bound*)

[*of* $_ \lambda x. f$ (*linepath* a b x) * *vector_derivative* (*linepath* a b) (*at* x *within* $\{0..1\}$)]]

show $\text{cmod } (f$ (*linepath* a b x) * *vector_derivative* (*linepath* a b) (*at* x *within* $\{0..1\}$))

$\leq B * \text{cmod}(b - a)$

if $x \in \text{cbox } 0 \ 1$ **for** $x::\text{real}$

using *that* *box_real(2)* *norm_mult*

by (*metis* B *linepath_in_path* *mult_right_mono* *norm_ge_zero* *vector_derivative_linepath_within*)

qed (*use* *assms* *has_contour_integral_def* **in** *auto*)

then show *?thesis*

by (*auto simp: content_real*)

qed

lemma *has_contour_integral_const_linepath*: ($(\lambda x. c)$ *has_contour_integral* $c*(b - a)$)(*linepath* a b)

unfolding *has_contour_integral_linepath*

by (*metis* *content_real* *diff_0_right* *has_integral_const_real* *lambda_one_of_real_1* *scaleR_conv_of_real* *zero_le_one*)

lemma *has_contour_integral_0*: ($(\lambda x. 0)$ *has_contour_integral* 0) g

by (*simp add: has_contour_integral_def*)

lemma *has_contour_integral_is_0*:

($\bigwedge z. z \in \text{path_image } g \implies f \ z = 0$) $\implies (f$ *has_contour_integral* 0) g

by (*rule* *has_contour_integral_eq* [*OF* *has_contour_integral_0*]) *auto*

lemma *has_contour_integral_sum*:

$\llbracket \text{finite } s; \bigwedge a. a \in s \implies (f \ a \ \text{has_contour_integral } i \ a) \ p \rrbracket$

$\implies ((\lambda x. \text{sum } (\lambda a. f \ a \ x) \ s) \ \text{has_contour_integral } \text{sum } i \ s) \ p$

by (*induction* s *rule: finite_induct*) (*auto simp: has_contour_integral_0* *has_contour_integral_add*)

1.9 Operations on path integrals

lemma *contour_integral_const_linepath* [*simp*]: *contour_integral* (*linepath* a b) $(\lambda x. c) = c*(b - a)$

by (*rule* *contour_integral_unique* [*OF* *has_contour_integral_const_linepath*])

lemma *contour_integral_neg*: *contour_integral* g $(\lambda z. -f \ z) = -\text{contour_integral}$

$g f$
by (*simp add: contour_integral_integral*)

lemma *contour_integral_add*:

$f1 \text{ contour_integrable_on } g \implies f2 \text{ contour_integrable_on } g \implies \text{contour_integral } g (\lambda x. f1 x + f2 x) =$
 $\text{contour_integral } g f1 + \text{contour_integral } g f2$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_add*)

lemma *contour_integral_diff*:

$f1 \text{ contour_integrable_on } g \implies f2 \text{ contour_integrable_on } g \implies \text{contour_integral } g (\lambda x. f1 x - f2 x) =$
 $\text{contour_integral } g f1 - \text{contour_integral } g f2$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_diff*)

lemma *contour_integral_lmul*:

shows $f \text{ contour_integrable_on } g$
 $\implies \text{contour_integral } g (\lambda x. c * f x) = c * \text{contour_integral } g f$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_lmul*)

lemma *contour_integral_rmul*:

shows $f \text{ contour_integrable_on } g$
 $\implies \text{contour_integral } g (\lambda x. f x * c) = \text{contour_integral } g f * c$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_rmul*)

lemma *contour_integral_div*:

shows $f \text{ contour_integrable_on } g$
 $\implies \text{contour_integral } g (\lambda x. f x / c) = \text{contour_integral } g f / c$

by (*simp add: contour_integral_unique has_contour_integral_integral has_contour_integral_div*)

lemma *contour_integral_eq*:

$(\bigwedge x. x \in \text{path_image } p \implies f x = g x) \implies \text{contour_integral } p f = \text{contour_integral } p g$

using *contour_integral_cong contour_integral_def* **by** *fastforce*

lemma *contour_integral_eq_0*:

$(\bigwedge z. z \in \text{path_image } g \implies f z = 0) \implies \text{contour_integral } g f = 0$

by (*simp add: has_contour_integral_is_0 contour_integral_unique*)

lemma *contour_integral_bound_linepath*:

shows
 $\llbracket f \text{ contour_integrable_on } (\text{linepath } a b);$
 $0 \leq B; \bigwedge x. x \in \text{closed_segment } a b \implies \text{norm}(f x) \leq B \rrbracket$
 $\implies \text{norm}(\text{contour_integral } (\text{linepath } a b) f) \leq B * \text{norm}(b - a)$

by (*meson has_contour_integral_bound_linepath has_contour_integral_integral*)

lemma *contour_integral_0 [simp]*: $\text{contour_integral } g (\lambda x. 0) = 0$

by (*simp add: contour_integral_unique has_contour_integral_0*)

lemma *contour_integral_sum*:

$\llbracket \text{finite } s; \bigwedge a. a \in s \implies (f\ a)\ \text{contour_integrable_on } p \rrbracket$
 $\implies \text{contour_integral } p\ (\lambda x. \text{sum } (\lambda a. f\ a\ x)\ s) = \text{sum } (\lambda a. \text{contour_integral } p\ (f\ a))\ s$
by (*auto simp: contour_integral_unique has_contour_integral_sum has_contour_integral_integral*)

lemma *contour_integrable_eq*:

$\llbracket f\ \text{contour_integrable_on } p; \bigwedge x. x \in \text{path_image } p \implies f\ x = g\ x \rrbracket \implies g\ \text{contour_integrable_on } p$
unfolding *contour_integrable_on_def*
by (*metis has_contour_integral_eq*)

1.10 Arithmetic theorems for path integrability

lemma *contour_integrable_neg*:

$f\ \text{contour_integrable_on } g \implies (\lambda x. -(f\ x))\ \text{contour_integrable_on } g$
using *has_contour_integral_neg contour_integrable_on_def* **by** *blast*

lemma *contour_integrable_add*:

$\llbracket f1\ \text{contour_integrable_on } g; f2\ \text{contour_integrable_on } g \rrbracket \implies (\lambda x. f1\ x + f2\ x)\ \text{contour_integrable_on } g$
using *has_contour_integral_add contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_diff*:

$\llbracket f1\ \text{contour_integrable_on } g; f2\ \text{contour_integrable_on } g \rrbracket \implies (\lambda x. f1\ x - f2\ x)\ \text{contour_integrable_on } g$
using *has_contour_integral_diff contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_lmul*:

$f\ \text{contour_integrable_on } g \implies (\lambda x. c * f\ x)\ \text{contour_integrable_on } g$
using *has_contour_integral_lmul contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_rmul*:

$f\ \text{contour_integrable_on } g \implies (\lambda x. f\ x * c)\ \text{contour_integrable_on } g$
using *has_contour_integral_rmul contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_div*:

$f\ \text{contour_integrable_on } g \implies (\lambda x. f\ x / c)\ \text{contour_integrable_on } g$
using *has_contour_integral_div contour_integrable_on_def*
by *fastforce*

lemma *contour_integrable_sum*:

$\llbracket \text{finite } s; \bigwedge a. a \in s \implies (f\ a)\ \text{contour_integrable_on } p \rrbracket$
 $\implies (\lambda x. \text{sum } (\lambda a. f\ a\ x)\ s)\ \text{contour_integrable_on } p$
unfolding *contour_integrable_on_def* **by** (*metis has_contour_integral_sum*)

lemma *contour_integrable_neg_iff*:
 $(\lambda x. -f x)$ *contour_integrable_on* $g \longleftrightarrow f$ *contour_integrable_on* g
using *contour_integrable_neg*[of $f g$] *contour_integrable_neg*[of $\lambda x. -f x g$] **by**
auto

lemma *contour_integrable_lmul_iff*:
 $c \neq 0 \implies (\lambda x. c * f x)$ *contour_integrable_on* $g \longleftrightarrow f$ *contour_integrable_on*
 g
using *contour_integrable_lmul*[of $f g c$] *contour_integrable_lmul*[of $\lambda x. c * f x$
 g *inverse* c]
by (*auto simp: field_simps*)

lemma *contour_integrable_rmul_iff*:
 $c \neq 0 \implies (\lambda x. f x * c)$ *contour_integrable_on* $g \longleftrightarrow f$ *contour_integrable_on*
 g
using *contour_integrable_rmul*[of $f g c$] *contour_integrable_rmul*[of $\lambda x. c * f x$
 g *inverse* c]
by (*auto simp: field_simps*)

lemma *contour_integrable_div_iff*:
 $c \neq 0 \implies (\lambda x. f x / c)$ *contour_integrable_on* $g \longleftrightarrow f$ *contour_integrable_on*
 g
using *contour_integrable_rmul_iff*[of *inverse* c] **by** (*simp add: field_simps*)

lemma *uniform_limit_contour_integral_linepath*:
assumes u : *uniform_limit* (*path_image* (*linepath* $a b$)) $f g F$
assumes c : $\bigwedge n. \text{continuous_on}$ (*path_image* (*linepath* $a b$)) ($f n$)
assumes [*simp*]: $F \neq \text{bot}$
obtains $I J$ **where**
 $\bigwedge n. (f n \text{ has_contour_integral } I n)$ (*linepath* $a b$)
 $(g \text{ has_contour_integral } J)$ (*linepath* $a b$)
 $(I \longrightarrow J) F$
proof (*rule uniform_limit_integral*)
note [*continuous_intros*] = *continuous_on_compose2*[*OF* c]

show *uniform_limit* $\{0..1\}$ $(\lambda x t. f x (\text{linepath } a b t) * (b - a))$
 $(\lambda t. g (\text{linepath } a b t) * (b - a)) F$
proof (*rule uniform_limit_intros*)
show *uniform_limit* $\{0..1\}$ $(\lambda x t. f x (\text{linepath } a b t))$
 $(\lambda t. g (\text{linepath } a b t)) F$
using u **unfolding** *path_image_def* **by** (*rule uniform_limit_compose'*) *auto*
qed

show *continuous_on* $\{0..1\}$ $(\lambda t. f n (\text{linepath } a b t) * (b - a))$ **for** n
by (*intro continuous_intros; unfold path_image_def*) *auto*

fix $I J$

```

assume  $I$ :  $\bigwedge n. ((\lambda t. f\ n\ (\text{linepath}\ a\ b\ t) * (b - a))\ \text{has\_integral}\ I\ n)\ \{0..1\}$ 
and  $J$ :  $((\lambda t. g\ (\text{linepath}\ a\ b\ t) * (b - a))\ \text{has\_integral}\ J)\ \{0..1\}$ 
and  $lim$ :  $(I \longrightarrow J)\ F$ 
show ?thesis
by (rule that[of I J]) (use I J lim in <auto simp: has_contour_integral>)
qed auto

```

lemma *contour_integral_sums_linepath*:

```

assumes  $u$ : uniform_limit (closed_segment  $a\ b$ )  $(\lambda N\ w. \sum_{n < N}. f\ n\ w)$  g sequentially

```

```

assumes  $c$ :  $\bigwedge n. \text{continuous\_on}\ (\text{closed\_segment}\ a\ b)\ (f\ n)$ 

```

```

obtains  $J$  where

```

```

   $(g\ \text{has\_contour\_integral}\ J)\ (\text{linepath}\ a\ b)$ 

```

```

   $(\lambda n. \text{contour\_integral}\ (\text{linepath}\ a\ b)\ (f\ n))\ \text{sums}\ J$ 

```

```

proof (rule uniform_limit_contour_integral_linepath)

```

```

show uniform_limit (path_image (linepath  $a\ b$ ))  $(\lambda N\ w. \sum_{n < N}. f\ n\ w)$  g sequentially

```

```

  using  $u$  by simp

```

```

next

```

```

show continuous_on (path_image (linepath  $a\ b$ ))  $(\lambda w. \sum_{n < N}. f\ n\ w)$  for  $N$ 

```

```

  by (intro continuous_intros continuous_on_subset[OF c]) simp_all

```

```

next

```

```

fix  $I\ J$ 

```

```

assume  $1$ :  $\bigwedge N. ((\lambda w. \sum_{n < N}. f\ n\ w)\ \text{has\_contour\_integral}\ I\ N)\ (\text{linepath}\ a\ b)$ 

```

```

assume  $2$ :  $(g\ \text{has\_contour\_integral}\ J)\ (\text{linepath}\ a\ b)$  and  $3$ :  $(I \longrightarrow J)$  sequentially

```

```

have  $4$ :  $I = (\lambda N. (\sum_{n < N}. \text{contour\_integral}\ (\text{linepath}\ a\ b)\ (f\ n)))$ 

```

```

proof

```

```

  fix  $N :: \text{nat}$ 

```

```

  have  $f\ n$  contour_integrable_on (linepath  $a\ b$ ) for  $n$ 

```

```

    by (intro contour_integrable_continuous_linepath assms)

```

```

  hence  $((\lambda w. \sum_{n < N}. f\ n\ w)\ \text{has\_contour\_integral}$ 

```

```

     $(\sum_{n < N}. \text{contour\_integral}\ (\text{linepath}\ a\ b)\ (f\ n)))\ (\text{linepath}\ a\ b)$ 

```

```

  using  $c$  by (intro has_contour_integral_sum) (simp_all add: has_contour_integral_integral)

```

```

  with  $1$ [of N] show  $I\ N = (\sum_{n < N}. \text{contour\_integral}\ (\text{linepath}\ a\ b)\ (f\ n))$ 

```

```

    using contour_integral_unique by metis

```

```

qed

```

```

have  $5$ :  $(\lambda n. \text{contour\_integral}\ (\text{linepath}\ a\ b)\ (f\ n))\ \text{sums}\ J$ 

```

```

  using  $1\ 2\ 3\ 4$  unfolding sums_def by blast

```

```

from that[OF 2 5] show ?thesis .

```

```

qed auto

```

lemma *contour_integral_linepath_same_Re*:

```

assumes  $\text{Re}\ z = c\ \text{Re}\ z' = c\ \text{Im}\ z = a\ \text{Im}\ z' = b\ a < b$ 

```

```

shows contour_integral (linepath  $z\ z'$ )  $f =$ 

```

```

   $i * \text{integral}\ \{a..b\}\ (\lambda x. f\ (\text{Complex}\ c\ x))$ 

```

```

proof -

```

```

have zz': z = Complex c a z' = Complex c b
  using assms by (auto simp: complex_eq_iff)
have contour_integral (linepath z z') f =
  (z' - z) * integral {0..1} (λx. f (linepath z z' x))
  by (simp add: contour_integral_integral)
also have z' - z = i * of_real (b - a)
  by (simp add: zz' Complex_eq algebra_simps)
also have integral {0..1} (λx. f (linepath z z' x)) =
  integral {0..1} (λx. f (Complex c (linepath a b x)))
  by (simp add: linepath_def Complex_eq scaleR_conv_of_real algebra_simps
zz')
also have ... = integral {0..(b - a) / (b - a)} (λx. f (Complex c (a + (b - a)
* x)))
  using ⟨a < b⟩ by (simp add: algebra_simps linepath_def)
also have {0..(b - a) / (b - a)} = (λx. x / (b - a)) ' {0..b - a}
  using ⟨a < b⟩ by simp
also have integral ... (λx. f (Complex c (a + (b - a) * x))) =
  integral {a..a..b-a} (λx. f (Complex c (x + a))) / of_real (b - a)
  using ⟨a < b⟩ by (subst integral_stretch_real) (auto simp: scaleR_conv_of_real
add_ac)
also have ... = integral {a..b} (λx. f (Complex c x)) / of_real (b - a)
  by (subst integral_shift_real_ivl) (rule refl)
finally show ?thesis
  using ⟨a < b⟩ by simp
qed

```

1.11 Reversing a path integral

```

lemma has_contour_integral_reverse_linepath:
  (f has_contour_integral i) (linepath a b)
  ⇒ (f has_contour_integral (-i)) (linepath b a)
  using has_contour_integral_reversepath valid_path_linepath by fastforce

```

```

lemma contour_integral_reverse_linepath:
  continuous_on (closed_segment a b) f ⇒ contour_integral (linepath a b) f =
  - (contour_integral (linepath b a) f)
  using contour_integral_reversepath by fastforce

```

Splitting a path integral in a flat way.*)

```

lemma has_contour_integral_split:
  assumes f: (f has_contour_integral i) (linepath a c) (f has_contour_integral j)
  (linepath c b)
  and k: 0 ≤ k k ≤ 1
  and c: c - a = k *R (b - a)
  shows (f has_contour_integral (i + j)) (linepath a b)
proof (cases k = 0 ∨ k = 1)
case True
then show ?thesis
  using assms by auto

```

```

next
case False
then have k: 0 < k k < 1
  using assms by auto
have c': c = k *R (b - a) + a
  by (metis diff_add_cancel c)
have bc: (b - c) = (1 - k) *R (b - a)
  by (simp add: algebra_simps c')
{ assume *: ((λx. f ((1 - x) *R a + x *R c) * (c - a)) has_integral i) {0..1}
  have ∧x. (x / k) *R a + ((k - x) / k) *R a = a
  using False by (simp add: field_split_simps flip: real_vector.scale_left_distrib)
  then have ∧x. ((k - x) / k) *R a + (x / k) *R c = (1 - x) *R a + x *R b
  using False by (simp add: c' algebra_simps)
  then have ((λx. f ((1 - x) *R a + x *R b) * (b - a)) has_integral i) {0..k}
  using k has_integral_affinity01 [OF *, of inverse k 0]
  by (force dest: has_integral_cmul [where c = inverse k]
      simp add: divide_simps mult.commute [of _ k] image_affinity_atLeastAtMost
  c)
} note fi = this
{ assume *: ((λx. f ((1 - x) *R c + x *R b) * (b - c)) has_integral j) {0..1}
  have **: ∧x. (((1 - x) / (1 - k)) *R c + ((x - k) / (1 - k)) *R b) = ((1 -
x) *R a + x *R b)
  using k
  apply (simp add: c' scaleR_conv_of_real divide_simps)
  apply (simp add: distrib_right distrib_left right_diff_distrib left_diff_distrib)
  done
  have ((λx. f ((1 - x) *R a + x *R b) * (b - a)) has_integral j) {k..1}
  using k has_integral_affinity01 [OF *, of inverse(1 - k) -(k/(1 - k))]
  apply (simp add: divide_simps mult.commute [of _ 1-k] image_affinity_atLeastAtMost
  ** bc)
  apply (auto dest: has_integral_cmul [where k = (1 - k) *R j and c =
inverse (1 - k)])
  done
}
then show ?thesis
  using f k unfolding has_contour_integral_linepath
  by (simp add: linepath_def has_integral_combine [OF _ _ fi])
qed

lemma continuous_on_closed_segment_transform:
  assumes f: continuous_on (closed_segment a b) f
  and k: 0 ≤ k k ≤ 1
  and c: c - a = k *R (b - a)
  shows continuous_on (closed_segment a c) f
proof -
  have c': c = (1 - k) *R a + k *R b
  using c by (simp add: algebra_simps)
  have closed_segment a c ⊆ closed_segment a b
  by (metis c' ends_in_segment(1) in_segment(1) k subset_closed_segment)

```

then show *continuous_on* (closed_segment a c) f
by (rule *continuous_on_subset* [OF f])
qed

lemma *contour_integral_split*:

assumes *f*: *continuous_on* (closed_segment a b) *f*
and *k*: $0 \leq k \leq 1$
and *c*: $c - a = k *_R (b - a)$
shows $\text{contour_integral}(\text{linepath } a \ b) \ f = \text{contour_integral}(\text{linepath } a \ c) \ f +$
 $\text{contour_integral}(\text{linepath } c \ b) \ f$
proof –
have *c'*: $c = (1 - k) *_R a + k *_R b$
using *c* **by** (*simp add: algebra_simps*)
have closed_segment a c \subseteq closed_segment a b
by (*metis c' ends_in_segment(1) in_segment(1) k subset_closed_segment*)
moreover have closed_segment c b \subseteq closed_segment a b
by (*metis c' ends_in_segment(2) in_segment(1) k subset_closed_segment*)
ultimately
have *continuous_on* (closed_segment a c) *f* *continuous_on* (closed_segment c
b) *f*
by (*auto intro: continuous_on_subset* [OF f])
then have (*f has_contour_integral*
 $\text{contour_integral}(\text{linepath } a \ c) \ f + \text{contour_integral}(\text{linepath } c \ b) \ f$)
(*linepath* a b)
by (*meson c contour_integrable_continuous_linepath*
has_contour_integral_integral has_contour_integral_split k)
then show ?thesis
by (*metis contour_integral_unique*)
qed

lemma *contour_integral_split_linepath*:

assumes *f*: *continuous_on* (closed_segment a b) *f*
and *c*: $c \in \text{closed_segment } a \ b$
shows $\text{contour_integral}(\text{linepath } a \ b) \ f = \text{contour_integral}(\text{linepath } a \ c) \ f +$
 $\text{contour_integral}(\text{linepath } c \ b) \ f$
using *c* **by** (*auto simp: closed_segment_def algebra_simps intro!: contour_integral_split*
[OF f])

1.12 Reversing the order in a double path integral

The condition is stronger than needed but it's often true in typical situations

lemma *fst_im_cbox* [*simp*]: $\text{cbox } c \ d \neq \{\}$ $\implies (\text{fst } \text{cbox } (a,c) \ (b,d)) = \text{cbox } a \ b$
by (*auto simp: cbox_Pair_eq*)

lemma *snd_im_cbox* [*simp*]: $\text{cbox } a \ b \neq \{\}$ $\implies (\text{snd } \text{cbox } (a,c) \ (b,d)) = \text{cbox } c \ d$
by (*auto simp: cbox_Pair_eq*)

proposition *contour_integral_swap*:


```

assumes fcon: continuous_on (path_image g × path_image h) (λ(y1,y2). f y1
y2)
and vp: valid_path g valid_path h
and gvcon: continuous_on {0..1} (λt. vector_derivative g (at t))
and hvcon: continuous_on {0..1} (λt. vector_derivative h (at t))
shows contour_integral g (λw. contour_integral h (f w)) =
contour_integral h (λz. contour_integral g (λw. f w z))
proof -
have gcon: continuous_on {0..1} g and hcon: continuous_on {0..1} h
using assms by (auto simp: valid_path_def piecewise_C1_differentiable_on_def)
have fgh1: λx. (λt. f (g x) (h t)) = (λ(y1,y2). f y1 y2) ∘ (λt. (g x, h t))
by (rule ext) simp
have fgh2: λx. (λt. f (g t) (h x)) = (λ(y1,y2). f y1 y2) ∘ (λt. (g t, h x))
by (rule ext) simp
have fcon_im1: λx. 0 ≤ x ⇒ x ≤ 1 ⇒ continuous_on ((λt. (g x, h t)) ‘
{0..1}) (λ(x, y). f x y)
by (rule continuous_on_subset [OF fcon]) (auto simp: path_image_def)
have fcon_im2: λx. 0 ≤ x ⇒ x ≤ 1 ⇒ continuous_on ((λt. (g t, h x)) ‘
{0..1}) (λ(x, y). f x y)
by (rule continuous_on_subset [OF fcon]) (auto simp: path_image_def)
have continuous_on (cbox (0, 0) (1, 1::real)) ((λx. vector_derivative g (at x))
∘ fst)
continuous_on (cbox (0, 0) (1::real, 1)) ((λx. vector_derivative h (at x)) ∘
snd)
by (rule continuous_intros | simp add: gvcon hvcon)+
then have gvcon': continuous_on (cbox (0, 0) (1, 1::real)) (λz. vector_derivative
g (at (fst z)))
and hvcon': continuous_on (cbox (0, 0) (1::real, 1)) (λx. vector_derivative
h (at (snd x)))
by auto
have continuous_on ((λx. (g (fst x), h (snd x))) ‘ cbox (0,0) (1,1)) (λ(y1, y2).
f y1 y2)
by (auto simp: path_image_def intro: continuous_on_subset [OF fcon])
then have continuous_on (cbox (0, 0) (1, 1)) ((λ(y1, y2). f y1 y2) ∘ (λw. ((g
∘ fst) w, (h ∘ snd) w)))
by (intro gcon hcon continuous_intros | simp)+
then have fgh: continuous_on (cbox (0, 0) (1, 1)) (λx. f (g (fst x)) (h (snd
x)))
by auto
have integral {0..1} (λx. contour_integral h (f (g x)) * vector_derivative g (at
x)) =
integral {0..1} (λx. contour_integral h (λy. f (g x) y * vector_derivative g
(at x)))
proof (rule integral_cong [OF contour_integral_rmul [symmetric]])
have λx. x ∈ {0..1} ⇒
continuous_on {0..1} (λxa. f (g x) (h xa))
by (subst fgh1) (rule fcon_im1 hcon continuous_intros | simp)+
then show λx. x ∈ {0..1} ⇒ f (g x) contour_integrable_on h
unfolding contour_integrable_on

```

```

    using continuous_on_mult hvcon integrable_continuous_real by blast
  qed
  also have ... = integral {0..1}
    (λy. contour_integral g (λx. f x (h y) * vector_derivative h (at
y)))
  unfolding contour_integral_integral
  apply (subst integral_swap_continuous [where 'a = real and 'b = real, of 0
0 1 1, simplified])
  subgoal
    by (rule fgh gvcon' hvcon' continuous_intros | simp add: split_def)+
    by (simp add: mult.commute mult.left_commute)
  also have ... = contour_integral h (λz. contour_integral g (λw. f w z))
  unfolding contour_integral_integral integral_mult_left [symmetric]
  by (simp add: algebra_simps)
  finally show ?thesis
  by (simp add: contour_integral_integral)
qed

```

```

lemma valid_path_negatepath: valid_path γ ⇒ valid_path (uminus ∘ γ)
  unfolding o_def using piecewise_C1_differentiable_neg valid_path_def by
blast

```

```

lemma has_contour_integral_negatepath:

```

```

  assumes γ: valid_path γ and cint: ((λz. f (- z)) has_contour_integral - i) γ
  shows (f has_contour_integral i) (uminus ∘ γ)

```

```

proof -

```

```

  obtain S where cont: continuous_on {0..1} γ and finite S and diff: γ C1_differentiable_on
{0..1} - S

```

```

  using γ by (auto simp: valid_path_def piecewise_C1_differentiable_on_def)

```

```

  have ((λx. - (f (- γ x) * vector_derivative γ (at x within {0..1}))) has_integral
i) {0..1}

```

```

  using cint by (auto simp: has_contour_integral_def dest: has_integral_neg)

```

```

then

```

```

  have ((λx. f (- γ x) * vector_derivative (uminus ∘ γ) (at x within {0..1}))
has_integral i) {0..1}

```

```

proof (rule rev_iffD1 [OF _ has_integral_spike_eq])

```

```

  show negligible S

```

```

  by (simp add: ⟨finite S⟩ negligible_finite)

```

```

  show f (- γ x) * vector_derivative (uminus ∘ γ) (at x within {0..1}) =
- (f (- γ x) * vector_derivative γ (at x within {0..1}))

```

```

  if x ∈ {0..1} - S for x

```

```

proof -

```

```

  have vector_derivative (uminus ∘ γ) (at x within cbox 0 1) = - vec-
tor_derivative γ (at x within cbox 0 1)

```

```

proof (rule vector_derivative_within_cbox)

```

```

  show (uminus ∘ γ has_vector_derivative - vector_derivative γ (at x within
cbox 0 1)) (at x within cbox 0 1)

```

```

  using that unfolding o_def

```

```

  by (metis C1_differentiable_on_eq UNIV_I diff differentiable_subset

```

```

has_vector_derivative_minus subsetI that vector_derivative_works)
  qed (use that in auto)
  then show ?thesis
    by simp
  qed
qed
then show ?thesis by (simp add: has_contour_integral_def)
qed

```

```

lemma contour_integrable_negatepath:
  assumes  $\gamma$ : valid_path  $\gamma$  and  $\pi$ :  $(\lambda z. f (-z))$  contour_integrable_on  $\gamma$ 
  shows  $f$  contour_integrable_on  $(\text{uminus} \circ \gamma)$ 
  by (metis  $\gamma$  add.inverse_inverse contour_integrable_on_def has_contour_integral_negatepath
   $\pi$ )

```

```

lemma C1_differentiable_polynomial_function:
  fixes  $p :: \text{real} \Rightarrow 'a::\text{euclidean\_space}$ 
  shows polynomial_function  $p \implies p$  C1_differentiable_on  $S$ 
  by (metis continuous_on_polynomial_function C1_differentiable_on_def has_vector_derivative_polynomial_function)

```

```

lemma valid_path_polynomial_function:
  fixes  $p :: \text{real} \Rightarrow 'a::\text{euclidean\_space}$ 
  shows polynomial_function  $p \implies \text{valid\_path } p$ 
  by (force simp: valid_path_def piecewise_C1_differentiable_on_def continuous_on_polynomial_function C1_differentiable_polynomial_function)

```

```

lemma valid_path_subpath_trivial [simp]:
  fixes  $g :: \text{real} \Rightarrow 'a::\text{euclidean\_space}$ 
  shows  $z \neq g\ x \implies \text{valid\_path } (\text{subpath } x\ x\ g)$ 
  by (simp add: subpath_def valid_path_polynomial_function)

```

1.13 Partial circle path

```

definition part_circlepath ::  $[\text{complex}, \text{real}, \text{real}, \text{real}, \text{real}] \Rightarrow \text{complex}$ 
  where part_circlepath  $z\ r\ s\ t \equiv \lambda x. z + \text{of\_real } r * \exp(i * \text{of\_real } (\text{linepath } s\ t\ x))$ 

```

```

lemma pathstart_part_circlepath [simp]:
  pathstart(part_circlepath  $z\ r\ s\ t$ ) =  $z + r * \exp(i * s)$ 
  by (metis part_circlepath_def pathstart_def pathstart_linepath)

```

```

lemma pathfinish_part_circlepath [simp]:
  pathfinish(part_circlepath  $z\ r\ s\ t$ ) =  $z + r * \exp(i * t)$ 
  by (metis part_circlepath_def pathfinish_def pathfinish_linepath)

```

```

lemma reversepath_part_circlepath [simp]:
  reversepath (part_circlepath  $z\ r\ s\ t$ ) = part_circlepath  $z\ r\ t\ s$ 
  unfolding part_circlepath_def reversepath_def linepath_def
  by (auto simp: algebra_simps)

```

```

lemma has_vector_derivative_part_circlepath [derivative_intros]:
  ((part_circlepath z r s t) has_vector_derivative
   (i * r * (of_real t - of_real s) * exp(i * linepath s t x)))
  (at x within X)
unfolding part_circlepath_def linepath_def scaleR_conv_of_real
by (rule has_vector_derivative_real_field derivative_eq_intros | simp)+

lemma differentiable_part_circlepath:
  part_circlepath c r a b differentiable at x within A
using has_vector_derivative_part_circlepath[of c r a b x A] differentiableI_vector
by blast

lemma vector_derivative_part_circlepath:
  vector_derivative (part_circlepath z r s t) (at x) =
    i * r * (of_real t - of_real s) * exp(i * linepath s t x)
using has_vector_derivative_part_circlepath vector_derivative_at by blast

lemma vector_derivative_part_circlepath01:
  [[0 ≤ x; x ≤ 1]]
  ⇒ vector_derivative (part_circlepath z r s t) (at x within {0..1}) =
    i * r * (of_real t - of_real s) * exp(i * linepath s t x)
using has_vector_derivative_part_circlepath
by (auto simp: vector_derivative_at_within_ivl)

lemma valid_path_part_circlepath [simp]: valid_path (part_circlepath z r s t)
unfolding valid_path_def
by (auto simp: C1_differentiable_on_eq vector_derivative_works vector_derivative_part_circlepath
  has_vector_derivative_part_circlepath
  intro!: C1_differentiable_imp_piecewise continuous_intros)

lemma path_part_circlepath [simp]: path (part_circlepath z r s t)
by (simp add: valid_path_imp_path)

proposition path_image_part_circlepath:
  assumes s ≤ t
  shows path_image (part_circlepath z r s t) = {z + r * exp(i * of_real x) | x.
s ≤ x ∧ x ≤ t}
proof -
  { fix z::real
    assume 0 ≤ z z ≤ 1
    with ⟨s ≤ t⟩ have ∃x. (exp (i * linepath s t z) = exp (i * of_real x)) ∧ s ≤ x
  ∧ x ≤ t
    apply (rule_tac x=(1 - z) * s + z * t in exI)
    apply (simp add: linepath_def scaleR_conv_of_real algebra_simps)
    by (metis (no_types) affine_ineq mult commute mult_left_mono)
  }
  moreover
  { fix z

```

```

assume  $s \leq z \leq t$ 
then have  $z + \text{of\_real } r * \exp(i * \text{of\_real } z) \in (\lambda x. z + \text{of\_real } r * \exp(i * \text{linepath } s \ t \ x)) \text{ ' } \{0..1\}$ 
apply (rule_tac  $x=(z - s)/(t - s)$  in image_eqI)
apply (simp add: linepath_def scaleR_conv_of_real divide_simps exp_eq)
apply (auto simp: field_split_simps)
done
}
ultimately show ?thesis
by (fastforce simp add: path_image_def part_circlepath_def)
qed

```

```

lemma path_image_part_circlepath':
   $\text{path\_image } (\text{part\_circlepath } z \ r \ s \ t) = (\lambda x. z + r * \text{cis } x) \text{ ' } \text{closed\_segment } s \ t$ 
by (metis (no_types, lifting) ext cis_conv_exp image_image linepath_image_01
  part_circlepath_def path_image_def)

```

```

lemma path_image_part_circlepath_subset:
   $\llbracket s \leq t; 0 \leq r \rrbracket \implies \text{path\_image}(\text{part\_circlepath } z \ r \ s \ t) \subseteq \text{sphere } z \ r$ 
by (auto simp: path_image_part_circlepath sphere_def dist_norm algebra_simps
  norm_mult)

```

```

lemma in_path_image_part_circlepath:
  assumes  $w \in \text{path\_image}(\text{part\_circlepath } z \ r \ s \ t)$   $s \leq t$   $0 \leq r$ 
  shows  $\text{norm}(w - z) = r$ 
by (smt (verit) assms dist_norm mem_Collect_eq norm_minus_commute path_image_part_circlepath_subset sphere_def subsetD)

```

```

lemma path_image_part_circlepath_subset':
  assumes  $r \geq 0$ 
  shows  $\text{path\_image } (\text{part\_circlepath } z \ r \ s \ t) \subseteq \text{sphere } z \ r$ 
by (smt (verit) assms path_image_part_circlepath_subset reversepath_part_circlepath
  reversepath_simps(2))

```

```

lemma part_circlepath_cnj:  $\text{cnj } (\text{part\_circlepath } c \ r \ a \ b \ x) = \text{part\_circlepath } (\text{cnj } c) \ r \ (-a) \ (-b) \ x$ 
by (simp add: part_circlepath_def exp_cnj linepath_def algebra_simps)

```

```

lemma contour_integrable_on_compose_cnj_iff:
  assumes valid_path  $\gamma$ 
  shows  $f \text{ contour\_integrable\_on } (\text{cnj } \circ \gamma) \iff (\text{cnj } \circ f \circ \text{cnj}) \text{ contour\_integrable\_on } \gamma$ 

```

proof –

```

from assms obtain  $S$  where  $S: \text{finite } S \ \gamma \ C1\_differentiable\_on \ \{0..1\} \ - \ S$ 
unfolding valid_path_def piecewise_C1_differentiable_on_def by blast
have  $f \text{ contour\_integrable\_on } (\text{cnj } \circ \gamma) \iff$ 
   $((\lambda t. \text{cnj } (\text{cnj } (f (\text{cnj } (\gamma \ t)))) * \text{vector\_derivative } \gamma \text{ (at } t))) \text{ integrable\_on } \{0..1\}$ 
unfolding contour_integrable_on_o_def

```

```

proof (intro integrable_spike_finite_eq [OF S(1)])
  fix t :: real assume t ∈ {0..1} - S
  hence  $\gamma$  differentiable at t
    using S(2) by (meson C1_differentiable_on_eq)
  hence vector_derivative ( $\lambda x.$  cnj ( $\gamma$  x)) (at t) = cnj (vector_derivative  $\gamma$  (at t))
    by (rule vector_derivative_cnj)
  thus f (cnj ( $\gamma$  t)) * vector_derivative ( $\lambda x.$  cnj ( $\gamma$  x)) (at t) =
    cnj (cnj (f (cnj ( $\gamma$  t))) * vector_derivative  $\gamma$  (at t))
    by simp
  qed
  also have ...  $\longleftrightarrow$  (( $\lambda t.$  cnj (f (cnj ( $\gamma$  t))) * vector_derivative  $\gamma$  (at t)) integrable_on {0..1})
    by (rule integrable_on_cnj_iff)
  also have ...  $\longleftrightarrow$  (cnj  $\circ$  f  $\circ$  cnj) contour_integrable_on  $\gamma$ 
    by (simp add: contour_integrable_on_o_def)
  finally show ?thesis .
qed

```

lemma contour_integral_cnj:

```

assumes valid_path  $\gamma$ 
shows contour_integral (cnj  $\circ$   $\gamma$ ) f = cnj (contour_integral  $\gamma$  (cnj  $\circ$  f  $\circ$  cnj))
proof -
  from assms obtain S where S: finite S  $\gamma$  C1_differentiable_on {0..1} - S
  unfolding valid_path_def piecewise_C1_differentiable_on_def by blast
  have contour_integral (cnj  $\circ$   $\gamma$ ) f =
    integral {0..1} ( $\lambda t.$  cnj (cnj (f (cnj ( $\gamma$  t))) * vector_derivative  $\gamma$  (at t)))
  unfolding contour_integral_integral
  proof (intro integral_spike)
    fix t assume t ∈ {0..1} - S
    hence  $\gamma$  differentiable at t
      using S(2) by (meson C1_differentiable_on_eq)
    hence vector_derivative ( $\lambda x.$  cnj ( $\gamma$  x)) (at t) = cnj (vector_derivative  $\gamma$  (at t))
      by (rule vector_derivative_cnj)
    thus cnj (cnj (f (cnj ( $\gamma$  t))) * vector_derivative  $\gamma$  (at t)) =
      f ((cnj  $\circ$   $\gamma$ ) t) * vector_derivative (cnj  $\circ$   $\gamma$ ) (at t)
      by (simp add: o_def)
    qed (use S(1) in auto)
  also have ... = cnj (integral {0..1} ( $\lambda t.$  cnj (f (cnj ( $\gamma$  t))) * vector_derivative  $\gamma$  (at t)))
    by (subst integral_cnj [symmetric]) auto
  also have ... = cnj (contour_integral  $\gamma$  (cnj  $\circ$  f  $\circ$  cnj))
    by (simp add: contour_integral_integral)
  finally show ?thesis .
qed

```

lemma contour_integral_negatepath:

```

assumes valid_path  $\gamma$ 

```

shows $\text{contour_integral } (\text{uminus } \circ \gamma) f = -(\text{contour_integral } \gamma (\lambda x. f (-x)))$
(is ?lhs = ?rhs)

proof (cases f $\text{contour_integrable_on } (\text{uminus } \circ \gamma)$)

case *True*

hence *: (f $\text{has_contour_integral } ?lhs$) ($\text{uminus } \circ \gamma$)

using $\text{has_contour_integral_integral}$ **by** *blast*

have $((\lambda z. f (-z)) \text{has_contour_integral } - \text{contour_integral } (\text{uminus } \circ \gamma) f)$
 $(\text{uminus } \circ (\text{uminus } \circ \gamma))$

by (rule $\text{has_contour_integral_negatepath}$) (use * *assms in auto*)

hence $((\lambda x. f (-x)) \text{has_contour_integral } - ?lhs) \gamma$

by (*simp add: o_def*)

thus *?thesis*

by (*simp add: contour_integral_unique*)

next

case *False*

hence $\neg(\lambda z. f (-z)) \text{contour_integrable_on } \gamma$

using $\text{contour_integrable_negatepath}$ [of γ f] *assms by auto*

with *False show ?thesis*

by (*simp add: not_integrable_contour_integral*)

qed

lemma $\text{contour_integral_bound_part_circlepath}$:

assumes $f \text{contour_integrable_on_part_circlepath } c r a b$

assumes $B \geq 0 r \geq 0 \wedge x. x \in \text{path_image } (\text{part_circlepath } c r a b) \implies \text{norm } (f x) \leq B$

shows $\text{norm } (\text{contour_integral } (\text{part_circlepath } c r a b) f) \leq B * r * |b - a|$

proof -

let $?I = \text{integral } \{0..1\} (\lambda x. f (\text{part_circlepath } c r a b x) * i * \text{of_real } (r * (b - a))) *$

$\text{exp } (i * \text{linepath } a b x)$

have $\text{norm } ?I \leq \text{integral } \{0..1\} (\lambda x::\text{real}. B * 1 * (r * |b - a|) * 1)$

proof (rule $\text{integral_norm_bound_integral}$, *goal_cases*)

case 1

with *assms(1) show ?case*

by (*simp add: contour_integrable_on vector_derivative_part_circlepath mult_ac*)

next

case ($\exists x$)

with *assms(2-)* **show** *?case unfolding norm_mult norm_of_real abs_mult*

by (*intro mult_mono*) (*auto simp: path_image_def*)

qed *auto*

also **have** $?I = \text{contour_integral } (\text{part_circlepath } c r a b) f$

by (*simp add: contour_integral_integral vector_derivative_part_circlepath mult_ac*)

finally **show** *?thesis by simp*

qed

lemma $\text{has_contour_integral_part_circlepath_iff}$:

assumes $a < b$

shows $(f \text{has_contour_integral } I) (\text{part_circlepath } c r a b) \longleftrightarrow$

```

      (( $\lambda t. f (c + r * cis t) * r * i * cis t$ ) has_integral I) {a..b}
proof -
  have (f has_contour_integral I) (part_circlepath c r a b)  $\longleftrightarrow$ 
    (( $\lambda x. f (part\_circlepath\ c\ r\ a\ b\ x) * vector\_derivative (part\_circlepath\ c\ r\ a\ b)$ )
      (at x within {0..1})) has_integral I {0..1}
  unfolding has_contour_integral_def ..
  also have ...  $\longleftrightarrow$  (( $\lambda x. f (part\_circlepath\ c\ r\ a\ b\ x) * r * (b - a) * i * cis (linepath\ a\ b\ x)$ )
    has_integral I) {0..1}
  by (intro has_integral_cong, subst vector_derivative_part_circlepath01)
    (simp_all add: cis_conv_exp)
  also have ...  $\longleftrightarrow$  (( $\lambda x. f (c + r * exp (i * linepath (of\_real\ a) (of\_real\ b) x))$ )
    *
      r * i * exp (i * linepath (of\_real a) (of\_real b) x) *
      vector_derivative (linepath (of\_real a) (of\_real b))
      (at x within {0..1})) has_integral I) {0..1}
  by (intro has_integral_cong, subst vector_derivative_linepath_within)
    (auto simp: part_circlepath_def cis_conv_exp of_real_linepath [symmetric])
  also have ...  $\longleftrightarrow$  (( $\lambda z. f (c + r * exp (i * z)) * r * i * exp (i * z)$ ) has_contour_integral
    I)
    (linepath (of\_real a) (of\_real b))
  by (simp add: has_contour_integral_def)
  also have ...  $\longleftrightarrow$  (( $\lambda t. f (c + r * cis t) * r * i * cis t$ ) has_integral I) {a..b}
using assms
  by (subst has_contour_integral_linepath_Reals_iff) (simp_all add: cis_conv_exp)
  finally show ?thesis .
qed

```

```

lemma contour_integrable_part_circlepath_iff:
  assumes a < b
  shows f contour_integrable_on (part_circlepath c r a b)  $\longleftrightarrow$ 
    ( $\lambda t. f (c + r * cis t) * r * i * cis t$ ) integrable_on {a..b}
  using assms by (auto simp: contour_integrable_on_def integrable_on_def
    has_contour_integral_part_circlepath_iff)

```

```

lemma contour_integral_part_circlepath_eq:
  assumes a < b
  shows contour_integral (part_circlepath c r a b) f =
    integral {a..b} ( $\lambda t. f (c + r * cis t) * r * i * cis t$ )
proof (cases f contour_integrable_on part_circlepath c r a b)
  case True
  hence ( $\lambda t. f (c + r * cis t) * r * i * cis t$ ) integrable_on {a..b}
  using assms by (simp add: contour_integrable_part_circlepath_iff)
  with True show ?thesis
  using has_contour_integral_part_circlepath_iff[OF assms]
    contour_integral_unique has_integral_integrable_integral by blast
next
  case False
  hence  $\neg(\lambda t. f (c + r * cis t) * r * i * cis t)$  integrable_on {a..b}

```



```

  using assms by (simp add: contour_integrable_part_circlepath_iff)
  with False show ?thesis
  by (simp add: not_integrable_contour_integral not_integrable_integral)
qed

```

```

lemma contour_integral_part_circlepath_reverse:
  contour_integral (part_circlepath c r a b) f = -contour_integral (part_circlepath
c r b a) f
  by (metis contour_integral_reversepath reversepath_part_circlepath valid_path_part_circlepath)

```

```

lemma contour_integral_part_circlepath_reverse':
  b < a  $\implies$  contour_integral (part_circlepath c r a b) f =
    -contour_integral (part_circlepath c r b a) f
  by (rule contour_integral_part_circlepath_reverse)

```

```

lemma finite_bounded_log: finite {z::complex. norm z  $\leq$  b  $\wedge$  exp z = w}
proof (cases w = 0)
  case True then show ?thesis by auto
next
  case False
  have *: finite {x. cmod ((2 * real_of_int x * pi) * i)  $\leq$  b + cmod (Ln w)}
  proof (simp add: norm_mult finite_int_iff_bounded_le)
    have abs ' {x. 2 * |real_of_int x| * pi  $\leq$  b + cmod (Ln w)}
       $\subseteq$  {..[(b + cmod (Ln w)) / (2 * pi)]}
    by (auto simp: field_split_simps le_floor_iff)
    then show  $\exists k. \text{abs ' } \{x. 2 * |\text{of\_int } x| * \text{pi} \leq b + \text{cmod } (\text{Ln } w)\} \subseteq \{..k\}$ 
      by blast
  qed
  have [simp]:  $\bigwedge P f. \{z. P z \wedge (\exists n. z = f n)\} = f ' \{n. P (f n)\}$ 
    by blast
  have finite {z. cmod z  $\leq$  b  $\wedge$  exp z = exp (Ln w)}
    using norm_add_leD by (fastforce intro: finite_subset [OF _ *] simp: exp_eq)
  then show ?thesis
    using False by auto
qed

```

```

lemma finite_bounded_log2:
  fixes a::complex
  assumes a  $\neq$  0
  shows finite {z. norm z  $\leq$  b  $\wedge$  exp(a*z) = w}
proof -
  have *: finite (( $\lambda z. z / a$ ) ' {z. cmod z  $\leq$  b * cmod a  $\wedge$  exp z = w})
    by (rule finite_imageI [OF finite_bounded_log])
  show ?thesis
    by (rule finite_subset [OF _ *]) (force simp: assms norm_mult)
qed

```

```

lemma has_contour_integral_bound_part_circlepath_strong:
  assumes fi: (f has_contour_integral i) (part_circlepath z r s t)

```

```

    and finite k and le: 0 ≤ B 0 < r s ≤ t
    and B:  $\bigwedge x. x \in \text{path\_image}(\text{part\_circlepath } z \ r \ s \ t) - k \implies \text{norm}(f \ x) \leq B$ 
    shows cmod i ≤ B * r * (t - s)
  proof -
    consider s = t | s < t using ⟨s ≤ t⟩ by linarith
    then show ?thesis
  proof cases
    case 1 with fi [unfolded has_contour_integral]
      have i = 0 by (simp add: vector_derivative_part_circlepath)
      with assms show ?thesis by simp
    next
      case 2
      have [simp]: |r| = r using ⟨r > 0⟩ by linarith
      have [simp]: cmod (of_real t - of_real s) = t - s
        by (metis 2 abs_of_pos diff_gt_0_iff_gt norm_of_real of_real_diff)
      have finite (part_circlepath z r s t - ' {y} ∩ {0..1}) if y ∈ k for y
      proof -
        let ?w = (y - z) / of_real r / exp(i * of_real s)
        have fin: finite (of_real - ' {z. cmod z ≤ 1 ∧ exp (i * of_real (t - s) * z) =
          ?w})
          using ⟨s < t⟩
          by (intro finite_vimageI [OF finite_bounded_log2]) (auto simp: inj_of_real)
        show ?thesis
          unfolding part_circlepath_def linepath_def vimage_def
          using le
          by (intro finite_subset [OF _ fin]) (auto simp: algebra_simps scaleR_conv_of_real
            exp_add exp_diff)
        qed
      then have fin01: finite ((part_circlepath z r s t) - ' k ∩ {0..1})
        by (rule finite_finite_vimage_IntI [OF ⟨finite k⟩])
      have **: (( $\lambda x. \text{if } (\text{part\_circlepath } z \ r \ s \ t \ x) \in k \text{ then } 0$ 
        else  $f(\text{part\_circlepath } z \ r \ s \ t \ x) *$ 
          vector_derivative (part_circlepath z r s t) (at x) has_integral
        i) {0..1})
        by (rule has_integral_spike [OF negligible_finite [OF fin01]]) (use fi has_contour_integral
          in auto)
      have *:  $\bigwedge x. \llbracket 0 \leq x; x \leq 1; \text{part\_circlepath } z \ r \ s \ t \ x \notin k \rrbracket \implies \text{cmod } (f$ 
        (part_circlepath z r s t x)) ≤ B
        by (auto intro!: B [unfolded path_image_def image_def])
      show ?thesis
        using has_integral_bound [where 'a=real, simplified, OF _ **]
        using assms le * 2 ⟨r > 0⟩ by (auto simp add: norm_mult vector_derivative_part_circlepath)
      qed
    qed
  qed

  corollary contour_integral_bound_part_circlepath_strong:
    assumes f contour_integrable_on part_circlepath z r s t
      and finite k and 0 ≤ B 0 < r s ≤ t
      and  $\bigwedge x. x \in \text{path\_image}(\text{part\_circlepath } z \ r \ s \ t) - k \implies \text{norm}(f \ x) \leq B$ 

```

shows $cmod (contour_integral (part_circlepath z r s t) f) \leq B * r * (t - s)$
using *assms has_contour_integral_bound_part_circlepath_strong has_contour_integral_integral*
by *blast*

lemma *has_contour_integral_bound_part_circlepath*:
 $\llbracket (f \text{ has_contour_integral } i) (part_circlepath z r s t);$
 $0 \leq B; 0 < r; s \leq t;$
 $\bigwedge x. x \in path_image(part_circlepath z r s t) \implies norm(f x) \leq B \rrbracket$
 $\implies norm i \leq B * r * (t - s)$
by (*auto intro: has_contour_integral_bound_part_circlepath_strong*)

lemma *contour_integrable_continuous_part_circlepath*:
 $continuous_on (path_image (part_circlepath z r s t)) f$
 $\implies f \text{ contour_integrable_on } (part_circlepath z r s t)$
unfolding *contour_integrable_on has_contour_integral_def vector_derivative_part_circlepath*
path_image_def
by (*best intro: integrable_continuous_real_path_part_circlepath [unfolded path_def]*)
continuous_intros
 $continuous_on_compose2 [where g=f, OF _ _ order_refl]$)

lemma *simple_path_part_circlepath*:
 $simple_path(part_circlepath z r s t) \longleftrightarrow (r \neq 0 \wedge s \neq t \wedge |s - t| \leq 2 * pi)$
proof (*cases r = 0 \vee s = t*)
case *True*
then show *?thesis*
unfolding *part_circlepath_def simple_path_def loop_free_def*
by (*rule disjE*) (*force intro: bexI [where x = 1/4] bexI [where x = 1/3]*)
next
case *False then have r \neq 0 s \neq t by auto*
have *: $\bigwedge x y z s t. i * ((1 - x) * s + x * t) = i * (((1 - y) * s + y * t)) + z \longleftrightarrow$
 $i * (x - y) * (t - s) = z$
by (*simp add: algebra_simps*)
have *abs01*: $\bigwedge x y :: real. 0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1$
 $\implies (x = y \vee x = 0 \wedge y = 1 \vee x = 1 \wedge y = 0 \longleftrightarrow |x - y| \in$
 $\{0, 1\})$
by *auto*
have **: $\bigwedge x y. (\exists n. (complex_of_real x - of_real y) * (of_real t - of_real s)$
 $= 2 * (of_int n * of_real pi)) \longleftrightarrow$
 $(\exists n. |x - y| * (t - s) = 2 * (of_int n * pi))$
by (*force simp: algebra_simps abs_if dest: arg_cong [where f=Re] arg_cong*)
 $[where f=complex_of_real]$
 $intro: exI [where x = -n for n]$)
have *1*: $|s - t| \leq 2 * pi$
if $\bigwedge x. 0 \leq x \wedge x \leq 1 \implies (\exists n. x * (t - s) = 2 * (real_of_int n * pi)) \longrightarrow x$
 $= 0 \vee x = 1$
proof (*rule ccontr*)
assume $\neg |s - t| \leq 2 * pi$
then have *: $\bigwedge n. t - s \neq of_int n * |s - t|$
using *False that [of 2*pi / |t - s|]*

```

    by (simp add: abs_minus_commute divide_simps)
  show False
    using * [of 1] * [of -1] by auto
qed
  have 2:  $|s - t| = |2 * (\text{real\_of\_int } n * \pi) / x|$  if  $x \neq 0$   $x * (t - s) = 2 * (\text{real\_of\_int } n * \pi)$  for  $x$ 
  proof -
    have  $t - s = 2 * (\text{real\_of\_int } n * \pi) / x$ 
      using that by (simp add: field_simps)
    then show ?thesis by (metis abs_minus_commute)
  qed
  have abs_away:  $\bigwedge P. (\forall x \in \{0..1\}. \forall y \in \{0..1\}. P |x - y|) \longleftrightarrow (\forall x :: \text{real}. 0 \leq x \wedge x < 1 \longrightarrow P x)$ 
  by force
  have  $\bigwedge x n. [|s \neq t; |s - t| \leq 2 * \pi; 0 \leq x; x < 1; x * (t - s) = 2 * (\text{real\_of\_int } n * \pi)] \implies x = 0$ 
  by (rule ccontr) (auto simp: 2 field_split_simps abs_mult dest: of_int_leD)
  then
  show ?thesis using False
    apply (simp add: simple_path_def loop_free_def)
    apply (simp add: part_circlepath_def linepath_def exp_eq ** abs01 del: Set.insert_iff)
    apply (subst abs_away)
    apply (auto simp: 1)
    done
qed

lemma arc_part_circlepath:
  assumes  $r \neq 0$   $s \neq t$   $|s - t| < 2 * \pi$ 
  shows arc (part_circlepath z r s t)
proof -
  have *:  $x = y$  if eq:  $i * (\text{linepath } s t x) = i * (\text{linepath } s t y) + 2 * \text{of\_int } n * \text{of\_real } \pi * i$ 
  and  $x: x \in \{0..1\}$  and  $y: y \in \{0..1\}$  for  $x y n$ 
  proof (rule ccontr)
    assume  $x \neq y$ 
    have  $(\text{linepath } s t x) = (\text{linepath } s t y) + 2 * \text{of\_int } n * \text{complex\_of\_real } \pi$ 
    by (metis add_divide_eq_iff complex_i_not_zero mult.commute nonzero_mult_div_cancel_left eq)
    then have  $s * y + t * x = s * x + (t * y + \text{of\_int } n * (\pi * 2))$ 
      by (force simp: algebra_simps linepath_def dest: arg_cong [where f=Re])
    with  $\langle x \neq y \rangle$  have st:  $s - t = (\text{of\_int } n * (\pi * 2) / (y - x))$ 
      by (force simp: field_simps)
    have  $|\text{real\_of\_int } n| < |y - x|$ 
      using assms  $\langle x \neq y \rangle$  by (simp add: st abs_mult field_simps)
    then show False
      using assms  $x y st$  by (auto dest: of_int_lessD)
  qed
qed

```

```

then have inj_on (part_circlepath z r s t) {0..1}
  using assms by (force simp add: part_circlepath_def inj_on_def exp_eq)
then show ?thesis
  by (simp add: arc_def)
qed

```

1.14 Special case of one complete circle

```

definition circlepath :: [complex, real, real]  $\Rightarrow$  complex
  where circlepath z r  $\equiv$  part_circlepath z r 0 (2*pi)

```

```

lemma circlepath: circlepath z r = ( $\lambda x$ . z + r * exp(2 * of_real pi * i * of_real x))
  by (simp add: circlepath_def part_circlepath_def linepath_def algebra_simps)

```

```

lemma pathstart_circlepath [simp]: pathstart (circlepath z r) = z + r
  by (simp add: circlepath_def)

```

```

lemma pathfinish_circlepath [simp]: pathfinish (circlepath z r) = z + r
  by (simp add: circlepath_def) (metis exp_two_pi_i mult.commute)

```

```

lemma circlepath_minus: circlepath z (-r) x = circlepath z r (x + 1/2)

```

proof –

```

  have z + of_real r * exp (2 * pi * i * (x + 1/2)) =
    z + of_real r * exp (2 * pi * i * x + pi * i)
    by (simp add: divide_simps) (simp add: algebra_simps)
  also have ... = z - r * exp (2 * pi * i * x)
    by (simp add: exp_add)
  finally show ?thesis
    by (simp add: circlepath_path_image_def sphere_def dist_norm)

```

qed

```

lemma circlepath_add1: circlepath z r (x+1) = circlepath z r x
  using circlepath_minus [of z r x+1/2] circlepath_minus [of z -r x]
  by (simp add: add.commute)

```

```

lemma circlepath_add_half: circlepath z r (x + 1/2) = circlepath z r (x - 1/2)
  using circlepath_add1 [of z r x-1/2]
  by (simp add: add.commute)

```

```

lemma path_image_circlepath_minus_subset:

```

```

  path_image (circlepath z (-r))  $\subseteq$  path_image (circlepath z r)

```

proof –

```

  have  $\exists x \in \{0..1\}$ . circlepath z r (y + 1/2) = circlepath z r x
    if  $0 \leq y \leq 1$  for y
  proof (cases y  $\leq$  1/2)
    case False
      with that show ?thesis
        by (force simp: circlepath_add_half)

```

```

qed (use that in force)
then show ?thesis
  by (auto simp add: path_image_def image_def circlepath_minus)
qed

```

```

lemma path_image_circlepath_minus: path_image (circlepath z (-r)) = path_image
(circlepath z r)
  using path_image_circlepath_minus_subset by fastforce

```

```

lemma has_vector_derivative_circlepath [derivative_intros]:
((circlepath z r) has_vector_derivative (2 * pi * i * r * exp (2 * of_real pi * i *
x)))
  (at x within X)
  unfolding circlepath_def scaleR_conv_of_real
  by (rule derivative_eq_intros) (simp add: algebra_simps)

```

```

lemma vector_derivative_circlepath:
vector_derivative (circlepath z r) (at x) =
  2 * pi * i * r * exp(2 * of_real pi * i * x)
  using has_vector_derivative_circlepath vector_derivative_at by blast

```

```

lemma vector_derivative_circlepath01:
[[0 ≤ x; x ≤ 1]]
  ⇒ vector_derivative (circlepath z r) (at x within {0..1}) =
    2 * pi * i * r * exp(2 * of_real pi * i * x)
  using has_vector_derivative_circlepath
  by (auto simp: vector_derivative_at_within_ivl)

```

```

lemma valid_path_circlepath [simp]: valid_path (circlepath z r)
  by (simp add: circlepath_def)

```

```

lemma path_circlepath [simp]: path (circlepath z r)
  by (simp add: valid_path_imp_path)

```

```

lemma path_image_circlepath_nonneg:
  assumes 0 ≤ r shows path_image (circlepath z r) = sphere z r
proof -
  have *: x ∈ (λu. z + (cmod (x - z)) * exp (i * (of_real u * (of_real pi * 2))))
  ‘{0..1} for x
  proof (cases x = z)
  case True then show ?thesis by force
  next
  case False
  define w where w = x - z
  then have w ≠ 0 by (simp add: False)
  have **: ∧t. [[Re w = cos t * cmod w; Im w = sin t * cmod w]] ⇒ w = of_real
(cmod w) * exp (i * t)
  using cis_conv_exp_complex_eq_iff by auto
  obtain t where 0 ≤ t t < 2*pi Re(w/norm w) = cos t Im(w/norm w) = sin t

```

```

  apply (rule sincos_total_2pi [of Re(w/(norm w)) Im(w/(norm w))])
  by (auto simp add: divide_simps ‹w ≠ 0› cmod_power2 [symmetric])
then
show ?thesis
  using False ** w_def ‹w ≠ 0›
  by (rule_tac x=t / (2*pi) in image_eqI) (auto simp add: field_simps)
qed
show ?thesis
  unfolding circlepath_path_image_def sphere_def dist_norm
  by (force simp: assms algebra_simps norm_mult norm_minus_commute intro:
*)
qed

```

```

lemma path_image_circlepath [simp]:
  path_image (circlepath z r) = sphere z |r|
  using path_image_circlepath_minus
  by (force simp: path_image_circlepath_nonneg abs_if)

```

```

lemma has_contour_integral_bound_circlepath_strong:
  [(f has_contour_integral i) (circlepath z r);
  finite k; 0 ≤ B; 0 < r;
  ∧x. [norm(x - z) = r; x ∉ k] ⇒ norm(f x) ≤ B]
  ⇒ norm i ≤ B*(2*pi*r)
  unfolding circlepath_def
  by (auto simp: algebra_simps in_path_image_part_circlepath dest!: has_contour_integral_bound_part_circlepath)

```

```

lemma has_contour_integral_bound_circlepath:
  [(f has_contour_integral i) (circlepath z r);
  0 ≤ B; 0 < r; ∧x. norm(x - z) = r ⇒ norm(f x) ≤ B]
  ⇒ norm i ≤ B*(2*pi*r)
  by (auto intro: has_contour_integral_bound_circlepath_strong)

```

```

lemma contour_integrable_continuous_circlepath:
  continuous_on (path_image (circlepath z r)) f
  ⇒ f contour_integrable_on (circlepath z r)
  by (simp add: circlepath_def contour_integrable_continuous_part_circlepath)

```

```

lemma simple_path_circlepath: simple_path(circlepath z r) ↔ (r ≠ 0)
  by (simp add: circlepath_def simple_path_part_circlepath)

```

```

lemma notin_path_image_circlepath [simp]: cmod (w - z) < r ⇒ w ∉ path_image
(circlepath z r)
  by (simp add: sphere_def dist_norm norm_minus_commute)

```

```

lemma contour_integral_circlepath:
  assumes r > 0
  shows contour_integral (circlepath z r) (λw. 1 / (w - z)) = 2 * of_real pi * i
  proof (rule contour_integral_unique)
  show ((λw. 1 / (w - z)) has_contour_integral 2 * of_real pi * i) (circlepath z

```

```

r)
  unfolding has_contour_integral_def using assms has_integral_const_real [of
_ 0 1]
  apply (subst has_integral_cong)
  apply (simp add: vector_derivative_circlepath01)
  apply (force simp: circlepath)
  done
qed

```

1.15 Uniform convergence of path integral

Uniform convergence when the derivative of the path is bounded, and in particular for the special case of a circle.

proposition *contour_integral_uniform_limit:*

```

assumes ev_fint: eventually ( $\lambda n::'a. (f\ n)$ ) contour_integrable_on  $\gamma$   $F$ 
and ul_f: uniform_limit (path_image  $\gamma$ )  $f$   $l$   $F$ 
and noleB:  $\bigwedge t. t \in \{0..1\} \implies \text{norm} (\text{vector\_derivative } \gamma (at\ t)) \leq B$ 
and  $\gamma$ : valid_path  $\gamma$ 
and [simp]:  $\neg$  trivial_limit  $F$ 
shows  $l$  contour_integrable_on  $\gamma$  (( $\lambda n. \text{contour\_integral } \gamma (f\ n)$ )  $\longrightarrow$  contour_integrable  $\gamma$   $l$ )  $F$ 

```

proof –

```

have  $0 \leq B$  by (meson noleB [of 0] atLeastAtMost_iff norm_ge_zero order_refl order_trans zero_le_one)

```

```

{ fix  $e::\text{real}$ 
assume  $0 < e$ 
then have  $0 < e / (|B| + 1)$  by simp
then have  $\S: \forall_F n \text{ in } F. \forall x \in \text{path\_image } \gamma. \text{cmod} (f\ n\ x - l\ x) < e / (|B| + 1)$ 
using ul_f [unfolded uniform_limit_iff dist_norm] by auto
obtain  $a$  where  $\text{fga}: \bigwedge x. x \in \{0..1\} \implies \text{cmod} (f\ a\ (\gamma\ x) - l\ (\gamma\ x)) < e / (|B| + 1)$ 
and  $\text{inta}: (\lambda t. f\ a\ (\gamma\ t) * \text{vector\_derivative } \gamma (at\ t)) \text{ integrable\_on } \{0..1\}$ 

```

```

using eventually_happens [OF eventually_conj [OF ev_fint  $\S$ ]]
by (fastforce simp: contour_integrable_on_path_image_def)
have  $\exists h. (\forall x \in \{0..1\}. \text{cmod} (l\ (\gamma\ x) * \text{vector\_derivative } \gamma (at\ x) - h\ x) \leq e) \wedge h \text{ integrable\_on } \{0..1\}$ 

```

```

proof (intro exI conjI ballI)
show  $\text{cmod} (l\ (\gamma\ x) * \text{vector\_derivative } \gamma (at\ x) - f\ a\ (\gamma\ x) * \text{vector\_derivative } \gamma (at\ x)) \leq e$ 

```

```

if  $x \in \{0..1\}$  for  $x$ 
proof –
have  $\text{cmod} (l\ (\gamma\ x) * \text{vector\_derivative } \gamma (at\ x) - f\ a\ (\gamma\ x) * \text{vector\_derivative } \gamma (at\ x)) \leq B * e / (|B| + 1)$ 

```

```

using noleB [OF that] fga [OF that]  $\langle 0 \leq B \rangle \langle 0 < e \rangle$ 
by (fastforce simp: mult_ac dest: mult_mono [OF less_imp_le] simp add: norm_mult left_diff_distrib [symmetric] norm_minus_commute divide_simps)
also have  $\dots \leq e$ 

```



```

      using ‹0 ≤ B› ‹0 < e› by (simp add: field_split_simps)
      finally show ?thesis .
    qed
  qed (rule inta)
}
then show lintg: l contour_integrable_on γ
  unfolding contour_integrable_on by (metis (mono_tags, lifting) integrable_uniform_limit_real)
{ fix e::real
  define B' where B' = B + 1
  have B': B' > 0 B' > B using ‹0 ≤ B› by (auto simp: B'_def)
  assume 0 < e
  then have ev_no': ∀F n in F. ∀ x ∈ path_image γ. 2 * cmod (f n x - l x) < e
/ B'
  using ul_f [unfolded uniform_limit_iff dist_norm, rule_format, of e / B'/2]
B'
  by (simp add: field_simps)
  have ie: integral {0..1::real} (λx. e/2) < e using ‹0 < e› by simp
  have *: cmod (f x (γ t) * vector_derivative γ (at t) - l (γ t) * vector_derivative
γ (at t)) ≤ e/2
    if t: t ∈ {0..1} and leB': 2 * cmod (f x (γ t) - l (γ t)) < e / B' for x t
  proof -
    have 2 * cmod (f x (γ t) - l (γ t)) * cmod (vector_derivative γ (at t)) ≤ e
* (B / B')
    using mult_mono [OF less_imp_le [OF leB'] noleB] B' ‹0 < e› t by auto
    also have ... < e
    by (simp add: B' ‹0 < e› mult_imp_div_pos_less)
    finally have 2 * cmod (f x (γ t) - l (γ t)) * cmod (vector_derivative γ (at
t)) < e .
  then show ?thesis
  by (simp add: left_diff_distrib [symmetric] norm_mult)
  qed
  have le_e: ∧x. [∀ u ∈ {0..1}. 2 * cmod (f x (γ u) - l (γ u)) < e / B'; f x
contour_integrable_on γ]
  ⇒ cmod (integral {0..1}
(λu. f x (γ u) * vector_derivative γ (at u) - l (γ u) * vector_derivative
γ (at u))) < e
  apply (rule le_less_trans [OF integral_norm_bound_integral ie])
  apply (simp add: lintg integrable_diff contour_integrable_on [symmetric])
  apply (blast intro: *)+
  done
  have ∀F x in F. dist (contour_integral γ (f x)) (contour_integral γ l) < e
  apply (rule eventually_mono [OF eventually_conj [OF ev_no' ev_fint]])
  apply (simp add: dist_norm contour_integrable_on path_image_def con-
tour_integral_integral)
  apply (simp add: lintg integral_diff [symmetric] contour_integrable_on
[symmetric] le_e)
  done
}
}
then show ((λn. contour_integral γ (f n)) → contour_integral γ l) F

```

by (rule tendstoI)
qed

corollary *contour_integral_uniform_limit_circlepath*:

assumes $\forall_F n::'a \text{ in } F. (f \ n) \text{ contour_integrable_on } (\text{circlepath } z \ r)$
and *uniform_limit* (sphere z r) f l F
and $\neg \text{trivial_limit } F \ 0 < r$
shows $l \text{ contour_integrable_on } (\text{circlepath } z \ r)$
 $((\lambda n. \text{contour_integral } (\text{circlepath } z \ r) (f \ n)) \longrightarrow \text{contour_integral } (\text{circlepath } z \ r) l) \ F$
using *assms* by (auto simp: vector_derivative_circlepath norm_mult intro!: contour_integral_uniform_limit)

lemma *has_contour_integral_linepath_same_Re_iff*:

assumes $\text{Re } z = c \ \text{Re } z' = c \ \text{Im } z = a \ \text{Im } z' = b \ a < b$
shows $(f \text{ has_contour_integral } I) (\text{linepath } z \ z') \longleftrightarrow$
 $((\lambda x. f (\text{Complex } c \ x)) \text{ has_integral } (-i * I)) \{a..b\}$

proof –

have $(f \text{ has_contour_integral } I) (\text{linepath } z \ z') \longleftrightarrow$
 $((\lambda x. f (\text{linepath } z \ z' \ x) * (z' - z)) \text{ has_integral } I) \{0..1\}$
by (subst *has_contour_integral_linepath*) *simp_all*
also have $\dots \longleftrightarrow ((\lambda x. f (c + (a + (b - a) * x) * i) * (i * (b - a)))$
has_integral I) $\{0..1\}$
using *assms*
by (intro *has_integral_cong_arg_cong2*[of $_ _ _ (*)$] *arg_cong*[of $_ _ f$])
(auto simp: *linepath_def* *complex_eq_iff_algebra_simps*)
also have $\{0..1\} = (\lambda x. x / (b - a)) \ ' \{0..b-a\}$
using *assms* by *simp*
also have $((\lambda x. f (c + (a + (b-a) * x) * i) * (i * (b-a))) \text{ has_integral } I) \dots$
 \longleftrightarrow
 $((\lambda x. f (c + (a + x) * i) * (i * (b-a))) \text{ has_integral } ((b-a) * i))$
 $\{0..b-a\}$
by (subst *has_integral_stretch_real_iff*) (use *assms* in *simp_all*)
also have $\dots \longleftrightarrow ((\lambda x. \text{of_real } (b-a) * i * (f (c + x * i))) \text{ has_integral } (b-a)$
 $* i) \{a..b\}$
by (subst *has_integral_shift_real_ivl_iff*[**where** $c = -a$])
(*simp_all* add: *scaleR_conv_of_real_mult_ac*)
also have $\dots \longleftrightarrow ((\lambda x. f (c + x * i) \text{ has_integral } (-i * I)) \{a..b\}$
by (subst *has_integral_mult_right_iff*) (use *assms* in \langle auto simp: *scaleR_conv_of_real* \rangle)
finally show *?thesis*
by (*simp* add: *scaleR_conv_of_real* *Complex_eq* *mult.commute*)
qed

end

2 Complex Path Integrals and Cauchy's Integral Theorem

By John Harrison et al. Ported from HOL Light by L C Paulson (2015)

theory *Cauchy_Integral_Theorem*

imports

HOL-Analysis.Analysis

Contour_Integration

begin

lemma *leibniz_rule_holomorphic*:

fixes $f::\text{complex} \Rightarrow 'b::\text{euclidean_space} \Rightarrow \text{complex}$

assumes $\bigwedge x t. x \in U \implies t \in \text{cbox } a \ b \implies ((\lambda x. f \ x \ t) \text{ has_field_derivative } f \ x \ t) \text{ (at } x \text{ within } U)$

assumes $\bigwedge x. x \in U \implies (f \ x) \text{ integrable_on } \text{cbox } a \ b$

assumes *continuous_on* $(U \times (\text{cbox } a \ b)) (\lambda(x, t). f \ x \ t)$

assumes *convex* U

shows $(\lambda x. \text{integral } (\text{cbox } a \ b) (f \ x)) \text{ holomorphic_on } U$

using *leibniz_rule_field_differentiable*[*OF* *assms*(1-3) *_* *assms*(4)]

by (*auto simp: holomorphic_on_def*)

lemma *Ln_measurable* [*measurable*]: $Ln \in \text{measurable borel borel}$

proof –

have $*$: $Ln \ (-\text{of_real } x) = \text{of_real } (\ln \ x) + i * \pi \text{ if } x > 0 \text{ for } x$

using *that* **by** (*subst Ln_minus*) (*auto simp: Ln_of_real*)

have $**$: $Ln \ (\text{of_real } x) = \text{of_real } (\ln \ (-x)) + i * \pi \text{ if } x < 0 \text{ for } x$

using $*$ [*of -x*] *that* **by** *simp*

have *cont*: $(\lambda x. \text{indicat_real } (-\mathbb{R}_{\leq 0}) \ x *_{\mathbb{R}} Ln \ x) \in \text{borel_measurable borel}$

by (*intro borel_measurable_continuous_on_indicator continuous_intros*) *auto*

have $(\lambda x. \text{if } x \in \mathbb{R}_{\leq 0} \text{ then } \ln \ (-\text{Re } x) + i * \pi \text{ else } \text{indicator } (-\mathbb{R}_{\leq 0}) \ x *_{\mathbb{R}} Ln \ x) \in \text{borel } \rightarrow_M \text{ borel}$

(*is ?f* \in *_*) **by** (*rule measurable>If_set*[*OF* *_* *cont*]) *auto*

hence $(\lambda x. \text{if } x = 0 \text{ then } Ln \ 0 \text{ else } ?f \ x) \in \text{borel } \rightarrow_M \text{ borel}$ **by** *measurable*

also have $(\lambda x. \text{if } x = 0 \text{ then } Ln \ 0 \text{ else } ?f \ x) = Ln$

by (*auto simp: fun_eq_iff ** nonpos_Reals_def*)

finally show *?thesis* .

qed

lemma *powr_complex_measurable* [*measurable*]:

assumes [*measurable*]: $f \in \text{measurable } M \text{ borel } g \in \text{measurable } M \text{ borel}$

shows $(\lambda x. f \ x \ \text{powr } g \ x :: \text{complex}) \in \text{measurable } M \text{ borel}$

using *assms* **by** (*simp add: powr_def*)

The special case of midpoints used in the main quadrisection

lemma *has_contour_integral_midpoint*:

assumes (*f has_contour_integral* *i*) (*linepath* *a* (*midpoint* *a* *b*))

(*f has_contour_integral* *j*) (*linepath* (*midpoint* *a* *b*) *b*)

shows (*f has_contour_integral* (*i* + *j*)) (*linepath* *a* *b*)

```

proof (rule has_contour_integral_split)
  show midpoint a b - a = (1/2) *R (b - a)
  using assms by (auto simp: midpoint_def scaleR_conv_of_real)
qed (use assms in auto)

```

```

lemma contour_integral_midpoint:
  assumes continuous_on (closed_segment a b) f
  shows contour_integral (linepath a b) f =
    contour_integral (linepath a (midpoint a b)) f + contour_integral (linepath
(midpoint a b) b) f
proof (rule contour_integral_split)
  show midpoint a b - a = (1/2) *R (b - a)
  using assms by (auto simp: midpoint_def scaleR_conv_of_real)
qed (use assms in auto)

```

A couple of special case lemmas that are useful below

```

lemma triangle_linear_has_chain_integral:
  (( $\lambda x. m*x + d$ ) has_contour_integral 0) (linepath a b +++ linepath b c +++
linepath c a)
proof (rule Cauchy_theorem_primitive)
  show  $\bigwedge x. x \in UNIV \implies ((\lambda x. m / 2 * x^2 + d * x)$  has_field_derivative  $m *
x + d$ ) (at x)
  by (auto intro!: derivative_eq_intros)
qed auto

```

```

lemma has_chain_integral_chain_integral3:
  assumes (f has_contour_integral i) (linepath a b +++ linepath b c +++ linepath
c d)
  (is (f has_contour_integral i) ?g)
  shows contour_integral (linepath a b) f + contour_integral (linepath b c) f +
contour_integral (linepath c d) f = i
  (is ?lhs = _)
proof -
  have f contour_integrable_on ?g
  using assms contour_integrable_on_def by blast
  then have ?lhs = contour_integral ?g f
  by (simp add: valid_path_join has_contour_integral_integrable)
  then show ?thesis
  using assms contour_integral_unique by blast
qed

```

```

lemma has_chain_integral_chain_integral4:
  assumes (f has_contour_integral i) (linepath a b +++ linepath b c +++ linepath
c d +++ linepath d e)
  (is (f has_contour_integral i) ?g)
  shows contour_integral (linepath a b) f + contour_integral (linepath b c) f +
contour_integral (linepath c d) f + contour_integral (linepath d e) f = i
  (is ?lhs = _)
proof -

```

```

have f contour_integrable_on ?g
  using assms contour_integrable_on_def by blast
then have ?lhs = contour_integral ?g f
  by (simp add: valid_path_join_has_contour_integral_integrable)
then show ?thesis
  using assms contour_integral_unique by blast
qed

```

2.1 The key quadrisection step

```

lemma norm_sum_half:
  assumes norm(a + b) ≥ e
  shows norm a ≥ e/2 ∨ norm b ≥ e/2
proof -
  have e ≤ norm (- a - b)
    by (simp add: add.commute assms norm_minus_commute)
  thus ?thesis
    using norm_triangle_ineq4 order_trans by fastforce
qed

```

```

lemma norm_sum_lemma:
  assumes e ≤ norm (a + b + c + d)
  shows e / 4 ≤ norm a ∨ e / 4 ≤ norm b ∨ e / 4 ≤ norm c ∨ e / 4 ≤ norm d
proof -
  have e ≤ norm ((a + b) + (c + d)) using assms
    by (simp add: algebra_simps)
  then show ?thesis
    by (auto dest!: norm_sum_half)
qed

```

```

lemma Cauchy_theorem_quadrisection:
  assumes f: continuous_on (convex hull {a,b,c}) f
  and dist: dist a b ≤ K dist b c ≤ K dist c a ≤ K
  and e: e * K^2 ≤
    norm (contour_integral(linepath a b) f + contour_integral(linepath b
c) f + contour_integral(linepath c a) f)
  shows ∃ a' b' c'.
    a' ∈ convex hull {a,b,c} ∧ b' ∈ convex hull {a,b,c} ∧ c' ∈ convex hull
{a,b,c} ∧
    dist a' b' ≤ K/2 ∧ dist b' c' ≤ K/2 ∧ dist c' a' ≤ K/2 ∧
    e * (K/2)^2 ≤ norm(contour_integral(linepath a' b') f + contour_integral(linepath
b' c') f + contour_integral(linepath c' a') f)
    (is ∃ x y z. ?Φ x y z)
proof -
  note divide_le_eq_numeral1 [simp del]
  define a' where a' = midpoint b c
  define b' where b' = midpoint c a
  define c' where c' = midpoint a b
  have fabc: continuous_on (closed_segment a b) f continuous_on (closed_segment

```

```

b c) f continuous_on (closed_segment c a) f
  using f continuous_on_subset segments_subset_convex_hull by metis+
have fcont': continuous_on (closed_segment c' b') f
  continuous_on (closed_segment a' c') f
  continuous_on (closed_segment b' a') f
  unfolding a'_def b'_def c'_def
  by (rule continuous_on_subset [OF f],
      metis midpoints_in_convex_hull convex_hull_subset hull_subset in-
sert_subset segment_convex_hull)+
define pathint where pathint x y  $\equiv$  contour_integral(linepath x y) f for x y
have *: pathint a b + pathint b c + pathint c a =
  (pathint a c' + pathint c' b' + pathint b' a) +
  (pathint a' c' + pathint c' b + pathint b a') +
  (pathint a' c + pathint c b' + pathint b' a') +
  (pathint a' b' + pathint b' c' + pathint c' a')
  unfolding pathint_def
  by (simp add: fcont' contour_integral_reverse_linepath) (simp add: a'_def
b'_def c'_def contour_integral_midpoint fabc)
have [simp]:  $\bigwedge x y. \text{cmod } (x * 2 - y * 2) = \text{cmod } (x - y) * 2$ 
  by (metis left_diff_distrib mult.commute norm_mult_numerals)
have [simp]:  $\bigwedge x y. \text{cmod } (x - y) = \text{cmod } (y - x)$ 
  by (simp add: norm_minus_commute)
consider e * K2 / 4  $\leq$  cmod (pathint a c' + pathint c' b' + pathint b' a) |
  e * K2 / 4  $\leq$  cmod (pathint a' c' + pathint c' b + pathint b a') |
  e * K2 / 4  $\leq$  cmod (pathint a' c + pathint c b' + pathint b' a') |
  e * K2 / 4  $\leq$  cmod (pathint a' b' + pathint b' c' + pathint c' a')
  using assms by (metis * norm_sum_lemma pathint_def)
then show ?thesis
proof cases
  case 1 then have ? $\Phi$  a c' b'
    using assms unfolding pathint_def [symmetric]
    apply (clarsimp simp: c'_def b'_def midpoints_in_convex_hull hull_subset
[THEN subsetD])
    apply (auto simp: midpoint_def dist_norm scaleR_conv_of_real field_split_simps)
    done
  then show ?thesis by blast
next
  case 2 then have ? $\Phi$  a' c' b
    using assms unfolding pathint_def [symmetric]
    apply (clarsimp simp: a'_def c'_def midpoints_in_convex_hull hull_subset
[THEN subsetD])
    apply (auto simp: midpoint_def dist_norm scaleR_conv_of_real field_split_simps)
    done
  then show ?thesis by blast
next
  case 3 then have ? $\Phi$  a' c b'
    using assms unfolding pathint_def [symmetric]
    apply (clarsimp simp: a'_def b'_def midpoints_in_convex_hull hull_subset
[THEN subsetD])

```

```

apply (auto simp: midpoint_def dist_norm scaleR_conv_of_real field_split_simps)
done
then show ?thesis by blast
next
case 4 then have ? $\Phi$  a' b' c'
  using assms unfolding pathint_def [symmetric]
  apply (clarsimp simp: a'_def c'_def b'_def midpoints_in_convex_hull
hull_subset [THEN subsetD])
  apply (auto simp: midpoint_def dist_norm scaleR_conv_of_real field_split_simps)
  done
then show ?thesis by blast
qed
qed

```

2.2 Cauchy's theorem for triangles

lemma triangle_points_closer:

```

fixes a::complex
shows  $\llbracket x \in \text{convex hull } \{a,b,c\}; y \in \text{convex hull } \{a,b,c\} \rrbracket$ 
 $\implies \text{norm}(x - y) \leq \text{norm}(a - b) \vee$ 
 $\text{norm}(x - y) \leq \text{norm}(b - c) \vee$ 
 $\text{norm}(x - y) \leq \text{norm}(c - a)$ 
using simplex_extremal_le [of {a,b,c}]
by (auto simp: norm_minus_commute)

```

lemma holomorphic_point_small_triangle:

```

assumes x:  $x \in S$ 
and f: continuous_on S f
and cd: f field_differentiable (at x within S)
and e:  $0 < e$ 
shows  $\exists k > 0. \forall a b c. \text{dist } a b \leq k \wedge \text{dist } b c \leq k \wedge \text{dist } c a \leq k \wedge$ 
 $x \in \text{convex hull } \{a,b,c\} \wedge \text{convex hull } \{a,b,c\} \subseteq S$ 
 $\longrightarrow \text{norm}(\text{contour\_integral}(\text{linepath } a b) f + \text{contour\_integral}(\text{linepath}$ 
 $b c) f +$ 
 $\text{contour\_integral}(\text{linepath } c a) f)$ 
 $\leq e * (\text{dist } a b + \text{dist } b c + \text{dist } c a)^2$ 
(is  $\exists k > 0. \forall a b c. \_ \longrightarrow ?\text{normle } a b c)$ 

```

proof –

```

have le_of_3:  $\bigwedge a x y z. \llbracket 0 \leq x*y; 0 \leq x*z; 0 \leq y*z; a \leq (e*(x + y + z))*x$ 
 $+ (e*(x + y + z))*y + (e*(x + y + z))*z \rrbracket$ 
 $\implies a \leq e*(x + y + z)^2$ 

```

by (simp add: algebra_simps power2_eq_square)

```

have disj_le:  $\llbracket x \leq a \vee x \leq b \vee x \leq c; 0 \leq a; 0 \leq b; 0 \leq c \rrbracket \implies x \leq a + b + c$ 
for x::real and a b c

```

by linarith

```

have abc: f contour_integrable_on linepath a b f contour_integrable_on linepath
b c f contour_integrable_on linepath c a
if convex_hull {a, b, c}  $\subseteq S$  for a b c

```

```

using segments_subset_convex_hull that
by (metis continuous_on_subset f contour_integrable_continuous_linepath)+
note path_bound = has_contour_integral_bound_linepath [simplified norm_minus_commute,
OF has_contour_integral_integral]
{ fix f' a b c d
  assume d: 0 < d
    and f':  $\bigwedge y. \llbracket c \text{ mod } (y - x) \leq d; y \in S \rrbracket \implies c \text{ mod } (f y - f x - f' * (y - x))$ 
  }
   $\leq e * c \text{ mod } (y - x)$ 
    and le:  $c \text{ mod } (a - b) \leq d \text{ mod } (b - c) \leq d \text{ mod } (c - a) \leq d$ 
    and xc:  $x \in \text{convex hull } \{a, b, c\}$ 
    and S:  $\text{convex hull } \{a, b, c\} \subseteq S$ 
  have pa:  $\text{contour\_integral } (\text{linepath } a \ b) \ f + \text{contour\_integral } (\text{linepath } b \ c) \ f$ 
+  $\text{contour\_integral } (\text{linepath } c \ a) \ f =$ 
     $\text{contour\_integral } (\text{linepath } a \ b) \ (\lambda y. f y - f x - f' * (y - x)) +$ 
     $\text{contour\_integral } (\text{linepath } b \ c) \ (\lambda y. f y - f x - f' * (y - x)) +$ 
     $\text{contour\_integral } (\text{linepath } c \ a) \ (\lambda y. f y - f x - f' * (y - x))$ 
  apply (simp add: contour_integral_diff contour_integral_lmul contour_integrable_lmul
contour_integrable_diff fabc [OF S])
  apply (simp add: field_simps)
  done
{ fix y
  assume yc:  $y \in \text{convex hull } \{a, b, c\}$ 
  have  $c \text{ mod } (f y - f x - f' * (y - x)) \leq e * \text{norm}(y - x)$ 
  proof (rule f')
    show  $c \text{ mod } (y - x) \leq d$ 
    by (metis triangle_points_closer [OF xc yc] le norm_minus_commute
order_trans)
  qed (use S yc in blast)
  also have  $\dots \leq e * (c \text{ mod } (a - b) + c \text{ mod } (b - c) + c \text{ mod } (c - a))$ 
  by (simp add: yc e xc disj_le [OF triangle_points_closer])
  finally have  $c \text{ mod } (f y - f x - f' * (y - x)) \leq e * (c \text{ mod } (a - b) + c \text{ mod } (b - c) + c \text{ mod } (c - a))$  .
  } note cm_le = this
  have ?normle a b c
  unfolding dist_norm pa
  using f' xc S e
  apply (intro le_of_3 norm_triangle_le add_mono path_bound)
  apply (simp_all add: contour_integral_diff contour_integral_lmul contour_integrable_lmul contour_integrable_diff fabc)
  apply (blast intro: cm_le elim: dest: segments_subset_convex_hull [THEN subsetD])+
  done
} note * = this
show ?thesis
  using cd e
  apply (simp add: field_differentiable_def has_field_derivative_def has_derivative_within_alt
approachable_lt_le2 Ball_def)
  apply (clarify dest!: spec mp)
  using * unfolding dist_norm

```



```

  apply blast
  done
qed

```

Hence the most basic theorem for a triangle.

```

locale Chain =
  fixes x0 At Follows
  assumes At0: At x0 0
    and AtSuc:  $\bigwedge x n. \text{At } x n \implies \exists x'. \text{At } x' (\text{Suc } n) \wedge \text{Follows } x' x$ 
begin
  primrec f where
    f 0 = x0
  | f (Suc n) = (SOME x. At x (Suc n) \wedge Follows x (f n))

  lemma At: At (f n) n
  proof (induct n)
    case 0 show ?case
    by (simp add: At0)
  next
    case (Suc n) show ?case
    by (metis (no_types, lifting) AtSuc [OF Suc] f.simps(2) someI_ex)
  qed

  lemma Follows: Follows (f(Suc n)) (f n)
  by (metis (no_types, lifting) AtSuc [OF At [of n]] f.simps(2) someI_ex)

  declare f.simps(2) [simp del]
end

```

```

lemma Chain3:
  assumes At0: At x0 y0 z0 0
    and AtSuc:  $\bigwedge x y z n. \text{At } x y z n \implies \exists x' y' z'. \text{At } x' y' z' (\text{Suc } n) \wedge \text{Follows } x' y' z' x y z$ 
  obtains f g h where
    f 0 = x0 g 0 = y0 h 0 = z0
     $\bigwedge n. \text{At } (f n) (g n) (h n) n$ 
     $\bigwedge n. \text{Follows } (f(\text{Suc } n)) (g(\text{Suc } n)) (h(\text{Suc } n)) (f n) (g n) (h n)$ 
  proof –
  interpret three: Chain (x0,y0,z0)  $\lambda(x,y,z). \text{At } x y z \lambda(x',y',z'). \lambda(x,y,z). \text{Follows } x' y' z' x y z$ 
  proof qed (use At0 AtSuc in auto)
  show ?thesis
  proof
    show  $\bigwedge n. \text{Follows } (fst (three.f (\text{Suc } n))) (fst (snd (three.f (\text{Suc } n))))$ 
      (snd (snd (three.f (\text{Suc } n)))) (fst (three.f n)))
      (fst (snd (three.f n)) (snd (snd (three.f n))))
     $\bigwedge n. \text{At } (fst (three.f n)) (fst (snd (three.f n))) (snd (snd (three.f n))) n$ 
    using three.At three.Follows
    by (simp_all add: split_beta')
  qed

```

qed auto
qed

proposition *Cauchy_theorem_triangle:*

assumes f holomorphic_on (convex hull $\{a,b,c\}$)

shows (f has_contour_integral 0) (linepath a b +++ linepath b c +++ linepath c a)

proof –

have $contf$: continuous_on (convex hull $\{a,b,c\}$) f

by (metis assms holomorphic_on_imp_continuous_on)

let $?pathint = \lambda x y. contour_integral(linepath\ x\ y)\ f$

{ fix $y::complex$

assume fy : (f has_contour_integral y) (linepath a b +++ linepath b c +++ linepath c a)

and ynz : $y \neq 0$

define K where $K = 1 + \max(\text{dist } a\ b) (\max(\text{dist } b\ c) (\text{dist } c\ a))$

define e where $e = \text{norm } y / K^2$

have $K1$: $K \geq 1$ by (simp add: K_def max.coboundedI1)

then have K : $K > 0$ by linarith

have [iff]: $\text{dist } a\ b \leq K \text{ dist } b\ c \leq K \text{ dist } c\ a \leq K$

by (simp_all add: K_def)

have e : $e > 0$

unfolding e_def using $ynz\ K1$ by simp

define At where $At\ x\ y\ z\ n \longleftrightarrow$

$\text{convex_hull } \{x,y,z\} \subseteq \text{convex_hull } \{a,b,c\} \wedge$

$\text{dist } x\ y \leq K/2^n \wedge \text{dist } y\ z \leq K/2^n \wedge \text{dist } z\ x \leq K/2^n \wedge$

$\text{norm}(?pathint\ x\ y + ?pathint\ y\ z + ?pathint\ z\ x) \geq e*(K/2^n)^2$

for $x\ y\ z\ n$

have $At0$: $At\ a\ b\ c\ 0$

using fy

by (simp add: At_def e_def has_chain_integral_chain_integral3)

{ fix $x\ y\ z\ n$

assume At : $At\ x\ y\ z\ n$

then have $contf'$: continuous_on (convex hull $\{x,y,z\}$) f

using $contf\ At_def$ continuous_on_subset by metis

have $\exists x'\ y'\ z'. At\ x'\ y'\ z' (Suc\ n) \wedge \text{convex_hull } \{x',y',z'\} \subseteq \text{convex_hull } \{x,y,z\}$

using At Cauchy_theorem_quadrisection [OF $contf'$, of $K/2^n\ e$]

apply (simp add: At_def algebra_simps)

apply (meson convex_hull_subset empty_subsetI insert_subset subsetCE)

done

} note $AtSuc = this$

obtain $fa\ fb\ fc$

where $f0$ [simp]: $fa\ 0 = a\ fb\ 0 = b\ fc\ 0 = c$

and $cosb$: $\bigwedge n. \text{convex_hull } \{fa\ n, fb\ n, fc\ n\} \subseteq \text{convex_hull } \{a,b,c\}$

and $dist$: $\bigwedge n. \text{dist } (fa\ n)\ (fb\ n) \leq K/2^n$

$\bigwedge n. \text{dist } (fb\ n)\ (fc\ n) \leq K/2^n$

$\bigwedge n. \text{dist } (fc\ n)\ (fa\ n) \leq K/2^n$

```

    and no:  $\bigwedge n. \text{norm}(\text{?pathint } (fa\ n) (fb\ n) +
      \text{?pathint } (fb\ n) (fc\ n) +
      \text{?pathint } (fc\ n) (fa\ n)) \geq e * (K/2^n)^2$ 
    and conv_le:  $\bigwedge n. \text{convex\_hull } \{fa(Suc\ n), fb(Suc\ n), fc(Suc\ n)\} \subseteq \text{convex\_hull } \{fa\ n, fb\ n, fc\ n\}$ 
  by (rule Chain3 [of At, OF At0 AtSuc]) (auto simp: At_def)
  obtain x where x:  $\bigwedge n. x \in \text{convex\_hull } \{fa\ n, fb\ n, fc\ n\}$ 
  proof (rule bounded_closed_nest)
    show  $\bigwedge n. \text{closed } (\text{convex\_hull } \{fa\ n, fb\ n, fc\ n\})$ 
      by (simp add: compact_imp_closed finite_imp_compact convex_hull)
    show  $\bigwedge m\ n. m \leq n \implies \text{convex\_hull } \{fa\ n, fb\ n, fc\ n\} \subseteq \text{convex\_hull } \{fa\ m, fb\ m, fc\ m\}$ 
      by (erule transitive_stepwise_le) (auto simp: conv_le)
  qed (fastforce intro: finite_imp_bounded_convex_hull)
  then have xin:  $x \in \text{convex\_hull } \{a,b,c\}$ 
  using assms f0 by blast
  then have fx:  $f \text{ field\_differentiable at } x \text{ within } (\text{convex\_hull } \{a,b,c\})$ 
  using assms holomorphic_on_def by blast
  { fix k n
    assume k:  $0 < k$ 
    and le:
       $\bigwedge x' y' z'. \llbracket \text{dist } x' y' \leq k; \text{dist } y' z' \leq k; \text{dist } z' x' \leq k;
        x \in \text{convex\_hull } \{x',y',z'\};
        \text{convex\_hull } \{x',y',z'\} \subseteq \text{convex\_hull } \{a,b,c\} \rrbracket$ 
       $\implies$ 
       $\text{cmod } (\text{?pathint } x' y' + \text{?pathint } y' z' + \text{?pathint } z' x') * 10$ 
       $\leq e * (\text{dist } x' y' + \text{dist } y' z' + \text{dist } z' x')^2$ 
    and Kk:  $K / k < 2^n$ 
    have  $K / 2^n < k$  using Kk k
    by (auto simp: field_simps)
    then have DD:  $\text{dist } (fa\ n) (fb\ n) \leq k \text{ dist } (fb\ n) (fc\ n) \leq k \text{ dist } (fc\ n) (fa\ n)$ 
     $\leq k$ 
    using dist [of n] k
    by linarith+
    have dle:  $(\text{dist } (fa\ n) (fb\ n) + \text{dist } (fb\ n) (fc\ n) + \text{dist } (fc\ n) (fa\ n))^2$ 
       $\leq (3 * K / 2^n)^2$ 
    using dist [of n] e K
    by (simp add: abs_le_square_iff [symmetric])
    have less10:  $\bigwedge x\ y::\text{real}. 0 < x \implies y \leq 9*x \implies y < x*10$ 
    by linarith
    have  $e * (\text{dist } (fa\ n) (fb\ n) + \text{dist } (fb\ n) (fc\ n) + \text{dist } (fc\ n) (fa\ n))^2 \leq e * (3 * K / 2^n)^2$ 
    using ynz dle e mult_le_cancel_left_pos by blast
    also have ... <
       $\text{cmod } (\text{?pathint } (fa\ n) (fb\ n) + \text{?pathint } (fb\ n) (fc\ n) + \text{?pathint } (fc\ n) (fa\ n)) * 10$ 
    using no [of n] e K
    by (simp add: e_def field_simps) (simp only: zero_less_norm_iff [symmetric])
  }

```

```

    finally have False
      using le [OF DD x cosb] by auto
  } then
  have ?thesis
    using holomorphic_point_small_triangle [OF xin contf fx, of e/10] e
    apply clarsimp
    apply (rule_tac y1=K/k in exE [OF real_arch_pow[of 2]], force+)
    done
  }
  moreover have f contour_integrable_on (linepath a b +++ linepath b c +++
linepath c a)
    by simp (meson contf continuous_on_subset contour_integrable_continuous_linepath
segments_subset_convex_hull(1)
      segments_subset_convex_hull(3) segments_subset_convex_hull(5))
  ultimately show ?thesis
    using has_contour_integral_integral by fastforce
qed

```

2.3 Version needing function holomorphic in interior only

lemma *Cauchy_theorem_flat_lemma*:

```

  assumes f: continuous_on (convex_hull {a,b,c}) f
    and c:  $c - a = k *_{\mathbb{R}} (b - a)$ 
    and k:  $0 \leq k$ 
  shows  $\text{contour\_integral } (\text{linepath } a \ b) \ f + \text{contour\_integral } (\text{linepath } b \ c) \ f +$ 
 $\text{contour\_integral } (\text{linepath } c \ a) \ f = 0$ 
proof -
  have fab: continuous_on (closed_segment a b) f continuous_on (closed_segment
b c) f continuous_on (closed_segment c a) f
    using f continuous_on_subset segments_subset_convex_hull by metis+
  show ?thesis
  proof (cases  $k \leq 1$ )
    case True show ?thesis
      by (simp add: contour_integral_split [OF fab(1) k True c] contour_integral_reverse_linepath
fab)
    next
    case False
      show ?thesis
      proof (subst contour_integral_split [symmetric])
        show  $b - a = (1/k) *_{\mathbb{R}} (c - a)$ 
          using False c by force
        show  $\text{contour\_integral } (\text{linepath } a \ c) \ f + \text{contour\_integral } (\text{linepath } c \ a) \ f =$ 
 $0$ 
          by (simp add: contour_integral_reverse_linepath fab(3))
        show continuous_on (closed_segment a c) f
          by (metis closed_segment_commute fab(3))
      qed (use False in auto)
    qed
  qed
qed

```

```

lemma Cauchy_theorem_flat:
  assumes f: continuous_on (convex_hull {a,b,c}) f
    and c: c - a = k *R (b - a)
    shows contour_integral (linepath a b) f +
      contour_integral (linepath b c) f +
      contour_integral (linepath c a) f = 0
proof (cases 0 ≤ k)
  case True with assms show ?thesis
    by (blast intro: Cauchy_theorem_flat_lemma)
next
  case False
    have continuous_on (closed_segment a b) f continuous_on (closed_segment b
c) f continuous_on (closed_segment c a) f
      using f continuous_on_subset segments_subset_convex_hull by metis+
    moreover have contour_integral (linepath b a) f + contour_integral (linepath
a c) f +
      contour_integral (linepath c b) f = 0
    proof (rule Cauchy_theorem_flat_lemma [of b a c f 1-k])
      show continuous_on (convex_hull {b, a, c}) f
        by (simp add: f_insert_commute)
      show c - b = (1 - k) *R (a - b)
        using c by (auto simp: algebra_simps)
    qed (use False in auto)
    ultimately show ?thesis
      by (metis (no_types, lifting) contour_integral_reverse_linepath eq_neg_iff_add_eq_0
minus_add_cancel)
  qed

```

```

proposition Cauchy_theorem_triangle_interior:
  assumes conf: continuous_on (convex_hull {a,b,c}) f
    and holf: f holomorphic_on_interior (convex_hull {a,b,c})
    shows (f has_contour_integral 0) (linepath a b +++ linepath b c +++ linepath
c a)
proof -
  define pathint where pathint ≡ λx y. contour_integral(linepath x y) f
  have fabc: continuous_on (closed_segment a b) f continuous_on (closed_segment
b c) f continuous_on (closed_segment c a) f
    using conf continuous_on_subset segments_subset_convex_hull by metis+
  have bounded (f ' (convex_hull {a,b,c}))
    by (simp add: compact_continuous_image compact_convex_hull compact_imp_bounded
conf)
  then obtain B where 0 < B and Bnf: ∧x. x ∈ convex_hull {a,b,c} ⇒ norm
(f x) ≤ B
    by (auto simp: dest!: bounded_pos [THEN iffD1])
  have bounded (convex_hull {a,b,c})
    by (simp add: bounded_convex_hull)
  then obtain C where C: 0 < C and Cno: ∧y. y ∈ convex_hull {a,b,c} ⇒

```

```

norm y < C
  using bounded_pos_less by blast
then have diff_2C: norm(x - y) ≤ 2*C
  if x: x ∈ convex hull {a, b, c} and y: y ∈ convex hull {a, b, c} for x y
proof -
  have cmod x ≤ C
  using x by (meson Cno not_le not_less_iff_gr_or_eq)
  hence cmod (x - y) ≤ C + C
  using y by (meson Cno add_mono_thms_linordered_field(4) less_eq_real_def
norm_triangle_ineq4 order_trans)
  thus cmod (x - y) ≤ 2 * C
  by (metis mult_2)
qed
have contf': continuous_on (convex hull {b,a,c}) f
  using contf by (simp add: insert_commute)
{ fix y::complex
  assume fy: (f has_contour_integral y) (linepath a b +++ linepath b c +++
linepath c a)
  and ynz: y ≠ 0
  have pi_eq_y: pathint a b + pathint b c + pathint c a = y
  unfolding pathint_def by (rule has_chain_integral_chain_integral3 [OF
fy])
  have ?thesis
  proof (cases c=a ∨ a=b ∨ b=c)
  case True then show ?thesis
    using Cauchy_theorem_flat [OF contf, of 0]
    using has_chain_integral_chain_integral3 [OF fy] ynz
    by (force simp: fabc contour_integral_reverse_linepath)
  next
  case False
  then have car3: card {a, b, c} = Suc (DIM(complex))
  by auto
  { assume interior(convex hull {a,b,c}) = {}
  then have collinear{a,b,c}
  using interior_convex_hull_eq_empty [OF car3]
  by (simp add: collinear_3_eq_affine_dependent)
  with False obtain d where c ≠ a a ≠ b b ≠ c c - b = d *R (a - b)
  by (auto simp: collinear_3 collinear_lemma)
  then have False
  using False Cauchy_theorem_flat [OF contf'] pi_eq_y ynz
  by (simp add: fabc add_eq_0_iff contour_integral_reverse_linepath
pathint_def)
  }
  then obtain d where d: d ∈ interior (convex hull {a, b, c})
  by blast
  { fix d1
  assume d1_pos: 0 < d1
  and d1: ∧x x'. [x ∈ convex hull {a, b, c}; x' ∈ convex hull {a, b, c}; cmod
(x' - x) < d1]

```

```

       $\implies \text{cmod } (f x' - f x) < \text{cmod } y / (24 * C)$ 
define e where e = min 1 (min (d1/(4*C)) ((norm y / 24 / C) / B))
define shrink where shrink x = x - e *R (x - d) for x
have e: 0 < e e ≤ 1 e ≤ d1 / (4 * C) e ≤ cmod y / 24 / C / B
  using d1_pos ‹C>0› ‹B>0› ynz by (simp_all add: e_def)
have e_le_d1: e * (4 * C) ≤ d1
  using e ‹C>0› by (simp add: field_simps)
have shrink a ∈ interior(convex hull {a,b,c})
  shrink b ∈ interior(convex hull {a,b,c})
  shrink c ∈ interior(convex hull {a,b,c})
using d e by (auto simp: hull_inc mem_interior_convex_shrink shrink_def)
then have fhp0: (f has_contour_integral 0)
  (linepath (shrink a) (shrink b) +++ linepath (shrink b) (shrink c)
  +++ linepath (shrink c) (shrink a))
  by (simp add: Cauchy_theorem_triangle holomorphic_on_subset [OF holf]
  hull_minimal)
  then have f_0_shrink: pathint (shrink a) (shrink b) + pathint (shrink b)
  (shrink c) + pathint (shrink c) (shrink a) = 0
  by (simp add: has_chain_integral_chain_integral3 pathint_def)
have fpi_abc: f contour_integrable_on linepath (shrink a) (shrink b)
  f contour_integrable_on linepath (shrink b) (shrink c)
  f contour_integrable_on linepath (shrink c) (shrink a)
using fhp0 by (auto simp: valid_path_join dest: has_contour_integral_integrable)
have cmod_shr:  $\bigwedge x y. \text{cmod } (\text{shrink } y - \text{shrink } x - (y - x)) = e * \text{cmod } (x - y)$ 
  using e by (simp add: shrink_def real_vector.scale_right_diff_distrib
  [symmetric])
have sh_eq:  $\bigwedge a b d::\text{complex}. (b - e *_{R} (b - d)) - (a - e *_{R} (a - d)) - (b - a) = e *_{R} (a - b)$ 
  by (simp add: algebra_simps)
have cmod_y / (24 * C) ≤ cmod y / cmod (b - a) / 12
  using False ‹C>0› diff_2C [of b a] ynz
  by (auto simp: field_split_simps hull_inc)
have less_C: x * cmod u < C if u ∈ convex hull {a,b,c} 0 ≤ x x ≤ 1 for
x u
proof (cases x=0)
case False
  with that show ?thesis
    using Cno [of u] mult_left_le_one_le [of cmod u x] le_less_trans
  norm_ge_zero by blast
qed (simp add: ‹0 < C›)
{ fix u v
  assume uv: u ∈ convex hull {a, b, c} v ∈ convex hull {a, b, c} u ≠ v
  and fpi_uv: f contour_integrable_on linepath (shrink u) (shrink v)
  have shr_uv: shrink u ∈ interior(convex hull {a,b,c})
  shrink v ∈ interior(convex hull {a,b,c})
  using d e uv
  by (auto simp: hull_inc mem_interior_convex_shrink shrink_def)
  have cmod_fuv:  $\bigwedge x. 0 \leq x \implies x \leq 1 \implies \text{cmod } (f (\text{linepath } (\text{shrink } u)$ 

```

```

(shrink v x) ≤ B
  using shr_uv by (blast intro: Bnflinepath_in_convex_hull_interior_subset
[THEN subsetD])
  { fix x::real assume x: 0 ≤ x x ≤ 1
    have |1 - x| * cmod u < C |x| * cmod v < C
      using uv x by (auto intro!: less_C)
    moreover have |x| * cmod d < C |1 - x| * cmod d < C
      using x d interior_subset by (auto intro!: less_C)
    ultimately
    have cmod_less_4C: cmod ((1 - x) *R u - (1 - x) *R d) + cmod (x
*_R v - x *_R d) < (C+C) + (C+C)
      by (metis add_strict_mono le_less_trans norm_scaleR norm_triangle_ineq4)
    have ll: linepath (shrink u) (shrink v) x - linepath u v x = -e * ((1 -
x) *_R (u - d) + x *_R (v - d))
      by (simp add: linepath_def shrink_def algebra_simps scaleR_conv_of_real)
    have cmod_less_dt: cmod (linepath (shrink u) (shrink v) x - linepath u
v x) < d1
      unfolding ll norm_mult scaleR_diff_right
      using ‹e>0› cmod_less_4C by (force intro: norm_triangle_lt
less_le_trans [OF _ e_le_d1])
    have cmod (f (linepath (shrink u) (shrink v) x)) * cmod (shrink v -
shrink u - (v - u)) +
      cmod (v - u) * cmod (f (linepath (shrink u) (shrink v) x) -
f (linepath u v x))
      ≤ B * (cmod y / 24 / C / B * 2 * C) + 2 * C * (cmod y /
24 / C)
    proof (intro add_mono [OF mult_mono])
      show cmod (f (linepath (shrink u) (shrink v) x)) ≤ B
        using cmod_fuv x by blast
      have B * (12 * (e * cmod (u - v))) ≤ 24 * e * C * B
        using e ‹B>0› diff_2C [of u v] uv by (auto simp: field_simps)
      also have ... ≤ cmod y
        using ‹C>0› ‹B>0› e by (simp add: field_simps)
      finally show cmod (shrink v - shrink u - (v - u)) ≤ cmod y / 24 /
C / B * 2 * C
        using ‹0 < B› ‹0 < C› by (simp add: cmod_shr mult_ac divide_simps)
      have cmod (f (linepath (shrink u) (shrink v) x) - f (linepath u v x))
< cmod y / (24 * C)
        using x uv shr_uv cmod_less_dt
        by (auto simp: hull_inc intro: d1 interior_subset [THEN subsetD]
linepath_in_convex_hull)
      also have ... ≤ cmod y / cmod (v - u) / 12
        using False uv ‹C>0› diff_2C [of v u] ynz
        by (auto simp: field_split_simps hull_inc)
      finally have cmod (f (linepath (shrink u) (shrink v) x) - f (linepath
u v x)) ≤ cmod y / cmod (v - u) / 12
        by simp
      then show cmod (v - u) * cmod (f (linepath (shrink u) (shrink v) x)
- f (linepath u v x))

```



```

      ≤ 2 * C * (cmod y / 24 / C)
    using uv C by (simp add: field_simps)
  qed (use ‹0 < B› in auto)
  also have ... ≤ cmod y / 6
    by simp
  finally have cmod (f (linepath (shrink u) (shrink v) x)) * cmod (shrink
v - shrink u - (v - u)) +
      cmod (v - u) * cmod (f (linepath (shrink u) (shrink v) x) -
f (linepath u v x))
      ≤ cmod y / 6 .
} note cmod_diff_le = this
have f_uv: continuous_on (closed_segment u v) f
by (blast intro: uv continuous_on_subset [OF contf closed_segment_subset convex_hull])
have **:  $\bigwedge f' x' f x :: \text{complex}. f' * x' - f * x = f' * (x' - x) + x * (f' - f)$ 
  by (simp add: algebra_simps)
have norm (pathint (shrink u) (shrink v) - pathint u v)
  ≤ (B * (norm y / 24 / C / B) * 2 * C + (2 * C) * (norm y / 24 / C)) * measure
lborel (cbox 0 (1 :: real))
  apply (rule has_integral_bound
    [of_  $\lambda x. f(\text{linepath } (\text{shrink } u) (\text{shrink } v) x) * (\text{shrink } v - \text{shrink }
u) - f(\text{linepath } u v x) * (v - u)$ 
      - 0 1])
  using ynz ‹0 < B› ‹0 < C›
  apply (simp_all add: pathint_def has_integral_diff has_contour_integral_linepath
[symmetric] has_contour_integral_integral
  fpi_uv f_uv contour_integrable_continuous_linepath del: le_divide_eq_numeral1)
  apply (auto simp: ** norm_triangle_le norm_mult cmod_diff_le simp
del: le_divide_eq_numeral1)
  done
  also have ... ≤ norm y / 6
    by simp
  finally have norm (pathint (shrink u) (shrink v) - pathint u v) ≤ norm
y / 6 .
} note * = this
have norm (pathint (shrink a) (shrink b) - pathint a b) ≤ norm y / 6
  using False fpi_abc by (rule_tac *) (auto simp: hull_inc)
moreover
have norm (pathint (shrink b) (shrink c) - pathint b c) ≤ norm y / 6
  using False fpi_abc by (rule_tac *) (auto simp: hull_inc)
moreover
have norm (pathint (shrink c) (shrink a) - pathint c a) ≤ norm y / 6
  using False fpi_abc by (rule_tac *) (auto simp: hull_inc)
ultimately
have norm((pathint (shrink a) (shrink b) - pathint a b) +
  (pathint (shrink b) (shrink c) - pathint b c) + (pathint (shrink c)
(shrink a) - pathint c a))
  ≤ norm y / 6 + norm y / 6 + norm y / 6
  by (metis norm_triangle_le add_mono)
  also have ... = norm y / 2

```

```

      by simp
      finally have norm((pathint (shrink a) (shrink b) + pathint (shrink b)
(shrink c) + pathint (shrink c) (shrink a)) -
      (pathint a b + pathint b c + pathint c a))
        ≤ norm y / 2
      by (simp add: algebra_simps)
    then
      have norm(pathint a b + pathint b c + pathint c a) ≤ norm y / 2
        by (simp add: f_0_shrink) (metis (mono_tags) add.commute mi-
nus_add_distrib norm_minus_cancel uminus_add_conv_diff)
      then have False
        using pi_eq_y ynz by auto
    }
    note * = this
    have uniformly_continuous_on (convex hull {a,b,c}) f
      by (simp add: contf compact_convex_hull_compact_uniformly_continuous)
    moreover have norm y / (24 * C) > 0
      using ynz ⟨C > 0⟩ by auto
    ultimately obtain δ where δ > 0 and
      ∀ x ∈ convex hull {a, b, c}. ∀ x' ∈ convex hull {a, b, c}.
        dist x' x < δ → dist (f x') (f x) < cmod y / (24 * C)
      using ⟨C > 0⟩ ynz unfolding uniformly_continuous_on_def dist_norm
by blast
      hence False using *[of δ] by (auto simp: dist_norm)
      then show ?thesis ..
    qed
  }
  moreover have f contour_integrable_on (linepath a b +++ linepath b c +++
linepath c a)
    using fabc contour_integrable_continuous_linepath by auto
  ultimately show ?thesis
    using has_contour_integral_integral by fastforce
qed

```

2.4 Version allowing finite number of exceptional points

proposition *Cauchy_theorem_triangle_cofinite:*

assumes *continuous_on (convex hull {a,b,c}) f*

and *finite S*

and $(\bigwedge x. x \in \text{interior}(\text{convex hull } \{a,b,c\}) - S \implies f \text{ field_differentiable (at } x))$

shows $(f \text{ has_contour_integral } 0) (\text{linepath } a \text{ } b \text{ } +++ \text{ linepath } b \text{ } c \text{ } +++ \text{ linepath } c \text{ } a)$

using *assms*

proof (*induction card S arbitrary: a b c S rule: less_induct*)

case (*less S a b c*)

show *?case*

proof (*cases S={}*)

case *True with less show ?thesis*

```

    by (fastforce simp: holomorphic_on_def field_differentiable_at_within Cauchy_theorem_triangle_interior)
  next
  case False
  then obtain d S' where d: S = insert d S' d ∉ S'
    by (meson Set.set_insert_all_not_in_conv)
  then show ?thesis
  proof (cases d ∈ convex_hull {a,b,c})
    case False
    show (f has_contour_integral 0) (linepath a b +++ linepath b c +++ linepath
c a)
    proof (rule less.hyps)
      show  $\bigwedge x. x \in \text{interior} (\text{convex\_hull} \{a, b, c\}) - S' \implies f \text{ field\_differentiable}$ 
at x
      using False d interior_subset by (auto intro!: less.prem)
    qed (use d less.prem in auto)
  next
  case True
  have *: convex_hull {a, b, d} ⊆ convex_hull {a, b, c}
    by (meson True hull_subset_insert_subset convex_hull_subset)
  have abd: (f has_contour_integral 0) (linepath a b +++ linepath b d +++
linepath d a)
  proof (rule less.hyps)
    show  $\bigwedge x. x \in \text{interior} (\text{convex\_hull} \{a, b, d\}) - S' \implies f \text{ field\_differentiable}$ 
at x
    using d not_in_interior_convex_hull_3
    by (clarsimp intro!: less.prem) (metis * insert_absorb insert_subset
interior_mono)
  qed (use d continuous_on_subset [OF _ *] less.prem in auto)
  have *: convex_hull {b, c, d} ⊆ convex_hull {a, b, c}
    by (meson True hull_subset_insert_subset convex_hull_subset)
  have bcd: (f has_contour_integral 0) (linepath b c +++ linepath c d +++
linepath d b)
  proof (rule less.hyps)
    show  $\bigwedge x. x \in \text{interior} (\text{convex\_hull} \{b, c, d\}) - S' \implies f \text{ field\_differentiable}$ 
at x
    using d not_in_interior_convex_hull_3
    by (clarsimp intro!: less.prem) (metis * insert_absorb insert_subset
interior_mono)
  qed (use d continuous_on_subset [OF _ *] less.prem in auto)
  have *: convex_hull {c, a, d} ⊆ convex_hull {a, b, c}
    by (meson True hull_subset_insert_subset convex_hull_subset)
  have cad: (f has_contour_integral 0) (linepath c a +++ linepath a d +++
linepath d c)
  proof (rule less.hyps)
    show  $\bigwedge x. x \in \text{interior} (\text{convex\_hull} \{c, a, d\}) - S' \implies f \text{ field\_differentiable}$ 
at x
    using d not_in_interior_convex_hull_3
    by (clarsimp intro!: less.prem) (metis * insert_absorb insert_subset
interior_mono)

```

```

qed (use d continuous_on_subset [OF _ *] less.premis in auto)
have f contour_integrable_on linepath a b
  using less.premis abd contour_integrable_joinD1 contour_integrable_on_def
by blast
moreover have f contour_integrable_on linepath b c
  using less.premis bcd contour_integrable_joinD1 contour_integrable_on_def
by blast
moreover have f contour_integrable_on linepath c a
  using less.premis cad contour_integrable_joinD1 contour_integrable_on_def
by blast
ultimately have fpi: f contour_integrable_on (linepath a b +++ linepath b
c +++ linepath c a)
  by auto
  { fix y::complex
    assume fy: (f has_contour_integral y) (linepath a b +++ linepath b c +++
linepath c a)
      and ynz: y ≠ 0
      have cont_ad: continuous_on (closed_segment a d) f
      by (meson * continuous_on_subset less.premis(1) segments_subset_convex_hull(3))
      have cont_bd: continuous_on (closed_segment b d) f
      by (meson True closed_segment_subset_convex_hull continuous_on_subset
hull_subset_insert_subset less.premis(1))
      have cont_cd: continuous_on (closed_segment c d) f
      by (meson * continuous_on_subset less.premis(1) segments_subset_convex_hull(2))
      have contour_integral (linepath a b) f = - (contour_integral (linepath b
d) f) + (contour_integral (linepath d a) f)
          contour_integral (linepath b c) f = - (contour_integral (linepath c d)
f) + (contour_integral (linepath d b) f)
          contour_integral (linepath c a) f = - (contour_integral (linepath a d)
f) + contour_integral (linepath d c) f)
      using has_chain_integral_chain_integral3 [OF abd]
          has_chain_integral_chain_integral3 [OF bcd]
          has_chain_integral_chain_integral3 [OF cad]
      by (simp_all add: algebra_simps add_eq_0_iff)
    then have ?thesis
      using cont_ad cont_bd cont_cd fy has_chain_integral_chain_integral3
contour_integral_reverse_linepath by fastforce
  }
  then show ?thesis
    using fpi contour_integrable_on_def by blast
qed
qed
qed

```

2.5 Cauchy's theorem for an open starlike set

lemma starlike_convex_subset:

assumes S : $a \in S$ closed_segment $b c \subseteq S$ and subs: $\bigwedge x. x \in S \implies$ closed_segment $a x \subseteq S$

```

shows convex hull {a,b,c}  $\subseteq$  S
proof -
  have convex hull {b, c}  $\subseteq$  S
    using assms(2) segment_convex_hull by auto
  then have  $\bigwedge u v d. [0 \leq u; 0 \leq v; u + v = 1; d \in \text{convex hull } \{b, c\}] \implies u$ 
  *R a + v *R d  $\in$  S
    by (meson subs convexD convex_closed_segment ends_in_segment subsetCE)
  then show ?thesis
    by (auto simp add: convex_hull_insert [of {b,c} a])
qed

lemma triangle_contour_integrals_starlike_primitive:
  assumes contf: continuous_on S f
    and S: a  $\in$  S open S
    and x: x  $\in$  S
    and subs:  $\bigwedge y. y \in S \implies \text{closed\_segment } a y \subseteq S$ 
    and zer:  $\bigwedge b c. \text{closed\_segment } b c \subseteq S$ 
       $\implies \text{contour\_integral (linepath a b) } f + \text{contour\_integral (linepath$ 
b c) } f +
      contour\_integral (linepath c a) } f = 0
  shows (( $\lambda x. \text{contour\_integral (linepath a x) } f$ ) has_field_derivative f x) (at x)
proof -
  let ?pathint =  $\lambda x y. \text{contour\_integral (linepath x y) } f$ 
  { fix e y
    assume e: 0 < e and bxe: ball x e  $\subseteq$  S and close: cmod (y - x) < e
    have y: y  $\in$  S
      using bxe close by (force simp: dist_norm norm_minus_commute)
    have cont_ayf: continuous_on (closed_segment a y) f
      using contf continuous_on_subset subs y by blast
    have xys: closed_segment x y  $\subseteq$  S
      by (metis bxe centre_in_ball close closed_segment_subset convex_ball dist_norm
dual_order.trans e mem_ball norm_minus_commute)
    have ?pathint a y - ?pathint a x = ?pathint x y
      using zer [OF xys] contour_integral_reverse_linepath [OF cont_ayf] add_eq_0_iff
    by force
  } note [simp] = this
  { fix e::real
    assume e: 0 < e
    have cont_atx: continuous (at x) f
      using x S contf continuous_on_eq_continuous_at by blast
    then obtain d1 where d1: d1 > 0 and d1_less:  $\bigwedge y. \text{cmod } (y - x) < d1 \implies$ 
cmod (f y - f x) < e/2
      unfolding continuous_at Lim_at dist_norm using e
    by (drule_tac x=e/2 in spec) force
    obtain d2 where d2: d2 > 0 ball x d2  $\subseteq$  S using <open S> x
      by (auto simp: open_contains_ball)
    have dpos: min d1 d2 > 0 using d1 d2 by simp
    { fix y
      assume yx: y  $\neq$  x and close: cmod (y - x) < min d1 d2

```

```

have  $y: y \in S$ 
  using  $d2$  close by ( $force$   $simp: dist\_norm$   $norm\_minus\_commute$ )
have  $closed\_segment\ x\ y \subseteq S$ 
  using  $close\ d2$  by ( $auto\ simp: dist\_norm$   $norm\_minus\_commute\ dest!$ :
 $segment\_bound(1)$ )
  then have  $fx: f\ contour\_integrable\_on\ linepath\ x\ y$ 
    by ( $metis\ contour\_integrable\_continuous\_linepath\ continuous\_on\_subset$ 
 $[OF\ contf]$ )
  then obtain  $i$  where  $i: (f\ has\_contour\_integral\ i)\ (linepath\ x\ y)$ 
    by ( $auto\ simp: contour\_integrable\_on\_def$ )
  then have  $((\lambda w. f\ w - f\ x)\ has\_contour\_integral\ (i - f\ x * (y - x)))$ 
 $(linepath\ x\ y)$ 
    by ( $rule\ has\_contour\_integral\_diff\ [OF\_has\_contour\_integral\_const\_linepath]$ )
  then have  $cmod\ (i - f\ x * (y - x)) \leq e / 2 * cmod\ (y - x)$ 
    proof ( $rule\ has\_contour\_integral\_bound\_linepath$ )
      show  $\bigwedge u. u \in closed\_segment\ x\ y \implies cmod\ (f\ u - f\ x) \leq e / 2$ 
        by ( $meson\ close\ d1\_less\ le\_less\_trans\ less\_imp\_le\ min.strict\_boundedE$ 
 $segment\_bound1$ )
      qed ( $use\ e\ in\ simp$ )
  also have  $\dots < e * cmod\ (y - x)$ 
    by ( $simp\ add: e\ yx$ )
  finally have  $cmod\ (?pathint\ x\ y - f\ x * (y - x)) / cmod\ (y - x) < e$ 
    using  $i\ yx$  by ( $simp\ add: contour\_integral\_unique\ divide\_less\_eq$ )
  }
  then have  $\exists d > 0. \forall y. y \neq x \wedge cmod\ (y - x) < d \implies cmod\ (?pathint\ x\ y - f$ 
 $x * (y - x)) / cmod\ (y - x) < e$ 
    using  $dpos$  by  $blast$ 
  }
  then have  $(\lambda y. (?pathint\ x\ y - f\ x * (y - x)) /_R\ cmod\ (y - x)) -x \rightarrow 0$ 
    by ( $simp\ add: Lim\_at\ dist\_norm\ inverse\_eq\_divide$ )
  then have  $(\lambda y. (1 / cmod\ (y - x)) *_R\ (?pathint\ a\ y - (?pathint\ a\ x + f\ x * (y$ 
 $- x)))) -x \rightarrow 0$ 
    using  $\langle open\ S \rangle\ x$ 
    by ( $force\ simp: dist\_norm\ open\_contains\_ball\ inverse\_eq\_divide\ [symmetric]$ 
 $eventually\_at\ intro: Lim\_transform\ [OF\_tendsto\_eventually]$ )
  then show  $?thesis$ 
    by ( $simp\ add: has\_field\_derivative\_def\ has\_derivative\_at2\ bounded\_linear\_mult\_right$ )
qed

```

Existence of a primitive

lemma $holomorphic_starlike_primitive$:

fixes $f :: complex \Rightarrow complex$

assumes $contf: continuous_on\ S\ f$

and $S: starlike\ S$ **and** $os: open\ S$

and $k: finite\ k$

and $fed: \bigwedge x. x \in S - k \implies f\ field_differentiable\ at\ x$

shows $\exists g. \forall x \in S. (g\ has_field_derivative\ f\ x)\ (at\ x)$

proof -

obtain a **where** $a: a \in S$ **and** $a_cs: \bigwedge x. x \in S \implies closed_segment\ a\ x \subseteq S$

```

    using S by (auto simp: starlike_def)
  { fix x b c
    assume x ∈ S closed_segment b c ⊆ S
    then have abcs: convex hull {a, b, c} ⊆ S
      by (simp add: a a_cs starlike_convex_subset)
    then have continuous_on (convex hull {a, b, c}) f
      by (simp add: continuous_on_subset [OF contf])
    then have (f has_contour_integral 0) (linepath a b +++ linepath b c +++
linepath c a)
      using abcs interior_subset by (force intro: fcd Cauchy_theorem_triangle_cofinite
[OF _ k])
    } note 0 = this
    show ?thesis
    proof (intro exI ballI)
      show  $\bigwedge x. x \in S \implies ((\lambda x. \text{contour\_integral } (\text{linepath } a \ x) \ f) \text{ has\_field\_derivative } f \ x) \text{ (at } x)$ 
        using 0 a a_cs contf has_chain_integral_chain_integral3 os triangle_contour_integrals_starlike_primitive
    by force
    qed
  qed

```

lemma *Cauchy_theorem_starlike:*

```

[[open S; starlike S; finite k; continuous_on S f;
 $\bigwedge x. x \in S - k \implies f \text{ field\_differentiable at } x;$ 
valid_path g; path_image g ⊆ S; pathfinish g = pathstart g]]
 $\implies (f \text{ has\_contour\_integral } 0) \ g$ 
by (metis holomorphic_starlike_primitive Cauchy_theorem_primitive at_within_open)

```

lemma *Cauchy_theorem_starlike_simple:*

```

[[open S; starlike S; f holomorphic_on S; valid_path g; path_image g ⊆ S; pathfinish g = pathstart g]]
 $\implies (f \text{ has\_contour\_integral } 0) \ g$ 
using Cauchy_theorem_starlike [OF _ _ finite.emptyI]
by (simp add: holomorphic_on_imp_continuous_on holomorphic_on_imp_differentiable_at)

```

2.6 Cauchy's theorem for a convex set

For a convex set we can avoid assuming openness and boundary analyticity

lemma *triangle_contour_integrals_convex_primitive:*

```

assumes contf: continuous_on S f
and S: a ∈ S convex S
and x: x ∈ S
and zer:  $\bigwedge b \ c. [b \in S; c \in S]$ 
 $\implies \text{contour\_integral } (\text{linepath } a \ b) \ f + \text{contour\_integral } (\text{linepath } b \ c) \ f +$ 
 $\text{contour\_integral } (\text{linepath } c \ a) \ f = 0$ 
shows (( $\lambda x. \text{contour\_integral}(\text{linepath } a \ x) \ f$ ) has_field_derivative f x) (at x within S)
proof -

```

```

let ?pathint =  $\lambda x y.$  contour_integral(linepath x y) f
{ fix y
  assume  $y: y \in S$ 
  have cont_ayf: continuous_on (closed_segment a y) f
  using S y by (meson contf continuous_on_subset convex_contains_segment)
  have xys: closed_segment x y  $\subseteq S$ 
  using convex_contains_segment S x y by auto
  have ?pathint a y - ?pathint a x = ?pathint x y
  using zer [OF x y] contour_integral_reverse_linepath [OF cont_ayf] add_eq_0_iff
by force
} note [simp] = this
{ fix e::real
  assume  $e: 0 < e$ 
  have cont_atx: continuous (at x within S) f
  using x S contf by (simp add: continuous_on_eq_continuous_within)
  then obtain d1 where  $d1: d1 > 0$  and  $d1\_less: \bigwedge y. \llbracket y \in S; cmod (y - x) < d1 \rrbracket \implies cmod (f y - f x) < e/2$ 
  unfolding continuous_within Lim_within dist_norm using e
  by (drule_tac  $x=e/2$  in spec) force
  { fix y
    assume  $yx: y \neq x$  and close:  $cmod (y - x) < d1$  and  $y: y \in S$ 
    have fxy: f contour_integrable_on linepath x y
    using convex_contains_segment S x y
    by (blast intro!: contour_integrable_continuous_linepath continuous_on_subset [OF contf])
    then obtain i where  $i: (f \text{ has\_contour\_integral } i) (linepath x y)$ 
    by (auto simp: contour_integrable_on_def)
    then have  $((\lambda w. f w - f x) \text{ has\_contour\_integral } (i - f x * (y - x)))$ 
    (linepath x y)
    by (rule has_contour_integral_diff [OF _ has_contour_integral_const_linepath])
    then have  $cmod (i - f x * (y - x)) \leq e / 2 * cmod (y - x)$ 
    proof (rule has_contour_integral_bound_linepath)
      show  $\bigwedge u. u \in \text{closed\_segment } x y \implies cmod (f u - f x) \leq e / 2$ 
      by (meson assms(3) close convex_contains_segment d1_less le_less_trans less_imp_le segment_bound1 subset_iff x y)
    qed (use e in simp)
    also have  $\dots < e * cmod (y - x)$ 
    by (simp add: e yx)
    finally have  $cmod (?pathint x y - f x * (y-x)) / cmod (y-x) < e$ 
    using i yx by (simp add: contour_integral_unique divide_less_eq)
  }
  then have  $\exists d > 0. \forall y \in S. y \neq x \wedge cmod (y-x) < d \longrightarrow cmod (?pathint x y - f x * (y-x)) / cmod (y-x) < e$ 
  using d1 by blast
}
then have  $((\lambda y. (?pathint x y - f x * (y - x)) /_R cmod (y - x)) \longrightarrow 0)$  (at x within S)
by (simp add: Lim_within dist_norm inverse_eq_divide)
then have  $((\lambda y. (1 / cmod (y - x)) *_R (?pathint a y - (?pathint a x + f x *$ 

```



```

(y - x)))  $\longrightarrow$  0)
  (at x within S)
  using linordered_field_no_ub
  by (force simp: inverse_eq_divide [symmetric] eventually_at intro: Lim_transform
[OF _ tendsto_eventually])
  then show ?thesis
  by (simp add: has_field_derivative_def has_derivative_within bounded_linear_mult_right)
qed

```

lemma *contour_integral_convex_primitive:*

```

assumes convex S continuous_on S f
   $\bigwedge a b c. \llbracket a \in S; b \in S; c \in S \rrbracket \implies (f \text{ has\_contour\_integral } 0) \text{ (linepath } a
b \text{ +++ linepath } b c \text{ +++ linepath } c a)$ 
  obtains g where  $\bigwedge x. x \in S \implies (g \text{ has\_field\_derivative } f x) \text{ (at } x \text{ within } S)$ 
proof (cases S={})
  case False
  with assms that show ?thesis
  by (blast intro: triangle_contour_integrals_convex_primitive has_chain_integral_chain_integral3)
qed auto

```

lemma *holomorphic_convex_primitive:*

```

fixes f :: complex  $\Rightarrow$  complex
assumes convex S finite K and contf: continuous_on S f
  and fd:  $\bigwedge x. x \in \text{interior } S - K \implies f \text{ field\_differentiable at } x$ 
  obtains g where  $\bigwedge x. x \in S \implies (g \text{ has\_field\_derivative } f x) \text{ (at } x \text{ within } S)$ 
proof (rule contour_integral_convex_primitive [OF <convex S> contf Cauchy_theorem_triangle_cofinite])
  have *: convex hull {a, b, c}  $\subseteq$  S if a  $\in$  S b  $\in$  S c  $\in$  S for a b c
  by (simp add: <convex S> hull_minimal that)
  show continuous_on (convex hull {a, b, c}) f if a  $\in$  S b  $\in$  S c  $\in$  S for a b c
  by (meson * contf continuous_on_subset that)
  show f field_differentiable at x if a  $\in$  S b  $\in$  S c  $\in$  S x  $\in$  interior (convex hull
{a, b, c}) - K for a b c x
  by (metis * DiffD1 DiffD2 DiffI fd interior_mono subsetCE that)
qed (use assms in <force+>)

```

lemma *holomorphic_convex_primitive':*

```

fixes f :: complex  $\Rightarrow$  complex
assumes convex S and open S and f holomorphic_on S
  obtains g where  $\bigwedge x. x \in S \implies (g \text{ has\_field\_derivative } f x) \text{ (at } x \text{ within } S)$ 
proof (rule holomorphic_convex_primitive)
  fix x assume x  $\in$  interior S - {}
  with assms show f field_differentiable at x
  by (auto intro!: holomorphic_on_imp_differentiable_at simp: interior_open)
qed (use assms in <auto intro: holomorphic_on_imp_continuous_on>)

```

corollary *Cauchy_theorem_convex:*

```

 $\llbracket \text{continuous\_on } S f; \text{convex } S; \text{finite } K;
\bigwedge x. x \in \text{interior } S - K \implies f \text{ field\_differentiable at } x;
\text{valid\_path } g; \text{path\_image } g \subseteq S; \text{pathfinish } g = \text{pathstart } g \rrbracket$ 

```

```

     $\implies$  (f has_contour_integral 0) g
  by (metis holomorphic_convex_primitive Cauchy_theorem_primitive)

corollary Cauchy_theorem_convex_simple:
  assumes holf: f holomorphic_on S
    and convex S valid_path g path_image g  $\subseteq$  S pathfinish g = pathstart g
  shows (f has_contour_integral 0) g
proof -
  have f holomorphic_on interior S
    by (meson holf holomorphic_on_subset interior_subset)
  with Cauchy_theorem_convex [where K = {}] show ?thesis
    using assms
    by (metis Diff_empty_finite.emptyI holomorphic_on_imp_continuous_on_holo-
        morphic_on_imp_differentiable_at_open_interior)
qed

```

In particular for a disc

```

corollary Cauchy_theorem_disc:
   $\llbracket$ finite K; continuous_on (cball a e) f;
   $\bigwedge x. x \in \text{ball } a \ e - K \implies$  f field_differentiable at x;
  valid_path g; path_image g  $\subseteq$  cball a e;
  pathfinish g = pathstart g  $\rrbracket \implies$  (f has_contour_integral 0) g
  by (auto intro: Cauchy_theorem_convex)

```

```

corollary Cauchy_theorem_disc_simple:
   $\llbracket$ f holomorphic_on (ball a e); valid_path g; path_image g  $\subseteq$  ball a e;
  pathfinish g = pathstart g  $\rrbracket \implies$  (f has_contour_integral 0) g
  by (simp add: Cauchy_theorem_convex_simple)

```

2.7 Generalize integrability to local primitives

```

lemma contour_integral_local_primitive_lemma:
  fixes f :: complex  $\Rightarrow$  complex
  assumes gpd: g piecewise_differentiable_on {a..b}
    and dh:  $\bigwedge x. x \in S \implies$  (f has_field_derivative f' x) (at x within S)
    and gs:  $\bigwedge x. x \in \{a..b\} \implies$  g x  $\in$  S
  shows ( $\lambda x. f' (g x) * \text{vector\_derivative } g$  (at x within {a..b})) integrable_on
    {a..b}
proof (cases cbox a b = {})
  case False
  then show ?thesis
    unfolding integrable_on_def by (auto intro: assms contour_integral_primitive_lemma)
qed auto

```

```

lemma contour_integral_local_primitive_any:
  fixes f :: complex  $\Rightarrow$  complex
  assumes gpd: g piecewise_differentiable_on {a..b}
    and dh:  $\bigwedge x. x \in S$ 
       $\implies \exists d \ h. 0 < d \wedge$ 

```

```

      ( $\forall y. \text{norm}(y - x) < d \longrightarrow (h \text{ has\_field\_derivative } f y) \text{ (at } y$ 
within  $S)$ )
    and gs:  $\bigwedge x. x \in \{a..b\} \implies g x \in S$ 
    shows  $(\lambda x. f(g x) * \text{vector\_derivative } g \text{ (at } x)) \text{ integrable\_on } \{a..b\}$ 
proof -
  { fix  $x$ 
    assume  $x: a \leq x \leq b$ 
    obtain  $d h$  where  $d: 0 < d$ 
      and  $h: (\bigwedge y. \text{norm}(y - g x) < d \implies (h \text{ has\_field\_derivative } f y) \text{ (at } y$ 
within  $S)$ )
    using  $x$  gs dh by (metis atLeastAtMost_iff)
    have continuous_on  $\{a..b\}$   $g$  using gpd piecewise_differentiable_on_def by
blast
    then obtain  $e$  where  $e: e > 0$  and lessd:  $\bigwedge x'. x' \in \{a..b\} \implies |x' - x| < e$ 
 $\implies \text{cmod } (g x' - g x) < d$ 
    using  $x$  d by (fastforce simp: dist_norm continuous_on_iff)
    have  $\exists e > 0. \forall u v. u \leq x \wedge x \leq v \wedge \{u..v\} \subseteq \text{ball } x e \wedge (u \leq v \longrightarrow a \leq u \wedge$ 
 $v \leq b) \longrightarrow$ 
       $(\lambda x. f(g x) * \text{vector\_derivative } g \text{ (at } x)) \text{ integrable\_on } \{u..v\}$ 
proof -
    have  $(\lambda x. f(g x) * \text{vector\_derivative } g \text{ (at } x \text{ within } \{u..v\})) \text{ integrable\_on}$ 
 $\{u..v\}$ 
    if  $u \leq x \leq v$  and ball:  $\{u..v\} \subseteq \text{ball } x e$  and auvb:  $u \leq v \implies a \leq u \wedge v$ 
 $\leq b$ 
    for  $u v$ 
    proof (rule contour_integral_local_primitive_lemma)
    show  $g$  piecewise_differentiable_on  $\{u..v\}$ 
    by (metis atLeastAtMost_subset_iff gpd piecewise_differentiable_on_subset
auvb)
    show  $\bigwedge x. x \in g^{-1} \{u..v\} \implies (h \text{ has\_field\_derivative } f x) \text{ (at } x \text{ within } g^{-1}$ 
 $\{u..v\})$ 
    using that by (force simp: ball_def dist_norm intro: lessd gs DERIV_subset
[OF h])
    qed auto
    then show ?thesis
    using  $e$  integrable_on_localized_vector_derivative by blast
    qed
  } then
show ?thesis
by (force simp: intro!: integrable_on_little_subintervals [of a b, simplified])
qed

lemma contour_integral_local_primitive:
fixes  $f :: \text{complex} \Rightarrow \text{complex}$ 
assumes  $g: \text{valid\_path } g \text{ path\_image } g \subseteq S$ 
and dh:  $\bigwedge x. x \in S$ 
 $\implies \exists d h. 0 < d \wedge$ 
 $(\forall y. \text{norm}(y - x) < d \longrightarrow (h \text{ has\_field\_derivative } f y) \text{ (at } y$ 
within  $S)$ )

```

```

  shows  $f$  contour_integrable_on  $g$ 
proof –
  have  $(\lambda x. f (g x) * \text{vector\_derivative } g \text{ (at } x))$  integrable_on  $\{0..1\}$ 
    using contour_integral_local_primitive_any [OF dh]  $g$ 
    unfolding path_image_def valid_path_def
    by (metis (no_types, lifting) image_subset_iff piecewise_C1_imp_differentiable)
  then show ?thesis
    using contour_integrable_on by presburger
qed

```

In particular if a function is holomorphic

lemma *contour_integrable_holomorphic*:

```

assumes contf: continuous_on  $S$   $f$ 
  and os: open  $S$ 
  and k: finite  $k$ 
  and g: valid_path  $g$  path_image  $g \subseteq S$ 
  and fed:  $\bigwedge x. x \in S - k \implies f$  field_differentiable at  $x$ 
shows  $f$  contour_integrable_on  $g$ 
proof –
  { fix  $z$ 
    assume  $z: z \in S$ 
    obtain  $d$  where  $d > 0$  and  $d: \text{ball } z \ d \subseteq S$  using  $\langle \text{open } S \rangle z$ 
      by (auto simp: open_contains_ball)
    then have contfb: continuous_on  $(\text{ball } z \ d)$   $f$ 
      using contf continuous_on_subset by blast
    obtain  $h$  where  $\forall y \in \text{ball } z \ d. (h \text{ has\_field\_derivative } f \ y)$  (at  $y$  within  $\text{ball } z \ d)$ 
      by (metis holomorphic_convex_primitive [OF convex_ball  $k$  contfb fed] d
interior_subset Diff_iff_subsetD)
    then have  $\forall y \in \text{ball } z \ d. (h \text{ has\_field\_derivative } f \ y)$  (at  $y$  within  $S$ )
      by (metis open_ball at_within_open  $d$  os subsetCE)
    then have  $\exists h. (\forall y. \text{cmod } (y - z) < d \implies (h \text{ has\_field\_derivative } f \ y))$  (at  $y$ 
within  $S)$ 
      by (force simp: dist_norm norm_minus_commute)
    then have  $\exists d \ h. 0 < d \wedge (\forall y. \text{cmod } (y - z) < d \implies (h \text{ has\_field\_derivative } f \ y))$ 
(at  $y$  within  $S)$ 
      using  $\langle 0 < d \rangle$  by blast
  }
then show ?thesis
  by (rule contour_integral_local_primitive [OF  $g$ ])
qed

```

lemma *contour_integrable_holomorphic_simple*:

```

assumes fh:  $f$  holomorphic_on  $S$ 
  and os: open  $S$ 
  and g: valid_path  $g$  path_image  $g \subseteq S$ 
shows  $f$  contour_integrable_on  $g$ 
proof –
  have  $\bigwedge x. x \in S \implies f$  field_differentiable at  $x$ 
    using fh holomorphic_on_imp_differentiable_at os by blast

```

```

moreover have continuous_on S f
  by (simp add: fh holomorphic_on_imp_continuous_on)
ultimately show ?thesis
  by (metis Diff_empty contour_integrable_holomorphic_finite.emptyI g os)
qed

```

```

lemma analytic_imp_contour_integrable:
  assumes f analytic_on path_image p valid_path p
  shows f contour_integrable_on p
  by (meson analytic_on_holomorphic assms contour_integrable_holomorphic_simple)

```

```

lemma continuous_on_inversediff:
  fixes z::'a::real_normed_field shows  $z \notin S \implies \text{continuous\_on } S (\lambda w. 1 / (w - z))$ 
  by (rule continuous_intros | force)+

```

```

lemma contour_integrable_inversediff:
  assumes g: valid_path g
  and notin: z \notin path_image g
  shows  $(\lambda w. 1 / (w - z)) \text{ contour\_integrable\_on } g$ 
proof (rule contour_integrable_holomorphic_simple)
  show  $(\lambda w. 1 / (w - z)) \text{ holomorphic\_on } UNIV - \{z\}$ 
  by (auto simp: holomorphic_on_open open_delete intro!: derivative_eq_intros)
qed (use assms in auto)

```

Key fact that path integral is the same for a "nearby" path. This is the main lemma for the homotopy form of Cauchy's theorem and is also useful if we want "without loss of generality" to assume some nice properties of a path (e.g. smoothness). It can also be used to define the integrals of analytic functions over arbitrary continuous paths. This is just done for winding numbers now.

A technical definition to avoid duplication of similar proofs, for paths joined at the ends versus looping paths

```

definition linked_paths :: bool  $\Rightarrow$  (real  $\Rightarrow$  'a)  $\Rightarrow$  (real  $\Rightarrow$  'a::topological_space)  $\Rightarrow$  bool
  where linked_paths attends g h ==
    (if attends then pathstart h = pathstart g \wedge pathfinish h = pathfinish g
     else pathfinish g = pathstart g \wedge pathfinish h = pathstart h)

```

This formulation covers two cases: g and h share their start and end points; g and h both loop upon themselves.

```

lemma contour_integral_nearby:
  assumes os: open S and p: path p path_image p \subseteq S
  shows  $\exists d. 0 < d \wedge$ 
     $(\forall g h. \text{valid\_path } g \wedge \text{valid\_path } h \wedge$ 
       $(\forall t \in \{0..1\}. \text{norm}(g\ t - p\ t) < d \wedge \text{norm}(h\ t - p\ t) < d) \wedge$ 
      linked_paths attends g h)

```

```

     $\longrightarrow \text{path\_image } g \subseteq S \wedge \text{path\_image } h \subseteq S \wedge$ 
     $(\forall f. f \text{ holomorphic\_on } S \longrightarrow \text{contour\_integral } h f = \text{contour\_integral } g f)$ 
  proof -
    have  $\forall z. \exists e. z \in \text{path\_image } p \longrightarrow 0 < e \wedge \text{ball } z e \subseteq S$ 
      using open_contains_ball os p(2) by blast
    then obtain ee where ee:  $\bigwedge z. z \in \text{path\_image } p \implies 0 < ee z \wedge \text{ball } z (ee z) \subseteq S$ 
      by metis
    define cover where cover =  $(\lambda z. \text{ball } z (ee z / 3)) \text{ ` } (\text{path\_image } p)$ 
    have compact (path_image p)
      by (metis p(1) compact_path_image)
    moreover have path_image p  $\subseteq (\bigcup c \in \text{path\_image } p. \text{ball } c (ee c / 3))$ 
      using ee by auto
    ultimately have  $\exists D \subseteq \text{cover}. \text{finite } D \wedge \text{path\_image } p \subseteq \bigcup D$ 
      by (simp add: compact_eq_Heine_Borel_cover_def)
    then obtain D where D:  $D \subseteq \text{cover}$  finite D path_image p  $\subseteq \bigcup D$ 
      by blast
    then obtain k where k:  $k \subseteq \{0..1\}$  finite k and D_eq:  $D = ((\lambda z. \text{ball } z (ee z / 3)) \circ p) \text{ ` } k$ 
      unfolding cover_def path_image_def image_comp
      by (meson finite_subset_image)
    then have kne:  $k \neq \{\}$ 
      using D by auto
    have pi:  $\bigwedge i. i \in k \implies p i \in \text{path\_image } p$ 
      using k by (auto simp: path_image_def)
    then have eepi:  $\bigwedge i. i \in k \implies 0 < ee(p i)$ 
      by (metis ee)
    define e where e =  $\text{Min}((ee \circ p) \text{ ` } k)$ 
    have fin_eep: finite ((ee \circ p) ` k)
      using k by blast
    have 0 < e
      using ee k by (simp add: kne e_def Min_gr_iff [OF fin_eep] eepi)
    have uniformly_continuous_on {0..1} p
      using p by (simp add: path_def compact_uniformly_continuous)
    then obtain d::real where d:  $d > 0$ 
      and de:  $\bigwedge x x'. |x' - x| < d \implies x \in \{0..1\} \implies x' \in \{0..1\} \implies \text{cmod } (p x' - p x) < e / 3$ 
      unfolding uniformly_continuous_on_def dist_norm real_norm_def
      by (metis divide_pos_pos <0 < e> zero_less_numeral)
    then obtain N::nat where N:  $N > 0$  inverse N < d
      using real_arch_inverse [of d] by auto
    show ?thesis
    proof (intro exI conjI allI; clarify?)
      show  $e / 3 > 0$ 
        using <0 < e> by simp
      fix g h
      assume g: valid_path g and ghp:  $\forall t \in \{0..1\}. \text{cmod } (g t - p t) < e / 3 \wedge \text{cmod } (h t - p t) < e / 3$ 

```

```

    and h: valid_path h
    and joins: linked_paths attends g h
  { fix t::real
    assume t: 0 ≤ t t ≤ 1
    then obtain u where u: u ∈ k and ptu: p t ∈ ball(p u) (ee(p u) / 3)
      using ‹path_image p ⊆ ⋃ D› D_eq by (force simp: path_image_def)
    then have ele: e ≤ ee (p u) using fn_eep
      by (simp add: e_def)
    have cmod (g t - p t) < e / 3 cmod (h t - p t) < e / 3
      using ghp t by auto
    with ele have cmod (g t - p t) < ee (p u) / 3
      cmod (h t - p t) < ee (p u) / 3
      by linarith+
    then have g t ∈ ball(p u) (ee(p u)) h t ∈ ball(p u) (ee(p u))
      using norm_diff_triangle_ineq [of g t p t p t p u]
      norm_diff_triangle_ineq [of h t p t p t p u] ptu eepi u
      by (force simp: dist_norm ball_def norm_minus_commute)+
    then have g t ∈ S h t ∈ S using ee u k
      by (auto simp: path_image_def ball_def)
  }
  then have ghs: path_image g ⊆ S path_image h ⊆ S
    by (auto simp: path_image_def)
  moreover
  { fix f
    assume fhols: f holomorphic_on S
    then have fpa: f contour_integrable_on g f contour_integrable_on h
      using g ghs h holomorphic_on_imp_continuous_on os contour_integrable_holomorphic_simple
      by blast+
    have contf: continuous_on S f
      by (simp add: fhols holomorphic_on_imp_continuous_on)
    { fix z
      assume z: z ∈ path_image p
      have f holomorphic_on ball z (ee z)
        using fhols ee z holomorphic_on_subset by blast
      then have ∃ ff. (∀ w ∈ ball z (ee z). (ff has_field_derivative f w) (at w))
        using holomorphic_convex_primitive [of ball z (ee z) {} f, simplified]
        by (metis open_ball at_within_open holomorphic_on_def holomorphic_on_imp_continuous_on mem_ball)
    }
    then obtain ff where ff:
      ⋀ z w. [z ∈ path_image p; w ∈ ball z (ee z)] ⇒ (ff z has_field_derivative
      f w) (at w)
      by metis
    { fix n
      assume n: n ≤ N
      then have contour_integral(subpath 0 (n/N) h) f - contour_integral(subpath
      0 (n/N) g) f =
        contour_integral(linepath (g(n/N)) (h(n/N))) f - con-
        tour_integral(linepath (g 0) (h 0)) f

```

```

proof (induct n)
  case 0 show ?case by simp
next
  case (Suc n)
  obtain t where t: t ∈ k and p (n/N) ∈ ball(p t) (ee(p t) / 3)
    using ⟨path_image p ⊆ ⋃ D⟩ [THEN subsetD, where c=p (n/N)] D_eq
N Suc.prem
    by (force simp: path_image_def)
  then have ptu: cmod (p t - p (n/N)) < ee (p t) / 3
    by (simp add: dist_norm)
  have e3le: e/3 ≤ ee (p t) / 3 using fin_eep t
    by (simp add: e_def)
  { fix x
    assume x: n/N ≤ x x ≤ (1 + n)/N
    then have nN01: 0 ≤ n/N (1 + n)/N ≤ 1
      using Suc.prem by auto
    then have x01: 0 ≤ x x ≤ 1
      using x by linarith+
    have cmod (p t - p x) < ee (p t) / 3 + e/3
    proof (rule norm_diff_triangle_less [OF ptu de])
      show |real n / real N - x| < d
        using x N by (auto simp: field_simps)
    qed (use x01 Suc.prem in auto)
    then have ptx: cmod (p t - p x) < 2*ee (p t)/3
      using e3le eepi [OF t] by simp
    have cmod (p t - g x) < 2*ee (p t)/3 + e/3
      using ghp x01
    by (force simp add: norm_minus_commute intro!: norm_diff_triangle_less
[OF ptx])
    also have ... ≤ ee (p t)
      using e3le eepi [OF t] by simp
    finally have gg: cmod (p t - g x) < ee (p t) .
    have cmod (p t - h x) < 2*ee (p t)/3 + e/3
      using ghp x01
    by (force simp add: norm_minus_commute intro!: norm_diff_triangle_less
[OF ptx])
    also have ... ≤ ee (p t)
      using e3le eepi [OF t] by simp
    finally have cmod (p t - g x) < ee (p t) cmod (p t - h x) < ee (p t)
      using gg by auto
  } note ptgh_ee = this
  have closed_segment (g (n/N)) (h (n/N)) = path_image (linepath (h
(n/N)) (g (n/N)))
    by (simp add: closed_segment_commute)
  also have pi_hgn: ... ⊆ ball (p t) (ee (p t))
    using ptgh_ee [of n/N] Suc.prem
  by (auto simp: field_simps dist_norm dest: segment_furthest_le [where
y=p t])
  finally have gh_ns: closed_segment (g (n/N)) (h (n/N)) ⊆ S

```



```

    using ee pi t by blast
  have pi_ghn': path_image (linepath (g ((1 + n) / N)) (h ((1 + n) / N)))
 $\subseteq$  ball (p t) (ee (p t))
    using ptgh_ee [of (1+n)/N] Suc.prem
  by (auto simp: field_simps dist_norm dest: segment_furthest_le [where
y=p t])
  then have gh_n's: closed_segment (g ((1 + n) / N)) (h ((1 + n) / N))
 $\subseteq$  S
    using <N>0 Suc.prem ee pi t
    by (auto simp: Path_Connected.path_image_join field_simps)
  have pi_subset_ball:
    path_image (subpath (n/N) ((1+n) / N) g +++ linepath (g ((1+n)
/ N)) (h ((1+n) / N)) +++
    subpath ((1+n) / N) (n/N) h +++ linepath (h (n/N)) (g
(n/N)))
 $\subseteq$  ball (p t) (ee (p t))
  proof (intro subset_path_image_join pi_hgn pi_ghn')
    show path_image (subpath (n/N) ((1+n) / N) g)  $\subseteq$  ball (p t) (ee (p t))
      path_image (subpath ((1+n) / N) (n/N) h)  $\subseteq$  ball (p t) (ee (p t))
    using <N>0 Suc.prem
    by (auto simp: path_image_subpath dist_norm field_simps ptgh_ee)
  qed
  have pi0: (f has_contour_integral 0)
    (subpath (n / N) ((Suc n)/N) g +++ linepath(g ((Suc n) / N))
(h((Suc n) / N)) +++
    subpath ((Suc n) / N) (n/N) h +++ linepath(h (n/N)) (g
(n/N)))
  proof (rule Cauchy_theorem_primitive)
    show  $\bigwedge x. x \in$  ball (p t) (ee (p t))
 $\implies$  (ff (p t) has_field_derivative f x) (at x within ball (p t) (ee
(p t)))
    by (metis ff_open_ball at_within_open pi t)
  qed (use Suc.prem pi_subset_ball in <simp_all add: valid_path_subpath
g h>)
  have fpa1: f contour_integrable_on subpath (n/N) (real (Suc n) / real N)
g
    using Suc.prem by (simp add: contour_integrable_subpath g fpa)
  have fpa2: f contour_integrable_on linepath (g (real (Suc n) / real N)) (h
(real (Suc n) / real N))
    using gh_n's
    by (auto intro!: contour_integrable_continuous_linepath continu-
ous_on_subset [OF contf])
  have fpa3: f contour_integrable_on linepath (h (n/N)) (g (n/N))
    using gh_ns
    by (auto simp: closed_segment_commute intro!: contour_integrable_continuous_linepath
continuous_on_subset [OF contf])
  have eq0: contour_integral (subpath (n/N) ((Suc n) / real N) g) f +
    contour_integral (linepath (g ((Suc n) / N)) (h ((Suc n) / N))) f
+

```

```

      contour_integral (subpath ((Suc n) / N) (n/N) h) f +
      contour_integral (linepath (h (n/N)) (g (n/N))) f = 0
    using contour_integral_unique [OF pi0] Suc.prem_s
  by (simp add: g h fpa valid_path_subpath contour_integrable_subpath
        fpa1 fpa2 fpa3 algebra_simps del: of_nat_Suc)
  have *:  $\bigwedge hn\ he\ hn'\ gn\ gd\ gn'\ hgn\ ghn\ gh0\ ghn'.$ 
     $[[hn - gn = ghn - gh0;$ 
       $gd + ghn' + he + hgn = (0::complex);$ 
       $hn - he = hn'; gn + gd = gn'; hgn = -ghn]] \implies hn' - gn' =$ 
     $ghn' - gh0$ 
  by (auto simp: algebra_simps)
  have contour_integral (subpath 0 (n/N) h) f - contour_integral (subpath
    ((Suc n) / N) (n/N) h) f =
    contour_integral (subpath 0 (n/N) h) f + contour_integral (subpath
    (n/N) ((Suc n) / N) h) f
  unfolding reversepath_subpath [symmetric, of ((Suc n) / N)]
  using Suc.prem_s by (simp add: h fpa contour_integral_reversepath
    valid_path_subpath contour_integrable_subpath)
  also have ... = contour_integral (subpath 0 ((Suc n) / N) h) f
  using Suc.prem_s by (simp add: contour_integral_subpath_combine h
    fpa)
  finally have pi0_eq:
    contour_integral (subpath 0 (n/N) h) f - contour_integral (subpath
    ((Suc n) / N) (n/N) h) f =
    contour_integral (subpath 0 ((Suc n) / N) h) f .
  show ?case
  proof (rule * [OF Suc.hyps eq0 pi0_eq])
    show contour_integral (subpath 0 (n/N) g) f +
      contour_integral (subpath (n/N) ((Suc n) / N) g) f =
      contour_integral (subpath 0 ((Suc n) / N) g) f
    using Suc.prem_s contour_integral_subpath_combine fpa(1) g by auto
    show contour_integral (linepath (h (n/N)) (g (n/N))) f = - contour_integral
      (linepath (g (n/N)) (h (n/N))) f
    by (metis contour_integral_unique fpa3 has_contour_integral_integral
      has_contour_integral_reverse_linepath)
  qed (use Suc.prem_s in auto)
  qed
} note ind = this
have contour_integral h f = contour_integral g f
  using ind [OF order_refl] N joins
by (simp add: linked_paths_def pathstart_def pathfinish_def split: if_split_asm)
}
ultimately
show path_image g  $\subseteq$  S  $\wedge$  path_image h  $\subseteq$  S  $\wedge$  ( $\forall f. f$  holomorphic_on S  $\longrightarrow$ 
contour_integral h f = contour_integral g f)
  by metis
qed
qed

```

lemma

assumes *open S path p path_image p* $\subseteq S$

shows *contour_integral_nearby_ends*:

$\exists d. 0 < d \wedge$

$(\forall g h. \text{valid_path } g \wedge \text{valid_path } h \wedge$

$(\forall t \in \{0..1\}. \text{norm}(g\ t - p\ t) < d \wedge \text{norm}(h\ t - p\ t) < d) \wedge$

$\text{pathstart } h = \text{pathstart } g \wedge \text{pathfinish } h = \text{pathfinish } g$

$\longrightarrow \text{path_image } g \subseteq S \wedge$

$\text{path_image } h \subseteq S \wedge$

$(\forall f. f \text{ holomorphic_on } S$

$\longrightarrow \text{contour_integral } h\ f = \text{contour_integral } g\ f))$

and *contour_integral_nearby_loops*:

$\exists d. 0 < d \wedge$

$(\forall g h. \text{valid_path } g \wedge \text{valid_path } h \wedge$

$(\forall t \in \{0..1\}. \text{norm}(g\ t - p\ t) < d \wedge \text{norm}(h\ t - p\ t) < d) \wedge$

$\text{pathfinish } g = \text{pathstart } g \wedge \text{pathfinish } h = \text{pathstart } h$

$\longrightarrow \text{path_image } g \subseteq S \wedge$

$\text{path_image } h \subseteq S \wedge$

$(\forall f. f \text{ holomorphic_on } S$

$\longrightarrow \text{contour_integral } h\ f = \text{contour_integral } g\ f))$

using *contour_integral_nearby* [*OF assms*, **where** *atends=True*]

using *contour_integral_nearby* [*OF assms*, **where** *atends=False*]

unfolding *linked_paths_def* **by** *simp_all*

lemma *contour_integral_bound_exists*:

assumes *S*: *open S*

and *g*: *valid_path g*

and *pag*: *path_image g* $\subseteq S$

shows $\exists L. 0 < L \wedge$

$(\forall f B. f \text{ holomorphic_on } S \wedge (\forall z \in S. \text{norm}(f\ z) \leq B)$

$\longrightarrow \text{norm}(\text{contour_integral } g\ f) \leq L*B)$

proof –

have *path g using g*

by (*simp add: valid_path_imp_path*)

then obtain *d::real* **and** *p*

where *d*: $0 < d$

and *p*: *polynomial_function p path_image p* $\subseteq S$

and *pi*: $\bigwedge f. f \text{ holomorphic_on } S \implies \text{contour_integral } g\ f = \text{contour_integral}$

p f

using *contour_integral_nearby_ends* [*OF S* \langle *path g* \rangle *pag*]

by (*metis cancel_comm_monoid_add_class.diff_cancel g norm_zero path_approx_polynomial_function valid_path_polynomial_function*)

then obtain *p'* **where** *p'*: *polynomial_function p'*

$\bigwedge x. (p \text{ has_vector_derivative } (p'\ x)) \text{ (at } x)$

by (*blast intro: has_vector_derivative_polynomial_function that*)

then have *bounded(p' ' {0..1})*

using *continuous_on_polynomial_function*

by (*force simp: intro!: compact_imp_bounded compact_continuous_image*)

```

then obtain  $L$  where  $L: L > 0$  and  $nop'$ :  $\bigwedge x. \llbracket 0 \leq x; x \leq 1 \rrbracket \implies \text{norm } (p' x) \leq L$ 
by (force simp: bounded_pos)
{ fix  $f B$ 
assume  $f: f$  holomorphic_on  $S$  and  $B: \bigwedge z. z \in S \implies \text{cmod } (f z) \leq B$ 
then have  $f$  contour_integrable_on  $p \wedge$  valid_path  $p$ 
using  $p S$ 
by (blast intro: valid_path_polynomial_function contour_integrable_holomorphic_simple
holomorphic_on_imp_continuous_on)
moreover have  $\text{cmod } (\text{vector\_derivative } p \text{ (at } x)) * \text{cmod } (f (p x)) \leq L * B$ 
if  $0 \leq x \leq 1$  for  $x$ 
proof (rule mult_mono)
show  $\text{cmod } (\text{vector\_derivative } p \text{ (at } x)) \leq L$ 
by (metis  $nop'$   $p'(2)$  that vector_derivative_at)
show  $\text{cmod } (f (p x)) \leq B$ 
by (metis  $B$  atLeastAtMost_iff imageI  $p(2)$  path_defs(4) subset_eq that)
qed (use  $\langle L > 0 \rangle$  in auto)
ultimately
have  $\text{cmod } (\text{integral } \{0..1\} (\lambda x. f (p x) * \text{vector\_derivative } p \text{ (at } x))) \leq L * B$ 
by (intro order_trans [OF integral_norm_bound_integral])
(auto simp: mult.commute norm_mult contour_integrable_on)
then have  $\text{cmod } (\text{contour\_integral } g f) \leq L * B$ 
using contour_integral_integral  $f pi$  by presburger
} then
show ?thesis using  $\langle L > 0 \rangle$ 
by (intro exI[of _  $L$ ]) auto
qed

```

2.8 Homotopy forms of Cauchy's theorem

lemma Cauchy_theorem_homotopic:

assumes hom : if $atends$ then homotopic_paths $S g h$ else homotopic_loops $S g h$

and open S **and** $f: f$ holomorphic_on S

and vpg : valid_path g **and** vph : valid_path h

shows contour_integral $g f =$ contour_integral $h f$

proof –

have $pathsf$: linked_paths $atends g h$

using hom **by** (auto simp: linked_paths_def homotopic_paths_imp_pathstart homotopic_paths_imp_pathfinish homotopic_loops_imp_loop)

obtain $k :: \text{real} \times \text{real} \Rightarrow \text{complex}$

where $contk$: continuous_on $(\{0..1\} \times \{0..1\}) k$

and ks : $k \text{ ' } (\{0..1\} \times \{0..1\}) \subseteq S$

and k [simp]: $\forall x. k (0, x) = g x \ \forall x. k (1, x) = h x$

and ksf : $\forall t \in \{0..1\}. \text{linked_paths } atends g (\lambda x. k (t, x))$

using $hom pathsf$ **by** (auto simp: linked_paths_def homotopic_paths_def homotopic_loops_def homotopic_with_def split: if_split_asm)

have $ucontk$: uniformly_continuous_on $(\{0..1\} \times \{0..1\}) k$

by (blast intro: compact_Times compact_uniformly_continuous [OF $contk$])

```

{ fix t::real assume t: t ∈ {0..1}
  have Pair t ' {0..1} ⊆ {0..1} × {0..1}
    using t by force
  then have pak: path (k ∘ (λu. (t, u)))
    unfolding path_def
    by (intro continuous_intros continuous_on_subset [OF contk])+
  have pik: path_image (k ∘ Pair t) ⊆ S
    using ks t by (auto simp: path_image_def)
  obtain e where e>0 and e:
    ∧g h. [valid_path g; valid_path h;
      ∀u∈{0..1}. cmod (g u - (k ∘ Pair t) u) < e ∧ cmod (h u - (k ∘
Pair t) u) < e;
      linked_paths attends g h]
    ⇒⇒ contour_integral h f = contour_integral g f
    using contour_integral_nearby [OF ‹open S› pak pik, of attends] f by metis
  obtain d where d>0 and d:
    ∧x x'. [x ∈ {0..1} × {0..1}; x' ∈ {0..1} × {0..1}; norm (x'-x) < d] ⇒⇒
norm (k x' - k x) < e/4
    by (rule uniformly_continuous_onE [OF ucontk, of e/4]) (auto simp: dist_norm
‹e>0)
    { fix t1 t2
      assume t1: 0 ≤ t1 t1 ≤ 1 and t2: 0 ≤ t2 t2 ≤ 1 and ltd: |t1 - t| < d |t2
- t| < d
      have no2: norm(g1 - kt) < e if norm(g1 - k1) < e/4 norm(k1 - kt) <
e/4 for g1 k1 kt :: complex
      proof (rule norm_triangle_half_1)
        show cmod (g1 - k1) < e/2 cmod (kt - k1) < e/2
          using ‹e > 0› that by (auto simp: norm_minus_commute intro: or-
der_less_trans)
        qed
      have ∃d>0. ∀g1 g2. valid_path g1 ∧ valid_path g2 ∧
        (∀u∈{0..1}. cmod (g1 u - k (t1, u)) < d ∧ cmod (g2 u - k
(t2, u)) < d) ∧
          linked_paths attends g1 g2 →
          contour_integral g2 f = contour_integral g1 f
      using t t1 t2 ltd ‹e > 0›
      by (rule_tac x=e/4 in exI) (auto intro!: e simp: d no2 simp del: less_divide_eq_numeral1)
    }
  then have ∃e. 0 < e ∧
    (∀t1 t2. t1 ∈ {0..1} ∧ t2 ∈ {0..1} ∧ |t1 - t| < e ∧ |t2 - t| < e
    → (∃d. 0 < d ∧
      (∀g1 g2. valid_path g1 ∧ valid_path g2 ∧
        (∀u ∈ {0..1}.
          norm(g1 u - k((t1,u))) < d ∧ norm(g2 u - k((t2,u))) < d) ∧
          linked_paths attends g1 g2
          → contour_integral g2 f = contour_integral g1 f)))
    by (rule_tac x=d in exI) (simp add: ‹d > 0›)
  }
}
then obtain ee where ee:

```

```


$$\bigwedge t. t \in \{0..1\} \implies ee\ t > 0 \wedge$$


$$(\forall t1\ t2. t1 \in \{0..1\} \longrightarrow t2 \in \{0..1\} \longrightarrow |t1 - t| < ee\ t \longrightarrow |t2 - t| <$$


$$ee\ t$$


$$\longrightarrow (\exists d. 0 < d \wedge$$


$$(\forall g1\ g2. \text{valid\_path } g1 \wedge \text{valid\_path } g2 \wedge$$


$$(\forall u \in \{0..1\}. \text{norm}(g1\ u - k((t1,u))) < d \wedge \text{norm}(g2\ u - k((t2,u))) < d) \wedge$$


$$\text{linked\_paths\ attends } g1\ g2$$


$$\longrightarrow \text{contour\_integral } g2\ f = \text{contour\_integral } g1\ f))$$

by metis
note ee_rule = ee [THEN conjunct2, rule_format, of 0 0 0]
define C where C = ( $\lambda t. \text{ball } t\ (ee\ t / 3)$ ) ‘{0..1}
obtain C' where C': C'  $\subseteq$  C finite C' and C'01: {0..1}  $\subseteq$   $\bigcup$  C'
proof (rule compactE [OF compact_interval])
  show {0..1}  $\subseteq$   $\bigcup$  C
    using ee [THEN conjunct1] by (auto simp: C_def dist_norm)
qed (use C_def in auto)
define kk where kk = {t  $\in$  {0..1}. ball t (ee t / 3)  $\in$  C'}
have kk01: kk  $\subseteq$  {0..1} by (auto simp: kk_def)
define e where e = Min (ee ‘kk)
have C'_eq: C' = ( $\lambda t. \text{ball } t\ (ee\ t / 3)$ ) ‘kk
  using C' by (auto simp: kk_def C_def)
have ee_pos[simp]:  $\bigwedge t. t \in \{0..1\} \implies ee\ t > 0$ 
  by (simp add: kk_def ee)
moreover have finite kk
  using <finite C'> kk01 by (force simp: C'_eq inj_on_def ball_eq_ball_iff dest:
ee_pos finite_imageD)
moreover have kk  $\neq$  {} using <{0..1}  $\subseteq$   $\bigcup$  C'> C'_eq by force
ultimately have e > 0
  using finite_less_Inf_iff [of ee ‘kk 0] kk01 by (force simp: e_def)
then obtain N::nat where N > 0 and N: 1/N < e/3
  by (meson divide_pos_pos nat_approx_posE zero_less_Suc zero_less_numeral)
have e_le_ee:  $\bigwedge i. i \in kk \implies e \leq ee\ i$ 
  using <finite kk> by (simp add: e_def Min_le_iff [of ee ‘kk])
have plus:  $\exists t \in kk. x \in \text{ball } t\ (ee\ t / 3)$  if x  $\in$  {0..1} for x
  using C' subsetD [OF C'01 that] unfolding C'_eq by blast
have [OF order_refl]:
  
$$\exists d. 0 < d \wedge (\forall j. \text{valid\_path } j \wedge (\forall u \in \{0..1\}. \text{norm}(j\ u - k(n/N, u)) <$$


$$d) \wedge \text{linked\_paths\ attends } g\ j$$


$$\longrightarrow \text{contour\_integral } j\ f = \text{contour\_integral } g\ f)$$

  if n  $\leq$  N for n
using that
proof (induct n)
  case 0 show ?case
    using ee_rule
    by clarsimp (metis diff_self norm_eq_zero vpg)
next
  case (Suc n)
  then have N01: n/N  $\in$  {0..1} (Suc n)/N  $\in$  {0..1} by auto

```

```

then obtain  $t$  where  $t \in \mathbb{K}$   $n/N \in \text{ball } t \text{ (} ee \ t \ / \ 3)$ 
  using plus [of n/N] by blast
then have  $nN\_less: |n/N - t| < ee \ t$ 
  by (simp add: dist_norm del: less_divide_eq_numeral1)
have  $n'N\_less: |\text{real } (Suc \ n) / \text{real } N - t| < ee \ t$ 
  using  $t \ N \ \langle N > 0 \rangle \ e\_le\_ee$  [of  $t$ ]
  by (simp add: dist_norm add_divide_distrib abs_diff_less_iff del: less_divide_eq_numeral1)
(simp add: field_simps)
have  $t01: t \in \{0..1\}$  using  $\langle \mathbb{K} \subseteq \{0..1\} \rangle \ \langle t \in \mathbb{K} \rangle$  by blast
obtain  $d1$  where  $d1 > 0$  and  $d1$ :
   $\bigwedge g1 \ g2. \llbracket \text{valid\_path } g1; \text{valid\_path } g2; \forall u \in \{0..1\}. \text{cmod } (g1 \ u - k \ (n/N, \ u)) < d1 \wedge \text{cmod } (g2 \ u - k \ (Suc \ n) / N, \ u)) < d1; \llbracket \text{linked\_paths \ attends } g1 \ g2 \rrbracket \implies \text{contour\_integral } g2 \ f = \text{contour\_integral } g1 \ f$ 
  using  $ee$  [THEN conjunct2, rule_format, OF t01 N01 nN_less n'N_less] by
fastforce
have  $n \leq N$  using Suc.prems by auto
with Suc.hyps
obtain  $d2$  where  $d2 > 0$ 
  and  $d2: \bigwedge j. \llbracket \text{valid\_path } j; \forall u \in \{0..1\}. \text{cmod } (j \ u - k \ (n/N, \ u)) < d2; \llbracket \text{linked\_paths \ attends } g \ j \rrbracket \implies \text{contour\_integral } j \ f = \text{contour\_integral } g \ f$ 
  by auto
have Pair  $(n / N) \ \{0..1\} \subseteq \{0..1\} \times \{0..1\}$ 
  using N01 by auto
then have continuous_on  $\{0..1\} \ (k \circ (\lambda u. \ (n/N, \ u)))$ 
  by (intro continuous_intros continuous_on_subset [OF contk])
then have  $pkn: \text{path } (\lambda u. \ k \ (n/N, \ u))$ 
  by (simp add: path_def)
have  $min12: \min \ d1 \ d2 > 0$  by (simp add: <0 < d1> <0 < d2>)
obtain  $p$  where polynomial_function  $p$ 
  and  $psf: \text{pathstart } p = \text{pathstart } (\lambda u. \ k \ (n/N, \ u))$ 
   $\text{pathfinish } p = \text{pathfinish } (\lambda u. \ k \ (n/N, \ u))$ 
  and  $pk\_le: \bigwedge t. \ t \in \{0..1\} \implies \text{cmod } (p \ t - k \ (n/N, \ t)) < \min \ d1 \ d2$ 
  using path_approx_polynomial_function [OF pkn min12] by blast
then have  $vpp: \text{valid\_path } p$  using valid_path_polynomial_function by blast
have  $lpa: \text{linked\_paths \ attends } g \ p$ 
  by (metis (mono_tags, lifting) N01(1) ksf linked_paths_def pathfinish_def pathstart_def psf)
show ?case
proof (intro exI; safe)
  fix  $j$ 
  assume valid_path  $j$  linked_paths attends  $g \ j$ 
  and  $\forall u \in \{0..1\}. \text{cmod } (j \ u - k \ (\text{real } (Suc \ n) / \text{real } N, \ u)) < \min \ d1 \ d2$ 
  then have  $\text{contour\_integral } j \ f = \text{contour\_integral } p \ f$ 
  using  $pk\_le \ N01(1) \ ksf$  by (force intro!: vpp d1 simp add: linked_paths_def psf)
  also have  $\dots = \text{contour\_integral } g \ f$ 

```

```

    using pk_le by (force intro!: vpp d2 lpa)
    finally show contour_integral j f = contour_integral g f .
  qed (simp add: ⟨0 < d1⟩ ⟨0 < d2⟩)
  qed
  then obtain d where 0 < d
    ∧ j. valid_path j ∧ (∀ u ∈ {0..1}. norm(j u - k (1,u)) < d) ∧
  linked_paths attends g j
    ⇒ contour_integral j f = contour_integral g f
    using ⟨N>0⟩ by auto
  then have linked_paths attends g h ⇒ contour_integral h f = contour_integral
  g f
    using ⟨N>0⟩ vph by fastforce
  then show ?thesis
    by (simp add: pathsf)
  qed

```

proposition *Cauchy_theorem_homotopic_paths:*

```

  assumes hom: homotopic_paths S g h
    and open S and f: f holomorphic_on S
    and vpg: valid_path g and vph: valid_path h
  shows contour_integral g f = contour_integral h f
  using Cauchy_theorem_homotopic [of True S g h] assms by simp

```

proposition *Cauchy_theorem_homotopic_loops:*

```

  assumes hom: homotopic_loops S g h
    and open S and f: f holomorphic_on S
    and vpg: valid_path g and vph: valid_path h
  shows contour_integral g f = contour_integral h f
  using Cauchy_theorem_homotopic [of False S g h] assms by simp

```

lemma *has_contour_integral_newpath:*

```

  [(f has_contour_integral y) h; f contour_integrable_on g; contour_integral g f
  = contour_integral h f]
  ⇒ (f has_contour_integral y) g
  using has_contour_integral_integral contour_integral_unique by auto

```

lemma *Cauchy_theorem_null_homotopic:*

```

  [(f holomorphic_on S; open S; valid_path g; homotopic_loops S g (linepath a
  a))]
  ⇒ (f has_contour_integral 0) g
  by (metis Cauchy_theorem_homotopic_loops contour_integrable_holomorphic_simple
  valid_path_linepath
  contour_integral_trivial has_contour_integral_integral homotopic_loops_imp_subset)

```

end

3 Winding numbers

theory *Winding_Numbers*

imports *Cauchy_Integral_Theorem*
begin

3.1 Definition

definition *winding_number_prop* :: $[real \Rightarrow complex, complex, real, real \Rightarrow complex, complex] \Rightarrow bool$ **where**

winding_number_prop $\gamma z e p n \equiv$
 $valid_path\ p \wedge z \notin path_image\ p \wedge$
 $pathstart\ p = pathstart\ \gamma \wedge$
 $pathfinish\ p = pathfinish\ \gamma \wedge$
 $(\forall t \in \{0..1\}. norm(\gamma\ t - p\ t) < e) \wedge$
 $contour_integral\ p\ (\lambda w. 1/(w - z)) = 2 * pi * i * n$

definition *winding_number*:: $[real \Rightarrow complex, complex] \Rightarrow complex$ **where**

winding_number $\gamma z \equiv SOME\ n. \forall e > 0. \exists p. winding_number_prop\ \gamma z e p n$

lemma *winding_number*:

assumes $path\ \gamma z \notin path_image\ \gamma\ 0 < e$

shows $\exists p. winding_number_prop\ \gamma z e p (winding_number\ \gamma z)$

proof –

have $path_image\ \gamma \subseteq UNIV - \{z\}$

using *assms* **by** *blast*

then obtain d

where $d: d > 0$

and $pi_eq: \wedge h1\ h2. valid_path\ h1 \wedge valid_path\ h2 \wedge$

$(\forall t \in \{0..1\}. cmod\ (h1\ t - \gamma\ t) < d \wedge cmod\ (h2\ t - \gamma\ t) < d) \wedge$

$pathstart\ h2 = pathstart\ h1 \wedge pathfinish\ h2 = pathfinish\ h1 \longrightarrow$

$path_image\ h1 \subseteq UNIV - \{z\} \wedge path_image\ h2 \subseteq UNIV -$

$\{z\} \wedge$

$(\forall f. f\ holomorphic_on\ UNIV - \{z\} \longrightarrow contour_integral\ h2\ f$

$= contour_integral\ h1\ f)$

using *contour_integral_nearby_ends* [*of* $UNIV - \{z\}$] γ] *assms* **by** (*auto simp: open_delete*)

then obtain h **where** $h: polynomial_function\ h \wedge pathstart\ h = pathstart\ \gamma \wedge pathfinish\ h = pathfinish\ \gamma \wedge$

$(\forall t \in \{0..1\}. norm(h\ t - \gamma\ t) < d/2)$

using *path_approx_polynomial_function* [*OF* $\langle path\ \gamma \rangle$, *of* $d/2$] d **by** (*metis half_gt_zero_iff*)

define nn **where** $nn = 1/(2 * pi * i) * contour_integral\ h\ (\lambda w. 1/(w - z))$

have $\exists n. \forall e > 0. \exists p. winding_number_prop\ \gamma z e p n$

proof (*rule_tac* $x=nn$ **in** *exI*, *clarify*)

fix $e::real$

assume $e: e > 0$

obtain p **where** $p: polynomial_function\ p \wedge$

$pathstart\ p = pathstart\ \gamma \wedge pathfinish\ p = pathfinish\ \gamma \wedge (\forall t \in \{0..1\}.$

$cmod\ (p\ t - \gamma\ t) < \min\ e\ (d/2))$

using *path_approx_polynomial_function* [*OF* $\langle path\ \gamma \rangle$, *of* $\min\ e\ (d/2)$] d

```

<0<e>
  by (metis min_less_iff_conj zero_less_divide_iff zero_less_numeral)
  have ( $\lambda w. 1 / (w - z)$ ) holomorphic_on UNIV - {z}
  by (auto simp: intro!: holomorphic_intros)
  then have winding_number_prop  $\gamma z e p nn$ 
  using pi_eq [of h p] h p d
  by (auto simp: valid_path_polynomial_function norm_minus_commute
nn_def winding_number_prop_def)
  then show  $\exists p. \text{winding\_number\_prop } \gamma z e p nn$ 
  by metis
qed
then show ?thesis
  unfolding winding_number_def by (rule someI2_ex) (blast intro: <0<e>)
qed

```

```

lemma winding_number_unique:
  assumes  $\gamma: \text{path } \gamma z \notin \text{path\_image } \gamma$ 
  and pi:  $\bigwedge e. e > 0 \implies \exists p. \text{winding\_number\_prop } \gamma z e p n$ 
  shows winding_number  $\gamma z = n$ 
proof -
  have path_image  $\gamma \subseteq \text{UNIV} - \{z\}$ 
  using assms by blast
  then obtain e
  where e:  $e > 0$ 
  and pi_eq:  $\bigwedge h1 h2 f. [\text{valid\_path } h1; \text{valid\_path } h2;
(\forall t \in \{0..1\}. \text{cmod } (h1 t - \gamma t) < e \wedge \text{cmod } (h2 t - \gamma t) < e);
\text{pathstart } h2 = \text{pathstart } h1; \text{pathfinish } h2 = \text{pathfinish } h1; f
\text{holomorphic\_on } \text{UNIV} - \{z\}] \implies
\text{contour\_integral } h2 f = \text{contour\_integral } h1 f$ 
  using contour_integral_nearby_ends [of UNIV - {z}  $\gamma$ ] assms by (auto simp:
open_delete)
  obtain p where p: winding_number_prop  $\gamma z e p n$ 
  using pi [OF e] by blast
  obtain q where q: winding_number_prop  $\gamma z e q (\text{winding\_number } \gamma z)$ 
  using winding_number [OF  $\gamma e$ ] by blast
  have  $2 * \text{complex\_of\_real } pi * i * n = \text{contour\_integral } p (\lambda w. 1 / (w - z))$ 
  using p by (auto simp: winding_number_prop_def)
  also have  $\dots = \text{contour\_integral } q (\lambda w. 1 / (w - z))$ 
  proof (rule pi_eq)
    show  $(\lambda w. 1 / (w - z)) \text{holomorphic\_on } \text{UNIV} - \{z\}$ 
    by (auto intro!: holomorphic_intros)
  qed (use p q in <auto simp: winding_number_prop_def norm_minus_commute>)
  also have  $\dots = 2 * \text{complex\_of\_real } pi * i * \text{winding\_number } \gamma z$ 
  using q by (auto simp: winding_number_prop_def)
  finally have  $2 * \text{complex\_of\_real } pi * i * n = 2 * \text{complex\_of\_real } pi * i *
\text{winding\_number } \gamma z .$ 
  then show ?thesis
  by simp
qed

```

lemma *winding_number_prop_reversepath*:
assumes *winding_number_prop* $\gamma z e p n$
shows *winding_number_prop* (*reversepath* γ) $z e$ (*reversepath* p) $(-n)$
proof –
have p : *valid_path* $p z \notin \text{path_image } p$ $\text{pathstart } p = \text{pathstart } \gamma$
 $\text{pathfinish } p = \text{pathfinish } \gamma \wedge t. t \in \{0..1\} \implies \text{norm } (\gamma t - p t) < e$
 $\text{contour_integral } p (\lambda w. 1 / (w - z)) = 2 * \text{complex_of_real } \pi i * i * n$
using *assms* **by** (*auto simp: winding_number_prop_def*)
show *?thesis*
unfolding *winding_number_prop_def*
proof (*intro conjI strip*)
show $\text{norm } (\text{reversepath } \gamma t - \text{reversepath } p t) < e$ **if** $t \in \{0..1\}$ **for** t
unfolding *reversepath_def* **using** $p(5)[\text{of } 1 - t]$ **that** **by** *auto*
show $\text{contour_integral } (\text{reversepath } p) (\lambda w. 1 / (w - z)) =$
 $\text{complex_of_real } (2 * \pi i) * i * -n$
using p **by** (*subst contour_integral_reversepath*) *auto*
qed (*use p in auto*)
qed

lemma *winding_number_prop_reversepath_iff*:
winding_number_prop (*reversepath* γ) $z e p n \longleftrightarrow \text{winding_number_prop } \gamma z e$
(*reversepath* p) $(-n)$
using *winding_number_prop_reversepath*[*of reversepath* $\gamma z e p n$]
winding_number_prop_reversepath[*of* $\gamma z e \text{reversepath } p -n$] **by** *auto*

lemma *winding_number_unique_loop*:
assumes γ : *path* $\gamma z \notin \text{path_image } \gamma$
and *loop*: *pathfinish* $\gamma = \text{pathstart } \gamma$
and π :
 $\wedge e. e > 0 \implies \exists p. \text{valid_path } p \wedge z \notin \text{path_image } p \wedge$
 $\text{pathfinish } p = \text{pathstart } p \wedge$
 $(\forall t \in \{0..1\}. \text{norm } (\gamma t - p t) < e) \wedge$
 $\text{contour_integral } p (\lambda w. 1 / (w - z)) = 2 * \pi i * i * n$
shows *winding_number* $\gamma z = n$
proof –
have $\text{path_image } \gamma \subseteq \text{UNIV} - \{z\}$
using *assms* **by** *blast*
then obtain e
where $e: e > 0$
and π : $\wedge h1 h2 f. [\text{valid_path } h1; \text{valid_path } h2;$
 $(\forall t \in \{0..1\}. \text{cmod } (h1 t - \gamma t) < e \wedge \text{cmod } (h2 t - \gamma t) < e);$
 $\text{pathfinish } h1 = \text{pathstart } h1; \text{pathfinish } h2 = \text{pathstart } h2; f$
holomorphic_on UNIV - {z}] \implies
 $\text{contour_integral } h2 f = \text{contour_integral } h1 f$
using *contour_integral_nearby_loops* [*of UNIV - {z}*] γ *assms* **by** (*auto simp:*
open_delete)
obtain p **where** p :

```

    valid_path p ∧ z ∉ path_image p ∧ pathfinish p = pathstart p ∧
    (∀ t ∈ {0..1}. norm (γ t - p t) < e) ∧
    contour_integral p (λw. 1/(w - z)) = 2 * pi * i * n
  using pi [OF e] by blast
  obtain q where q: winding_number_prop γ z e q (winding_number γ z)
  using winding_number [OF γ e] by blast
  have 2 * complex_of_real pi * i * n = contour_integral p (λw. 1 / (w - z))
  using p by auto
  also have ... = contour_integral q (λw. 1 / (w - z))
  proof (rule pi_eq)
    show (λw. 1 / (w - z)) holomorphic_on UNIV - {z}
    by (auto intro!: holomorphic_intros)
  qed (use p q loop in ⟨auto simp: winding_number_prop_def norm_minus_commute⟩)
  also have ... = 2 * complex_of_real pi * i * winding_number γ z
  using q by (auto simp: winding_number_prop_def)
  finally have 2 * complex_of_real pi * i * n = 2 * complex_of_real pi * i *
winding_number γ z .
  then show ?thesis
  by simp
qed

```

proposition *winding_number_valid_path*:

```

  assumes valid_path γ z ∉ path_image γ
  shows winding_number γ z = 1/(2*pi*i) * contour_integral γ (λw. 1/(w - z))
  by (rule winding_number_unique)
  (use assms in ⟨auto simp: valid_path_imp_path winding_number_prop_def⟩)

```

proposition *has_contour_integral_winding_number*:

```

  assumes γ: valid_path γ z ∉ path_image γ
  shows ((λw. 1/(w - z)) has_contour_integral (2*pi*i*winding_number γ z))
  γ
  by (simp add: winding_number_valid_path has_contour_integral_integral con-
tour_integrable_inversediff assms)

```

lemma *winding_number_trivial* [simp]: $z \neq a \implies \text{winding_number}(\text{linepath } a \ a) z = 0$

```

  by (simp add: winding_number_valid_path)

```

lemma *winding_number_subpath_trivial* [simp]: $z \neq g \ x \implies \text{winding_number}(\text{subpath } x \ g) z = 0$

```

  by (simp add: path_image_subpath winding_number_valid_path)

```

lemma *winding_number_join*:

```

  assumes γ1: path γ1 z ∉ path_image γ1
    and γ2: path γ2 z ∉ path_image γ2
    and pathfinish γ1 = pathstart γ2
  shows winding_number(γ1 +++ γ2) z = winding_number γ1 z + wind-
ing_number γ2 z
  proof (rule winding_number_unique)

```

```

show  $\exists p. \text{winding\_number\_prop } (\gamma 1 \text{ +++ } \gamma 2) z e p$ 
       $(\text{winding\_number } \gamma 1 z + \text{winding\_number } \gamma 2 z) \text{ if } e > 0 \text{ for } e$ 
proof –
  obtain  $p1$  where  $\text{winding\_number\_prop } \gamma 1 z e p1$   $(\text{winding\_number } \gamma 1 z)$ 
    using  $\langle 0 < e \rangle \gamma 1 \text{ winding\_number by blast}$ 
  moreover
  obtain  $p2$  where  $\text{winding\_number\_prop } \gamma 2 z e p2$   $(\text{winding\_number } \gamma 2 z)$ 
    using  $\langle 0 < e \rangle \gamma 2 \text{ winding\_number by blast}$ 
  ultimately
  have  $\text{winding\_number\_prop } (\gamma 1 \text{ +++ } \gamma 2) z e (p1 \text{ +++ } p2)$   $(\text{winding\_number } \gamma 1 z + \text{winding\_number } \gamma 2 z)$ 
    using  $\text{assms}$ 
  apply  $(\text{simp add: winding\_number\_prop\_def not\_in\_path\_image\_join contour\_integrable\_inversediff algebra\_simps})$ 
  apply  $(\text{auto simp: joinpaths\_def})$ 
  done
  then show  $?thesis$ 
    by  $\text{blast}$ 
qed
qed  $(\text{use assms in } \langle \text{auto simp: not\_in\_path\_image\_join} \rangle)$ 

```

```

lemma  $\text{winding\_number\_reversepath}$ :
  assumes  $\text{path } \gamma z \notin \text{path\_image } \gamma$ 
  shows  $\text{winding\_number}(\text{reversepath } \gamma) z = - (\text{winding\_number } \gamma z)$ 
proof  $(\text{rule winding\_number\_unique})$ 
  show  $\exists p. \text{winding\_number\_prop } (\text{reversepath } \gamma) z e p$   $(- \text{winding\_number } \gamma z)$ 
if  $e > 0$  for  $e$ 
  proof –
  obtain  $p$  where  $\text{winding\_number\_prop } \gamma z e p$   $(\text{winding\_number } \gamma z)$ 
    using  $\langle 0 < e \rangle \text{assms winding\_number by blast}$ 
  then have  $\text{winding\_number\_prop } (\text{reversepath } \gamma) z e (\text{reversepath } p)$   $(- \text{winding\_number } \gamma z)$ 
    using  $\text{assms unfolding winding\_number\_prop\_def}$ 
  apply  $(\text{simp add: contour\_integral\_reversepath contour\_integrable\_inversediff valid\_path\_imp\_reverse})$ 
  apply  $(\text{auto simp: reversepath\_def})$ 
  done
  then show  $?thesis$ 
    by  $\text{blast}$ 
qed
qed  $(\text{use assms in auto})$ 

```

```

lemma  $\text{winding\_number\_shiftpath}$ :
  assumes  $\gamma: \text{path } \gamma z \notin \text{path\_image } \gamma$ 
  and  $\text{pathfinish } \gamma = \text{pathstart } \gamma a \in \{0..1\}$ 
  shows  $\text{winding\_number}(\text{shiftpath } a \gamma) z = \text{winding\_number } \gamma z$ 
proof  $(\text{rule winding\_number\_unique\_loop})$ 
  show  $\exists p. \text{valid\_path } p \wedge z \notin \text{path\_image } p \wedge \text{pathfinish } p = \text{pathstart } p \wedge$ 
     $(\forall t \in \{0..1\}. \text{cmod } (\text{shiftpath } a \gamma t - p t) < e) \wedge$ 

```

```

      contour_integral p ( $\lambda w. 1 / (w - z)$ ) =
      2 * pi * i * winding_number  $\gamma z$ 
    if  $e > 0$  for  $e$ 
  proof -
    obtain  $p$  where winding_number_prop  $\gamma z e p$  (winding_number  $\gamma z$ )
      using  $\langle 0 < e \rangle$  assms winding_number by blast
    then show ?thesis
      apply (rule_tac  $x = \text{shiftpath } a p$  in exI)
      using assms that
      apply (auto simp: winding_number_prop_def path_image_shiftpath pathfinish_shiftpath pathstart_shiftpath contour_integral_shiftpath)
      apply (simp add: shiftpath_def)
    done
  qed
qed (use assms in  $\langle \text{auto simp: path_shiftpath path_image_shiftpath pathfinish_shiftpath pathstart_shiftpath} \rangle$ )

```

lemma *winding_number_split_linepath*:

```

  assumes  $c \in \text{closed\_segment } a b$   $z \notin \text{closed\_segment } a b$ 
  shows  $\text{winding\_number}(\text{linepath } a b) z = \text{winding\_number}(\text{linepath } a c) z +$ 
 $\text{winding\_number}(\text{linepath } c b) z$ 
  proof -
    have  $z \notin \text{closed\_segment } a c$   $z \notin \text{closed\_segment } c b$ 
      using assms by (meson convex_contains_segment convex_segment_ends_in_segment subsetCE)+
    then show ?thesis
      using assms
      by (simp add: winding_number_valid_path contour_integral_split_linepath [symmetric] continuous_on_inversediff field_simps)
  qed

```

lemma *winding_number_cong*:

```

  ( $\bigwedge t. [0 \leq t; t \leq 1] \implies p t = q t$ )  $\implies \text{winding\_number } p z = \text{winding\_number } q z$ 
  by (simp add: winding_number_def winding_number_prop_def pathstart_def pathfinish_def)

```

lemma *winding_number_constI*:

```

  assumes  $c \neq z$  and  $g: \bigwedge t. [0 \leq t; t \leq 1] \implies g t = c$ 
  shows  $\text{winding\_number } g z = 0$ 
  proof -
    have  $\text{winding\_number } g z = \text{winding\_number } (\text{linepath } c c) z$ 
      using  $g$  winding_number_cong by fastforce
    moreover have  $\text{winding\_number } (\text{linepath } c c) z = 0$ 
      using  $\langle c \neq z \rangle$  by auto
    ultimately show ?thesis by auto
  qed

```

lemma *winding_number_offset*: $\text{winding_number } p z = \text{winding_number } (\lambda w. p$

```

w - z) 0
  unfolding winding_number_def
proof (intro ext arg_cong [where f = Eps] arg_cong [where f = All] imp_cong
  refl, safe)
  fix n e g
  assume 0 < e and g: winding_number_prop p z e g n
  then show  $\exists r. \text{winding\_number\_prop } (\lambda w. p w - z) 0 e r n$ 
    by (rule_tac x= $\lambda t. g t - z$  in exI)
      (force simp: winding_number_prop_def contour_integral_integral_valid_path_def
  path_defs
      vector_derivative_def has_vector_derivative_diff_const piecewise_C1_differentiable_diff C1_differentiable_imp_piecewise)
  next
  fix n e g
  assume 0 < e and g: winding_number_prop ( $\lambda w. p w - z$ ) 0 e g n
  then have winding_number_prop p z e ( $\lambda t. g t + z$ ) n
    apply (simp add: winding_number_prop_def contour_integral_integral_valid_path_def
  path_defs
      piecewise_C1_differentiable_add vector_derivative_def has_vector_derivative_add_const
  C1_differentiable_imp_piecewise)
    apply (force simp: algebra_simps)
    done
  then show  $\exists r. \text{winding\_number\_prop } p z e r n$ 
    by metis
qed

lemma winding_number_offset_NO_MATCH:
  NO_MATCH 0 z  $\implies \text{winding\_number } p z = \text{winding\_number } (\lambda w. p w - z) 0$ 
  using winding_number_offset by metis

lemma winding_number_negatepath:
  assumes  $\gamma: \text{valid\_path } \gamma$  and 0:  $0 \notin \text{path\_image } \gamma$ 
  shows  $\text{winding\_number}(uminus \circ \gamma) 0 = \text{winding\_number } \gamma 0$ 
proof -
  have (/) 1 contour_integrable_on  $\gamma$ 
    using 0  $\gamma$  contour_integrable_inversediff by fastforce
  then have (( $\lambda z. 1/z$ ) has_contour_integral contour_integral  $\gamma$  ((/) 1))  $\gamma$ 
    by (rule has_contour_integral_integral)
  then have (( $\lambda z. 1 / - z$ ) has_contour_integral - contour_integral  $\gamma$  ((/) 1))
 $\gamma$ 
    using has_contour_integral_neg by auto
  then have contour_integral (uminus  $\circ \gamma$ ) ((/) 1) =
    contour_integral  $\gamma$  ((/) 1)
    using  $\gamma$  by (simp add: contour_integral_unique has_contour_integral_negatepath)
  then show ?thesis
    using assms by (simp add: winding_number_valid_path_valid_path_negatepath
  image_def path_defs)
qed

```

```

lemma winding_number_cnj:
  assumes path  $\gamma$   $z \notin \text{path\_image } \gamma$ 
  shows winding_number (cnj  $\circ$   $\gamma$ ) (cnj  $z$ ) = -cnj (winding_number  $\gamma$   $z$ )
proof (rule winding_number_unique)
  show  $\exists p. \text{winding\_number\_prop (cnj } \circ \gamma) (cnj z) e p (-cnj (\text{winding\_number } \gamma z))$ 
  if  $e > 0$  for  $e$ 
  proof -
    from winding_number[OF assms(1,2)  $\langle e > 0 \rangle$ ]
    obtain  $p$  where winding_number_prop  $\gamma z e p (\text{winding\_number } \gamma z)$ 
    by blast
    then have  $p: \text{valid\_path } p z \notin \text{path\_image } p$ 
      pathstart  $p = \text{pathstart } \gamma$ 
      pathfinish  $p = \text{pathfinish } \gamma$ 
       $\bigwedge t. t \in \{0..1\} \implies \text{cmod } (\gamma t - p t) < e$  and
       $p\_cont: \text{contour\_integral } p (\lambda w. 1 / (w - z)) =$ 
        complex_of_real (2 * pi) * i * winding_number  $\gamma z$ 
    unfolding winding_number_prop_def by auto

    have valid_path (cnj  $\circ$   $p$ )
      using  $p(1)$  by (subst valid_path_cnj) auto
    moreover have  $cnj z \notin \text{path\_image (cnj } \circ p)$ 
      using  $p(2)$  by (auto simp: path_image_def)
    moreover have pathstart (cnj  $\circ$   $p$ ) = pathstart (cnj  $\circ$   $\gamma$ )
      using  $p(3)$  by (simp add: pathstart_compose)
    moreover have pathfinish (cnj  $\circ$   $p$ ) = pathfinish (cnj  $\circ$   $\gamma$ )
      using  $p(4)$  by (simp add: pathfinish_compose)
    moreover have  $\text{cmod } ((cnj \circ \gamma) t - (cnj \circ p) t) < e$ 
      if  $t \in \{0..1\}$  for  $t$ 
    proof -
      have  $(cnj \circ \gamma) t - (cnj \circ p) t = cnj (\gamma t - p t)$ 
        by simp
      also have  $\text{norm } \dots = \text{norm } (\gamma t - p t)$ 
        by (subst complex_mod_cnj) auto
      also have  $\dots < e$ 
        using  $p(5)[OF t]$  by simp
      finally show ?thesis .
    qed
    moreover have  $\text{contour\_integral (cnj } \circ p) (\lambda w. 1 / (w - cnj z)) =$ 
       $cnj (\text{complex\_of\_real (2 * pi) * i * winding\_number } \gamma z)$  (is ?L=?R)
    proof -
      have ?L =  $\text{contour\_integral (cnj } \circ p) (cnj \circ (\lambda w. 1 / (cnj w - z)))$ 
        by (simp add: o_def)
      also have  $\dots = cnj (\text{contour\_integral } p (\lambda x. 1 / (x - z)))$ 
        using  $p(1)$  by (subst contour_integral_cnj) (auto simp: o_def)
      also have  $\dots = ?R$ 
        using  $p\_cont$  by simp
      finally show ?thesis .
    qed

```



```

ultimately show ?thesis
  by (intro exI[of _ cnj ∘ p]) (auto simp: winding_number_prop_def)
qed
show path (cnj ∘ γ)
  by (intro path_continuous_image continuous_intros) (use assms in auto)
show cnj z ∉ path_image (cnj ∘ γ)
  using ⟨z ∉ path_image γ⟩ unfolding path_image_def by auto
qed

```

A combined theorem deducing several things piecewise.

lemma *winding_number_join_pos_combined*:

```

[[valid_path γ1; z ∉ path_image γ1; 0 < Re(winding_number γ1 z);
  valid_path γ2; z ∉ path_image γ2; 0 < Re(winding_number γ2 z); pathfinish
γ1 = pathstart γ2]]
  ⇒ valid_path(γ1 +++ γ2) ∧ z ∉ path_image(γ1 +++ γ2) ∧ 0 <
Re(winding_number(γ1 +++ γ2) z)
  by (simp add: valid_path_join path_image_join winding_number_join valid_path_imp_path)

```

3.1.1 Useful sufficient conditions for the winding number to be positive

lemma *Re_winding_number*:

```

[[valid_path γ; z ∉ path_image γ]]
  ⇒ Re(winding_number γ z) = Im(contour_integral γ (λw. 1/(w - z))) /
(2*πi)
  by (simp add: winding_number_valid_path field_simps Re_divide_power2_eq_square)

```

lemma *winding_number_pos_le*:

```

assumes γ: valid_path γ z ∉ path_image γ
  and ge: ∧x. [[0 < x; x < 1]] ⇒ 0 ≤ Im (vector_derivative γ (at x) * cnj(γ
x - z))
  shows 0 ≤ Re(winding_number γ z)

```

proof –

```

have ge0: 0 ≤ Im (vector_derivative γ (at x) / (γ x - z)) if x: 0 < x x < 1
for x

```

```

  using ge by (simp add: Complex.Im_divide algebra_simps x)

```

```

let ?vd = λx. 1 / (γ x - z) * vector_derivative γ (at x)

```

```

let ?int = λz. contour_integral γ (λw. 1 / (w - z))

```

```

have 0 ≤ Im (?int z)

```

```

proof (rule has_integral_component_nonneg [of i, simplified])

```

```

  show ∧x. x ∈ cbox 0 1 ⇒ 0 ≤ Im (if 0 < x ∧ x < 1 then ?vd x else 0)

```

```

  by (force simp: ge0)

```

```

  have ((λa. 1 / (a - z)) has_contour_integral contour_integral γ (λw. 1 / (w
- z))) γ

```

```

  using γ by (simp flip: add: contour_integrable_inversediff has_contour_integral_integral)

```

```

  then have hi: (?vd has_integral ?int z) (cbox 0 1)

```

```

  using has_contour_integral by auto

```

```

  show ((λx. if 0 < x ∧ x < 1 then ?vd x else 0) has_integral ?int z) (cbox 0 1)

```

```

  by (rule has_integral_spike_interior [OF hi]) simp

```

```

qed
then show ?thesis
  by (simp add: Re_winding_number [OF  $\gamma$ ] field_simps)
qed

lemma winding_number_pos_lt_lemma:
  assumes  $\gamma$ : valid_path  $\gamma$   $z \notin \text{path\_image } \gamma$ 
    and  $e$ :  $0 < e$ 
    and  $ge$ :  $\bigwedge x. [0 < x; x < 1] \implies e \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) / (\gamma x - z))$ 
  shows  $0 < \text{Re}(\text{winding\_number } \gamma z)$ 
proof -
  let ?vd =  $\lambda x. 1 / (\gamma x - z) * \text{vector\_derivative } \gamma \text{ (at } x)$ 
  let ?int =  $\lambda z. \text{contour\_integral } \gamma (\lambda w. 1 / (w - z))$ 
  have  $e \leq \text{Im } (\text{contour\_integral } \gamma (\lambda w. 1 / (w - z)))$ 
  proof (rule has_integral_component_le [of  $i \lambda x. i * e \ i * e \ \{0..1\}$ , simplified])
    have  $((\lambda a. 1 / (a - z)) \text{ has\_contour\_integral } \text{contour\_integral } \gamma (\lambda w. 1 / (w - z))) \ \gamma$ 
    thm has_integral_component_le [of  $i \lambda x. i * e \ i * e \ \{0..1\}$ , simplified]
    using  $\gamma$  by (simp add: contour_integrable_inversediff has_contour_integral_integral)
    then have  $hi$ :  $(?vd \text{ has\_integral } ?int z) \text{ (cbox } 0 \ 1)$ 
    using has_contour_integral by auto
    show  $((\lambda x. \text{if } 0 < x \wedge x < 1 \text{ then } ?vd \ x \ \text{else } i * e) \text{ has\_integral } ?int z) \ \{0..1\}$ 
    by (rule has_integral_spike_interior [OF  $hi$ , simplified box_real]) (use  $e$  in simp)
  show  $\bigwedge x. 0 \leq x \wedge x \leq 1 \implies e \leq \text{Im } (\text{if } 0 < x \wedge x < 1 \text{ then } ?vd \ x \ \text{else } i * e)$ 
    by (simp add:  $ge$ )
  qed (use has_integral_const_real [of  $\_ \ 0 \ 1$ ] in auto)
  with  $e$  show ?thesis
  by (simp add: Re_winding_number [OF  $\gamma$ ] field_simps)
qed

```

```

lemma winding_number_pos_lt:
  assumes  $\gamma$ : valid_path  $\gamma$   $z \notin \text{path\_image } \gamma$ 
    and  $e$ :  $0 < e$ 
    and  $ge$ :  $\bigwedge x. [0 < x; x < 1] \implies e \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) * \text{cnj}(\gamma x - z))$ 
  shows  $0 < \text{Re} (\text{winding\_number } \gamma z)$ 
proof -
  have  $bm$ : bounded  $((\lambda w. w - z) \text{ ' (path\_image } \gamma))$ 
  using bounded_translation [of  $\_ -z$ ]  $\gamma$  by (simp add: bounded_valid_path_image)
  then obtain  $B$  where  $B$ :  $B > 0$  and  $Bno$ :  $\bigwedge x. x \in (\lambda w. w - z) \text{ ' (path\_image } \gamma) \implies \text{norm } x \leq B$ 
  using bounded_pos [THEN iffD1, OF  $bm$ ] by blast
  { fix  $x::\text{real}$  assume  $x$ :  $0 < x < 1$ 
    then have  $B2$ :  $\text{cmod } (\gamma x - z)^2 \leq B^2$  using  $Bno$  [of  $\gamma x - z$ ]
    by (simp add: path_image_def power2_eq_square mult_mono')
    with  $x$  have  $\gamma x \neq z$  using  $\gamma$ 
    using path_image_def by fastforce
  }

```

```

then have  $e / B^2 \leq e / (\text{cmod } (\gamma x - z))^2$ 
using  $B B^2 e$  by (auto simp: divide_left_mono)
also have  $\dots \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) * \text{cnj } (\gamma x - z)) / (\text{cmod } (\gamma x - z))^2$ 
using  $ge [OF x]$  by (auto simp: divide_right_mono)
finally have  $e / B^2 \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) * \text{cnj } (\gamma x - z)) / (\text{cmod } (\gamma x - z))^2$  .
then have  $e / B^2 \leq \text{Im } (\text{vector\_derivative } \gamma \text{ (at } x) / (\gamma x - z))$ 
by (simp add: complex_div_cnj [of _  $\gamma x - z$  for  $x$ ] del: complex_cnj_diff times_complex.sel)
} note  $*$  = this
show ?thesis
using  $e B$  by (simp add: * winding_number_pos_lt_lemma [OF  $\gamma$ , of  $e/B^2$ ])
qed

```

3.2 The winding number is an integer

Proof from the book Complex Analysis by Lars V. Ahlfors, Chapter 4, section 2.1, Also on page 134 of Serge Lang's book with the name title, etc.

lemma *exp_fg*:

```

fixes  $z::\text{complex}$ 
assumes  $g: (g \text{ has\_vector\_derivative } g') \text{ (at } x \text{ within } s)$ 
and  $f: (f \text{ has\_vector\_derivative } (g' / (g x - z))) \text{ (at } x \text{ within } s)$ 
and  $z: g x \neq z$ 
shows  $((\lambda x. \text{exp}(-f x) * (g x - z)) \text{ has\_vector\_derivative } 0) \text{ (at } x \text{ within } s)$ 
proof -
have  $*$ :  $(\text{exp} \circ (\lambda x. (- f x))) \text{ has\_vector\_derivative } - (g' / (g x - z)) * \text{exp}(-f x)$   $\text{ (at } x \text{ within } s)$ 
using assms unfolding has_vector_derivative_def scaleR_conv_of_real
by (auto intro!: derivative_eq_intros)
show ?thesis
using  $z$  by (auto intro!: derivative_eq_intros * [unfolded o_def] g)
qed

```

lemma *winding_number_exp_integral*:

```

fixes  $z::\text{complex}$ 
assumes  $\gamma: \gamma \text{ piecewise\_C1\_differentiable\_on } \{a..b\}$ 
and  $ab: a \leq b$ 
and  $z: z \notin \gamma ' \{a..b\}$ 
shows  $(\lambda x. \text{vector\_derivative } \gamma \text{ (at } x) / (\gamma x - z)) \text{ integrable\_on } \{a..b\}$ 
(is ?thesis1)
 $\text{exp}(-(\text{integral } \{a..b\} (\lambda x. \text{vector\_derivative } \gamma \text{ (at } x) / (\gamma x - z)))) * (\gamma b - z) = \gamma a - z$ 
(is ?thesis2)
proof -
let  $?D\gamma = \lambda x. \text{vector\_derivative } \gamma \text{ (at } x)$ 
have [simp]:  $\bigwedge x. a \leq x \implies x \leq b \implies \gamma x \neq z$ 
using  $z$  by force
have con_g: continuous_on  $\{a..b\} \gamma$ 

```

```

    using  $\gamma$  by (simp add: piecewise_C1_differentiable_on_def)
  obtain  $k$  where fink: finite  $k$  and  $g\_C1\_diff: \gamma$  C1_differentiable_on  $\{a..b\} - k$ 
  using  $\gamma$  by (force simp: piecewise_C1_differentiable_on_def)
  have  $\circ: open \{a<..b\} - k$ 
    using  $\langle$ finite  $k\rangle$  by (simp add: finite_imp_closed open_Diff)
  moreover have  $\{a<..b\} - k \subseteq \{a..b\} - k$ 
    by force
  ultimately have  $g\_diff\_at: \bigwedge x. \llbracket x \notin k; x \in \{a<..b\} \rrbracket \implies \gamma$  differentiable at  $x$ 
    by (metis Diff_iff differentiable_on_subset C1_diff_imp_diff [OF  $g\_C1\_diff$ ] differentiable_on_def at_within_open)
  { fix  $w$ 
    assume  $w \neq z$ 
    have continuous_on (ball  $w$  (cmod  $(w - z)$ ))  $(\lambda w. 1 / (w - z))$ 
      by (auto simp: dist_norm intro!: continuous_intros)
    moreover have  $\bigwedge x. cmod (w - x) < cmod (w - z) \implies \exists f'. ((\lambda w. 1 / (w - z)))$  has_field_derivative  $f'$  (at  $x$ )
      by (auto simp: intro!: derivative_eq_intros)
    ultimately have  $\exists h. \forall y. norm(y - w) < norm(w - z) \longrightarrow (h$  has_field_derivative  $1/(y - z))$  (at  $y$ )
      using holomorphic_convex_primitive [of ball  $w$  (norm $(w - z)$ )  $\{\}$   $\lambda w. 1/(w - z)$ ]
    by (force simp: field_differentiable_def Ball_def dist_norm at_within_open NO_MATCH norm_minus_commute)
  }
  then obtain  $h$  where  $h: \bigwedge w y. w \neq z \implies norm(y - w) < norm(w - z) \implies (h$  has_field_derivative  $1/(y - z))$  (at  $y$ )
    by meson
  have  $exy: \exists y. ((\lambda x. inverse (\gamma x - z) * ?D\gamma x)$  has_integral  $y)$   $\{a..b\}$ 
    unfolding integrable_on_def [symmetric]
  proof (rule contour_integral_local_primitive_any [OF piecewise_C1_imp_differentiable [OF  $\gamma$ ]])
    show  $\exists d h. 0 < d \wedge (\forall y. cmod (y - w) < d \longrightarrow (h$  has_field_derivative  $inverse (y - z))(at y$  within  $-\{z\}))$ 
      if  $w \in -\{z\}$  for  $w$ 
      using that inverse_eq_divide has_field_derivative_at_within  $h$ 
    by (metis Compl_insert DiffD2 insertCI right_minus_eq zero_less_norm_iff)
  qed simp
  have  $vg\_int: (\lambda x. ?D\gamma x / (\gamma x - z))$  integrable_on  $\{a..b\}$ 
    unfolding box_real [symmetric] divide_inverse_commute
  by (auto intro!: exy integrable_subinterval simp add: integrable_on_def ab)
  with  $ab$  show ?thesis1
    by (simp add: divide_inverse_commute integral_def integrable_on_def)
  { fix  $t$ 
    assume  $t: t \in \{a..b\}$ 
    have  $cball: continuous_on$  (ball  $(\gamma t)$  (dist  $(\gamma t)$   $z$ ))  $(\lambda x. inverse (x - z))$ 
      using  $\gamma$  by (auto intro!: continuous_intros simp: dist_norm)

```

```

have icd:  $\bigwedge x. \text{cmod } (\gamma t - x) < \text{cmod } (\gamma t - z) \implies (\lambda w. \text{inverse } (w - z))$ 
field_differentiable at x
unfolding field_differentiable_def by (force simp: intro!: derivative_eq_intros)
obtain h where h:  $\bigwedge x. \text{cmod } (\gamma t - x) < \text{cmod } (\gamma t - z) \implies$ 
 $(h \text{ has\_field\_derivative } \text{inverse } (x - z)) \text{ (at } x \text{ within } \{y. \text{cmod } (\gamma t - y) < \text{cmod } (\gamma t - z)\})$ 
using holomorphic_convex_primitive [where f =  $\lambda w. \text{inverse}(w - z)$ , OF
convex_ball finite.emptyI cball icd]
by simp (auto simp: ball_def dist_norm that)
have exp ( $-\text{integral } \{a..t\} (\lambda x. ?D\gamma x / (\gamma x - z))) * (\gamma t - z) = \gamma a - z$ 
proof (rule has_derivative_zero_unique_strong_interval [of {a,b}  $\cup$  k a b])
show continuous_on  $\{a..b\}$  ( $\lambda b. \text{exp } (-\text{integral } \{a..b\} (\lambda x. ?D\gamma x / (\gamma x - z))) * (\gamma b - z)$ )
by (auto intro!: continuous_intros con_g indefinite_integral_continuous_1
[OF vg_int])
show ( $(\lambda b. \text{exp } (-\text{integral } \{a..b\} (\lambda x. ?D\gamma x / (\gamma x - z))) * (\gamma b - z))$ 
has_derivative ( $\lambda h. 0$ ))
 $(\text{at } x \text{ within } \{a..b\})$ 
if  $x \in \{a..b\} - (\{a, b\} \cup k)$  for x
proof -
have x:  $x \notin k \wedge a < x < b$ 
using that by auto
then have x  $\in \text{interior } (\{a..b\} - k)$ 
using open_subset_interior [OF o] by fastforce
then have con: isCont  $?D\gamma x$ 
using g_C1_diff x by (auto simp: C1_differentiable_on_eq intro: continuous_on_interior)
then have con_vd: continuous ( $\text{at } x \text{ within } \{a..b\}$ ) ( $\lambda x. ?D\gamma x$ )
by (rule continuous_at_imp_continuous_within)
have gdx:  $\gamma$  differentiable at x
using x by (simp add: g_diff_at)
then obtain d where d: ( $\gamma$  has_derivative ( $\lambda x. x *_R d$ )) (at x)
by (auto simp add: differentiable_iff_scaleR)
show ( $(\lambda c. \text{exp } (-\text{integral } \{a..c\} (\lambda x. ?D\gamma x / (\gamma x - z))) * (\gamma c - z))$ 
has_derivative ( $\lambda h. 0$ ))
 $(\text{at } x \text{ within } \{a..b\})$ 
proof (rule exp_fg [unfolded has_vector_derivative_def, simplified])
show ( $\gamma$  has_derivative ( $\lambda c. c *_R d$ )) ( $\text{at } x \text{ within } \{a..b\}$ )
using d by (blast intro: has_derivative_at_withinI)
have ( $(\lambda x. \text{integral } \{a..x\} (\lambda x. ?D\gamma x / (\gamma x - z)))$  has_vector_derivative
 $d / (\gamma x - z)$ )
 $(\text{at } x \text{ within } \{a..b\})$ 
proof (rule has_vector_derivative_eq_rhs [OF integral_has_vector_derivative_continuous_at
[where S =  $\{\}$ , simplified]])
show continuous ( $\text{at } x \text{ within } \{a..b\}$ ) ( $\lambda x. \text{vector\_derivative } \gamma \text{ (at } x) /$ 
 $(\gamma x - z)$ )
using continuous_at_imp_continuous_at_within differentiable_imp_continuous_within
gdx x
by (intro con_vd continuous_intros) (force+)

```

```

      show vector_derivative  $\gamma$  (at  $x$ ) / ( $\gamma x - z$ ) =  $d$  / ( $\gamma x - z$ )
      using d vector_derivative_at
      by (simp add: vector_derivative_at has_vector_derivative_def)
    qed (use x vg_int in auto)
  then show (( $\lambda x$ . integral {a..x} ( $\lambda x$ . ?D $\gamma x$  / ( $\gamma x - z$ ))) has_derivative
( $\lambda c$ .  $c *_R (d / (\gamma x - z))$ ))
    (at  $x$  within {a..b})
    by (auto simp: has_vector_derivative_def)
  qed (use x in auto)
qed
qed
}
with ab show ?thesis2
  by (simp add: divide_inverse_commute_integral_def)
qed

```

lemma winding_number_exp_2pi:

```

  [[path  $p$ ;  $z \notin \text{path\_image } p$ ]
   $\implies \text{pathfinish } p - z = \exp (2 * \pi * i * \text{winding\_number } p z) * (\text{pathstart } p - z)$ 
  using winding_number [of  $p z 1$ ] unfolding valid_path_def path_image_def pathstart_def pathfinish_def winding_number_prop_def
  by (force dest: winding_number_exp_integral(2) [of _ 0 1  $z$ ] simp: field_simps contour_integral_integral_exp_minus)

```

lemma integer_winding_number_eq:

```

  assumes  $\gamma$ : path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$ 
  shows winding_number  $\gamma z \in \mathbb{Z} \iff \text{pathfinish } \gamma = \text{pathstart } \gamma$ 
  proof -
    obtain  $p$  where  $p$ : valid_path  $p z \notin \text{path\_image } p$ 
      pathstart  $p = \text{pathstart } \gamma$  pathfinish  $p = \text{pathfinish } \gamma$ 
      and eq: contour_integral  $p (\lambda w. 1 / (w - z)) = \text{complex\_of\_real } (2 * \pi) * i * \text{winding\_number } \gamma z$ 
    using winding_number [OF assms, of 1] unfolding winding_number_prop_def
    by auto
    then have wneq: winding_number  $\gamma z = \text{winding\_number } p z$ 
      using eq winding_number_valid_path by force
    have iff: (winding_number  $\gamma z \in \mathbb{Z}$ )  $\iff (\exp (\text{contour\_integral } p (\lambda w. 1 / (w - z))) = 1)$ 
      using eq by (simp add: exp_eq_1 complex_is_Int_iff)
    have  $\gamma 0 \neq z$ 
      by (metis pathstart_def pathstart_in_path_image  $z$ )
    then have exp (contour_integral  $p (\lambda w. 1 / (w - z))) = (\gamma 1 - z) / (\gamma 0 - z)$ 
      using  $p$  winding_number_exp_integral(2) [of  $p 0 1 z$ ]
      by (simp add: valid_path_def path_defs contour_integral_integral_exp_minus field_split_simps)
    then have winding_number  $p z \in \mathbb{Z} \iff \text{pathfinish } p = \text{pathstart } p$ 
      using  $p$  wneq iff by (auto simp: path_defs)
    then show ?thesis using  $p$  eq
  end

```

by (auto simp: winding_number_valid_path)
qed

theorem integer_winding_number:

$\llbracket \text{path } \gamma; \text{pathfinish } \gamma = \text{pathstart } \gamma; z \notin \text{path_image } \gamma \rrbracket \implies \text{winding_number } \gamma$
 $z \in \mathbb{Z}$

by (metis integer_winding_number_eq)

If the winding number's magnitude is at least one, then the path must contain points in every direction.*) We can thus bound the winding number of a path that doesn't intersect a given ray.

lemma winding_number_pos_meets:

fixes $z::\text{complex}$

assumes $\gamma: \text{valid_path } \gamma$ and $z: z \notin \text{path_image } \gamma$ and $1: \text{Re } (\text{winding_number } \gamma z) \geq 1$

and $w: w \neq z$

shows $\exists a::\text{real}. 0 < a \wedge z + \text{of_real } a * (w - z) \in \text{path_image } \gamma$

proof –

have [simp]: $\bigwedge x. 0 \leq x \implies x \leq 1 \implies \gamma x \neq z$

using z by (auto simp: path_image_def)

have [simp]: $z \notin \gamma \text{ ` } \{0..1\}$

using path_image_def z by auto

have gpd: $\gamma \text{ piecewise_C1_differentiable_on } \{0..1\}$

using $\gamma \text{ valid_path_def}$ by blast

define r where $r = (w - z) / (\gamma 0 - z)$

have [simp]: $r \neq 0$

using $w z$ by (auto simp: r_def)

have cont: continuous_on $\{0..1\}$

$(\lambda x. \text{Im } (\text{integral } \{0..x\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z))))$

by (intro continuous_intros indefinite_integral_continuous_1 winding_number_exp_integral [OF gpd]; simp)

have $\text{Arg}2\pi r \leq 2*\pi$

by (simp add: Arg2pi_less_eq_real_def)

also have $\dots \leq \text{Im } (\text{integral } \{0..1\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z)))$

using 1

by (simp add: winding_number_valid_path [OF γz] contour_integral_integral Complex.Re_divide field_simps power2_eq_square)

finally have $\text{Arg}2\pi r \leq \text{Im } (\text{integral } \{0..1\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z)))$.

then have $\exists t. t \in \{0..1\} \wedge \text{Im}(\text{integral } \{0..t\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z))) = \text{Arg}2\pi r$

by (simp add: Arg2pi_ge_0 cont IVT')

then obtain t where $t: t \in \{0..1\}$

and eqArg: $\text{Im } (\text{integral } \{0..t\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z))) = \text{Arg}2\pi r$

by blast

define i where $i = \text{integral } \{0..t\} (\lambda x. \text{vector_derivative } \gamma (at x) / (\gamma x - z))$

have gpd: $\gamma \text{ piecewise_C1_differentiable_on } \{0..t\}$

```

  by (metis atLeastAtMost_iff atLeastatMost_subset_iff order_refl piecewise_C1_differentiable_on_subset
  gpd t)
  have  $\exp(-i) * (\gamma t - z) = \gamma 0 - z$ 
    unfolding  $i\_def$ 
  proof (rule winding_number_exp_integral [OF gpd])
    show  $z \notin \gamma \text{ ` } \{0..t\}$ 
      using  $t z$  unfolding  $path\_image\_def$  by force
  qed (use  $t$  in auto)
  then have  $*: \gamma t - z = \exp i * (\gamma 0 - z)$ 
    by (simp add:  $\exp\_minus\_field\_simps$ )
  then have  $(w - z) = r * (\gamma 0 - z)$ 
    by (simp add:  $r\_def$ )
  moreover have  $z + \exp(\operatorname{Re} i) * (\exp(i * \operatorname{Im} i) * (\gamma 0 - z)) = \gamma t$ 
    using  $*$  by (simp add:  $\exp\_eq\_polar\_field\_simps$ )
  moreover have  $\operatorname{Arg} 2\pi r = \operatorname{Im} i$ 
    using  $\operatorname{eqArg}$  by (simp add:  $i\_def$ )
  ultimately have  $z + \operatorname{complex\_of\_real}(\exp(\operatorname{Re} i)) * (w - z) / \operatorname{complex\_of\_real}(cmod r) = \gamma t$ 
    using  $\operatorname{Complex\_Transcendental.Arg} 2\pi\_eq$  [of  $r$ ]  $\langle r \neq 0 \rangle$ 
    by (metis  $\operatorname{mult.left\_commute\_nonzero\_mult\_div\_cancel\_left}$   $\operatorname{norm\_eq\_zero}$ 
     $\operatorname{of\_real\_0}$   $\operatorname{of\_real\_eq\_iff\_times\_divide\_eq}$   $\operatorname{left}$ )
  with  $t$  show ?thesis
    by (rule_tac  $x = \exp(\operatorname{Re} i) / \operatorname{norm} r$  in  $\operatorname{exI}$ ) (auto simp:  $path\_image\_def$ )
qed

lemma winding_number_big_meets:
  fixes  $z::\operatorname{complex}$ 
  assumes  $\gamma$ :  $\operatorname{valid\_path} \gamma$  and  $z$ :  $z \notin \operatorname{path\_image} \gamma$  and  $|\operatorname{Re}(\operatorname{winding\_number} \gamma z)| \geq 1$ 
  and  $w$ :  $w \neq z$ 
  shows  $\exists a::\operatorname{real}. 0 < a \wedge z + \operatorname{of\_real} a * (w - z) \in \operatorname{path\_image} \gamma$ 
proof -
  { assume  $\operatorname{Re}(\operatorname{winding\_number} \gamma z) \leq -1$ 
    then have  $\operatorname{Re}(\operatorname{winding\_number}(\operatorname{reversepath} \gamma) z) \geq 1$ 
      by (simp add:  $\gamma \operatorname{valid\_path\_imp\_path} \operatorname{winding\_number\_reversepath} z$ )
    moreover have  $\operatorname{valid\_path}(\operatorname{reversepath} \gamma)$ 
      using  $\gamma \operatorname{valid\_path\_imp\_reverse}$  by auto
    moreover have  $z \notin \operatorname{path\_image}(\operatorname{reversepath} \gamma)$ 
      by (simp add:  $z$ )
    ultimately have  $\exists a::\operatorname{real}. 0 < a \wedge z + \operatorname{of\_real} a * (w - z) \in \operatorname{path\_image}(\operatorname{reversepath} \gamma)$ 
      using  $\operatorname{winding\_number\_pos\_meets} w$  by blast
    then have ?thesis
      by simp
  }
  then show ?thesis
    using  $\operatorname{assms}$ 
    by (simp add:  $\operatorname{abs\_if\_winding\_number\_pos\_meets}$   $\operatorname{split: if\_split\_asm}$ )
qed

```



```

lemma winding_number_less_1:
  fixes z::complex
  shows
    [[valid_path  $\gamma$ ;  $z \notin \text{path\_image } \gamma$ ;  $w \neq z$ ;
      $\bigwedge a::\text{real}. 0 < a \implies z + \text{of\_real } a * (w - z) \notin \text{path\_image } \gamma$ ]
      $\implies \text{Re}(\text{winding\_number } \gamma z) < 1$ 
  by (auto simp: not_less dest: winding_number_big_meets)

```

One way of proving that $\text{WN}=1$ for a loop.

```

lemma winding_number_eq_1:
  fixes z::complex
  assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and loop: pathfinish  $\gamma =$ 
  pathstart  $\gamma$ 
  and 0:  $0 < \text{Re}(\text{winding\_number } \gamma z)$  and 2:  $\text{Re}(\text{winding\_number } \gamma z) < 2$ 
  shows winding_number  $\gamma z = 1$ 
proof -
  have winding_number  $\gamma z \in \text{Ints}$ 
  by (simp add:  $\gamma$  integer_winding_number loop valid_path_imp_path z)
  then show ?thesis
  using 0 2 by (auto simp: Ints_def)
qed

```

3.3 Continuity of winding number and invariance on connected sets

```

theorem continuous_at_winding_number:
  fixes z::complex
  assumes  $\gamma$ : path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$ 
  shows continuous (at  $z$ ) (winding_number  $\gamma$ )
proof -
  obtain  $e$  where  $e > 0$  and  $cbg$ :  $\text{cball } z e \subseteq - \text{path\_image } \gamma$ 
  using open_contains_cball [of  $- \text{path\_image } \gamma$ ]  $z$ 
  by (force simp: closed_def [symmetric] closed_path_image [OF  $\gamma$ ])
  then have  $ppag$ :  $\text{path\_image } \gamma \subseteq - \text{cball } z (e/2)$ 
  by (force simp: cball_def dist_norm)
  have  $oc$ : open  $(- \text{cball } z (e/2))$ 
  by (simp add: closed_def [symmetric])
  obtain  $d$  where  $d > 0$  and  $pi\_eq$ :
     $\bigwedge h1 h2. [\text{valid\_path } h1; \text{valid\_path } h2;$ 
       $(\forall t \in \{0..1\}. \text{cmod } (h1 t - \gamma t) < d \wedge \text{cmod } (h2 t - \gamma t) < d);$ 
       $\text{pathstart } h2 = \text{pathstart } h1; \text{pathfinish } h2 = \text{pathfinish } h1]$ 
     $\implies$ 
       $\text{path\_image } h1 \subseteq - \text{cball } z (e/2) \wedge$ 
       $\text{path\_image } h2 \subseteq - \text{cball } z (e/2) \wedge$ 
       $(\forall f. f \text{ holomorphic\_on } - \text{cball } z (e/2) \longrightarrow \text{contour\_integral } h2 f =$ 
       $\text{contour\_integral } h1 f)$ 
  using contour_integral_nearby_ends [OF  $oc \ \gamma \ ppag$ ] by metis
  obtain  $p$  where valid_path  $p \ z \notin \text{path\_image } p$ 

```

```

    and p: pathstart p = pathstart  $\gamma$  pathfinish p = pathfinish  $\gamma$ 
    and pg:  $\bigwedge t. t \in \{0..1\} \implies \text{cmod } (\gamma t - p t) < \min d e/2$ 
    and pi: contour_integral p ( $\lambda x. 1 / (x - z)$ ) = complex_of_real (2 *
pi) * i * winding_number  $\gamma z$ 
    using winding_number [OF  $\gamma z$ , of min d e/2] <d>0> <e>0> by (auto simp:
winding_number_prop_def)
  { fix w
    assume d2: cmod (w - z) < d/2 and e2: cmod (w - z) < e/2
    have wnotp: w  $\notin$  path_image p
    proof (clarsimp simp add: path_image_def)
      show False if w: w = p x and 0  $\leq$  x  $\leq$  1 for x
      proof -
        have cmod ( $\gamma x - p x$ ) < min d e/2
          using pg that by auto
        then have cmod (z -  $\gamma x$ ) < e
          by (metis e2 less_divide_eq_numeral1(1) min_less_iff_conj norm_minus_commute
norm_triangle_half_l w)
        then show ?thesis
          using cbg that by (auto simp add: path_image_def cball_def dist_norm
less_eq_real_def)
      qed
    qed
    have wnotg: w  $\notin$  path_image  $\gamma$ 
      using cbg e2 <e>0> by (force simp: dist_norm norm_minus_commute)
    { fix k::real
      assume k: k > 0
      then obtain q where q: valid_path q w  $\notin$  path_image q
        pathstart q = pathstart  $\gamma$   $\wedge$  pathfinish q = pathfinish  $\gamma$ 
        and qg:  $\bigwedge t. t \in \{0..1\} \implies \text{cmod } (\gamma t - q t) < \min k (\min d e)$ 
/ 2
        and qi: contour_integral q ( $\lambda u. 1 / (u - w)$ ) = complex_of_real
(2 * pi) * i * winding_number  $\gamma w$ 
        using winding_number [OF  $\gamma$  wnotg, of min k (min d e) / 2] <d>0> <e>0>
k
        by (force simp: min_divide_distrib_right winding_number_prop_def)
      moreover have  $\bigwedge t. t \in \{0..1\} \implies \text{cmod } (q t - \gamma t) < d \wedge \text{cmod } (p t - \gamma
t) < d$ 
        using pg qg <0 < d> by (fastforce simp add: norm_minus_commute)
      moreover have ( $\lambda u. 1 / (u - w)$ ) holomorphic_on - cball z (e/2)
        using e2 by (auto simp: dist_norm norm_minus_commute intro!: holo-
morphic_intros)
      ultimately have contour_integral p ( $\lambda u. 1 / (u - w)$ ) = contour_integral q
( $\lambda u. 1 / (u - w)$ )
        by (metis p <valid_path p> pi_eq)
      then have contour_integral p ( $\lambda x. 1 / (x - w)$ ) = complex_of_real (2 * pi)
* i * winding_number  $\gamma w$ 
        by (simp add: pi qi)
    } note pip = this
    have path p

```

```

    by (simp add: ‹valid_path p› valid_path_imp_path)
  moreover have  $\bigwedge e. e > 0 \implies \text{winding\_number\_prop } p \ w \ e \ p \ (\text{winding\_number } \gamma \ w)$ 
    by (simp add: ‹valid_path p› pip_winding_number_prop_def wnotp)
  ultimately have  $\text{winding\_number } p \ w = \text{winding\_number } \gamma \ w$ 
    using winding_number_unique wnotp by blast
} note wnw = this
obtain pe where  $pe > 0$  and cbp:  $\text{cball } z \ (3 / 4 * pe) \subseteq - \text{path\_image } p$ 
using ‹valid_path p› ‹ $z \notin \text{path\_image } p$ › open_contains_cball [of - path_image p]
by (force simp: closed_def [symmetric] closed_path_image [OF valid_path_imp_path])
obtain L
where  $L > 0$ 
and L:  $\bigwedge f \ B. \llbracket f \text{ holomorphic\_on } - \text{cball } z \ (3 / 4 * pe); \forall z \in - \text{cball } z \ (3 / 4 * pe). \text{cmod } (f \ z) \leq B \rrbracket \implies \text{cmod } (\text{contour\_integral } p \ f) \leq L * B$ 
using contour_integral_bound_exists [of - cball z (3/4*pe) p] cbp ‹valid_path p› by blast
{ fix e::real and w::complex
  assume  $e: 0 < e$  and  $w: \text{cmod } (w - z) < pe/4$   $\text{cmod } (w - z) < e * pe^2 / (8 * L)$ 
  then have [simp]:  $w \notin \text{path\_image } p$ 
    using cbp p(2) ‹ $0 < pe$ ›
  by (force simp: dist_norm norm_minus_commute path_image_def cball_def)
  have [simp]:  $\text{contour\_integral } p \ (\lambda x. 1/(x - w)) - \text{contour\_integral } p \ (\lambda x. 1/(x - z)) =$ 
 $\text{contour\_integral } p \ (\lambda x. 1/(x - w) - 1/(x - z))$ 
  by (simp add: ‹valid_path p› ‹ $z \notin \text{path\_image } p$ › contour_integrable_inversediff contour_integral_diff)
  { fix x
    assume  $pe: 3/4 * pe < \text{cmod } (z - x)$ 
    have  $\text{cmod } (w - x) < pe/4 + \text{cmod } (z - x)$ 
    by (meson add_less_cancel_right norm_diff_triangle_le order_refl order_trans_rules(21) w(1))
    then have  $wx: \text{cmod } (w - x) < 4/3 * \text{cmod } (z - x)$  using pe by simp
    have  $\text{cmod } (z - x) \leq \text{cmod } (z - w) + \text{cmod } (w - x)$ 
    using norm_diff_triangle_le by blast
    also have  $\dots < pe/4 + \text{cmod } (w - x)$ 
    using w by (simp add: norm_minus_commute)
    finally have  $pe/2 < \text{cmod } (w - x)$ 
    using pe by auto
    then have  $pe\_less: (pe/2)^2 < \text{cmod } (w - x)^2$ 
    by (simp add: ‹ $0 < pe$ › less_eq_real_def power_strict_mono)
    then have  $pe2: pe^2 < 4 * \text{cmod } (w - x)^2$ 
    by (simp add: power_divide)
    have  $8 * L * \text{cmod } (w - z) < e * pe^2$ 
    using w ‹ $L > 0$ › by (simp add: field_simps)
    also have  $\dots < e * 4 * \text{cmod } (w - x) * \text{cmod } (w - x)$ 
    using pe2 ‹ $e > 0$ › by (simp add: power2_eq_square)

```

```

also have ... < e * 4 * cmod (w - x) * (4/3 * cmod (z - x))
  using <0 < pe> pe_less e_less_eq_real_def wx by fastforce
finally have L * cmod (w - z) < 2/3 * e * cmod (w - x) * cmod (z - x)
  by simp
also have ... ≤ e * cmod (w - x) * cmod (z - x)
  using e by simp
finally have Lwz: L * cmod (w - z) < e * cmod (w - x) * cmod (z - x) .
have L * cmod (1 / (x - w) - 1 / (x - z)) ≤ e
proof (cases x=z ∨ x=w)
  case True
  with pe <pe>0> w <L>0>
  show ?thesis
    by (force simp: norm_minus_commute)
  next
  case False
  with wx w(2) <L>0> pe pe2 Lwz
  show ?thesis
    by (auto simp: divide_simps mult_less_0_iff norm_minus_commute
norm_divide norm_mult_power2_eq_square)
  qed
} note L_cmod_le = this
let ?f = (λx. 1 / (x - w) - 1 / (x - z))
have cmod (contour_integral p ?f) ≤ L * (e * pe2 / L / 4 * (inverse (pe /
2))2)
proof (rule L)
  show ?f holomorphic_on - cball z (3 / 4 * pe)
  using <pe>0> w
  by (force simp: dist_norm norm_minus_commute intro!: holomorphic_intros)
  show ∀ u ∈ - cball z (3 / 4 * pe). cmod (?f u) ≤ e * pe2 / L / 4 * (inverse
(pe / 2))2
  using <pe>0> w <L>0>
  by (auto simp: cball_def dist_norm field_simps L_cmod_le simp del:
less_divide_eq_numeral1 le_divide_eq_numeral1)
  qed
also have ... < 2*e
  using <L>0> e by (force simp: field_simps)
finally have cmod (winding_number p w - winding_number p z) < e
  using pi_ge_two e
  by (force simp: winding_number_valid_path <valid_path p> <z ∉ path_image
p> field_simps norm_divide norm_mult intro: less_le_trans)
} note cmod_wn_diff = this
have isCont (winding_number p) z
proof (clarsimp simp add: continuous_at_eps_delta)
  fix e::real assume e>0
  show ∃ d>0. ∀ x'. dist x' z < d → dist (winding_number p x') (winding_number
p z) < e
  using <pe>0> <L>0> <e>0>
  by (rule_tac x=min (pe/4) (e/2*pe2/L/4) in exI) (simp add: dist_norm
cmod_wn_diff)

```

```

qed
then show ?thesis
  apply (rule continuous_transform_within [where  $\delta = \min d e/2$ ])
  apply (auto simp:  $\langle d > 0 \rangle \langle e > 0 \rangle \text{dist\_norm wnw}$ )
  done
qed

```

corollary *continuous_on_winding_number*:

```

path  $\gamma \implies \text{continuous\_on } (- \text{path\_image } \gamma) (\lambda w. \text{winding\_number } \gamma w)$ 
by (simp add: continuous_at_imp_continuous_on continuous_at_winding_number)

```

3.4 The winding number is constant on a connected region

lemma *winding_number_constant*:

```

assumes  $\gamma$ : path  $\gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$  and cs: connected  $S$ 
and sg:  $S \cap \text{path\_image } \gamma = \{\}$ 

```

```

shows winding_number  $\gamma$  constant_on  $S$ 

```

proof –

```

have *:  $1 \leq \text{cmod } (\text{winding\_number } \gamma y - \text{winding\_number } \gamma z)$ 

```

```

  if ne: winding_number  $\gamma y \neq \text{winding\_number } \gamma z$  and  $y \in S z \in S$  for  $y z$ 

```

proof –

```

  have winding_number  $\gamma y \in \mathbb{Z}$  winding_number  $\gamma z \in \mathbb{Z}$ 

```

```

  using that integer_winding_number [OF  $\gamma$  loop] sg  $\langle y \in S \rangle$  by auto

```

```

  with ne show ?thesis

```

```

  by (auto simp: Ints_def simp flip: of_int_diff)

```

qed

```

have cont: continuous_on  $S$  ( $\lambda w. \text{winding\_number } \gamma w$ )

```

```

  using continuous_on_winding_number [OF  $\gamma$ ] sg

```

```

  by (meson continuous_on_subset disjoint_eq_subset_Compl)

```

show ?thesis

```

  using * zero_less_one

```

```

  by (blast intro: continuous_discrete_range_constant [OF cs cont])

```

qed

lemma *winding_number_eq*:

```

[[path  $\gamma$ ; pathfinish  $\gamma = \text{pathstart } \gamma$ ;  $w \in S$ ;  $z \in S$ ; connected  $S$ ;  $S \cap \text{path\_image } \gamma = \{\}$ ]

```

```

 $\implies \text{winding\_number } \gamma w = \text{winding\_number } \gamma z$ 

```

```

  using winding_number_constant by (metis constant_on_def)

```

lemma *open_winding_number_levelsets*:

```

assumes  $\gamma$ : path  $\gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 

```

```

shows open  $\{z. z \notin \text{path\_image } \gamma \wedge \text{winding\_number } \gamma z = k\}$ 

```

proof (clarsimp simp: open_dist)

```

fix  $z$  assume  $z$ :  $z \notin \text{path\_image } \gamma$  and  $k$ :  $k = \text{winding\_number } \gamma z$ 

```

```

have open  $(- \text{path\_image } \gamma)$ 

```

```

  by (simp add: closed_path_image  $\gamma$  open_Compl)

```

```

then obtain  $e$  where  $e > 0$  ball  $z e \subseteq - \text{path\_image } \gamma$ 

```

```

  using open_contains_ball [of  $- \text{path\_image } \gamma$ ]  $z$  by blast

```

```

then show  $\exists e > 0. \forall y. \text{dist } y \ z < e \longrightarrow y \notin \text{path\_image } \gamma \wedge \text{winding\_number}$ 
 $\gamma \ y = \text{winding\_number } \gamma \ z$ 
using  $\langle e > 0 \rangle$  by (force simp: norm_minus_commute dist_norm intro: wind-
ing_number_eq [OF assms, where  $S = \text{ball } z \ e$ ])
qed

```

3.5 Winding number is zero "outside" a curve

proposition *winding_number_zero_in_outside:*

assumes γ : path γ **and** loop: pathfinish $\gamma = \text{pathstart } \gamma$ **and** z : $z \in \text{outside}$
(path_image γ)

shows $\text{winding_number } \gamma \ z = 0$

proof –

obtain $B::\text{real}$ **where** $0 < B$ **and** B : path_image $\gamma \subseteq \text{ball } 0 \ B$

using bounded_subset_ballD [OF bounded_path_image [OF γ]] **by** auto

obtain $w::\text{complex}$ **where** $w \notin \text{ball } 0 \ (B + 1)$

by (metis abs_of_nonneg le_less less_irrefl mem_ball_0 norm_of_real)

have $-\text{ball } 0 \ (B + 1) \subseteq \text{outside} \ (\text{path_image } \gamma)$

using B subset_ball **by** (intro outside_subset_convex) auto

then have wout: $w \in \text{outside} \ (\text{path_image } \gamma)$

using w **by** blast

moreover have $\text{winding_number } \gamma \ \text{constant_on } \text{outside} \ (\text{path_image } \gamma)$

using winding_number_constant [OF γ loop, of outside(path_image γ)] con-
nected_outside

by (metis DIM_complex bounded_path_image_dual_order_refl γ outside_no_overlap)

ultimately have $\text{winding_number } \gamma \ z = \text{winding_number } \gamma \ w$

by (metis (no_types, opaque_lifting) constant_on_def z)

also have $\dots = 0$

proof –

have wnot: $w \notin \text{path_image } \gamma$ **using** wout **by** (simp add: outside_def)

{ **fix** $e::\text{real}$ **assume** $0 < e$

obtain p **where** p : polynomial_function p pathstart $p = \text{pathstart } \gamma$ pathfinish
 $p = \text{pathfinish } \gamma$

and $pg1$: $(\bigwedge t. [0 \leq t; t \leq 1] \Longrightarrow \text{cmod } (p \ t - \gamma \ t) < 1)$

and $pg2$: $(\bigwedge t. [0 \leq t; t \leq 1] \Longrightarrow \text{cmod } (p \ t - \gamma \ t) < e)$

using path_approx_polynomial_function [OF γ , of min 1 e] $\langle e > 0 \rangle$

by (metis atLeastAtMost_iff min_less_iff_conj zero_less_one)

have $\exists p. \text{valid_path } p \wedge w \notin \text{path_image } p \wedge$

$\text{pathstart } p = \text{pathstart } \gamma \wedge \text{pathfinish } p = \text{pathfinish } \gamma \wedge$

$(\forall t \in \{0..1\}. \text{cmod } (\gamma \ t - p \ t) < e) \wedge \text{contour_integral } p \ (\lambda wa. 1$

$/ (wa - w)) = 0$

proof (intro exI conjI)

have $\bigwedge x. [0 \leq x; x \leq 1] \Longrightarrow \text{cmod } (p \ x) < B + 1$

using B unfolding image_subset_iff path_image_def

by (meson add_strict_mono atLeastAtMost_iff le_less_trans mem_ball_0
norm_triangle_sub $pg1$)

then have pip : path_image $p \subseteq \text{ball } 0 \ (B + 1)$

by (auto simp add: path_image_def dist_norm ball_def)

then show $w \notin \text{path_image } p$ **using** w **by** blast

```

show vap: valid_path p
  by (simp add: p(1) valid_path_polynomial_function)
show  $\forall t \in \{0..1\}. \text{cmod } (\gamma t - p t) < e$ 
  by (metis atLeastAtMost_iff norm_minus_commute pge)
show contour_integral p ( $\lambda w a. 1 / (w a - w)$ ) = 0
proof (rule contour_integral_unique [OF Cauchy_theorem_convex_simple
[OF _convex_ball [of 0 B+1]]])
  have  $\bigwedge z. \text{cmod } z < B + 1 \implies z \neq w$ 
    using mem_ball_0 w by blast
  then show ( $\lambda z. 1 / (z - w)$ ) holomorphic_on ball 0 (B + 1)
    by (intro holomorphic_intros; simp add: dist_norm)
  qed (use p vap pip loop in auto)
qed (use p in auto)
}
then show ?thesis
by (auto intro: winding_number_unique [OF  $\gamma$ ] simp add: winding_number_prop_def
wnot)
qed
finally show ?thesis .
qed

```

```

corollary winding_number_zero_const: a  $\neq$  z  $\implies$  winding_number ( $\lambda t. a$ ) z =
0
by (rule winding_number_zero_in_outside)
(auto simp: pathfinish_def pathstart_def path_polynomial_function)

```

```

corollary winding_number_zero_outside:
 $\llbracket \text{path } \gamma; \text{convex } s; \text{pathfinish } \gamma = \text{pathstart } \gamma; z \notin s; \text{path\_image } \gamma \subseteq s \rrbracket \implies$ 
winding_number  $\gamma$  z = 0
by (meson convex_in_outside outside_mono subsetCE winding_number_zero_in_outside)

```

```

lemma winding_number_zero_at_infinity:
assumes  $\gamma: \text{path } \gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
shows  $\exists B. \forall z. B \leq \text{norm } z \longrightarrow \text{winding\_number } \gamma z = 0$ 
proof –
obtain B::real where  $0 < B$  and B: path_image  $\gamma \subseteq \text{ball } 0 B$ 
using bounded_subset_ballD [OF bounded_path_image [OF  $\gamma$ ]] by auto
have winding_number  $\gamma$  z = 0 if  $B + 1 \leq \text{cmod } z$  for z
proof (rule winding_number_zero_outside [OF  $\gamma$  convex_cball loop])
show  $z \notin \text{cball } 0 B$ 
using that by auto
show path_image  $\gamma \subseteq \text{cball } 0 B$ 
using B order.trans by blast
qed
then show ?thesis
by metis
qed

```

```

lemma bounded_winding_number_nz:

```

```

assumes path g pathfinish g = pathstart g
shows bounded {z. winding_number g z ≠ 0}
proof –
obtain B where  $\bigwedge x. \text{norm } x \geq B \implies \text{winding\_number } g \ x = 0$ 
using winding_number_zero_at_infinity[OF assms] by auto
thus ?thesis
unfolding bounded_iff by (intro exI[of _ B + 1]) force
qed

```

```

lemma winding_number_zero_point:
   $\llbracket \text{path } \gamma; \text{convex } S; \text{pathfinish } \gamma = \text{pathstart } \gamma; \text{open } S; \text{path\_image } \gamma \subseteq S \rrbracket$ 
   $\implies \exists z. z \in S \wedge \text{winding\_number } \gamma \ z = 0$ 
using outside_compact_in_open [of path_image  $\gamma$  S] path_image_nonempty
winding_number_zero_in_outside
by (fastforce simp add: compact_path_image)

```

If a path winds round a set, it winds rounds its inside.

```

lemma winding_number_around_inside:
assumes  $\gamma$ : path  $\gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
and cls: closed S and cos: connected S and S_disj:  $S \cap \text{path\_image } \gamma = \{\}$ 
and z:  $z \in S$  and wn_nz: winding_number  $\gamma \ z \neq 0$  and w:  $w \in S \cup \text{inside } S$ 
shows winding_number  $\gamma \ w = \text{winding\_number } \gamma \ z$ 
proof –
have ssb:  $S \subseteq \text{inside}(\text{path\_image } \gamma)$ 
proof
fix x :: complex
assume  $x \in S$ 
hence  $x \notin \text{path\_image } \gamma$ 
by (meson disjoint_iff_not_equal S_disj)
thus  $x \in \text{inside}(\text{path\_image } \gamma)$ 
by (metis Compl_iff S_disj UnE  $\gamma \ \langle x \in S \rangle$  cos inside_outside loop winding_number_eq winding_number_zero_in_outside wn_nz z)
qed
show ?thesis
proof (rule winding_number_eq [OF  $\gamma$  loop w])
show  $z \in S \cup \text{inside } S$ 
using z by blast
show connected (S  $\cup$  inside S)
by (simp add: cls connected_with_inside cos)
show  $(S \cup \text{inside } S) \cap \text{path\_image } \gamma = \{\}$ 
unfolding disjoint_iff Un_iff
by (metis ComplD UnI1  $\gamma$  cls compact_path_image connected_path_image inside_Un_outside inside_inside_compact_connected ssb subsetD)
qed
qed

```

Bounding a WN by 1/2 for a path and point in opposite halfspaces.

```

lemma winding_number_subpath_continuous:

```



```

assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$ 
shows continuous_on  $\{0..1\}$   $(\lambda x. \text{winding\_number}(\text{subpath } 0 \ x \ \gamma) \ z)$ 
proof (rule continuous_on_eq)
let  $?f = \lambda x. \text{integral } \{0..x\} (\lambda t. \text{vector\_derivative } \gamma \ (\text{at } t) / (\gamma \ t - z))$ 
show continuous_on  $\{0..1\}$   $(\lambda x. 1 / (2 * \text{pi} * i) * ?f \ x)$ 
proof (intro indefinite_integral_continuous_1 winding_number_exp_integral
continuous_intros)
show  $\gamma$  piecewise_C1_differentiable_on  $\{0..1\}$ 
using  $\gamma$  valid_path_def by blast
qed (use path_image_def z in auto)
show  $1 / (2 * \text{pi} * i) * ?f \ x = \text{winding\_number} \ (\text{subpath } 0 \ x \ \gamma) \ z$ 
if  $x \in \{0..1\}$  for  $x$ 
proof -
have  $1 / (2 * \text{pi} * i) * ?f \ x = 1 / (2 * \text{pi} * i) * \text{contour\_integral} \ (\text{subpath } 0 \ x \ \gamma)$ 
 $(\lambda w. 1 / (w - z))$ 
using assms x
by (simp add: contour_integral_subcontour_integral [OF contour_integrable_inversediff])
also have  $\dots = \text{winding\_number} \ (\text{subpath } 0 \ x \ \gamma) \ z$ 
proof (subst winding_number_valid_path)
show  $z \notin \text{path\_image} \ (\text{subpath } 0 \ x \ \gamma)$ 
using assms x atLeastAtMost_iff path_image_subpath_subset by force
qed (use assms x valid_path_subpath in <force+>)
finally show ?thesis .
qed
qed

```

lemma *winding_number_ivt_pos*:

```

assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and  $0 \leq w \leq \text{Re}(\text{winding\_number}$ 
 $\gamma \ z)$ 
shows  $\exists t \in \{0..1\}. \text{Re}(\text{winding\_number}(\text{subpath } 0 \ t \ \gamma) \ z) = w$ 
proof -
have continuous_on  $\{0..1\}$   $(\lambda x. \text{winding\_number} \ (\text{subpath } 0 \ x \ \gamma) \ z)$ 
using  $\gamma$  winding_number_subpath_continuous z by blast
moreover have  $\text{Re} \ (\text{winding\_number} \ (\text{subpath } 0 \ 0 \ \gamma) \ z) \leq w \leq \text{Re} \ (\text{winding\_number}$ 
 $(\text{subpath } 0 \ 1 \ \gamma) \ z)$ 
using assms by (auto simp: path_image_def image_def)
ultimately show ?thesis
using ivt_increasing_component_on_1 [of 0 1, where ?k = 1] by force
qed

```

lemma *winding_number_ivt_neg*:

```

assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and  $\text{Re}(\text{winding\_number}$ 
 $\gamma \ z) \leq w \leq 0$ 
shows  $\exists t \in \{0..1\}. \text{Re}(\text{winding\_number}(\text{subpath } 0 \ t \ \gamma) \ z) = w$ 
proof -
have continuous_on  $\{0..1\}$   $(\lambda x. \text{winding\_number} \ (\text{subpath } 0 \ x \ \gamma) \ z)$ 
using  $\gamma$  winding_number_subpath_continuous z by blast
moreover have  $\text{Re} \ (\text{winding\_number} \ (\text{subpath } 0 \ 0 \ \gamma) \ z) \geq w \geq \text{Re} \ (\text{winding\_number}$ 
 $(\text{subpath } 0 \ 1 \ \gamma) \ z)$ 

```

```

    using assms by (auto simp: path_image_def image_def)
  ultimately show ?thesis
    using ivt_decreasing_component_on_1 [of 0 1, where ?k = 1] by force
qed

```

lemma *winding_number_ivt_abs*:

```

  assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and  $0 \leq w \leq |\text{Re}(\text{winding\_number } \gamma z)|$ 
  shows  $\exists t \in \{0..1\}. |\text{Re}(\text{winding\_number } (\text{subpath } 0 t \gamma) z)| = w$ 
  using assms winding_number_ivt_pos [of  $\gamma z w$ ] winding_number_ivt_neg [of  $\gamma z -w$ ]
  by force

```

lemma *winding_number_lt_half_lemma*:

```

  assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin \text{path\_image } \gamma$  and  $az$ :  $a \cdot z \leq b$  and  $pag$ :
  path_image  $\gamma \subseteq \{w. a \cdot w > b\}$ 

```

```

  shows  $\text{Re}(\text{winding\_number } \gamma z) < 1/2$ 

```

proof –

```

  { assume  $\text{Re}(\text{winding\_number } \gamma z) \geq 1/2$ 
    then obtain  $t::\text{real}$  where  $t: 0 \leq t \leq 1$  and  $sub12$ :  $\text{Re}(\text{winding\_number } (\text{subpath } 0 t \gamma) z) = 1/2$ 
      using winding_number_ivt_pos [OF  $\gamma z$ , of  $1/2$ ] by auto
    have  $gt$ :  $\gamma t - z = -(\text{of\_real } (\exp(- (2 * \pi * \text{Im}(\text{winding\_number } (\text{subpath } 0 t \gamma) z)))) * (\gamma 0 - z))$ 
      using winding_number_exp_2pi [of subpath 0 t  $\gamma z$ ]
      apply (simp add: t  $\gamma$  valid_path_imp_path)
      using closed_segment_eq_real_ivl path_image_def t z by (fastforce simp: path_image_subpath Euler sub12)
    have  $b < a \cdot \gamma 0$ 
      proof –
        have  $\gamma 0 \in \{c. b < a \cdot c\}$ 
          by (metis (no_types) pag atLeastAtMost_iff image_subset_iff order_refl path_image_def zero_le_one)
        thus ?thesis
          by blast
      qed
    moreover have  $b < a \cdot \gamma t$ 
      by (metis atLeastAtMost_iff image_eqI mem_Collect_eq pag path_image_def subset_iff t)
    ultimately have  $0 < a \cdot (\gamma 0 - z) \wedge 0 < a \cdot (\gamma t - z)$  using  $az$ 
      by (simp add: inner_diff_right)+
    then have False
      by (simp add: gt inner_mult_right mult_less_0_iff)
  }
  then show ?thesis by force
qed

```

lemma *winding_number_lt_half*:

```

  assumes valid_path  $\gamma$   $a \cdot z \leq b$  path_image  $\gamma \subseteq \{w. a \cdot w > b\}$ 

```

```

  shows  $|Re (winding\_number \ \gamma \ z)| < 1/2$ 
proof -
  have  $z \notin path\_image \ \gamma$  using assms by auto
  with assms have  $Re (winding\_number \ \gamma \ z) < 0 \implies - Re (winding\_number \ \gamma \ z) < 1/2$ 
  by (metis complex_inner_1_right winding_number_lt_half_lemma [OF valid_path_imp_reverse, of  $\gamma \ z \ a \ b$ ])
  winding_number_reversepath valid_path_imp_path inner_minus_left path_image_reversepath)
  with assms show ?thesis
  using  $\langle z \notin path\_image \ \gamma \rangle$  winding_number_lt_half_lemma by fastforce
qed

```

lemma *winding_number_le_half*:

```

  assumes  $\gamma$ : valid_path  $\gamma$  and  $z$ :  $z \notin path\_image \ \gamma$ 
  and anz:  $a \neq 0$  and azb:  $a \cdot z \leq b$  and pag:  $path\_image \ \gamma \subseteq \{w. a \cdot w \geq b\}$ 
  shows  $|Re (winding\_number \ \gamma \ z)| \leq 1/2$ 
proof -
  { assume wnz_12:  $|Re (winding\_number \ \gamma \ z)| > 1/2$ 
  have isCont (winding_number  $\gamma$ )  $z$ 
  by (metis continuous_at_winding_number_valid_path_imp_path  $\gamma \ z$ )
  then obtain  $d$  where  $d > 0$  and  $d$ :  $\bigwedge x'. dist \ x' \ z < d \implies dist (winding\_number \ \gamma \ x') (winding\_number \ \gamma \ z) < |Re(winding\_number \ \gamma \ z)| - 1/2$ 
  using continuous_at_eps_delta_wnz_12_diff_gt_0_iff_gt by blast
  define  $z'$  where  $z' = z - (d / (2 * cmod \ a)) *_R \ a$ 
  have  $a \cdot z * 6 \leq d * cmod \ a + b * 6$ 
  by (metis  $\langle 0 < d \rangle$  add_increasing azb less_eq_real_def mult_nonneg_nonneg mult_right_mono norm_ge_zero norm_numeral)
  with anz have  $*$ :  $a \cdot z' \leq b - d / 3 * cmod \ a$ 
  unfolding  $z'_def$  inner_mult_right' divide_inverse
  by (simp add: field_split_simps algebra_simps dot_square_norm power2_eq_square)
  have cmod (winding_number  $\gamma \ z' - winding\_number \ \gamma \ z$ )  $< |Re (winding\_number \ \gamma \ z)| - 1/2$ 
  using  $d$  [of  $z'$ ] anz  $\langle d > 0 \rangle$  by (simp add: dist_norm  $z'_def$ )
  then have  $1/2 < |Re (winding\_number \ \gamma \ z)| - cmod (winding\_number \ \gamma \ z' - winding\_number \ \gamma \ z)$ 
  by simp
  then have  $1/2 < |Re (winding\_number \ \gamma \ z)| - |Re (winding\_number \ \gamma \ z') - Re (winding\_number \ \gamma \ z)|$ 
  using abs_Re_le_cmod [of  $winding\_number \ \gamma \ z' - winding\_number \ \gamma \ z$ ]
  by simp
  then have wnz_12':  $|Re (winding\_number \ \gamma \ z')| > 1/2$ 
  by linarith
  moreover have  $|Re (winding\_number \ \gamma \ z')| < 1/2$ 
  proof (rule winding_number_lt_half [OF  $\gamma \ *$ ])
  show  $path\_image \ \gamma \subseteq \{w. b - d / 3 * cmod \ a < a \cdot w\}$ 
  using azb  $\langle d > 0 \rangle$  pag by (auto simp: add_strict_increasing anz field_split_simps dest!: subsetD)
  qed
  ultimately have False

```

```

    by simp
  }
  then show ?thesis by force
qed

```

```

lemma winding_number_lt_half_linepath:
  assumes  $z \notin \text{closed\_segment } a \ b$  shows  $|\text{Re}(\text{winding\_number}(\text{linepath } a \ b) \ z)| < 1/2$ 
proof -
  obtain  $u \ v$  where  $u \cdot z \leq v$  and  $uv: \bigwedge x. x \in \text{closed\_segment } a \ b \implies \text{inner } u \ x > v$ 
  using separating_hyperplane_closed_point assms closed_segment convex_closed_segment less_eq_real_def by metis
  moreover have  $\text{path\_image}(\text{linepath } a \ b) \subseteq \{w. v < u \cdot w\}$ 
  using in_segment(1) uv by auto
  ultimately show ?thesis
  using winding_number_lt_half by auto
qed

```

Positivity of WN for a linepath.

```

lemma winding_number_linepath_pos_lt:
  assumes  $0 < \text{Im}((b - a) * \text{cnj}(b - z))$ 
  shows  $0 < \text{Re}(\text{winding\_number}(\text{linepath } a \ b) \ z)$ 
proof -
  have  $z: z \notin \text{path\_image}(\text{linepath } a \ b)$ 
  using assms
  by (simp add: closed_segment_def) (force simp: algebra_simps)
  show ?thesis
  by (intro winding_number_pos_lt [OF valid_path_linepath z assms]) (simp add: linepath_def algebra_simps)
qed

```

```

lemma winding_number_linepath_neg_lt:
  assumes  $0 < \text{Im}((a - b) * \text{cnj}(a - z))$ 
  shows  $\text{Re}(\text{winding\_number}(\text{linepath } a \ b) \ z) < 0$ 
proof -
  have  $z: z \notin \text{path\_image}(\text{linepath } a \ b)$ 
  using assms
  by (simp add: closed_segment_def) (force simp: algebra_simps)
  have  $\text{Re}(\text{winding\_number}(\text{linepath } a \ b) \ z) = -\text{Re}(\text{winding\_number}(\text{reversepath}(\text{linepath } a \ b)) \ z)$ 
  by (subst winding_number_reversepath) (use z in auto)
  also have  $\dots = -\text{Re}(\text{winding\_number}(\text{linepath } b \ a) \ z)$ 
  by simp
  finally have  $\text{Re}(\text{winding\_number}(\text{linepath } a \ b) \ z) = -\text{Re}(\text{winding\_number}(\text{linepath } b \ a) \ z)$ 
  moreover have  $0 < \text{Re}(\text{winding\_number}(\text{linepath } b \ a) \ z)$ 
  using winding_number_linepath_pos_lt[OF assms] .
  ultimately show ?thesis by auto

```

qed

3.6 More winding number properties

including the fact that it's $+1$ inside a simple closed curve.

lemma *winding_number_homotopic_paths*:

assumes *homotopic_paths* $(-\{z\})$ *g h*

shows *winding_number* *g z* = *winding_number* *h z*

proof –

have *path* *g path h* **using** *homotopic_paths_imp_path* [*OF assms*] **by** *auto*

moreover have *pag*: $z \notin \text{path_image } g$ **and** *pah*: $z \notin \text{path_image } h$

using *homotopic_paths_imp_subset* [*OF assms*] **by** *auto*

ultimately obtain *d e* **where** $d > 0$ $e > 0$

and *d*: $\bigwedge p. \llbracket \text{path } p; \text{pathstart } p = \text{pathstart } g; \text{pathfinish } p = \text{pathfinish } g; \forall t \in \{0..1\}. \text{norm } (p \ t - g \ t) < d \rrbracket$

$\implies \text{homotopic_paths } (-\{z\})$ *g p*

and *e*: $\bigwedge q. \llbracket \text{path } q; \text{pathstart } q = \text{pathstart } h; \text{pathfinish } q = \text{pathfinish } h; \forall t \in \{0..1\}. \text{norm } (q \ t - h \ t) < e \rrbracket$

$\implies \text{homotopic_paths } (-\{z\})$ *h q*

using *homotopic_nearby_paths* [*of g* $-\{z\}$] *homotopic_nearby_paths* [*of h* $-\{z\}$] **by** *force*

obtain *p* **where** *p*:

valid_path *p z* $\notin \text{path_image } p$

pathstart *p* = *pathstart* *g pathfinish* *p* = *pathfinish* *g*

and *gp_less*: $\forall t \in \{0..1\}. \text{cmod } (g \ t - p \ t) < d$

and *pap*: *contour_integral* *p* $(\lambda w. 1 / (w - z)) = \text{complex_of_real } (2 * \pi i)$

* i * *winding_number* *g z*

using *winding_number* [*OF* $\langle \text{path } g \rangle \text{ pag } \langle 0 < d \rangle$] **unfolding** *winding_number_prop_def* **by** *blast*

obtain *q* **where** *q*:

valid_path *q z* $\notin \text{path_image } q$

pathstart *q* = *pathstart* *h pathfinish* *q* = *pathfinish* *h*

and *hq_less*: $\forall t \in \{0..1\}. \text{cmod } (h \ t - q \ t) < e$

and *paq*: *contour_integral* *q* $(\lambda w. 1 / (w - z)) = \text{complex_of_real } (2 * \pi i)$

* i * *winding_number* *h z*

using *winding_number* [*OF* $\langle \text{path } h \rangle \text{ pah } \langle 0 < e \rangle$] **unfolding** *winding_number_prop_def* **by** *blast*

have *homotopic_paths* $(-\{z\})$ *g p*

by (*simp add*: *d p valid_path_imp_path norm_minus_commute gp_less*)

moreover have *homotopic_paths* $(-\{z\})$ *h q*

by (*simp add*: *e q valid_path_imp_path norm_minus_commute hq_less*)

ultimately have *homotopic_paths* $(-\{z\})$ *p q*

by (*blast intro*: *homotopic_paths_trans homotopic_paths_sym assms*)

then have *contour_integral* *p* $(\lambda w. 1 / (w - z)) = \text{contour_integral } q$ $(\lambda w. 1 / (w - z))$

by (*rule Cauchy_theorem_homotopic_paths*) (*auto intro*!: *holomorphic_intros simp*: *p q*)

then show *?thesis*

by (*simp add*: *pap paq*)

qed

lemma winding_number_homotopic_loops:

assumes homotopic_loops $(-\{z\})$ g h
 shows winding_number g z = winding_number h z

proof -

have path g path h using homotopic_loops_imp_path [OF assms] by auto

moreover have pag: $z \notin \text{path_image } g$ and pah: $z \notin \text{path_image } h$

using homotopic_loops_imp_subset [OF assms] by auto

moreover have gloop: pathfinish g = pathstart g and hloop: pathfinish h = pathstart h

using homotopic_loops_imp_loop [OF assms] by auto

ultimately obtain d e where $d > 0$ $e > 0$

and d : $\bigwedge p. [\text{path } p; \text{pathfinish } p = \text{pathstart } p; \forall t \in \{0..1\}. \text{norm } (p \ t - g \ t) < d]$

$\implies \text{homotopic_loops } (-\{z\})$ g p

and e : $\bigwedge q. [\text{path } q; \text{pathfinish } q = \text{pathstart } q; \forall t \in \{0..1\}. \text{norm } (q \ t - h \ t) < e]$

$\implies \text{homotopic_loops } (-\{z\})$ h q

using homotopic_nearby_loops [of g $-\{z\}$] homotopic_nearby_loops [of h $-\{z\}$] by force

obtain p where p :

valid_path p $z \notin \text{path_image } p$

pathstart p = pathstart g pathfinish p = pathfinish g

and gp_less: $\forall t \in \{0..1\}. \text{cmod } (g \ t - p \ t) < d$

and pap: contour_integral p $(\lambda w. 1 / (w - z)) = \text{complex_of_real } (2 * \pi i)$

* i * winding_number g z

using winding_number [OF $\langle \text{path } g \rangle$ pag $\langle 0 < d \rangle$] unfolding winding_number_prop_def by blast

obtain q where q :

valid_path q $z \notin \text{path_image } q$

pathstart q = pathstart h pathfinish q = pathfinish h

and hq_less: $\forall t \in \{0..1\}. \text{cmod } (h \ t - q \ t) < e$

and paq: contour_integral q $(\lambda w. 1 / (w - z)) = \text{complex_of_real } (2 * \pi i)$

* i * winding_number h z

using winding_number [OF $\langle \text{path } h \rangle$ pah $\langle 0 < e \rangle$] unfolding winding_number_prop_def by blast

have gp: homotopic_loops $(-\{z\})$ g p

by (simp add: gloop d gp_less norm_minus_commute p valid_path_imp_path)

have hq: homotopic_loops $(-\{z\})$ h q

by (simp add: e hloop hq_less norm_minus_commute q valid_path_imp_path)

have contour_integral p $(\lambda w. 1 / (w - z)) = \text{contour_integral } q$ $(\lambda w. 1 / (w - z))$

proof (rule Cauchy_theorem_homotopic_loops)

show homotopic_loops $(-\{z\})$ p q

by (blast intro: homotopic_loops_trans homotopic_loops_sym gp hq assms)

qed (auto intro!: holomorphic_intros simp: p q)

then show ?thesis

by (simp add: pap paq)

qed

lemma *winding_number_paths_linear_eq*:
 $\llbracket \text{path } g; \text{ path } h; \text{ pathstart } h = \text{pathstart } g; \text{ pathfinish } h = \text{pathfinish } g; \wedge t. t \in \{0..1\} \implies z \notin \text{closed_segment } (g \ t) \ (h \ t) \rrbracket$
 $\implies \text{winding_number } h \ z = \text{winding_number } g \ z$
by (*blast intro: sym homotopic_paths_linear winding_number_homotopic_paths*)

lemma *winding_number_loops_linear_eq*:
 $\llbracket \text{path } g; \text{ path } h; \text{ pathfinish } g = \text{pathstart } g; \text{ pathfinish } h = \text{pathstart } h; \wedge t. t \in \{0..1\} \implies z \notin \text{closed_segment } (g \ t) \ (h \ t) \rrbracket$
 $\implies \text{winding_number } h \ z = \text{winding_number } g \ z$
by (*blast intro: sym homotopic_loops_linear winding_number_homotopic_loops*)

lemma *winding_number_nearby_paths_eq*:
 $\llbracket \text{path } g; \text{ path } h; \text{ pathstart } h = \text{pathstart } g; \text{ pathfinish } h = \text{pathfinish } g; \wedge t. t \in \{0..1\} \implies \text{norm}(h \ t - g \ t) < \text{norm}(g \ t - z) \rrbracket$
 $\implies \text{winding_number } h \ z = \text{winding_number } g \ z$
by (*metis segment_bound(2) norm_minus_commute not_le winding_number_paths_linear_eq*)

lemma *winding_number_nearby_loops_eq*:
 $\llbracket \text{path } g; \text{ path } h; \text{ pathfinish } g = \text{pathstart } g; \text{ pathfinish } h = \text{pathstart } h; \wedge t. t \in \{0..1\} \implies \text{norm}(h \ t - g \ t) < \text{norm}(g \ t - z) \rrbracket$
 $\implies \text{winding_number } h \ z = \text{winding_number } g \ z$
by (*metis segment_bound(2) norm_minus_commute not_le winding_number_loops_linear_eq*)

lemma *winding_number_subpath_combine*:
assumes *path g and notin: z ∉ path_image g and u ∈ {0..1} v ∈ {0..1} w ∈ {0..1}*
shows $\text{winding_number } (\text{subpath } u \ v \ g) \ z + \text{winding_number } (\text{subpath } v \ w \ g) \ z =$
 $\text{winding_number } (\text{subpath } u \ w \ g) \ z$ (**is** *?lhs = ?rhs*)

proof –
have *?lhs = winding_number (subpath u v g) + winding_number (subpath v w g) z*
using *assms*
by (*metis (no_types, lifting) path_image_subpath_subset path_subpath pathfinish_subpath pathstart_subpath subsetD winding_number_join*)
also have *... = ?rhs*
by (*meson assms homotopic_join_subpaths subset_Compl_singleton winding_number_homotopic_paths*)
finally show *?thesis* .
qed

Winding numbers of circular contours

proposition *winding_number_part_circlepath_pos_less*:
assumes *s < t and no: norm(w - z) < r*
shows $0 < \text{Re } (\text{winding_number}(\text{part_circlepath } z \ r \ s \ t) \ w)$
proof –
have $0 < r$ **by** (*meson no norm_not_less_zero not_le order.strict_trans2*)

```

note valid_path_part_circlepath
moreover have  $w \notin \text{path\_image } (\text{part\_circlepath } z \ r \ s \ t)$ 
  using assms by (auto simp: path_image_def image_def part_circlepath_def
norm_mult linepath_def)
moreover have  $0 < r * (t - s) * (r - \text{cmod } (w - z))$ 
  using assms by (metis <0 < r> diff_gt_0_iff_gt mult_pos_pos)
ultimately show ?thesis
  apply (rule winding_number_pos_lt [where e = r*(t - s)*(r - norm(w - z))])
  apply (simp add: vector_derivative_part_circlepath right_diff_distrib [symmetric]
mult_ac mult_le_cancel_left_pos assms <0 < r>)
  using Re_Im_le_cmod [of w-z linepath s t x for x]
  by (simp add: exp_Euler cos_of_real sin_of_real part_circlepath_def algebra_simps cos_squared_eq [unfolded power2_eq_square])
qed

```

```

lemma winding_number_circlepath_centre:  $0 < r \implies \text{winding\_number } (\text{circlepath } z \ r) \ z = 1$ 
  apply (rule winding_number_unique_loop)
  apply (simp_all add: sphere_def valid_path_imp_path)
  apply (rule_tac x=circlepath z r in exI)
  apply (simp add: sphere_def contour_integral_circlepath)
done

```

```

proposition winding_number_circlepath:
  assumes  $\text{norm}(w - z) < r$  shows  $\text{winding\_number}(\text{circlepath } z \ r) \ w = 1$ 
proof (cases w = z)
  case True then show ?thesis
    using assms winding_number_circlepath_centre by auto
next
  case False
    have [simp]:  $r > 0$ 
      using assms le_less_trans norm_ge_zero by blast
    define  $r'$  where  $r' = \text{norm}(w - z)$ 
    have  $r' < r$ 
      by (simp add: assms r'_def)
    have disjo: cball z r' ∩ sphere z r = {}
      using  $\langle r' < r \rangle$  by (force simp: cball_def sphere_def)
    have  $\text{winding\_number}(\text{circlepath } z \ r) \ w = \text{winding\_number}(\text{circlepath } z \ r) \ z$ 
    proof (rule winding_number_around_inside [where S = cball z r'])
      show  $\text{winding\_number } (\text{circlepath } z \ r) \ z \neq 0$ 
        by (simp add: winding_number_circlepath_centre)
      show  $\text{cball } z \ r' \cap \text{path\_image } (\text{circlepath } z \ r) = \{\}$ 
        by (simp add: disjo less_eq_real_def)
    qed (auto simp: r'_def dist_norm norm_minus_commute)
    also have  $\dots = 1$ 
      by (simp add: winding_number_circlepath_centre)
    finally show ?thesis .
qed

```



```

lemma no_bounded_connected_component_imp_winding_number_zero:
  assumes g: path g path_image g  $\subseteq$  S pathfinish g = pathstart g z  $\notin$  S
    and nb:  $\bigwedge z$ . bounded (connected_component_set (- S) z)  $\implies$  z  $\in$  S
  shows winding_number g z = 0
proof -
  have z  $\in$  outside (path_image g)
    by (metis nb [of z]  $\langle$ path_image g  $\subseteq$  S $\rangle$   $\langle$ z  $\notin$  S $\rangle$  subsetD mem_Collect_eq
    outside outside_mono)
  then show ?thesis
    by (simp add: g winding_number_zero_in_outside)
qed

lemma no_bounded_path_component_imp_winding_number_zero:
  assumes g: path g path_image g  $\subseteq$  S pathfinish g = pathstart g z  $\notin$  S
    and nb:  $\bigwedge z$ . bounded (path_component_set (- S) z)  $\implies$  z  $\in$  S
  shows winding_number g z = 0
by (simp add: bounded_subset nb path_component_subset_connected_component

    no_bounded_connected_component_imp_winding_number_zero [OF
g])

```

3.7 Winding number for a triangle

```

lemma wn_triangle1:
  assumes 0  $\in$  interior(convex hull {a,b,c})
  shows  $\neg$  (Im(a/b)  $\leq$  0  $\wedge$  0  $\leq$  Im(b/c))
proof -
  { assume 0: Im(a/b)  $\leq$  0 0  $\leq$  Im(b/c)
    have 0  $\notin$  interior (convex hull {a,b,c})
    proof (cases a=0  $\vee$  b=0  $\vee$  c=0)
      case True then show ?thesis
        by (auto simp: not_in_interior_convex_hull_3)
    next
      case False
      then have b  $\neq$  0 by blast
      { fix x y::complex and u::real
        assume eq_f': Im x * Re b  $\leq$  Im b * Re x Im y * Re b  $\leq$  Im b * Re y 0  $\leq$ 
u u  $\leq$  1
        then have ((1 - u) * Im x) * Re b  $\leq$  Im b * ((1 - u) * Re x)
          by (simp add: mult_left_mono mult.assoc mult.left_commute [of Im b])
        then have ((1 - u) * Im x + u * Im y) * Re b  $\leq$  Im b * ((1 - u) * Re x
+ u * Re y)
          using eq_f' ordered_comm_semiring_class.comm_mult_left_mono
          by (fastforce simp add: algebra_simps)
        }
      then have convex {z. Im z * Re b  $\leq$  Im b * Re z}
        by (auto simp: algebra_simps convex_alt)
      with False 0 have convex hull {a,b,c}  $\leq$  {z. Im z * Re b  $\leq$  Im b * Re z}

```

```

    by (simp add: subset_hull mult.commute Complex.Im_divide divide_simps
        complex_neq_0 [symmetric])
    moreover have  $0 \notin \text{interior}(\{z. \text{Im } z * \text{Re } b \leq \text{Im } b * \text{Re } z\})$ 
    proof
      assume  $0 \in \text{interior } \{z. \text{Im } z * \text{Re } b \leq \text{Im } b * \text{Re } z\}$ 
      then obtain  $e$  where  $e > 0$  and  $e: \text{ball } 0 \ e \subseteq \{z. \text{Im } z * \text{Re } b \leq \text{Im } b * \text{Re } z\}$ 
      by (meson mem_interior)
      define  $z$  where  $z \equiv - \text{sgn } (\text{Im } b) * (e/3) + \text{sgn } (\text{Re } b) * (e/3) * i$ 
      have  $\text{cmod } z = \text{cmod } (e/3) * \text{cmod } (- \text{sgn } (\text{Im } b) + \text{sgn } (\text{Re } b) * i)$ 
      unfolding  $z\_def$  norm_mult [symmetric] by (simp add: algebra_simps)
      also have  $\dots < e$ 
      using  $\langle e > 0 \rangle$  by (auto simp: norm_mult intro: le_less_trans [OF
        norm_triangle_ineq4])
      finally have  $z \in \text{ball } 0 \ e$ 
      using  $\langle e > 0 \rangle$  by simp
      then have  $\text{Im } z * \text{Re } b \leq \text{Im } b * \text{Re } z$ 
      using  $e$  by blast
      then have  $b: 0 < \text{Re } b \ 0 < \text{Im } b$  and  $\text{disj}: e * \text{Re } b < - (\text{Im } b * e) \vee e * \text{Re } b = - (\text{Im } b * e)$ 
      using  $\langle e > 0 \rangle \ \langle b \neq 0 \rangle$ 
      by (auto simp add:  $z\_def$  dist_norm sgn_if_less_eq_real_def mult_less_0_iff
        complex.expand split: if_split_asm)
      show False — or just one smt line
      using disj
      proof
        assume  $e * \text{Re } b < - (\text{Im } b * e)$ 
        with  $b \ \langle 0 < e \rangle$  less_trans [of_ 0] show False
      by (metis (no_types) mult_pos_pos neg_less_0_iff_less not_less_iff_gr_or_eq)
      next
        assume  $e * \text{Re } b = - (\text{Im } b * e)$ 
        with  $b \ \langle 0 < e \rangle$  show False
      by (metis mult_pos_pos neg_less_0_iff_less not_less_iff_gr_or_eq)
      qed
    qed
    ultimately show ?thesis
    using interior_mono by blast
  qed
} with assms show ?thesis by blast
qed

```

lemma $wn_triangle2_0$:

assumes $0 \in \text{interior}(\text{convex hull } \{a, b, c\})$

shows

$$\begin{aligned}
 & 0 < \text{Im}((b - a) * \text{cnj } (b)) \wedge \\
 & 0 < \text{Im}((c - b) * \text{cnj } (c)) \wedge \\
 & 0 < \text{Im}((a - c) * \text{cnj } (a)) \\
 & \vee \\
 & \text{Im}((b - a) * \text{cnj } (b)) < 0 \wedge
 \end{aligned}$$

$$0 < \text{Im}((b - c) * \text{cnj } (b)) \wedge$$

$$0 < \text{Im}((a - b) * \text{cnj } (a)) \wedge$$

$$0 < \text{Im}((c - a) * \text{cnj } (c))$$
proof –**have** [simp]: $\{b, c, a\} = \{a, b, c\}$ $\{c, a, b\} = \{a, b, c\}$ **by** auto**show** ?thesis**using** wn_triangle1 [OF assms] wn_triangle1 [of b c a] wn_triangle1 [of c a b] assms**by** (auto simp: algebra_simps Im_complex_div_gt_0 Im_complex_div_lt_0 not_le not_less)**qed****lemma** wn_triangle2:**assumes** $z \in \text{interior}(\text{convex hull } \{a, b, c\})$ **shows** $0 < \text{Im}((b - a) * \text{cnj } (b - z)) \wedge$ $0 < \text{Im}((c - b) * \text{cnj } (c - z)) \wedge$ $0 < \text{Im}((a - c) * \text{cnj } (a - z))$ \vee $\text{Im}((b - a) * \text{cnj } (b - z)) < 0 \wedge$ $0 < \text{Im}((b - c) * \text{cnj } (b - z)) \wedge$ $0 < \text{Im}((a - b) * \text{cnj } (a - z)) \wedge$ $0 < \text{Im}((c - a) * \text{cnj } (c - z))$ **proof** –**have** 0: $0 \in \text{interior}(\text{convex hull } \{a-z, b-z, c-z\})$ **using** assms convex_hull_translation [of $-z$ $\{a, b, c\}$] interior_translation [of $-z$]**by** (simp cong: image_cong_simp)**show** ?thesis **using** wn_triangle2_0 [OF 0]**by** simp**qed****lemma** wn_triangle3:**assumes** $z: z \in \text{interior}(\text{convex hull } \{a, b, c\})$ **and** $0 < \text{Im}((b-a) * \text{cnj } (b-z))$ $0 < \text{Im}((c-b) * \text{cnj } (c-z))$ $0 < \text{Im}((a-c) * \text{cnj } (a-z))$ **shows** $\text{winding_number } (\text{linepath } a \ b \ +++ \ \text{linepath } b \ c \ +++ \ \text{linepath } c \ a) \ z = 1$ **proof** –**have** znot[simp]: $z \notin \text{closed_segment } a \ b$ $z \notin \text{closed_segment } b \ c$ $z \notin \text{closed_segment } c \ a$ **using** z interior_of_triangle [of a b c]**by** (auto simp: closed_segment_def)**have** gt0: $0 < \text{Re } (\text{winding_number } (\text{linepath } a \ b \ +++ \ \text{linepath } b \ c \ +++ \ \text{linepath } c \ a) \ z)$ **using** assms**by** (simp add: winding_number_linepath_pos_lt_path_image_join winding_number_join_pos_combined)**have** lt2: $\text{Re } (\text{winding_number } (\text{linepath } a \ b \ +++ \ \text{linepath } b \ c \ +++ \ \text{linepath } c \ a) \ z) < 2$

```

    using winding_number_lt_half_linepath [of _ a b]
    using winding_number_lt_half_linepath [of _ b c]
    using winding_number_lt_half_linepath [of _ c a] znot
    by (fastforce simp add: winding_number_join_path_image_join)
  show ?thesis
    by (rule winding_number_eq_1) (simp_all add: path_image_join gt0 lt2)
qed

```

proposition *winding_number_triangle:*

```

  assumes z: z ∈ interior(convex hull {a,b,c})
  shows winding_number(linepath a b +++ linepath b c +++ linepath c a) z =
    (if 0 < Im((b - a) * conj (b - z)) then 1 else -1)

```

proof –

```

  have [simp]: {a,c,b} = {a,b,c} by auto
  have znot[simp]: z ∉ closed_segment a b z ∉ closed_segment b c z ∉ closed_segment
  c a
    using z_interior_of_triangle [of a b c]
    by (auto simp: closed_segment_def)
  then have [simp]: z ∉ closed_segment b a z ∉ closed_segment c b z ∉ closed_segment
  a c
    using closed_segment_commute by blast+
  have *: winding_number (linepath a b +++ linepath b c +++ linepath c a) z =
    winding_number (reversepath (linepath a c +++ linepath c b +++
  linepath b a)) z
    by (simp add: reversepath_joinpaths winding_number_join_not_in_path_image_join)
  show ?thesis
    apply (rule disjE [OF wn_triangle2 [OF z]])
    subgoal
      by (simp add: wn_triangle3 z)
    subgoal
      by (simp add: path_image_join winding_number_reversepath * wn_triangle3
  z)
    done
qed

```

3.8 Winding numbers for simple closed paths

lemma *winding_number_from_innerpath:*

```

  assumes simple_path c1 and c1: pathstart c1 = a pathfinish c1 = b
    and simple_path c2 and c2: pathstart c2 = a pathfinish c2 = b
    and simple_path c and c: pathstart c = a pathfinish c = b
    and c1c2: path_image c1 ∩ path_image c2 = {a,b}
    and c1c: path_image c1 ∩ path_image c = {a,b}
    and c2c: path_image c2 ∩ path_image c = {a,b}
    and ne_12: path_image c ∩ inside(path_image c1 ∪ path_image c2) ≠ {}
    and z: z ∈ inside(path_image c1 ∪ path_image c)
    and wn_d: winding_number (c1 +++ reversepath c) z = d
    and a ≠ b d ≠ 0
  obtains z ∈ inside(path_image c1 ∪ path_image c2) winding_number (c1 +++

```

```

reversepath c2) z = d
proof -
  obtain 0: inside(path_image c1  $\cup$  path_image c)  $\cap$  inside(path_image c2  $\cup$ 
path_image c) = {}
  and 1: inside(path_image c1  $\cup$  path_image c)  $\cup$  inside(path_image c2  $\cup$ 
path_image c)  $\cup$ 
    (path_image c - {a,b}) = inside(path_image c1  $\cup$  path_image c2)
  by (rule split_inside_simple_closed_curve
    [OF  $\langle$ simple_path c1 $\rangle$  c1  $\langle$ simple_path c2 $\rangle$  c2  $\langle$ simple_path c $\rangle$  c  $\langle$ a  $\neq$ 
b $\rangle$  c1c2 c1c c2c ne_12])
  have znot: z  $\notin$  path_image c z  $\notin$  path_image c1 z  $\notin$  path_image c2
  using union_with_outside z 1 by auto
  then have zout: z  $\in$  outside (path_image c  $\cup$  path_image c2)
  by (metis 0 CompII UnE disjoint_iff_not_equal sup.commute union_with_inside
z)
  have wn_cc2: winding_number (c +++ reversepath c2) z = 0
  by (rule winding_number_zero_in_outside; simp add: zout  $\langle$ simple_path c2 $\rangle$ 
c c2  $\langle$ simple_path c $\rangle$  simple_path_imp_path path_image_join)
  show ?thesis
  proof
    show z  $\in$  inside (path_image c1  $\cup$  path_image c2)
    using 1 z by blast
    have winding_number c1 z - winding_number c z = d
    using assms znot
    by (metis wn_d winding_number_join simple_path_imp_path winding_number_reversepath
add.commute path_image_reversepath path_reversepath pathstart_reversepath minus
add_conv_diff)
    then show winding_number (c1 +++ reversepath c2) z = d
    using wn_cc2 by (simp add: winding_number_join simple_path_imp_path
assms znot winding_number_reversepath)
  qed
qed

lemma simple_closed_path_wn1:
  fixes a::complex and e::real
  assumes 0 < e
  and sp_pl: simple_path(p +++ linepath (a - e) (a + e)) (is simple_path
?pae)
  and psp: pathstart p = a + e
  and pfp: pathfinish p = a - e
  and disj: ball a e  $\cap$  path_image p = {}
obtains z where z  $\in$  inside (path_image ?pae) cmod (winding_number ?pae z)
= 1
proof -
  have arc p and arc_lp: arc (linepath (a - e) (a + e))
  and pap: path_image p  $\cap$  path_image (linepath (a - e) (a + e))  $\subseteq$  {pathstart
p, a-e}
  using simple_path_join_loop_eq [of linepath (a - e) (a + e) p] assms by auto
  have mid_eq_a: midpoint (a - e) (a + e) = a

```

```

  by (simp add: midpoint_def)
with assms have a ∈ path_image ?pae
  by (simp add: assms path_image_join) (metis midpoint_in_closed_segment)
then have a ∈ frontier(inside (path_image ?pae))
  using assms by (simp add: Jordan_inside_outside )
with ‹0 < e› obtain c where c: c ∈ inside (path_image ?pae)
  and dac: dist a c < e
  by (auto simp: frontier_straddle)
then have c ∉ path_image ?pae
  using inside_no_overlap by blast
then have c ∉ path_image p c ∉ closed_segment (a - e) (a + e)
  by (simp_all add: assms path_image_join)
with ‹0 < e› dac have c ∉ affine_hull {a - e, a + e}
  by (simp add: segment_as_ball not_le)
with ‹0 < e› have *: ¬ collinear {a - e, c, a + e}
  using collinear_3_affine_hull [of a-e a+e] by (auto simp: insert_commute)
have 13: 1/3 + 1/3 + 1/3 = (1::real) by simp
have (1/3) *R (a - of_real e) + (1/3) *R c + (1/3) *R (a + of_real e) ∈
interior(convex_hull {a - e, c, a + e})
  using interior_convex_hull_3_minimal [OF * DIM_complex]
  by clarsimp (metis 13 zero_less_divide_1 iff_zero_less_numeral)
then obtain z where z: z ∈ interior(convex_hull {a - e, c, a + e}) by force
have [simp]: z ∉ closed_segment (a - e) c
  by (metis DIM_complex Diff_iff IntD2 inf_sup_absorb interior_of_triangle
z)
have [simp]: z ∉ closed_segment (a + e) (a - e)
  by (metis DIM_complex DiffD2 Un_iff interior_of_triangle z)
have [simp]: z ∉ closed_segment c (a + e)
  by (metis (no_types, lifting) DIM_complex DiffD2 Un_insert_right inf_sup_aci(5)
insertCI interior_of_triangle mk_disjoint_insert z)
show thesis
proof
  have norm (winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z) = 1
    using winding_number_triangle [OF z] by simp
  have zin: z ∈ inside (path_image (linepath (a + e) (a - e)) ∪ path_image p)
    and zeq: winding_number (linepath (a + e) (a - e) +++ reversepath p) z =
      winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z
  proof (rule winding_number_from_innerpath
[of linepath (a + e) (a - e) a+e a-e p
linepath (a + e) c +++ linepath c (a - e) z
winding_number (linepath (a - e) c +++ linepath c (a + e) +++
linepath (a + e) (a - e)) z])
    show sp_aec: simple_path (linepath (a + e) c +++ linepath c (a - e))
  proof (rule arc_imp_simple_path [OF arc_join])
    show arc (linepath (a + e) c)
  by (metis ‹c ∉ path_image p› arc_linepath pathstart_in_path_image psp)
    show arc (linepath c (a - e))

```

```

    by (metis ‹ $c \notin \text{path\_image } p$ › arc_linepath pathfinish_in_path_image pfp)
    show path_image (linepath (a + e) c)  $\cap$  path_image (linepath c (a - e))
 $\subseteq$  {pathstart (linepath c (a - e))}
    by clarsimp (metis * IntI Int_closed_segment closed_segment_commute
singleton_iff)
  qed auto
  show simple_path p
    using ‹arc p› arc_simple_path by blast
  show sp_ae2: simple_path (linepath (a + e) (a - e))
    using ‹arc p› arc_distinct_ends pfp psp by fastforce
  show pa: pathfinish (linepath (a + e) (a - e)) = a - e
    pathstart (linepath (a + e) c +++ linepath c (a - e)) = a + e
    pathfinish (linepath (a + e) c +++ linepath c (a - e)) = a - e
    pathstart p = a + e pathfinish p = a - e
    pathstart (linepath (a + e) (a - e)) = a + e
  by (simp_all add: assms)
  show 1: path_image (linepath (a + e) (a - e))  $\cap$  path_image p = {a + e,
a - e}
  proof
    show path_image (linepath (a + e) (a - e))  $\cap$  path_image p  $\subseteq$  {a + e, a
- e}
      using pap closed_segment_commute psp segment_convex_hull by fastforce
    show {a + e, a - e}  $\subseteq$  path_image (linepath (a + e) (a - e))  $\cap$  path_image
p
      using pap pathfinish_in_path_image pathstart_in_path_image pfp psp
  by fastforce
  qed
  show 2: path_image (linepath (a + e) (a - e))  $\cap$  path_image (linepath (a +
e) c +++ linepath c (a - e)) =
    {a + e, a - e} (is ?lhs = ?rhs)
  proof
    have  $\neg$  collinear {c, a + e, a - e}
      using * by (simp add: insert_commute)
    then have convex_hull {a + e, a - e}  $\cap$  convex_hull {a + e, c} = {a + e}
      convex_hull {a + e, a - e}  $\cap$  convex_hull {c, a - e} = {a - e}
    by (metis (full_types) Int_closed_segment insert_commute segment_convex_hull)+
    then show ?lhs  $\subseteq$  ?rhs
      by (metis Int_Un_distrib equalityD1 insert_is_Un path_image_join
path_image_linepath path_join_eq path_linepath segment_convex_hull simple_path_def
sp_aec)
    show ?rhs  $\subseteq$  ?lhs
      using segment_convex_hull by (simp add: path_image_join)
  qed
  have path_image p  $\cap$  path_image (linepath (a + e) c)  $\subseteq$  {a + e}
  proof (clarsimp simp: path_image_join)
    fix x
    assume x  $\in$  path_image p and x_ac: x  $\in$  closed_segment (a + e) c
    then have dist x a  $\geq$  e
      by (metis IntI all_not_in_conv disj_dist_commute mem_ball not_less)

```

```

    with  $x\_ac$   $dac \langle e > 0 \rangle$  show  $x = a + e$ 
  by (auto simp: norm_minus_commute dist_norm closed_segment_eq_open
dest: open_segment_furthest_le [where  $y=a$ ])
qed
moreover
have  $path\_image\ p \cap path\_image\ (linepath\ c\ (a - e)) \subseteq \{a - e\}$ 
proof (clarsimp simp: path_image_join)
  fix  $x$ 
  assume  $x \in path\_image\ p$  and  $x\_ac: x \in closed\_segment\ c\ (a - e)$ 
  then have  $dist\ x\ a \geq e$ 
    by (metis IntI all_not_in_conv disj dist_commute mem_ball not_less)
  with  $x\_ac$   $dac \langle e > 0 \rangle$  show  $x = a - e$ 
  by (auto simp: norm_minus_commute dist_norm closed_segment_eq_open
dest: open_segment_furthest_le [where  $y=a$ ])
qed
ultimately
have  $path\_image\ p \cap path\_image\ (linepath\ (a + e)\ c\ +++\ linepath\ c\ (a - e)) \subseteq \{a + e, a - e\}$ 
  by (force simp: path_image_join)
then show  $\exists: path\_image\ p \cap path\_image\ (linepath\ (a + e)\ c\ +++\ linepath\ c\ (a - e)) = \{a + e, a - e\}$ 
  using 1 2 by blast
show 4:  $path\_image\ (linepath\ (a + e)\ c\ +++\ linepath\ c\ (a - e)) \cap inside\ (path\_image\ (linepath\ (a + e)\ (a - e)) \cup path\_image\ p) \neq \{\}$ 
apply (clarsimp simp: path_image_join segment_convex_hull disjoint_iff_not_equal)
by (metis (no_types, opaque_lifting) UnI1 Un_commute c closed_segment_commute ends_in_segment(1) path_image_join
  path_image_linepath pathstart_linepath pfp segment_convex_hull)
show  $z \in inside\ (path\_image\ (linepath\ (a + e)\ (a - e)) \cup path\_image\ (linepath\ (a + e)\ c\ +++\ linepath\ c\ (a - e)))$ 
proof (simp add: path_image_join)
  show  $z \in inside\ (closed\_segment\ (a + e)\ (a - e) \cup (closed\_segment\ (a + e)\ c \cup closed\_segment\ c\ (a - e)))$ 
    by (metis z inside_of_triangle DIM_complex Un_commute closed_segment_commute)
  qed
show 5: winding_number
  ( $linepath\ (a + e)\ (a - e)\ +++\ reversepath\ (linepath\ (a + e)\ c\ +++\ linepath\ c\ (a - e))$ )  $z =$ 
  winding_number ( $linepath\ (a - e)\ c\ +++\ linepath\ c\ (a + e)\ +++\ linepath\ (a + e)\ (a - e)$ )  $z$ 
  by (simp add: reversepath_joinpaths path_image_join winding_number_join)
show 6: winding_number ( $linepath\ (a - e)\ c\ +++\ linepath\ c\ (a + e)\ +++\ linepath\ (a + e)\ (a - e)$ )  $z \neq 0$ 
  by (simp add: winding_number_triangle z)
show winding_number ( $linepath\ (a + e)\ (a - e)\ +++\ reversepath\ p$ )  $z =$ 
  winding_number ( $linepath\ (a - e)\ c\ +++\ linepath\ c\ (a + e)\ +++\ linepath\ (a + e)\ (a - e)$ )  $z$ 
by (metis 1 2 3 4 5 6 pa sp_aec sp_ae2  $\langle arc\ p \rangle \langle simple\_path\ p \rangle arc\_distinct\_ends$ 

```



```

winding_number_from_innerpath zin_inside)
  qed (use assms ⟨e > 0⟩ in auto)
  show z_inside: z ∈ inside (path_image ?pae)
  using zin by (simp add: assms path_image_join Un_commute closed_segment_commute)
  have cmod (winding_number ?pae z) = cmod ((winding_number(reversepath
?pae) z))
  proof (subst winding_number_reversepath)
    show path ?pae
      using simple_path_imp_path sp_pl by blast
    show z ∉ path_image ?pae
      by (metis IntI emptyE inside_no_overlap z_inside)
    qed (simp add: inside_def)
    also have ... = cmod (winding_number(linepath (a + e) (a - e) +++ re-
versepath p) z)
      by (simp add: pfp reversepath_joinpaths)
    also have ... = cmod (winding_number (linepath (a - e) c +++ linepath c (a
+ e) +++ linepath (a + e) (a - e)) z)
      by (simp add: zeq)
    also have ... = 1
      using z by (simp add: interior_of_triangle winding_number_triangle)
    finally show cmod (winding_number ?pae z) = 1 .
  qed
qed

```

lemma simple_closed_path_wn2:

```

  fixes a::complex and d e::real
  assumes 0 < d 0 < e
    and sp_pl: simple_path(p +++ linepath (a - d) (a + e))
    and psp: pathstart p = a + e
    and pfp: pathfinish p = a - d
  obtains z where z ∈ inside (path_image (p +++ linepath (a - d) (a + e)))
    cmod (winding_number (p +++ linepath (a - d) (a + e)) z) = 1
  proof -
    have [simp]: a + of_real x ∈ closed_segment (a - α) (a - β) ⟷ x ∈
closed_segment (-α) (-β) for x α β::real
      using closed_segment_translation_eq [of a]
      by (metis (no_types, opaque_lifting) add_uminus_conv_diff of_real_minus
of_real_closed_segment)
    have [simp]: a - of_real x ∈ closed_segment (a + α) (a + β) ⟷ -x ∈
closed_segment α β for x α β::real
      by (metis closed_segment_translation_eq diff_conv_add_uminus of_real_closed_segment
of_real_minus)
    have arc p and arc_lp: arc (linepath (a - d) (a + e)) and path p
      and pap: path_image p ∩ closed_segment (a - d) (a + e) ⊆ {a+e, a-d}
      using simple_path_join_loop_eq [of linepath (a - d) (a + e) p] assms
arc_imp_path by auto
    have 0 ∈ closed_segment (-d) e
      using ⟨0 < d⟩ ⟨0 < e⟩ closed_segment_eq_real_ivl by auto
    then have a ∈ path_image (linepath (a - d) (a + e))

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    using of_real_closed_segment [THEN iffD2]
    by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
of_real_closed_segment)
  then have  $a \notin \text{path\_image } p$ 
    using  $\langle 0 < d \rangle \langle 0 < e \rangle \text{ pap}$  by auto
  then obtain  $k$  where  $0 < k$  and  $k: \text{ball } a \ k \cap (\text{path\_image } p) = \{\}$ 
    using  $\langle 0 < e \rangle \langle \text{path } p \rangle \text{ not\_on\_path\_ball}$  by blast
  define  $kde$  where  $kde \equiv (\min k (\min d e)) / 2$ 
  have  $0 < kde \quad kde < k \quad kde < d \quad kde < e$ 
    using  $\langle 0 < k \rangle \langle 0 < d \rangle \langle 0 < e \rangle$  by (auto simp: kde_def)
  let  $?q = \text{linepath } (a + kde) (a + e) \text{ +++ } p \text{ +++ } \text{linepath } (a - d) (a - kde)$ 
  have  $-kde \in \text{closed\_segment } (-d) \ e$ 
    using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle kde < e \rangle \text{ closed\_segment\_eq\_real\_ivl}$  by auto
  then have  $a\_diff\_kde: a - kde \in \text{closed\_segment } (a - d) (a + e)$ 
    using of_real_closed_segment [THEN iffD2]
    by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
of_real_closed_segment)
  then have  $\text{clsub2}: \text{closed\_segment } (a - d) (a - kde) \subseteq \text{closed\_segment } (a - d) (a + e)$ 
    by (simp add: subset_closed_segment)
  then have  $\text{path\_image } p \cap \text{closed\_segment } (a - d) (a - kde) \subseteq \{a + e, a - d\}$ 
    using pap by force
  moreover
  have  $a + e \notin \text{path\_image } p \cap \text{closed\_segment } (a - d) (a - kde)$ 
    using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle 0 < e \rangle$  by (auto simp: closed_segment_eq_real_ivl)
  ultimately have  $\text{sub\_a\_diff\_d}: \text{path\_image } p \cap \text{closed\_segment } (a - d) (a - kde) \subseteq \{a - d\}$ 
    by blast
  have  $kde \in \text{closed\_segment } (-d) \ e$ 
    using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle kde < e \rangle \text{ closed\_segment\_eq\_real\_ivl}$  by auto
  then have  $a\_diff\_kde: a + kde \in \text{closed\_segment } (a - d) (a + e)$ 
    using of_real_closed_segment [THEN iffD2]
    by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
of_real_closed_segment)
  then have  $\text{clsub1}: \text{closed\_segment } (a + kde) (a + e) \subseteq \text{closed\_segment } (a - d) (a + e)$ 
    by (simp add: subset_closed_segment)
  then have  $\text{closed\_segment } (a + kde) (a + e) \cap \text{path\_image } p \subseteq \{a + e, a - d\}$ 
    using pap by force
  moreover
  have  $\text{closed\_segment } (a + kde) (a + e) \cap \text{closed\_segment } (a - d) (a - kde) = \{\}$ 
    proof (clarsimp intro!: equals0I)
      fix  $y$ 
      assume  $y1: y \in \text{closed\_segment } (a + kde) (a + e)$ 
        and  $y2: y \in \text{closed\_segment } (a - d) (a - kde)$ 
      obtain  $u::\text{real}$  where  $u: y = a + u$  and  $0 < u$ 

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proof –
  obtain  $\xi$  where  $\xi: y = (1 - \xi) *_R (a + kde) + \xi *_R (a + e)$  and  $0 \leq \xi \leq 1$ 
    using y1 by (auto simp: in_segment)
  show thesis
  proof
    show  $y = a + ((1 - \xi) * kde + \xi * e)$ 
      using  $\xi$  by (auto simp: scaleR_conv_of_real algebra_simps)
    have  $(1 - \xi) * kde + \xi * e \geq 0$ 
      using  $\langle 0 < kde \rangle \langle 0 \leq \xi \rangle \langle \xi \leq 1 \rangle \langle 0 < e \rangle$  by auto
    moreover have  $(1 - \xi) * kde + \xi * e \neq 0$ 
      using  $\langle 0 < kde \rangle \langle 0 \leq \xi \rangle \langle \xi \leq 1 \rangle \langle 0 < e \rangle$  by (auto simp: add_nonneg_eq_0_iff)
    ultimately show  $(1 - \xi) * kde + \xi * e > 0$  by simp
  qed
moreover
obtain v::real where  $v: y = a + v$  and  $v \leq 0$ 
  proof –
    obtain  $\xi$  where  $\xi: y = (1 - \xi) *_R (a - d) + \xi *_R (a - kde)$  and  $0 \leq \xi \leq 1$ 
      using y2 by (auto simp: in_segment)
    show thesis
    proof
      show  $y = a + (-((1 - \xi) * d + \xi * kde))$ 
        using  $\xi$  by (force simp: scaleR_conv_of_real algebra_simps)
      show  $-((1 - \xi) * d + \xi * kde) \leq 0$ 
        using  $\langle 0 < kde \rangle \langle 0 \leq \xi \rangle \langle \xi \leq 1 \rangle \langle 0 < d \rangle$ 
        by (metis add.left_neutral add_nonneg_nonneg le_diff_eq less_eq_real_def neg_le_0_iff_le_split_mult_pos_le)
    qed
  qed
  ultimately show False
    by auto
qed
moreover have  $a - d \notin \text{closed\_segment } (a + kde) (a + e)$ 
  using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle 0 < e \rangle$  by (auto simp: closed_segment_eq_real_ivl)
  ultimately have sub_a_plus_e:
     $\text{closed\_segment } (a + kde) (a + e) \cap (\text{path\_image } p \cup \text{closed\_segment } (a - d) (a - kde)) \subseteq \{a + e\}$ 
    by auto
  have  $kde \in \text{closed\_segment } (-kde) e$ 
    using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle kde < e \rangle$  closed_segment_eq_real_ivl by auto
  then have a_add_kde:  $a + kde \in \text{closed\_segment } (a - kde) (a + e)$ 
    using of_real_closed_segment [THEN iffD2]
    by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del: of_real_closed_segment)
  have  $\text{closed\_segment } (a - kde) (a + kde) \cap \text{closed\_segment } (a + kde) (a + e) = \{a + kde\}$ 

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    by (metis a_add_kde Int_closed_segment)
  moreover
  have path_image p  $\cap$  closed_segment (a - kde) (a + kde) = {}
  proof (rule equals0I, clarify)
    fix y assume y  $\in$  path_image p y  $\in$  closed_segment (a - kde) (a + kde)
    with equals0D [OF k, of y]  $\langle 0 < kde \rangle \langle kde < k \rangle$  show False
    by (auto simp: dist_norm dest: dist_decreases_closed_segment [where c=a])
  qed
  moreover
  have -kde  $\in$  closed_segment (-d) kde
    using  $\langle 0 < kde \rangle \langle kde < d \rangle \langle kde < e \rangle$  closed_segment_eq_real_ivl by auto
  then have a_diff_kde': a - kde  $\in$  closed_segment (a - d) (a + kde)
    using of_real_closed_segment [THEN iffD2]
    by (force dest: closed_segment_translation_eq [of a, THEN iffD2] simp del:
of_real_closed_segment)
  then have closed_segment (a - d) (a - kde)  $\cap$  closed_segment (a - kde) (a +
kde) = {a - kde}
    by (metis Int_closed_segment)
  ultimately
  have pa_subset_pm_kde: path_image ?q  $\cap$  closed_segment (a - kde) (a + kde)
 $\subseteq$  {a - kde, a + kde}
    by (auto simp: path_image_join assms)
  have ge_kde1:  $\exists y. x = a + \text{of\_real } y \wedge y \geq kde$  if x: x  $\in$  closed_segment (a +
kde) (a + e) for x
  proof -
    obtain u where  $0 \leq u \leq 1$  and u:  $x = (1 - u) *_R (a + kde) + u *_R (a
+ e)$ 
    using x by (auto simp: in_segment)
    then have kde  $\leq (1 - u) * kde + u * e$ 
      using  $\langle kde < e \rangle$  segment_bound_lemma by auto
    moreover have  $x = a + ((1 - u) * kde + u * e)$ 
      by (fastforce simp: u_algebra_simps scaleR_conv_of_real)
    ultimately
    show ?thesis by blast
  qed
  have ge_kde2:  $\exists y. x = a + \text{of\_real } y \wedge y \leq -kde$  if x: x  $\in$  closed_segment (a
- d) (a - kde) for x
  proof -
    obtain u where  $0 \leq u \leq 1$  and u:  $x = (1 - u) *_R (a - d) + u *_R (a -
kde)$ 
    using x by (auto simp: in_segment)
    then have kde  $\leq ((1 - u) * d + u * kde)$ 
      using  $\langle kde < d \rangle$  segment_bound_lemma by auto
    moreover have  $x = a - ((1 - u) * d + u * kde)$ 
      by (fastforce simp: u_algebra_simps scaleR_conv_of_real)
    ultimately show ?thesis
      by (metis add_uminus_conv_diff neg_le_iff_le of_real_minus)
  qed
  have notin_paq: x  $\notin$  path_image ?q if dist a x < kde for x

```

```

proof -
  have  $x \notin \text{path\_image } p$ 
    using  $k \langle kde < k \rangle$  that by force
  then show  $?thesis$ 
    using that assms by (auto simp: path_image_join dist_norm dest!: ge_kde1 ge_kde2)
  qed
  obtain  $z$  where  $zin: z \in \text{inside } (\text{path\_image } (?q +++ \text{linepath } (a - kde) (a + kde)))$ 
    and  $z1: \text{cmod } (\text{winding\_number } (?q +++ \text{linepath } (a - kde) (a + kde))) z = 1$ 
  proof (rule simple_closed_path_wn1 [of kde ?q a])
    show  $\text{simple\_path } (?q +++ \text{linepath } (a - kde) (a + kde))$ 
    proof (intro simple_path_join_loop conjI)
      show  $\text{arc } ?q$ 
      proof (rule arc_join)
        show  $\text{arc } (\text{linepath } (a + kde) (a + e))$ 
          using  $\langle kde < e \rangle \langle \text{arc } p \rangle$  by (force simp: pfp)
        show  $\text{arc } (p +++ \text{linepath } (a - d) (a - kde))$ 
          using  $\langle kde < d \rangle \langle kde < e \rangle \langle \text{arc } p \rangle \text{sub\_a\_diff\_d}$  by (force simp: pfp intro: arc_join)
        qed (auto simp: psp pfp path_image_join sub_a_plus_e)
        show  $\text{arc } (\text{linepath } (a - kde) (a + kde))$ 
          using  $\langle 0 < kde \rangle$  by auto
        qed (use pa_subset_pm_kde in auto)
        qed (use  $\langle 0 < kde \rangle$  notin_paq in auto)
        have  $\text{eq: path\_image } (?q +++ \text{linepath } (a - kde) (a + kde)) = \text{path\_image } (p +++ \text{linepath } (a - d) (a + e))$ 
          (is  $?lhs = ?rhs$ )
      proof
        show  $?lhs \subseteq ?rhs$ 
          using clsub1 clsub2 apply (auto simp: path_image_join assms)
          by (meson subsetCE subset_closed_segment)
        show  $?rhs \subseteq ?lhs$ 
          apply (simp add: path_image_join assms Un_ac)
          by (metis Un_closed_segment Un_assoc a_diff_kde a_diff_kde' le_supI2 subset_refl)
        qed
        show thesis
        proof
          show  $zzin: z \in \text{inside } (\text{path\_image } (p +++ \text{linepath } (a - d) (a + e)))$ 
            by (metis eq zin)
          then have  $znotin: z \notin \text{path\_image } p$ 
            by (metis ComplD Un_iff inside_Un_outside path_image_join pathfinish_linepath pathstart_reversepath pfp reversepath_linepath)
          have  $znotin\_d\_kde: z \notin \text{closed\_segment } (a - d) (a + kde)$ 
            by (metis ComplD Un_iff Un_closed_segment a_diff_kde inside_Un_outside path_image_join path_image_linepath pathstart_linepath pfp zzin)
          have  $znotin\_d\_e: z \notin \text{closed\_segment } (a - d) (a + e)$ 

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  by (metis IntI UnCI emptyE inside_no_overlap path_image_join path_image_linepath
  pathstart_linepath pfp zzin)
  have znotin_kde_e:  $z \notin \text{closed\_segment } (a + kde) (a + e)$  and znotin_d_kde':
 $z \notin \text{closed\_segment } (a - d) (a - kde)$ 
  using clsub1 clsub2 zzin
  by (metis (no_types, opaque_lifting) IntI Un_iff emptyE inside_no_overlap
  path_image_join path_image_linepath pathstart_linepath pfp subsetD)+
  have winding_number (linepath (a - d) (a + e)) z =
    winding_number (linepath (a - d) (a + kde)) z + winding_number
  (linepath (a + kde) (a + e)) z
  proof (rule winding_number_split_linepath)
    show  $a + \text{complex\_of\_real } kde \in \text{closed\_segment } (a - d) (a + e)$ 
    by (simp add: a_diff_kde)
    show  $z \notin \text{closed\_segment } (a - d) (a + e)$ 
    by (metis ComplD Un_iff inside_Un_outside path_image_join path_image_linepath
  pathstart_linepath pfp zzin)
  qed
  also have ... = winding_number (linepath (a + kde) (a + e)) z +
    (winding_number (linepath (a - d) (a - kde)) z + winding_number
  (linepath (a - kde) (a + kde)) z)
  by (simp add: winding_number_split_linepath [of a-kde, symmetric] znotin_d_kde
  a_diff_kde')
  finally have winding_number (p +++ linepath (a - d) (a + e)) z =
    winding_number p z + winding_number (linepath (a + kde) (a +
  e)) z +
    (winding_number (linepath (a - d) (a - kde)) z +
    winding_number (linepath (a - kde) (a + kde)) z)
  by (metis (no_types, lifting) ComplD Un_iff ⟨arc p⟩ add.assoc arc_imp_path
  eq path_image_join path_join_path_ends path_linepath simple_path_imp_path
  sp_pl union_with_outside winding_number_join zin)
  also have ... = winding_number (linepath (a + kde) (a + e)) z
    + winding_number (p +++ linepath (a - d) (a - kde)) z
    + winding_number (linepath (a - kde) (a + kde)) z
  using ⟨path p⟩ znotin assms
  by simp (metis Un_iff Un_closed_segment a_diff_kde' path_image_linepath
  path_linepath pathstart_linepath winding_number_join znotin_d_kde)
  also have ... = winding_number ?q z + winding_number (linepath (a - kde)
  (a + kde)) z
  using ⟨path p⟩ znotin assms by (simp add: path_image_join winding_number_join
  znotin_kde_e znotin_d_kde')
  also have ... = winding_number (?q +++ linepath (a - kde) (a + kde)) z
  by (metis (mono_tags, lifting) ComplD UnCI ⟨path p⟩ eq inside_outside
  path_image_join path_join_eq path_linepath pathfinish_join pathfinish_linepath
  pathstart_join pathstart_linepath pfp psp winding_number_join zzin)
  finally have winding_number (p +++ linepath (a - d) (a + e)) z =
    winding_number (?q +++ linepath (a - kde) (a + kde)) z .
  then show cmod (winding_number (p +++ linepath (a - d) (a + e)) z) = 1
  by (simp add: z1)
  qed

```

qed

lemma *simple_closed_path_wn3*:

fixes $p :: \text{real} \Rightarrow \text{complex}$

assumes *simple_path* p **and** *loop*: $\text{pathfinish } p = \text{pathstart } p$

obtains z **where** $z \in \text{inside } (\text{path_image } p)$ $\text{cmod } (\text{winding_number } p \ z) = 1$

proof –

have ins : $\text{inside}(\text{path_image } p) \neq \{\}$ $\text{open}(\text{inside}(\text{path_image } p))$
 $\text{connected}(\text{inside}(\text{path_image } p))$

and out : $\text{outside}(\text{path_image } p) \neq \{\}$ $\text{open}(\text{outside}(\text{path_image } p))$
 $\text{connected}(\text{outside}(\text{path_image } p))$

and bo : $\text{bounded}(\text{inside}(\text{path_image } p)) \neg \text{bounded}(\text{outside}(\text{path_image } p))$

and ins_out : $\text{inside}(\text{path_image } p) \cap \text{outside}(\text{path_image } p) = \{\}$
 $\text{inside}(\text{path_image } p) \cup \text{outside}(\text{path_image } p) = - \text{path_image } p$

and fro : $\text{frontier}(\text{inside}(\text{path_image } p)) = \text{path_image } p$
 $\text{frontier}(\text{outside}(\text{path_image } p)) = \text{path_image } p$

using *Jordan_inside_outside* [*OF assms*] **by** *auto*

obtain a **where** $a \in \text{inside}(\text{path_image } p)$

using $\langle \text{inside } (\text{path_image } p) \neq \{\} \rangle$ **by** *blast*

obtain $d :: \text{real}$ **where** $0 < d$ **and** d_fro : $a - d \in \text{frontier } (\text{inside } (\text{path_image } p))$

and d_int : $\bigwedge \varepsilon. \llbracket 0 \leq \varepsilon; \varepsilon < d \rrbracket \implies (a - \varepsilon) \in \text{inside } (\text{path_image } p)$

using *ray_to_frontier* [*of inside (path_image p) a -1*] $\text{bo ins a interior_eq}$

by (*metis ab_group_add_class.ab_diff_conv_add_uminus of_real_def scale_minus_right zero_neq_neg_one*)

obtain $e :: \text{real}$ **where** $0 < e$ **and** e_fro : $a + e \in \text{frontier } (\text{inside } (\text{path_image } p))$

and e_int : $\bigwedge \varepsilon. \llbracket 0 \leq \varepsilon; \varepsilon < e \rrbracket \implies (a + \varepsilon) \in \text{inside } (\text{path_image } p)$

using *ray_to_frontier* [*of inside (path_image p) a 1*] $\text{bo ins a interior_eq}$

by (*metis of_real_def zero_neq_one*)

obtain $t0$ **where** $0 \leq t0$ $t0 \leq 1$ **and** pt : $p \ t0 = a - d$

using $a \ d_fro \ \text{fro}$ **by** (*auto simp: path_image_def*)

obtain q **where** *simple_path* q **and** q_ends : $\text{pathstart } q = a - d$ $\text{pathfinish } q = a - d$

and q_eq_p : $\text{path_image } q = \text{path_image } p$

and $\text{wn}_q \ \text{eq} \ \text{wn}_p$: $\bigwedge z. z \in \text{inside}(\text{path_image } p) \implies \text{winding_number } q \ z = \text{winding_number } p \ z$

proof

show *simple_path* (*shiftpath* $t0$ p)

by (*simp add: pathstart_shiftpath pathfinish_shiftpath*

simple_path_shiftpath path_image_shiftpath $\langle 0 \leq t0 \rangle \langle t0 \leq 1 \rangle$ *assms*)

show $\text{pathstart } (\text{shiftpath } t0 \ p) = a - d$

using pt **by** (*simp add:* $\langle t0 \leq 1 \rangle$ *pathstart_shiftpath*)

show $\text{pathfinish } (\text{shiftpath } t0 \ p) = a - d$

by (*simp add:* $\langle 0 \leq t0 \rangle$ *loop pathfinish_shiftpath* pt)

show $\text{path_image } (\text{shiftpath } t0 \ p) = \text{path_image } p$

by (*simp add:* $\langle 0 \leq t0 \rangle \langle t0 \leq 1 \rangle$ *loop path_image_shiftpath*)

show $\text{winding_number } (\text{shiftpath } t0 \ p) \ z = \text{winding_number } p \ z$

if $z \in \text{inside } (\text{path_image } p)$ **for** z

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    by (metis ComplD Un_iff ‹0 ≤ t0› ‹t0 ≤ 1› ‹simple_path p› atLeastAt-
Most_iff inside_Un_outside
    loop simple_path_imp_path that winding_number_shiftpath)
  qed
  have ad_not_ae: a - d ≠ a + e
  by (metis ‹0 < d› ‹0 < e› add.left_inverse add_left_cancel add_uminus_conv_diff
    le_add_same_cancel2 less_eq_real_def not_less of_real_add of_real_def
of_real_eq_0_iff pt)
  have ad_ae_q: {a - d, a + e} ⊆ path_image q
  using ‹path_image q = path_image p› d_fro e_fro fro(1) by auto
  have ada: open_segment (a - d) a ⊆ inside (path_image p)
  proof (clarsimp simp: in_segment)
    fix u::real assume 0 < u < 1
    with d_int have a - (1 - u) * d ∈ inside (path_image p)
    by (metis ‹0 < d› add commute diff_add_cancel left_diff_distrib' less_add_same_cancel2
less_eq_real_def mult.left_neutral zero_less_mult_iff)
    then show (1 - u) *R (a - d) + u *R a ∈ inside (path_image p)
    by (simp add: diff_add_eq of_real_def real_vector.scale_right_diff_distrib)
  qed
  have aae: open_segment a (a + e) ⊆ inside (path_image p)
  proof (clarsimp simp: in_segment)
    fix u::real assume 0 < u < 1
    with e_int have a + u * e ∈ inside (path_image p)
    by (meson ‹0 < e› less_eq_real_def mult_less_cancel_right2 not_less
zero_less_mult_iff)
    then show (1 - u) *R a + u *R (a + e) ∈ inside (path_image p)
    by (metis (mono_tags, lifting) add.assoc of_real_mult pth_6 scaleR_collapse
scaleR_conv_of_real)
  qed
  have complex_of_real (d * d + (e * e + d * (e + e))) ≠ 0
  using ad_not_ae
  by (metis ‹0 < d› ‹0 < e› add_strict_left_mono less_add_same_cancel1
not_sum_squares_lt_zero
of_real_eq_0_iff zero_less_double_add_iff zero_less_single_add zero_less_mult_iff)
  moreover have ∃ u > 0. u < 1 ∧ d = u * d + u * e
  using ‹0 < d› ‹0 < e›
  by (rule_tac x=d / (d+e) in exI) (simp add: divide_simps scaleR_conv_of_real
flip: distrib_left)
  ultimately have a_in_de: a ∈ open_segment (a - d) (a + e)
  using ad_not_ae by (simp add: in_segment algebra_simps scaleR_conv_of_real
flip: of_real_add of_real_mult)
  then have open_segment (a - d) (a + e) ⊆ inside (path_image p)
  using ada aae Un_open_segment [of a -d a+e] by blast
  then have path_image q ∩ open_segment (a - d) (a + e) = {}
  using inside_no_overlap by (fastforce simp: q_eq_p)
  with ad_ae_q have paq_Int_cs: path_image q ∩ closed_segment (a - d) (a +
e) = {a - d, a + e}
  by (simp add: closed_segment_eq_open)
  obtain t where 0 ≤ t ≤ 1 and qt: q t = a + e

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    using a_e_fro fro ad_ae_q by (auto simp: path_defs)
  then have t ≠ 0
    by (metis ad_not_ae pathstart_def q_ends(1))
  then have t ≠ 1
    by (metis ad_not_ae pathfinish_def q_ends(2) qt)
  have q01: q 0 = a - d q 1 = a - d
    using q_ends by (auto simp: pathstart_def pathfinish_def)
  obtain z where zin: z ∈ inside (path_image (subpath t 0 q +++ linepath (a -
d) (a + e)))
    and z1: cmod (winding_number (subpath t 0 q +++ linepath (a - d)
(a + e)) z) = 1
  proof (rule simple_closed_path_wn2 [of d e subpath t 0 q a], simp_all add: q01)
    show simple_path (subpath t 0 q +++ linepath (a - d) (a + e))
    proof (rule simple_path_join_loop, simp_all add: qt q01)
      have inj_on q (closed_segment t 0)
        using ‹0 ≤ t› ‹simple_path q› ‹t ≤ 1› ‹t ≠ 0› ‹t ≠ 1›
      by (fastforce simp: simple_path_def loop_free_def inj_on_def closed_segment_eq_real_ivl)
      then show arc (subpath t 0 q)
        using ‹0 ≤ t› ‹simple_path q› ‹t ≤ 1› ‹t ≠ 0›
        by (simp add: arc_subpath_eq simple_path_imp_path)
      show arc (linepath (a - d) (a + e))
        by (simp add: ad_not_ae)
      show path_image (subpath t 0 q) ∩ closed_segment (a - d) (a + e) ⊆ {a +
e, a - d}
        using qt paq_Int_cs ‹simple_path q› ‹0 ≤ t› ‹t ≤ 1›
      by (force simp: dest: rev_subsetD [OF path_image_subpath_subset] intro:
simple_path_imp_path)
    qed
  qed (auto simp: ‹0 < d› ‹0 < e› qt)
  have pa01_Un: path_image (subpath 0 t q) ∪ path_image (subpath 1 t q) =
path_image q
    unfolding path_image_subpath
    using ‹0 ≤ t› ‹t ≤ 1› by (force simp: path_image_def image_iff)
  with paq_Int_cs have pa_01q:
    (path_image (subpath 0 t q) ∪ path_image (subpath 1 t q)) ∩ closed_segment
(a - d) (a + e) = {a - d, a + e}
    by metis
  have z_notin_ed: z ∉ closed_segment (a + e) (a - d)
    using zin q01 by (simp add: path_image_join closed_segment_commute in-
side_def)
  have z_notin_0t: z ∉ path_image (subpath 0 t q)
    by (metis (no_types, opaque_lifting) IntI Un_upper1 subsetD empty_iff in-
side_no_overlap path_image_join
    path_image_subpath_commute pathfinish_subpath pathstart_def pathstart_linepath
q_ends(1) qt subpath_trivial zin)
  have [simp]: - winding_number (subpath t 0 q) z = winding_number (subpath
0 t q) z
    by (metis ‹0 ≤ t› ‹simple_path q› ‹t ≤ 1› atLeastAtMost_iff zero_le_one
    path_image_subpath_commute path_subpath_real_eq_0_iff_le_ge_0

```

```

    reversepath_subpath simple_path_imp_path winding_number_reversepath
z_notin_0t)
  obtain z_in_q: z ∈ inside(path_image q)
    and wn_q: winding_number (subpath 0 t q +++ subpath t 1 q) z = -
winding_number (subpath t 0 q +++ linepath (a - d) (a + e)) z
  proof (rule winding_number_from_innerpath
    [of subpath 0 t q a-d a+e subpath 1 t q linepath (a - d) (a + e)
      z - winding_number (subpath t 0 q +++ linepath (a - d) (a + e)) z],
    simp_all add: q01 qt pa01_Un reversepath_subpath)
    show simple_path (subpath 0 t q) simple_path (subpath 1 t q)
      by (simp_all add: <0 ≤ t> <simple_path q> <t ≤ 1> <t ≠ 0> <t ≠ 1> sim-
ple_path_subpath)
    show simple_path (linepath (a - d) (a + e))
      using ad_not_ae by blast
    show path_image (subpath 0 t q) ∩ path_image (subpath 1 t q) = {a - d, a +
e} (is ?lhs = ?rhs)
  proof
    show ?lhs ⊆ ?rhs
      using <0 ≤ t> <simple_path q> <t ≤ 1> <t ≠ 1> q_ends qt q01
      by (force simp: pathfinish_def qt simple_path_def loop_free_def path_image_subpath)
    show ?rhs ⊆ ?lhs
      using <0 ≤ t> <t ≤ 1> q01 qt by (force simp: path_image_subpath)
  qed
  show path_image (subpath 0 t q) ∩ closed_segment (a - d) (a + e) = {a -
d, a + e} (is ?lhs = ?rhs)
  proof
    show ?lhs ⊆ ?rhs using paq_Int_cs pa01_Un by fastforce
    show ?rhs ⊆ ?lhs using <0 ≤ t> <t ≤ 1> q01 qt by (force simp: path_image_subpath)
  qed
  show path_image (subpath 1 t q) ∩ closed_segment (a - d) (a + e) = {a -
d, a + e} (is ?lhs = ?rhs)
  proof
    show ?lhs ⊆ ?rhs by (auto simp: pa_01q [symmetric])
    show ?rhs ⊆ ?lhs using <0 ≤ t> <t ≤ 1> q01 qt by (force simp: path_image_subpath)
  qed
  show closed_segment (a - d) (a + e) ∩ inside (path_image q) ≠ {}
    using a_a_in_de open_closed_segment pa01_Un q_eq_p by fastforce
  show z ∈ inside (path_image (subpath 0 t q) ∪ closed_segment (a - d) (a +
e))
    by (metis path_image_join path_image_linepath path_image_subpath_commute
pathfinish_subpath pathstart_linepath q01(1) zin)
  show winding_number (subpath 0 t q +++ linepath (a + e) (a - d)) z =
- winding_number (subpath t 0 q +++ linepath (a - d) (a + e)) z
    using z_notin_ed z_notin_0t <0 ≤ t> <simple_path q> <t ≤ 1>
    by (simp add: simple_path_imp_path qt q01 path_image_subpath_commute
closed_segment_commute winding_number_join winding_number_reversepath [symmetric])
  show - d ≠ e
    using ad_not_ae by auto
  show winding_number (subpath t 0 q +++ linepath (a - d) (a + e)) z ≠ 0

```

```

    using z1 by auto
  qed
  show ?thesis
  proof
    show  $z \in \text{inside}(\text{path\_image } p)$ 
      using  $q\_eq\_p \ z\_in\_q$  by auto
    then have [simp]:  $z \notin \text{path\_image } q$ 
      by (metis disjoint_iff_not_equal inside_no_overlap q_eq_p)
    have [simp]:  $z \notin \text{path\_image}(\text{subpath } 1 \ t \ q)$ 
      using inside_def pa01_Un z_in_q by fastforce
    have  $\text{winding\_number}(\text{subpath } 0 \ t \ q \ +++ \ \text{subpath } t \ 1 \ q) \ z = \text{winding\_number}(\text{subpath } 0 \ 1 \ q) \ z$ 
      using  $z \ \text{notin } 0t \ \langle 0 \leq t \rangle \ \langle \text{simple\_path } q \rangle \ \langle t \leq 1 \rangle$ 
      by (simp add: simple_path_imp_path qt_path_image_subpath_commute winding_number_join winding_number_subpath_combine)
    with wn_q have  $\text{winding\_number}(\text{subpath } t \ 0 \ q \ +++ \ \text{linepath } (a - d) \ (a + e)) \ z = - \text{winding\_number } q \ z$ 
      by auto
    with z1 have  $\text{cmod}(\text{winding\_number } q \ z) = 1$ 
      by simp
    with z1 wn_q_eq_wn_p show  $\text{cmod}(\text{winding\_number } p \ z) = 1$ 
      using z1 wn_q_eq_wn_p by (simp add:  $\langle z \in \text{inside}(\text{path\_image } p) \rangle$ )
  qed
  qed

```

proposition *simple_closed_path_winding_number_inside:*

```

  assumes simple_path  $\gamma$ 
  obtains  $\bigwedge z. z \in \text{inside}(\text{path\_image } \gamma) \implies \text{winding\_number } \gamma \ z = 1$ 
    |  $\bigwedge z. z \in \text{inside}(\text{path\_image } \gamma) \implies \text{winding\_number } \gamma \ z = -1$ 
  proof (cases pathfinish  $\gamma = \text{pathstart } \gamma$ )
  case True
  have path  $\gamma$ 
    by (simp add: assms simple_path_imp_path)
  then have const:  $\text{winding\_number } \gamma \ \text{constant\_on } \text{inside}(\text{path\_image } \gamma)$ 
  proof (rule winding_number_constant)
    show connected ( $\text{inside}(\text{path\_image } \gamma)$ )
      by (simp add: Jordan_inside_outside True assms)
  qed (use inside_no_overlap True in auto)
  obtain z where zin:  $z \in \text{inside}(\text{path\_image } \gamma)$  and z1:  $\text{cmod}(\text{winding\_number } \gamma \ z) = 1$ 
    using simple_closed_path_wn3 [of  $\gamma$ ] True assms by blast
  have winding_number  $\gamma \ z \in \mathbb{Z}$ 
    using zin integer_winding_number [OF  $\langle \text{path } \gamma \rangle$  True] inside_def by blast
  moreover have  $\text{real\_of\_int } x = -1 \iff x = -1$  for x
    by linarith
  ultimately consider  $\text{winding\_number } \gamma \ z = 1 \mid \text{winding\_number } \gamma \ z = -1$ 
    using z1 by (auto simp: Ints_def abs_if_split: if_split_asm)
  with that const zin show ?thesis
    unfolding constant_on_def by metis

```

```

next
  case False
  then show ?thesis
    using inside_simple_curve_imp_closed assms that(2) by blast
qed

lemma simple_closed_path_abs_winding_number_inside:
  assumes simple_path  $\gamma$   $z \in \text{inside}(\text{path\_image } \gamma)$ 
  shows  $|\text{Re}(\text{winding\_number } \gamma z)| = 1$ 
  by (metis assms norm_minus_cancel norm_one one_complex.simps(1) real_norm_def
simple_closed_path_winding_number_inside uminus_complex.simps(1))

lemma simple_closed_path_norm_winding_number_inside:
  assumes simple_path  $\gamma$   $z \in \text{inside}(\text{path\_image } \gamma)$ 
  shows  $\text{norm}(\text{winding\_number } \gamma z) = 1$ 
proof -
  have pathfinish  $\gamma = \text{pathstart } \gamma$ 
  using assms inside_simple_curve_imp_closed by blast
  with assms integer_winding_number have winding_number  $\gamma z \in \mathbb{Z}$ 
  by (simp add: inside_def simple_path_def)
  then show ?thesis
  by (metis assms norm_minus_cancel norm_one simple_closed_path_winding_number_inside)
qed

lemma simple_closed_path_winding_number_cases:
  assumes simple_path  $\gamma$  pathfinish  $\gamma = \text{pathstart } \gamma$   $z \notin \text{path\_image } \gamma$ 
  shows winding_number  $\gamma z \in \{-1, 0, 1\}$ 
proof -
  consider  $z \in \text{inside}(\text{path\_image } \gamma) \mid z \in \text{outside}(\text{path\_image } \gamma)$ 
  by (metis ComplI UnE assms(3) inside_Un_outside)
  then show ?thesis
  proof cases
    case 1
    then show ?thesis
    using assms simple_closed_path_winding_number_inside by auto
  next
    case 2
    then show ?thesis
    using assms simple_path_def winding_number_zero_in_outside by blast
  qed
qed

lemma simple_closed_path_winding_number_pos:
  [[simple_path  $\gamma$ ; pathfinish  $\gamma = \text{pathstart } \gamma$ ;  $z \notin \text{path\_image } \gamma$ ;  $0 < \text{Re}(\text{winding\_number } \gamma z)$ ]
   $\implies \text{winding\_number } \gamma z = 1$ 
using simple_closed_path_winding_number_cases
  by fastforce

```

lemma *simple_closed_path_winding_number_neg*:
 $\llbracket \text{simple_path } \gamma; \text{pathfinish } \gamma = \text{pathstart } \gamma; z \notin \text{path_image } \gamma; \text{Re } (\text{winding_number } \gamma z) < 0 \rrbracket$
 $\implies \text{winding_number } \gamma z = -1$
using *simple_closed_path_winding_number_cases* **by** *fastforce*

3.9 Winding number for rectangular paths

proposition *winding_number_rectpath*:
assumes $z \in \text{box } a1 a3$
shows $\text{winding_number } (\text{rectpath } a1 a3) z = 1$
proof –
from *assms* **have** *less*: $\text{Re } a1 < \text{Re } a3$ $\text{Im } a1 < \text{Im } a3$
by (*auto simp: in_box_complex_iff*)
define $a2 a4$ **where** $a2 = \text{Complex } (\text{Re } a3) (\text{Im } a1)$ **and** $a4 = \text{Complex } (\text{Re } a1) (\text{Im } a3)$
let $?l1 = \text{linepath } a1 a2$ **and** $?l2 = \text{linepath } a2 a3$
and $?l3 = \text{linepath } a3 a4$ **and** $?l4 = \text{linepath } a4 a1$
from *assms* **and** *less* **have** $z \notin \text{path_image } (\text{rectpath } a1 a3)$
by (*auto simp: path_image_rectpath_cbox_minus_box*)
also **have** $\text{path_image } (\text{rectpath } a1 a3) =$
 $\text{path_image } ?l1 \cup \text{path_image } ?l2 \cup \text{path_image } ?l3 \cup \text{path_image } ?l4$
by (*simp add: rectpath_def Let_def path_image_join Un_assoc a2_def a4_def*)
finally **have** $z \notin \dots$
moreover **have** $\forall l \in \{?l1, ?l2, ?l3, ?l4\}. \text{Re } (\text{winding_number } l z) > 0$
unfolding *ball_simps HOL.simp_thms a2_def a4_def*
by (*intro conjI; (rule winding_number_linepath_pos_lt;*
(insert assms, auto simp: a2_def a4_def in_box_complex_iff mult_neg_neg)))+
ultimately **have** $\text{Re } (\text{winding_number } (\text{rectpath } a1 a3) z) > 0$
by (*simp add: winding_number_join path_image_join rectpath_def Let_def a2_def a4_def*)
thus $\text{winding_number } (\text{rectpath } a1 a3) z = 1$ **using** *assms less*
by (*intro simple_closed_path_winding_number_pos simple_path_rectpath*)
(auto simp: path_image_rectpath_cbox_minus_box)
qed

proposition *winding_number_rectpath_outside*:
assumes $\text{Re } a1 \leq \text{Re } a3$ $\text{Im } a1 \leq \text{Im } a3$
assumes $z \notin \text{cbox } a1 a3$
shows $\text{winding_number } (\text{rectpath } a1 a3) z = 0$
using *assms* **by** (*intro winding_number_zero_outside[OF ___ assms(3)]*
 $\text{path_image_rectpath_subset_cbox}$) *simp_all*

A per-function version for continuous logs, a kind of monodromy

proposition *winding_number_compose_exp*:
assumes *path p*
shows $\text{winding_number } (\text{exp } \circ p) 0 = (\text{pathfinish } p - \text{pathstart } p) / (2 * \text{of_real } \pi * i)$

```

proof –
  obtain  $e$  where  $0 < e$  and  $e: \bigwedge t. t \in \{0..1\} \implies e \leq \text{norm}(\text{exp}(p\ t))$ 
  proof
    have  $\text{closed}(\text{path\_image}(\text{exp} \circ p))$ 
      by (simp add: assms closed_path_image holomorphic_on_exp holomorphic_on_imp_continuous_on path_continuous_image)
    then show  $0 < \text{setdist}\ \{0\}\ (\text{path\_image}(\text{exp} \circ p))$ 
      by (metis exp_not_eq_zero imageE image_comp infdist_eq_setdist infdist_pos_not_in_closed path_defs(4) path_image_nonempty)
  next
    fix  $t::\text{real}$ 
    assume  $t \in \{0..1\}$ 
    have  $\text{setdist}\ \{0\}\ (\text{path\_image}(\text{exp} \circ p)) \leq \text{dist}\ 0\ (\text{exp}\ (p\ t))$ 
    proof (rule setdist_le_dist)
      show  $\text{exp}\ (p\ t) \in \text{path\_image}(\text{exp} \circ p)$ 
        using  $\langle t \in \{0..1\} \rangle$  path_image_def by fastforce+
      qed auto
    then show  $\text{setdist}\ \{0\}\ (\text{path\_image}(\text{exp} \circ p)) \leq \text{cmod}(\text{exp}\ (p\ t))$ 
      by simp
    qed
  have  $\text{bounded}(\text{path\_image}\ p)$ 
    by (simp add: assms bounded_path_image)
  then obtain  $B$  where  $0 < B$  and  $B: \text{path\_image}\ p \subseteq \text{cball}\ 0\ B$ 
    by (meson bounded_pos mem_cball_0 subsetI)
  let  $?B = \text{cball}\ (0::\text{complex})\ (B+1)$ 
  have  $\text{uniformly\_continuous\_on}\ ?B\ \text{exp}$ 
    using holomorphic_on_exp holomorphic_on_imp_continuous_on
    by (force intro: compact_uniformly_continuous)
  then obtain  $d$  where  $d > 0$ 
    and  $d: \bigwedge x\ x'. \llbracket x \in ?B; x' \in ?B; \text{dist}\ x'\ x < d \rrbracket \implies \text{norm}(\text{exp}\ x' - \text{exp}\ x) < e$ 
    using  $\langle e > 0 \rangle$  by (auto simp: uniformly_continuous_on_def dist_norm)
  then have  $\text{min}\ 1\ d > 0$ 
    by force
  then obtain  $g$  where pfg: polynomial_function  $g$  and  $g\ 0 = p\ 0\ g\ 1 = p\ 1$ 
    and gless:  $\bigwedge t. t \in \{0..1\} \implies \text{norm}(g\ t - p\ t) < \text{min}\ 1\ d$ 
    using path_approx_polynomial_function [OF  $\langle \text{path}\ p \rangle$   $\langle d > 0 \rangle$   $\langle 0 < e \rangle$ ]
    unfolding pathfinish_def pathstart_def by blast
  have  $\text{winding\_number}(\text{exp} \circ p)\ 0 = \text{winding\_number}(\text{exp} \circ g)\ 0$ 
  proof (rule winding_number_nearby_paths_eq [symmetric])
    show  $\text{path}(\text{exp} \circ p)\ \text{path}(\text{exp} \circ g)$ 
    by (simp_all add: pfg assms holomorphic_on_exp holomorphic_on_imp_continuous_on path_continuous_image path_polynomial_function)
  next
    fix  $t::\text{real}$ 
    assume  $t \in \{0..1\}$ 
    with gless have  $\text{norm}(g\ t - p\ t) < 1$ 
      using min_less_iff_conj by blast
    moreover have  $\text{pt}B: \text{norm}(p\ t) \leq B$ 
      using  $B\ t$  by (force simp: path_image_def)

```

```

ultimately have  $cmod (g t) \leq B + 1$ 
  by (meson add_mono_thms_linordered_field(4) le_less_trans less_imp_le
norm_triangle_sub)
with ptB gless t have  $cmod ((exp \circ g) t - (exp \circ p) t) < e$ 
  by (auto simp: dist_norm d)
with e t show  $cmod ((exp \circ g) t - (exp \circ p) t) < cmod ((exp \circ p) t - 0)$ 
  by fastforce
qed (use ⟨g 0 = p 0⟩ ⟨g 1 = p 1⟩ in ⟨auto simp: pathfinish_def pathstart_def⟩)
also have ... =  $1 / (of\_real (2 * pi) * i) * contour\_integral (exp \circ g) (\lambda w. 1 / (w - 0))$ 
  proof (rule winding_number_valid_path)
    have continuous_on (path_image g) (deriv exp)
      by (metis DERIV_exp DERIV_imp_deriv continuous_on_cong holomorphic_on_exp holomorphic_on_imp_continuous_on)
    then show valid_path (exp \circ g)
      by (simp add: field_differentiable_within_exp pfg valid_path_compose valid_path_polynomial_function)
    show  $0 \notin path\_image (exp \circ g)$ 
      by (auto simp: path_image_def)
  qed
also have ... =  $1 / (of\_real (2 * pi) * i) * integral \{0..1\} (\lambda x. vector\_derivative g (at x))$ 
  proof (simp add: contour_integral_integral, rule integral_cong)
    fix t :: real
    assume t:  $t \in \{0..1\}$ 
    show  $vector\_derivative (exp \circ g) (at t) / exp (g t) = vector\_derivative g (at t)$ 
      proof -
        have (exp \circ g has_vector_derivative vector_derivative (exp \circ g) (at t)) (at t)
          by (meson DERIV_exp differentiable_def field_vector_diff_chain_at has_vector_derivative_def
has_vector_derivative_polynomial_function pfg vector_derivative_works)
        moreover have (exp \circ g has_vector_derivative vector_derivative g (at t) * exp (g t)) (at t)
          by (metis DERIV_exp field_vector_diff_chain_at has_vector_derivative_polynomial_function pfg vector_derivative_at)
        ultimately show ?thesis
          by (simp add: divide_simps, rule vector_derivative_unique_at)
      qed
  qed
also have ... = (pathfinish p - pathstart p) / (2 * of_real pi * i)
  proof -
    have (( $\lambda x. vector\_derivative g (at x)$ ) has_integral g 1 - g 0) {0..1}
      by (meson differentiable_at_polynomial_function fundamental_theorem_of_calculus
has_vector_derivative_at_within pfg vector_derivative_works
zero_le_one)
    then show ?thesis
      unfolding pathfinish_def pathstart_def
      using ⟨g 0 = p 0⟩ ⟨g 1 = p 1⟩ by auto
  qed

```

finally show ?thesis .

qed

end

4 Cauchy's Integral Formula

theory Cauchy_Integral_Formula

imports Winding_Numbers

begin

4.1 Proof

lemma Cauchy_integral_formula_weak:

assumes S : convex S and finite k and conf: continuous_on S f

and fcd: $(\bigwedge x. x \in \text{interior } S - k \implies f \text{ field_differentiable at } x)$

and z : $z \in \text{interior } S - k$ and vpg: valid_path γ

and pasz: path_image $\gamma \subseteq S - \{z\}$ and loop: pathfinish $\gamma = \text{pathstart } \gamma$

shows $((\lambda w. f w / (w - z)) \text{ has_contour_integral } (2 * \pi * i * \text{winding_number } \gamma z * f z)) \gamma$

proof -

let ?fz = $\lambda w. (f w - f z) / (w - z)$

obtain f' where f' : $(f \text{ has_field_derivative } f') (at z)$

using fcd [OF z] by (auto simp: field_differentiable_def)

have pas: path_image $\gamma \subseteq S$ and znotin: $z \notin \text{path_image } \gamma$ using pasz by blast+

have c : continuous (at x within S) $(\lambda w. \text{if } w = z \text{ then } f' \text{ else } (f w - f z) / (w - z))$

if $x \in S$ for x

proof (cases $x = z$)

case True then show ?thesis

using LIM_equal [of z ?fz $\lambda w. \text{if } w = z \text{ then } f' \text{ else } ?fz w$] has_field_derivativeD [OF f']

by (force simp add: continuous_within Lim_at_imp_Lim_at_within)

next

case False

then have dxz: dist $x z > 0$ by auto

have cf : continuous (at x within S) f

using conf continuous_on_eq_continuous_within that by blast

have continuous (at x within S) $(\lambda w. (f w - f z) / (w - z))$

by (rule cf continuous_intros | simp add: False)+

then show ?thesis

using continuous_transform_within [OF dxz that] by (force simp: dist_commute)

qed

have fink': finite (insert $z k$) using ⟨finite k ⟩ by blast

have *: $((\lambda w. \text{if } w = z \text{ then } f' \text{ else } ?fz w) \text{ has_contour_integral } 0) \gamma$

proof (rule Cauchy_theorem_convex [OF S fink' vpg pas loop])

show $(\lambda w. \text{if } w = z \text{ then } f' \text{ else } ?fz w) \text{ field_differentiable at } w$

if $w \in \text{interior } S - \text{insert } z k$ for w

proof (rule field_differentiable_transform_within)


```

    show ( $\lambda w. ?fz w$ ) field_differentiable at w
      using that by (intro derivative_intros fcd; simp)
    qed (use that in <auto simp add: dist_pos_lt dist_commute>)
    qed (use c in <force simp: continuous_on_eq continuous_within>)
    note ** = has_contour_integral_add [OF has_contour_integral_lmul [OF has_contour_integral_winding_
[OF vpg znotin], of f z] *]
    show ?thesis
      apply (rule has_contour_integral_eq)
      using znotin ** apply (auto simp: ac_simps divide_simps)
    done
  qed

```

```

theorem Cauchy_integral_formula_convex_simple:
  assumes convex S and holf: f holomorphic_on S and z  $\in$  interior S valid_path
   $\gamma$  path_image  $\gamma \subseteq S - \{z\}$ 
    pathfinish  $\gamma =$  pathstart  $\gamma$ 
  shows (( $\lambda w. f w / (w - z)$ ) has_contour_integral (2 * pi * i * winding_number
 $\gamma z * f z$ ))  $\gamma$ 
proof -
  have  $\bigwedge x. x \in$  interior S  $\implies$  f field_differentiable at x
    using holf at_within_interior holomorphic_onD interior_subset by fastforce
  then show ?thesis
    using assms
    by (intro Cauchy_integral_formula_weak [where k = {}]) (auto simp: holo-
morphic_on_imp_continuous_on)
  qed

```

Hence the Cauchy formula for points inside a circle.

```

theorem Cauchy_integral_circlepath:
  assumes contf: continuous_on (cball z r) f and holf: f holomorphic_on (ball z
r) and wz: norm(w - z) < r
  shows (( $\lambda u. f u / (u - w)$ ) has_contour_integral (2 * of_real pi * i * f w))
    (circlepath z r)
proof -
  have r > 0
    using assms le_less_trans norm_ge_zero by blast
  have (( $\lambda u. f u / (u - w)$ ) has_contour_integral (2 * pi) * i * winding_number
(circlepath z r) w * f w)
    (circlepath z r)
  proof (rule Cauchy_integral_formula_weak [where S = cball z r and k = {}])
    show  $\bigwedge x. x \in$  interior (cball z r) - {w}  $\implies$ 
      f field_differentiable at x
      using holf holomorphic_on_imp_differentiable_at by auto
    have w  $\notin$  sphere z r
      by simp (metis dist_commute dist_norm not_le order_refl wz)
    then show path_image (circlepath z r)  $\subseteq$  cball z r - {w}
      using <r > 0> by (auto simp add: cball_def sphere_def)
  qed (use wz in <simp_all add: dist_norm norm_minus_commute contf>)
  then show ?thesis

```

by (simp add: winding_number_circlepath assms)
qed

corollary *Cauchy_integral_circlepath_simple*:

assumes f holomorphic_on cball z r $\text{norm}(w-z) < r$
shows $((\lambda u. f u / (u-w))$ has_contour_integral $(2 * \text{of_real } \pi * i * f w)$
(circlepath z r)

using *assms* by (force simp: holomorphic_on_imp_continuous_on holomorphic_on_subset
Cauchy_integral_circlepath)

4.2 General stepping result for derivative formulas

lemma *Cauchy_next_derivative*:

fixes $f' :: \text{complex} \Rightarrow \text{complex}$

defines $h \equiv \lambda k w u. f' u / (u-w) \hat{\sim} k$

assumes continuous_on (path_image γ) f'

and $leB: \bigwedge t. t \in \{0..1\} \implies \text{norm}(\text{vector_derivative } \gamma \text{ (at } t)) \leq B$

and $int: \bigwedge w. w \in S - \text{path_image } \gamma \implies (h k w \text{ has_contour_integral } f w) \gamma$

and $k: k \neq 0$

and open S

and $\gamma: \text{valid_path } \gamma$

and $w: w \in S - \text{path_image } \gamma$

shows h (Suc k) w contour_integrable_on γ

and $(f \text{ has_field_derivative } (k * \text{contour_integral } \gamma \text{ (} h \text{ (Suc } k) \text{ } w))) \text{ (at } w)$

(is *?thes2*)

proof –

have open $(S - \text{path_image } \gamma)$ using $\langle \text{open } S \rangle$ closed_valid_path_image γ by
blast

then obtain d where $d > 0$ and $d: \text{ball } w \ d \subseteq S - \text{path_image } \gamma$ using w

using open_contains_ball by *blast*

have [simp]: $\bigwedge n. \text{cmod}(1 + \text{of_nat } n) = 1 + \text{of_nat } n$

by (metis norm_of_nat_of_nat_Suc)

have $cint: (\lambda z. (h k x z - h k w z) / (x * k - w * k))$ contour_integrable_on γ

if $x \neq w$ $\text{cmod}(x-w) < d$ for $x::\text{complex}$

proof –

have $x \in S - \text{path_image } \gamma$

by (metis d dist_commute dist_norm mem_ball subsetD that(2))

then show *?thesis*

using contour_integrable_diff contour_integrable_div contour_integrable_on_def

int w

by *meson*

qed

then have $1: \forall_F x$ in at $w. (\lambda z. (h k x z - h k w z) / (x-w) / \text{of_nat } k)$
contour_integrable_on γ

unfolding eventually_at

by (force intro: *exI* [where $x=d$] simp add: $\langle d > 0 \rangle$ dist_norm field_simps)

have *bim_g*: bounded (image f' (path_image γ))

by (simp add: compact_imp_bounded compact_continuous_image compact_valid_path_image
assms)

```

then obtain  $C$  where  $C > 0$  and  $C: \bigwedge x. [0 \leq x; x \leq 1] \implies \text{cmod } (f' (\gamma x)) \leq C$ 
by (force simp: bounded_pos path_image_def)
have twom:  $\forall_F n$  in at w.
   $\forall x \in \text{path\_image } \gamma.$ 
   $\text{cmod } ((\text{inverse } (x-n) \wedge k - \text{inverse } (x-w) \wedge k) / (n-w) / k -$ 
inverse  $(x-w) \wedge \text{Suc } k) < e$ 
  if  $0 < e$  for  $e$ 
proof -
  have *:  $\text{cmod } ((\text{inverse } (x-u) \wedge k - \text{inverse } (x-w) \wedge k) / ((u-w) * k) -$ 
inverse  $(x-w) \wedge \text{Suc } k) < e$ 
  if  $x \in \text{path\_image } \gamma$  and  $u \neq w$  and uwd:  $\text{cmod } (u-w) < d/2$ 
  and uw_less:  $\text{cmod } (u-w) < e * (d/2) \wedge (k+2) / (1 + \text{real } k)$ 
  for  $u x$ 
proof -
define ff where [abs_def]:
  ff  $n w =$ 
  (if  $n = 0$  then  $\text{inverse}(x-w) \wedge k$ 
  else if  $n = 1$  then  $k / (x-w) \wedge (\text{Suc } k)$ 
  else  $(k * \text{of\_real}(\text{Suc } k)) / (x-w) \wedge (k + 2))$  for  $n :: \text{nat}$  and  $w$ 
have km1:  $\bigwedge z :: \text{complex}. z \neq 0 \implies z \wedge (k - \text{Suc } 0) = z \wedge k / z$ 
by (simp add: field_simps) (metis Suc_pred <k ≠ 0> neq0_conv power_Suc)
have ff1: (ff  $i$  has field_derivative ff  $(\text{Suc } i) z$ ) (at  $z$  within ball  $w (d/2)$ )
  if  $z \in \text{ball } w (d/2)$   $i \leq 1$  for  $i z$ 
proof -
have  $z \notin \text{path\_image } \gamma$ 
  using  $\langle x \in \text{path\_image } \gamma \rangle d$  that ball_divide_subset_numeral by blast
then have xz[simp]:  $x \neq z$  using  $\langle x \in \text{path\_image } \gamma \rangle$  by blast
then have neq:  $x * x + z * z \neq x * (z * 2)$ 
  by (blast intro: dest!: sum_sqs_eq)
with xz have  $\bigwedge v. v \neq 0 \implies (x * x + z * z) * v \neq (x * (z * 2)) * v$  by
auto
then have neqq:  $\bigwedge v. v \neq 0 \implies x * (x * v) + z * (z * v) \neq x * (z * (2 * v))$ 
  by (simp add: algebra_simps)
show ?thesis using  $\langle i \leq 1 \rangle$ 
  apply (simp add: ff_def dist_norm Nat.le_Suc_eq, safe)
  apply (rule derivative_eq_intros | simp add: km1 | simp add: field_simps)
neq neqq+)
  done
qed
{ fix  $a :: \text{real}$  and  $b :: \text{real}$  assume  $ab: a > 0 b > 0$ 
  then have  $k * (1 + \text{real } k) * (1 / a) \leq k * (1 + \text{real } k) * (4 / b) \iff b \leq 4 * a$ 
  by (subst mult_le_cancel_left_pos)
  (use <k ≠ 0> in <auto simp: divide_simps>)
  with ab have  $\text{real } k * (1 + \text{real } k) / a \leq (\text{real } k * 4 + \text{real } k * \text{real } k * 4) / b \iff b \leq 4 * a$ 
  by (simp add: field_simps)

```

```

} note canc = this
have ff2:  $\text{cmod } (\text{ff } (\text{Suc } 1) v) \leq \text{real } (k * (k + 1)) / (d/2) ^ (k + 2)$ 
  if  $v \in \text{ball } w (d/2)$  for  $v$ 
proof -
  have lessd:  $\bigwedge z. \text{cmod } (\gamma z - v) < d/2 \implies \text{cmod } (w - \gamma z) < d$ 
    by (metis that norm_minus_commute norm_triangle_half_r dist_norm
mem_ball)
  have  $d/2 \leq \text{cmod } (x-v)$  using  $d x$  that
    using lessd  $d x$  unfolding path_image_def
    by (smt (verit, best) dist_norm imageE insert_Diff mem_ball subset_Diff_insert)
  then have  $d \leq \text{cmod } (x-v) * 2$ 
    by (simp add: field_split_simps)
  then have dpow_le:  $d ^ (k+2) \leq (\text{cmod } (x-v) * 2) ^ (k+2)$ 
    using  $\langle 0 < d \rangle$  order_less_imp_le power_mono by blast
  have  $x \neq v$  using that
    using  $\langle x \in \text{path\_image } \gamma \rangle$  ball_divide_subset_numeral  $d$  by fastforce
  then show ?thesis
    using  $\langle d > 0 \rangle$  apply (simp add: ff_def norm_mult norm_divide norm_power
dist_norm canc)
    using dpow_le apply (simp add: field_split_simps)
  done
qed
have ub:  $u \in \text{ball } w (d/2)$ 
  using uwd by (simp add: dist_commute dist_norm)
  have  $\text{cmod } (\text{inverse } (x-u) ^ k - (\text{inverse } (x-w) ^ k + \text{of\_nat } k * (u-w) / ((x-w) * (x-w) ^ k)))$ 
     $\leq (\text{real } k * 4 + \text{real } k * \text{real } k * 4) * (\text{cmod } (u-w) * \text{cmod } (u-w))$ 
  /  $(d * (d * (d/2) ^ k))$ 
  using complex_Taylor [OF __ ff1 ff2 __ ub, of  $w$ , simplified]
  by (simp add: ff_def  $\langle 0 < d \rangle$ )
  then have  $\text{cmod } (\text{inverse } (x-u) ^ k - (\text{inverse } (x-w) ^ k + \text{of\_nat } k * (u-w) / ((x-w) * (x-w) ^ k)))$ 
     $\leq (\text{cmod } (u-w) * \text{real } k) * (1 + \text{real } k) * \text{cmod } (u-w) / (d/2) ^ (k+2)$ 
  by (simp add: field_simps)
  then have  $\text{cmod } (\text{inverse } (x-u) ^ k - (\text{inverse } (x-w) ^ k + \text{of\_nat } k * (u-w) / ((x-w) * (x-w) ^ k)))$ 
    /  $(\text{cmod } (u-w) * \text{real } k)$ 
     $\leq (1 + \text{real } k) * \text{cmod } (u-w) / (d/2) ^ (k+2)$ 
  using  $\langle k \neq 0 \rangle$   $\langle u \neq w \rangle$  by (simp add: mult_ac zero_less_mult_iff
pos_divide_le_eq)
  also have  $\dots < e$ 
    using uw_less  $\langle 0 < d \rangle$  by (simp add: mult_ac divide_simps)
  finally have  $e: \text{cmod } (\text{inverse } (x-u) ^ k - (\text{inverse } (x-w) ^ k + \text{of\_nat } k * (u-w) / ((x-w) * (x-w) ^ k)))$ 
    /  $\text{cmod } ((u-w) * \text{real } k) < e$ 
  by (simp add: norm_mult)
have  $x \neq u$ 

```

```

using uwd ⟨0 < d⟩ x d by (force simp: dist_norm ball_def norm_minus_commute)
show ?thesis
  using ⟨k ≠ 0⟩ ⟨x ≠ w⟩ ⟨u ≠ w⟩ le_less_trans [OF _ e]
  by (simp add: field_simps flip: norm_divide)
qed
show ?thesis
  unfolding eventually_at
  apply (rule_tac x = min (d/2) ((e*(d/2)^(k+2))/(Suc k)) in exI)
  apply (force simp: ⟨d > 0⟩ dist_norm that simp del: power_Suc intro: *)
  done
qed
have 2: uniform_limit (path_image γ) (λx z. (h k x z - h k w z) / (x-w) / k)
(h (Suc k) w) (at w)
  unfolding uniform_limit_iff dist_norm
  proof clarify
    fix e::real
    assume 0 < e
    have *: cmod ((h k x (γ u) - h k w (γ u)) / ((x-w) * k) - h (Suc k) w (γ
u)) < e
      if ec: cmod ((inverse (γ u - x) ^ k - inverse (γ u - w) ^ k) / ((x-w) *
k) -
          inverse (γ u - w) * inverse (γ u - w) ^ k) < e / C
        and x: 0 ≤ u u ≤ 1
      for x u
    proof (cases (f' (γ u)) = 0)
      case True then show ?thesis by (simp add: ⟨0 < e⟩ h_def)
    next
      case False
      have cmod ((h k x (γ u) - h k w (γ u)) / ((x-w) * k) - h (Suc k) w (γ u))
=
          cmod (f' (γ u) * ((inverse (γ u - x) ^ k - inverse (γ u - w) ^ k) /
((x-w) * k) -
              inverse (γ u - w) * inverse (γ u - w) ^ k))
        by (simp add: h_def field_simps)
      also have ... = cmod (f' (γ u)) *
          cmod ((inverse (γ u - x) ^ k - inverse (γ u - w) ^ k) / ((x-w)
* k) -
              inverse (γ u - w) * inverse (γ u - w) ^ k)
        by (simp add: norm_mult)
      also have ... < cmod (f' (γ u)) * (e/C)
        using False mult_strict_left_mono [OF ec] by force
      also have ... ≤ e using C
        by (metis False ⟨0 < e⟩ frac_le less_eq_real_def mult.commute pos_le_divide_eq
x zero_less_norm_iff)
    finally show ?thesis .
  qed
show ∀_F u in at w.
  ∀x∈path_image γ.
  cmod ((h k u x - h k w x) / (u-w) / of_nat k - h (Suc k) w x) < e

```

```

    using twom [OF divide_pos_pos [OF ‹0 < e› ‹C > 0›]] *
    unfolding path_image_def h_def
    by (force elim: eventually_mono)
qed
show h (Suc k) w contour_integrable_on γ
  using contour_integral_uniform_limit [OF 1 2 leB γ] by (simp add: h_def)
have *: (λu. contour_integral γ (λx. (h k u x - h k w x) / (u-w) / k))
  -w → contour_integral γ (h (Suc k) w)
  by (rule contour_integral_uniform_limit [OF 1 2 leB γ]) auto
have **: contour_integral γ (λx. (h k u x - h k w x) / ((u-w) * k)) =
  (f u - f w) / (u-w) / k
  if dist u w < d for u
proof -
  have u ∈ S - path_image γ
    by (metis subsetD d dist_commute mem_ball that)
  then have (h k u has_contour_integral f u) γ (h k w has_contour_integral f
w) γ
    using w by (simp_all add: field_simps int)
  then show ?thesis
    using contour_integral_unique has_contour_integral_diff
    has_contour_integral_div by force
qed
show ?thes2
  unfolding has_field_derivative_iff
  by (simp add: ‹k ≠ 0› ** Lim_transform_within [OF tendsto_mult_left [OF
*] ‹0 < d›])
qed

lemma Cauchy_next_derivative_circlepath:
  assumes contf: continuous_on (path_image (circlepath z r)) f
  and int: ∧w. w ∈ ball z r ⇒ ((λu. f u / (u-w) ^ k) has_contour_integral g
w) (circlepath z r)
  and k: k ≠ 0
  and w: w ∈ ball z r
  shows (λu. f u / (u-w) ^ (Suc k)) contour_integrable_on (circlepath z r)
    (is ?thes1)
  and (g has_field_derivative (k * contour_integral (circlepath z r) (λu. f
u / (u-w) ^ (Suc k)))) (at w)
    (is ?thes2)
proof -
  have r > 0 using w
  using ball_eq_empty by fastforce
  have wim: w ∈ ball z r - path_image (circlepath z r)
  using w by (auto simp: dist_norm)
  show ?thes1 ?thes2
  by (rule Cauchy_next_derivative [OF contf __ int k open_ball valid_path_circlepath
wim, where B = 2 * pi * |r|];
  auto simp: vector_derivative_circlepath norm_mult)+
qed

```

In particular, the first derivative formula.

lemma *Cauchy_derivative_integral_circlepath:*

assumes *contf*: *continuous_on* (*cball* *z* *r*) *f*
and *holf*: *f* *holomorphic_on* *ball* *z* *r*
and *w*: $w \in \text{ball } z \ r$
shows $(\lambda u. f \ u / (u-w)^2)$ *contour_integrable_on* (*circlepath* *z* *r*)
(is ?thes1)
and (*f* *has_field_derivative* ($1 / (2 * \text{of_real } \pi * i) * \text{contour_integral}(\text{circlepath } z \ r) (\lambda u. f \ u / (u-w)^2)$)) (*at* *w*)
(is ?thes2)

proof –

have [*simp*]: $r \geq 0$ **using** *w*
using *ball_eq_empty* **by** *fastforce*
have *f*: *continuous_on* (*path_image* (*circlepath* *z* *r*)) *f*
by (*rule* *continuous_on_subset* [*OF contf*]) (*force simp: cball_def sphere_def*)
have *int*: $\bigwedge w. \text{dist } z \ w < r \implies$
 $(\lambda u. f \ u / (u-w)) \text{ has_contour_integral } (\lambda x. 2 * \text{of_real } \pi * i * f$
 $x) \ w) (\text{circlepath } z \ r)$
by (*rule* *Cauchy_integral_circlepath* [*OF contf holf*]) (*simp add: dist_norm*
norm_minus_commute)
show ?*thes1*
unfolding *power2_eq_square*
using *int_Cauchy_next_derivative_circlepath* [*OF f* __ *w*, **where** $k=1$]
by *fastforce*
have $(\lambda x. 2 * \text{of_real } \pi * i * f \ x) \text{ has_field_derivative } \text{contour_integral}$
 $(\text{circlepath } z \ r) (\lambda u. f \ u / (u-w)^2)$ (*at* *w*)
unfolding *power2_eq_square*
using *int_Cauchy_next_derivative_circlepath* [*OF f* __ *w*, **where** $k=1$ **and**
 $g = \lambda x. 2 * \text{of_real } \pi * i * f \ x]$
by *fastforce*
then have *fder*: (*f* *has_field_derivative* *contour_integral* (*circlepath* *z* *r*) $(\lambda u. f$
 $u / (u-w)^2) / (2 * \text{of_real } \pi * i)$) (*at* *w*)
by (*rule* *DERIV_cdivide* [**where** $f = \lambda x. 2 * \text{of_real } \pi * i * f \ x$ **and** $c = 2 * \text{of_real } \pi * i$, *simplified*])
show ?*thes2*
by *simp* (*rule fder*)
qed

4.3 Existence of all higher derivatives

proposition *derivative_is_holomorphic:*

assumes *open* *S*
and *fder*: $\bigwedge z. z \in S \implies (f \text{ has_field_derivative } f' \ z) (\text{at } z)$
shows *f'* *holomorphic_on* *S*

proof –

have *: $\exists h. (f' \text{ has_field_derivative } h) (\text{at } z)$ **if** $z \in S$ **for** *z*

proof –

obtain *r* **where** $r > 0$ **and** *r*: $\text{cball } z \ r \subseteq S$

using *open_contains_cball* $\langle z \in S \rangle \langle \text{open } S \rangle$ **by** *blast*

```

then have holf_cball:  $f$  holomorphic_on cball  $z$   $r$ 
  unfolding holomorphic_on_def
  using field_differentiable_at_within field_differentiable_def fder by fastforce
then have continuous_on (path_image (circlepath  $z$   $r$ ))  $f$ 
  using  $\langle r > 0 \rangle$  by (force elim: holomorphic_on_subset [THEN holomorphic_on_imp_continuous_on])
  then have contfpi: continuous_on (path_image (circlepath  $z$   $r$ ))  $(\lambda x. 1/(2 * of\_real\ pi * i) * f\ x)$ 
    by (auto intro: continuous_intros)+
  have contf_cball: continuous_on (cball  $z$   $r$ )  $f$  using holf_cball
  by (simp add: holomorphic_on_imp_continuous_on holomorphic_on_subset)
  have holf_ball:  $f$  holomorphic_on ball  $z$   $r$  using holf_cball
  using ball_subset_cball holomorphic_on_subset by blast
  { fix  $w$  assume  $w \in$  ball  $z$   $r$ 
    have intf:  $(\lambda u. f\ u / (u-w)^2)$  contour_integrable_on circlepath  $z$   $r$ 
      by (blast intro:  $w$  Cauchy_derivative_integral_circlepath [OF contf_cball holf_ball])
    have fder':  $(f$  has_field_derivative  $1 / (2 * of\_real\ pi * i) * contour\_integral$ 
      (circlepath  $z$   $r$ )  $(\lambda u. f\ u / (u-w)^2)$ )
      (at  $w$ )
      by (blast intro:  $w$  Cauchy_derivative_integral_circlepath [OF contf_cball holf_ball])
    have f'_eq:  $f'$   $w = contour\_integral$  (circlepath  $z$   $r$ )  $(\lambda u. f\ u / (u-w)^2) / (2 * of\_real\ pi * i)$ 
      using fder' ball_subset_cball  $r$   $w$  by (force intro: DERIV_unique [OF fder])
    have  $((\lambda u. f\ u / (u-w)^2) / (2 * of\_real\ pi * i))$  has_contour_integral
      contour_integral (circlepath  $z$   $r$ )  $(\lambda u. f\ u / (u-w)^2) / (2 * of\_real\ pi * i)$ 
      (circlepath  $z$   $r$ )
      by (rule has_contour_integral_div [OF has_contour_integral_integral [OF intf]])
    then have  $((\lambda u. f\ u / (2 * of\_real\ pi * i * (u-w)^2))$  has_contour_integral
      contour_integral (circlepath  $z$   $r$ )  $(\lambda u. f\ u / (u-w)^2) / (2 * of\_real\ pi * i)$ 
      (circlepath  $z$   $r$ )
      by (simp add: algebra_simps)
    then have  $((\lambda u. f\ u / (2 * of\_real\ pi * i * (u-w)^2))$  has_contour_integral
       $f'$   $w$ ) (circlepath  $z$   $r$ )
      by (simp add: f'_eq)
    } note  $*$  = this
  show ?thesis
    using Cauchy_next_derivative_circlepath [OF contfpi, of 2  $f'$ ]  $\langle 0 < r \rangle$  *
    using centre_in_ball mem_ball by force
  qed
show ?thesis
  by (simp add: holomorphic_on_open [OF  $\langle$ open  $S$  $\rangle$ ]) *
qed

```

lemma holomorphic_deriv [holomorphic_intros]:

$\llbracket f \text{ holomorphic_on } S; \text{ open } S \rrbracket \implies (\text{deriv } f) \text{ holomorphic_on } S$
by (metis DERIV_deriv_iff_field_differentiable_at_within_open derivative_is_holomorphic holomorphic_on_def)

lemma holomorphic_deriv_compose:

assumes $g: g \text{ holomorphic_on } B$ **and** $f: f \text{ holomorphic_on } A$ **and** $f' A \subseteq B$
open B

shows $(\lambda x. \text{deriv } g (f x)) \text{ holomorphic_on } A$

using holomorphic_on_compose_gen [OF f holomorphic_deriv[OF g]] *assms*

by (auto simp: o_def)

lemma analytic_deriv [analytic_intros]: $f \text{ analytic_on } S \implies (\text{deriv } f) \text{ analytic_on } S$

using analytic_on_holomorphic holomorphic_deriv **by** auto

lemma holomorphic_higher_deriv [holomorphic_intros]: $\llbracket f \text{ holomorphic_on } S; \text{ open } S \rrbracket \implies (\text{deriv } \hat{\hat{}} n) f \text{ holomorphic_on } S$

by (induction n) (auto simp: holomorphic_deriv)

lemma analytic_higher_deriv [analytic_intros]: $f \text{ analytic_on } S \implies (\text{deriv } \hat{\hat{}} n) f \text{ analytic_on } S$

unfolding analytic_on_def **using** holomorphic_higher_deriv **by** blast

lemma has_field_derivative_higher_deriv:

$\llbracket f \text{ holomorphic_on } S; \text{ open } S; x \in S \rrbracket$

$\implies ((\text{deriv } \hat{\hat{}} n) f \text{ has_field_derivative } (\text{deriv } \hat{\hat{}} (\text{Suc } n)) f x) (at x)$

using holomorphic_derivI holomorphic_higher_deriv **by** fastforce

lemma higher_deriv_cmult:

assumes $f \text{ holomorphic_on } A$ $x \in A$ *open A*

shows $(\text{deriv } \hat{\hat{}} j) (\lambda x. c * f x) x = c * (\text{deriv } \hat{\hat{}} j) f x$

using *assms*

proof (induction j arbitrary: f x)

case (Suc j f x)

have $\text{deriv } ((\text{deriv } \hat{\hat{}} j) (\lambda x. c * f x)) x = \text{deriv } (\lambda x. c * (\text{deriv } \hat{\hat{}} j) f x) x$

using eventually_nhds_in_open[of A x] *assms*(2,3) *Suc.prem*s

by (intro deriv_cong_ev refl) (auto elim!: eventually_mono simp: *Suc.IH*)

also have $\dots = c * \text{deriv } ((\text{deriv } \hat{\hat{}} j) f) x$ **using** *Suc.prem*s *assms*(2,3)

by (intro deriv_cmult holomorphic_on_imp_differentiable_at holomorphic_higher_deriv) *auto*

finally show ?*case* **by** *simp*

qed *simp_all*

lemma higher_deriv_cmult':

assumes $f \text{ analytic_on } \{x\}$

shows $(\text{deriv } \hat{\hat{}} j) (\lambda x. c * f x) x = c * (\text{deriv } \hat{\hat{}} j) f x$

using *assms* higher_deriv_cmult[of f _ x j c] *assms*

using analytic_at_two **by** blast

```

lemma deriv_cmult':
  assumes f analytic_on {x}
  shows deriv ( $\lambda x. c * f x$ ) x = c * deriv f x
  using higher_deriv_cmult'[OF assms, of 1 c] by simp

lemma analytic_derivI:
  assumes f analytic_on {z}
  shows (f has_field_derivative (deriv f z)) (at z within A)
  using assms holomorphic_derivI[of f _ z] analytic_at by blast

lemma deriv_compose_analytic:
  fixes f g :: complex  $\Rightarrow$  complex
  assumes f analytic_on {g z} g analytic_on {z}
  shows deriv ( $\lambda x. f (g x)$ ) z = deriv f (g z) * deriv g z
proof -
  have ((f  $\circ$  g) has_field_derivative (deriv f (g z) * deriv g z)) (at z)
    by (intro DERIV_chain analytic_derivI assms)
  thus ?thesis
    by (auto intro!: DERIV_imp_deriv simp add: o_def)
qed

lemma valid_path_compose_holomorphic:
  assumes valid_path g f holomorphic_on S and open S path_image g  $\subseteq$  S
  shows valid_path (f  $\circ$  g)
  by (meson assms holomorphic_deriv holomorphic_on_imp_continuous_on holomorphic_on_imp_differentiable_at
    holomorphic_on_subset subsetD valid_path_compose)

lemma valid_path_compose_analytic:
  assumes valid_path g and holo:f analytic_on S and path_image g  $\subseteq$  S
  shows valid_path (f  $\circ$  g)
proof (rule valid_path_compose[OF  $\langle$ valid_path g $\rangle$ ])
  fix x assume x  $\in$  path_image g
  then show f field_differentiable at x
    using analytic_on_imp_differentiable_at analytic_on_open assms holo by
blast
next
  show continuous_on (path_image g) (deriv f)
    by (intro holomorphic_on_imp_continuous_on analytic_imp_holomorphic analytic_intros
    analytic_on_subset[OF holo] assms)
qed

lemma analytic_on_deriv [analytic_intros]:
  assumes f analytic_on g ' A
  assumes g analytic_on A
  shows ( $\lambda x. deriv f (g x)$ ) analytic_on A
proof -
  have (deriv f  $\circ$  g) analytic_on A

```

```

  by (rule analytic_on_compose_gen[OF assms(2) analytic_deriv[OF assms(1)]])
  auto
  thus ?thesis
    by (simp add: o_def)
qed

```

```

lemma contour_integral_comp_analyticW:
  assumes f analytic_on s valid_path  $\gamma$  path_image  $\gamma \subseteq s$ 
  shows contour_integral (f  $\circ$   $\gamma$ ) g = contour_integral  $\gamma$  ( $\lambda w$ . deriv f w * g (f w))
proof -
  obtain spikes where finite_spikes and  $\gamma\_diff$ :  $\gamma$  C1_differentiable_on {0..1}
  - spikes
  using <valid_path  $\gamma$ > unfolding valid_path_def piecewise_C1_differentiable_on_def
by auto
  show contour_integral (f  $\circ$   $\gamma$ ) g
    = contour_integral  $\gamma$  ( $\lambda w$ . deriv f w * g (f w))
  unfolding contour_integral_integral
proof (rule integral_spike[rule_format,OF negligible_finite[OF <finite_spikes>]])
  fix t::real assume t:t  $\in$  {0..1} - spikes
  then have  $\gamma$  differentiable_at t
    using  $\gamma\_diff$  unfolding C1_differentiable_on_eq by auto
  moreover have f field_differentiable_at ( $\gamma$  t)
proof -
  have  $\gamma$  t  $\in$  s using t assms unfolding path_image_def by auto
  thus ?thesis
    using <f analytic_on s> analytic_on_imp_differentiable_at by blast
  qed
  ultimately show deriv f ( $\gamma$  t) * g (f ( $\gamma$  t)) * vector_derivative  $\gamma$  (at t) =
    g ((f  $\circ$   $\gamma$ ) t) * vector_derivative (f  $\circ$   $\gamma$ ) (at t)
  by (subst vector_derivative_chain_at_general) (simp_all add:field_simps)
qed
qed

```

4.4 Morera's theorem

```

lemma Morera_local_triangle_ball:
  assumes  $\bigwedge z. z \in S$ 
     $\implies \exists e a. 0 < e \wedge z \in \text{ball } a \ e \wedge \text{continuous\_on } (\text{ball } a \ e) \ f \wedge$ 
       $(\forall b \ c. \text{closed\_segment } b \ c \subseteq \text{ball } a \ e$ 
         $\longrightarrow \text{contour\_integral } (\text{linepath } a \ b) \ f +$ 
         $\text{contour\_integral } (\text{linepath } b \ c) \ f +$ 
         $\text{contour\_integral } (\text{linepath } c \ a) \ f = 0)$ 
  shows f analytic_on S
proof -
  { fix z assume z  $\in$  S
    with assms obtain e a where
      0 < e and z: z  $\in$  ball a e and contf: continuous_on (ball a e) f
      and 0:  $\bigwedge b \ c. \text{closed\_segment } b \ c \subseteq \text{ball } a \ e$ 
         $\implies \text{contour\_integral } (\text{linepath } a \ b) \ f +$ 

```

```

      contour_integral (linepath b c) f +
      contour_integral (linepath c a) f = 0
    by blast
  have az: dist a z < e using mem_ball z by blast
  have  $\exists e > 0. f \text{ holomorphic\_on ball } z e$ 
  proof (intro exI conjI)
    show f holomorphic_on ball z (e - dist a z)
    proof (rule holomorphic_on_subset)
      show ball z (e - dist a z)  $\subseteq$  ball a e
      by (simp add: dist_commute ball_subset_ball_iff)
      have sub_ball:  $\bigwedge y. \text{dist } a y < e \implies \text{closed\_segment } a y \subseteq \text{ball } a e$ 
      by (meson  $\langle 0 < e \rangle$  centre_in_ball convex_ball convex_contains_segment
      mem_ball)
      show f holomorphic_on ball a e
      using triangle_contour_integrals_starlike_primitive [OF contf_open_ball,
      of a]
      derivative_is_holomorphic[OF open_ball]
      by (force simp add:  $0 < 0 < e$  sub_ball)
    qed
  qed (simp add: az)
}
then show ?thesis
  by (simp add: analytic_on_def)
qed

```

lemma Morera_local_triangle:

```

  assumes  $\bigwedge z. z \in S$ 
     $\implies \exists t. \text{open } t \wedge z \in t \wedge \text{continuous\_on } t f \wedge$ 
       $(\forall a b c. \text{convex\_hull } \{a,b,c\} \subseteq t$ 
         $\implies \text{contour\_integral (linepath } a b) f +$ 
           $\text{contour\_integral (linepath } b c) f +$ 
           $\text{contour\_integral (linepath } c a) f = 0)$ 
      )
  shows f analytic_on S
  proof -
    { fix z assume z  $\in S$ 
      with assms obtain t where
        open t and z: z  $\in t$  and contf: continuous_on t f
        and 0:  $\bigwedge a b c. \text{convex\_hull } \{a,b,c\} \subseteq t$ 
           $\implies \text{contour\_integral (linepath } a b) f +$ 
             $\text{contour\_integral (linepath } b c) f +$ 
             $\text{contour\_integral (linepath } c a) f = 0$ 
        by force
      then obtain e where e > 0 and e: ball z e  $\subseteq t$ 
        using open_contains_ball by blast
      have [simp]: continuous_on (ball z e) f using contf
        using continuous_on_subset e by blast
      have eq0:  $\bigwedge b c. \text{closed\_segment } b c \subseteq \text{ball } z e \implies$ 
         $\text{contour\_integral (linepath } z b) f +$ 
           $\text{contour\_integral (linepath } b c) f +$ 
    
```

```

      contour_integral (linepath c z) f = 0
    by (meson 0 z ⟨0 < e⟩ centre_in_ball closed_segment_subset convex_ball
dual_order.trans e starlike_convex_subset)
    have ∃ e a. 0 < e ∧ z ∈ ball a e ∧ continuous_on (ball a e) f ∧
      (∀ b c. closed_segment b c ⊆ ball a e →
        contour_integral (linepath a b) f + contour_integral (linepath b
c) f + contour_integral (linepath c a) f = 0)
      using ⟨e > 0⟩ eq0 by force
  }
  then show ?thesis
    by (simp add: Morera_local_triangle_ball)
qed

```

proposition *Morera_triangle*:

```

[[continuous_on S f; open S;
  ∧ a b c. convex_hull {a,b,c} ⊆ S
  → contour_integral (linepath a b) f +
    contour_integral (linepath b c) f +
    contour_integral (linepath c a) f = 0]]
⇒ f analytic_on S
using Morera_local_triangle by blast

```

4.5 Combining theorems for higher derivatives including Leibniz rule

lemma *higher_deriv_linear* [simp]:

```

(deriv ~ n) (λw. c*w) = (λz. if n = 0 then c*z else if n = 1 then c else 0)
by (induction n) auto

```

lemma *higher_deriv_const* [simp]: (deriv ~ n) (λw. c) = (λw. if n=0 then c else 0)

by (induction n) auto

lemma *higher_deriv_ident* [simp]:

```

(deriv ~ n) (λw. w) z = (if n = 0 then z else if n = 1 then 1 else 0)

```

proof (induction n)

case (Suc n)

then show ?case by (metis higher_deriv_linear lambda_one)

qed auto

lemma *higher_deriv_id* [simp]:

```

(deriv ~ n) id z = (if n = 0 then z else if n = 1 then 1 else 0)

```

by (simp add: id_def)

lemma *has_complex_derivative_funpow_1*:

```

[[f has_field_derivative 1] (at z); f z = z]] ⇒ (f ~ n has_field_derivative 1)
(at z)

```

proof (induction n)

case 0

```

then show ?case
  by (simp add: id_def)
next
  case (Suc n)
  then show ?case
    by (metis DERIV_chain funpow_Suc_right mult.right_neutral)
qed

lemma higher_deriv_uminus:
  assumes f holomorphic_on S open S and z: z ∈ S
  shows (deriv  $\hat{\hat{}}$  n) (λw. -(f w)) z = - ((deriv  $\hat{\hat{}}$  n) f z)
using z
proof (induction n arbitrary: z)
  case 0 then show ?case by simp
next
  case (Suc n z)
  have *: ((deriv  $\hat{\hat{}}$  n) f has_field_derivative deriv ((deriv  $\hat{\hat{}}$  n) f) z) (at z)
    using Suc.premss assms has_field_derivative_higher_deriv by auto
  have  $\bigwedge x. x \in S \implies - (deriv \hat{\hat{}} n) f x = (deriv \hat{\hat{}} n) (\lambda w. - f w) x$ 
    by (auto simp add: Suc)
  then have ((deriv  $\hat{\hat{}}$  n) (λw. - f w) has_field_derivative - deriv ((deriv  $\hat{\hat{}}$  n)
f) z) (at z)
    using has_field_derivative_transform_within_open [of λw. -((deriv  $\hat{\hat{}}$  n) f
w)]
    using * DERIV_minus Suc.premss ⟨open S⟩ by blast
  then show ?case
    by (simp add: DERIV_imp_deriv)
qed

lemma higher_deriv_add:
  fixes z::complex
  assumes f holomorphic_on S g holomorphic_on S open S and z: z ∈ S
  shows (deriv  $\hat{\hat{}}$  n) (λw. f w + g w) z = (deriv  $\hat{\hat{}}$  n) f z + (deriv  $\hat{\hat{}}$  n) g z
using z
proof (induction n arbitrary: z)
  case 0 then show ?case by simp
next
  case (Suc n z)
  have *: ((deriv  $\hat{\hat{}}$  n) f has_field_derivative deriv ((deriv  $\hat{\hat{}}$  n) f) z) (at z)
    ((deriv  $\hat{\hat{}}$  n) g has_field_derivative deriv ((deriv  $\hat{\hat{}}$  n) g) z) (at z)
    using Suc.premss assms has_field_derivative_higher_deriv by auto
  have  $\bigwedge x. x \in S \implies (deriv \hat{\hat{}} n) f x + (deriv \hat{\hat{}} n) g x = (deriv \hat{\hat{}} n) (\lambda w. f
w + g w) x$ 
    by (auto simp add: Suc)
  then have ((deriv  $\hat{\hat{}}$  n) (λw. f w + g w) has_field_derivative
deriv ((deriv  $\hat{\hat{}}$  n) f) z + deriv ((deriv  $\hat{\hat{}}$  n) g) z) (at z)
    using has_field_derivative_transform_within_open [of λw. (deriv  $\hat{\hat{}}$  n) f w
+ (deriv  $\hat{\hat{}}$  n) g w]
    using * Deriv.field_differentiable_add Suc.premss ⟨open S⟩ by blast

```

```

then show ?case
  by (simp add: DERIV_imp_deriv)
qed

```

```

lemma higher_deriv_diff:
  fixes z::complex
  assumes f holomorphic_on S g holomorphic_on S open S z ∈ S
  shows (deriv  $\hat{\sim}$  n) (λw. f w - g w) z = (deriv  $\hat{\sim}$  n) f z - (deriv  $\hat{\sim}$  n) g z
  unfolding diff_conv_add_uminus_higher_deriv_add
  using asms higher_deriv_add higher_deriv_uminus holomorphic_on_minus
by presburger

```

```

lemma Suc_choose: Suc n choose k = (n choose k) + (if k = 0 then 0 else (n
choose (k-1)))
  by (cases k) simp_all

```

```

lemma higher_deriv_mult:
  fixes z::complex
  assumes f holomorphic_on S g holomorphic_on S open S and z: z ∈ S
  shows (deriv  $\hat{\sim}$  n) (λw. f w * g w) z =
    (∑ i = 0..n. of_nat (n choose i) * (deriv  $\hat{\sim}$  i) f z * (deriv  $\hat{\sim}$  (n-i)) g z)
using z
proof (induction n arbitrary: z)
  case 0 then show ?case by simp
next
  case (Suc n z)
  have *: ∧n. ((deriv  $\hat{\sim}$  n) f has_field_derivative deriv ((deriv  $\hat{\sim}$  n) f) z) (at z)
    ∧n. ((deriv  $\hat{\sim}$  n) g has_field_derivative deriv ((deriv  $\hat{\sim}$  n) g) z) (at z)
  using Suc.premss asms has_field_derivative_higher_deriv by auto
  have sumeq: (∑ i = 0..n.
    of_nat (n choose i) * (deriv ((deriv  $\hat{\sim}$  i) f) z * (deriv  $\hat{\sim}$  (n-i)) g z
  + deriv ((deriv  $\hat{\sim}$  (n-i)) g) z * (deriv  $\hat{\sim}$  i) f z)) =
    g z * deriv ((deriv  $\hat{\sim}$  n) f) z + (∑ i = 0..n. (deriv  $\hat{\sim}$  i) f z * (of_nat
(Suc n choose i) * (deriv  $\hat{\sim}$  (Suc n - i)) g z))
  apply (simp add: Suc_choose algebra_simps sum.distrib)
  apply (subst (4) sum_Suc_reindex)
  apply (auto simp: algebra_simps Suc_diff_le intro: sum.cong)
  done
  have ((deriv  $\hat{\sim}$  n) (λw. f w * g w) has_field_derivative
    (∑ i = 0..Suc n. (Suc n choose i) * (deriv  $\hat{\sim}$  i) f z * (deriv  $\hat{\sim}$  (Suc n -
i)) g z))
    (at z)
  apply (rule has_field_derivative_transform_within_open
    [of λw. (∑ i = 0..n. of_nat (n choose i) * (deriv  $\hat{\sim}$  i) f w * (deriv  $\hat{\sim}$ 
(n-i)) g w) _ _ S])
  apply (simp add: algebra_simps)
  apply (rule derivative_eq_intros | simp)+
  apply (auto intro: DERIV_mult * ⟨open S⟩ Suc.premss Suc.IH [symmetric])
by (metis (no_types, lifting) mult.commute sum.cong sumeq)

```

```

then show ?case
  unfolding funpow.simps o_apply
  by (simp add: DERIV_imp_deriv)
qed

```

lemma *higher_deriv_transform_within_open*:

```

fixes z::complex
assumes f holomorphic_on S g holomorphic_on S open S and z: z ∈ S
  and fg:  $\bigwedge w. w \in S \implies f w = g w$ 
  shows (deriv  $\hat{\sim}$  i) f z = (deriv  $\hat{\sim}$  i) g z
using z
by (induction i arbitrary: z)
  (auto simp: fg intro: complex_derivative_transform_within_open holomorphic_higher_deriv
  assms)

```

lemma *higher_deriv_compose_linear'*:

```

fixes z::complex
assumes f: f holomorphic_on T and S: open S and T: open T and z: z ∈ S
  and fg:  $\bigwedge w. w \in S \implies u*w + c \in T$ 
  shows (deriv  $\hat{\sim}$  n) ( $\lambda w. f (u*w + c)$ ) z =  $u^n * (deriv \hat{\sim} n) f (u*z + c)$ 
using z
proof (induction n arbitrary: z)
  case 0 then show ?case by simp
next
  case (Suc n z)
  have holo0: f holomorphic_on ( $\lambda w. u * w + c$ ) ' S
    by (meson fg f holomorphic_on_subset image_subset_iff)
  have holo2: (deriv  $\hat{\sim}$  n) f holomorphic_on ( $\lambda w. u * w + c$ ) ' S
    by (meson f fg holomorphic_higher_deriv holomorphic_on_subset image_subset_iff
  T)
  have holo3: ( $\lambda z. u^n * (deriv \hat{\sim} n) f (u * z + c)$ ) holomorphic_on S
    by (intro holo2 holomorphic_on_compose [where g=(deriv  $\hat{\sim}$  n) f, unfolded
  o_def] holomorphic_intros)
  have ( $\lambda w. u * w + c$ ) holomorphic_on S f holomorphic_on ( $\lambda w. u * w + c$ ) ' S
    by (rule holo0 holomorphic_intros)+
  then have holo1: ( $\lambda w. f (u * w + c)$ ) holomorphic_on S
    by (rule holomorphic_on_compose [where g=f, unfolded o_def])
  have deriv ((deriv  $\hat{\sim}$  n) ( $\lambda w. f (u * w + c)$ )) z = deriv ( $\lambda z. u^n * (deriv \hat{\sim} n)$ 
  f (u*z+c)) z
proof (rule complex_derivative_transform_within_open [OF _ holo3 S Suc.prem])
  show (deriv  $\hat{\sim}$  n) ( $\lambda w. f (u * w + c)$ ) holomorphic_on S
    by (rule holomorphic_higher_deriv [OF holo1 S])
qed (simp add: Suc.IH)
also have ... =  $u^n * deriv (\lambda z. (deriv \hat{\sim} n) f (u * z + c)) z$ 
proof -
  have (deriv  $\hat{\sim}$  n) f analytic_on T
    by (simp add: analytic_on_open f holomorphic_higher_deriv T)
  then have ( $\lambda w. (deriv \hat{\sim} n) f (u * w + c)$ ) analytic_on S
    using holomorphic_on_compose[OF _ holo2]  $\langle (\lambda w. u * w + c)$  holomor-

```



```

phic_on S
  by (simp add: S analytic_on_open o_def)
  then show ?thesis
    by (intro deriv_cmult analytic_on_imp_differentiable_at [OF _ Suc.prem])
  qed
also have ... = u * u ^ n * deriv ((deriv ^ n) f) (u * z+c)
proof -
  have (deriv ^ n) f field_differentiable at (u * z+c)
  using Suc.prem T ffg holomorphic_higher_deriv holomorphic_on_imp_differentiable_at
by blast
  then show ?thesis
    by (simp add: deriv_compose_linear')
  qed
  finally show ?case
    by simp
qed

```

lemma *higher_deriv_compose_linear*:

```

fixes z::complex
assumes f: f holomorphic_on T and S: open S and T: open T and z: z ∈ S
  and fg:  $\bigwedge w. w \in S \implies u * w \in T$ 
  shows (deriv ^ n) ( $\lambda w. f (u * w)$ ) z = u ^ n * (deriv ^ n) f (u * z)
using higher_deriv_compose_linear' [where c=0] assms by simp

```

lemma *higher_deriv_add_at*:

```

assumes f analytic_on {z} g analytic_on {z}
  shows (deriv ^ n) ( $\lambda w. f w + g w$ ) z = (deriv ^ n) f z + (deriv ^ n) g z
using analytic_at_two assms higher_deriv_add by blast

```

lemma *higher_deriv_diff_at*:

```

assumes f analytic_on {z} g analytic_on {z}
  shows (deriv ^ n) ( $\lambda w. f w - g w$ ) z = (deriv ^ n) f z - (deriv ^ n) g z
using analytic_at_two assms higher_deriv_diff by blast

```

lemma *higher_deriv_uminus_at*:

```

f analytic_on {z}  $\implies$  (deriv ^ n) ( $\lambda w. -(f w)$ ) z = - ((deriv ^ n) f z)
using higher_deriv_uminus by (auto simp: analytic_at)

```

lemma *higher_deriv_mult_at*:

```

assumes f analytic_on {z} g analytic_on {z}
  shows (deriv ^ n) ( $\lambda w. f w * g w$ ) z =
  ( $\sum i = 0..n. of\_nat (n choose i) * (deriv ^ i) f z * (deriv ^ (n-i)) g z$ )
using analytic_at_two assms higher_deriv_mult by blast

```

Nonexistence of isolated singularities and a stronger integral formula.

proposition *no_isolated_singularity*:

```

fixes z::complex
assumes f: continuous_on S f and holf: f holomorphic_on (S-K) and S: open
S and K: finite K

```

```

    shows f holomorphic_on S
  proof -
    { fix z
      assume z ∈ S and cdf:  $\bigwedge x. x \in S - K \implies f \text{ field\_differentiable at } x$ 
      have f field\_differentiable at z
      proof (cases z ∈ K)
        case False then show ?thesis by (blast intro: cdf ⟨z ∈ S⟩)
      next
        case True
          with finite_set_avoid [OF K, of z]
          obtain d where d > 0 and d:  $\bigwedge x. [x \in K; x \neq z] \implies d \leq \text{dist } z \ x$ 
            by blast
          obtain e where e > 0 and e: ball z e ⊆ S
            using S ⟨z ∈ S⟩ by (force simp: open_contains_ball)
          have fde: continuous_on (ball z (min d e)) f
            by (metis Int_iff_ball_min_Int continuous_on_subset e f subsetI)
          have cont:  $\{a, b, c\} \subseteq \text{ball } z \ (\text{min } d \ e) \implies \text{continuous\_on } (\text{convex hull } \{a, b, c\}) \ f$ 
            for a b c
            by (simp add: hull_minimal continuous_on_subset [OF fde])
          have fd:  $[\{a, b, c\} \subseteq \text{ball } z \ (\text{min } d \ e); x \in \text{interior } (\text{convex hull } \{a, b, c\}) - K]$ 
             $\implies f \text{ field\_differentiable at } x$  for a b c x
            by (metis cdf Diff_iff_Int_iff_ball_min_Int subsetD convex_ball e interior_mono interior_subset subset_hull)
          obtain g where  $\bigwedge w. w \in \text{ball } z \ (\text{min } d \ e) \implies (g \text{ has\_field\_derivative } f \ w)$ 
            (at w within ball z (min d e))
            apply (rule contour_integral_convex_primitive
              [OF convex_ball fde Cauchy_theorem_triangle_cofinite [OF _
                K]])
            using cont fd by auto
          then have f holomorphic_on ball z (min d e)
            by (metis open_ball at_within_open derivative_is_holomorphic)
          then show ?thesis
            unfolding holomorphic_on_def
            by (metis open_ball ⟨0 < d⟩ ⟨0 < e⟩ at_within_open centre_in_ball min_less_iff_conj)
        qed
      }
    with holf S K show ?thesis
      by (simp add: holomorphic_on_open open_Diff finite_imp_closed field_differentiable_def
        [symmetric])
    qed
  lemma no_isolated_singularity':
    fixes z::complex
    assumes f:  $\bigwedge z. z \in K \implies (f \longrightarrow f \ z)$  (at z within S)
      and holf: f holomorphic_on (S - K) and S: open S and K: finite K
    shows f holomorphic_on S
  proof (rule no_isolated_singularity[OF _ assms(2-)])

```

```

have continuous_on (S-K) f
  using holf holomorphic_on_imp_continuous_on by auto
then show continuous_on S f
  by (metis Diff_iff K S at_within_open continuous_on_eq_continuous_at
    continuous_within f finite_imp_closed open_Diff)
qed

```

proposition *Cauchy_integral_formula_convex:*

```

assumes S: convex S and K: finite K and contf: continuous_on S f
  and fcd: ( $\bigwedge x. x \in \text{interior } S - K \implies f \text{ field\_differentiable at } x$ )
  and z: z  $\in$  interior S and vpg: valid_path  $\gamma$ 
  and pasz: path_image  $\gamma \subseteq S - \{z\}$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
shows (( $\lambda w. f w / (w-z)$ ) has_contour_integral (2*pi * i * winding_number  $\gamma$ 
  z * f z))  $\gamma$ 
proof -
  have *:  $\bigwedge x. x \in \text{interior } S \implies f \text{ field\_differentiable at } x$ 
    unfolding holomorphic_on_open [symmetric] field_differentiable_def
    using no_isolated_singularity [where S = interior S]
  by (meson K contf continuous_on_subset fcd field_differentiable_def open_interior
    has_field_derivative_at_within holomorphic_derivI holomorphic_onI interior_subset)
  show ?thesis
    by (rule Cauchy_integral_formula_weak [OF S finite.emptyI contf]) (use *
    assms in auto)
qed

```

Formula for higher derivatives.

lemma *Cauchy_has_contour_integral_higher_derivative_circlepath:*

```

assumes contf: continuous_on (cball z r) f
  and holf: f holomorphic_on ball z r
  and w: w  $\in$  ball z r
shows (( $\lambda u. f u / (u-w) \wedge (\text{Suc } k)$ ) has_contour_integral ((2 * pi * i) / (fact
  k) * (deriv  $\sim k$ ) f w))
  (circlepath z r)
using w
proof (induction k arbitrary: w)
  case 0 then show ?case
    using assms by (auto simp: Cauchy_integral_circlepath dist_commute dist_norm)
  next
    case (Suc k)
    have [simp]: r > 0 using w
      using ball_eq_empty by fastforce
    have f: continuous_on (path_image (circlepath z r)) f
      by (rule continuous_on_subset [OF contf]) (force simp: cball_def sphere_def
      less_imp_le)
    obtain X where X: (( $\lambda u. f u / (u-w) \wedge \text{Suc } (\text{Suc } k)$ ) has_contour_integral X)
      (circlepath z r)
      using Cauchy_next_derivative_circlepath(1) [OF f Suc.IH _ Suc.prem]
    by (auto simp: contour_integrable_on_def)

```

```

then have con: contour_integral (circlepath z r) (( $\lambda u. f u / (u-w) \wedge \text{Suc } (\text{Suc } k)$ )) = X
  by (rule contour_integral_unique)
have  $\wedge n. ((\text{deriv } \wedge n) f \text{ has\_field\_derivative } \text{deriv } ((\text{deriv } \wedge n) f) w) (at w)$ 
  using Suc.premss assms has_field_derivative_higher_deriv by auto
then have dnf_diff:  $\wedge n. (\text{deriv } \wedge n) f \text{ field\_differentiable } (at w)$ 
  by (force simp: field_differentiable_def)
have deriv ( $\lambda w. \text{complex\_of\_real } (2 * \text{pi}) * i / (\text{fact } k) * (\text{deriv } \wedge k) f w$ ) w =
  of_nat (Suc k) * contour_integral (circlepath z r) ( $\lambda u. f u / (u-w) \wedge \text{Suc } (\text{Suc } k)$ )
  by (force intro!: DERIV_imp_deriv Cauchy_next_derivative_circlepath [OF f Suc.IH _ Suc.premss])
also have ... = of_nat (Suc k) * X
  by (simp only: con)
finally have deriv ( $\lambda w. ((2 * \text{pi}) * i / (\text{fact } k)) * (\text{deriv } \wedge k) f w$ ) w = of_nat (Suc k) * X .
  then have  $((2 * \text{pi}) * i / (\text{fact } k)) * \text{deriv } (\lambda w. (\text{deriv } \wedge k) f w) w = \text{of\_nat } (\text{Suc } k) * X$ 
  by (metis deriv_cmult dnf_diff)
then have deriv ( $\lambda w. (\text{deriv } \wedge k) f w$ ) w = of_nat (Suc k) * X /  $((2 * \text{pi}) * i / (\text{fact } k))$ 
  by (simp add: field_simps)
then show ?case
using of_nat_eq_0_iff X by fastforce
qed

```

lemma *Cauchy_higher_derivative_integral_circlepath*:

```

assumes contf: continuous_on (cball z r) f
  and holf: f holomorphic_on ball z r
  and w:  $w \in \text{ball } z r$ 
shows  $(\lambda u. f u / (u-w) \wedge (\text{Suc } k)) \text{ contour\_integrable\_on } (\text{circlepath } z r)$ 
  (is ?thes1)
  and  $(\text{deriv } \wedge k) f w = (\text{fact } k) / (2 * \text{pi} * i) * \text{contour\_integral}(\text{circlepath } z r) (\lambda u. f u / (u-w) \wedge (\text{Suc } k))$ 
  (is ?thes2)
proof -
  have *:  $((\lambda u. f u / (u-w) \wedge \text{Suc } k) \text{ has\_contour\_integral } (2 * \text{pi}) * i / (\text{fact } k) * (\text{deriv } \wedge k) f w)$ 
  (circlepath z r)
  using Cauchy_has_contour_integral_higher_derivative_circlepath [OF assms]
  by simp
show ?thes1 using *
  using contour_integrable_on_def by blast
show ?thes2
  unfolding contour_integral_unique [OF *] by (simp add: field_split_simps)
qed

```

corollary *Cauchy_contour_integral_circlepath*:

```

assumes continuous_on (cball z r) f f holomorphic_on ball z r  $w \in \text{ball } z r$ 

```

shows $\text{contour_integral}(\text{circlepath } z \ r) (\lambda u. f \ u / (u-w)^{\wedge}(\text{Suc } k)) = (2 * \text{pi} * i) * (\text{deriv } \widehat{\sim} k) f \ w / (\text{fact } k)$
by (*simp add: Cauchy_higher_derivative_integral_circlepath [OF assms]*)

lemma *Cauchy_contour_integral_circlepath_2:*

assumes *continuous_on (cball z r) f f holomorphic_on ball z r w ∈ ball z r*
shows $\text{contour_integral}(\text{circlepath } z \ r) (\lambda u. f \ u / (u-w)^{\wedge}2) = (2 * \text{pi} * i) * \text{deriv } f \ w$
using *Cauchy_contour_integral_circlepath [OF assms, of 1]*
by (*simp add: power2_eq_square*)

4.6 A holomorphic function is analytic, i.e. has local power series

theorem *holomorphic_power_series:*

assumes *holf: f holomorphic_on ball z r*
and *w: w ∈ ball z r*
shows $((\lambda n. (\text{deriv } \widehat{\sim} n) f \ z / (\text{fact } n) * (w-z)^{\wedge}n) \ \text{sums } f \ w)$

proof –

— Replacing r and the original (weak) premises with stronger ones

obtain r **where** $r > 0$ **and** *holfc: f holomorphic_on cball z r* **and** *w: w ∈ ball z r*

proof

have $\text{cball } z \ ((r + \text{dist } w \ z) / 2) \subseteq \text{ball } z \ r$

using w **by** (*simp add: dist_commute field_sum_of_halves subset_eq*)

then show $f \ \text{holomorphic_on } \text{cball } z \ ((r + \text{dist } w \ z) / 2)$

by (*rule holomorphic_on_subset [OF holf]*)

have $r > 0$

by (*metis w dist_norm mem_ball norm_ge_zero not_less_iff_gr_or_eq order_less_le_trans*)

then show $0 < (r + \text{dist } w \ z) / 2$

by *simp (use zero_le_dist [of w z] in linarith)*

qed (*use w in <auto simp: dist_commute>*)

then have *holfc: f holomorphic_on ball z r*

using *ball_subset_cball holomorphic_on_subset* **by** *blast*

have *contf: continuous_on (cball z r) f*

by (*simp add: holfc holomorphic_on_imp_continuous_on*)

have *cint: $\bigwedge k. (\lambda u. f \ u / (u-z)^{\wedge} \text{Suc } k) \ \text{contour_integrable_on } \text{circlepath } z \ r$*

by (*rule Cauchy_higher_derivative_integral_circlepath [OF contf holfc]*) (*simp add: <0 < r>*)

obtain B **where** $0 < B$ **and** $B: \bigwedge u. u \in \text{cball } z \ r \implies \text{norm}(f \ u) \leq B$

by (*metis (no_types) bounded_pos compact_cball compact_continuous_image compact_imp_bounded contf image_eqI*)

obtain k **where** $0 < k \ k \leq r$ **and** *wz_eq: norm(w-z) = r - k*

and *kle: $\bigwedge u. \text{norm}(u-z) = r \implies k \leq \text{norm}(u-w)$*

proof

show $\bigwedge u. \text{cmod } (u-z) = r \implies r - \text{dist } z \ w \leq \text{cmod } (u-w)$

by (*metis add_diff_eq diff_add_cancel dist_norm norm_diff_ineq*)

qed (*use w in <auto simp: dist_norm norm_minus_commute>*)

```

have ul: uniform_limit (sphere z r) ( $\lambda n x. (\sum k < n. (w-z) ^ k * (f x / (x-z) ^ Suc k))$ ) ( $\lambda x. f x / (x-w)$ ) sequentially
unfolding uniform_limit_iff dist_norm
proof clarify
fix e::real
assume  $0 < e$ 
have rr:  $0 \leq (r-k) / r$   $(r-k) / r < 1$  using k by auto
obtain n where n:  $((r-k) / r) ^ n < e / B * k$ 
using real_arch_pow_inv [of  $e/B*k$   $(r-k)/r$ ]  $\langle 0 < e \rangle$   $\langle 0 < B \rangle$  k by force
have norm  $((\sum k < N. (w-z) ^ k * f u / (u-z) ^ Suc k) - f u / (u-w)) < e$ 
if  $n \leq N$  and r:  $r = dist\ z\ u$  for N u
proof -
have N:  $((r-k) / r) ^ N < e / B * k$ 
using le_less_trans [OF power_decreasing n]
using  $\langle n \leq N \rangle$  k by auto
have u [simp]:  $(u \neq z) \wedge (u \neq w)$ 
using  $\langle 0 < r \rangle$  r w by auto
have wzu_not1:  $(w-z) / (u-z) \neq 1$ 
by (metis (no_types) dist_norm divide_eq_1_iff less_irrefl mem_ball
norm_minus_commute r w)
have norm  $((\sum k < N. (w-z) ^ k * f u / (u-z) ^ Suc k) * (u-w) - f u$ 
 $= norm ((\sum k < N. (((w-z) / (u-z)) ^ k)) * f u * (u-w) / (u-z) - f u$ 
unfolding sum_distrib_right sum_divide_distrib power_divide by (simp
add: algebra_simps)
also have  $\dots = norm (((w-z) / (u-z)) ^ N - 1) * (u-w) / (((w-z) /$ 
 $(u-z) - 1) * (u-z)) - 1) * norm (f u)$ 
using  $\langle 0 < B \rangle$ 
apply (simp add: geometric_sum [OF wzu_not1] flip: norm_mult)
apply (simp add: field_simps)
done
also have  $\dots = norm ((u-z) ^ N * (w-u) - ((w-z) ^ N - (u-z) ^ N) *$ 
 $(u-w)) / (r ^ N * norm (u-w)) * norm (f u)$ 
using  $\langle 0 < r \rangle$  r by (simp add: divide_simps norm_mult norm_divide
norm_power dist_norm norm_minus_commute)
also have  $\dots = norm ((w-z) ^ N * (w-u)) / (r ^ N * norm (u-w)) *$ 
 $norm (f u)$ 
by (simp add: algebra_simps)
also have  $\dots = norm (w-z) ^ N * norm (f u) / r ^ N$ 
by (simp add: norm_mult norm_power norm_minus_commute)
also have  $\dots \leq (((r-k)/r)^N) * B$ 
using  $\langle 0 < r \rangle$  w k
by (simp add: B divide_simps mult_mono r wz_eq)
also have  $\dots < e * k$ 
using  $\langle 0 < B \rangle$  N by (simp add: divide_simps)
also have  $\dots \leq e * norm (u-w)$ 
using r kle  $\langle 0 < e \rangle$  by (simp add: dist_commute dist_norm)
finally show ?thesis
by (simp add: field_split_simps norm_divide del: power_Suc)
qed

```

```

with  $\langle 0 < r \rangle$  show  $\forall_F n$  in sequentially.  $\forall x \in \text{sphere } z \ r.$ 
      norm  $((\sum_{k < n}. (w-z) \wedge k * (f x / (x-z) \wedge \text{Suc } k)) - f x / (x-w))$ 
< e
  by (auto simp: mult_ac less_imp_le eventually_sequentially Ball_def)
qed
have  $\S: \bigwedge x \ k. k \in \{..<x\} \implies$ 
       $(\lambda u. (w-z) \wedge k * (f u / (u-z) \wedge \text{Suc } k))$  contour_integrable_on circlepath
z r
  using contour_integrable_lmul [OF cint, of  $(w-z) \wedge a$  for a] by (simp add:
field_simps)
  have eq:  $\bigwedge n.$ 
       $(\sum_{k < n}. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u-z) \wedge \text{Suc } k) * (w-z) \wedge k) =$ 
       $\text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. \sum_{k < n}. (w-z) \wedge k * (f u / (u-z) \wedge \text{Suc } k))$ 
  apply (subst contour_integral_sum)
  apply (simp_all only:  $\S$  finite_lessThan contour_integral_lmul cint algebra_simps)
done
  have  $\bigwedge u \ k. k \in \{..<u\} \implies (\lambda x. f x / (x-z) \wedge \text{Suc } k)$  contour_integrable_on
circlepath z r
  using  $\langle 0 < r \rangle$  by (force intro!: Cauchy_higher_derivative_integral_circlepath
[OF contf holf])
  then have  $\bigwedge u. (\lambda y. \sum_{k < u}. (w-z) \wedge k * (f y / (y-z) \wedge \text{Suc } k))$  contour_integrable_on
circlepath z r
  by (intro contour_integrable_sum contour_integrable_lmul, simp)
  then have  $(\lambda k. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u-z) \wedge (\text{Suc } k)) * (w-z) \wedge k)$ 
      sums contour_integral (circlepath z r)  $(\lambda u. f u / (u-w))$ 
  unfolding sums_def eq
  using  $\langle 0 < r \rangle$  contour_integral_uniform_limit_circlepath [OF eventuallyI ul]

  by fastforce
  then have  $(\lambda k. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u-z) \wedge (\text{Suc } k)) * (w-z) \wedge k)$ 
      sums  $(2 * \text{of\_real } \pi * i * f w)$ 
  using w by (auto simp: dist_commute dist_norm contour_integral_unique [OF
Cauchy_integral_circlepath_simple [OF holfc]])
  then have  $(\lambda k. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u-z) \wedge \text{Suc } k) * (w-z) \wedge k / (i * (\text{of\_real } \pi * 2)))$ 
      sums  $((2 * \text{of\_real } \pi * i * f w) / (i * (\text{complex\_of\_real } \pi * 2)))$ 
  by (rule sums_divide)
  then have  $(\lambda n. (w-z) \wedge n * \text{contour\_integral } (\text{circlepath } z \ r) (\lambda u. f u / (u-z) \wedge \text{Suc } n) / (i * (\text{of\_real } \pi * 2)))$ 
      sums f w
  by (simp add: field_simps)
then show ?thesis
  by (simp add: field_simps  $\langle 0 < r \rangle$  Cauchy_higher_derivative_integral_circlepath
[OF contf holf])

```

qed

4.7 The Liouville theorem and the Fundamental Theorem of Algebra

These weak Liouville versions don't even need the derivative formula.

lemma *Liouville_weak_0*:
assumes *holf*: *f holomorphic_on UNIV* **and** *inf*: (*f* \longrightarrow 0) *at_infinity*
shows *f z = 0*
proof (*rule ccontr*)
assume *fz*: *f z \neq 0*
with *inf* [*unfolded Lim_at_infinity*, *rule_format*, *of norm(f z)/2*]
obtain *B* **where** *B*: $\bigwedge x. B \leq \text{cmod } x \implies \text{norm } (f x) * 2 < \text{cmod } (f z)$
by (*auto simp: dist_norm*)
define *R* **where** $R = 1 + |B| + \text{norm } z$
have $R > 0$
unfolding *R_def* **by** (*smt (verit) norm_ge_zero*)
have *: $((\lambda u. f u / (u - z)) \text{ has_contour_integral } 2 * \text{complex_of_real } \pi * i * f z)$ (*circlepath z R*)
using *continuous_on_subset holf holomorphic_on_subset* $\langle 0 < R \rangle$
by (*force intro: holomorphic_on_imp_continuous_on Cauchy_integral_circlepath*)
have $\text{cmod } (x - z) = R \implies \text{cmod } (f x) * 2 < \text{cmod } (f z)$ **for** *x*
unfolding *R_def* **by** (*rule B*) (*use norm_triangle_ineq4 [of x z] in auto*)
with $\langle R > 0 \rangle$ *fz* **show** *False*
using *has_contour_integral_bound_circlepath [OF *, of norm(f z)/2/R]*
by (*auto simp: less_imp_le norm_mult norm_divide field_split_simps*)
 qed

proposition *Liouville_weak*:
assumes *f holomorphic_on UNIV* **and** (*f* \longrightarrow *l*) *at_infinity*
shows *f z = l*
using *Liouville_weak_0* [*of* $\lambda z. f z - l$]
by (*simp add: assms holomorphic_on_diff LIM_zero*)

proposition *Liouville_weak_inverse*:
assumes *f holomorphic_on UNIV* **and** *unbounded*: $\bigwedge B. \text{eventually } (\lambda x. \text{norm } (f x) \geq B)$ *at_infinity*
obtains *z* **where** *f z = 0*
proof –
 { **assume** *f*: $\bigwedge z. f z \neq 0$
have 1: $(\lambda x. 1 / f x)$ *holomorphic_on UNIV*
by (*simp add: holomorphic_on_divide assms f*)
have 2: $(\lambda x. 1 / f x) \longrightarrow 0$ *at_infinity*
proof (*rule tendstoI [OF eventually_mono]*)
fix *e*::*real*
assume $e > 0$
show *eventually* $(\lambda x. 2/e \leq \text{cmod } (f x))$ *at_infinity*
by (*rule_tac B=2/e in unbounded*)
 qed (*simp add: dist_norm norm_divide field_split_simps*)


```

    have False
      using Liouville_weak_0 [OF 1 2] f by simp
  }
  then show ?thesis
    using that by blast
qed

```

In particular we get the Fundamental Theorem of Algebra.

```

theorem fundamental_theorem_of_algebra:
  fixes a :: nat  $\Rightarrow$  complex
  assumes a 0 = 0  $\vee$  ( $\exists i \in \{1..n\}. a\ i \neq 0$ )
  obtains z where ( $\sum_{i \leq n}. a\ i * z^i$ ) = 0
using assms
proof (elim disjE bexE)
  assume a 0 = 0 then show ?thesis
    by (auto simp: that [of 0])
next
  fix i
  assume i: i  $\in \{1..n\}$  and nz: a i  $\neq 0$ 
  have 1: (\lambda z. \sum_{i \leq n}. a i * z^i) holomorphic_on UNIV
    by (rule holomorphic_intros)+
  show thesis
proof (rule Liouville_weak_inverse [OF 1])
  show  $\forall_F x$  in at_infinity. B  $\leq$  cmod ( $\sum_{i \leq n}. a\ i * x^i$ ) for B
    using i nz by (intro polyfun_extremal exI [of _ i]) auto
qed (use that in auto)
qed

```

4.8 Weierstrass convergence theorem

```

lemma holomorphic_uniform_limit:
  assumes cont: eventually ( $\lambda n. continuous\_on (cball\ z\ r) (f\ n) \wedge (f\ n)$ ) holomorphic_on ball z r F
  and ulim: uniform_limit (cball z r) f g F
  and F:  $\neg$  trivial_limit F
  obtains continuous_on (cball z r) g g holomorphic_on ball z r
proof (cases r 0::real rule: linorder_cases)
  case less then show ?thesis by (force simp: ball_empty less_imp_le continuous_on_def holomorphic_on_def intro: that)
next
  case equal then show ?thesis
    by (force simp: holomorphic_on_def intro: that)
next
  case greater
  have contg: continuous_on (cball z r) g
    using cont uniform_limit_theorem [OF eventually_mono ulim F] by blast
  have path_image (circlepath z r)  $\subseteq$  cball z r
    using  $\langle 0 < r \rangle$  by auto
  then have 1: continuous_on (path_image (circlepath z r)) ( $\lambda x. 1 / (2 * com-$ 

```

```

plex_of_real pi * i) * g x)
  by (intro continuous_intros continuous_on_subset [OF contg])
  have 2: (( $\lambda u. 1 / (2 * \text{of\_real } \pi * i) * g u / (u-w) ^ 1$ ) has_contour_integral
g w) (circlepath z r)
  if w: w  $\in$  ball z r for w
  proof -
    define d where d = (r - norm(w-z))
    have 0 < d d  $\leq$  r using w by (auto simp: norm_minus_commute d_def
dist_norm)
    have dle:  $\bigwedge u. \text{cmod } (z-u) = r \implies d \leq \text{cmod } (u-w)$ 
    unfolding d_def by (metis add_diff_eq diff_add_cancel norm_diff_ineq
norm_minus_commute)
    have ev_int:  $\forall_F n \text{ in } F. (\lambda u. f n u / (u-w)) \text{ contour\_integrable\_on circlepath }
z r$ 
    using w
    by (auto intro: eventually_mono [OF cont] Cauchy_higher_derivative_integral_circlepath
[where k=0, simplified])
    have  $\bigwedge e. [0 < r; 0 < d; 0 < e]$ 
 $\implies \forall_F n \text{ in } F.$ 
 $\forall x \in \text{sphere } z r.$ 
 $x \neq w \longrightarrow$ 
 $\text{cmod } (f n x - g x) < e * \text{cmod } (x-w)$ 
    apply (rule_tac e1=e * d in eventually_mono [OF uniform_limitD [OF
ulim]])
    apply (force simp: dist_norm intro: dle mult_left_mono less_le_trans)+
    done
    then have ul_less: uniform_limit (sphere z r) ( $\lambda n x. f n x / (x-w)$ ) ( $\lambda x. g x /
(x-w)$ ) F
    using greater <0 < d>
    by (auto simp add: uniform_limit_iff dist_norm norm_divide diff_divide_distrib
[symmetric] divide_simps)
    have g_cint: ( $\lambda u. g u / (u-w)$ ) contour_integrable_on circlepath z r
    by (rule contour_integral_uniform_limit_circlepath [OF ev_int ul_less F <0
< r>])
    have cif_tends_cig: (( $\lambda n. \text{contour\_integral}(\text{circlepath } z r) (\lambda u. f n u / (u-w))$ )
 $\longrightarrow \text{contour\_integral}(\text{circlepath } z r) (\lambda u. g u / (u-w))$ ) F
    by (rule contour_integral_uniform_limit_circlepath [OF ev_int ul_less F <0
< r>])
    have f_tends_cig: (( $\lambda n. 2 * \text{of\_real } \pi * i * f n w$ )  $\longrightarrow \text{contour\_integral}
(\text{circlepath } z r) (\lambda u. g u / (u-w))$ ) F
    proof (rule Lim_transform_eventually)
      show  $\forall_F x \text{ in } F. \text{contour\_integral } (\text{circlepath } z r) (\lambda u. f x u / (u-w))$ 
 $= 2 * \text{of\_real } \pi * i * f x w$ 
      using w <0 < d> d_def
      by (auto intro: eventually_mono [OF cont contour_integral_unique [OF
Cauchy_integral_circlepath]])
    qed (auto simp: cif_tends_cig)
    have  $\bigwedge e. 0 < e \implies \forall_F n \text{ in } F. \text{dist } (f n w) (g w) < e$ 
    by (rule eventually_mono [OF uniform_limitD [OF ulim]]) (use w in auto)

```

```

then have (( $\lambda n. 2 * \text{of\_real } \pi * i * f n w$ )  $\longrightarrow$   $2 * \text{of\_real } \pi * i * g w$ )  $F$ 
  by (rule tendsto_mult_left [OF tendstoI])
then have (( $\lambda u. g u / (u-w)$ ) has_contour_integral  $2 * \text{of\_real } \pi * i * g w$ )
  (circlepath z r)
  using has_contour_integral_integral [OF g_cint] tendsto_unique [OF F
f_tends_cig] w
  by fastforce
then have (( $\lambda u. g u / (2 * \text{of\_real } \pi * i * (u-w))$ ) has_contour_integral g
w) (circlepath z r)
  using has_contour_integral_div [where c =  $2 * \text{of\_real } \pi * i$ ]
  by (force simp: field_simps)
then show ?thesis
  by (simp add: dist_norm)
qed
show ?thesis
  using Cauchy_next_derivative_circlepath(2) [OF 1 2, simplified]
  by (fastforce simp add: holomorphic_on_open contg intro: that)
qed

```

lemma higher_deriv_complex_uniform_limit:

```

assumes ulim: uniform_limit A f g F
  and f_holo: eventually ( $\lambda n. f n$  holomorphic_on A) F
  and F:  $F \neq \text{bot}$ 
  and A: open A z  $\in$  A
shows (( $\lambda n. (\text{deriv } \sim m) (f n) z$ )  $\longrightarrow$  ( $\text{deriv } \sim m) g z$ ) F
proof -
obtain r where r:  $r > 0$  cball z r  $\subseteq$  A
  using A by (meson open_contains_cball)
have r': ball z r  $\subseteq$  A
  using r by auto
define h where h = ( $\lambda n z. f n z - g z$ )
define c where c =  $\text{of\_real } (2 * \pi) * i / \text{fact } m$ 
have [simp]:  $c \neq 0$ 
  by (simp add: c_def)
have g holomorphic_on ball z r  $\wedge$  continuous_on (cball z r) g
proof (rule holomorphic_uniform_limit)
  show uniform_limit (cball z r) f g F
    by (rule uniform_limit_on_subset [OF ulim r(2)])
  show  $\forall_F n$  in F. continuous_on (cball z r) (f n)  $\wedge$  f n holomorphic_on ball z
r using f_holo
    by eventually_elim
      (use holomorphic_on_subset [OF _ r(2)] holomorphic_on_subset [OF _ r'])

  in <auto intro!: holomorphic_on_imp_continuous_on>
qed (use F in auto)
hence g_holo: g holomorphic_on ball z r and g_cont: continuous_on (cball z
r) g
  by blast+

```

```

have ulim': uniform_limit (sphere z r) ( $\lambda n x. h n x / (x - z) \wedge (Suc m)$ ) ( $\lambda_. 0$ ) F
proof -
  have uniform_limit (sphere z r) ( $\lambda n x. f n x / (x - z) \wedge Suc m$ ) ( $\lambda x. g x / (x - z) \wedge Suc m$ ) F
  proof (intro uniform_lim_divide uniform_limit_intros uniform_limit_on_subset[OF ulim])
    have compact (g ` sphere z r)
    by (intro compact_continuous_image continuous_on_subset[OF g_cont])
  auto
  thus bounded (g ` sphere z r)
  by (rule compact_imp_bounded)
  show  $r \wedge Suc m \leq norm ((x - z) \wedge Suc m)$  if  $x \in sphere z r$  for  $x$  unfolding
  norm_power
  by (intro power_mono) (use that r(1) in  $\langle auto simp: dist_norm norm_minus_commute \rangle$ )
  qed (use r in auto)
  hence uniform_limit (sphere z r) ( $\lambda n x. f n x / (x - z) \wedge Suc m - g x / (x - z) \wedge Suc m$ )
  ( $\lambda x. g x / (x - z) \wedge Suc m - g x / (x - z) \wedge Suc m$ ) F
  by (intro uniform_limit_intros)
  thus ?thesis
  by (simp add: h_def diff_divide_distrib)
qed

have has_integral: eventually ( $\lambda n. ((\lambda u. h n u / (u - z) \wedge Suc m)$  has_contour_integral
   $c * (deriv \sim m) (h n) z) (circlepath z r)$ ) F
  using f_holo
proof eventually_elim
  case (elim n)
  show ?case
  unfolding c_def
proof (rule Cauchy_has_contour_integral_higher_derivative_circlepath)
  show continuous_on (cball z r) (h n) unfolding h_def
  by (intro continuous_intros g_cont holomorphic_on_imp_continuous_on
  holomorphic_on_subset[OF elim] r)
  show h n holomorphic_on ball z r
  unfolding h_def by (intro holomorphic_intros g_holo holomorphic_on_subset[OF
  elim] r')
  qed (use r(1) in auto)
qed

have ( $(\lambda n. contour\_integral (circlepath z r) (\lambda u. h n u / (u - z) \wedge Suc m))$ 
 $\longrightarrow$ 
   $contour\_integral (circlepath z r) (\lambda u. 0 / (u - z) \wedge Suc m)$ ) F
proof (rule contour_integral_uniform_limit_circlepath)
  show  $\forall_F n$   $in F. (\lambda u. h n u / (u - z) \wedge Suc m)$  contour_integrable_on
  circlepath z r
  using has_integral by eventually_elim (blast intro: has_contour_integral_integrable)

```

```

qed (use r(1) ‹F ≠ bot› ulim' in simp_all)
hence ((λn. contour_integral (circlepath z r) (λu. h n u / (u - z) ^ Suc m))
  → 0) F
  by simp
also have ?this ‹→ ((λn. c * (deriv ~ m) (h n) z) → 0) F
proof (rule tendsto_cong)
  show ∀F x in F. contour_integral (circlepath z r) (λu. h x u / (u - z) ^ Suc
m) =
      c * (deriv ~ m) (h x) z
    using has_integral by eventually_elim (simp add: contour_integral_unique)
qed
finally have ((λn. (deriv ~ m) g z + c * (deriv ~ m) (h n) z / c) →
  (deriv ~ m) g z + 0 / c) F
  by (intro tendsto_intros) auto
also have ?this ‹→ ((λn. (deriv ~ m) (f n) z) → (deriv ~ m) g z) F
proof (intro filterlim_cong)
  show ∀F n in F. (deriv ~ m) g z + c * (deriv ~ m) (h n) z / c = (deriv ~
m) (f n) z
    using f_holo
  proof eventually_elim
    case (elim n)
    have (deriv ~ m) (h n) z = (deriv ~ m) (f n) z - (deriv ~ m) g z
unfolding h_def
    by (rule higher_deriv_diff holomorphic_on_subset[OF elim r'] g_holo A)+
    (use r(1) in auto)
    thus ?case
    by simp
  qed
qed auto
finally show ?thesis .
qed

```

lemma deriv_complex_uniform_limit:

```

assumes ulim: uniform_limit A f g F
  and f_holo: eventually (λn. f n holomorphic_on A) F
  and F: F ≠ bot
  and A: open A z ∈ A
  shows ((λn. deriv (f n) z) → deriv g z) F
using higher_deriv_complex_uniform_limit[OF assms, of 1] by simp

```

lemma logderiv_producing_complex_uniform_limit:

```

fixes f :: nat ⇒ complex ⇒ complex
assumes lim: uniform_limit A (λn x. ∏ k < n. f k x) P sequentially
assumes holo: ∧k. f k holomorphic_on A
assumes nz: P z ≠ 0
assumes A: open A z ∈ A
shows (λk. deriv (f k) z / f k z) sums (deriv P z / P z)
proof -
  define f' where f' = (λk. deriv (f k))

```

```

note [derivative_intros] = has_field_derivative_prod'
have [derivative_intros]:
  (f k has_field_derivative f' k z) (at z within B) if z ∈ A for B z k
  using that holomorphic_derivI[OF holo[of k], of z B] A unfolding f'_def by
  auto
have lim': (λn. ∏ k<n. f k z) ⟶ P z
  using lim by (rule tendsto_uniform_limitI) fact+

have nz': f k z ≠ 0 for k
proof
  assume f k z = 0
  have eventually (λn. (∏ k<n. f k z) = 0) sequentially
    using eventually_gt_at_top[of k] by eventually_elim (use ⟨f k z = 0⟩ in
  auto)
  hence (λn. (∏ k<n. f k z)) ⟶ 0
    by (rule tendsto_eventually)
  with lim' have P z = 0
    using tendsto_unique sequentially_bot by blast
  with nz show False
    by simp
qed

from lim have (λn. deriv (λx. ∏ k<n. f k x) z) ⟶ deriv P z
  by (rule deriv_complex_uniform_limit)
  (use A in ⟨auto intro!: always_eventually_holomorphic_intros holo⟩)
also have (λn. deriv (λx. ∏ k<n. f k x) z) = (λn. (∏ k<n. f k z) * (∑ k<n. f'
  k z / f k z))
  using ⟨z ∈ A⟩ by (auto intro!: ext DERIV_imp_deriv derivative_eq_intros
  simp: nz')
  finally have (λn. (∏ k<n. f k z) * (∑ k<n. f' k z / f k z)) ⟶ deriv P z .
  hence (λn. (∏ k<n. f k z) * (∑ k<n. f' k z / f k z) / (∏ k<n. f k z)) ⟶
  deriv P z / P z
    by (intro tendsto_intros) (use nz lim' in auto)
  also have (λn. (∏ k<n. f k z) * (∑ k<n. f' k z / f k z) / (∏ k<n. f k z)) =
    (λn. (∑ k<n. f' k z / f k z))
    by (simp add: nz')
  finally show (λk. f' k z / f k z) sums (deriv P z / P z)
    unfolding sums_def .
qed

```

Version showing that the limit is the limit of the derivatives.

```

proposition has_complex_derivative_uniform_limit:
  fixes z::complex
  assumes cont: eventually (λn. continuous_on (cball z r) (f n) ∧
    (∀ w ∈ ball z r. ((f n) has_field_derivative (f' n w)) (at
  w))) F
  and ulim: uniform_limit (cball z r) f g F
  and F: ¬ trivial_limit F and 0 < r
  obtains g' where

```

```

    continuous_on (cball z r) g
     $\bigwedge w. w \in \text{ball } z \ r \implies (g \text{ has\_field\_derivative } (g' w)) \text{ (at } w) \wedge ((\lambda n. f' n w) \longrightarrow g' w) \ F$ 
  proof -
    let ?conint = contour_integral (circlepath z r)
    have g: continuous_on (cball z r) g g holomorphic_on ball z r
      by (rule holomorphic_uniform_limit [OF eventually_mono [OF cont] ulim F];
          auto simp: holomorphic_on_open field_differentiable_def)
    then obtain g' where g':  $\bigwedge x. x \in \text{ball } z \ r \implies (g \text{ has\_field\_derivative } g' x) \text{ (at } x)$ 
      using DERIV_deriv_iff_has_field_derivative
      by (fastforce simp add: holomorphic_on_open)
    then have derg:  $\bigwedge x. x \in \text{ball } z \ r \implies \text{deriv } g \ x = g' x$ 
      by (simp add: DERIV_imp_deriv)
    have tends_f'n_g':  $((\lambda n. f' n w) \longrightarrow g' w) \ F$  if w:  $w \in \text{ball } z \ r$  for w
      proof -
        have eq_f': ?conint  $(\lambda x. f \ n \ x / (x-w)^2) - ?conint (\lambda x. g \ x / (x-w)^2) = (f' n w - g' w) * (2 * \text{of\_real } \pi * i)$ 
          if cont_fn: continuous_on (cball z r) (f n)
          and fnd:  $\bigwedge w. w \in \text{ball } z \ r \implies (f \ n \ \text{has\_field\_derivative } f' n w) \text{ (at } w)$ 
        for n
          proof -
            have hol_fn: f n holomorphic_on ball z r
              using fnd by (force simp: holomorphic_on_open)
            have (f n has_field_derivative 1 / (2 * of_real pi * i) * ?conint  $(\lambda u. f \ n \ u / (u-w)^2)$ ) (at w)
              by (rule Cauchy_derivative_integral_circlepath [OF cont_fn hol_fn w])
            then have f':  $f' n w = 1 / (2 * \text{of\_real } \pi * i) * ?conint (\lambda u. f \ n \ u / (u-w)^2)$ 
              using DERIV_unique [OF fnd] w by blast
            show ?thesis
              by (simp add: f' Cauchy_contour_integral_circlepath_2 [OF g w] derg [OF w] field_split_simps)
          qed
        define d where  $d = (r - \text{norm}(w-z))^2$ 
        have  $d > 0$ 
          using w by (simp add: dist_commute dist_norm d_def)
        have dle:  $d \leq \text{cmod } ((y-w)^2)$  if  $r = \text{cmod } (z-y)$  for y
          by (smt (verit, best) d_def diff_add_cancel diff_diff_eq2 dist_norm mem_ball norm_minus_commute norm_power norm_triangle_ineq2 power_mono)
        that w
        have 1:  $\forall_F n \ \text{in } F. (\lambda x. f \ n \ x / (x-w)^2) \text{ contour\_integrable\_on } \text{circlepath } z \ r$ 
          by (force simp: holomorphic_on_open intro: w Cauchy_derivative_integral_circlepath eventually_mono [OF cont])
        have 2: uniform_limit (sphere z r)  $(\lambda n \ x. f \ n \ x / (x-w)^2) (\lambda x. g \ x / (x-w)^2)$ 
          F
          unfolding uniform_limit_iff
        proof clarify
          fix e::real
          assume  $e > 0$ 

```

```

with ⟨ $r > 0$ ⟩
have  $\forall_F n \text{ in } F. \forall x. x \neq w \longrightarrow \text{cmod } (z-x) = r \longrightarrow \text{cmod } (f \ n \ x - g \ x) < e * \text{cmod } ((x-w)^2)$ 
by (force simp: ⟨ $0 < d$ ⟩ dist_norm dle intro: less_le_trans eventually_mono [OF uniform_limitD [OF ulim], of e*d])
with ⟨ $r > 0$ ⟩ ⟨ $e > 0$ ⟩
show  $\forall_F n \text{ in } F. \forall x \in \text{sphere } z \ r. \text{dist } (f \ n \ x / (x-w)^2) (g \ x / (x-w)^2) < e$ 
by (simp add: norm_divide field_split_simps sphere_def dist_norm)
qed
have (( $\lambda n. \text{contour\_integral } (\text{circlepath } z \ r) (\lambda x. f \ n \ x / (x-w)^2)$ )
 $\longrightarrow \text{contour\_integral } (\text{circlepath } z \ r) ((\lambda x. g \ x / (x-w)^2))$ ) F
by (rule contour_integral_uniform_limit_circlepath [OF 1 2 F ⟨ $0 < r$ ⟩])
then have tendsto_0: (( $\lambda n. 1 / (2 * \text{of\_real } \pi * i) * (?conint (\lambda x. f \ n \ x / (x-w)^2) - ?conint (\lambda x. g \ x / (x-w)^2))$ )  $\longrightarrow 0$ ) F
using Lim_null by (force intro!: tendsto_mult_right_zero)
have (( $\lambda n. f' \ n \ w - g' \ w$ )  $\longrightarrow 0$ ) F
by (force simp: divide_simps
intro: eq_f' eventually_mono [OF cont] Lim_transform_eventually [OF tendsto_0])
then show ?thesis using Lim_null by blast
qed
obtain  $g'$  where  $\bigwedge w. w \in \text{ball } z \ r \implies (g \text{ has\_field\_derivative } (g' \ w)) (at \ w) \wedge ((\lambda n. f' \ n \ w) \longrightarrow g' \ w)$  F
by (blast intro: tends_f'n_g' g')
then show ?thesis using  $g$ 
using that by blast
qed

```

4.9 Some more simple/convenient versions for applications

lemma holomorphic_uniform_sequence:

```

assumes  $S$ : open  $S$ 
and hol_fn:  $\bigwedge n. (f \ n)$  holomorphic_on  $S$ 
and ulim_g:  $\bigwedge x. x \in S \implies \exists d. 0 < d \wedge \text{cball } x \ d \subseteq S \wedge \text{uniform\_limit } (\text{cball } x \ d) \ f \ g$  sequentially
shows  $g$  holomorphic_on  $S$ 
proof -
have  $\exists f'. (g \text{ has\_field\_derivative } f') (at \ z)$  if  $z \in S$  for  $z$ 
proof -
obtain  $r$  where  $0 < r$  and  $r$ :  $\text{cball } z \ r \subseteq S$ 
and  $ul$ : uniform_limit (cball  $z \ r$ )  $f \ g$  sequentially
using ulim_g [OF ⟨ $z \in S$ ⟩] by blast
have *:  $\forall_F n \text{ in } \text{sequentially. continuous\_on } (\text{cball } z \ r) (f \ n) \wedge f \ n$  holomorphic_on ball  $z \ r$ 
by (smt (verit, best) ball_subset_cball hol_fn holomorphic_on_imp_continuous_on
holomorphic_on_subset not_eventuallyD  $r$ )
show ?thesis
using ⟨ $0 < r$ ⟩ centre_in_ball  $ul$ 

```



```

    by (auto simp: holomorphic_on_open intro: holomorphic_uniform_limit [OF
*])
  qed
  with S show ?thesis
    by (simp add: holomorphic_on_open)
  qed

```

lemma *has_complex_derivative_uniform_sequence:*

```

  fixes S :: complex set
  assumes S: open S
    and hfd:  $\bigwedge n x. x \in S \implies ((f\ n)\ \text{has\_field\_derivative}\ f'\ n\ x)\ (at\ x)$ 
    and ulim_g:  $\bigwedge x. x \in S \implies \exists d. 0 < d \wedge cball\ x\ d \subseteq S \wedge \text{uniform\_limit}\ (cball\ x\ d)\ f\ g\ \text{sequentially}$ 
  shows  $\exists g'. \forall x \in S. (g\ \text{has\_field\_derivative}\ g'\ x)\ (at\ x) \wedge ((\lambda n. f'\ n\ x) \longrightarrow g'\ x)\ \text{sequentially}$ 
  proof -
    have y:  $\exists y. (g\ \text{has\_field\_derivative}\ y)\ (at\ z) \wedge (\lambda n. f'\ n\ z) \longrightarrow y$  if  $z \in S$ 
  for z
    proof -
      obtain r where  $0 < r$  and  $r: cball\ z\ r \subseteq S$ 
        and ul: uniform_limit (cball z r) f g sequentially
        using ulim_g [OF ‹z ∈ S›] by blast
      have *:  $\forall_F n\ \text{in}\ \text{sequentially. continuous\_on}\ (cball\ z\ r)\ (f\ n) \wedge$ 
         $(\forall w \in ball\ z\ r. ((f\ n)\ \text{has\_field\_derivative}\ (f'\ n\ w))$ 
      (at w))
    proof (intro eventuallyI conjI ballI)
      show continuous_on (cball z r) (f x) for x
      by (meson S continuous_on_subset hfd holomorphic_on_imp_continuous_on
holomorphic_on_open r)
      show  $w \in ball\ z\ r \implies (f\ x\ \text{has\_field\_derivative}\ f'\ x\ w)\ (at\ w)$  for w x
        using ball_subset_cball hfd r by blast
    qed
    show ?thesis
      by (rule has_complex_derivative_uniform_limit [OF *, of g]) (use ‹0 < r›
ul in ‹force+›)
  qed
  show ?thesis
    by (rule bchoice) (blast intro: y)
  qed

```

4.10 On analytic functions defined by a series

lemma *series_and_derivative_comparison:*

```

  fixes S :: complex set
  assumes S: open S
    and h: summable h
    and hfd:  $\bigwedge n x. x \in S \implies (f\ n\ \text{has\_field\_derivative}\ f'\ n\ x)\ (at\ x)$ 
    and to_g:  $\forall_F n\ \text{in}\ \text{sequentially. } \forall x \in S. \text{norm}\ (f\ n\ x) \leq h\ n$ 
  obtains  $g\ g'$  where  $\forall x \in S. ((\lambda n. f\ n\ x)\ \text{sums}\ g\ x) \wedge ((\lambda n. f'\ n\ x)\ \text{sums}\ g'\ x)$ 

```

$\wedge (g \text{ has_field_derivative } g' x) (at x)$
proof –
obtain g **where** g : *uniform_limit* $S (\lambda n x. \sum_{i < n}. f i x)$ *g sequentially*
using *Weierstrass_m_test_ev* [*OF to_g h*] **by force**
have $*$: $\exists d > 0. cball\ x\ d \subseteq S \wedge \text{uniform_limit } (cball\ x\ d) (\lambda n x. \sum_{i < n}. f i x)$
g sequentially
if $x \in S$ **for** x
using *open_contains_cball* [*of S*] $\langle x \in S \rangle S$ *g uniform_limit_on_subset* **by blast**
have $\bigwedge x. x \in S \implies (\lambda n. \sum_{i < n}. f i x) \longrightarrow g x$
by (*metis tendsto_uniform_limitI* [*OF g*])
moreover have $\exists g'. \forall x \in S. (g \text{ has_field_derivative } g' x) (at x) \wedge (\lambda n. \sum_{i < n}. f' i x) \longrightarrow g' x$
by (*rule has_complex_derivative_uniform_sequence* [*OF S*]) (*auto intro: * hfd DERIV_sum*)
ultimately show *?thesis*
by (*metis sums_def that*)
qed

A version where we only have local uniform/comparative convergence.

lemma *series_and_derivative_comparison_local*:

fixes $S :: \text{complex set}$
assumes S : *open S*
and *hfd*: $\bigwedge n x. x \in S \implies (f n \text{ has_field_derivative } f' n x) (at x)$
and *to_g*: $\bigwedge x. x \in S \implies \exists d h. 0 < d \wedge \text{summable } h \wedge (\forall_F n \text{ in sequentially. } \forall y \in \text{ball } x\ d \cap S. \text{norm } (f n y) \leq h n)$
shows $\exists g g'. \forall x \in S. ((\lambda n. f n x) \text{ sums } g x) \wedge ((\lambda n. f' n x) \text{ sums } g' x) \wedge (g \text{ has_field_derivative } g' x) (at x)$
proof –
have $\exists y. (\lambda n. f n z) \text{ sums } (\sum n. f n z) \wedge (\lambda n. f' n z) \text{ sums } y \wedge ((\lambda x. \sum n. f n x) \text{ has_field_derivative } y) (at z)$
if $z \in S$ **for** z
proof –
obtain $d h$ **where** $0 < d$ *summable h* **and** *le_h*: $\forall_F n \text{ in sequentially. } \forall y \in \text{ball } z\ d \cap S. \text{norm } (f n y) \leq h n$
using *to_g* $\langle z \in S \rangle$ **by meson**
then obtain r **where** $r > 0$ **and** r : $\text{ball } z\ r \subseteq \text{ball } z\ d \cap S$ **using** $\langle z \in S \rangle S$
by (*metis Int_iff_open_ball_centre_in_ball open_Int open_contains_ball_eq*)
have 1 : *open* ($\text{ball } z\ d \cap S$)
by (*simp add: open_Int*)
have 2 : $\bigwedge n x. x \in \text{ball } z\ d \cap S \implies (f n \text{ has_field_derivative } f' n x) (at x)$
by (*auto simp: hfd*)
obtain $g g'$ **where** gg' : $\forall x \in \text{ball } z\ d \cap S. ((\lambda n. f n x) \text{ sums } g x) \wedge ((\lambda n. f' n x) \text{ sums } g' x) \wedge (g \text{ has_field_derivative } g' x) (at x)$
by (*auto intro: le_h series_and_derivative_comparison* [*OF 1*] $\langle \text{summable } h \rangle$ *hfd*)
then have $(\lambda n. f' n z) \text{ sums } g' z$
by (*meson* $\langle 0 < r \rangle \text{centre_in_ball contra_subsetD } r$)

```

moreover have  $(\lambda n. f\ n\ z)$  sums  $(\sum n. f\ n\ z)$ 
  using summable_sums centre_in_ball  $\langle 0 < d \rangle$  summable_h le_h
  by (metis (full_types) Int_iff gg' summable_def that)
moreover have  $((\lambda x. \sum n. f\ n\ x)$  has_field_derivative  $g'\ z)$  (at  $z$ )
by (metis (no_types, lifting) 1 r  $\langle 0 < r \rangle$  gg' has_field_derivative_transform_within_open

      open_contains_ball_eq sums_unique)
ultimately show ?thesis by auto
qed
then show ?thesis
  by meson
qed

```

Sometimes convenient to compare with a complex series of positive reals.
(?)

```

lemma series_and_derivative_comparison_complex:
  fixes  $S :: \text{complex set}$ 
  assumes  $S$ : open  $S$ 
    and hfd:  $\bigwedge n\ x. x \in S \implies (f\ n\ \text{has\_field\_derivative } f'\ n\ x)$  (at  $x$ )
    and to_g:  $\bigwedge x. x \in S \implies \exists d\ h. 0 < d \wedge \text{summable } h \wedge \text{range } h \subseteq \mathbb{R}_{\geq 0} \wedge (\forall_F$ 
n in sequentially.  $\forall y \in \text{ball } x\ d \cap S. \text{cmod}(f\ n\ y) \leq \text{cmod}(h\ n))$ )
  shows  $\exists g\ g'. \forall x \in S. ((\lambda n. f\ n\ x)$  sums  $g\ x) \wedge ((\lambda n. f'\ n\ x)$  sums  $g'\ x) \wedge (g$ 
has_field_derivative  $g'\ x)$  (at  $x$ )
  apply (rule series_and_derivative_comparison_local [OF S hfd], assumption)
  apply (rule ex_forward [OF to_g], assumption)
  apply (erule exE)
  apply (rule_tac  $x = \text{Re} \circ h$  in exI)
  apply (force simp: summable_Re o_def nonneg_Reals_cmod_eq_Re image_subset_iff)
  done

```

Sometimes convenient to compare with a complex series of positive reals.
(?)

```

lemma series_differentiable_comparison_complex:
  fixes  $S :: \text{complex set}$ 
  assumes  $S$ : open  $S$ 
    and hfd:  $\bigwedge n\ x. x \in S \implies f\ n$  field_differentiable (at  $x$ )
    and to_g:  $\bigwedge x. x \in S \implies \exists d\ h. 0 < d \wedge \text{summable } h \wedge \text{range } h \subseteq \mathbb{R}_{\geq 0} \wedge (\forall_F$ 
n in sequentially.  $\forall y \in \text{ball } x\ d \cap S. \text{cmod}(f\ n\ y) \leq \text{cmod}(h\ n))$ )
  obtains  $g$  where  $\forall x \in S. ((\lambda n. f\ n\ x)$  sums  $g\ x) \wedge g$  field_differentiable (at  $x$ )
  proof –
    have hfd':  $\bigwedge n\ x. x \in S \implies (f\ n\ \text{has\_field\_derivative } \text{deriv } (f\ n)\ x)$  (at  $x$ )
      using hfd field_differentiable_derivI by blast
    show ?thesis
      by (metis field_differentiable_def that series_and_derivative_comparison_complex
        [OF S hfd' to_g])
  qed

```

In particular, a power series is analytic inside circle of convergence.

```

lemma power_series_and_derivative_0:

```

```

fixes  $a :: \text{nat} \Rightarrow \text{complex}$  and  $r :: \text{real}$ 
assumes  $\text{summable } (\lambda n. a\ n * r^{\wedge} n)$ 
shows  $\exists g\ g'. \forall z. \text{cmod } z < r \longrightarrow$ 
 $((\lambda n. a\ n * z^{\wedge} n) \text{ sums } g\ z) \wedge ((\lambda n. \text{of\_nat } n * a\ n * z^{\wedge}(n-1)) \text{ sums } g'$ 
 $z) \wedge (g \text{ has\_field\_derivative } g' z) (at\ z)$ 
proof (cases  $0 < r$ )
  case True
    have  $\text{der}: \bigwedge n\ z. ((\lambda x. a\ n * x^{\wedge} n) \text{ has\_field\_derivative } \text{of\_nat } n * a\ n * z^{\wedge}$ 
 $(n-1)) (at\ z)$ 
      by (rule derivative_eq_intros | simp) +
    have  $y\_le: \text{cmod } y \leq \text{cmod } (\text{of\_real } r + \text{of\_real } (\text{cmod } z)) / 2$ 
      if  $\text{cmod } (z-y) * 2 < r - \text{cmod } z$  for  $z\ y$ 
      by (smt (verit, best) field_sum_of_halves norm_minus_commute norm_of_real
norm_triangle_ineq2 of_real_add that)
    have  $\text{summable } (\lambda n. a\ n * \text{complex\_of\_real } r^{\wedge} n)$ 
      using assms  $\langle r > 0 \rangle$  by simp
    moreover have  $\bigwedge z. \text{cmod } z < r \implies \text{cmod } ((\text{of\_real } r + \text{of\_real } (\text{cmod } z)) /$ 
 $2) < \text{cmod } (\text{of\_real } r)$ 
      using  $\langle r > 0 \rangle$ 
      by (simp flip: of_real_add)
    ultimately have  $\text{sum}: \bigwedge z. \text{cmod } z < r \implies \text{summable } (\lambda n. \text{of\_real } (\text{cmod } (a$ 
 $n)) * ((\text{of\_real } r + \text{complex\_of\_real } (\text{cmod } z)) / 2)^{\wedge} n)$ 
      by (rule power_series_conv_imp_absconv_weak)
    have  $\exists g\ g'. \forall z \in \text{ball } 0\ r. (\lambda n. (a\ n) * z^{\wedge} n) \text{ sums } g\ z \wedge$ 
 $(\lambda n. \text{of\_nat } n * (a\ n) * z^{\wedge}(n-1)) \text{ sums } g' z \wedge (g \text{ has\_field\_derivative}$ 
 $g' z) (at\ z)$ 
      apply (rule series_and_derivative_comparison_complex [OF open_ball der])
      apply (rule tac x=(r - norm z)/2 in exI)
      apply (rule tac x= $\lambda n. \text{of\_real}(norm(a\ n)*((r + norm z)/2)^{\wedge} n$ ) in exI)
      using  $\langle r > 0 \rangle$ 
      apply (auto simp: sum_eventually_sequentially norm_mult norm_power
dist_norm intro!: mult_left_mono power_mono y_le)
    done
  then show ?thesis
    by (simp add: ball_def)
next
  case False then show ?thesis
    unfolding not_less using less_le_trans norm_not_less_zero by blast
qed

```

proposition *power_series_and_derivative*:

```

fixes  $a :: \text{nat} \Rightarrow \text{complex}$  and  $r :: \text{real}$ 
assumes  $\text{summable } (\lambda n. a\ n * r^{\wedge} n)$ 
obtains  $g\ g'$  where  $\forall z \in \text{ball } w\ r.$ 
 $((\lambda n. a\ n * (z-w)^{\wedge} n) \text{ sums } g\ z) \wedge ((\lambda n. \text{of\_nat } n * a\ n * (z-w)^{\wedge}$ 
 $(n-1)) \text{ sums } g' z) \wedge$ 
 $(g \text{ has\_field\_derivative } g' z) (at\ z)$ 
using power_series_and_derivative_0 [OF assms]
apply clarify

```

```

apply (rule_tac g=( $\lambda z. g(z-w)$ ) in that)
using DERIV_shift [where  $z=-w$ ]
apply (auto simp: norm_minus_commute Ball_def dist_norm)
done

```

proposition power_series_holomorphic:

```

assumes  $\bigwedge w. w \in \text{ball } z \ r \implies ((\lambda n. a \ n * (w-z)^{\wedge} n) \text{ sums } f \ w)$ 
shows  $f \text{ holomorphic\_on ball } z \ r$ 

```

proof –

```

have  $\exists f'. (f \text{ has\_field\_derivative } f') \text{ (at } w) \text{ if } w: \text{dist } z \ w < r$  for  $w$ 

```

proof –

```

have  $wz: \text{cmod } (w-z) < r$  using  $w$ 

```

```

by (auto simp: field_split_simps dist_norm norm_minus_commute)

```

```

then have  $0 \leq r$ 

```

```

by (meson less_eq_real_def norm_ge_zero order_trans)

```

```

have  $\text{inb}: z + \text{complex\_of\_real } ((\text{dist } z \ w + r) / 2) \in \text{ball } z \ r$ 

```

```

using  $w$  by (simp add: dist_norm ‹ $0 \leq r$ › flip: of_real_add)

```

```

have  $\text{sum}: \text{summable } (\lambda n. a \ n * \text{of\_real } (((\text{cmod } (z-w) + r) / 2)^{\wedge} n))$ 

```

```

using  $\text{assms}$  [OF inb] by (force simp: summable_def dist_norm)

```

```

obtain  $g \ g'$  where  $gg': \bigwedge u. u \in \text{ball } z \ ((\text{cmod } (z-w) + r) / 2) \implies$ 

```

```

 $(\lambda n. a \ n * (u-z)^{\wedge} n) \text{ sums } g \ u \wedge$ 

```

```

 $(\lambda n. \text{of\_nat } n * a \ n * (u-z)^{\wedge} (n-1)) \text{ sums } g' \ u \wedge (g$ 

```

```

 $\text{has\_field\_derivative } g' \ u) \text{ (at } u)$ 

```

```

by (rule power_series_and_derivative [OF sum, of z]) fastforce

```

```

have [simp]:  $g \ u = f \ u$  if  $\text{cmod } (u-w) < (r - \text{cmod } (z-w)) / 2$  for  $u$ 

```

proof –

```

have  $\text{less}: \text{cmod } (z-u) * 2 < \text{cmod } (z-w) + r$ 

```

```

using that dist_triangle2 [of z u w]

```

```

by (simp add: dist_norm [symmetric] algebra_simps)

```

```

have  $(\lambda n. a \ n * (u-z)^{\wedge} n) \text{ sums } g \ u \ (\lambda n. a \ n * (u-z)^{\wedge} n) \text{ sums } f \ u$ 

```

```

using  $gg'$  [of u] less  $w$  by (auto simp: assms dist_norm)

```

```

then show ?thesis

```

```

by (metis sums_unique2)

```

qed

```

have  $(f \text{ has\_field\_derivative } g' \ w) \text{ (at } w)$ 

```

```

proof (rule has_field_derivative_transform_within [where  $d=(r - \text{norm}(z-w))/2$ ])

```

```

qed (use  $w \ gg'$  [of w] in ‹(force simp: dist_norm)+›)

```

```

then show ?thesis ..

```

qed

```

then show ?thesis by (simp add: holomorphic_on_open)

```

qed

corollary holomorphic_iff_power_series:

```

 $f \text{ holomorphic\_on ball } z \ r \iff$ 

```

```

 $(\forall w \in \text{ball } z \ r. (\lambda n. (\text{deriv } \wedge n) f \ z / (\text{fact } n) * (w-z)^{\wedge} n) \text{ sums } f \ w)$ 

```

```

using power_series_holomorphic [where  $a = \lambda n. (\text{deriv } \wedge n) f \ z / (\text{fact } n)$ ]

```

holomorphic_power_series

by blast

lemma *power_series_analytic*:

$(\bigwedge w. w \in \text{ball } z \ r \implies (\lambda n. a \ n * (w-z)^{\wedge} n) \text{ sums } f \ w) \implies f \text{ analytic_on_ball } z \ r$
by (*force simp: analytic_on_open intro!: power_series_holomorphic*)

lemma *analytic_iff_power_series*:

$f \text{ analytic_on_ball } z \ r \longleftrightarrow$
 $(\forall w \in \text{ball } z \ r. (\lambda n. (\text{deriv } \sim n) f \ z / (\text{fact } n) * (w-z)^{\wedge} n) \text{ sums } f \ w)$
by (*simp add: analytic_on_open holomorphic_iff_power_series*)

4.11 Equality between holomorphic functions, on open ball then connected set

lemma *holomorphic_fun_eq_on_ball*:

$\llbracket f \text{ holomorphic_on_ball } z \ r; g \text{ holomorphic_on_ball } z \ r;$
 $w \in \text{ball } z \ r;$
 $\bigwedge n. (\text{deriv } \sim n) f \ z = (\text{deriv } \sim n) g \ z \rrbracket$
 $\implies f \ w = g \ w$
by (*auto simp: holomorphic_iff_power_series sums_unique2 [of $\lambda n. (\text{deriv } \sim n) f \ z / (\text{fact } n) * (w-z)^{\wedge} n$]*)

lemma *holomorphic_fun_eq_0_on_ball*:

$\llbracket f \text{ holomorphic_on_ball } z \ r; w \in \text{ball } z \ r;$
 $\bigwedge n. (\text{deriv } \sim n) f \ z = 0 \rrbracket$
 $\implies f \ w = 0$
using *holomorphic_fun_eq_on_ball [where $g = \lambda z. 0$]* **by** *simp*

lemma *holomorphic_fun_eq_0_on_connected*:

assumes *holf*: $f \text{ holomorphic_on } S$ **and** *open* S
and *cons*: *connected* S
and *der*: $\bigwedge n. (\text{deriv } \sim n) f \ z = 0$
and $z \in S \ w \in S$
shows $f \ w = 0$
proof –
have $*$: $\text{ball } x \ e \subseteq (\bigcap n. \{w \in S. (\text{deriv } \sim n) f \ w = 0\})$
if $\forall u. (\text{deriv } \sim u) f \ x = 0 \ \text{ball } x \ e \subseteq S$ **for** $x \ e$
proof –
have $(\text{deriv } \sim m) ((\text{deriv } \sim n) f) \ x = 0$ **for** $m \ n$
by (*metis funpow_add o_apply that(1)*)
then have $\bigwedge x' \ n. \text{dist } x \ x' < e \implies (\text{deriv } \sim n) f \ x' = 0$
using $\langle \text{open } S \rangle$
by (*meson holf holomorphic_fun_eq_0_on_ball holomorphic_higher_deriv holomorphic_on_subset mem_ball that(2)*)
with that show *?thesis* **by** *auto*
qed
obtain e **where** $e > 0$ **and** $e: \text{ball } w \ e \subseteq S$ **using** *openE [OF $\langle \text{open } S \rangle \langle w \in S \rangle$]*
then have *holfb*: $f \text{ holomorphic_on_ball } w \ e$
using *holf holomorphic_on_subset* **by** *blast*
have *open* $(\bigcap n. \{w \in S. (\text{deriv } \sim n) f \ w = 0\})$

```

using ⟨open S⟩
apply (simp add: open_contains_ball Ball_def image_iff)
by (metis (mono_tags) * mem_Collect_eq)
then have openin (top_of_set S) ( $\bigcap n. \{w \in S. (\text{deriv } \sim n) f w = 0\}$ )
by (force intro: open_subset)
moreover have closedin (top_of_set S) ( $\bigcap n. \{w \in S. (\text{deriv } \sim n) f w = 0\}$ )
using assms
by (auto intro: continuous_closedin_preimage_constant holomorphic_on_imp_continuous_on
holomorphic_higher_deriv)
moreover have ( $\bigcap n. \{w \in S. (\text{deriv } \sim n) f w = 0\} = S \implies f w = 0$ )
using ⟨e>0⟩ e by (force intro: holomorphic_fun_eq_0_on_ball [OF holfb])
ultimately show ?thesis
using cons der ⟨z ∈ S⟩
by (auto simp add: connected_clopen)
qed

```

lemma holomorphic_fun_eq_on_connected:

```

assumes f holomorphic_on S g holomorphic_on S and open S connected S
and  $\bigwedge n. (\text{deriv } \sim n) f z = (\text{deriv } \sim n) g z$ 
and  $z \in S w \in S$ 
shows  $f w = g w$ 
proof (rule holomorphic_fun_eq_0_on_connected [of  $\lambda x. f x - g x$  S z, simplified])
show ( $\lambda x. f x - g x$ ) holomorphic_on S
by (intro assms holomorphic_intros)
show  $\bigwedge n. (\text{deriv } \sim n) (\lambda x. f x - g x) z = 0$ 
using assms higher_deriv_diff by auto
qed (use assms in auto)

```

lemma holomorphic_fun_eq_const_on_connected:

```

assumes holf: f holomorphic_on S and open S
and cons: connected S
and der:  $\bigwedge n. 0 < n \implies (\text{deriv } \sim n) f z = 0$ 
and  $z \in S w \in S$ 
shows  $f w = f z$ 
proof (rule holomorphic_fun_eq_0_on_connected [of  $\lambda w. f w - f z$  S z, simplified])
show ( $\lambda w. f w - f z$ ) holomorphic_on S
by (intro assms holomorphic_intros)
show  $\bigwedge n. (\text{deriv } \sim n) (\lambda w. f w - f z) z = 0$ 
by (subst higher_deriv_diff) (use assms in ⟨auto intro: holomorphic_intros⟩)
qed (use assms in auto)

```

4.12 Some basic lemmas about poles/singularities

lemma pole_lemma:

```

assumes holf: f holomorphic_on S and a:  $a \in \text{interior } S$ 
shows ( $\lambda z. \text{if } z = a \text{ then deriv } f a$ 
else  $(f z - f a) / (z - a)$ ) holomorphic_on S (is ?F holomorphic_on S)

```

```

proof -
  have *: ?F field_differentiable (at u within S) if  $u \in S$   $u \neq a$  for  $u$ 
  proof -
    have fcd:  $f$  field_differentiable at  $u$  within  $S$ 
    using hol_f holomorphic_on_def by (simp add:  $\langle u \in S \rangle$ )
    have cd:  $(\lambda z. (f z - f a) / (z-a))$  field_differentiable at  $u$  within  $S$ 
    by (rule fcd derivative_intros | simp add: that)+
    have  $0 < \text{dist } a \ u$  using that dist_nz by blast
    then show ?thesis
    by (rule field_differentiable_transform_within [OF _ _ _ cd]) (auto simp:
 $\langle u \in S \rangle$ )
  qed
  moreover
  have ?F field_differentiable at  $a$  if  $0 < e$  ball  $a \ e \subseteq S$  for  $e$ 
  proof -
    have holfb:  $f$  holomorphic_on ball  $a \ e$ 
    by (rule holomorphic_on_subset [OF hol_f  $\langle \text{ball } a \ e \subseteq S \rangle$ ])
    have 2: ?F holomorphic_on ball  $a \ e - \{a\}$ 
    using mem_ball that
    by (auto simp add: holomorphic_on_def simp flip: field_differentiable_def
intro: * field_differentiable_within_subset)
    have isCont  $(\lambda z. \text{if } z = a \text{ then deriv } f \ a \ \text{else } (f z - f a) / (z-a)) \ x$ 
    if  $\text{dist } a \ x < e$  for  $x$ 
    proof (cases  $x=a$ )
      case True
      then have  $f$  field_differentiable at  $a$ 
      using holfb  $\langle 0 < e \rangle$  holomorphic_on_imp_differentiable_at by auto
      with True show ?thesis
      by (smt (verit) DERIV_deriv_iff_field_differentiable LIM_equal continuous_at has_field_derivativeD)
    next
      case False with 2 that show ?thesis
      by (simp add: field_differentiable_imp_continuous_at holomorphic_on_imp_differentiable_at
open_Diff)
    qed
    then have 1: continuous_on (ball  $a \ e$ ) ?F
    by (clarsimp simp: continuous_on_eq_continuous_at)
    have ?F holomorphic_on ball  $a \ e$ 
    by (auto intro: no_isolated_singularity [OF 1 2])
    with that show ?thesis
    by (simp add: holomorphic_on_imp_differentiable_at)
  qed
  ultimately show ?thesis
  by (metis (lifting) a_at_within_interior holomorphic_onI mem_interior)
qed

lemma pole_theorem:
  assumes holg:  $g$  holomorphic_on  $S$  and  $a: a \in \text{interior } S$ 
  and eq:  $\bigwedge z. z \in S - \{a\} \implies g \ z = (z-a) * f \ z$ 

```



```

shows ( $\lambda z$ . if  $z = a$  then  $\text{deriv } g \ a$ 
         else  $f \ z - g \ a / (z - a)$ ) holomorphic_on S
using pole_lemma [OF holg a]
by (rule holomorphic_transform) (simp add: eq field_split_simps)

lemma pole_lemma_open:
assumes  $f$  holomorphic_on S open S
shows ( $\lambda z$ . if  $z = a$  then  $\text{deriv } f \ a$  else  $(f \ z - f \ a) / (z - a)$ ) holomorphic_on S
proof (cases a ∈ S)
case True with assms interior_eq pole_lemma
show ?thesis by fastforce
next
case False
then have ( $\lambda z$ .  $(f \ z - f \ a) / (z - a)$ ) field_differentiable at x within S
if  $x \in S$  for  $x$ 
using assms that unfolding holomorphic_on_def
by (intro derivative_intros) auto
with False show ?thesis
using holomorphic_on_def holomorphic_transform by presburger
qed

lemma pole_theorem_open:
assumes holg: g holomorphic_on S and S: open S
and eq:  $\bigwedge z. z \in S - \{a\} \implies g \ z = (z - a) * f \ z$ 
shows ( $\lambda z$ . if  $z = a$  then  $\text{deriv } g \ a$ 
         else  $f \ z - g \ a / (z - a)$ ) holomorphic_on S
using pole_lemma_open [OF holg S]
by (rule holomorphic_transform) (auto simp: eq divide_simps)

lemma pole_theorem_0:
assumes holg: g holomorphic_on S and a: a ∈ interior S
and eq:  $\bigwedge z. z \in S - \{a\} \implies g \ z = (z - a) * f \ z$ 
and [simp]:  $f \ a = \text{deriv } g \ a \ g \ a = 0$ 
shows  $f$  holomorphic_on S
using pole_theorem [OF holg a eq]
by (rule holomorphic_transform) (auto simp: eq field_split_simps)

lemma pole_theorem_open_0:
assumes holg: g holomorphic_on S and S: open S
and eq:  $\bigwedge z. z \in S - \{a\} \implies g \ z = (z - a) * f \ z$ 
and [simp]:  $f \ a = \text{deriv } g \ a \ g \ a = 0$ 
shows  $f$  holomorphic_on S
using pole_theorem_open [OF holg S eq]
by (rule holomorphic_transform) (auto simp: eq field_split_simps)

lemma pole_theorem_analytic:
assumes  $g$ :  $g$  analytic_on S
and eq:  $\bigwedge z. z \in S$ 
          $\implies \exists d. 0 < d \wedge (\forall w \in \text{ball } z \ d - \{a\}. g \ w = (w - a) * f \ w)$ 

```

shows $(\lambda z. \text{if } z = a \text{ then deriv } g \ a \ \text{else } f \ z - g \ a / (z - a)) \text{ analytic_on } S$ (**is** $?F$ *analytic_on* S)
unfolding *analytic_on_def*
proof
fix x
assume $x \in S$
with g **obtain** e **where** $0 < e$ **and** $e: g \text{ holomorphic_on } \text{ball } x \ e$
by (*auto simp add: analytic_on_def*)
obtain d **where** $0 < d$ **and** $d: \bigwedge w. w \in \text{ball } x \ d - \{a\} \implies g \ w = (w - a) * f \ w$
using $\langle x \in S \rangle$ **eq** **by** *blast*
have $?F \text{ holomorphic_on } \text{ball } x \ (\min \ d \ e)$
using $d \ e \ \langle x \in S \rangle$ **by** (*fastforce simp: holomorphic_on_subset subset_ball intro!: pole_theorem_open*)
then show $\exists e > 0. ?F \text{ holomorphic_on } \text{ball } x \ e$
using $\langle 0 < d \rangle \ \langle 0 < e \rangle$ **not_le** **by** *fastforce*
qed

lemma *pole_theorem_analytic_0*:

assumes $g: g \text{ analytic_on } S$
and $eq: \bigwedge z. z \in S \implies \exists d. 0 < d \wedge (\forall w \in \text{ball } z \ d - \{a\}. g \ w = (w - a) * f \ w)$
and [*simp*]: $f \ a = \text{deriv } g \ a \ g \ a = 0$
shows $f \text{ analytic_on } S$
proof –
have [*simp*]: $(\lambda z. \text{if } z = a \text{ then deriv } g \ a \ \text{else } f \ z - g \ a / (z - a)) = f$
by *auto*
show $?thesis$
using *pole_theorem_analytic* [*OF* $g \ eq$] **by** *simp*
qed

lemma *pole_theorem_analytic_open_superset*:

assumes $g: g \text{ analytic_on } S$ **and** $S \subseteq T$ *open* T
and $eq: \bigwedge z. z \in T - \{a\} \implies g \ z = (z - a) * f \ z$
shows $(\lambda z. \text{if } z = a \text{ then deriv } g \ a \ \text{else } f \ z - g \ a / (z - a)) \text{ analytic_on } S$
proof (*rule pole_theorem_analytic* [*OF* g])
fix z
assume $z \in S$
then obtain e **where** $0 < e$ **and** $e: \text{ball } z \ e \subseteq T$
using *assms openE* **by** *blast*
then show $\exists d > 0. \forall w \in \text{ball } z \ d - \{a\}. g \ w = (w - a) * f \ w$
using eq **by** *auto*
qed

lemma *pole_theorem_analytic_open_superset_0*:

assumes $g: g \text{ analytic_on } S$ $S \subseteq T$ *open* T $\bigwedge z. z \in T - \{a\} \implies g \ z = (z - a) * f \ z$
and [*simp*]: $f \ a = \text{deriv } g \ a \ g \ a = 0$
shows $f \text{ analytic_on } S$
proof –

```

have [simp]: (λz. if z = a then deriv g a else f z - g a / (z-a)) = f
  by auto
have (λz. if z = a then deriv g a else f z - g a / (z-a)) analytic_on S
  by (rule pole_theorem_analytic_open_superset [OF g])
then show ?thesis by simp
qed

```

4.13 General, homology form of Cauchy's theorem

Proof is based on Dixon's, as presented in Lang's "Complex Analysis" book (page 147).

lemma *contour_integral_continuous_on_linepath_2D*:

assumes *open U* **and** *cont_dw*: $\bigwedge w. w \in U \implies F w$ *contour_integrable_on* (*linepath a b*)

and *cond_uu*: *continuous_on* ($U \times U$) ($\lambda(x,y). F x y$)

and *abu*: *closed_segment a b* $\subseteq U$

shows *continuous_on U* ($\lambda w. \text{contour_integral } (\text{linepath } a \ b) \ (F \ w)$)

proof –

have $\exists d > 0. \forall x' \in U. \text{dist } x' \ w < d \implies$

$$\begin{aligned} & \text{dist } (\text{contour_integral } (\text{linepath } a \ b) \ (F \ x')) \\ & \quad (\text{contour_integral } (\text{linepath } a \ b) \ (F \ w)) \leq \varepsilon \end{aligned}$$

if $w \in U$ $0 < \varepsilon$ $a \neq b$ **for** $w \ \varepsilon$

proof –

obtain δ **where** $\delta > 0$ **and** δ : *cball w δ* $\subseteq U$ **using** *open_contains_cball* $\langle \text{open } U \rangle$ $\langle w \in U \rangle$ **by** *force*

let $?TZ = \text{cball } w \ \delta \times \text{closed_segment } a \ b$

have *uniformly_continuous_on* $?TZ$ ($\lambda(x,y). F x y$)

by (*metis Sigma_mono δ abu compact_Times compact_cball compact_segment compact_uniformly_continuous*

cond_uu continuous_on_subset)

then obtain η **where** $\eta > 0$

and η : $\bigwedge x \ x'. \llbracket x \in ?TZ; x' \in ?TZ; \text{dist } x' \ x < \eta \rrbracket \implies$
 $\text{dist } ((\lambda(x,y). F x y) \ x') \ ((\lambda(x,y). F x y) \ x) < \varepsilon / \text{norm}(b-a)$

using $\langle 0 < \varepsilon \rangle$ $\langle a \neq b \rangle$

by (*auto elim: uniformly_continuous_onE* [**where** $\varepsilon = \varepsilon / \text{norm}(b-a)$])

have η : $\llbracket \text{norm } (w - x1) \leq \delta; x2 \in \text{closed_segment } a \ b;$

$\text{norm } (w - x1') \leq \delta; x2' \in \text{closed_segment } a \ b; \text{norm } ((x1', x2') - (x1, x2)) < \eta \rrbracket$

$$\implies \text{norm } (F \ x1' \ x2' - F \ x1 \ x2) \leq \varepsilon / \text{cmod } (b-a)$$

for $x1 \ x2 \ x1' \ x2'$

using η [*of* $(x1, x2) \ (x1', x2')$] **by** (*force simp: dist_norm*)

have *le_ee*: $\text{cmod } (\text{contour_integral } (\text{linepath } a \ b) \ (\lambda x. F \ x' \ x - F \ w \ x)) \leq \varepsilon$

if $x' \in U$ $\text{cmod } (x' - w) < \delta$ $\text{cmod } (x' - w) < \eta$ **for** x'

proof –

have $(\lambda x. F \ x' \ x - F \ w \ x)$ *contour_integrable_on* *linepath a b*

by (*simp add:* $\langle w \in U \rangle$ *cont_dw contour_integrable_diff that*)

then have $\text{cmod } (\text{contour_integral } (\text{linepath } a \ b) \ (\lambda x. F \ x' \ x - F \ w \ x)) \leq$

$\varepsilon / \text{norm}(b-a) * \text{norm}(b-a)$

using *has_contour_integral_bound_linepath* [*OF has_contour_integral_integral*

```

_ η]
  using ⟨0 < ε⟩ ⟨0 < δ⟩ that by (force simp: norm_minus_commute)
  also have ... = ε using ⟨a ≠ b⟩ by simp
  finally show ?thesis .
qed
show ?thesis
  using ⟨0 < δ⟩ ⟨0 < η⟩ ⟨w ∈ U⟩
  apply (intro exI[where x=min δ η])
  by (auto simp: dist_norm contour_integral_diff [OF cont_dw cont_dw,
symmetric] intro: le_ee)
  qed
  then show ?thesis
  by (metis (no_types, lifting) continuous_onI continuous_on_iff
contour_integral_trivial dist_self)
qed

```

This version has *polynomial_function* γ as an additional assumption.

lemma *Cauchy_integral_formula_global_weak*:

```

  assumes open U and holf: f holomorphic_on U
  and z: z ∈ U and γ: polynomial_function γ
  and pasz: path_image γ ⊆ U - {z} and loop: pathfinish γ = pathstart γ
  and zero: ∧w. w ∉ U ⇒ winding_number γ w = 0
  shows ((λw. f w / (w-z)) has_contour_integral (2*pi * i * winding_number
γ z * f z)) γ
proof -
  obtain γ' where pfγ': polynomial_function γ' and γ': ∧x. (γ has_vector_derivative
(γ' x)) (at x)
  using has_vector_derivative_polynomial_function [OF γ] by blast
  then have bounded(path_image γ')
  by (simp add: path_image_def compact_imp_bounded compact_continuous_image
continuous_on_polynomial_function)
  then obtain B where B>0 and B: ∧x. x ∈ path_image γ' ⇒ norm x ≤ B
  using bounded_pos by force
  define d where [abs_def]: d z w = (if w = z then deriv f z else (f w - f z)/(w-z))
for z w
  define v where v = {w. w ∉ path_image γ ∧ winding_number γ w = 0}
  have path γ valid_path γ using γ
  by (auto simp: path_polynomial_function valid_path_polynomial_function)
  then have ov: open v
  by (simp add: v_def open_winding_number_levelsets loop)
  have uv_Un: U ∪ v = UNIV
  using pasz zero by (auto simp: v_def)
  have conf: continuous_on U f
  by (metis holf holomorphic_on_imp_continuous_on)
  have hol_d: (d y) holomorphic_on U if y ∈ U for y
proof -
  have *: (λc. if c = y then deriv f y else (f c - f y) / (c-y)) holomorphic_on U
  by (simp add: holf pole_lemma_open ⟨open U⟩)
  then have isCont (λx. if x = y then deriv f y else (f x - f y) / (x-y)) y

```

```

    using at_within_open field_differentiable_imp_continuous_at holomor-
    phic_on_def that ‹open U› by fastforce
    then have continuous_on U (d y)
    using * d_def holomorphic_on_imp_continuous_on by auto
    moreover have d y holomorphic_on U - {y}
    proof -
    have (λw. if w = y then deriv f y else (f w - f y) / (w - y)) field_differentiable
    at w
    if w ∈ U - {y} for w
    proof (rule field_differentiable_transform_within)
    show (λw. (f w - f y) / (w - y)) field_differentiable at w
    using that ‹open U› holf
    by (auto intro!: holomorphic_on_imp_differentiable_at derivative_intros)
    show dist w y > 0
    using that by auto
    qed (auto simp: dist_commute)
    then show ?thesis
    unfolding field_differentiable_def by (simp add: d_def holomorphic_on_open
    ‹open U› open_delete)
    qed
    ultimately show ?thesis
    by (rule no_isolated_singularity) (auto simp: ‹open U›)
    qed
    have cint_fxy: (λx. (f x - f y) / (x - y)) contour_integrable_on γ if y ∉
    path_image γ for y
    proof (rule contour_integrable_holomorphic_simple [where S = U - {y}])
    show (λx. (f x - f y) / (x - y)) holomorphic_on U - {y}
    by (force intro: holomorphic_intros holomorphic_on_subset [OF holf])
    show path_image γ ⊆ U - {y}
    using pasz that by blast
    qed (auto simp: ‹open U› open_delete ‹valid_path γ›)
    define h where
    h z = (if z ∈ U then contour_integral γ (d z) else contour_integral γ (λw. f
    w / (w - z))) for z
    have U: ((d z) has_contour_integral h z) γ if z ∈ U for z
    proof -
    have d z holomorphic_on U
    by (simp add: hol_d that)
    with that show ?thesis
    by (metis Diff_subset ‹valid_path γ› ‹open U› contour_integrable_holomorphic_simple
    h_def
    has_contour_integral_integral pasz subset_trans)
    qed
    have V: ((λw. f w / (w - z)) has_contour_integral h z) γ if z: z ∈ v for z
    proof -
    have 0: 0 = (f z) * 2 * of_real (2 * pi) * i * winding_number γ z
    using v_def z by auto
    then have ((λx. 1 / (x - z)) has_contour_integral 0) γ
    using z v_def has_contour_integral_winding_number [OF ‹valid_path γ›]

```

```

by fastforce
  then have (( $\lambda x. f z * (1 / (x-z))$ ) has_contour_integral 0)  $\gamma$ 
    using has_contour_integral_lmul by fastforce
  then have (( $\lambda x. f z / (x-z)$ ) has_contour_integral 0)  $\gamma$ 
    by (simp add: field_split_simps)
  moreover have (( $\lambda x. (f x - f z) / (x-z)$ ) has_contour_integral contour_integral
 $\gamma (d z)$ )  $\gamma$ 
    by (metis (no_types, lifting) z cint_fxy contour_integral_eq d_def has_contour_integral_integral
mem_Collect_eq v_def)
  ultimately have *: (( $\lambda x. f z / (x-z) + (f x - f z) / (x-z)$ ) has_contour_integral
(0 + contour_integral  $\gamma (d z)$ ))  $\gamma$ 
    by (rule has_contour_integral_add)
  have (( $\lambda w. f w / (w-z)$ ) has_contour_integral contour_integral  $\gamma (d z)$ )  $\gamma$ 
    if  $z \in U$ 
    using * by (auto simp: divide_simps has_contour_integral_eq)
  moreover have (( $\lambda w. f w / (w-z)$ ) has_contour_integral contour_integral  $\gamma$ 
( $\lambda w. f w / (w-z)$ ))  $\gamma$ 
    if  $z \notin U$ 
  proof (rule has_contour_integral_integral [OF contour_integrable_holomorphic_simple
[where  $S=U$ ]])
    show ( $\lambda w. f w / (w-z)$ ) holomorphic_on U
      by (rule holomorphic_intros assms | use that in force)+
    qed (use ⟨open U⟩ pasz ⟨valid_path  $\gamma$ ⟩ in auto)
    ultimately show ?thesis
      using z by (simp add: h_def)
    qed
  have znot:  $z \notin \text{path\_image } \gamma$ 
    using pasz by blast
  obtain  $d0$  where  $d0 > 0$  and  $d0$ :  $\bigwedge x y. x \in \text{path\_image } \gamma \implies y \in - U \implies d0$ 
 $\leq \text{dist } x y$ 
    using separate_compact_closed [of path_image  $\gamma - U$ ] pasz ⟨open U⟩ ⟨path  $\gamma$ ⟩
compact_path_image
    by blast
  obtain  $dd$  where  $0 < dd$  and  $dd$ :  $\{y + k \mid y k. y \in \text{path\_image } \gamma \wedge k \in \text{ball } 0$ 
 $dd\} \subseteq U$ 
  proof
    show  $0 < d0 / 2$  using ⟨ $0 < d0$ ⟩ by auto
  qed (use ⟨ $0 < d0$ ⟩  $d0$  in ⟨force simp: dist_norm⟩)
  define T where  $T \equiv \{y + k \mid y k. y \in \text{path\_image } \gamma \wedge k \in \text{cball } 0 (dd / 2)\}$ 
  have  $\bigwedge x x'. \llbracket x \in \text{path\_image } \gamma; \text{dist } x x' * 2 < dd \rrbracket$ 
 $\implies \exists y k. x' = y + k \wedge y \in \text{path\_image } \gamma \wedge \text{dist } 0 k * 2 \leq dd$ 
    by (metis add.commute diff_add_cancel dist_0_norm dist_commute dist_norm
less_eq_real_def)
  then have subt: path_image  $\gamma \subseteq \text{interior } T$ 
    using ⟨ $0 < dd$ ⟩
    apply (clarsimp simp add: mem_interior T_def)
    apply (rule_tac  $x=dd/2$  in exI, auto)
    done
  have compact T

```

```

  unfolding T_def
  using ‹valid_path  $\gamma$ › compact_cball compact_sums compact_valid_path_image
by blast
  have T:  $T \subseteq U$ 
  unfolding T_def using ‹0 < dd› dd by fastforce
  obtain L where L>0
    and L:  $\bigwedge f B. \llbracket f \text{ holomorphic\_on interior } T; \bigwedge z. z \in \text{interior } T \implies cmod$ 
    ( $f z$ )  $\leq B$   $\implies$ 
       $cmod (\text{contour\_integral } \gamma f) \leq L * B$ 
  using contour_integral_bound_exists [OF open_interior ‹valid_path  $\gamma$ › subt]
  by blast
  have bounded(f ‘ T)
  by (meson ‹compact T› compact_continuous_image compact_imp_bounded
conf continuous_on_subset T)
  then obtain D where D>0 and D:  $\bigwedge x. x \in T \implies \text{norm } (f x) \leq D$ 
  by (auto simp: bounded_pos)
  obtain C where C>0 and C:  $\bigwedge x. x \in T \implies \text{norm } x \leq C$ 
  using ‹compact T› bounded_pos compact_imp_bounded by force
  have dist (h y)  $0 \leq e$  if  $0 < e$  and le:  $D * L / e + C \leq cmod y$  for  $e y$ 
  proof -
  have  $D * L / e > 0$  using ‹D>0› ‹L>0› ‹e>0› by simp
  with le have ybig:  $\text{norm } y > C$  by force
  with C have  $y \notin T$  by force
  then have ynot:  $y \notin \text{path\_image } \gamma$ 
  using subt_interior_subset by blast
  have [simp]:  $\text{winding\_number } \gamma y = 0$ 
  proof (rule winding_number_zero_outside)
  show  $\text{path\_image } \gamma \subseteq \text{cball } 0 C$ 
  by (meson C interior_subset mem_cball_0 subset_eq subt)
  qed (use ybig loop ‹path  $\gamma$ › in auto)
  have [simp]:  $h y = \text{contour\_integral } \gamma (\lambda w. f w / (w - y))$ 
  by (rule contour_integral_unique [symmetric]) (simp add: v_def ynot V)
  have holint:  $(\lambda w. f w / (w - y)) \text{ holomorphic\_on interior } T$ 
  proof (intro holomorphic_intros)
  show  $f \text{ holomorphic\_on interior } T$ 
  using holf holomorphic_on_subset interior_subset T by blast
  qed (use ‹ $y \notin T$ › interior_subset in auto)
  have leD:  $cmod (f z / (z - y)) \leq D * (e / L / D)$  if  $z: z \in \text{interior } T$  for  $z$ 
  proof -
  have  $D * L / e + cmod z \leq cmod y$ 
  using le C [of z] z using interior_subset by force
  then have DL2:  $D * L / e \leq cmod (z - y)$ 
  using norm_triangle_ineq2 [of y z] by (simp add: norm_minus_commute)
  have  $cmod (f z / (z - y)) = cmod (f z) * \text{inverse } (cmod (z - y))$ 
  by (simp add: norm_mult norm_inverse Fields.field_class.field_divide_inverse)
  also have  $\dots \leq D * (e / L / D)$ 
  proof (rule mult_mono)
  show  $cmod (f z) \leq D$ 
  using D interior_subset z by blast

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    show inverse (cmod (z-y)) ≤ e / L / D D ≥ 0
    using ‹L>0› ‹e>0› ‹D>0› DL2 by (auto simp: norm_divide field_split_simps)
  qed auto
  finally show ?thesis .
qed
have dist (h y) 0 = cmod (contour_integral γ (λw. f w / (w-y)))
  by (simp add: dist_norm)
also have ... ≤ L * (D * (e / L / D))
  by (rule L [OF holint leD])
also have ... = e
  using ‹L>0› ‹0 < D› by auto
finally show ?thesis .
qed
then have (h → 0) at_infinity
  by (meson Lim_at_infinityI)
moreover have h holomorphic_on UNIV
proof -
  have con_ff: continuous (at (x,z)) (λ(x,y). (f y - f x) / (y-x))
    if x ∈ U z ∈ U x ≠ z for x z
  using that conf
  apply (simp add: split_def continuous_on_eq_continuous_at ‹open U›)
  apply (simp | rule continuous_intros continuous_within_compose2 [where
g=f])+
  done
  have con_fstsnd: continuous_on UNIV (λx. (fst x - snd x) :: complex)
    by (rule continuous_intros)+
  have open_uu_Id: open (U × U - Id)
proof (rule open_Diff)
  show open (U × U)
    by (simp add: open_Times ‹open U›)
  show closed (Id :: complex rel)
  using continuous_closed_preimage_constant [OF con_fstsnd closed_UNIV,
of 0]
  by (auto simp: Id_fstsnd_eq algebra_simps)
qed
have con_derf: continuous (at z) (deriv f) if z ∈ U for z
  by (meson analytic_at analytic_at_imp_isCont assms(1) holf holomor-
phic_deriv that)
  have tendsto_f': ((λ(x,y). if y = x then deriv f (x)
    else (f (y) - f (x)) / (y-x)) → deriv f x)
    (at (x, x) within U × U) if x ∈ U for x
proof (rule Lim_withinI)
  fix e::real assume 0 < e
  obtain k1 where k1>0 and k1: ∧x'. norm (x' - x) ≤ k1 ⇒ norm (deriv
f x' - deriv f x) < e
  using ‹0 < e› continuous_within_E [OF con_derf [OF ‹x ∈ U›]]
  by (metis UNIV_I dist_norm)
  obtain k2 where k2>0 and k2: ball x k2 ⊆ U
  by (blast intro: openE [OF ‹open U›] ‹x ∈ U›)

```



```

have neq: norm ((f z' - f x') / (z' - x') - deriv f x) ≤ e
      if z' ≠ x' and less_k1: norm (x'-x, z'-x) < k1 and less_k2:
norm (x'-x, z'-x) < k2
      for x' z'
proof -
  have cs_less: w ∈ closed_segment x' z' ⇒ cmod (w-x) ≤ norm (x'-x,
z'-x) for w
    using segment_furthest_le [of w x' z' x]
    by (metis (no_types) dist_commute dist_norm norm_fst_le norm_snd_le
order_trans)
  have derf_le: w ∈ closed_segment x' z' ⇒ z' ≠ x' ⇒ cmod (deriv f w -
deriv f x) ≤ e for w
    by (blast intro: cs_less less_k1 k1 [unfolded divide_const_simps dist_norm]
less_imp_le le_less_trans)
  have f_has_der: ∧x. x ∈ U ⇒ (f has_field_derivative deriv f x) (at x
within U)
    by (metis DERIV_deriv_iff_field_differentiable at_within_open holf
holomorphic_on_def ‹open U›)
  have closed_segment x' z' ⊆ U
    by (rule order_trans [OF _ k2]) (simp add: cs_less le_less_trans [OF _
less_k2] dist_complex_def norm_minus_commute subset_iff)
  then have cint_derf: (deriv f has_contour_integral f z' - f x') (linepath x'
z')
    using contour_integral_primitive [OF f_has_der valid_path_linepath]
pasz by simp
  then have *: ((λx. deriv f x / (z' - x')) has_contour_integral (f z' - f x')
/ (z' - x')) (linepath x' z')
    by (rule has_contour_integral_div)
  have norm ((f z' - f x') / (z' - x') - deriv f x) ≤ e/norm(z' - x') *
norm(z' - x')
    apply (rule has_contour_integral_bound_linepath [OF has_contour_integral_diff
[OF *]])
    using has_contour_integral_div [where c = z' - x', OF has_contour_integral_const_linepath
[of deriv f x z' x']]
      ‹e > 0› ‹z' ≠ x'›
    apply (auto simp: norm_divide divide_simps derf_le)
    done
  also have ... ≤ e using ‹0 < e› by simp
  finally show ?thesis .
qed
show ∃ d>0. ∀ xa∈U × U.
  0 < dist xa (x, x) ∧ dist xa (x, x) < d ⇒
  dist (case xa of (x, y) ⇒ if y = x then deriv f x else (f y - f x) /
(y-x)) (deriv f x) ≤ e
    apply (rule_tac x=min k1 k2 in exI)
    using ‹k1>0› ‹k2>0› ‹e>0›
    by (force simp: dist_norm neq intro: dual_order.strict_trans2 k1 less_imp_le
norm_fst_le)
qed

```

```

have con_pa_f: continuous_on (path_image  $\gamma$ ) f
by (meson holf holomorphic_on_imp_continuous_on holomorphic_on_subset
interior_subset subt T)
have le_B:  $\bigwedge T. T \in \{0..1\} \implies \text{cmod} (\text{vector\_derivative } \gamma \text{ (at } T)) \leq B$ 
using  $\gamma' B$  by (simp add: path_image_def vector_derivative_at rev_image_eqI)
have f_has_cint:  $\bigwedge w. w \in v - \text{path\_image } \gamma \implies ((\lambda u. f u / (u-w)) \hat{=} 1)$ 
has_contour_integral h w)  $\gamma$ 
by (simp add: V)
have  $\bigwedge x y. \llbracket x \in U; y \in U; y \neq x \rrbracket \implies (\lambda(x, y). d x y) - (x, y) \rightarrow (f y - f x) / (y - x)$ 
unfolding d_def
apply (rule Lim_transform_within_open [OF _ open_uu_Id, where  $f =$ 
 $(\lambda(x, y). (f y - f x) / (y-x))$ ])
using con_ff by (auto simp: continuous_within)
then have cond_uu: continuous_on  $(U \times U)$   $(\lambda(x, y). d x y)$ 
unfolding continuous_on_eq_continuous_within continuous_within_d_def
by (fastforce simp add: tendsto_f' intro: Lim_at_imp_Lim_at_within)
have hol_dw:  $(\lambda z. d z w)$  holomorphic_on U if  $w \in U$  for w
proof -
have continuous_on U  $((\lambda(x, y). d x y) \circ (\lambda z. (w, z)))$ 
by (rule continuous_on_compose continuous_intros continuous_on_subset
[OF cond_uu] | force intro: that)+
then have *: continuous_on U  $(\lambda z. \text{if } w = z \text{ then deriv } f z \text{ else } (f w - f z) / (w-z))$ 
by (rule rev_iffD1 [OF _ continuous_on_cong [OF refl]]) (simp add: d_def
field_simps)
have **:  $(\lambda z. \text{if } w = z \text{ then deriv } f z \text{ else } (f w - f z) / (w-z))$  field_differentiable
at x
if  $x \in U$   $x \neq w$  for x
proof (rule_tac  $f = \lambda x. (f w - f x) / (w-x)$  and  $d = \text{dist } x w$  in field_differentiable_transform_within)
show  $(\lambda x. (f w - f x) / (w-x))$  field_differentiable at x
using that  $\langle \text{open } U \rangle$ 
by (intro derivative_intros holomorphic_on_imp_differentiable_at [OF
holf]; force)
qed (use that  $\langle \text{open } U \rangle$  in  $\langle \text{auto simp: dist\_commute} \rangle$ )
show ?thesis
unfolding d_def
proof (rule no_isolated_singularity [OF * _  $\langle \text{open } U \rangle$ ])
show  $(\lambda z. \text{if } w = z \text{ then deriv } f z \text{ else } (f w - f z) / (w-z))$  holomorphic_on
 $U - \{w\}$ 
by (auto simp: field_differentiable_def [symmetric] holomorphic_on_open
open_Diff  $\langle \text{open } U \rangle$  **)
qed auto
qed
{ fix a b
assume abu: closed_segment a b  $\subseteq U$ 
have cont_cint_d: continuous_on U  $(\lambda w. \text{contour\_integral} (\text{linepath } a b)$ 
 $(\lambda z. d z w))$ 
proof (rule contour_integral_continuous_on_linepath_2D [OF  $\langle \text{open } U \rangle$  _
```

```

_ abu])
  show  $\bigwedge w. w \in U \implies (\lambda z. d z w) \text{ contour\_integrable\_on } (\text{linepath } a \ b)$ 
  by (metis abu hol_dw continuous_on_subset contour_integrable_continuous_linepath
holomorphic_on_imp_continuous_on)
  show continuous_on (U × U) ( $\lambda(x, y). d y x$ )
  by (auto intro: continuous_on_swap_args cond_uu)
qed
  have cont_cint_d $\gamma$ : continuous_on {0..1} (( $\lambda w. \text{ contour\_integral } (\text{linepath } a \ b) (\lambda z. d z w) \circ \gamma$ )
  by (metis Diff_subset ⟨path  $\gamma$ ⟩ cont_cint_d continuous_on_compose continuous_on_subset
pasz path_def path_image_def)
  have continuous_on {0..1} ( $\lambda x. \text{ vector\_derivative } \gamma \text{ (at } x)$ )
  using pf $\gamma'$  by (simp add: continuous_on_polynomial_function vector_derivative_at
[OF  $\gamma'$ ])
  then have cint_cint: ( $\lambda w. \text{ contour\_integral } (\text{linepath } a \ b) (\lambda z. d z w)$ )
contour_integrable_on  $\gamma$ 
  apply (simp add: contour_integrable_on)
  apply (rule integrable_continuous_real)
  by (rule continuous_on_mult [OF cont_cint_d $\gamma$  [unfolded o_def]])
  have contour_integral (linepath a b) h = contour_integral (linepath a b) ( $\lambda z. \text{ contour\_integral } \gamma \text{ (d } z)$ )
  using abu by (force simp: h_def intro: contour_integral_eq)
  also have ... = contour_integral  $\gamma$  ( $\lambda w. \text{ contour\_integral } (\text{linepath } a \ b) (\lambda z. d z w)$ )
  proof (rule contour_integral_swap)
  show continuous_on (path_image (linepath a b) × path_image  $\gamma$ ) ( $\lambda(y1, y2). d y1 y2$ )
  using abu pasz by (auto intro: continuous_on_subset [OF cond_uu])
  show continuous_on {0..1} ( $\lambda t. \text{ vector\_derivative } (\text{linepath } a \ b) \text{ (at } t)$ )
  by (auto intro!: continuous_intros)
  show continuous_on {0..1} ( $\lambda t. \text{ vector\_derivative } \gamma \text{ (at } t)$ )
  by (metis  $\gamma'$  continuous_on_eq path_def path_polynomial_function pf $\gamma'$ 
vector_derivative_at)
  qed (use ⟨valid_path  $\gamma$ ⟩ in auto)
  finally have cint_h_eq:
    contour_integral (linepath a b) h =
      contour_integral  $\gamma$  ( $\lambda w. \text{ contour\_integral } (\text{linepath } a \ b) (\lambda z. d z w)$ ) .
  note cint_cint cint_h_eq
} note cint_h = this
have conth_u: continuous_on U h
proof (simp add: continuous_on_sequentially, clarify)
fix a x
assume x:  $x \in U$  and au:  $\forall n. a \ n \in U$  and ax:  $a \longrightarrow x$ 
then have A1:  $\forall_F n$  in sequentially. d (a n) contour_integrable_on  $\gamma$ 
  by (meson U contour_integrable_on_def eventuallyI)
  obtain dd where dd>0 and dd: cball x dd  $\subseteq U$  using open_contains_cball
⟨open U⟩ x by force
  have A2: uniform_limit (path_image  $\gamma$ ) ( $\lambda n. d (a \ n)$ ) (d x) sequentially

```

```

unfolding uniform_limit_iff dist_norm
proof clarify
  fix ee::real
  assume 0 < ee
  show  $\forall_F n$  in sequentially.  $\forall \xi \in \text{path\_image } \gamma. \text{cmod } (d (a n) \xi - d x \xi) <$ 
ee
  proof -
    let ?ddpa =  $\{(w,z) \mid w z. w \in \text{cball } x \text{ dd} \wedge z \in \text{path\_image } \gamma\}$ 
    have uniformly_continuous_on ?ddpa  $(\lambda(x,y). d x y)$ 
    proof (rule compact_uniformly_continuous [OF continuous_on_subset [OF
cond_uu]])
      show compact  $\{(w, z) \mid w z. w \in \text{cball } x \text{ dd} \wedge z \in \text{path\_image } \gamma\}$ 
      using <valid_path  $\gamma$ >
      by (auto simp: compact_Times compact_valid_path_image simp del:
mem_cball)
    qed (use dd pasz in auto)
    then obtain kk where kk > 0
    and kk:  $\bigwedge x x'. \llbracket x \in ?ddpa; x' \in ?ddpa; \text{dist } x' x < kk \rrbracket \implies$ 
 $\text{dist } ((\lambda(x,y). d x y) x') ((\lambda(x,y). d x y) x) < ee$ 
    by (rule uniformly_continuous_onE [where  $\varepsilon = ee$ ]) (use <0 < ee> in
auto)
    have kk:  $\llbracket \text{norm } (w-x) \leq dd; z \in \text{path\_image } \gamma; \text{norm } ((w, z) - (x, z))$ 
< kk  $\rrbracket \implies \text{norm } (d w z - d x z) < ee$ 
    for w z
    using <dd > 0> kk [of (x,z) (w,z)] by (force simp: norm_minus_commute
dist_norm)
    obtain no where  $\forall n \geq no. \text{dist } (a n) x < \min dd kk$ 
    using ax unfolding lim_sequentially
    by (meson <0 < dd> <0 < kk> min_less_iff_conj)
    then show ?thesis
    using <dd > 0> <kk > 0> by (fastforce simp: eventually_sequentially kk
dist_norm)
    qed
  qed
  have  $(\lambda n. \text{contour\_integral } \gamma (d (a n))) \longrightarrow \text{contour\_integral } \gamma (d x)$ 
by (rule contour_integral_uniform_limit [OF A1 A2 le_B]) (auto simp:
<valid_path  $\gamma$ >)
  then have tendsto_hx:  $(\lambda n. \text{contour\_integral } \gamma (d (a n))) \longrightarrow h x$ 
by (simp add: h_def x)
  then show  $(h \circ a) \longrightarrow h x$ 
by (simp add: h_def x au o_def)
qed
show ?thesis
proof (simp add: holomorphic_on_open field_differentiable_def [symmetric],
clarify)
  fix z0
  consider z0  $\in v \mid z0 \in U$  using uv_Un by blast
  then show h field_differentiable at z0
proof cases

```

```

    assume z0 ∈ v then show ?thesis
      using Cauchy_next_derivative [OF con_pa_f le_B f_has_cint __ ov] V
f_has_cint ⟨valid_path γ⟩
      by (auto simp: field_differentiable_def v_def)
    next
      assume z0 ∈ U then
        obtain e where e > 0 and e: ball z0 e ⊆ U by (blast intro: openE [OF
⟨open U⟩])
        have *: contour_integral (linepath a b) h + contour_integral (linepath b c)
h + contour_integral (linepath c a) h = 0
          if abc_subset: convex_hull {a, b, c} ⊆ ball z0 e for a b c
        proof -
          have *: ∧x1 x2 z. z ∈ U ⇒ closed_segment x1 x2 ⊆ U ⇒ (λw. d w z)
contour_integrable_on linepath x1 x2
            using hol_dw holomorphic_on_imp_continuous_on ⟨open U⟩
            by (auto intro!: contour_integrable_holomorphic_simple)
          have abc: closed_segment a b ⊆ U closed_segment b c ⊆ U closed_segment
c a ⊆ U
            using that e segments_subset_convex_hull by fastforce+
          have eq0: ∧w. w ∈ U ⇒ contour_integral (linepath a b +++ linepath b
c +++ linepath c a) (λz. d z w) = 0
            proof (rule contour_integral_unique [OF Cauchy_theorem_triangle])
              show ∧w. w ∈ U ⇒ (λz. d z w) holomorphic_on convex_hull {a, b, c}
                using e abc_subset by (auto intro: holomorphic_on_subset [OF
hol_dw])
            qed
          have ∧z. z ∈ path_image γ ⇒
            contour_integral (linepath a b) (λx. d x z) +
            (contour_integral (linepath b c) (λx. d x z) +
            contour_integral (linepath c a) (λx. d x z)) = 0
            using abc pasz U * eq0 by auto
          then show ?thesis
            by (simp add: contour_integral_eq_0 cint_h abc contour_integrable_add
contour_integral_add [symmetric] add_ac)
          qed
          show ?thesis
            using e ⟨e > 0⟩
            by (auto intro!: holomorphic_on_imp_differentiable_at [OF __ open_ball]
analytic_imp_holomorphic
            Morera_triangle_continuous_on_subset [OF conthu] *)
          qed
        qed
      ultimately have [simp]: h z = 0 for z
        by (meson Liouville_weak)
      have ((λw. 1 / (w - z)) has_contour_integral complex_of_real (2 * pi) * i *
winding_number γ z) γ
        by (rule has_contour_integral_winding_number [OF ⟨valid_path γ⟩ znot])
      then have ((λw. f z * (1 / (w - z))) has_contour_integral complex_of_real (2

```

```

* pi) * i * winding_number  $\gamma$  z * f z)  $\gamma$ 
  by (metis mult.commute has_contour_integral_lm1)
  then have 1: (( $\lambda w. f z / (w-z)$ ) has_contour_integral complex_of_real (2 *
pi) * i * winding_number  $\gamma$  z * f z)  $\gamma$ 
  by (simp add: field_split_simps)
  moreover have 2: (( $\lambda w. (f w - f z) / (w-z)$ ) has_contour_integral 0)  $\gamma$ 
  using U [OF z] pasz d_def by (force elim: has_contour_integral_eq [where
g =  $\lambda w. (f w - f z)/(w-z)$ ])
  show ?thesis
  using has_contour_integral_add [OF 1 2] by (simp add: diff_divide_distrib)
qed

```

theorem *Cauchy_integral_formula_global*:

```

  assumes S: open S and holf: f holomorphic_on S
    and z: z  $\in$  S and vpg: valid_path  $\gamma$ 
    and pasz: path_image  $\gamma \subseteq S - \{z\}$  and loop: pathfinish  $\gamma =$  pathstart  $\gamma$ 
    and zero:  $\bigwedge w. w \notin S \implies$  winding_number  $\gamma$  w = 0
  shows (( $\lambda w. f w / (w-z)$ ) has_contour_integral (2*pi * i * winding_number
 $\gamma$  z * f z))  $\gamma$ 
  proof -
    have path  $\gamma$  using vpg by (blast intro: valid_path_imp_path)
    have hols: ( $\lambda w. f w / (w-z)$ ) holomorphic_on S - {z} ( $\lambda w. 1 / (w-z)$ ) holo-
morphic_on S - {z}
    by (rule holomorphic_intros holomorphic_on_subset [OF holf] | force)+
    then have cint_fw: ( $\lambda w. f w / (w-z)$ ) contour_integrable_on  $\gamma$ 
    by (meson contour_integrable_holomorphic_simple holomorphic_on_imp_continuous_on
open_delete S vpg pasz)
    obtain d where d>0
    and d:  $\bigwedge g h. \llbracket$ valid_path g; valid_path h;  $\forall t \in \{0..1\}. cmod (g t - \gamma t) < d$ 
 $\wedge cmod (h t - \gamma t) < d;$ 
    pathstart h = pathstart g  $\wedge$  pathfinish h = pathfinish g $\rrbracket$ 
     $\implies$  path_image h  $\subseteq S - \{z\} \wedge (\forall f. f$  holomorphic_on S - {z}
 $\longrightarrow$  contour_integral h f = contour_integral g f)
    using contour_integral_nearby_ends [OF _  $\langle$ path  $\gamma \rangle$  pasz] S by (simp add:
open_Diff) metis
    obtain p where polyp: polynomial_function p
    and ps: pathstart p = pathstart  $\gamma$  and pf: pathfinish p = pathfinish  $\gamma$ 
  and led:  $\forall t \in \{0..1\}. cmod (p t - \gamma t) < d$ 
    using path_approx_polynomial_function [OF  $\langle$ path  $\gamma \rangle$   $\langle$ d > 0 $\rangle$ ] by metis
    then have ploop: pathfinish p = pathstart p using loop by auto
    have vpp: valid_path p using polyp valid_path_polynomial_function by blast
    have [simp]: z  $\notin$  path_image  $\gamma$  using pasz by blast
    have paps: path_image p  $\subseteq S - \{z\}$  and cint_eq: ( $\bigwedge f. f$  holomorphic_on S -
{z}  $\implies$  contour_integral p f = contour_integral  $\gamma$  f)
    using pf ps led d [OF vpg vpp]  $\langle$ d > 0 $\rangle$  by auto
    have wn_eq: winding_number p z = winding_number  $\gamma$  z
    using vpp paps
    by (simp add: subset_Diff_insert vpg valid_path_polynomial_function wind-
ing_number_valid_path_cint_eq hols)

```

```

have winding_number p w = winding_number  $\gamma$  w if  $w \notin S$  for w
proof -
  have hol:  $(\lambda v. 1 / (v-w))$  holomorphic_on  $S - \{z\}$ 
  using that by (force intro: holomorphic_intros holomorphic_on_subset [OF
hol])
  have  $w \notin \text{path\_image } p$   $w \notin \text{path\_image } \gamma$  using paps pasz that by auto
  then show ?thesis
  using vpp vpg by (simp add: subset_Diff_insert valid_path_polynomial_function
winding_number_valid_path cint_eq [OF hol])
qed
then have wn0:  $\bigwedge w. w \notin S \implies \text{winding\_number } p \ w = 0$ 
  by (simp add: zero)
show ?thesis
  using Cauchy_integral_formula_global_weak [OF S holf z polyp paps ploop
wn0] hols
  by (metis wn_eq cint_eq has_contour_integral_eqpath cint_fw cint_eq)
qed

```

theorem Cauchy_theorem_global:

```

assumes S: open S and hol: f holomorphic_on S
  and vpg: valid_path  $\gamma$  and loop: pathfinish  $\gamma = \text{pathstart } \gamma$ 
  and pas: path_image  $\gamma \subseteq S$ 
  and zero:  $\bigwedge w. w \notin S \implies \text{winding\_number } \gamma \ w = 0$ 
shows (f has_contour_integral 0)  $\gamma$ 
proof -
  have path_image  $\gamma \neq S$ 
  by (metis compact_valid_path_image vpg compact_open path_image_nonempty
S)
  then obtain z where  $z \in S$  and znot:  $z \notin \text{path\_image } \gamma$  and pasz: path_image
 $\gamma \subseteq S - \{z\}$ 
  using pas by blast
  have hol:  $(\lambda w. (w-z) * f \ w)$  holomorphic_on S
  by (rule holomorphic_intros holf)+
  show ?thesis
  using Cauchy_integral_formula_global [OF S hol  $\langle z \in S \rangle$  vpg pasz loop zero]
  by (auto simp: znot elim!: has_contour_integral_eq)
qed

```

corollary Cauchy_theorem_global_outside:

```

assumes open S f holomorphic_on S valid_path  $\gamma$  pathfinish  $\gamma = \text{pathstart } \gamma$ 
path_image  $\gamma \subseteq S$ 
   $\bigwedge w. w \notin S \implies w \in \text{outside}(\text{path\_image } \gamma)$ 
shows (f has_contour_integral 0)  $\gamma$ 
by (metis Cauchy_theorem_global assms winding_number_zero_in_outside valid_path_imp_path)

```

lemma simply_connected_imp_winding_number_zero:

```

assumes simply_connected S path g
  path_image g  $\subseteq S$  pathfinish g = pathstart g  $z \notin S$ 
shows winding_number g z = 0

```

proof –

have *hom*: *homotopic_loops* *S* *g* (*linepath* (*pathstart* *g*) (*pathstart* *g*))
by (*meson* *assms* *homotopic_paths_imp_homotopic_loops* *pathfinish_linepath*
simply_connected_eq_contractible_path)
then have *homotopic_paths* ($-\{z\}$) *g* (*linepath* (*pathstart* *g*) (*pathstart* *g*))
by (*meson* $\langle z \notin S \rangle$ *homotopic_loops_imp_homotopic_paths_null* *homotopic_paths_subset*
subset_Cmpl_singleton)
then have *winding_number* *g* *z* = *winding_number*(*linepath* (*pathstart* *g*) (*pathstart*
g)) *z*
by (*rule* *winding_number_homotopic_paths*)
also have $\dots = 0$
using *assms* **by** (*force* *intro*: *winding_number_trivial*)
finally show *?thesis* .
qed

lemma *Cauchy_theorem_simply_connected*:

assumes *open* *S* *simply_connected* *S* *f* *holomorphic_on* *S* *valid_path* *g*
path_image *g* \subseteq *S* *pathfinish* *g* = *pathstart* *g*
shows (*f* *has_contour_integral* 0) *g*
by (*meson* *assms* *Cauchy_theorem_global_simply_connected_imp_winding_number_zero*
valid_path_imp_path)

proposition *holomorphic_logarithm_exists*:

assumes *A*: *convex* *A* *open* *A*
and *f*: *f* *holomorphic_on* *A* $\bigwedge x. x \in A \implies f\ x \neq 0$
and *z0*: *z0* \in *A*
obtains *g* **where** *g* *holomorphic_on* *A* **and** $\bigwedge x. x \in A \implies \exp (g\ x) = f\ x$
proof –
note *f'* = *holomorphic_derivI* [*OF* *f*(1) *A*(2)]
obtain *g* **where** *g*: $\bigwedge x. x \in A \implies (g\ \text{has_field_derivative}\ \text{deriv}\ f\ x / f\ x)\ (\text{at}\ x)$
proof (*rule* *holomorphic_convex_primitive'* [*OF* *A*])
show $(\lambda x. \text{deriv}\ f\ x / f\ x)$ *holomorphic_on* *A*
by (*intro* *holomorphic_intros* *f* *A*)
qed (*auto simp*: *A* *at_within_open*[*of* $-\ A$])
define *h* **where** *h* = $(\lambda x. -g\ z0 + \ln (f\ z0) + g\ x)$
from *g* **and** *A* **have** *g_holo*: *g* *holomorphic_on* *A*
by (*auto simp*: *holomorphic_on_def* *at_within_open*[*of* $-\ A$] *field_differentiable_def*)
hence *h_holo*: *h* *holomorphic_on* *A*
by (*auto simp*: *h_def* *intro!*: *holomorphic_intros*)
note [*simp*] = *at_within_open*[*OF* $-\ \langle \text{open}\ A \rangle$]
have $\exists c. \forall x \in A. f\ x / \exp (h\ x) - 1 = c$
using $\langle \text{convex}\ A \rangle$ *z0* *f*
by (*force simp*: *h_def* *exp_diff_field_simps* *intro!*: *has_field_derivative_zero_constant*
derivative_eq_intros *g* *f'*)
then obtain *c* **where** *c*: $\bigwedge x. x \in A \implies f\ x / \exp (h\ x) - 1 = c$
by *blast*
from *c*[*OF* *z0*] **and** *z0* **and** *f* **have** *c* = 0
by (*simp add*: *h_def*)
with *c* **have** $\bigwedge x. x \in A \implies \exp (h\ x) = f\ x$ **by** *simp*

from that[OF h_holo this] show ?thesis .
qed

4.14 Cauchy's inequality and more versions of Liouville

lemma Cauchy_higher_deriv_bound:

```

  assumes holf: f holomorphic_on (ball z r)
    and contf: continuous_on (cball z r) f
    and fin :  $\bigwedge w. w \in \text{ball } z \ r \implies f \ w \in \text{ball } y \ B0$ 
    and 0 < r and 0 < n
  shows cmod ((deriv  $\hat{\sim}$  n) f z)  $\leq$  (fact n) * B0 / r $\hat{\sim}$ n
proof -
  have 0 < B0
  using <0 < r> fin [of z] by (metis ball_eq_empty ex_in_conv fin not_less)
  have le_B0: cmod (f w - y)  $\leq$  B0 if cmod (w-z)  $\leq$  r for w
  proof (rule continuous_on_closure_norm_le [of ball z r  $\lambda w. f \ w - y$ ], use <0
  < r> in simp_all)
    show continuous_on (cball z r) ( $\lambda w. f \ w - y$ )
    by (intro continuous_intros contf)
  show dist z w  $\leq$  r
  by (simp add: dist_commute dist_norm that)
  qed (use fin in <auto simp: dist_norm less_eq_real_def norm_minus_commute>)
  have (deriv  $\hat{\sim}$  n) f z = (deriv  $\hat{\sim}$  n) ( $\lambda w. f \ w$ ) z - (deriv  $\hat{\sim}$  n) ( $\lambda w. y$ ) z
  using <0 < n> by simp
  also have ... = (deriv  $\hat{\sim}$  n) ( $\lambda w. f \ w - y$ ) z
  using <0 < r> higher_deriv_diff holf by auto
  finally have (deriv  $\hat{\sim}$  n) f z = (deriv  $\hat{\sim}$  n) ( $\lambda w. f \ w - y$ ) z .
  have contf': continuous_on (cball z r) ( $\lambda u. f \ u - y$ )
  by (rule contf continuous_intros)+
  have holf': ( $\lambda u. (f \ u - y)$ ) holomorphic_on (ball z r)
  by (simp add: holf holomorphic_on_diff)
  define a where a = (2 * pi)/(fact n)
  have 0 < a by (simp add: a_def)
  have B0/r $\hat{\sim}$ (Suc n)*2 * pi * r = a*((fact n)*B0/r $\hat{\sim}$ n)
  using <0 < r> by (simp add: a_def field_split_simps)
  have der_dif: (deriv  $\hat{\sim}$  n) ( $\lambda w. f \ w - y$ ) z = (deriv  $\hat{\sim}$  n) f z
  using <0 < r> <0 < n>
  by (auto simp: higher_deriv_diff [OF holf holomorphic_on_const])
  have norm ((2 * of_real pi * i)/(fact n) * (deriv  $\hat{\sim}$  n) ( $\lambda w. f \ w - y$ ) z)
   $\leq$  (B0/r $\hat{\sim}$ (Suc n)) * (2 * pi * r)
  apply (rule has_contour_integral_bound_circlepath [of ( $\lambda u. (f \ u - y)/(u-z)$ ) $\hat{\sim}$ (Suc
  n)) _ z])
  using Cauchy_has_contour_integral_higher_derivative_circlepath [OF contf'
  holf]
  using <0 < B0> <0 < r>
  apply (auto simp: norm_divide norm_mult norm_power divide_simps le_B0)
  done
  then show ?thesis
  using <0 < r>

```

by (auto simp: norm_divide norm_mult norm_power field_simps der_diff_le_B0)
qed

lemma *Cauchy_inequality:*

assumes *holf*: f holomorphic_on (ball ξ r)
and *contf*: continuous_on (cball ξ r) f
and $0 < r$
and *nof*: $\bigwedge x. \text{norm}(\xi - x) = r \implies \text{norm}(f x) \leq B$
shows $\text{norm}((\text{deriv} \hat{\sim} n) f \xi) \leq (\text{fact } n) * B / r^n$

proof –

obtain x where $\text{norm}(\xi - x) = r$
by (metis $\langle 0 < r \rangle$ dist_norm_order_less_imp_le vector_choose_dist)
then have $0 \leq B$
by (metis *nof* norm_not_less_zero not_le_order_trans)
have $\xi \in \text{ball } \xi r$
using $\langle 0 < r \rangle$ by simp
then have $((\lambda u. f u / (u - \xi)^{\wedge \text{Suc } n}) \text{ has_contour_integral } (2 * \pi) * i / \text{fact } n * (\text{deriv} \hat{\sim} n) f \xi)$
(circlepath ξr)
by (rule *Cauchy_has_contour_integral_higher_derivative_circlepath* [OF *contf holf*])
have $\text{norm}((2 * \pi * i) / (\text{fact } n) * (\text{deriv} \hat{\sim} n) f \xi) \leq (B / r^{\wedge \text{Suc } n}) * (2 * \pi * r)$
proof (rule *has_contour_integral_bound_circlepath*)
have $\xi \in \text{ball } \xi r$
using $\langle 0 < r \rangle$ by simp
then show $((\lambda u. f u / (u - \xi)^{\wedge \text{Suc } n}) \text{ has_contour_integral } (2 * \pi) * i / \text{fact } n * (\text{deriv} \hat{\sim} n) f \xi)$
(circlepath ξr)
by (rule *Cauchy_has_contour_integral_higher_derivative_circlepath* [OF *contf holf*])
show $\bigwedge x. \text{cmod}(x - \xi) = r \implies \text{cmod}(f x / (x - \xi)^{\wedge \text{Suc } n}) \leq B / r^{\wedge \text{Suc } n}$
using $\langle 0 \leq B \rangle \langle 0 < r \rangle$
by (simp add: norm_divide norm_power *nof* frac_le norm_minus_commute del: power_Suc)
qed (use $\langle 0 \leq B \rangle \langle 0 < r \rangle$ in auto)
then show ?thesis using $\langle 0 < r \rangle$
by (simp add: norm_divide norm_mult field_simps)
qed

lemma *Liouville_polynomial:*

assumes *holf*: f holomorphic_on UNIV
and *nof*: $\bigwedge z. A \leq \text{norm } z \implies \text{norm}(f z) \leq B * \text{norm } z^{\wedge n}$
shows $f \xi = (\sum_{k \leq n}. (\text{deriv} \hat{\sim} k) f 0 / \text{fact } k * \xi^{\wedge k})$

proof (cases rule: le_less_linear [THEN disjE])

assume $B \leq 0$

then have $\bigwedge z. A \leq \text{norm } z \implies \text{norm}(f z) = 0$

by (metis *nof* less_le_trans zero_less_mult_iff neqE norm_not_less_zero)

```

norm_power not_le)
  then have f0: (f ⟶ 0) at_infinity
    using Lim_at_infinity by force
  then have [simp]: f = (λw. 0)
    using Liouville_weak [OF holf, of 0]
    by (simp add: eventually_at_infinity f0) meson
  show ?thesis by simp
next
  assume 0 < B
  have ((λk. (deriv ~ k) f 0 / (fact k) * (ξ - 0) ^ k) sums f ξ)
  proof (rule holomorphic_power_series [where r = norm ξ + 1])
    show f holomorphic_on ball 0 (cmod ξ + 1) ξ ∈ ball 0 (cmod ξ + 1)
      using holf holomorphic_on_subset by auto
  qed
  then have sumsf: ((λk. (deriv ~ k) f 0 / (fact k) * ξ ^ k) sums f ξ) by simp
  have (deriv ~ k) f 0 / fact k * ξ ^ k = 0 if k > n for k
  proof (cases (deriv ~ k) f 0 = 0)
    case True then show ?thesis by simp
  next
    case False
      define w where w = complex_of_real (fact k * B / cmod ((deriv ~ k) f 0)
+ (|A| + 1))
      have 1 ≤ abs (fact k * B / cmod ((deriv ~ k) f 0) + (|A| + 1))
        using ‹0 < B› by simp
      then have wge1: 1 ≤ norm w
        by (metis norm_of_real w_def)
      then have w ≠ 0 by auto
      have kB: 0 < fact k * B
        using ‹0 < B› by simp
      then have 0 ≤ fact k * B / cmod ((deriv ~ k) f 0)
        by simp
      then have wgeA: A ≤ cmod w
        by (simp only: w_def norm_of_real)
      have fact k * B / cmod ((deriv ~ k) f 0) < abs (fact k * B / cmod ((deriv ~
k) f 0) + (|A| + 1))
        using ‹0 < B› by simp
      then have wge: fact k * B / cmod ((deriv ~ k) f 0) < norm w
        by (metis norm_of_real w_def)
      then have fact k * B / norm w < cmod ((deriv ~ k) f 0)
        using False by (simp add: field_split_simps mult.commute split: if_split_asm)
      also have ... ≤ fact k * (B * norm w ^ n) / norm w ^ k
      proof (rule Cauchy_inequality)
        show f holomorphic_on ball 0 (cmod w)
          using holf holomorphic_on_subset by force
        show continuous_on (cball 0 (cmod w)) f
          using holf holomorphic_on_imp_continuous_on holomorphic_on_subset
      by blast
      show ∧x. cmod (0 - x) = cmod w ⟹ cmod (f x) ≤ B * cmod w ^ n
        by (metis nof wgeA dist_0_norm dist_norm)

```

```

qed (use ⟨w ≠ 0⟩ in auto)
also have ... = fact k * B / cmod w ^ (k-n)
  using ⟨k>n⟩ by (simp add: divide_simps flip: power_add)
finally have fact k * B / cmod w < fact k * B / cmod w ^ (k-n) .
then have 1 / cmod w < 1 / cmod w ^ (k-n)
  by (metis kB divide_inverse inverse_eq_divide mult_less_cancel_left_pos)
then have cmod w ^ (k-n) < cmod w
  by (smt (verit, best) ⟨w ≠ 0⟩ frac_le zero_less_norm_iff)
with self_le_power [OF wge1] show ?thesis
  by (meson diff_is_0_eq not_gr0 not_le that)
qed
then have (deriv ^^ (k + Suc n)) f 0 / fact (k + Suc n) * ξ ^ (k + Suc n) =
0 for k
  using not_less_eq by blast
then have (λi. (deriv ^^ (i + Suc n)) f 0 / fact (i + Suc n) * ξ ^ (i + Suc n))
sums 0
  by (rule sums_0)
with sums_split_initial_segment [OF sumsf, where n = Suc n]
show ?thesis
  using atLeast0AtMost lessThan_Suc_atMost sums_unique2 by fastforce
qed

```

Every bounded entire function is a constant function.

```

theorem Liouville_theorem:
  assumes holf: f holomorphic_on UNIV
    and bf: bounded (range f)
  shows f constant_on UNIV
  using Liouville_polynomial [OF holf, of 0 0, simplified]
  by (metis bf bounded_iff constant_on_def rangeI)

```

A holomorphic function f has only isolated zeros unless f is 0.

```

lemma powser_0_nonzero:
  fixes a :: nat ⇒ 'a::{real_normed_field,banach}
  assumes r: 0 < r
    and sm: ∧x. norm (x-ξ) < r ⇒ (λn. a n * (x-ξ) ^ n) sums (f x)
    and [simp]: f ξ = 0
    and m0: a m ≠ 0 and m>0
  obtains s where 0 < s and ∧z. z ∈ cball ξ s - {ξ} ⇒ f z ≠ 0
proof -
  have r < conv_radius a
    using sm sums_summable by (auto simp: le_conv_radius_iff [where ξ=ξ])
  obtain m where am: a m ≠ 0 and az [simp]: (∧n. n<m ⇒ a n = 0)
proof
  show a (LEAST n. a n ≠ 0) ≠ 0
    by (metis (mono_tags, lifting) m0 LeastI)
qed (fastforce dest!: not_less_Least)
define b where b i = a (i+m) / a m for i
define g where g x = suminf (λi. b i * (x-ξ) ^ i) for x
have [simp]: b 0 = 1

```

```

  by (simp add: am b_def)
{ fix x::'a
  assume norm (x-ξ) < r
  then have (λn. (a m * (x-ξ) ^ m) * (b n * (x-ξ) ^ n)) sums (f x)
    using am az sm sums_zero_iff_shift [of m (λn. a n * (x-ξ) ^ n) f x]
    by (simp add: b_def monoid_mult_class.power_add algebra_simps)
  then have x ≠ ξ ⇒ (λn. b n * (x-ξ) ^ n) sums (f x / (a m * (x-ξ) ^ m))
    using am by (simp add: sums_mult_D)
} note bsums = this
then have norm (x-ξ) < r ⇒ summable (λn. b n * (x-ξ) ^ n) for x
  using sums_summable by (cases x=ξ) auto
then have r ≤ conv_radius b
  by (simp add: le_conv_radius_iff [where ξ=ξ])
then have r/2 < conv_radius b
  using not_le order_trans r by fastforce
then have continuous_on (cball ξ (r/2)) g
  using powner_continuous_suminf [of r/2 b ξ] by (simp add: g_def)
then obtain s where s>0 ∧ x. [norm (x-ξ) ≤ s; norm (x-ξ) ≤ r/2] ⇒ dist
(g x) (g ξ) < 1/2
proof (rule continuous_onE)
  show ξ ∈ cball ξ (r / 2) 1/2 > (0::real)
    using r by auto
qed (auto simp: dist_commute dist_norm)
moreover have g ξ = 1
  by (simp add: g_def)
ultimately have gnz: ∧x. [norm (x-ξ) ≤ s; norm (x-ξ) ≤ r/2] ⇒ (g x) ≠ 0
  by fastforce
show ?thesis
proof
  have *: f x ≠ 0 if x ≠ ξ norm (x-ξ) ≤ s norm (x-ξ) ≤ r/2 for x
    using bsums [of x] that gnz [of x] r sums_iff unfolding g_def by fastforce
  show ∧z. z ∈ cball ξ (min s (r / 2)) - {ξ} ⇒ f z ≠ 0
    by (simp add: * dist_norm norm_minus_commute)
qed (use ‹0 < r› ‹0 < s› in auto)
qed

```

4.15 Complex functions and power series

The following defines the power series expansion of a complex function at a given point (assuming that it is analytic at that point).

definition *fps_expansion* :: (complex ⇒ complex) ⇒ complex ⇒ complex *fps*
where

$$fps_expansion\ f\ z0 = Abs_fps\ (\lambda n. (deriv \hat{\sim} n)\ f\ z0 / fact\ n)$$

lemma *fps_expansion_cong*:

assumes $\forall_F w$ in *nhds* *x*. $f\ w = g\ w$

shows $fps_expansion\ f\ x = fps_expansion\ g\ x$

unfolding *fps_expansion_def* **using** *assms higher_deriv_cong_ev* **by** *fastforce*

```

lemma
  fixes r :: ereal
  assumes f holomorphic_on eball z0 r
  shows conv_radius_fps_expansion: fps_conv_radius (fps_expansion f z0) ≥ r
    and eval_fps_expansion:  $\bigwedge z. z \in \text{eball } z0 \ r \implies \text{eval\_fps } (\text{fps\_expansion } f \ z0) \ (z - z0) = f \ z$ 
    and eval_fps_expansion':  $\bigwedge z. \text{norm } z < r \implies \text{eval\_fps } (\text{fps\_expansion } f \ z0) \ z = f \ (z0 + z)$ 
  proof -
    have  $(\lambda n. \text{fps\_nth } (\text{fps\_expansion } f \ z0) \ n * (z - z0) ^ n)$  sums f z
      if  $z \in \text{ball } z0 \ r' \ \text{ereal } r' < r$  for  $z \ r'$ 
    proof -
      have f holomorphic_on ball z0 r'
        using holomorphic_on_subset[OF _ ball_eball_mono] assms that by force
      then show ?thesis
        using fps_expansion_def holomorphic_power_series that by auto
    qed
    hence *:  $(\lambda n. \text{fps\_nth } (\text{fps\_expansion } f \ z0) \ n * (z - z0) ^ n)$  sums f z
      if  $z \in \text{eball } z0 \ r$  for  $z$ 
      using that by (subst (asm) eball_conv_UNION_balls) blast
    show fps_conv_radius (fps_expansion f z0) ≥ r unfolding fps_conv_radius_def
    proof (rule conv_radius_geI_ex)
      fix r' :: real assume r':  $r' > 0 \ \text{ereal } r' < r$ 
      thus  $\exists z. \text{norm } z = r' \wedge \text{summable } (\lambda n. \text{fps\_nth } (\text{fps\_expansion } f \ z0) \ n * z ^ n)$ 
        using *[of z0 + of_real r']
        by (intro exI[of _ of_real r']) (auto simp: summable_def dist_norm)
    qed
    show eval_fps (fps_expansion f z0) (z - z0) = f z if  $z \in \text{eball } z0 \ r$  for  $z$ 
      using *[OF that] by (simp add: eval_fps_def sums_iff)
    show eval_fps (fps_expansion f z0) z = f (z0 + z) if  $\text{ereal } (\text{norm } z) < r$  for  $z$ 
      using *[of z0 + z] and that by (simp add: eval_fps_def sums_iff dist_norm)
  qed

```

We can now show several more facts about power series expansions (at least in the complex case) with relative ease that would have been trickier without complex analysis.

```

lemma
  fixes f :: complex fps and r :: ereal
  assumes  $\bigwedge z. \text{ereal } (\text{norm } z) < r \implies \text{eval\_fps } f \ z \neq 0$ 
  shows fps_conv_radius_inverse:  $\text{fps\_conv\_radius } (\text{inverse } f) \geq \min r \ (\text{fps\_conv\_radius } f)$ 
    and eval_fps_inverse:  $\bigwedge z. \text{ereal } (\text{norm } z) < \text{fps\_conv\_radius } f \implies \text{ereal } (\text{norm } z) < r \implies$ 
       $\text{eval\_fps } (\text{inverse } f) \ z = \text{inverse } (\text{eval\_fps } f \ z)$ 
  proof -
    define R where  $R = \min (\text{fps\_conv\_radius } f) \ r$ 
    have *:  $\text{fps\_conv\_radius } (\text{inverse } f) \geq \min r \ (\text{fps\_conv\_radius } f) \wedge$ 
       $(\forall z \in \text{eball } 0 \ (\min (\text{fps\_conv\_radius } f) \ r). \text{eval\_fps } (\text{inverse } f) \ z = \text{inverse } (\text{eval\_fps } f \ z))$ 

```

```

(eval_fps f z)
proof (cases min r (fps_conv_radius f) > 0)
  case True
  define f' where f' = fps_expansion (λz. inverse (eval_fps f z)) 0
  have holo: (λz. inverse (eval_fps f z)) holomorphic_on eball 0 (min r (fps_conv_radius
f))
  using assms by (intro holomorphic_intros) auto
  from holo have radius: fps_conv_radius f' ≥ min r (fps_conv_radius f)
  unfolding f'_def by (rule conv_radius_fps_expansion)
  have eval_f': eval_fps f' z = inverse (eval_fps f z)
  if norm z < fps_conv_radius f norm z < r for z
  using that unfolding f'_def by (subst eval_fps_expansion'[OF holo]) auto

  have f * f' = 1
  proof (rule eval_fps_eqD)
    from radius and True have 0 < min (fps_conv_radius f) (fps_conv_radius
f')
    by (auto simp: min_def split: if_splits)
    also have ... ≤ fps_conv_radius (f * f') by (rule fps_conv_radius_mult)
    finally show ... > 0 .
  next
  from True have R > 0 by (auto simp: R_def)
  hence eventually (λz. z ∈ eball 0 R) (nhds 0)
  by (intro eventually_nhds_in_open) (auto simp: zero_ereal_def)
  thus eventually (λz. eval_fps (f * f') z = eval_fps 1 z) (nhds 0)
  proof eventually_elim
    case (elim z)
    hence eval_fps (f * f') z = eval_fps f z * eval_fps f' z
    using radius by (intro eval_fps_mult)
    (auto simp: R_def min_def split: if_splits intro: less_trans)
    also have eval_fps f' z = inverse (eval_fps f z)
    using elim by (intro eval_f') (auto simp: R_def)
    also from elim have eval_fps f z ≠ 0
    by (intro assms) (auto simp: R_def)
    hence eval_fps f z * inverse (eval_fps f z) = eval_fps 1 z
    by simp
    finally show eval_fps (f * f') z = eval_fps 1 z .
  qed
  qed simp_all
  hence f' = inverse f
  by (intro fps_inverse_unique [symmetric]) (simp_all add: mult_ac)
  with eval_f' and radius show ?thesis by simp
next
case False
  hence *: eball 0 R = {}
  by (intro eball_empty) (auto simp: R_def min_def split: if_splits)
  show ?thesis
proof
  from False have min r (fps_conv_radius f) ≤ 0

```

```

    by (simp add: min_def)
  also have  $0 \leq \text{fps\_conv\_radius } (\text{inverse } f)$ 
    by (simp add: fps_conv_radius_def conv_radius_nonneg)
  finally show  $\min r (\text{fps\_conv\_radius } f) \leq \dots$  .
qed (unfold * [unfolded R_def], auto)
qed

show  $\text{fps\_conv\_radius } (\text{inverse } f) \geq \min r (\text{fps\_conv\_radius } f)$ 
  using * by blast
show  $\text{eval\_fps } (\text{inverse } f) z = \text{inverse } (\text{eval\_fps } f z)$ 
  if  $\text{ereal } (\text{norm } z) < \text{fps\_conv\_radius } f$   $\text{ereal } (\text{norm } z) < r$  for  $z$ 
  using that * by auto
qed

lemma
  fixes  $f g :: \text{complex fps}$  and  $r :: \text{ereal}$ 
  defines  $R \equiv \text{Min } \{r, \text{fps\_conv\_radius } f, \text{fps\_conv\_radius } g\}$ 
  assumes  $\text{fps\_conv\_radius } f > 0$   $\text{fps\_conv\_radius } g > 0$   $r > 0$ 
  assumes  $\text{nz}: \bigwedge z. z \in \text{eball } 0 r \implies \text{eval\_fps } g z \neq 0$ 
  shows  $\text{fps\_conv\_radius\_divide}': \text{fps\_conv\_radius } (f / g) \geq R$ 
    and  $\text{eval\_fps\_divide}':$ 
       $\text{ereal } (\text{norm } z) < R \implies \text{eval\_fps } (f / g) z = \text{eval\_fps } f z / \text{eval\_fps } g z$ 
proof -
  from  $\text{nz}[\text{of } 0]$  and  $\langle r > 0 \rangle$  have  $\text{nz}': \text{fps\_nth } g 0 \neq 0$ 
    by (auto simp: eval_fps_at_0 zero_ereal_def)
  have  $R \leq \min r (\text{fps\_conv\_radius } g)$ 
    by (auto simp: R_def intro: min.coboundedI2)
  also have  $\min r (\text{fps\_conv\_radius } g) \leq \text{fps\_conv\_radius } (\text{inverse } g)$ 
    by (intro fps_conv_radius_inverse assms) (auto simp: zero_ereal_def)
  finally have  $\text{radius}: \text{fps\_conv\_radius } (\text{inverse } g) \geq R$  .
  have  $R \leq \min (\text{fps\_conv\_radius } f) (\text{fps\_conv\_radius } (\text{inverse } g))$ 
    by (intro radius min.boundedI) (auto simp: R_def intro: min.coboundedI1
min.coboundedI2)
  also have  $\dots \leq \text{fps\_conv\_radius } (f * \text{inverse } g)$ 
    by (rule fps_conv_radius_mult)
  also have  $f * \text{inverse } g = f / g$ 
    by (intro fps_divide_unit [symmetric] nz')
  finally show  $\text{fps\_conv\_radius } (f / g) \geq R$  .

  assume  $z: \text{ereal } (\text{norm } z) < R$ 
  have  $\text{eval\_fps } (f * \text{inverse } g) z = \text{eval\_fps } f z * \text{eval\_fps } (\text{inverse } g) z$ 
    using radius by (intro eval_fps_mult less_le_trans[OF z])
    (auto simp: R_def intro: min.coboundedI1 min.coboundedI2)
  also have  $\text{eval\_fps } (\text{inverse } g) z = \text{inverse } (\text{eval\_fps } g z)$  using  $\langle r > 0 \rangle$ 
    by (intro eval_fps_inverse[where  $r = r$ ] less_le_trans[OF z] nz)
    (auto simp: R_def intro: min.coboundedI1 min.coboundedI2)
  also have  $f * \text{inverse } g = f / g$  by fact
  finally show  $\text{eval\_fps } (f / g) z = \text{eval\_fps } f z / \text{eval\_fps } g z$ 
    by (simp add: field_split_simps)

```


qed

lemma

```

fixes  $f\ g :: \text{complex fps}$  and  $r :: \text{ereal}$ 
defines  $R \equiv \text{Min } \{r, \text{fps\_conv\_radius } f, \text{fps\_conv\_radius } g\}$ 
assumes  $\text{subdegree } g \leq \text{subdegree } f$ 
assumes  $\text{fps\_conv\_radius } f > 0$   $\text{fps\_conv\_radius } g > 0$   $r > 0$ 
assumes  $\bigwedge z. z \in \text{eball } 0\ r \implies z \neq 0 \implies \text{eval\_fps } g\ z \neq 0$ 
shows  $\text{fps\_conv\_radius\_divide: fps\_conv\_radius } (f / g) \geq R$ 
and  $\text{eval\_fps\_divide:}$ 
 $\text{ereal } (\text{norm } z) < R \implies c = \text{fps\_nth } f\ (\text{subdegree } g) / \text{fps\_nth } g\ (\text{subdegree } g) \implies$ 
 $\text{eval\_fps } (f / g)\ z = (\text{if } z = 0 \text{ then } c \text{ else } \text{eval\_fps } f\ z / \text{eval\_fps } g\ z)$ 
proof -
define  $f'\ g'$  where  $f' = \text{fps\_shift } (\text{subdegree } g)\ f$  and  $g' = \text{fps\_shift } (\text{subdegree } g)\ g$ 
have  $f\_eq: f = f' * \text{fps\_X}^{\wedge} \text{subdegree } g$  and  $g\_eq: g = g' * \text{fps\_X}^{\wedge} \text{subdegree } g$ 
unfolding  $f'\_def\ g'\_def$  by (rule  $\text{subdegree\_decompose' le\_refl | fact}$ )+
have  $\text{subdegree: subdegree } f' = \text{subdegree } f - \text{subdegree } g$   $\text{subdegree } g' = 0$ 
using  $\text{assms}(2)$  by ( $\text{simp\_all add: } f'\_def\ g'\_def$ )
have [ $\text{simp}$ ]:  $\text{fps\_conv\_radius } f' = \text{fps\_conv\_radius } f$   $\text{fps\_conv\_radius } g' = \text{fps\_conv\_radius } g$ 
by ( $\text{simp\_all add: } f'\_def\ g'\_def$ )
have [ $\text{simp}$ ]:  $\text{fps\_nth } f'\ 0 = \text{fps\_nth } f\ (\text{subdegree } g)$ 
 $\text{fps\_nth } g'\ 0 = \text{fps\_nth } g\ (\text{subdegree } g)$  by ( $\text{simp\_all add: } f'\_def\ g'\_def$ )
have  $g\_nz: g \neq 0$ 
proof -
define  $z :: \text{complex}$  where  $z = (\text{if } r = \infty \text{ then } 1 \text{ else } \text{of\_real } (\text{real\_of\_ereal } r / 2))$ 
have  $z \in \text{eball } 0\ r$ 
using  $\langle r > 0 \rangle$   $\text{ereal\_less\_real\_iff } z\_def$  by  $\text{fastforce}$ 
moreover have  $z \neq 0$  using  $\langle r > 0 \rangle$ 
by ( $\text{cases } r$ ) ( $\text{auto simp: } z\_def$ )
ultimately have  $\text{eval\_fps } g\ z \neq 0$  by (rule  $\text{assms}(6)$ )
thus  $g \neq 0$  by  $\text{auto}$ 
qed
have  $fg: f / g = f' * \text{inverse } g'$ 
by ( $\text{subst } f\_eq, \text{subst } (2)\ g\_eq$ ) ( $\text{insert } g\_nz, \text{simp add: fps\_divide\_unit}$ )
have  $g'\_nz: \text{eval\_fps } g'\ z \neq 0$  if  $z: \text{norm } z < \text{min } r\ (\text{fps\_conv\_radius } g)$  for  $z$ 
proof ( $\text{cases } z = 0$ )
case  $\text{False}$ 
with  $\text{assms}$  and  $z$  have  $\text{eval\_fps } g\ z \neq 0$  by  $\text{auto}$ 
also from  $z$  have  $\text{eval\_fps } g\ z = \text{eval\_fps } g'\ z * z^{\wedge} \text{subdegree } g$ 
by ( $\text{subst } g\_eq$ ) ( $\text{auto simp: eval\_fps\_mult}$ )
finally show  $?thesis$  by  $\text{auto}$ 
qed ( $\text{use } \langle g \neq 0 \rangle$  in  $\langle \text{auto simp: } g'\_def\ \text{eval\_fps\_at\_0} \rangle$ )

```

```

have  $R \leq \min (\min r (\text{fps\_conv\_radius } g)) (\text{fps\_conv\_radius } g')$ 
  by (auto simp:  $R\_def$   $\text{min.coboundedI1}$   $\text{min.coboundedI2}$ )
also have  $\dots \leq \text{fps\_conv\_radius } (\text{inverse } g')$ 
  using  $g'\_nz$  by (rule  $\text{fps\_conv\_radius\_inverse}$ )
finally have  $\text{conv\_radius\_inv}: R \leq \text{fps\_conv\_radius } (\text{inverse } g')$  .
hence  $R \leq \text{fps\_conv\_radius } (f' * \text{inverse } g')$ 
  by (intro  $\text{order.trans}[OF\_ \text{fps\_conv\_radius\_mult}]$ )
  (auto simp:  $R\_def$   $\text{intro: min.coboundedI1 min.coboundedI2}$ )
thus  $\text{fps\_conv\_radius } (f / g) \geq R$  by (simp add:  $fg$ )

fix  $z c :: \text{complex}$  assume  $z: \text{ereal } (\text{norm } z) < R$ 
assume  $c: c = \text{fps\_nth } f (\text{subdegree } g) / \text{fps\_nth } g (\text{subdegree } g)$ 
show  $\text{eval\_fps } (f / g) z = (\text{if } z = 0 \text{ then } c \text{ else } \text{eval\_fps } f z / \text{eval\_fps } g z)$ 
proof (cases  $z = 0$ )
  case False
  from  $z$  and  $\text{conv\_radius\_inv}$  have  $\text{ereal } (\text{norm } z) < \text{fps\_conv\_radius } (\text{inverse } g')$ 
  by simp
  with  $z$  have  $\text{eval\_fps } (f / g) z = \text{eval\_fps } f' z * \text{eval\_fps } (\text{inverse } g') z$ 
  unfolding  $fg$  by (subst  $\text{eval\_fps\_mult}$ ) (auto simp:  $R\_def$ )
  also have  $\text{eval\_fps } (\text{inverse } g') z = \text{inverse } (\text{eval\_fps } g' z)$ 
  using  $z$  by (intro  $\text{eval\_fps\_inverse}[of \text{min } r (\text{fps\_conv\_radius } g')]$   $g'\_nz$ )
  (auto simp:  $R\_def$ )
  also have  $\text{eval\_fps } f' z * \dots = \text{eval\_fps } f z / \text{eval\_fps } g z$ 
  using  $z$  False  $\text{assms}(2)$  by (simp add:  $f'\_def$   $g'\_def$   $\text{eval\_fps\_shift } R\_def$ )
  finally show ?thesis using False by simp
qed (simp_all add:  $\text{eval\_fps\_at\_0}$   $fg$   $\text{field\_sims } c$ )
qed

lemma  $\text{has\_fps\_expansion\_fps\_expansion}$  [intro]:
  assumes  $\text{open } A$   $0 \in A$   $f$   $\text{holomorphic\_on } A$ 
  shows  $f$   $\text{has\_fps\_expansion } \text{fps\_expansion } f$   $0$ 
proof -
  from  $\text{assms}$  obtain  $r$  where  $r > 0$  and  $r: \text{ball } 0 r \subseteq A$ 
  by (auto simp:  $\text{open\_contains\_ball}$ )
  with  $\text{assms}$  have  $\text{holo}: f$   $\text{holomorphic\_on } \text{eball } 0 (\text{ereal } r)$ 
  by auto
  have  $r \leq \text{fps\_conv\_radius } (\text{fps\_expansion } f$   $0)$ 
  using  $\text{holo}$  by (intro  $\text{conv\_radius\_fps\_expansion}$ ) auto
  then have  $\dots > 0$ 
  by (simp add:  $\text{ereal\_le\_less } \langle r > 0 \rangle$   $\text{zero\_ereal\_def}$ )
  moreover have  $\text{eventually } (\lambda z. z \in \text{ball } 0 r)$   $(\text{nhds } 0)$ 
  using  $\langle r > 0 \rangle$  by (intro  $\text{eventually\_nhds\_in\_open}$ ) auto
  hence  $\text{eventually } (\lambda z. \text{eval\_fps } (\text{fps\_expansion } f$   $0) z = f z)$   $(\text{nhds } 0)$ 
  by  $\text{eventually\_elim}$  (subst  $\text{eval\_fps\_expansion}'[OF \text{holo}]$ , auto)
  ultimately show ?thesis
  using  $\langle r > 0 \rangle$  by (auto simp:  $\text{has\_fps\_expansion\_def}$ )
qed

```

```

lemma fps_conv_radius_tan:
  fixes c :: complex
  assumes c ≠ 0
  shows fps_conv_radius (fps_tan c) ≥ pi / (2 * norm c)
proof -
  have fps_conv_radius (fps_tan c) ≥
    Min {pi / (2 * norm c), fps_conv_radius (fps_sin c), fps_conv_radius
(fps_cos c)}
  unfolding fps_tan_def
  proof (rule fps_conv_radius_divide)
    fix z :: complex assume z ∈ eball 0 (pi / (2 * norm c))
    with cos_eq_zero_imp_norm_ge[of c*z] assms
    show eval_fps (fps_cos c) z ≠ 0 by (auto simp: norm_mult field_simps)
  qed (insert assms, auto)
  thus ?thesis by (simp add: min_def)
qed

lemma eval_fps_tan:
  fixes c :: complex
  assumes norm z < pi / (2 * norm c)
  shows eval_fps (fps_tan c) z = tan (c * z)
proof (cases c = 0)
  case False
  show ?thesis unfolding fps_tan_def
  proof (subst eval_fps_divide'[where r = pi / (2 * norm c)])
    fix z :: complex assume z ∈ eball 0 (pi / (2 * norm c))
    with cos_eq_zero_imp_norm_ge[of c*z] assms
    show eval_fps (fps_cos c) z ≠ 0 using False by (auto simp: norm_mult
field_simps)
  qed (use False assms in ‹auto simp: field_simps tan_def›)
qed simp_all

end

```

5 Conformal Mappings and Consequences of Cauchy's Integral Theorem

By John Harrison et al. Ported from HOL Light by L C Paulson (2016)

Also Cauchy's residue theorem by Wenda Li (2016)

```

theory Conformal_Mappings
imports Cauchy_Integral_Formula

begin

```

5.1 Analytic continuation

proposition *isolated_zeros*:

assumes *holf*: f holomorphic_on S
and *open* S **connected** S $\xi \in S$ $f \xi = 0$ $\beta \in S$ $f \beta \neq 0$
obtains r **where** $0 < r$ **and** $\text{ball } \xi \ r \subseteq S$ **and**
 $\bigwedge z. z \in \text{ball } \xi \ r - \{\xi\} \implies f z \neq 0$

proof –

obtain r **where** $0 < r$ **and** r : $\text{ball } \xi \ r \subseteq S$
using $\langle \text{open } S \rangle \langle \xi \in S \rangle$ *open_contains_ball_eq* **by** *blast*
have *powf*: $((\lambda n. (\text{deriv } \sim n) f \xi / (\text{fact } n) * (z - \xi) ^ n) \text{ sums } f z)$ **if** $z \in \text{ball } \xi \ r$ **for** z
by (*intro holomorphic_power_series [OF _ that] holomorphic_on_subset [OF holf r]*)
obtain m **where** m : $(\text{deriv } \sim m) f \xi / (\text{fact } m) \neq 0$
using *holomorphic_fun_eq_0_on_connected [OF holf \langle open S \rangle \langle connected S \rangle \langle \xi \in S \rangle \langle \beta \in S \rangle \langle f \beta \neq 0 \rangle]*
by *auto*
then **have** $m \neq 0$ **using** *assms(5) funpow_0* **by** *fastforce*
obtain s **where** $0 < s$ **and** s : $\bigwedge z. z \in \text{cball } \xi \ s - \{\xi\} \implies f z \neq 0$
using *power_0_nonzero [OF \langle 0 < r \rangle powf \langle f \xi = 0 \rangle m]*
by (*metis \langle m \neq 0 \rangle dist_norm mem_ball norm_minus_commute not_gr_zero*)
have $0 < \min r \ s$ **by** (*simp add: \langle 0 < r \rangle \langle 0 < s \rangle*)
then **show** *thesis*
proof *qed* (*use r s in auto*)
qed

proposition *analytic_continuation*:

assumes *holf*: f holomorphic_on S
and *open* S **and** *connected* S
and $U \subseteq S$ **and** $\xi \in S$
and ξ *islimpt* U
and $fU0$ [*simp*]: $\bigwedge z. z \in U \implies f z = 0$
and $w \in S$
shows $f w = 0$

proof –

obtain e **where** $0 < e$ **and** e : $\text{cball } \xi \ e \subseteq S$
using $\langle \text{open } S \rangle \langle \xi \in S \rangle$ *open_contains_cball_eq* **by** *blast*
define T **where** $T = \text{cball } \xi \ e \cap U$
have *contf*: *continuous_on* (*closure* T) f
by (*metis T_def closed_cball closure_minimal e holf holomorphic_on_imp_continuous_on holomorphic_on_subset inf.cobounded1*)
have $fT0$ [*simp*]: $\bigwedge x. x \in T \implies f x = 0$
by (*simp add: T_def*)
have $\bigwedge r. [\forall e > 0. \exists x' \in U. x' \neq \xi \wedge \text{dist } x' \ \xi < e; 0 < r] \implies \exists x' \in \text{cball } \xi \ e \cap U. x' \neq \xi \wedge \text{dist } x' \ \xi < r$
by (*metis \langle 0 < e \rangle IntI dist_commute less_eq_real_def mem_cball min_less_iff_conj*)
then **have** ξ *islimpt* T **using** $\langle \xi$ *islimpt* $U \rangle$
by (*auto simp: T_def islimpt_approachable*)
then **have** $\xi \in \text{closure } T$

```

  by (simp add: closure_def)
then have  $f \xi = 0$ 
  by (auto simp: continuous_constant_on_closure [OF contf])
moreover have  $\bigwedge r. [0 < r; \bigwedge z. z \in \text{ball } \xi \ r - \{\xi\} \implies f z \neq 0] \implies \text{False}$ 
  by (metis open_ball  $\langle \xi \text{ islimpt } T \rangle$  centre_in_ball fT0 insertE insert_Diff
islimptE)
ultimately show ?thesis
  by (metis  $\langle \text{open } S \rangle$   $\langle \text{connected } S \rangle$   $\langle \xi \in S \rangle$   $\langle w \in S \rangle$  holf isolated_zeros)
qed

```

```

corollary analytic_continuation_open:
  assumes open s and open s' and  $s \neq \{\}$  and connected s'
    and  $s \subseteq s'$ 
  assumes f holomorphic_on s' and g holomorphic_on s'
    and  $\bigwedge z. z \in s \implies f z = g z$ 
  assumes  $z \in s'$ 
  shows  $f z = g z$ 
proof -
  from  $\langle s \neq \{\} \rangle$  obtain  $\xi$  where  $\xi \in s$  by auto
  with  $\langle \text{open } s \rangle$  have  $\xi: \xi \text{ islimpt } s$ 
    by (intro interior_limit_point) (auto simp: interior_open)
  have  $f z - g z = 0$ 
    by (rule analytic_continuation[of  $\lambda z. f z - g z \ s' \ s \ \xi$ ])
      (insert assms  $\langle \xi \in s \rangle \ \xi$ , auto intro: holomorphic_intros)
  thus ?thesis by simp
qed

```

```

corollary analytic_continuation':
  assumes f holomorphic_on S open S connected S
    and  $U \subseteq S \ \xi \in S \ \xi \text{ islimpt } U$ 
    and f constant_on U
  shows f constant_on S
proof -
  obtain c where  $c: \bigwedge x. x \in U \implies f x - c = 0$ 
    by (metis  $\langle f \text{ constant_on } U \rangle$  constant_on_def diff_self)
  have  $(\lambda z. f z - c)$  holomorphic_on S
    using assms by (intro holomorphic_intros)
  with c analytic_continuation assms have  $\bigwedge x. x \in S \implies f x - c = 0$ 
    by blast
  then show ?thesis
    unfolding constant_on_def by force
qed

```

```

lemma holomorphic_compact_finite_zeros:
  assumes S: f holomorphic_on S open S connected S
    and compact K  $K \subseteq S$ 
    and  $\neg f \text{ constant_on } S$ 
  shows finite  $\{z \in K. f z = 0\}$ 
proof (rule ccontr)

```

```

assume infinite { $z \in K. f z = 0$ }
then obtain  $z$  where  $z \in K$  and  $z: z \text{ islimpt } \{z \in K. f z = 0\}$ 
  using  $\langle \text{compact } K \rangle$  by (auto simp: compact_eq_Bolzano_Weierstrass)
moreover have  $\{z \in K. f z = 0\} \subseteq S$ 
  using  $\langle K \subseteq S \rangle$  by blast
  ultimately show False
  using assms analytic_continuation [OF S] unfolding constant_on_def
  by blast
qed

```

lemma *holomorphic_countable_zeros*:

```

assumes  $S: f \text{ holomorphic\_on } S$   $open\ S$   $connected\ S$  and  $f \text{ sigma } S$ 
  and  $\neg f \text{ constant\_on } S$ 
shows countable { $z \in S. f z = 0$ }
proof –
  obtain  $F::nat \Rightarrow \text{complex set}$ 
    where  $F: range\ F \subseteq \text{Collect compact}$  and  $Seq: S = (\bigcup i. F i)$ 
    using  $\langle f \text{ sigma } S \rangle$  by (meson fsigma_Union_compact)
  have  $fin: \text{finite } \{z \in F i. f z = 0\}$  for  $i$ 
    using holomorphic_compact_finite_zeros assms F Seq Union_iff by blast
  have  $\{z \in S. f z = 0\} = (\bigcup i. \{z \in F i. f z = 0\})$ 
    using Seq by auto
  with  $fin$  show ?thesis
    by (simp add: countable_finite)
qed

```

lemma *holomorphic_countable_equal*:

```

assumes  $f \text{ holomorphic\_on } S$   $g \text{ holomorphic\_on } S$   $open\ S$   $connected\ S$  and  $f \text{ sigma } S$ 
  and  $eq: \text{uncountable } \{z \in S. f z = g z\}$ 
shows  $S \subseteq \{z \in S. f z = g z\}$ 
proof –
  obtain  $z$  where  $z: z \in S\ f z = g z$ 
    using eq not_finite_existsD uncountable_infinite by blast
  have  $(\lambda x. f x - g x) \text{ holomorphic\_on } S$ 
    by (simp add: assms holomorphic_on_diff)
  then have  $(\lambda x. f x - g x) \text{ constant\_on } S$ 
    using holomorphic_countable_zeros assms by force
  with  $z$  have  $\bigwedge x. x \in S \implies f x - g x = 0$ 
    unfolding constant_on_def by force
  then show ?thesis
    by auto
qed

```

lemma *holomorphic_countable_equal_UNIV*:

```

assumes  $fg: f \text{ holomorphic\_on } UNIV$   $g \text{ holomorphic\_on } UNIV$ 
  and  $eq: \text{uncountable } \{z. f z = g z\}$ 
shows  $f=g$ 
using holomorphic_countable_equal [OF fg] eq by fastforce

```

5.2 Open mapping theorem

lemma *holomorphic_contract_to_zero*:

assumes *contf*: *continuous_on* (cball ξ r) *f*
 and *holf*: *f* *holomorphic_on* ball ξ r
 and $0 < r$
 and *norm_less*: $\bigwedge z. \text{norm}(\xi - z) = r \implies \text{norm}(f \xi) < \text{norm}(f z)$
 obtains *z* where $z \in \text{ball } \xi \ r \ f z = 0$

proof –

```

{ assume fnz:  $\bigwedge w. w \in \text{ball } \xi \ r \implies f w \neq 0$ 
  then have  $0 < \text{norm} (f \xi)$ 
    by (simp add:  $\langle 0 < r \rangle$ )
  have fnz':  $\bigwedge w. w \in \text{cball } \xi \ r \implies f w \neq 0$ 
    using dist_complex_def fnz norm_less order_le_less by fastforce
  have frontier(cball  $\xi$   $r$ )  $\neq \{\}$ 
    using  $\langle 0 < r \rangle$  by simp
  define g where [abs_def]:  $g z = \text{inverse} (f z)$  for z
  have contg: continuous_on (cball  $\xi$   $r$ ) g
    unfolding g_def using contf continuous_on_inverse fnz' by blast
  have holg: g holomorphic_on ball  $\xi$   $r$ 
    unfolding g_def using fnz holf holomorphic_on_inverse by blast
  have frontier (cball  $\xi$   $r$ )  $\subseteq$  cball  $\xi$   $r$ 
    by (simp add: subset_iff)
  then have contf': continuous_on (frontier (cball  $\xi$   $r$ )) f
    and contg': continuous_on (frontier (cball  $\xi$   $r$ )) g
    by (blast intro: contf contg continuous_on_subset)+
  have froc: frontier(cball  $\xi$   $r$ )  $\neq \{\}$ 
    using  $\langle 0 < r \rangle$  by simp
  moreover have continuous_on (frontier (cball  $\xi$   $r$ )) (norm  $\circ$  f)
    using contf' continuous_on_compose continuous_on_norm_id by blast
  ultimately obtain w where  $w \in \text{frontier}(\text{cball } \xi \ r)$ 
    and now:  $\bigwedge x. x \in \text{frontier}(\text{cball } \xi \ r) \implies \text{norm} (f w) \leq \text{norm}$ 
(f x)
    using continuous_attains_inf [OF compact_frontier [OF compact_cball]]
    by (metis comp_apply)
  then have fw:  $0 < \text{norm} (f w)$ 
    by (simp add: fnz')
  have continuous_on (frontier (cball  $\xi$   $r$ )) (norm  $\circ$  g)
    using contg' continuous_on_compose continuous_on_norm_id by blast
  then obtain v where  $v \in \text{frontier}(\text{cball } \xi \ r)$ 
    and nov:  $\bigwedge x. x \in \text{frontier}(\text{cball } \xi \ r) \implies \text{norm} (g v) \geq \text{norm} (g x)$ 
    using continuous_attains_sup [OF compact_frontier [OF compact_cball]]
froc] by force
  then have fv:  $0 < \text{norm} (f v)$ 
    by (simp add: fnz')
  have  $\text{norm} ((\text{deriv } \hat{\sim} 0) \ g \ \xi) \leq \text{fact } 0 * \text{norm} (g v) / r \wedge 0$ 
    by (rule Cauchy_inequality [OF holg contg  $\langle 0 < r \rangle$ ]) (simp add: dist_norm
nov)
  then have cmod (g  $\xi$ )  $\leq$  cmod (g  $v$ )
    by simp

```

```

moreover have cmod ( $\xi - w$ ) = r
  by (metis (no_types) dist_norm frontier_cball mem_sphere w)
ultimately obtain wr: norm ( $\xi - w$ ) = r and nfw: norm (f w)  $\leq$  norm (f  $\xi$ )
  unfolding g_def
  by (smt (verit, del_insts)  $\langle 0 < \text{cmod } (f \xi) \rangle$  inverse_le_imp_le norm_inverse
now v)
  with fw have False
  using norm_less by force
}
with that show ?thesis by blast
qed

```

theorem open_mapping_thm:

```

assumes holf: f holomorphic_on S
  and S: open S and connected S
  and open U and U  $\subseteq$  S
  and fne:  $\neg$  f constant_on S
shows open (f ' U)
proof -
  have *: open (f ' U)
    if U  $\neq$  {} and U: open U connected U and f holomorphic_on U and
fneU:  $\bigwedge x. \exists y \in U. f y \neq x$ 
    for U
  proof (clarsimp simp: open_contains_ball)
  fix  $\xi$  assume  $\xi \in U$ 
  show  $\exists e > 0. \text{ball } (f \xi) e \subseteq f ' U$ 
  proof -
  have hol: ( $\lambda z. f z - f \xi$ ) holomorphic_on U
    by (rule holomorphic_intros that)+
  obtain s where  $0 < s$  and sbU: ball  $\xi$  s  $\subseteq$  U
    and sne:  $\bigwedge z. z \in \text{ball } \xi s - \{\xi\} \implies (\lambda z. f z - f \xi) z \neq 0$ 
    using isolated_zeros [OF hol U  $\xi$ ] by (metis fneU right_minus_eq)
  obtain r where  $0 < r$  and r: cball  $\xi$  r  $\subseteq$  ball  $\xi$  s
    using  $\langle 0 < s \rangle$  by (rule_tac r=s/2 in that) auto
  have cball  $\xi$  r  $\subseteq$  U
    using sbU r by blast
  then have frsbU: frontier (cball  $\xi$  r)  $\subseteq$  U
    using Diff_subset frontier_def order_trans by fastforce
  then have cof: compact (frontier (cball  $\xi$  r))
    by blast
  have frne: frontier (cball  $\xi$  r)  $\neq$  {}
    using  $\langle 0 < r \rangle$  by auto
  have contfr: continuous_on (frontier (cball  $\xi$  r)) ( $\lambda z. \text{norm } (f z - f \xi)$ )
    by (metis continuous_on_norm continuous_on_subset frsbU hol holomorphic_on_imp_continuous_on)
  obtain w where norm ( $\xi - w$ ) = r
    and w: ( $\bigwedge z. \text{norm } (\xi - z) = r \implies \text{norm } (f w - f \xi) \leq \text{norm } (f z - f \xi)$ )
    using continuous_attains_inf [OF cof frne contfr] by (auto simp: dist_norm)

```



```

moreover define  $\varepsilon$  where  $\varepsilon \equiv \text{norm } (f w - f \xi) / 3$ 
ultimately have  $0 < \varepsilon$ 
  using  $\langle 0 < r \rangle \text{ dist\_complex\_def } r \text{ sne by auto}$ 
have  $\text{ball } (f \xi) \varepsilon \subseteq f' U$ 
proof
  fix  $\gamma$ 
  assume  $\gamma: \gamma \in \text{ball } (f \xi) \varepsilon$ 
  have  $*$ :  $\text{cmod } (\gamma - f \xi) < \text{cmod } (\gamma - f z)$  if  $\text{cmod } (\xi - z) = r$  for  $z$ 
  proof -
    have  $lt$ :  $\text{cmod } (f w - f \xi) / 3 < \text{cmod } (\gamma - f z)$ 
    using  $w [OF \text{ that}] \gamma$ 
    using  $\text{dist\_triangle2 [of } f \xi \gamma f z] \text{ dist\_triangle2 [of } f \xi f z \gamma]$ 
    by  $(\text{simp add: } \varepsilon\_def \text{ dist\_norm norm\_minus\_commute})$ 
  show  $?thesis$ 
    by  $(\text{metis } \varepsilon\_def \text{ dist\_commute dist\_norm less\_trans } lt \text{ mem\_ball } \gamma)$ 
  qed
have  $\text{continuous\_on } (\text{cball } \xi r) (\lambda z. \gamma - f z)$ 
  using  $\langle \text{cball } \xi r \subseteq U \rangle \langle f \text{ holomorphic\_on } U \rangle$ 
by  $(\text{force intro: continuous\_intros continuous\_on\_subset holomorphic\_on\_imp\_continuous\_on})$ 
moreover have  $(\lambda z. \gamma - f z) \text{ holomorphic\_on ball } \xi r$ 
  using  $\langle \text{cball } \xi r \subseteq U \rangle \text{ ball\_subset\_cball holomorphic\_on\_subset that}(4)$ 
  by  $(\text{intro holomorphic\_intros}) \text{ blast}$ 
ultimately obtain  $z$  where  $z \in \text{ball } \xi r \wedge \gamma - f z = 0$ 
  using  $*$   $\langle 0 < r \rangle \text{ holomorphic\_contract\_to\_zero by blast}$ 
then show  $\gamma \in f' U$ 
  using  $\langle \text{cball } \xi r \subseteq U \rangle \text{ by fastforce}$ 
qed
then show  $?thesis \text{ using } \langle 0 < \varepsilon \rangle \text{ by blast}$ 
qed
qed
have  $\text{open } (f' X)$  if  $X \in \text{components } U$  for  $X$ 
proof -
  have  $\text{hol}fU: f \text{ holomorphic\_on } U$ 
  using  $\langle U \subseteq S \rangle \text{ hol}f \text{ holomorphic\_on\_subset by blast}$ 
have  $X \neq \{\}$ 
  using  $\text{that by (simp add: in\_components\_nonempty)}$ 
moreover have  $\text{open } X$ 
  using  $\text{that } \langle \text{open } U \rangle \text{ open\_components by auto}$ 
moreover have  $\text{connected } X$ 
  using  $\text{that in\_components\_maximal by blast}$ 
moreover have  $f \text{ holomorphic\_on } X$ 
  by  $(\text{meson that hol}fU \text{ holomorphic\_on\_subset in\_components\_maximal})$ 
moreover have  $\exists y \in X. f y \neq x$  for  $x$ 
proof  $(\text{rule ccontr})$ 
  assume  $\text{not: } \neg (\exists y \in X. f y \neq x)$ 
  have  $X \subseteq S$ 
  using  $\langle U \subseteq S \rangle \text{ in\_components\_subset that by blast}$ 
obtain  $w$  where  $w \in X$  using  $\langle X \neq \{\} \rangle \text{ by blast}$ 
have  $wis: w \text{ islimpt } X$ 

```

```

    using w ⟨open X⟩ interior_eq by auto
  have hol: (λz. f z - x) holomorphic_on S
    by (simp add: holf holomorphic_on_diff)
  with fne [unfolded constant_on_def]
    analytic_continuation[OF hol S ⟨connected S⟩ ⟨X ⊆ S⟩ _ wis] not ⟨X ⊆
S⟩ w
  show False by auto
qed
ultimately show ?thesis
  by (rule *)
qed
then have open (f ' ⋃ (components U))
  by (metis (no_types, lifting) imageE image_Union open_Union)
then show ?thesis
  by force
qed

```

No need for S to be connected. But the nonconstant condition is stronger.

corollary *open_mapping_thm2*:

```

  assumes holf: f holomorphic_on S
    and S: open S
    and open U U ⊆ S
    and fnc: ⋀X. [⋈open X; X ⊆ S; X ≠ {}] ⇒ ¬ f constant_on X
  shows open (f ' U)
proof -
  have S = ⋃ (components S) by simp
  with ⟨U ⊆ S⟩ have U = (⋃ C ∈ components S. C ∩ U) by auto
  then have f ' U = (⋃ C ∈ components S. f ' (C ∩ U))
    using image_UN by fastforce
  moreover
  { fix C assume C ∈ components S
    with S ⟨C ∈ components S⟩ open_components in_components_connected
    have C: open C connected C by auto
    have C ⊆ S
      by (metis ⟨C ∈ components S⟩ in_components_maximal)
    have nf: ¬ f constant_on C
      using ⟨open C⟩ ⟨C ∈ components S⟩ ⟨C ⊆ S⟩ fnc in_components_nonempty
  by blast
  have f holomorphic_on C
    by (metis holf holomorphic_on_subset ⟨C ⊆ S⟩)
  then have open (f ' (C ∩ U))
    by (meson C ⟨open U⟩ inf_le1 nf open_Int open_mapping_thm)
  } ultimately show ?thesis
  by force
qed

```

corollary *open_mapping_thm3*:

```

  assumes f holomorphic_on S
    and open S and inj_on f S

```

shows $open (f \text{ ` } S)$
by (*meson* *assms inj_on_subset injective_not_constant open_mapping_thm2 order.refl*)

5.3 Maximum modulus principle

If f is holomorphic, then its norm (modulus) cannot exhibit a true local maximum that is properly within the domain of f .

proposition *maximum_modulus_principle*:

assumes *holf*: f *holomorphic_on* S
and S : *open* S **and** *connected* S
and *open* U **and** $U \subseteq S$ **and** $\xi \in U$
and *no*: $\bigwedge z. z \in U \implies norm(f z) \leq norm(f \xi)$
shows f *constant_on* S

proof (*rule ccontr*)

assume $\neg f$ *constant_on* S

then have *open* $(f \text{ ` } U)$

using *open_mapping_thm assms* **by** *blast*

moreover have \neg *open* $(f \text{ ` } U)$

proof –

have $\exists t. cmod (f \xi - t) < e \wedge t \notin f \text{ ` } U$ **if** $0 < e$ **for** e

using *that*

apply (*rule_tac* $x = \text{if } 0 < \text{Re}(f \xi) \text{ then } f \xi + (e/2) \text{ else } f \xi - (e/2)$ **in** *exI*)

apply (*simp add: dist_norm*)

apply (*fastforce simp: cmod_Re_le_iff dest!: no dest: sym*)

done

then show *?thesis*

unfolding *open_contains_ball* **by** (*metis* $\langle \xi \in U \rangle$ *contra_subsetD dist_norm imageI mem_ball*)

qed

ultimately show *False*

by *blast*

qed

proposition *maximum_modulus_frontier*:

assumes *holf*: f *holomorphic_on* (*interior* S)

and *contf*: *continuous_on* (*closure* S) f

and *bos*: *bounded* S

and *leB*: $\bigwedge z. z \in \text{frontier } S \implies norm(f z) \leq B$

and $\xi \in S$

shows $norm(f \xi) \leq B$

proof –

have *compact* (*closure* S) **using** *bos*

by (*simp add: bounded_closure compact_eq_bounded_closed*)

moreover have *continuous_on* (*closure* S) (*cmod* $\circ f$)

using *contf continuous_on_compose continuous_on_norm_id* **by** *blast*

ultimately obtain z **where** $z \in \text{closure } S$ **and** $z: \bigwedge y. y \in \text{closure } S \implies (cmod \circ f) y \leq (cmod \circ f) z$

using *continuous_attains_sup* [*of closure S norm o f*] $\langle \xi \in S \rangle$ **by** *auto*

```

then consider  $z \in \text{frontier } S \mid z \in \text{interior } S$  using frontier_def by auto
then have  $\text{norm}(f z) \leq B$ 
proof cases
  case 1 then show ?thesis using leB by blast
next
  case 2
  have  $f \text{ constant\_on } (\text{connected\_component\_set } (\text{interior } S) z)$ 
  proof (rule maximum_modulus_principle)
    show  $f \text{ holomorphic\_on } \text{connected\_component\_set } (\text{interior } S) z$ 
    by (metis connected_component_subset holf holomorphic_on_subset)
    show  $z \in \text{connected\_component\_set } (\text{interior } S) z$ 
    by (simp add: 2)
    show  $\bigwedge W. W \in \text{connected\_component\_set } (\text{interior } S) z \implies \text{cmod } (f W) \leq \text{cmod } (f z)$ 
    using closure_def connected_component_subset z by fastforce
  qed (auto simp: open_connected_component)
  then obtain  $c$  where  $c: \bigwedge w. w \in \text{connected\_component\_set } (\text{interior } S) z \implies f w = c$ 
  by (auto simp: constant_on_def)
  have  $f \text{ `closure}(\text{connected\_component\_set } (\text{interior } S) z) \subseteq \{c\}$ 
  proof (rule image_closure_subset)
    show  $\text{continuous\_on } (\text{closure } (\text{connected\_component\_set } (\text{interior } S) z)) f$ 
    by (meson closure_mono connected_component_subset contf continuous_on_subset interior_subset)
  qed (use c in auto)
  then have  $cc: \bigwedge w. w \in \text{closure}(\text{connected\_component\_set } (\text{interior } S) z) \implies f w = c$  by blast
  have  $\text{connected\_component } (\text{interior } S) z z$ 
  by (simp add: 2)
  moreover have  $\text{connected\_component\_set } (\text{interior } S) z \neq \text{UNIV}$ 
  by (metis bos bounded_interior connected_component_eq_UNIV not_bounded_UNIV)
  ultimately have  $\text{frontier}(\text{connected\_component\_set } (\text{interior } S) z) \neq \{\}$ 
  by (meson 2 connected_component_eq_empty frontier_not_empty)
  then obtain  $w$  where  $w: w \in \text{frontier}(\text{connected\_component\_set } (\text{interior } S) z)$ 
  by auto
  then have  $\text{norm } (f z) = \text{norm } (f w)$  by (simp add: 2 c cc frontier_def)
  also have  $\dots \leq B$ 
  using  $w \text{ frontier\_interior\_subset frontier\_of\_connected\_component\_subset}$ 
  by (blast intro: leB)
  finally show ?thesis .
qed
then show ?thesis
  using  $z \in \{ \xi \in S \mid \text{closure\_subset } \}$  by fastforce
qed

corollary maximum_real_frontier:
  assumes  $\text{holf}: f \text{ holomorphic\_on } (\text{interior } S)$ 
  and  $\text{contf}: \text{continuous\_on } (\text{closure } S) f$ 

```

```

and bos: bounded S
and leB:  $\bigwedge z. z \in \text{frontier } S \implies \text{Re}(f z) \leq B$ 
and  $\xi \in S$ 
shows  $\text{Re}(f \xi) \leq B$ 
using maximum_modulus_frontier [of exp o f S exp B]
  Transcendental.continuous_on_exp holomorphic_on_compose holomorphic_on_exp
  assms
by auto

```

5.4 Factoring out a zero according to its order

lemma holomorphic_factor_order_of_zero:

```

assumes holf: f holomorphic_on S
and os: open S
and  $\xi \in S$   $0 < n$ 
and dnz:  $(\text{deriv } \sim^n) f \xi \neq 0$ 
and dfz:  $\bigwedge i. \llbracket 0 < i; i < n \rrbracket \implies (\text{deriv } \sim^i) f \xi = 0$ 
obtains g r where  $0 < r$ 
  g holomorphic_on ball  $\xi$  r
   $\bigwedge w. w \in \text{ball } \xi$  r  $\implies f w - f \xi = (w - \xi)^n * g w$ 
   $\bigwedge w. w \in \text{ball } \xi$  r  $\implies g w \neq 0$ 

```

proof –

```

obtain r where  $r > 0$  and r: ball  $\xi$  r  $\subseteq S$  using assms by (blast elim!: openE)
then have holfb: f holomorphic_on ball  $\xi$  r
using holf holomorphic_on_subset by blast
define g where  $g w = \text{suminf } (\lambda i. (\text{deriv } \sim^{i+n}) f \xi / (\text{fact}(i+n)) * (w - \xi)^i)$  for w
have sumsg:  $(\lambda i. (\text{deriv } \sim^{i+n}) f \xi / (\text{fact}(i+n)) * (w - \xi)^i)$  sums g w
and feq:  $f w - f \xi = (w - \xi)^n * g w$ 
  if w: w  $\in$  ball  $\xi$  r for w
proof –
define powf where  $\text{powf} = (\lambda i. (\text{deriv } \sim^i) f \xi / (\text{fact } i) * (w - \xi)^i)$ 
have [simp]: powf 0 = f  $\xi$ 
by (simp add: powf_def)
have sing:  $\{..<n\} - \{i. \text{powf } i = 0\} = (\text{if } f \xi = 0 \text{ then } \{\} \text{ else } \{0\})$ 
unfolding powf_def using  $\langle 0 < n \rangle$  dfz by (auto simp: dfz; metis funpow_0
not_gr0)

```

```

have powf sums f w
unfolding powf_def by (rule holomorphic_power_series [OF holfb w])
moreover have  $(\sum i < n. \text{powf } i) = f \xi$ 
by (subst sum.setdiff_irrelevant [symmetric]; simp add: dfz sing)
ultimately have fsums:  $(\lambda i. \text{powf } (i+n))$  sums (f w - f  $\xi$ )
using w sums_iff_shift' by metis
then have *: summable  $(\lambda i. (w - \xi)^n * ((\text{deriv } \sim^{i+n}) f \xi * (w - \xi)^i / \text{fact } (i+n)))$ 
unfolding powf_def using sums_summable
by (auto simp: power_add mult_ac)
have summable  $(\lambda i. (\text{deriv } \sim^{i+n}) f \xi * (w - \xi)^i / \text{fact } (i+n))$ 
proof (cases w= $\xi$ )

```

```

    case False then show ?thesis
      using summable_mult [OF *, of 1 / (w - ξ) ^ n] by simp
    next
      case True then show ?thesis
        by (auto simp: Power.semiring_1_class.power_0_left intro!: summable_finite
            [of {0}])
            split: if_split_asm)
      qed
      then show sumsg: (λi. (deriv ^^ (i + n)) f ξ / (fact(i + n)) * (w - ξ) ^ i)
sums g w
        by (simp add: summable_sums_iff g_def)
      show f w - f ξ = (w - ξ) ^ n * g w
        using sums_mult [OF sumsg, of (w - ξ) ^ n]
        by (intro sums_unique2 [OF fsums]) (auto simp: power_add mult_ac powf_def)
      qed
      then have holg: g holomorphic_on ball ξ r
        by (meson sumsg power_series_holomorphic)
      then have contg: continuous_on (ball ξ r) g
        by (blast intro: holomorphic_on_imp_continuous_on)
      have g ξ ≠ 0
        using dnz_unfolding g_def
        by (subst suminf_finite [of {0}]) auto
      obtain d where 0 < d and d: ∧w. w ∈ ball ξ d ⇒ g w ≠ 0
        using ⟨0 < r⟩ continuous_on_avoid [OF contg _ ⟨g ξ ≠ 0⟩]
        by (metis centre_in_ball le_cases mem_ball mem_ball_leI)
      show ?thesis
      proof
        show g holomorphic_on ball ξ (min r d)
          using holg by (auto simp: feq holomorphic_on_subset subset_ball d)
        qed (use ⟨0 < r⟩ ⟨0 < d⟩ in ⟨auto simp: feq d⟩)
      qed

```

lemma *holomorphic_factor_order_of_zero_strong*:

assumes *holf*: f holomorphic_on S open S ξ ∈ S 0 < n

and (deriv ^^ n) f ξ ≠ 0

and ∧i. [0 < i; i < n] ⇒ (deriv ^^ i) f ξ = 0

obtains g r where 0 < r

g holomorphic_on ball ξ r

∧w. w ∈ ball ξ r ⇒ f w - f ξ = ((w - ξ) * g w) ^ n

∧w. w ∈ ball ξ r ⇒ g w ≠ 0

proof –

obtain g r where 0 < r

and *holg*: g holomorphic_on ball ξ r

and *feq*: ∧w. w ∈ ball ξ r ⇒ f w - f ξ = (w - ξ) ^ n * g w

and *gne*: ∧w. w ∈ ball ξ r ⇒ g w ≠ 0

by (auto intro: holomorphic_factor_order_of_zero [OF *assms*])

have *con*: continuous_on (ball ξ r) (λz. deriv g z / g z)

by (rule continuous_intros) (auto simp: gne holg holomorphic_deriv holomor-

```

phic_on_imp_continuous_on)
  have cd: ( $\lambda z. \text{deriv } g \ z / g \ z$ ) field_differentiable at  $x$  if  $\text{dist } \xi \ x < r$  for  $x$ 
  proof (intro derivative_intros)
    show  $\text{deriv } g$  field_differentiable at  $x$ 
      using that holg mem_ball by (blast intro: holomorphic_deriv holomorphic_on_imp_differentiable_at)
    show  $g$  field_differentiable at  $x$ 
      by (metis that open_ball at_within_open holg holomorphic_on_def mem_ball)
    qed (simp add: gne that)
  obtain  $h$  where  $h: \bigwedge x. x \in \text{ball } \xi \ r \implies (h \text{ has\_field\_derivative } \text{deriv } g \ x / g \ x)$  (at  $x$ )
    using holomorphic_convex_primitive [of ball  $\xi \ r$   $\{\}$   $\lambda z. \text{deriv } g \ z / g \ z$ ]
      by (metis (no_types, lifting) Diff_empty_at_within_interior cd convex_ball_infinite_imp_nonempty_interior_ball mem_ball)
  then have continuous_on (ball  $\xi \ r$ )  $h$ 
    by (metis open_ball holomorphic_on_imp_continuous_on holomorphic_on_open)
  then have con: continuous_on (ball  $\xi \ r$ ) ( $\lambda x. \text{exp } (h \ x) / g \ x$ )
    by (auto intro!: continuous_intros simp add: holg holomorphic_on_imp_continuous_on gne)
  have gfd:  $\text{dist } \xi \ x < r \implies g$  field_differentiable at  $x$  if  $\text{dist } \xi \ x < r$  for  $x$ 
    using holg holomorphic_on_imp_differentiable_at by auto
  have 0:  $\text{dist } \xi \ x < r \implies ((\lambda x. \text{exp } (h \ x) / g \ x) \text{ has\_field\_derivative } 0)$  (at  $x$ )
  for  $x$ 
    by (rule gfd h derivative_eq_intros DERIV_deriv_iff_field_differentiable [THEN iffD2] | simp add: gne)+
  obtain  $c$  where  $c: \bigwedge x. x \in \text{ball } \xi \ r \implies \text{exp } (h \ x) / g \ x = c$ 
    by (rule DERIV_zero_connected_constant [of ball  $\xi \ r$   $\{\}$   $\lambda x. \text{exp } (h \ x) / g \ x$ ])
    (auto simp: con 0)
  have hol: ( $\lambda z. \text{exp } ((Ln (\text{inverse } c) + h \ z) / \text{of\_nat } n)$ ) holomorphic_on ball  $\xi \ r$ 
    proof (intro holomorphic_intros holomorphic_on_compose [unfolded o_def,
  where  $g = \text{exp}$ ])
      show  $h$  holomorphic_on ball  $\xi \ r$ 
        using  $h$  holomorphic_on_open by blast
      qed (use  $\langle 0 < n \rangle$  in auto)
    show ?thesis
      proof
        show  $\bigwedge w. w \in \text{ball } \xi \ r \implies f \ w - f \ \xi = ((w - \xi) * \text{exp } ((Ln (\text{inverse } c) + h \ w) / \text{of\_nat } n)) ^ n$ 
          using  $\langle 0 < n \rangle$ 
          by (auto simp: feq power_mult_distrib exp_divide_power_eq exp_add gne simp flip: c)
        qed (use hol  $\langle 0 < r \rangle$  in auto)
      qed
  qed

```

lemma

```

fixes  $k :: 'a::wellorder$ 
assumes  $a\_def: a == \text{LEAST } x. P \ x$  and  $P: P \ k$ 
shows  $\text{def\_LeastI}: P \ a$  and  $\text{def\_Least\_le}: a \leq k$ 

```

```

unfolding a_def
by (rule LeastI Least_le; rule P)+

lemma holomorphic_factor_zero_nonconstant:
assumes holf: f holomorphic_on S and S: open S connected S
and  $\xi \in S$  f  $\xi = 0$ 
and nonconst:  $\neg$  f constant_on S
obtains g r n
where  $0 < n$   $0 < r$  ball  $\xi$  r  $\subseteq S$ 
g holomorphic_on ball  $\xi$  r
 $\bigwedge w. w \in \text{ball } \xi \text{ r} \implies f w = (w - \xi)^{\wedge n} * g w$ 
 $\bigwedge w. w \in \text{ball } \xi \text{ r} \implies g w \neq 0$ 
proof (cases  $\forall n > 0. (\text{deriv } \wedge n) f \xi = 0$ )
case True then show ?thesis
using holomorphic_fun_eq_const_on_connected [OF holf S  $\langle \xi \in S \rangle$ ] non-
const by (simp add: constant_on_def)
next
case False
then obtain n0 where n0 > 0 and n0: (deriv  $\wedge$  n0) f  $\xi \neq 0$  by blast
obtain r0 where r0 > 0 ball  $\xi$  r0  $\subseteq S$  using S openE  $\langle \xi \in S \rangle$  by auto
define n where n  $\equiv$  LEAST n. (deriv  $\wedge$  n) f  $\xi \neq 0$ 
have n_ne: (deriv  $\wedge$  n) f  $\xi \neq 0$ 
by (rule def_LeastI [OF n_def]) (rule n0)
then have  $0 < n$  using  $\langle f \xi = 0 \rangle$ 
using funpow_0 by fastforce
have n_min:  $\bigwedge k. k < n \implies (\text{deriv } \wedge k) f \xi = 0$ 
using def_Least_le [OF n_def] not_le by blast
then obtain g r1
where g:  $0 < r1$  g holomorphic_on ball  $\xi$  r1
and geq:  $\bigwedge w. w \in \text{ball } \xi \text{ r1} \implies f w = (w - \xi)^{\wedge n} * g w$ 
and g0:  $\bigwedge w. w \in \text{ball } \xi \text{ r1} \implies g w \neq 0$ 
by (auto intro: holomorphic_factor_order_of_zero [OF holf  $\langle \text{open } S \rangle \langle \xi \in S \rangle$ 
 $\langle n > 0 \rangle$  n_ne] simp:  $\langle f \xi = 0 \rangle$ )
show ?thesis
proof
show g holomorphic_on ball  $\xi$  (min r0 r1)
using g by auto
show  $\bigwedge w. w \in \text{ball } \xi \text{ (min r0 r1)} \implies f w = (w - \xi)^{\wedge n} * g w$ 
by (simp add: geq)
qed (use  $\langle 0 < n \rangle \langle 0 < r0 \rangle \langle 0 < r1 \rangle \langle \text{ball } \xi \text{ r0} \subseteq S \rangle g0$  in auto)
qed

```

```

lemma holomorphic_lower_bound_difference:
assumes holf: f holomorphic_on S and S: open S connected S
and  $\xi \in S$  and  $\varphi \in S$ 
and fne: f  $\varphi \neq f \xi$ 
obtains k n r
where  $0 < k$   $0 < r$ 

```



```

      ball ξ r ⊆ S
      ∧ w. w ∈ ball ξ r ⇒ k * norm(w - ξ) ^ n ≤ norm(f w - f ξ)
proof -
  define n where n = (LEAST n. 0 < n ∧ (deriv ~ n) f ξ ≠ 0)
  obtain n0 where 0 < n0 and n0: (deriv ~ n0) f ξ ≠ 0
    using fne holomorphic_fun_eq_const_on_connected [OF hol f S] ⟨ξ ∈ S⟩ ⟨φ ∈ S⟩ by blast
  then have 0 < n and n_ne: (deriv ~ n) f ξ ≠ 0
    unfolding n_def by (metis (mono_tags, lifting) LeastI)+
  have n_min: ∧k. [0 < k; k < n] ⇒ (deriv ~ k) f ξ = 0
    unfolding n_def by (blast dest: not_less_Least)
  then obtain g r
    where 0 < r and holg: g holomorphic_on ball ξ r
      and fne: ∧w. w ∈ ball ξ r ⇒ f w - f ξ = (w - ξ) ^ n * g w
      and gnz: ∧w. w ∈ ball ξ r ⇒ g w ≠ 0
      by (auto intro: holomorphic_factor_order_of_zero [OF hol f ⟨open S⟩ ⟨ξ ∈ S⟩ ⟨n > 0⟩ n_ne])
  obtain e where e > 0 and e: ball ξ e ⊆ S using assms by (blast elim!: openE)
  then have holfb: f holomorphic_on ball ξ e
    using hol f holomorphic_on_subset by blast
  define d where d = (min e r) / 2
  have 0 < d using ⟨0 < r⟩ ⟨0 < e⟩ by (simp add: d_def)
  have d < r
    using ⟨0 < r⟩ by (auto simp: d_def)
  then have cbb: cball ξ d ⊆ ball ξ r
    by (auto simp: cball_subset_ball_iff)
  then have g holomorphic_on cball ξ d
    by (rule holomorphic_on_subset [OF holg])
  then have closed (g ' cball ξ d)
    by (simp add: compact_imp_closed compact_continuous_image holomorphic_on_imp_continuous_on)
  moreover have g ' cball ξ d ≠ {}
    using ⟨0 < d⟩ by auto
  ultimately obtain x where x: x ∈ g ' cball ξ d and ∧y. y ∈ g ' cball ξ d ⇒
    dist 0 x ≤ dist 0 y
    by (rule distance_attains_inf) blast
  then have leg: ∧w. w ∈ cball ξ d ⇒ norm x ≤ norm (g w)
    by auto
  have ball ξ d ⊆ cball ξ d by auto
  also have ... ⊆ ball ξ e using ⟨0 < d⟩ d_def by auto
  also have ... ⊆ S by (rule e)
  finally have dS: ball ξ d ⊆ S .
  have x ≠ 0 using gnz x ⟨d < r⟩ by auto
  show thesis
proof
  show ∧w. w ∈ ball ξ d ⇒ cmod x * cmod (w - ξ) ^ n ≤ cmod (f w - f ξ)
    using ⟨d < r⟩ leg by (auto simp: fne norm_mult norm_power algebra_simps
    mult_right_mono)
  qed (use dS ⟨x ≠ 0⟩ ⟨d > 0⟩ in auto)
qed

```

lemma

assumes *holg*: f holomorphic_on $(S - \{\xi\})$ **and** $\xi: \xi \in \text{interior } S$

shows *holomorphic_on_extend_lim*:

$(\exists g. g \text{ holomorphic_on } S \wedge (\forall z \in S - \{\xi\}. g z = f z)) \longleftrightarrow$
 $((\lambda z. (z - \xi) * f z) \longrightarrow 0) \text{ (at } \xi)$
(is $?P = ?Q)$

and *holomorphic_on_extend_bounded*:

$(\exists g. g \text{ holomorphic_on } S \wedge (\forall z \in S - \{\xi\}. g z = f z)) \longleftrightarrow$
 $(\exists B. \text{eventually } (\lambda z. \text{norm}(f z) \leq B) \text{ (at } \xi))$
(is $?P = ?R)$

proof –

obtain δ **where** $0 < \delta$ **and** $\delta: \text{ball } \xi \ \delta \subseteq S$

using $\xi \text{ mem_interior}$ **by** *blast*

have $?R$ **if** *holg*: g holomorphic_on S **and** *gf*: $\bigwedge z. z \in S - \{\xi\} \implies g z = f z$

for g

proof –

have \S : $\text{cmod } (f x) \leq \text{cmod } (g \ \xi) + 1$ **if** $x \neq \xi$ $\text{dist } x \ \xi < \delta$ $\text{dist } (g \ x) \ (g \ \xi) < 1$ **for** x

proof –

have $x \in S$

by (*metis* $\delta \ \text{dist_commute mem_ball subsetD that}(2)$)

with *that gf [of x]* **show** $?thesis$

using *norm_triangle_ineq2 [of f x g \xi]* *dist_complex_def* **by** *auto*

qed

then have $*$: $\forall_F z \text{ in at } \xi. \text{dist } (g z) \ (g \ \xi) < 1 \longrightarrow \text{cmod } (f z) \leq \text{cmod } (g \ \xi) + 1$

using $\langle 0 < \delta \rangle$ *eventually_at* **by** *blast*

have *continuous_on* (*interior* S) g

by (*meson continuous_on_subset holg holomorphic_on_imp_continuous_on interior_subset*)

then have $\bigwedge x. x \in \text{interior } S \implies (g \longrightarrow g \ x) \text{ (at } x)$

using *continuous_on_interior continuous_within holg holomorphic_on_imp_continuous_on*

by *blast*

then have $(g \longrightarrow g \ \xi) \text{ (at } \xi)$

by (*simp add: \xi*)

then have $\forall_F z \text{ in at } \xi. \text{cmod } (f z) \leq \text{cmod } (g \ \xi) + 1$

by (*rule eventually_mp [OF * tendstoD [where e=1]]*, *auto*)

then show $?thesis$

by *blast*

qed

moreover have $?Q$ **if** $\forall_F z \text{ in at } \xi. \text{cmod } (f z) \leq B$ **for** B

by (*rule lim_null_mult_right_bounded [OF _ that]*) (*simp add: LIM_zero*)

moreover have $?P$ **if** $(\lambda z. (z - \xi) * f z) \xrightarrow{-\xi} 0$

proof –

define h **where** [*abs_def*]: $h z = (z - \xi)^{\wedge 2} * f z$ **for** z

have $(\lambda y. (y - \xi)^2 * f y / (y - \xi)) \xrightarrow{-\xi} 0$

by (*simp add: LIM_cong power2_eq_square that*)

then have $h0$: (*h has_field_derivative* 0) (at ξ)

```

    by (simp add: h_def has_field_derivative_iff)
  have holh: h holomorphic_on S
  proof (simp add: holomorphic_on_def, clarify)
    fix z assume z ∈ S
    show h field_differentiable at z within S
    proof (cases z = ξ)
      case True then show ?thesis
        using field_differentiable_at_within field_differentiable_def h0 by blast
      next
      case False
      then have f field_differentiable at z within S
        using holomorphic_onD [OF holh, of z] ⟨z ∈ S⟩
        unfolding field_differentiable_def has_field_derivative_iff
        by (force intro: exI [where x=dist ξ z] elim: Lim_transform_within_set
[unfolded eventually_at])
      then show ?thesis
        by (simp add: h_def power2_eq_square derivative_intros)
    qed
  qed
  define g where [abs_def]: g z = (if z = ξ then deriv h ξ else (h z - h ξ) / (z
- ξ)) for z
  have §: ∀ z ∈ S - {ξ}. (g z - g ξ) / (z - ξ) = f z
  using h0 by (auto simp: g_def power2_eq_square divide_simps DERIV_imp_deriv
h_def)
  have g holomorphic_on S
  unfolding g_def by (rule pole_lemma [OF holh ξ])
  then have (λz. if z = ξ then deriv g ξ else (g z - g ξ) / (z - ξ)) holomorphic_on
S
  using ξ pole_lemma by blast
  then show ?thesis
  using § by (smt (verit, best) DiffD2 singletonI)
  qed
  ultimately show ?P = ?Q and ?P = ?R
  by meson+
  qed

lemma pole_at_infinity:
  assumes holh: f holomorphic_on UNIV and lim: ((inverse o f) ⟶ l) at_infinity
  obtains a n where ∧z. f z = (∑ i ≤ n. a i * z ^ i)
  proof (cases l = 0)
    case False
    show thesis
    proof
      show f z = (∑ i ≤ 0. inverse l * z ^ i) for z
      using False tendsto_inverse [OF lim] by (simp add: Liouville_weak [OF
holh])
    qed
  next
  case True

```

```

then have [simp]:  $l = 0$  .
show ?thesis
proof (cases  $\exists r. 0 < r \wedge (\forall z \in \text{ball } 0 \ r - \{0\}. f(\text{inverse } z) \neq 0)$ )
  case True
    then obtain  $r$  where  $0 < r$  and  $r: \bigwedge z. z \in \text{ball } 0 \ r - \{0\} \implies f(\text{inverse } z) \neq 0$ 
      by auto
      have 1:  $\text{inverse} \circ f \circ \text{inverse}$  holomorphic_on ball 0 r - {0}
      by (rule holomorphic_on_compose holomorphic_intros holomorphic_on_subset [OF holf] | force simp: r)+
      have 2:  $0 \in \text{interior } (\text{ball } 0 \ r)$ 
      using  $\langle 0 < r \rangle$  by simp
      obtain  $g$  where holg:  $g$  holomorphic_on ball 0 r
      and geq:  $\bigwedge z. z \in \text{ball } 0 \ r - \{0\} \implies g \ z = (\text{inverse} \circ f \circ \text{inverse}) \ z$ 
      using tendstoD [OF lim [unfolded lim_at_infinity_0] zero_less_one] holomorphic_on_extend_bounded [OF 1 2]
      by (smt (verit, del_insts)  $\langle l = 0 \rangle$  eventually_mono norm_conv_dist)
      have ifi0:  $(\text{inverse} \circ f \circ \text{inverse}) - 0 \rightarrow 0$ 
      using  $\langle l = 0 \rangle$  lim_lim_at_infinity_0 by blast
      have g2g0:  $g - 0 \rightarrow g \ 0$ 
      using  $\langle 0 < r \rangle$  centre_in_ball continuous_at_continuous_on_eq_continuous_at holg
      by (blast intro: holomorphic_on_imp_continuous_on)
      have g2g1:  $g - 0 \rightarrow 0$ 
      proof (rule Lim_transform_within_open [OF ifi0 open_ball])
        show  $\bigwedge x. \llbracket x \in \text{ball } 0 \ r; x \neq 0 \rrbracket \implies (\text{inverse} \circ f \circ \text{inverse}) \ x = g \ x$ 
        by (auto simp: geq)
      qed (auto simp:  $\langle 0 < r \rangle$ )
      have [simp]:  $g \ 0 = 0$ 
      by (rule tendsto_unique [OF _ g2g0 g2g1]) simp
      have ball 0 r - {0::complex}  $\neq \{\}$ 
      using  $\langle 0 < r \rangle$  by (metis 2 Diff_iff insert_Diff interior_ball interior_singleton)
      then obtain  $w::\text{complex}$  where  $w \neq 0$  and  $w: \text{norm } w < r$  by force
      then have  $g \ w \neq 0$  by (simp add: geq r)
      obtain  $B \ n \ e$  where  $0 < B \ 0 < e \ e \leq r$ 
      and leg:  $\bigwedge w. \text{norm } w < e \implies B * \text{cmod } w ^ n \leq \text{cmod } (g \ w)$ 
      proof (rule holomorphic_lower_bound_difference [OF holg open_ball connected_ball])
        show  $g \ w \neq g \ 0$ 
        by (simp add:  $\langle g \ w \neq 0 \rangle$ )
        show  $w \in \text{ball } 0 \ r$ 
        using mem_ball_0 w by blast
      qed (use  $\langle 0 < r \rangle$  in  $\langle \text{auto simp: ball_subset_ball_iff} \rangle$ )
      have  $\S: \text{cmod } (f \ z) \leq \text{cmod } z ^ n / B$  if  $2/e \leq \text{cmod } z$  for  $z$ 
      proof -
        have ize:  $\text{inverse } z \in \text{ball } 0 \ e - \{0\}$  using that  $\langle 0 < e \rangle$ 
        by (auto simp: norm_divide field_split_simps algebra_simps)
        then have [simp]:  $z \neq 0$  and izr:  $\text{inverse } z \in \text{ball } 0 \ r - \{0\}$  using  $\langle e \leq r \rangle$ 

```

```

    by auto
  then have [simp]: f z ≠ 0
    using r [of inverse z] by simp
  have [simp]: f z = inverse (g (inverse z))
    using izr geq [of inverse z] by simp
  show ?thesis using ize leg [of inverse z] ‹0 < B› ‹0 < e›
    by (simp add: field_split_simps norm_divide algebra_simps)
qed
show thesis
proof
  show f z = (∑ i≤n. (deriv ^ i) f 0 / fact i * z ^ i) for z
    using § by (rule_tac A = 2/e and B = 1/B in Liouville_polynomial
[OF holf], simp)
  qed
next
case False
  then have fi0: ∧r. r > 0 ⇒ ∃ z∈ball 0 r - {0}. f (inverse z) = 0
    by simp
  have fz0: f z = 0 if 0 < r and lt1: ∧x. x ≠ 0 ⇒ cmod x < r ⇒ inverse
(cmod (f (inverse x))) < 1
    for z r
  proof -
    have f0: (f ⟶ 0) at_infinity
    proof -
      have DIM_complex[intro]: 2 ≤ DIM(complex) — should not be necessary!
        by simp
      have f (inverse x) ≠ 0 ⇒ x ≠ 0 ⇒ cmod x < r ⇒ 1 < cmod (f (inverse
x)) for x
        using lt1[of x] by (auto simp: field_simps)
      then have **: cmod (f (inverse x)) ≤ 1 ⇒ x ≠ 0 ⇒ cmod x < r ⇒ f
(inverse x) = 0 for x
        by force
      then have *: (f ∘ inverse) ‘ (ball 0 r - {0}) ⊆ {0} ∪ - ball 0 1
        by force
      have continuous_on (inverse ‘ (ball 0 r - {0})) f
        using continuous_on_subset holf holomorphic_on_imp_continuous_on
by blast
      then have connected ((f ∘ inverse) ‘ (ball 0 r - {0}))
        using connected_punctured_ball
        by (intro connected_continuous_image continuous_intros; force)
      then have {0} ∩ (f ∘ inverse) ‘ (ball 0 r - {0}) = {} ∨ - ball 0 1 ∩ (f ∘
inverse) ‘ (ball 0 r - {0}) = {}
        by (rule connected_closedD) (use * in auto)
      then have f (inverse w) = 0 if w ≠ 0 cmod w < r for w
        using **[of w] fi0 ‹0 < r› that by force
      then show ?thesis
        unfolding lim_at_infinity_0
        using eventually_at ‹r > 0› by (force simp: intro: tendsto_eventually)
    qed

```

```

obtain  $w$  where  $w \in \text{ball } 0 \ r - \{0\}$  and  $f (\text{inverse } w) = 0$ 
using  $\text{False } \langle 0 < r \rangle$  by blast
then show ?thesis
by (auto simp: f0 Liouville_weak [OF holf, of 0])
qed
show thesis
proof
show  $\bigwedge z. f z = (\sum_{i \leq 0}. 0 * z^i)$ 
using lim
apply (simp add: lim_at_infinity_0 Lim_at_dist_norm norm_inverse)
using fz0 zero_less_one by blast
qed
qed
qed

```

5.5 Entire proper functions are precisely the non-trivial polynomials

```

lemma proper_map_polyfun:
fixes  $c :: \text{nat} \Rightarrow 'a::\{\text{real\_normed\_div\_algebra, heine\_borel}\}$ 
assumes closed S and compact K and  $c: c \ i \neq 0 \ 1 \leq i \leq n$ 
shows compact (S ∩ {z. (∑_{i ≤ n}. c i * z^i) ∈ K})
proof -
obtain  $B$  where  $B > 0$  and  $B: \bigwedge x. x \in K \implies \text{norm } x \leq B$ 
by (metis compact_imp_bounded <compact K> bounded_pos)
have  $*$ :  $\text{norm } x \leq b$ 
if  $\bigwedge x. b \leq \text{norm } x \implies B + 1 \leq \text{norm } (\sum_{i \leq n}. c \ i * x^i)$ 
 $(\sum_{i \leq n}. c \ i * x^i) \in K$  for  $b \ x$ 
proof -
have  $\text{norm } (\sum_{i \leq n}. c \ i * x^i) \leq B$ 
using  $B$  that by blast
moreover have  $\neg B + 1 \leq B$ 
by simp
ultimately show  $\text{norm } x \leq b$ 
using that by (metis (no_types) less_eq_real_def not_less order_trans)
qed
have bounded {z. (∑_{i ≤ n}. c i * z^i) ∈ K}
using Limits.polyfun_extremal [where c=c and B=B+1, OF c]
by (auto simp: bounded_pos eventually_at_infinity_pos *)
moreover have closed ((λz. (∑_{i ≤ n}. c i * z^i)) -' K)
using <compact K> compact_eq_bounded_closed
by (intro allI continuous_closed_vimage continuous_intros; force)
ultimately show ?thesis
using closed_Int_compact [OF <closed S>] compact_eq_bounded_closed
by (auto simp add: vimage_def)
qed

```

```

lemma proper_map_polyfun_univ:
fixes  $c :: \text{nat} \Rightarrow 'a::\{\text{real\_normed\_div\_algebra, heine\_borel}\}$ 

```

```

assumes compact  $K$   $c\ i \neq 0$   $1 \leq i \leq n$ 
shows compact  $\{z. (\sum_{i \leq n}. c\ i * z^i) \in K\}$ 
using proper_map_polyfun [of UNIV  $K\ c\ i\ n$ ] assms by simp

lemma proper_map_polyfun_eq:
assumes  $f$  holomorphic_on UNIV
shows  $(\forall k. \text{compact } k \longrightarrow \text{compact } \{z. f\ z \in k\}) \longleftrightarrow$ 
 $(\exists c\ n. 0 < n \wedge (c\ n \neq 0) \wedge f = (\lambda z. \sum_{i \leq n}. c\ i * z^i))$ 
(is ?lhs = ?rhs)

proof
assume compf [rule_format]: ?lhs
have 2:  $\exists k. 0 < k \wedge a\ k \neq 0 \wedge f = (\lambda z. \sum_{i \leq k}. a\ i * z^i)$ 
if  $\bigwedge z. f\ z = (\sum_{i \leq n}. a\ i * z^i)$  for  $a\ n$ 
proof (cases  $\forall i \leq n. 0 < i \longrightarrow a\ i = 0$ )
case True
then have [simp]:  $\bigwedge z. f\ z = a\ 0$ 
by (simp add: that sum.atMost_shift)
have False using compf [of  $\{a\ 0\}$ ] by simp
then show ?thesis ..
next
case False
then obtain  $k$  where  $k: 0 < k\ k \leq n\ a\ k \neq 0$  by force
define  $m$  where  $m = (\text{GREATEST } k. k \leq n \wedge a\ k \neq 0)$ 
have  $m: m \leq n \wedge a\ m \neq 0$ 
unfolding m_def
using GreatestI_nat [where  $b = n$ ]  $k$  by (metis (mono_tags, lifting))
have [simp]:  $a\ i = 0$  if  $m < i \leq n$  for  $i$ 
using Greatest_le_nat [where  $b = n$  and  $P = \lambda k. k \leq n \wedge a\ k \neq 0$ ]
using m_def not_le that by auto
have  $k \leq m$ 
unfolding m_def
using Greatest_le_nat [where  $b = n$ ]  $k$  by (metis (mono_tags, lifting))
with  $k\ m$  show ?thesis
by (rule_tac  $x=m$  in exI) (auto simp: that comm_monoid_add_class.sum_mono_neutral_right)
qed
have *:  $((\text{inverse} \circ f) \longrightarrow 0)$  at_infinity
proof (rule Lim_at_infinityI)
fix  $e::\text{real}$  assume  $0 < e$ 
with compf [of cball 0 (inverse e)]
show  $\exists B. \forall x. B \leq \text{cmod } x \longrightarrow \text{dist } ((\text{inverse} \circ f)\ x)\ 0 \leq e$ 
apply (clarsimp simp: compact_eq_bounded_closed norm_divide divide_simps
mult.commute elim!: bounded_normE_less)
by (meson linorder_not_le nle_le)
qed
then obtain  $a\ n$  where  $\bigwedge z. f\ z = (\sum_{i \leq n}. a\ i * z^i)$ 
using assms pole_at_infinity by blast
with * 2 show ?rhs by blast
next
assume ?rhs

```

then obtain $c\ n$ where $0 < n\ c\ n \neq 0\ f = (\lambda z. \sum_{i \leq n}. c\ i * z^i)$ by *blast*
then have *compact* $\{z. f\ z \in k\}$ if *compact* k for k
by (*auto intro: proper_map_polyfun_univ [OF that]*)
then show *?lhs* by *blast*
qed

5.6 Relating invertibility and nonvanishing of derivative

lemma *has_complex_derivative_locally_injective*:

assumes *holf*: f holomorphic_on S
and S : $\xi \in S$ open S
and *dnz*: $\text{deriv } f\ \xi \neq 0$
obtains r where $r > 0$ ball $\xi\ r \subseteq S$ *inj_on* f (ball $\xi\ r$)
proof –
have $*$: $\exists d > 0. \forall x. \text{dist } \xi\ x < d \longrightarrow \text{onorm } (\lambda v. \text{deriv } f\ x * v - \text{deriv } f\ \xi * v) < e$ if $e > 0$ for e
proof –
have *contdf*: *continuous_on* S (*deriv* f)
by (*simp add: holf holomorphic_deriv holomorphic_on_imp_continuous_on* \langle open S \rangle)
obtain δ where $\delta > 0$ and δ : $\bigwedge x. \llbracket x \in S; \text{dist } x\ \xi \leq \delta \rrbracket \implies \text{cmod } (\text{deriv } f\ x - \text{deriv } f\ \xi) \leq e/2$
using *continuous_onE* [*OF contdf* $\langle \xi \in S \rangle$, *of* $e/2$] $\langle 0 < e \rangle$
by (*metis dist_complex_def half_gt_zero less_imp_le*)
have \S : $\bigwedge \zeta z. \llbracket \zeta \in S; \text{dist } \xi\ \zeta < \delta \rrbracket \implies \text{cmod } (\text{deriv } f\ \zeta - \text{deriv } f\ \xi) * \text{cmod } z \leq e/2 * \text{cmod } z$
by (*intro mult_right_mono [OF δ]*) (*auto simp: dist_commute*)
obtain ε where $\varepsilon > 0$ ball $\xi\ \varepsilon \subseteq S$
by (*metis openE [OF \langle open S $\rangle \langle \xi \in S \rangle$]*)
with $\langle \delta > 0 \rangle$ have $\exists \delta > 0. \forall x. \text{dist } \xi\ x < \delta \longrightarrow \text{onorm } (\lambda v. \text{deriv } f\ x * v - \text{deriv } f\ \xi * v) \leq e/2$
using \S
apply (*rule_tac* $x = \text{min } \delta\ \varepsilon$ in *exI*)
apply (*intro conjI allI impI Operator_Norm.onorm_le*)
apply (*force simp: mult_right_mono norm_mult [symmetric] left_diff_distrib* δ)
done
with $\langle e > 0 \rangle$ show *?thesis* by *force*
qed
have *inj* ($*$) (*deriv* $f\ \xi$)
using *dnz* by *simp*
then obtain g' where g' : *linear* $g'\ g' \circ (*)$ (*deriv* $f\ \xi$) = *id*
using *linear_injective_left_inverse* [*of* $*$] (*deriv* $f\ \xi$)] by *auto*

have *fder*: $\bigwedge x. x \in S \implies (f \text{ has_derivative } (*)$ (*deriv* $f\ x$)) (*at* x)
using \langle open S \rangle *has_field_derivative_imp_has_derivative holf holomorphic_derivI*
by *blast*
show *?thesis*
apply (*rule has_derivative_locally_injective [OF S, where f=f and f' = λz*

*h. deriv f z * h and g' = g')*
using *g' * by (simp_all add: fder linear_conv_bounded_linear that)*
qed

lemma *has_complex_derivative_locally_invertible:*

assumes *holf: f holomorphic_on S*
and *S: $\xi \in S$ open S*
and *dnz: deriv f $\xi \neq 0$*
obtains *r where $r > 0$ ball ξ r $\subseteq S$ open ($f' (ball \xi r)$) inj_on f (ball ξ r)*
proof –
obtain *r where $r > 0$ ball ξ r $\subseteq S$ inj_on f (ball ξ r)*
by (*blast intro: that has_complex_derivative_locally_injective [OF assms]*)
then have *$\xi: \xi \in ball \xi r$ by simp*
then have *nc: \neg f constant_on ball ξ r*
using *$\langle inj_on\ f\ (ball\ \xi\ r) \rangle$ injective_not_constant by fastforce*
have *holf': f holomorphic_on ball ξ r*
using *$\langle ball\ \xi\ r \subseteq S \rangle$ holf holomorphic_on_subset by blast*
have *open ($f' ball \xi r$)*
by (*simp add: $\langle inj_on\ f\ (ball\ \xi\ r) \rangle$ holf' open_mapping_thm3*)
then show *?thesis*
using *$\langle 0 < r \rangle \langle ball\ \xi\ r \subseteq S \rangle \langle inj_on\ f\ (ball\ \xi\ r) \rangle$ that by blast*
qed

lemma *holomorphic_injective_imp_regular:*

assumes *holf: f holomorphic_on S*
and *open S and injf: inj_on f S*
and *$\xi \in S$*
shows *deriv f $\xi \neq 0$*
proof –
obtain *r where $r > 0$ and $r: ball \xi r \subseteq S$ using assms by (blast elim!: openE)*
have *holf': f holomorphic_on ball ξ r*
using *$\langle ball \xi r \subseteq S \rangle$ holf holomorphic_on_subset by blast*
show *?thesis*
proof (*cases $\forall n > 0. (deriv \wedge n) f \xi = 0$*)
case *True*
have *fcon: $f w = f \xi$ if $w \in ball \xi r$ for w*
by (*meson open_ball True $\langle 0 < r \rangle$ centre_in_ball connected_ball holf'*
holomorphic_fun_eq_const_on_connected that)
have *False*
using *fcon [of $\xi + r/2$] $\langle 0 < r \rangle$ r injf unfolding inj_on_def*
by (*metis $\langle \xi \in S \rangle$ contra_subsetD dist_commute fcon mem_ball perfect_choose_dist*)
then show *?thesis ..*
next
case *False*
then obtain *n0 where n0: $n0 > 0 \wedge (deriv \wedge n0) f \xi \neq 0$ by blast*
define *n where [abs_def]: $n = (LEAST n. n > 0 \wedge (deriv \wedge n) f \xi \neq 0)$*
have *n_ne: $n > 0$ (deriv $\wedge n$) f $\xi \neq 0$*
using *def_LeastI [OF n_def n0] by auto*
have *n_min: $\bigwedge k. 0 < k \implies k < n \implies (deriv \wedge k) f \xi = 0$*

```

    using def_Least_le [OF n_def] not_le by auto
  obtain g δ where 0 < δ
    and holg: g holomorphic_on ball ξ δ
    and fd:  $\bigwedge w. w \in \text{ball } \xi \ \delta \implies f w - f \xi = ((w - \xi) * g w) ^ n$ 
    and gnz:  $\bigwedge w. w \in \text{ball } \xi \ \delta \implies g w \neq 0$ 
  by (blast intro: n_min holomorphic_factor_order_of_zero_strong [OF holg
    ‹open S› ‹ξ ∈ S› n_ne])
  show ?thesis
  proof (cases n=1)
    case True
      with n_ne show ?thesis by auto
    next
      case False
        have g holomorphic_on ball ξ (min r δ)
          using holg by (simp add: holomorphic_on_subset subset_ball)
        then have holgw:  $(\lambda w. (w - \xi) * g w)$  holomorphic_on ball ξ (min r δ)
          by (intro holomorphic_intros)
        have gd:  $\bigwedge w. \text{dist } \xi \ w < \delta \implies (g \text{ has\_field\_derivative } \text{deriv } g \ w) \text{ (at } w)$ 
          using holg
          by (simp add: DERIV_deriv_iff_field_differentiable holomorphic_on_def
            at_within_open_NO_MATCH)
        have *:  $\bigwedge w. w \in \text{ball } \xi \ (\text{min } r \ \delta) \implies ((\lambda w. (w - \xi) * g w) \text{ has\_field\_derivative } ((w - \xi) * \text{deriv } g \ w + g \ w))$ 
          (at w)
          by (rule gd derivative_eq_intros | simp)+
        have [simp]:  $\text{deriv } (\lambda w. (w - \xi) * g w) \ \xi \neq 0$ 
          using * [of ξ] ‹0 < δ› ‹0 < r› by (simp add: DERIV_imp_deriv gnz)
        obtain T where ξ ∈ T open T and Tsb:  $T \subseteq \text{ball } \xi \ (\text{min } r \ \delta)$  and oimT:
          open  $((\lambda w. (w - \xi) * g w) ^ T)$ 
          using ‹0 < r› ‹0 < δ› has_complex_derivative_locally_invertible [OF
            holgw, of ξ]
          by force
        define U where  $U = (\lambda w. (w - \xi) * g w) ^ T$ 
        have open U by (metis oimT U_def)
        moreover have 0 ∈ U
          using ‹ξ ∈ T› by (auto simp: U_def intro: image_eqI [where x = ξ])
        ultimately obtain ε where ε > 0 and ε:  $\text{cball } 0 \ \varepsilon \subseteq U$ 
          using ‹open U› open_contains_cball by blast
        then have ε *  $\exp(2 * \text{of\_real } \pi * i * (0/n)) \in \text{cball } 0 \ \varepsilon$ 
          and ε *  $\exp(2 * \text{of\_real } \pi * i * (1/n)) \in \text{cball } 0 \ \varepsilon$ 
          by (auto simp: norm_mult)
        with ε have ε *  $\exp(2 * \text{of\_real } \pi * i * (0/n)) \in U$ 
          and ε *  $\exp(2 * \text{of\_real } \pi * i * (1/n)) \in U$  by blast+
        then obtain y0 y1 where y0 ∈ T and y0:  $(y0 - \xi) * g y0 = \varepsilon * \exp(2 * \text{of\_real } \pi * i * (0/n))$ 
          and y1 ∈ T and y1:  $(y1 - \xi) * g y1 = \varepsilon * \exp(2 * \text{of\_real } \pi * i * (1/n))$ 
          by (auto simp: U_def)

```

```

then have  $y0 \in \text{ball } \xi \ \delta \ y1 \in \text{ball } \xi \ \delta$  using Tsb by auto
then have  $f \ y0 - f \ \xi = ((y0 - \xi) * g \ y0) \wedge^n \ f \ y1 - f \ \xi = ((y1 - \xi) * g \ y1) \wedge^n$ 
using fd by blast+
moreover have  $y0 \neq y1$ 
using  $y0 \ y1 \ \langle \varepsilon > 0 \rangle$  complex_root_unity_eq_1 [of n 1]  $\langle n > 0 \rangle$  False by
auto
moreover have  $T \subseteq S$ 
by (meson Tsb min.cobounded1 order_trans r subset_ball)
ultimately have False
using inj_onD [OF injf, of y0 y1]  $\langle y0 \in T \rangle \ \langle y1 \in T \rangle$ 
using complex_root_unity [of n 1]
by (auto simp: y0 y1 power_mult_distrib diff_eq_eq n_ne)
then show ?thesis ..
qed
qed
qed

```

5.6.1 Hence a nice clean inverse function theorem

lemma *has_field_derivative_inverse_strong*:

```

fixes  $f :: 'a :: \{\text{euclidean\_space, real\_normed\_field}\} \Rightarrow 'a$ 
shows  $\llbracket \text{DERIV } f \ x \ :> f'; f' \neq 0; \text{open } S; x \in S; \text{continuous\_on } S \ f; \wedge z. z \in S \implies g \ (f \ z) = z \rrbracket$ 
 $\implies \text{DERIV } g \ (f \ x) \ :> \text{inverse } (f')$ 
unfolding has_field_derivative_def
by (rule has_derivative_inverse_strong [of S x f g]) auto

```

lemma *has_field_derivative_inverse_strong_x*:

```

fixes  $f :: 'a :: \{\text{euclidean\_space, real\_normed\_field}\} \Rightarrow 'a$ 
shows  $\llbracket \text{DERIV } f \ (g \ y) \ :> f'; f' \neq 0; \text{open } S; \text{continuous\_on } S \ f; g \ y \in S; f(g \ y) = y; \wedge z. z \in S \implies g \ (f \ z) = z \rrbracket$ 
 $\implies \text{DERIV } g \ y \ :> \text{inverse } (f')$ 
unfolding has_field_derivative_def
by (rule has_derivative_inverse_strong_x [of S g y f]) auto

```

proposition *holomorphic_has_inverse*:

```

assumes holf: f holomorphic_on S
and open S and injf: inj_on f S
obtains  $g$  where  $g$  holomorphic_on ( $f \ ' \ S$ )
 $\wedge z. z \in S \implies \text{deriv } f \ z * \text{deriv } g \ (f \ z) = 1$ 
 $\wedge z. z \in S \implies g(f \ z) = z$ 

```

proof –

```

have ofs: open (f \ ' S)
by (rule open_mapping_thm3 [OF assms])
have contf: continuous_on S f
by (simp add: holf holomorphic_on_imp_continuous_on)
have  $*$ : (the_inv_into S f has_field_derivative_inverse (deriv f z)) (at (f z)) if

```

```

z ∈ S for z
proof -
  have 1: (f has_field_derivative deriv f z) (at z)
    using DERIV_deriv_iff_field_differentiable ⟨z ∈ S⟩ ⟨open S⟩ holf holomor-
    phic_on_imp_differentiable_at
    by blast
  have 2: deriv f z ≠ 0
    using ⟨z ∈ S⟩ ⟨open S⟩ holf holomorphic_injective_imp_regular injf by blast
  show ?thesis
  proof (rule has_field_derivative_inverse_strong [OF 1 2 ⟨open S⟩ ⟨z ∈ S⟩])
    show continuous_on S f
      by (simp add: holf holomorphic_on_imp_continuous_on)
    show ∧z. z ∈ S ⇒ the_inv_into S f (f z) = z
      by (simp add: injf the_inv_into_f_f)
  qed
qed
show ?thesis
proof
  show the_inv_into S f holomorphic_on f ` S
    by (simp add: holomorphic_on_open ofs) (blast intro: *)
next
fix z assume z ∈ S
have deriv f z ≠ 0
  using ⟨z ∈ S⟩ ⟨open S⟩ holf holomorphic_injective_imp_regular injf by
blast
then show deriv f z * deriv (the_inv_into S f) (f z) = 1
  using * [OF ⟨z ∈ S⟩] by (simp add: DERIV_imp_deriv)
next
fix z assume z ∈ S
show the_inv_into S f (f z) = z
  by (simp add: ⟨z ∈ S⟩ injf the_inv_into_f_f)
qed
qed

```

5.6.2 Holomorphism of covering maps and lifts.

lemma covering_space_lift_is_holomorphic:

```

assumes cov: covering_space C p S
  and C: open C p holomorphic_on C
  and holf: f holomorphic_on U and fim: f ∈ U → S and gim: g ∈ U → C
  and contg: continuous_on U g and pg_f: ∧x. x ∈ U ⇒ p(g x) = f x
shows g holomorphic_on U
unfolding holomorphic_on_def
proof (intro strip)
  fix z
  assume z ∈ U
  with fim have f z ∈ S by blast
  then obtain T V where f z ∈ T and opeT: openin (top_of_set S) T
    and UV: ∪ V = C ∩ p ` T

```

```

    and  $\bigwedge W. W \in \mathcal{V} \implies \text{openin } (\text{top\_of\_set } C) W$ 
    and  $\text{disV}: \text{pairwise disjoint } \mathcal{V}$  and  $\text{homeV}: \bigwedge W. W \in \mathcal{V} \implies \exists q. \text{homeomor-}$ 
    phism  $W T p q$ 
    using cov fim unfolding covering_space_def by meson
    then have  $\bigwedge W. W \in \mathcal{V} \implies \text{open } W \wedge W \subseteq C$ 
    by (metis  $\langle \text{open } C \rangle \text{inf\_le1 open\_Int openin\_open}$ )
    then obtain  $V$  where  $V \in \mathcal{V}$   $g z \in V$   $\text{open } V$   $V \subseteq C$ 
    by (metis IntI UnionE image_subset_iff vimageI UV  $\langle f z \in T \rangle \langle z \in U \rangle$  gim
    pg_f image_subset_iff_funcset)
    have holp:  $p$  holomorphic_on  $V$ 
    using  $\langle V \subseteq C \rangle \langle p \text{ holomorphic\_on } C \rangle$  holomorphic_on_subset by blast
    moreover have injp: inj_on  $p$   $V$ 
    by (metis  $\langle V \in \mathcal{V} \rangle$  homeV homeomorphism_def inj_on_inverseI)
    ultimately
    obtain  $p'$  where holp':  $p'$  holomorphic_on  $(p^{-1} V)$  and pp':  $\bigwedge z. z \in V \implies p'(p$ 
     $z) = z$ 
    using  $\langle \text{open } V \rangle$  holomorphic_has_inverse by metis
    have  $z \in U \cap g^{-1} V$ 
    using  $\langle g z \in V \rangle \langle z \in U \rangle$  by blast
    moreover have openin  $(\text{top\_of\_set } U)$   $(U \cap g^{-1} V)$ 
    using continuous_openin_preimage [OF contg gim]
    by (meson  $\langle \text{open } V \rangle$  contg continuous_openin_preimage_eq)
    ultimately obtain  $\varepsilon$  where  $\varepsilon > 0$  and e:  $\text{ball } z \varepsilon \cap U \subseteq g^{-1} V$ 
    by (force simp: openin_contains_ball)
    show  $g$  field_differentiable at  $z$  within  $U$ 
    proof (rule field_differentiable_transform_within)
    show  $(0::\text{real}) < \varepsilon$ 
    by (simp add:  $\langle 0 < \varepsilon \rangle$ )
    show  $z \in U$ 
    by (simp add:  $\langle z \in U \rangle$ )
    show  $(p' \circ f) x' = g x'$  if  $x' \in U$  dist  $x' z < \varepsilon$  for  $x'$ 
    using that
    by (metis Int_iff comp_apply dist_commute e mem_ball pg_f pp' subsetD
    vimage_eq)
    have open  $(p^{-1} V)$ 
    using  $\langle \text{open } V \rangle$  holp injp open_mapping_thm3 by force
    moreover have  $f z \in p^{-1} V$ 
    by (metis  $\langle z \in U \rangle$  image_iff pg_f  $\langle g z \in V \rangle$ )
    ultimately have  $p'$  field_differentiable at  $(f z)$ 
    using holomorphic_on_imp_differentiable_at holp' by blast
    moreover have  $f$  field_differentiable at  $z$  within  $U$ 
    by (metis (no_types)  $\langle z \in U \rangle$  holf holomorphic_on_def)
    ultimately show  $(p' \circ f)$  field_differentiable at  $z$  within  $U$ 
    by (metis (no_types) field_differentiable_at_within field_differentiable_compose_within)
  qed
qed

```

lemma *covering_space_lift_holomorphic*:

assumes *cov*: *covering_space* $C p S$

```

    and C: open C p holomorphic_on C
    and f: f holomorphic_on U f ∈ U → S
    and U: simply_connected U locally_path_connected U
  obtains g where g holomorphic_on U g ∈ U → C ∧ y. y ∈ U ⇒ p(g y) = f y
proof -
  obtain g where continuous_on U g g ∈ U → C ∧ y. y ∈ U ⇒ p(g y) = f y
  using covering_space_lift [OF cov U] f holomorphic_on_imp_continuous_on
by blast
  then show ?thesis
  by (metis C cov covering_space_lift_is_holomorphic f image_subset_iff_funcset
that)
qed

```

5.7 The Schwarz Lemma

lemma Schwarz1:

```

  assumes holf: f holomorphic_on S
    and conf: continuous_on (closure S) f
    and S: open S connected S
    and boS: bounded S
    and S ≠ {}
  obtains w where w ∈ frontier S
    ∧ z. z ∈ closure S ⇒ norm (f z) ≤ norm (f w)
proof -
  have connf: continuous_on (closure S) (norm o f)
  using conf continuous_on_compose continuous_on_norm_id by blast
  have coc: compact (closure S)
  by (simp add: ⟨bounded S⟩ bounded_closure compact_eq_bounded_closed)
  then obtain x where x: x ∈ closure S and xmax: ∧ z. z ∈ closure S ⇒ norm(f
z) ≤ norm(f x)
  using continuous_attains_sup [OF _ _ connf] ⟨S ≠ {}⟩ by auto
  then show ?thesis
proof (cases x ∈ frontier S)
  case True
  then show ?thesis using that xmax by blast
next
  case False
  then have x ∈ S
  using ⟨open S⟩ frontier_def interior_eq x by auto
  then have f_constant_on S
proof (rule maximum_modulus_principle [OF holf S ⟨open S⟩ order_refl])
  show ∧ z. z ∈ S ⇒ cmod (f z) ≤ cmod (f x)
  using closure_subset by (blast intro: xmax)
qed
  then have f_constant_on (closure S)
  by (rule constant_on_closureI [OF _ conf])
  then obtain c where c: ∧ x. x ∈ closure S ⇒ f x = c
  by (meson constant_on_def)
  obtain w where w ∈ frontier S

```

```

    by (metis coc_all_not_in_conv assms(6) closure_UNIV frontier_eq_empty
not_compact_UNIV)
    then show ?thesis
      by (simp add: c_frontier_def that)
  qed
qed

```

lemma Schwarz2:

```

[[f holomorphic_on ball 0 r;
  0 < s; ball w s ⊆ ball 0 r;
  ∧z. norm (w-z) < s ⇒ norm(f z) ≤ norm(f w)]
⇒ f constant_on ball 0 r
by (rule maximum_modulus_principle [where U = ball w s and ξ = w]) (simp_all
add: dist_norm)

```

lemma Schwarz3:

```

assumes holf: f holomorphic_on (ball 0 r) and [simp]: f 0 = 0
obtains h where h holomorphic_on (ball 0 r) and ∧z. norm z < r ⇒ f z =
z * (h z) and deriv f 0 = h 0
proof -
  define h where h z = (if z = 0 then deriv f 0 else f z / z) for z
  have d0: deriv f 0 = h 0
    by (simp add: h_def)
  moreover have h holomorphic_on (ball 0 r)
    by (rule pole_theorem_open_0 [OF holf, of 0]) (auto simp: h_def)
  moreover have norm z < r ⇒ f z = z * h z for z
    by (simp add: h_def)
  ultimately show ?thesis
    using that by blast
qed

```

proposition Schwarz_Lemma:

```

assumes holf: f holomorphic_on (ball 0 1) and [simp]: f 0 = 0
and no: ∧z. norm z < 1 ⇒ norm (f z) < 1
and ξ: norm ξ < 1
shows norm (f ξ) ≤ norm ξ and norm(deriv f 0) ≤ 1
and ((∃z. norm z < 1 ∧ z ≠ 0 ∧ norm(f z) = norm z)
  ∨ norm(deriv f 0) = 1)
  ⇒ ∃α. (∀z. norm z < 1 → f z = α * z) ∧ norm α = 1
(is ?P ⇒ ?Q)
proof -
  obtain h where holh: h holomorphic_on (ball 0 1)
    and fz_eq: ∧z. norm z < 1 ⇒ f z = z * (h z) and df0: deriv f 0 = h
0
  by (rule Schwarz3 [OF holf]) auto
  have noh_le: norm (h z) ≤ 1 if z: norm z < 1 for z
  proof -
    have norm (h z) < a if a: 1 < a for a
    proof -

```

```

have max (inverse a) (norm z) < 1
  using z a by (simp_all add: inverse_less_1_iff)
then obtain r where r: max (inverse a) (norm z) < r and r < 1
  using Rats_dense_in_real by blast
then have nzs: norm z < r and ira: inverse r < a
  using z a less_imp_inverse_less by force+
then have 0 < r
  by (meson norm_not_less_zero not_le order.strict_trans2)
have holh': h holomorphic_on ball 0 r
  by (meson holh ⟨r < 1⟩ holomorphic_on_subset less_eq_real_def subset_ball)
have conth': continuous_on (cball 0 r) h
  by (meson ⟨r < 1⟩ dual_order.trans holh holomorphic_on_imp_continuous_on holomorphic_on_subset mem_ball_0 mem_cball_0 not_less subsetI)
obtain w where w: norm w = r and lenw:  $\bigwedge z. \text{norm } z < r \implies \text{norm}(h z) \leq \text{norm}(h w)$ 
  using conth' ⟨0 < r⟩ by (auto simp add: intro: Schwarz1 [OF holh'])
have h w = f w / w using fz_eq ⟨r < 1⟩ nzs w by auto
then have cmod (h z) < inverse r
  by (metis ⟨0 < r⟩ ⟨r < 1⟩ divide_strict_right_mono inverse_eq_divide le_less_trans lenw no_norm_divide nzs w)
then show ?thesis using ira by linarith
qed
then show norm (h z) ≤ 1
  using not_le by blast
qed
show cmod (f ξ) ≤ cmod ξ
proof (cases ξ = 0)
  case True then show ?thesis by auto
next
  case False
  then show ?thesis
    by (simp add: noh_le fz_eq ξ mult_left_le norm_mult)
qed
show no_df0: norm(deriv f 0) ≤ 1
  by (simp add: ⟨ $\bigwedge z. \text{cmod } z < 1 \implies \text{cmod}(h z) \leq 1$ ⟩ df0)
show ?Q if ?P
  using that
proof
  assume  $\exists z. \text{cmod } z < 1 \wedge z \neq 0 \wedge \text{cmod}(f z) = \text{cmod } z$ 
  then obtain γ where γ: cmod γ < 1 γ ≠ 0 cmod (f γ) = cmod γ by blast
  then have [simp]: norm (h γ) = 1
    by (simp add: fz_eq norm_mult)
  have §: ball γ (1 - cmod γ) ⊆ ball 0 1
    by (simp add: ball_subset_ball_iff)
  moreover have  $\bigwedge z. \text{cmod}(\gamma - z) < 1 - \text{cmod } \gamma \implies \text{cmod}(h z) \leq \text{cmod}(h \gamma)$ 
    by (metis ⟨cmod (h γ) = 1⟩ § dist_0_norm dist_complex_def in_mono mem_ball noh_le)

```



```

ultimately obtain c where c:  $\bigwedge z. \text{norm } z < 1 \implies h z = c$ 
  using Schwarz2 [OF holh, of 1 - norm  $\gamma$ , unfolded constant_on_def]  $\gamma$ 
by auto
  then have norm c = 1
    using  $\gamma$  by force
  with c show ?thesis
    using fz_eq by auto
next
  assume [simp]: cmod (deriv f 0) = 1
  then obtain c where c:  $\bigwedge z. \text{norm } z < 1 \implies h z = c$ 
    using Schwarz2 [OF holh zero_less_one, of 0, unfolded constant_on_def]
df0 noh_le
  by auto
  moreover have norm c = 1 using df0 c by auto
  ultimately show ?thesis
    using fz_eq by auto
qed
qed

```

corollary Schwarz_Lemma':

```

assumes holh: f holomorphic_on (ball 0 1) and [simp]: f 0 = 0
  and no:  $\bigwedge z. \text{norm } z < 1 \implies \text{norm } (f z) < 1$ 
  shows (( $\forall \xi. \text{norm } \xi < 1 \longrightarrow \text{norm } (f \xi) \leq \text{norm } \xi$ )
     $\wedge \text{norm } (\text{deriv } f 0) \leq 1$ )
     $\wedge ((\exists z. \text{norm } z < 1 \wedge z \neq 0 \wedge \text{norm } (f z) = \text{norm } z)$ 
       $\vee \text{norm } (\text{deriv } f 0) = 1)$ 
     $\longrightarrow (\exists \alpha. (\forall z. \text{norm } z < 1 \longrightarrow f z = \alpha * z) \wedge \text{norm } \alpha = 1))$ 
  using Schwarz_Lemma [OF assms]
  by (metis (no_types) norm_eq_zero zero_less_one)

```

5.8 The Schwarz reflection principle

lemma hol_pal_lem0:

assumes $d \cdot a \leq k \leq d \cdot b$

obtains c where

```

  c  $\in$  closed_segment a b  $\wedge d \cdot c = k$ 
   $\bigwedge z. z \in \text{closed\_segment } a c \implies d \cdot z \leq k$ 
   $\bigwedge z. z \in \text{closed\_segment } c b \implies k \leq d \cdot z$ 

```

proof -

obtain c where cin: $c \in \text{closed_segment } a b$ and keq: $k = d \cdot c$

using connected_ivt_hyperplane [of closed_segment a b a b d k]

by (auto simp: assms)

have closed_segment a c $\subseteq \{z. d \cdot z \leq k\}$ closed_segment c b $\subseteq \{z. k \leq d \cdot z\}$

unfolding segment_convex_hull using assms keq

by (auto simp: convex_halfspace_le convex_halfspace_ge hull_minimal)

then show ?thesis using cin that by fastforce

qed

lemma hol_pal_lem1:

```

assumes convex  $S$  open  $S$ 
and abc:  $a \in S$   $b \in S$   $c \in S$ 
            $d \neq 0$  and lek:  $d \cdot a \leq k$   $d \cdot b \leq k$   $d \cdot c \leq k$ 
and holf1:  $f$  holomorphic_on  $\{z. z \in S \wedge d \cdot z < k\}$ 
and contf: continuous_on  $S$   $f$ 
shows contour_integral (linepath  $a$   $b$ )  $f$  +
        contour_integral (linepath  $b$   $c$ )  $f$  +
        contour_integral (linepath  $c$   $a$ )  $f$  = 0
proof -
have interior (convex hull  $\{a, b, c\}$ )  $\subseteq$  interior( $S \cap \{x. d \cdot x \leq k\}$ )
proof (intro interior_mono hull_minimal)
show  $\{a, b, c\} \subseteq S \cap \{x. d \cdot x \leq k\}$ 
by (simp add: abc lek)
show convex ( $S \cap \{x. d \cdot x \leq k\}$ )
by (rule convex_Int [OF  $\langle$ convex  $S$  $\rangle$  convex_halfspace_le])
qed
also have  $\dots \subseteq \{z \in S. d \cdot z < k\}$ 
by (force simp: interior_open [OF  $\langle$ open  $S$  $\rangle$ ]  $\langle d \neq 0 \rangle$ )
finally have *: interior (convex hull  $\{a, b, c\}$ )  $\subseteq \{z \in S. d \cdot z < k\}$  .
have continuous_on (convex hull  $\{a, b, c\}$ )  $f$ 
using  $\langle$ convex  $S$  $\rangle$  contf abc continuous_on_subset subset_hull
by fastforce
moreover have  $f$  holomorphic_on interior (convex hull  $\{a, b, c\}$ )
by (rule holomorphic_on_subset [OF holf1 *])
ultimately show ?thesis
using Cauchy_theorem_triangle_interior_has_chain_integral_chain_integral3
by blast
qed

lemma hol_pal_lem2:
assumes  $S$ : convex  $S$  open  $S$ 
and abc:  $a \in S$   $b \in S$   $c \in S$ 
and  $d \neq 0$  and lek:  $d \cdot a \leq k$   $d \cdot b \leq k$ 
and holf1:  $f$  holomorphic_on  $\{z. z \in S \wedge d \cdot z < k\}$ 
and holf2:  $f$  holomorphic_on  $\{z. z \in S \wedge k < d \cdot z\}$ 
and contf: continuous_on  $S$   $f$ 
shows contour_integral (linepath  $a$   $b$ )  $f$  +
        contour_integral (linepath  $b$   $c$ )  $f$  +
        contour_integral (linepath  $c$   $a$ )  $f$  = 0
proof (cases  $d \cdot c \leq k$ )
case True show ?thesis
by (rule hol_pal_lem1 [OF  $S$  abc  $\langle d \neq 0 \rangle$  lek True holf1 contf])
next
case False
then have  $d \cdot c > k$  by force
obtain  $a'$  where  $a'$ :  $a' \in$  closed_segment  $b$   $c$  and  $d \cdot a' = k$ 
and  $ba'$ :  $\bigwedge z. z \in$  closed_segment  $b$   $a' \implies d \cdot z \leq k$ 
and  $a'c$ :  $\bigwedge z. z \in$  closed_segment  $a'$   $c \implies k \leq d \cdot z$ 
using False hol_pal_lem0 [of  $d$   $b$   $k$   $c$ , OF  $\langle d \cdot b \leq k \rangle$ ] by auto

```

```

obtain  $b'$  where  $b': b' \in \text{closed\_segment } a \ c$  and  $d \cdot b' = k$ 
  and  $ab': \bigwedge z. z \in \text{closed\_segment } a \ b' \implies d \cdot z \leq k$ 
  and  $b'c: \bigwedge z. z \in \text{closed\_segment } b' \ c \implies k \leq d \cdot z$ 
  using  $\text{False hol\_pal\_lem0 [of } d \ a \ k \ c, OF \langle d \cdot a \leq k \rangle]$  by auto
have  $a'b': a' \in S \wedge b' \in S$ 
  using  $a' \ abc \ b' \ \text{convex\_contains\_segment } \langle \text{convex } S \rangle$  by auto
have  $\text{continuous\_on } (\text{closed\_segment } c \ a) \ f$ 
  by (meson  $abc \ \text{contf } \text{continuous\_on\_subset } \text{convex\_contains\_segment } \langle \text{convex } S \rangle$ )
then have 1:  $\text{contour\_integral } (\text{linepath } c \ a) \ f =$ 
   $\text{contour\_integral } (\text{linepath } c \ b') \ f + \text{contour\_integral } (\text{linepath } b' \ a) \ f$ 
  using  $b' \ \text{closed\_segment\_commute } \text{contour\_integral\_split\_linepath}$  by blast
have  $\text{continuous\_on } (\text{closed\_segment } b \ c) \ f$ 
  by (meson  $abc \ \text{contf } \text{continuous\_on\_subset } \text{convex\_contains\_segment } \langle \text{convex } S \rangle$ )
then have 2:  $\text{contour\_integral } (\text{linepath } b \ c) \ f =$ 
   $\text{contour\_integral } (\text{linepath } b \ a') \ f + \text{contour\_integral } (\text{linepath } a' \ c) \ f$ 
  by (rule  $\text{contour\_integral\_split\_linepath } [OF \_ a']$ )
have 3:  $\text{contour\_integral } (\text{reversepath } (\text{linepath } b' \ a')) \ f =$ 
   $-\ \text{contour\_integral } (\text{linepath } b' \ a') \ f$ 
  by (rule  $\text{contour\_integral\_reversepath } [OF \ \text{valid\_path\_linepath}]$ )
have  $\text{fcd\_le: } f \ \text{field\_differentiable at } x$ 
  if  $x \in \text{interior } S \wedge x \in \text{interior } \{x. d \cdot x \leq k\}$  for  $x$ 
proof -
  have  $f \ \text{holomorphic\_on } S \cap \{c. d \cdot c < k\}$ 
  by (metis (no_types)  $\text{Collect\_conj\_eq } \text{Collect\_mem\_eq } \text{hol}f1$ )
  then have  $\exists C \ D. x \in \text{interior } C \cap \text{interior } D \wedge f \ \text{holomorphic\_on } \text{interior } C$ 
 $\cap \text{interior } D$ 
  using that
  by (metis  $\text{Collect\_mem\_eq } \text{Int\_Collect } \langle d \neq 0 \rangle \ \text{interior\_halfspace\_le } \text{interior\_open } \langle \text{open } S \rangle$ )
  then show  $f \ \text{field\_differentiable at } x$ 
  by (metis  $\text{at\_within\_interior } \text{holomorphic\_on\_def } \text{interior\_Int } \text{interior\_interior}$ )
qed
have  $ab\_le: \bigwedge x. x \in \text{closed\_segment } a \ b \implies d \cdot x \leq k$ 
proof -
  fix  $x :: \text{complex}$ 
  assume  $x \in \text{closed\_segment } a \ b$ 
  then have  $\bigwedge C. x \in C \vee b \notin C \vee a \notin C \vee \neg \text{convex } C$ 
  by (meson  $\text{contra\_subsetD } \text{convex\_contains\_segment}$ )
  then show  $d \cdot x \leq k$ 
  by (metis  $\text{lek } \text{convex\_halfspace\_le } \text{mem\_Collect\_eq}$ )
qed
have  $cs: \text{closed\_segment } a' \ b' \subseteq \{x. d \cdot x \leq k\} \wedge \text{closed\_segment } b' \ a \subseteq \{x. d \cdot x \leq k\}$ 
  by (simp  $\text{add: } \langle d \cdot a' = k \rangle \langle d \cdot b' = k \rangle \ \text{closed\_segment\_subset } \text{convex\_halfspace\_le } \text{lek}(1)$ )
have  $\text{continuous\_on } (S \cap \{x. d \cdot x \leq k\}) \ f$  using contf
  by (simp  $\text{add: } \text{continuous\_on\_subset}$ )

```

```

then have (f has_contour_integral 0)
  (linepath a b +++ linepath b a' +++ linepath a' b' +++ linepath b' a)
apply (rule Cauchy_theorem_convex [where K = {}])
by (simp_all add: path_image_join convex_Int convex_halfspace_le ‹convex
S› fcd_le ab_le
      closed_segment_subset abc a'b' ba' cs)
then have 4: contour_integral (linepath a b) f +
  contour_integral (linepath b a') f +
  contour_integral (linepath a' b') f +
  contour_integral (linepath b' a) f = 0
by (rule has_chain_integral_chain_integral4)
have fcd_ge: f field_differentiable at x
  if x ∈ interior S ∧ x ∈ interior {x. k ≤ d · x} for x
proof -
have f2: f holomorphic_on S ∩ {c. k < d · c}
  by (metis (full_types) Collect_conj_eq Collect_mem_eq holf2)
have f3: interior S = S
  by (simp add: interior_open ‹open S›)
then have x ∈ S ∩ interior {c. k ≤ d · c}
  using that by simp
then show f field_differentiable at x
  using f3 f2 unfolding holomorphic_on_def
  by (metis (no_types) ‹d ≠ 0› at_within_interior interior_Int interior_halfspace_ge
interior_interior)
qed
have cs: closed_segment c b' ⊆ {x. k ≤ d · x} ∧ closed_segment b' a' ⊆ {x. k
≤ d · x}
  by (simp add: ‹d · a' = k› b'c closed_segment_subset convex_halfspace_ge)
have continuous_on (S ∩ {x. k ≤ d · x}) f using contf
  by (simp add: continuous_on_subset)
then have (f has_contour_integral 0) (linepath a' c +++ linepath c b' +++
linepath b' a')
  apply (rule Cauchy_theorem_convex [where K = {}])
  by (simp_all add: path_image_join convex_Int convex_halfspace_ge ‹convex
S›
      fcd_ge closed_segment_subset abc a'b' a'c cs)
then have 5: contour_integral (linepath a' c) f + contour_integral (linepath c
b') f + contour_integral (linepath b' a') f = 0
  by (rule has_chain_integral_chain_integral3)
show ?thesis
  using 1 2 3 4 5 by (metis add.assoc eq_neg_iff_add_eq_0 reversepath_linepath)
qed

```

lemma hol_pal_lem3:

```

assumes S: convex S open S
  and abc: a ∈ S b ∈ S c ∈ S
  and d ≠ 0 and lek: d · a ≤ k
  and holf1: f holomorphic_on {z. z ∈ S ∧ d · z < k}
  and holf2: f holomorphic_on {z. z ∈ S ∧ k < d · z}

```

```

    and conf: continuous_on S f
  shows contour_integral (linepath a b) f +
        contour_integral (linepath b c) f +
        contour_integral (linepath c a) f = 0
proof (cases  $d \cdot b \leq k$ )
  case True show ?thesis
    by (rule hol_pal_lem2 [OF S abc  $d \neq 0$ ] lek True holf1 holf2 conf)
next
  case False
  show ?thesis
  proof (cases  $d \cdot c \leq k$ )
    case True
    have contour_integral (linepath c a) f +
          contour_integral (linepath a b) f +
          contour_integral (linepath b c) f = 0
    by (rule hol_pal_lem2 [OF S  $\langle c \in S \rangle \langle a \in S \rangle \langle b \in S \rangle \langle d \neq 0 \rangle \langle d \cdot c \leq k \rangle$ ]
      lek holf1 holf2 conf)
    then show ?thesis
      by (simp add: algebra_simps)
  next
    case False
    have contour_integral (linepath b c) f +
          contour_integral (linepath c a) f +
          contour_integral (linepath a b) f = 0
    using hol_pal_lem2 [OF S  $\langle b \in S \rangle \langle c \in S \rangle \langle a \in S \rangle$ , of  $-d -k$ ]
    using  $\langle d \neq 0 \rangle \langle \neg d \cdot b \leq k \rangle$  False by (simp_all add: holf1 holf2 conf)
    then show ?thesis
      by (simp add: algebra_simps)
  qed
qed

```

lemma hol_pal_lem4:

```

  assumes S: convex S open S
    and abc:  $a \in S \ b \in S \ c \in S$  and  $d \neq 0$ 
    and holf1:  $f$  holomorphic_on  $\{z. z \in S \wedge d \cdot z < k\}$ 
    and holf2:  $f$  holomorphic_on  $\{z. z \in S \wedge k < d \cdot z\}$ 
    and conf: continuous_on S f
  shows contour_integral (linepath a b) f +
        contour_integral (linepath b c) f +
        contour_integral (linepath c a) f = 0
proof (cases  $d \cdot a \leq k$ )
  case True show ?thesis
    by (rule hol_pal_lem3 [OF S abc  $d \neq 0$ ] True holf1 holf2 conf)
next
  case False
  show ?thesis
    using  $\langle d \neq 0 \rangle$  hol_pal_lem3 [OF S abc, of  $-d -k$ ] False
    by (simp_all add: holf1 holf2 conf)
qed

```

lemma *holomorphic_on_paste_across_line*:

assumes S : open S **and** $d \neq 0$
and $holf1$: f holomorphic_on $(S \cap \{z. d \cdot z < k\})$
and $holf2$: f holomorphic_on $(S \cap \{z. k < d \cdot z\})$
and $contf$: continuous_on S f
shows f holomorphic_on S

proof –

have *: $\exists t. \text{open } t \wedge p \in t \wedge \text{continuous_on } t \ f \wedge$
 $(\forall a \ b \ c. \text{convex_hull } \{a, b, c\} \subseteq t \longrightarrow$
 $\text{contour_integral } (\text{linepath } a \ b) \ f +$
 $\text{contour_integral } (\text{linepath } b \ c) \ f +$
 $\text{contour_integral } (\text{linepath } c \ a) \ f = 0)$
if $p \in S$ **for** p

proof –

obtain e **where** $e > 0$ **and** e : ball p $e \subseteq S$
using $\langle p \in S \rangle$ openE S **by** blast
then have continuous_on (ball p e) f
using $contf$ continuous_on_subset **by** blast
moreover
have $\{z. \text{dist } p \ z < e \wedge d \cdot z < k\} \subseteq S \cap \{z. d \cdot z < k\}$
 $\{z. \text{dist } p \ z < e \wedge k < d \cdot z\} \subseteq S \cap \{z. k < d \cdot z\}$
using e **by** auto
then have f holomorphic_on $\{z. \text{dist } p \ z < e \wedge d \cdot z < k\}$
 f holomorphic_on $\{z. \text{dist } p \ z < e \wedge k < d \cdot z\}$
using holomorphic_on_subset $holf1$ $holf2$ **by** presburger+
ultimately show ?thesis
apply (rule_tac $x = \text{ball } p \ e$ **in** exI)
using $\langle e > 0 \rangle$ $e \langle d \neq 0 \rangle$ hol_pal_lem4 [of ball p e ___ d k]
by (force simp: subset_hull)

qed

show ?thesis

by (blast intro: * Morera_local_triangle_analytic_imp_holomorphic)

qed

proposition *Schwarz_reflection*:

assumes open S **and** $cnjs$: $cnj \ ' S \subseteq S$
and $holf$: f holomorphic_on $(S \cap \{z. 0 < \text{Im } z\})$
and $contf$: continuous_on $(S \cap \{z. 0 \leq \text{Im } z\})$ f
and f : $\bigwedge z. \llbracket z \in S; z \in \mathbb{R} \rrbracket \implies (f \ z) \in \mathbb{R}$
shows $(\lambda z. \text{if } 0 \leq \text{Im } z \text{ then } f \ z \text{ else } cnj(f(cnj \ z)))$ holomorphic_on S

proof –

have 1: $(\lambda z. \text{if } 0 \leq \text{Im } z \text{ then } f \ z \text{ else } cnj(f(cnj \ z)))$ holomorphic_on $(S \cap \{z. 0 < \text{Im } z\})$

by (force intro: iffD1 [OF holomorphic_cong [OF refl] $holf$])

have cont_cfc: continuous_on $(S \cap \{z. \text{Im } z \leq 0\})$ $(cnj \ o \ f \ o \ cnj)$

using $cnjs$

by (intro continuous_intros continuous_on_compose continuous_on_subset [OF $contf$]) auto

```

have cnj ∘ f ∘ cnj field_differentiable at x within  $S \cap \{z. \text{Im } z < 0\}$ 
  if  $x \in S$  Im  $x < 0$  f field_differentiable at (cnj x) within  $S \cap \{z. 0 < \text{Im } z\}$ 
for x
  using that
  apply (clarsimp simp add: field_differentiable_def has_field_derivative_iff
Lim_within dist_norm)
  apply (rule_tac x=cnj f' in exI)
  apply (elim_all_forward ex_forward conj_forward imp_forward asm_rl, clarify)
  apply (drule_tac x=cnj xa in bspec)
  using cnjs apply force
  apply (metis complex_cnj_cnj complex_cnj_diff complex_cnj_divide complex_mod_cnj)
  done
  then have hol_cfc: (cnj ∘ f ∘ cnj) holomorphic_on ( $S \cap \{z. \text{Im } z < 0\}$ )
    using holf cnjs
    by (force simp: holomorphic_on_def)
  have 2: ( $\lambda z. \text{if } 0 \leq \text{Im } z \text{ then } f z \text{ else } \text{cnj } (f (\text{cnj } z))$ ) holomorphic_on ( $S \cap \{z. \text{Im } z < 0\}$ )
    by (smt (verit) Int_Collect comp_def hol_cfc holomorphic_cong)
  have [simp]: ( $S \cap \{z. 0 \leq \text{Im } z\} \cup (S \cap \{z. \text{Im } z \leq 0\}) = S$ )
    by force
  have eq:  $\bigwedge z. \llbracket z \in S; \text{Im } z \leq 0; 0 \leq \text{Im } z \rrbracket \implies f z = \text{cnj } (f (\text{cnj } z))$ 
    using f Reals_cnj_iff complex_is_Real_iff by auto
  have continuous_on ( $(S \cap \{z. 0 \leq \text{Im } z\}) \cup (S \cap \{z. \text{Im } z \leq 0\})$ )
    ( $\lambda z. \text{if } 0 \leq \text{Im } z \text{ then } f z \text{ else } \text{cnj } (f (\text{cnj } z))$ )
    apply (rule continuous_on_cases_local)
    using cont_cfc contf
  by (simp_all add: closedin_closed_Int closed_halfspace_Im_le closed_halfspace_Im_ge
eq)
  then have 3: continuous_on  $S$  ( $\lambda z. \text{if } 0 \leq \text{Im } z \text{ then } f z \text{ else } \text{cnj } (f (\text{cnj } z))$ )
    by force
  show ?thesis
    using holomorphic_on_paste_across_line [OF ‹open S›, of - i _ 0]
    using 1 2 3 by auto
qed

```

5.9 Bloch's theorem

lemma Bloch_lemma_0:

```

assumes holf: f holomorphic_on cball  $0 r$  and  $0 < r$ 
  and [simp]:  $f 0 = 0$ 
  and le:  $\bigwedge z. \text{norm } z < r \implies \text{norm}(deriv f z) \leq 2 * \text{norm}(deriv f 0)$ 
  shows ball  $0 ((3 - 2 * \text{sqrt } 2) * r * \text{norm}(deriv f 0)) \subseteq f \text{ ` ball } 0 r$ 

```

proof –

```

have sqrt 2 < 3/2
  by (rule real_less_sqrt) (auto simp: power2_eq_square)
then have sq3:  $0 < 3 - 2 * \text{sqrt } 2$  by simp
show ?thesis

```

```

proof (cases deriv f 0 = 0)
  case True then show ?thesis by simp
next
  case False
  define C where C = 2 * norm(deriv f 0)
  have 0 < C using False by (simp add: C_def)
  have holf': f holomorphic_on ball 0 r using holf
    using ball_subset_cball holomorphic_on_subset by blast
  then have holdf': deriv f holomorphic_on ball 0 r
    by (rule holomorphic_deriv [OF _ open_ball])
  have Le1: norm(deriv f z - deriv f 0) ≤ norm z / (r - norm z) * C
    if norm z < r for z
proof -
  have T1: norm(deriv f z - deriv f 0) ≤ norm z / (R - norm z) * C
    if R: norm z < R R < r for R
proof -
  have 0 < R using R
    by (metis less_trans norm_zero zero_less_norm_iff)
  have df_le:  $\bigwedge x. \text{norm } x < r \implies \text{norm } (\text{deriv } f \ x) \leq C$ 
    using le by (simp add: C_def)
  have hol_df: deriv f holomorphic_on cball 0 R
    using R holdf' holomorphic_on_subset by auto
  have *: (( $\lambda w. \text{deriv } f \ w / (w - z)$ ) has_contour_integral 2 * pi * i * deriv
f z) (circlepath 0 R)
    if norm z < R for z
    using <0 < R> that Cauchy_integral_formula_convex_simple [OF con-
vex_cball hol_df, of _ circlepath 0 R]
    by (force simp: winding_number_circlepath)
  have **: (( $\lambda x. \text{deriv } f \ x / (x - z) - \text{deriv } f \ x / x$ ) has_contour_integral
of_real (2 * pi) * i * (deriv f z - deriv f 0))
(circlepath 0 R)
    using has_contour_integral_diff [OF * [of z] * [of 0]] <0 < R> that
by (simp add: algebra_simps)
  have [simp]:  $\bigwedge x. \text{norm } x = R \implies x \neq z$  using that(1) by blast
  have norm (deriv f x / (x - z) - deriv f x / x)
≤ C * norm z / (R * (R - norm z))
if norm x = R for x
proof -
  have [simp]: norm (deriv f x * x - deriv f x * (x - z)) =
norm (deriv f x) * norm z
by (simp add: norm_mult right_diff_distrib')
  show ?thesis
using <0 < R> <0 < C> R that
by (auto simp add: norm_mult norm_divide divide_simps df_le
mult_mono norm_triangle_ineq2)
qed
then show ?thesis
using has_contour_integral_bound_circlepath
[OF **, of C * norm z / (R * (R - norm z))]

```



```

      ‹0 < R› ‹0 < C› R
    apply (simp add: norm_mult norm_divide)
    apply (simp add: divide_simps mult.commute)
    done
  qed
  obtain r' where r': norm z < r' r' < r
    using Rats_dense_in_real [of norm z r] ‹norm z < r› by blast
  then have [simp]: closure {r'<..2 / (r - norm z)) * norm (deriv f 0) ≤ norm (f z)
    if r: norm z < r for z
  proof -
    have 1: ∧x. x ∈ ball 0 r ⇒
      ((λz. f z - deriv f 0 * z) has_field_derivative deriv f x - deriv f 0)
      (at x within ball 0 r)
      by (rule derivative_eq_intros holomorphic_derivI holf' | simp)+
    have 2: closed_segment 0 z ⊆ ball 0 r
      by (metis ‹0 < r› convex_ball convex_contains_segment dist_self mem_ball
        mem_ball_0 that)
    have 4: norm (deriv f (x *R z) - deriv f 0) * norm z ≤ norm z * norm z *
      x * C / (r - norm z)
      if x: 0 ≤ x x ≤ 1 for x
    proof -
      have [simp]: x * norm z < r
      using r x by (meson le_less_trans mult_le_cancel_right2 norm_not_less_zero)
      then have cmod (x *R z) < r
        by (simp add: x)
      then have norm (deriv f (x *R z) - deriv f 0) ≤ norm (x *R z) / (r -
        norm (x *R z)) * C
        by (metis Le1)
      also have ... ≤ norm (x *R z) / (r - norm z) * C
        using r x ‹0 < r› ‹0 < C› by (simp add: frac_le mult_left_le_one_le)
      finally have norm (deriv f (x *R z) - deriv f 0) * norm z ≤ norm (x *R
        z) / (r - norm z) * C * norm z
        by (rule mult_right_mono) simp
      with x show ?thesis by (simp add: algebra_simps)
    qed
  qed
  have le_norm: abc ≤ norm d - e ⇒ norm (f - d) ≤ e ⇒ abc ≤ norm f
  for abc d e and f::complex
  by (metis add_diff_cancel_left' add_diff_eq diff_left_mono norm_diff_ineq
    order_trans)
  have norm (integral {0..1} (λx. (deriv f (x *R z) - deriv f 0) * z))
    ≤ integral {0..1} (λt. (norm z)2 * t / (r - norm z) * C)

```

```

proof (rule integral_norm_bound_integral)
  show ( $\lambda x. (\text{deriv } f (x *_{\mathbb{R}} z) - \text{deriv } f 0) * z$ ) integrable_on {0..1}
    using contour_integral_primitive [OF 1, of linepath 0 z] 2
  by (simp add: has_contour_integral_linepath has_integral_integrable_integral)
  have (*) (( $\text{cmod } z$ )2) integrable_on {0..1}
  by (metis ident_integrable_on integrable_0 integrable_eq integrable_on_cmult_iff
lambda_zero)
  then show ( $\lambda t. (\text{norm } z)^2 * t / (r - \text{norm } z) * C$ ) integrable_on {0..1}
    using integrable_on_cmult_right[where 'b=real, simplified] integrable_on_divide
[where 'b=real, simplified]
    by blast
  qed (simp add: norm_mult_power2_eq_square 4)
  then have int_le:  $\text{norm } (f z - \text{deriv } f 0 * z) \leq (\text{norm } z)^2 * \text{norm}(\text{deriv } f 0)$ 
/ (( $r - \text{norm } z$ ))
    using contour_integral_primitive [OF 1, of linepath 0 z] 2
  by (simp add: has_contour_integral_linepath has_integral_integrable_integral
C_def)
  have  $\text{norm } z * (\text{norm } (\text{deriv } f 0) * (r - \text{norm } z - \text{norm } z)) \leq \text{norm } z * (\text{norm } (\text{deriv } f 0) * (r - \text{norm } z) - \text{norm } (\text{deriv } f 0) * \text{norm } z)$ 
    by (simp add: algebra_simps)
  then have §:  $(\text{norm } z * (r - \text{norm } z) - \text{norm } z * \text{norm } z) * \text{norm } (\text{deriv } f 0) \leq \text{norm } (\text{deriv } f 0) * \text{norm } z * (r - \text{norm } z) - \text{norm } z * \text{norm } z * \text{norm } (\text{deriv } f 0)$ 
    by (simp add: algebra_simps)
  show ?thesis
    using ‹ $\text{norm } z < r$ ›
    by (force simp add: power2_eq_square divide_simps C_def norm_mult §
intro!: le_norm [OF _ int_le])
  qed
  have sq201 [simp]:  $0 < (1 - \sqrt{2} / 2) (1 - \sqrt{2} / 2) < 1$ 
    by (auto simp: sqrt2_less_2)
  have 1: continuous_on (closure (ball 0 (( $1 - \sqrt{2} / 2$ ) * r))) f
  proof (rule continuous_on_subset [OF holomorphic_on_imp_continuous_on
[OF holf]])
    show closure (ball 0 (( $1 - \sqrt{2} / 2$ ) * r))  $\subseteq$  cball 0 r
    proof -
      have  $(1 - \sqrt{2} / 2) * r \leq r$ 
        by (simp add: ‹ $0 < r$ ›)
      then show ?thesis
        by (meson ball_subset_cball closed_cball closure_minimal dual_order.trans
subset_ball)
    qed
  qed
  have 2: open (f ` interior (ball 0 (( $1 - \sqrt{2} / 2$ ) * r)))
  proof (rule open_mapping_thm [OF holf' open_ball_connected_ball])
    show interior (ball 0 (( $1 - \sqrt{2} / 2$ ) * r))  $\subseteq$  ball (0::complex) r
      using ‹ $0 < r$ › mult_pos_pos sq201 by (simp add: ball_subset_ball_iff)
    show  $\neg$  f constant_on ball 0 r
      using False ‹ $0 < r$ › centre_in_ball holf' holomorphic_nonconstant by blast

```

```

qed auto
have ball 0 ((3 - 2 * sqrt 2) * r * norm (deriv f 0)) =
  ball (f 0) ((3 - 2 * sqrt 2) * r * norm (deriv f 0))
  by simp
also have ...  $\subseteq$  f ' ball 0 ((1 - sqrt 2 / 2) * r)
proof -
  have  $\exists$ : (3 - 2 * sqrt 2) * r * norm (deriv f 0)  $\leq$  norm (f z)
    if norm z = (1 - sqrt 2 / 2) * r for z
  proof (rule order_trans [OF _ *])
    show (3 - 2 * sqrt 2) * r * cmod (deriv f 0)
       $\leq$  (cmod z - (cmod z)2 / (r - cmod z)) * cmod (deriv f 0)
    by (simp add: le_less algebra_simps divide_simps power2_eq_square that)
  qed (use <0 < r> that in auto)
  show ?thesis
    using <0 < r> sq201  $\exists$  C_def <0 < C> sq3
    by (intro ball_subset_open_map_image [OF 1 2 _ bounded_ball]) auto
qed
also have ...  $\subseteq$  f ' ball 0 r
proof -
  have  $\bigwedge$ x. (1 - sqrt 2 / 2) * r  $\leq$  r
    using <0 < r> by (auto simp: field_simps)
  then show ?thesis
    by auto
qed
finally show ?thesis .
qed
qed

```

lemma Bloch_lemma:

```

assumes holf: f holomorphic_on cball a r and 0 < r
  and le:  $\bigwedge$ z. z  $\in$  ball a r  $\implies$  norm(deriv f z)  $\leq$  2 * norm(deriv f a)
shows ball (f a) ((3 - 2 * sqrt 2) * r * norm(deriv f a))  $\subseteq$  f ' ball a r (is ?lhs
 $\subseteq$  ?rhs)
proof -
  have fz: ( $\lambda$ z. f (a + z)) = f o ( $\lambda$ z. (a + z))
    by (simp add: o_def)
  have hol0: ( $\lambda$ z. f (a + z)) holomorphic_on cball 0 r
    unfolding fz by (intro holomorphic_intros holf holomorphic_on_compose |
simp)+
  then have [simp]:  $\bigwedge$ x. norm x < r  $\implies$  ( $\lambda$ z. f (a + z)) field_differentiable at x
    by (metis open_ball at_within_open ball_subset_cball diff_0 dist_norm holomorphic_on_def holomorphic_on_subset mem_ball norm_minus_cancel)
  have [simp]:  $\bigwedge$ z. norm z < r  $\implies$  f field_differentiable at (a + z)
    by (metis holf open_ball add_diff_cancel_left' dist_complex_def holomorphic_on_imp_differentiable_at holomorphic_on_subset interior_cball interior_subset mem_ball norm_minus_commute)
  then have [simp]: f field_differentiable at a
    by (metis add.comm_neutral <0 < r> norm_eq_zero)
  have hol1: ( $\lambda$ z. f (a + z) - f a) holomorphic_on cball 0 r

```

```

  by (intro holomorphic_intros hol0)
  moreover have  $\bigwedge z. \text{cmod } z < r \implies$ 
     $\text{cmod } (\text{deriv } (\lambda z. f (a + z)) z) \leq 2 * \text{cmod } (\text{deriv } (\lambda z. f (a + z)) 0)$ 
  by (simp add: fz deriv_chain dist_norm le)
  ultimately have  $\S: \text{ball } 0 ((3 - 2 * \text{sqrt } 2) * r * \text{norm } (\text{deriv } (\lambda z. f (a + z))$ 
   $- f a 0))$ 
     $\subseteq (\lambda z. f (a + z) - f a) ' \text{ball } 0 r$ 
  using  $\langle 0 < r \rangle$  by (intro Bloch_lemma_0) auto
  show ?thesis
  proof clarify
    fix x
    assume  $x \in ?lhs$ 
    with subsetD [OF  $\S$ , of  $x - f a$ ] show  $x \in ?rhs$ 
    by (force simp: fz  $\langle 0 < r \rangle$  dist_norm deriv_chain field_differentiable_compose)
  qed
qed

```

proposition *Bloch_unit*:

assumes *holf*: *f* holomorphic_on ball a 1 and [*simp*]: *deriv f a* = 1
 obtains *b r* where $1/12 < r$ and $\text{ball } b r \subseteq f ' (\text{ball } a 1)$

proof –

```

  define r :: real where  $r = 249/256$ 
  have  $0 < r r < 1$  by (auto simp: r_def)
  define g where  $g z = \text{deriv } f z * \text{of\_real}(r - \text{norm}(z - a))$  for z
  have deriv f holomorphic_on ball a 1
    by (rule holomorphic_deriv [OF holf open_ball])
  then have continuous_on (ball a 1) (deriv f)
    using holomorphic_on_imp_continuous_on by blast
  then have continuous_on (cball a r) (deriv f)
    by (rule continuous_on_subset) (simp add: cball_subset_ball_iff  $\langle r < 1 \rangle$ )
  then have continuous_on (cball a r) g
    by (simp add: g_def continuous_intros)
  then have 1: compact ( $g ' \text{cball } a r$ )
    by (rule compact_continuous_image [OF _ compact_cball])
  have 2:  $g ' \text{cball } a r \neq \{\}$ 
    using  $\langle r > 0 \rangle$  by auto
  obtain p where pr:  $p \in \text{cball } a r$ 
    and pge:  $\bigwedge y. y \in \text{cball } a r \implies \text{norm } (g y) \leq \text{norm } (g p)$ 
    using distance_attains_sup [OF 1 2, of 0] by force
  define t where  $t = (r - \text{norm}(p - a)) / 2$ 
  have  $\text{norm } (p - a) \neq r$ 
    using pge [of a]  $\langle r > 0 \rangle$  by (auto simp: g_def norm_mult)
  then have  $\text{norm } (p - a) < r$  using pr
    by (simp add: norm_minus_commute dist_norm)
  then have  $0 < t$ 
    by (simp add: t_def)
  have cpt:  $\text{cball } p t \subseteq \text{ball } a r$ 
    using  $\langle 0 < t \rangle$  by (simp add: cball_subset_ball_iff dist_norm t_def field_simps)
  have gen_le_dfp:  $\text{norm } (\text{deriv } f y) * (r - \text{norm } (y - a)) / (r - \text{norm } (p - a))$ 

```

```

≤ norm (deriv f p)
  if y ∈ cball a r for y
proof -
  have [simp]: norm (y - a) ≤ r
    using that by (simp add: dist_norm norm_minus_commute)
  have norm (g y) ≤ norm (g p)
    using pge [OF that] by simp
  then have norm (deriv f y) * abs (r - norm (y - a)) ≤ norm (deriv f p) *
    abs (r - norm (p - a))
    by (simp only: dist_norm g_def norm_mult norm_of_real)
  with that ⟨norm (p - a) < r⟩ show ?thesis
    by (simp add: dist_norm field_split_simps)
qed
have le_norm_dfp: r / (r - norm (p - a)) ≤ norm (deriv f p)
  using gen_le_dfp [of a] ⟨r > 0⟩ by auto
have 1: f holomorphic_on cball p t
  using cpt ⟨r < 1⟩ order_subst1 subset_ball
  by (force simp: intro!: holomorphic_on_subset [OF holf])
have 2: norm (deriv f z) ≤ 2 * norm (deriv f p) if z ∈ ball p t for z
proof -
  have z: z ∈ cball a r
    by (meson ball_subset_cball subsetD cpt that)
  then have norm(z - a) < r
  by (metis ball_subset_cball contra_subsetD cpt dist_norm mem_ball norm_minus_commute
that)
  have norm (deriv f z) * (r - norm (z - a)) / (r - norm (p - a)) ≤ norm
(deriv f p)
    using gen_le_dfp [OF z] by simp
  with ⟨norm (z - a) < r⟩ ⟨norm (p - a) < r⟩
  have norm (deriv f z) ≤ (r - norm (p - a)) / (r - norm (z - a)) * norm
(deriv f p)
    by (simp add: field_simps)
  also have ... ≤ 2 * norm (deriv f p)
  proof (rule mult_right_mono)
    show (r - cmod (p - a)) / (r - cmod (z - a)) ≤ 2
      using that ⟨norm (p - a) < r⟩ ⟨norm(z - a) < r⟩ dist_triangle3 [of z a p]
      by (simp add: field_simps t_def dist_norm [symmetric])
    qed auto
  finally show ?thesis .
qed
have sqrt2: sqrt 2 < 2113/1494
  by (rule real_less_sqrt) (auto simp: power2_eq_square)
then have sq3: 0 < 3 - 2 * sqrt 2 by simp
have 1 / 12 / ((3 - 2 * sqrt 2) / 2) < r
  using sq3 sqrt2 by (auto simp: field_simps r_def)
also have ... ≤ cmod (deriv f p) * (r - cmod (p - a))
  using ⟨norm (p - a) < r⟩ le_norm_dfp by (simp add: pos_divide_le_eq)
finally have 1 / 12 < cmod (deriv f p) * (r - cmod (p - a)) * ((3 - 2 * sqrt
2) / 2)

```

```

    using pos_divide_less_eq half_gt_zero_iff sq3 by blast
  then have **:  $1 / 12 < (3 - 2 * \text{sqrt } 2) * t * \text{norm } (\text{deriv } f p)$ 
    using sq3 by (simp add: mult.commute t_def)
  have ball (f p)  $((3 - 2 * \text{sqrt } 2) * t * \text{norm } (\text{deriv } f p)) \subseteq f \text{ ' ball } p t$ 
    by (rule Bloch_lemma [OF 1 ‹0 < t› 2])
  also have ...  $\subseteq f \text{ ' ball } a 1$ 
    by (meson ‹r < 1› ball_subset_cball cpt_dual_order.trans image_mono less_le_not_le
subset_ball)
  finally have ball (f p)  $((3 - 2 * \text{sqrt } 2) * t * \text{norm } (\text{deriv } f p)) \subseteq f \text{ ' ball } a 1 .$ 
  with ** show ?thesis
    by (rule that)
qed

```

theorem Bloch:

```

  assumes holf:  $f \text{ holomorphic\_on ball } a r$  and  $0 < r$ 
    and r':  $r' \leq r * \text{norm } (\text{deriv } f a) / 12$ 
  obtains b where  $\text{ball } b r' \subseteq f \text{ ' (ball } a r)$ 
proof (cases  $\text{deriv } f a = 0$ )
  case True with r' show ?thesis
    using ball_eq_empty that by fastforce
next
  case False
  define C where  $C = \text{deriv } f a$ 
  have  $0 < \text{norm } C$  using False by (simp add: C_def)
  have dfa:  $f \text{ field\_differentiable\_at } a$ 
    using ‹0 < r› holomorphic_on_imp_differentiable_at [OF holf] by auto
  have fo:  $(\lambda z. f (a + \text{of\_real } r * z)) = f o (\lambda z. (a + \text{of\_real } r * z))$ 
    by (simp add: o_def)
  have holf':  $f \text{ holomorphic\_on } (\lambda z. a + \text{complex\_of\_real } r * z) \text{ ' ball } 0 1$ 
    using ‹0 < r› holomorphic_on_subset [OF holf] by (force simp: dist_norm
norm_mult)
  have 1:  $(\lambda z. f (a + r * z) / (C * r)) \text{ holomorphic\_on ball } 0 1$ 
    using ‹0 < r› ‹0 < norm C›
    by (intro holomorphic_intros holomorphic_on_compose holf'; simp add: fo)+
  have  $((\lambda z. f (a + \text{of\_real } r * z) / (C * \text{of\_real } r)) \text{ has\_field\_derivative}$ 
     $(\text{deriv } f (a + \text{of\_real } r * z) / C)) \text{ (at } z)$ 
    if  $\text{norm } z < 1$  for  $z$ 
  proof -
    have fd:  $f \text{ field\_differentiable\_at } (a + \text{complex\_of\_real } r * z)$ 
      using ‹0 < r› by (simp_all add: dist_norm norm_mult holomorphic_on_imp_differentiable_at
[OF holf] that)
    have *:  $((\lambda x. f (a + \text{of\_real } r * x)) \text{ has\_field\_derivative}$ 
       $(\text{deriv } f (a + \text{of\_real } r * z) * \text{of\_real } r)) \text{ (at } z)$ 
      by (rule fd DERIV_chain [OF field_differentiable_derivI] derivative_eq_intros
| simp add: fo)+
    show ?thesis
      apply (rule derivative_eq_intros * | simp)+
      using ‹0 < r› by (auto simp: C_def False)
  qed
qed

```

```

obtain  $f'$  where ( $f$  has_field_derivative  $f'$ ) (at  $a$ )
  using dfa field_differentiable_def by blast
then have  $\exists c. ((\lambda c. f (a + \text{complex\_of\_real } r * c)) \text{ has\_field\_derivative } c)$  (at
0)
  by (metis (no_types) DERIV_chain2 add_cancel_left_right field_differentiable_add_const

      field_differentiable_def field_differentiable_linear_mult_eq_0_iff)
then have ( $\lambda w. f (a + \text{complex\_of\_real } r * w)$ ) field_differentiable at 0
  by (simp add: field_differentiable_def)
then have deriv ( $\lambda z. f (a + \text{of\_real } r * z) / (C * \text{of\_real } r)$ ) 0
      = deriv ( $\lambda z. f (a + \text{of\_real } r * z)$ ) 0 / (C * of_real r)
  by (rule deriv_cdivide_right)
also have ... = 1
  using  $\langle 0 < r \rangle$  by (simp add: C_def False fo_derivative_intros dfa deriv_chain)
finally have 2: deriv ( $\lambda z. f (a + \text{of\_real } r * z) / (C * \text{of\_real } r)$ ) 0 = 1 .
have sb1:  $(*) (C * r) \text{ ' } (\lambda z. f (a + \text{of\_real } r * z) / (C * r)) \text{ ' ball } 0 \ 1$ 
       $\subseteq f \text{ ' ball } a \ r$ 
  using  $\langle 0 < r \rangle$  by (auto simp: dist_norm norm_mult C_def False)
have sb2: ball (C * r * b) r'  $\subseteq (*) (C * r) \text{ ' ball } b \ t$ 
      if  $1 / 12 < t$  for  $b \ t$ 
proof -
  have *:  $r * \text{cmod} (deriv f a) / 12 \leq r * (t * \text{cmod} (deriv f a))$ 
  using that  $\langle 0 < r \rangle$  less_eq_real_def mult.commute mult.right_neutral mult_left_mono
norm_ge_zero times_divide_eq_right
  by auto
  show ?thesis
  apply clarify
  apply (rule_tac  $x=x / (C * r)$  in image_eqI)
  using  $\langle 0 < r \rangle$  apply (simp_all add: dist_norm norm_mult norm_divide
C_def False field_simps)
  using * r' by linarith
qed
show ?thesis
  apply (rule Bloch_unit [OF 1 2])
  using image_mono sb1 sb2 that by fastforce
qed

corollary Bloch_general:
  assumes holf:  $f$  holomorphic_on  $S$  and  $a \in S$ 
  and tle:  $\bigwedge z. z \in \text{frontier } S \implies t \leq \text{dist } a \ z$ 
  and rle:  $r \leq t * \text{norm}(deriv f a) / 12$ 
  obtains  $b$  where ball  $b \ r \subseteq f \text{ ' } S$ 
proof -
consider  $r \leq 0 \mid 0 < t * \text{norm}(deriv f a) / 12$  using rle by force
then show ?thesis
proof cases
  case 1 then show ?thesis
  by (simp add: ball_empty that)
next

```

```

case 2
show ?thesis
proof (cases deriv f a = 0)
  case True then show ?thesis
    using rle by (simp add: ball_empty that)
  next
  case False
  then have t > 0
    using 2 by (force simp: zero_less_mult_iff)
  have  $\neg \text{ball } a \ t \subseteq S \implies \text{ball } a \ t \cap \text{frontier } S \neq \{\}$ 
    by (metis Diff_eq_empty_iff <0 < t> <a ∈ S> closure_Int_ball_not_empty
closure_subset connected_Int_frontier_connected_ball_inf.commute)
  with tle have *:  $\text{ball } a \ t \subseteq S$  by fastforce
  then have 1:  $f \text{ holomorphic\_on } \text{ball } a \ t$ 
    using holf using holomorphic_on_subset by blast
  show ?thesis
    using Bloch [OF 1 <t > 0> rle] * by (metis image_mono order_trans that)
qed
qed
qed
end

```

6 The Great Picard Theorem and its Applications

Ported from HOL Light (cauchy.ml) by L C Paulson, 2017

```

theory Great_Picard
  imports Conformal_Mappings
begin

```

6.1 Schottky's theorem

lemma Schottky_lemma0:

```

  assumes holf:  $f \text{ holomorphic\_on } S$  and cons:  $\text{contractible } S$  and  $a \in S$ 
  and f:  $\bigwedge z. z \in S \implies f z \neq 1 \wedge f z \neq -1$ 
  obtains g where  $g \text{ holomorphic\_on } S$ 
     $\text{norm}(g a) \leq 1 + \text{norm}(f a) / 3$ 
     $\bigwedge z. z \in S \implies f z = \cos(\text{of\_real } \pi * g z)$ 

```

proof –

```

  obtain g where holg:  $g \text{ holomorphic\_on } S$  and g:  $\text{norm}(g a) \leq \pi + \text{norm}(f a)$ 
  and f_eq_cos:  $\bigwedge z. z \in S \implies f z = \cos(g z)$ 
  using contractible_imp_holomorphic_arccos_bounded [OF assms]
  by blast
  show ?thesis
  proof
    show  $(\lambda z. g z / \pi) \text{ holomorphic\_on } S$ 
      by (auto intro: holomorphic_intros holg)
    have  $3 \leq \pi$ 

```



```

    using pi_approx by force
  have  $3 * \text{norm}(g a) \leq 3 * (\text{pi} + \text{norm}(f a))$ 
    using g by auto
  also have  $\dots \leq \text{pi} * 3 + \text{pi} * \text{cmod}(f a)$ 
    using  $\langle 3 \leq \text{pi} \rangle$  by (simp add: mult_right_mono algebra_simps)
  finally show  $\text{cmod}(g a / \text{complex\_of\_real pi}) \leq 1 + \text{cmod}(f a) / 3$ 
    by (simp add: field_simps norm_divide)
  show  $\bigwedge z. z \in S \implies f z = \cos(\text{complex\_of\_real pi} * (g z / \text{complex\_of\_real pi}))$ 
    by (simp add: f_eq_cos)
  qed
qed

```

lemma Schottky_lemma1:

```

  fixes n::nat
  assumes  $0 < n$ 
  shows  $0 < n + \text{sqrt}(\text{real } n^2 - 1)$ 
proof -
  have  $0 < n * n$ 
    by (simp add: assms)
  then show ?thesis
    by (metis add.commute add.right_neutral add_pos_nonneg assms diff_ge_0_iff_ge
      nat_less_real_le of_nat_0 of_nat_0_less_iff of_nat_power power2_eq_square
      real_sqrt_ge_0_iff)
qed

```

lemma Schottky_lemma2:

```

  fixes x::real
  assumes  $0 \leq x$ 
  obtains n where  $0 < n \wedge |x - \ln(\text{real } n + \text{sqrt}((\text{real } n)^2 - 1)) / \text{pi}| < 1/2$ 
proof -
  obtain n0::nat where  $0 < n0 \wedge \ln(n0 + \text{sqrt}(\text{real } n0^2 - 1)) / \text{pi} \leq x$ 
  proof
    show  $\ln(\text{real } 1 + \text{sqrt}(\text{real } 1^2 - 1)) / \text{pi} \leq x$ 
      by (auto simp: assms)
  qed auto
  moreover
  obtain M::nat where  $\bigwedge n. \llbracket 0 < n; \ln(n + \text{sqrt}(\text{real } n^2 - 1)) / \text{pi} \leq x \rrbracket \implies n \leq M$ 
  proof
    fix n::nat
    assume  $0 < n \wedge \ln(n + \text{sqrt}((\text{real } n)^2 - 1)) / \text{pi} \leq x$ 
    then have  $\ln(n + \text{sqrt}((\text{real } n)^2 - 1)) \leq x * \text{pi}$ 
      by (simp add: field_split_simps)
    then have  $*: \exp(\ln(n + \text{sqrt}((\text{real } n)^2 - 1))) \leq \exp(x * \text{pi})$ 
      by blast
    have 0:  $0 \leq \text{sqrt}((\text{real } n)^2 - 1)$ 

```

```

    using ‹0 < n› by auto
    have n + sqrt ((real n)2 - 1) = exp (ln (n + sqrt ((real n)2 - 1)))
    by (simp add: Suc_leI ‹0 < n› add_pos_nonneg real_of_nat_ge_one_iff)
    also have ... ≤ exp (x * pi)
    using * by blast
    finally have real n ≤ exp (x * pi)
    using 0 by linarith
    then show n ≤ nat (ceiling (exp(x * pi)))
    by linarith
qed
ultimately obtain n where
  0 < n and le_x: ln(n + sqrt(real n ^ 2 - 1)) / pi ≤ x
  and le_n:  $\bigwedge k. \llbracket 0 < k; \ln(k + \text{sqrt}(\text{real } k^2 - 1)) / \text{pi} \leq x \rrbracket \implies k \leq$ 
n
  using bounded_Max_nat [of  $\lambda n. 0 < n \wedge \ln(n + \text{sqrt}(\text{real } n^2 - 1)) / \text{pi} \leq$ 
x] by metis
define a where a  $\equiv \ln(n + \text{sqrt}(\text{real } n^2 - 1)) / \text{pi}$ 
define b where b  $\equiv \ln(1 + \text{real } n + \text{sqrt}((1 + \text{real } n)^2 - 1)) / \text{pi}$ 
have le_xa: a ≤ x
and le_na:  $\bigwedge k. \llbracket 0 < k; \ln(k + \text{sqrt}(\text{real } k^2 - 1)) / \text{pi} \leq x \rrbracket \implies k \leq n$ 
  using le_x le_n by (auto simp: a_def)
moreover have x < b
  using le_n [of Suc n] by (force simp: b_def)
moreover have b - a < 1
proof -
  have ln (1 + real n + sqrt ((1 + real n)2 - 1)) - ln (real n + sqrt ((real n)2 - 1)) =
    ln ((1 + real n + sqrt ((1 + real n)2 - 1)) / (real n + sqrt ((real n)2 - 1)))
  by (simp add: ‹0 < n› Schottky_lemma1 add_pos_nonneg ln_divide_pos
[symmetric])
  also have ... ≤ 3
proof (cases n = 1)
  case True
  have sqrt 3 ≤ 2
  by (simp add: real_le_sqrt)
  then have (2 + sqrt 3) ≤ 4
  by simp
  also have ... ≤ exp 3
  using exp_ge_add_one_self [of 3::real] by simp
  finally have ln (2 + sqrt 3) ≤ 3
  by (metis add_nonneg_nonneg add_pos_nonneg dbl_def dbl_inc_def
dbl_inc_simps(3)
  dbl_simps(3) exp_gt_zero ln_exp ln_le_cancel_iff real_sqrt_ge_0_iff
zero_le_one zero_less_one)
  then show ?thesis
  by (simp add: True)
next
  case False with ‹0 < n› have 1 < n 2 ≤ n

```

```

    by linarith+
  then have 1:  $1 \leq \text{real } n * \text{real } n$ 
  by (metis less_imp_le_nat one_le_power power2_eq_square real_of_nat_ge_one_iff)
  have *:  $4 + (m+2) * 2 \leq (m+2) * ((m+2) * 3)$  for  $m::\text{nat}$ 
    by simp
  have  $4 + n * 2 \leq n * (n * 3)$ 
    using * [of  $n-2$ ]  $\langle 2 \leq n \rangle$ 
    by (metis le_add_diff_inverse2)
  then have **:  $4 + \text{real } n * 2 \leq \text{real } n * (\text{real } n * 3)$ 
  by (metis (mono_tags, opaque_lifting) of_nat_le_iff of_nat_add of_nat_mult
of_nat_numeral)
  have  $\text{sqrt } ((1 + \text{real } n)^2 - 1) \leq 2 * \text{sqrt } ((\text{real } n)^2 - 1)$ 
    by (auto simp: real_le_sqrt power2_eq_square algebra_simps 1 **)
  then
  have  $((1 + \text{real } n + \text{sqrt } ((1 + \text{real } n)^2 - 1)) / (\text{real } n + \text{sqrt } ((\text{real } n)^2 - 1))) \leq 2$ 
    using Schottky_lemma1  $\langle 0 < n \rangle$  by (simp add: field_split_simps)
  then have  $\ln ((1 + \text{real } n + \text{sqrt } ((1 + \text{real } n)^2 - 1)) / (\text{real } n + \text{sqrt } ((\text{real } n)^2 - 1))) \leq \ln 2$ 
    using Schottky_lemma1 [of  $n$ ]  $\langle 0 < n \rangle$ 
    by (simp add: field_split_simps add_pos_nonneg)
  also have  $\dots \leq 3$ 
    using ln_add_one_self_le_self [of 1] by auto
  finally show ?thesis .
qed
also have  $\dots < \pi$ 
  using pi_approx by simp
finally show ?thesis
  by (simp add: a_def b_def field_split_simps)
qed
ultimately have  $|x - a| < 1/2 \vee |x - b| < 1/2$ 
  by (auto simp: abs_if)
then show thesis
proof
  assume  $|x - a| < 1/2$ 
  then show ?thesis
    by (rule_tac  $n=n$  in that) (auto simp: a_def  $\langle 0 < n \rangle$ )
next
  assume  $|x - b| < 1/2$ 
  then show ?thesis
    by (rule_tac  $n=\text{Suc } n$  in that) (auto simp: b_def  $\langle 0 < n \rangle$ )
qed
qed

```

lemma Schottky_lemma3:

```

  fixes  $z::\text{complex}$ 
  assumes  $z \in (\bigcup m \in \text{Ints. } \bigcup n \in \{0 < ..\}. \{\text{Complex } m (\ln(n + \text{sqrt}(\text{real } n ^ 2 - 1)) / \pi)\})$ 

```

```

    ∪ (∪ m ∈ Ints. ∪ n ∈ {0<..}. {Complex m (-ln(n + sqrt(real n ^ 2 -
1)) / pi)})
  shows cos(pi * cos(pi * z)) = 1 ∨ cos(pi * cos(pi * z)) = -1
proof -
  have sqrt2 [simp]: complex_of_real (sqrt x) * complex_of_real (sqrt x) = x if x
≥ 0 for x::real
  by (metis abs_of_nonneg of_real_mult real_sqrt_mult_self that)
  define plusi where plusi (e::complex) ≡ e + inverse e for e
  have 1: ∃ k. plusi (exp (i * (of_int m * complex_of_real pi) - ln (real n + sqrt
((real n)2 - 1)))) = of_int k * 2
    (is ∃ k. ?Φ k)
    if n > 0 for m n
proof -
  have eqq: e ≠ 0 ⇒ plusi e = n ↔ (inverse e) ^ 2 = n/e - 1 for n e::complex
    by (auto simp: plusi_def field_simps power2_eq_square)
  have [simp]: 1 ≤ real n * real n
    using nat_0_less_mult_iff nat_less_real_le that by force
  consider odd m | even m
  by blast
  then have ∃ k. ?Φ k
proof cases
  case 1
  then have ?Φ (- n)
    using Schottky_lemma1 [OF that]
    by (simp add: eqq) (simp add: power2_eq_square exp_diff exp_Euler
exp_of_real algebra_simps sin_int_times_real cos_int_times_real)
  then show ?thesis ..
next
  case 2
  then have ?Φ n
    using Schottky_lemma1 [OF that]
    by (simp add: eqq) (simp add: power2_eq_square exp_diff exp_Euler
exp_of_real algebra_simps)
  then show ?thesis ..
qed
  then show ?thesis by blast
qed
  have 2: ∃ k. plusi (exp (i * (of_int m * complex_of_real pi) +
(ln (real n + sqrt ((real n)2 - 1)))))) = of_int k * 2
    (is ∃ k. ?Φ k)
    if n > 0 for m n
proof -
  have eqq: e ≠ 0 ⇒ plusi e = n ↔ e2 - n*e + 1 = 0 for n e::complex
    by (auto simp: plusi_def field_simps power2_eq_square)
  have [simp]: 1 ≤ real n * real n
    by (metis One_nat_def add.commute nat_less_real_le of_nat_1 of_nat_Suc
one_le_power power2_eq_square that)
  consider odd m | even m
  by blast

```

```

then have  $\exists k. ?\Phi k$ 
proof cases
  case 1
  then have  $?\Phi (-n)$ 
    using Schottky_lemma1 [OF that]
    by (simp add: eeq) (simp add: power2_eq_square exp_add exp_Euler
exp_of_real algebra_simps sin_int_times_real cos_int_times_real)
  then show ?thesis ..
next
  case 2
  then have  $?\Phi n$ 
    using Schottky_lemma1 [OF that]
    by (simp add: eeq) (simp add: power2_eq_square exp_add exp_Euler
exp_of_real algebra_simps)
  then show ?thesis ..
qed
then show ?thesis by blast
qed
have  $\exists x. \cos(\text{complex\_of\_real } \pi * z) = \text{of\_int } x$ 
  using assms
  apply (auto simp: Ints_def cos_exp_eq exp_minus Complex_eq simp flip:
plusi_def)
  apply (auto simp: algebra_simps dest: 1 2)
done
then have  $\sin(\pi * \cos(\pi * z)) ^ 2 = 0$ 
  by (simp add: Complex_Transcendental.sin_eq_0)
then have  $1 - \cos(\pi * \cos(\pi * z)) ^ 2 = 0$ 
  by (simp add: sin_squared_eq)
then show ?thesis
  using power2_eq_1_iff by auto
qed

```

theorem Schottky:

```

assumes holf:  $f$  holomorphic_on cball 0 1
  and nof0:  $\text{norm}(f 0) \leq r$ 
  and not01:  $\bigwedge z. z \in \text{cball } 0 1 \implies \neg(f z = 0 \vee f z = 1)$ 
  and 0 < t t < 1  $\text{norm } z \leq t$ 
  shows  $\text{norm}(f z) \leq \exp(\pi * \exp(\pi * (2 + 2 * r + 12 * t / (1 - t))))$ 
proof -
  obtain h where holf:  $h$  holomorphic_on cball 0 1
    and nh0:  $\text{norm}(h 0) \leq 1 + \text{norm}(2 * f 0 - 1) / 3$ 
    and h:  $\bigwedge z. z \in \text{cball } 0 1 \implies 2 * f z - 1 = \cos(\text{of\_real } \pi * h z)$ 
  proof (rule Schottky_lemma0 [of  $\lambda z. 2 * f z - 1$  cball 0 1 0])
    show  $(\lambda z. 2 * f z - 1)$  holomorphic_on cball 0 1
      by (intro holomorphic_intros holf)
    show contractible (cball (0::complex) 1)
      by (auto simp: convex_imp_contractible)
    show  $\bigwedge z. z \in \text{cball } 0 1 \implies 2 * f z - 1 \neq 1 \wedge 2 * f z - 1 \neq -1$ 

```

```

    using not01 by force
  qed auto
  obtain g where holg: g holomorphic_on cball 0 1
    and ng0: norm(g 0) ≤ 1 + norm(h 0) / 3
    and g:  $\bigwedge z. z \in \text{cball } 0 \ 1 \implies h \ z = \cos(\text{of\_real } \pi * g \ z)$ 
  proof (rule Schottky_lemma0 [OF holf convex_imp_contractible, of 0])
    show  $\bigwedge z. z \in \text{cball } 0 \ 1 \implies h \ z \neq 1 \wedge h \ z \neq -1$ 
      using h not01 by fastforce+
  qed auto
  have g0_2_f0: norm(g 0) ≤ 2 + norm(f 0)
  proof -
    have cmod (2 * f 0 - 1) ≤ cmod (2 * f 0) + 1
      by (metis norm_one norm_triangle_ineq4)
    also have ... ≤ 6 + 9 * cmod (f 0)
      by auto
    finally have 1 + norm(2 * f 0 - 1) / 3 ≤ (2 + norm(f 0) - 1) * 3
      by (simp add: divide_simps)
    with nh0 have norm(h 0) ≤ (2 + norm(f 0) - 1) * 3
      by linarith
    then have 1 + norm(h 0) / 3 ≤ 2 + norm(f 0)
      by simp
    with ng0 show ?thesis
      by auto
  qed
  have z ∈ ball 0 1
    using assms by auto
  have norm_g_12: norm(g z - g 0) ≤ (12 * t) / (1 - t)
  proof -
    obtain g' where g':  $\bigwedge x. x \in \text{cball } 0 \ 1 \implies (g \text{ has\_field\_derivative } g' \ x)$  (at x
    within cball 0 1)
      using holg [unfolded holomorphic_on_def field_differentiable_def] by metis
    have int_g': (g' has_contour_integral g z - g 0) (linepath 0 z)
      using contour_integral_primitive [OF g' valid_path_linepath, of 0 z]
      using ⟨z ∈ ball 0 1⟩ segment_bound1 by fastforce
    have cmod (g' w) ≤ 12 / (1 - t) if w ∈ closed_segment 0 z for w
  proof -
    have w: w ∈ ball 0 1
      using segment_bound [OF that] ⟨z ∈ ball 0 1⟩ by simp
    have *:  $\llbracket \bigwedge b. (\exists w \in T \cup U. w \in \text{ball } b \ 1); \bigwedge x. x \in D \implies g \ x \notin T \cup U \rrbracket$ 
     $\implies \nexists b. \text{ball } b \ 1 \subseteq g \ ' \ D$  for T U D
      by force
    have ttt: 1 - t ≤ dist w u if cmod u = 1 for u
      using ⟨norm z ≤ t⟩ segment_bound1 [OF ⟨w ∈ closed_segment 0 z⟩]
      norm_triangle_ineq2 [of u w] that
      by (simp add: dist_norm norm_minus_commute)
    have  $\nexists b. \text{ball } b \ 1 \subseteq g \ ' \ \text{cball } 0 \ 1$ 
  proof (rule *)
    show  $(\exists w \in (\bigcup m \in \text{Ints. } \bigcup n \in \{0 <.. \}. \{\text{Complex } m \ (\ln(n + \text{sqrt}(\text{real } n$ 
    ^ 2 - 1)) / \pi i\}) \cup

```

```

      ( $\bigcup m \in \text{Ints. } \bigcup n \in \{0 < ..\}. \{ \text{Complex } m (-\ln(n + \sqrt{(\text{real } n)^2 - 1})) / \pi i \}$ ).  $w \in \text{ball } b \ 1$ ) for  $b$ 
    proof -
      obtain  $m$  where  $m: m \in \mathbb{Z} \mid \text{Re } b - m \leq 1/2$ 
      by (metis Ints_of_int abs_minus_commute of_int_round_abs_le)
      show ?thesis
      proof (cases 0::real Im b rule: le_cases)
        case le
          then obtain  $n$  where  $0 < n$  and  $n: |\text{Im } b - \ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i| < 1/2$ 
          using Schottky_lemma2 [of Im b] by blast
          have  $\text{dist } b (\text{Complex } m (\text{Im } b)) \leq 1/2$ 
          by (metis cancel_comm_monoid_add_class.diff_cancel cmod_eq_Re complex.sel(1) complex.sel(2) dist_norm m(2) minus_complex.code)
          moreover
            have  $\text{dist} (\text{Complex } m (\text{Im } b)) (\text{Complex } m (\ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i)) < 1/2$ 
            using  $n$  by (simp add: complex_norm cmod_eq_Re complex_diff dist_norm del: Complex_eq)
            ultimately have  $\text{dist } b (\text{Complex } m (\ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i)) < 1$ 
            by (simp add: dist_triangle_lt [of b Complex m (Im b)] dist_commute)
            with  $le \ m \ \langle 0 < n \rangle$  show ?thesis
            apply (rule_tac  $x = \text{Complex } m (\ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i)$  in bexI)
            by (force simp del: Complex_eq greaterThan_0)+
          next
            case ge
              then obtain  $n$  where  $0 < n$  and  $n: |-\text{Im } b - \ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i| < 1/2$ 
              using Schottky_lemma2 [of - Im b] by auto
              have  $\text{dist } b (\text{Complex } m (\text{Im } b)) \leq 1/2$ 
              by (metis cancel_comm_monoid_add_class.diff_cancel cmod_eq_Re complex.sel(1) complex.sel(2) dist_norm m(2) minus_complex.code)
              moreover
                have  $\text{dist} (\text{Complex } m (-\ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i)) (\text{Complex } m (\text{Im } b))$ 
                   $= |-\text{Im } b - \ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i|$ 
                by (simp add: complex_norm dist_norm cmod_eq_Re complex_diff)
                ultimately have  $\text{dist } b (\text{Complex } m (-\ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i)) < 1$ 
                using  $n$  by (simp add: dist_triangle_lt [of b Complex m (Im b)] dist_commute)
                with  $ge \ m \ \langle 0 < n \rangle$  show ?thesis
                by (rule_tac  $x = \text{Complex } m (-\ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i)$  in bexI) auto
              qed
            qed
          show  $g \ v \notin (\bigcup m \in \text{Ints. } \bigcup n \in \{0 < ..\}. \{ \text{Complex } m (\ln(n + \sqrt{(\text{real } n)^2 - 1}) / \pi i) \})$ 

```

```

2 - 1)) / pi)) ∪
  (∪ m ∈ Ints. ∪ n ∈ {0<..}. {Complex m (-ln(n + sqrt(real n ^ 2
- 1)) / pi)})
  if v ∈ cball 0 1 for v
  using not01 [OF that]
  by (force simp: g [OF that, symmetric] h [OF that, symmetric] dest!:
Schottky_lemma3 [of g v])
qed
then have 12: (1 - t) * cmod (deriv g w) / 12 < 1
  using Bloch_general [OF holg __ ttt, of 1] w by force
have g field_differentiable at w within cball 0 1
  using holg w by (simp add: holomorphic_on_def)
then have g field_differentiable at w within ball 0 1
  using ball_subset_cball field_differentiable_within_subset by blast
with w have der_gw: (g has_field_derivative deriv g w) (at w)
by (simp add: field_differentiable_within_open [of _ ball 0 1] field_differentiable_derivI)
with DERIV_unique [OF der_gw] g' w have deriv g w = g' w
by (metis open_ball at_within_open ball_subset_cball has_field_derivative_subset
subsetCE)
then show cmod (g' w) ≤ 12 / (1 - t)
  using g' 12 ‹t < 1› by (simp add: field_simps)
qed
then have cmod (g z - g 0) ≤ 12 / (1 - t) * cmod z
  using has_contour_integral_bound_linepath [OF int_g', of 12/(1 - t)] assms
  by simp
with ‹cmod z ≤ t› ‹t < 1› show ?thesis
  by (simp add: field_split_simps)
qed
have fz: f z = (1 + cos(of_real pi * h z)) / 2
  using h ‹z ∈ ball 0 1› by (auto simp: field_simps)
have cmod (f z) ≤ exp (cmod (complex_of_real pi * h z))
  by (simp add: fz mult.commute norm_cos_plus1_le)
also have ... ≤ exp (pi * exp (pi * (2 + 2 * r + 12 * t / (1 - t))))
proof (simp add: norm_mult)
  have cmod (g z - g 0) ≤ 12 * t / (1 - t)
    using norm_g 12 ‹t < 1› by (simp add: norm_mult)
  then have cmod (g z) - cmod (g 0) ≤ 12 * t / (1 - t)
    using norm_triangle_ineq2 order_trans by blast
  then have *: cmod (g z) ≤ 2 + 2 * r + 12 * t / (1 - t)
    using g0_2_f0 norm_ge_zero [of f 0] nof0
    by linarith
  have cmod (h z) ≤ exp (cmod (complex_of_real pi * g z))
    using ‹z ∈ ball 0 1› by (simp add: g norm_cos_le)
  also have ... ≤ exp (pi * (2 + 2 * r + 12 * t / (1 - t)))
    using ‹t < 1› nof0 * by (simp add: norm_mult)
  finally show cmod (h z) ≤ exp (pi * (2 + 2 * r + 12 * t / (1 - t))) .
qed
finally show ?thesis .
qed

```


6.2 The Little Picard Theorem

theorem *Landau_Picard*:

obtains *R*

where $\bigwedge z. 0 < R z$

$\bigwedge f. \llbracket f \text{ holomorphic_on } \text{cball } 0 (R(f\ 0));$

$\bigwedge z. \text{norm } z \leq R(f\ 0) \implies f z \neq 0 \wedge f z \neq 1 \rrbracket \implies \text{norm}(\text{deriv } f\ 0)$

< 1

proof –

define *R* **where** $R \equiv \lambda z. 3 * \exp(\text{pi} * \exp(\text{pi} * (2 + 2 * \text{cmod } z + 12)))$

show *?thesis*

proof

show *Rpos*: $\bigwedge z. 0 < R z$

by (*auto simp: R_def*)

show $\text{norm}(\text{deriv } f\ 0) < 1$

if *holf*: $f \text{ holomorphic_on } \text{cball } 0 (R(f\ 0))$

and *Rf*: $\bigwedge z. \text{norm } z \leq R(f\ 0) \implies f z \neq 0 \wedge f z \neq 1$ **for** *f*

proof –

let *?r* = $R(f\ 0)$

define *g* **where** $g \equiv f \circ (\lambda z. \text{of_real } ?r * z)$

have $0 < ?r$

using *Rpos* **by** *blast*

have *holg*: $g \text{ holomorphic_on } \text{cball } 0 1$

unfolding *g_def*

proof (*intro holomorphic_intros holomorphic_on_compose holomorphic_on_subset [OF holf]*)

show (*) (*complex_of_real* ($R(f\ 0)$)) ‘ $\text{cball } 0 1 \subseteq \text{cball } 0 (R(f\ 0))$ ’

using *Rpos* **by** (*auto simp: less_imp_le norm_mult*)

qed

have *: $\text{norm}(g z) \leq \exp(\text{pi} * \exp(\text{pi} * (2 + 2 * \text{norm}(f\ 0) + 12 * t / (1 - t))))$

if $0 < t < 1$ $\text{norm } z \leq t$ **for** *t z*

proof (*rule Schottky [OF holg]*)

show $\text{cmod}(g\ 0) \leq \text{cmod}(f\ 0)$

by (*simp add: g_def*)

show $\bigwedge z. z \in \text{cball } 0 1 \implies \neg (g z = 0 \vee g z = 1)$

using *Rpos* **by** (*simp add: g_def less_imp_le norm_mult Rf*)

qed (*auto simp: that*)

have *C1*: $g \text{ holomorphic_on } \text{ball } 0 (1/2)$

by (*rule holomorphic_on_subset [OF holg]*) *auto*

have *C2*: $\text{continuous_on } (\text{cball } 0 (1/2))\ g$

by (*meson cball_divide_subset_numeral holg holomorphic_on_imp_continuous_on holomorphic_on_subset*)

have *C3*: $\text{cmod}(g z) \leq R(f\ 0) / 3$ **if** $\text{cmod}(0 - z) = 1/2$ **for** *z*

proof –

have $\text{norm}(g z) \leq \exp(\text{pi} * \exp(\text{pi} * (2 + 2 * \text{norm}(f\ 0) + 12)))$

using * [*of 1/2*] **that** **by** *simp*

also have ... = *?r / 3*

by (*simp add: R_def*)

finally show *?thesis* .

```

qed
then have cmod_g'_le: cmod (deriv g 0) * 3 ≤ R (f 0) * 2
  using Cauchy_inequality [OF C1 C2 _ C3, of 1] by simp
have holf': f holomorphic_on ball 0 (R(f 0))
  by (rule holomorphic_on_subset [OF holf]) auto
then have fd0: f field_differentiable at 0
  by (rule holomorphic_on_imp_differentiable_at [OF _ open_ball])
  (auto simp: Rpos [of f 0])
have g_eq: deriv g 0 = of_real ?r * deriv f 0
  unfolding g_def
  by (metis DERIV_imp_deriv DERIV_chain DERIV_cmult_Id fd0 field_differentiable_derivI
mult.commute mult_zero_right)
show ?thesis
  using cmod_g'_le Rpos [of f 0] by (simp add: g_eq norm_mult)
qed
qed
qed

```

lemma little_Picard_01:

```

assumes holf: f holomorphic_on UNIV and f01:  $\bigwedge z. f z \neq 0 \wedge f z \neq 1$ 
obtains c where f = ( $\lambda x. c$ )
proof -
  obtain R
    where Rpos:  $\bigwedge z. 0 < R z$ 
    and R:  $\bigwedge h. \llbracket h \text{ holomorphic\_on } cball\ 0\ (R(h\ 0));$ 
     $\bigwedge z. norm\ z \leq R(h\ 0) \implies h\ z \neq 0 \wedge h\ z \neq 1 \rrbracket \implies norm(deriv$ 
h 0) < 1
  using Landau_Picard by metis
  have contf: continuous_on UNIV f
    by (simp add: holf holomorphic_on_imp_continuous_on)
  show ?thesis
  proof (cases  $\forall x. deriv\ f\ x = 0$ )
    case True
      have (f has_field_derivative 0) (at x) for x
        by (metis True UNIV_I holf holomorphic_derivI open_UNIV)
      then obtain c where  $\bigwedge x. f(x) = c$ 
        by (meson UNIV_I DERIV_zero_connected_constant [OF connected_UNIV
open_UNIV finite.emptyI contf])
      then show ?thesis
        using that by auto
    next
      case False
      then obtain w where w: deriv f w ≠ 0 by auto
      define fw where fw ≡ (f ∘ ( $\lambda z. w + z / deriv\ f\ w$ ))
      have norm_let1: norm(deriv fw 0) < 1
      proof (rule R)
        show fw holomorphic_on cball 0 (R (fw 0))
          unfolding fw_def
        by (intro holomorphic_intros holomorphic_on_compose w holomorphic_on_subset

```

```

[OF holf] subset_UNIV)
  show  $fw\ z \neq 0 \wedge fw\ z \neq 1$  if  $cmod\ z \leq R$  ( $fw\ 0$ ) for  $z$ 
    using f01 by (simp add: fw_def)
  qed
  have (fw has_field_derivative deriv f w * inverse (deriv f w)) (at 0)
    unfolding fw_def
    apply (intro DERIV_chain derivative_eq_intros w)+
    using holf holomorphic_derivI by (force simp: field_simps)+
  then show ?thesis
    using norm_let1 w by (simp add: DERIV_imp_deriv)
  qed
qed

```

```

theorem little_Picard:
  assumes holf:  $f$  holomorphic_on UNIV
    and  $a \neq b$  range  $f \cap \{a,b\} = \{\}$ 
    obtains  $c$  where  $f = (\lambda x. c)$ 
proof -
  let ?g =  $\lambda x. 1/(b-a)*(f\ x - b) + 1$ 
  obtain  $c$  where ?g =  $(\lambda x. c)$ 
  proof (rule little_Picard_01)
    show ?g holomorphic_on UNIV
      by (intro holomorphic_intros holf)
    show  $\bigwedge z. ?g\ z \neq 0 \wedge ?g\ z \neq 1$ 
      using assms by (auto simp: field_simps)
  qed auto
  then have ?g  $x = c$  for  $x$ 
    by meson
  then have  $f\ x = c * (b-a) + a$  for  $x$ 
    using assms by (auto simp: field_simps)
  then show ?thesis
    using that by blast
qed

```

A couple of little applications of Little Picard

```

lemma holomorphic_periodic_fixpoint:
  assumes holf:  $f$  holomorphic_on UNIV
    and  $p \neq 0$  and per:  $\bigwedge z. f(z + p) = f\ z$ 
    obtains  $x$  where  $f\ x = x$ 
proof -
  have False if non:  $\bigwedge x. f\ x \neq x$ 
  proof -
    obtain  $c$  where  $(\lambda z. f\ z - z) = (\lambda z. c)$ 
    proof (rule little_Picard)
      show  $(\lambda z. f\ z - z)$  holomorphic_on UNIV
        by (simp add: holf holomorphic_on_diff)
      show range  $(\lambda z. f\ z - z) \cap \{p,0\} = \{\}$ 
        using assms non by auto (metis add.commute diff_eq_eq)
    end
  end

```

```

    qed (auto simp: assms)
  with per show False
    by (metis add.commute add_cancel_left_left ⟨p ≠ 0⟩ diff_add_cancel)
  qed
  then show ?thesis
    using that by blast
  qed

lemma holomorphic_involution_point:
  assumes holfU: f holomorphic_on UNIV and non:  $\bigwedge a. f \neq (\lambda x. a + x)$ 
  obtains x where  $f(f x) = x$ 
proof -
  { assume non_ff [simp]:  $\bigwedge x. f(f x) \neq x$ 
    then have non_fp [simp]:  $f z \neq z$  for z
      by metis
    have holf: f holomorphic_on X for X
      using assms holomorphic_on_subset by blast
    obtain c where c:  $(\lambda x. (f(f x) - x)/(f x - x)) = (\lambda x. c)$ 
    proof (rule little_Picard_01)
      show  $(\lambda x. (f(f x) - x)/(f x - x))$  holomorphic_on UNIV
        using non_fp
      by (intro holomorphic_intros holf holomorphic_on_compose [unfolded o_def,
        OF holf]) auto
    qed auto
    then obtain c ≠ 0 c ≠ 1
      by (metis (no_types, opaque_lifting) non_ff diff_zero divide_eq_0_iff_right_inverse_eq)
    have eq:  $f(f x) - c * f x = x * (1 - c)$  for x
      using fun_cong [OF c, of x] by (simp add: field_simps)
    have df_times_dff:  $\text{deriv } f z * (\text{deriv } f (f z) - c) = 1 * (1 - c)$  for z
      proof (rule DERIV_unique)
        show  $(\lambda x. f (f x) - c * f x)$  has_field_derivative
           $\text{deriv } f z * (\text{deriv } f (f z) - c)$  (at z)
          by (rule derivative_eq_intros holomorphic_derivI [OF holfU]
            DERIV_chain [unfolded o_def, where f=f and g=f] | simp add:
            algebra_simps)+
        show  $(\lambda x. f (f x) - c * f x)$  has_field_derivative  $1 * (1 - c)$  (at z)
          by (simp add: eq mult_commute_abs)
      qed
    { fix z::complex
      obtain k where k:  $\text{deriv } f \circ f = (\lambda x. k)$ 
      proof (rule little_Picard)
        show  $(\text{deriv } f \circ f)$  holomorphic_on UNIV
          by (meson holfU holomorphic_deriv holomorphic_on_compose holomor-
            phic_on_subset open_UNIV subset_UNIV)
        obtain  $\text{deriv } f (f x) \neq 0$   $\text{deriv } f (f x) \neq c$  for x
          using df_times_dff ⟨c ≠ 1⟩ eq_iff_diff_eq_0
          by (metis lambda_one mult_zero_left mult_zero_right)
        then show  $\text{range } (\text{deriv } f \circ f) \cap \{0, c\} = \{\}$ 

```

```

    by force
  qed (use ‹c ≠ 0› in auto)
  have ¬ f constant_on UNIV
    by (meson UNIV_I non_ff constant_on_def)
  with holf open_mapping_thm have open(range f)
    by blast
  obtain l where l:  $\bigwedge x. f x - k * x = l$ 
  proof (rule DERIV_zero_connected_constant [of UNIV {}]  $\lambda x. f x - k * x$ ),
  simp_all)
    have deriv f w - k = 0 for w
      proof (rule analytic_continuation [OF open_UNIV connected_UNIV
subset_UNIV, of  $\lambda z. deriv f z - k f z$  range f w])
        show  $(\lambda z. deriv f z - k)$  holomorphic_on UNIV
          by (intro holomorphic_intros holf open_UNIV)
        show f z islimpt range f
          by (metis (no_types, lifting) IntI UNIV_I ‹open (range f)› im-
age_eqI inf.absorb_iff2 inf_aci(1) islimpt_UNIV islimpt_eq_acc_point open_Int
top_greatest)
        show  $\bigwedge z. z \in range f \implies deriv f z - k = 0$ 
          by (metis comp_def diff_self image_iff k)
      qed auto
    moreover
    have  $((\lambda x. f x - k * x)$  has_field_derivative deriv f x - k) (at x) for x
      by (metis DERIV_cmult_Id Deriv.field_differentiable_diff UNIV_I
field_differentiable_derivI holf holomorphic_on_def)
    ultimately
    show  $\forall x. ((\lambda x. f x - k * x)$  has_field_derivative 0) (at x)
      by auto
    show continuous_on UNIV  $(\lambda x. f x - k * x)$ 
      by (simp add: continuous_on_diff holf holomorphic_on_imp_continuous_on)
  qed (auto simp: connected_UNIV)
  have False
  proof (cases k=1)
    case True
      then have  $\exists x. k * x + l \neq a + x$  for a
        using l non [of a] ext [of f (+) a]
        by (metis add commute diff_eq_eq)
      with True show ?thesis by auto
    next
    case False
      have  $\bigwedge x. (1 - k) * x \neq f 0$ 
        using l [of 0]
        by (simp add: algebra_simps) (metis diff_add_cancel l mult.commute
non_fp)
      then show False
        by (metis False eq_iff_diff_eq_0 mult.commute nonzero_mult_div_cancel_right
times_divide_eq_right)
  qed
}

```

```

}
then show thesis
  using that by blast
qed

```

6.3 The Arzelà–Ascoli theorem

lemma *subsequence_diagonalization_lemma*:

```

fixes P :: nat ⇒ (nat ⇒ 'a) ⇒ bool
assumes sub:  $\bigwedge i r. \exists k. \text{strict\_mono } (k :: nat \Rightarrow nat) \wedge P i (r \circ k)$ 
  and P_P:  $\bigwedge i r :: nat \Rightarrow 'a. \bigwedge k1 k2 N. \llbracket P i (r \circ k1); \bigwedge j. N \leq j \implies \exists j'. j \leq j' \wedge k2 j = k1 j \rrbracket \implies P i (r \circ k2)$ 
obtains k where  $\text{strict\_mono } (k :: nat \Rightarrow nat) \wedge \bigwedge i. P i (r \circ k)$ 
proof –
  obtain kk where  $\bigwedge i r. \text{strict\_mono } (kk i r :: nat \Rightarrow nat) \wedge P i (r \circ (kk i r))$ 
    using sub by metis
  then have sub_kk:  $\bigwedge i r. \text{strict\_mono } (kk i r)$  and P_kk:  $\bigwedge i r. P i (r \circ (kk i r))$ 
    by auto
  define rr where  $rr \equiv \text{rec\_nat } (kk 0 r) (\lambda n x. x \circ kk (Suc n) (r \circ x))$ 
  then have [simp]:  $rr 0 = kk 0 r \wedge n. rr (Suc n) = rr n \circ kk (Suc n) (r \circ rr n)$ 
    by auto
  show thesis
  proof
    have sub_rr:  $\text{strict\_mono } (rr i)$  for i
      using sub_kk by (induction i) (auto simp: strict_mono_def o_def)
    have P_rr:  $P i (r \circ rr i)$  for i
      using P_kk by (induction i) (auto simp: o_def)
    have  $i \leq i+d \implies rr i n \leq rr (i+d) n$  for d i n
    proof (induction d)
      case 0 then show ?case
        by simp
      next
        case (Suc d) then show ?case
          using seq_suble [OF sub_kk] strict_mono_less_eq [OF sub_rr]
            by (simp add: order_subst1)
    qed
    then have  $\bigwedge i j n. i \leq j \implies rr i n \leq rr j n$ 
      by (metis le_iff_add)
    show  $\text{strict\_mono } (\lambda n. rr n n)$ 
      unfolding strict_mono_Suc_iff
        by (simp add: Suc_le_lessD strict_monoD strict_mono_imp_increasing sub_kk sub_rr)
    have  $\exists j. i \leq j \wedge rr (n+d) i = rr n j$  for d n i
    proof (induction d arbitrary: i)
      case (Suc d)
      then show ?case
        using seq_suble [OF sub_kk] by simp (meson order_trans)

```

```

qed auto
then have  $\bigwedge m n i. n \leq m \implies \exists j. i \leq j \wedge rr\ m\ i = rr\ n\ j$ 
  by (metis le_iff_add)
then show  $P\ i\ (r \circ (\lambda n. rr\ n\ n))$  for  $i$ 
  by (meson P_rr P_P)
qed
qed

lemma function_convergent_subsequence:
  fixes  $f :: [nat, 'a] \Rightarrow 'b :: \{real\_normed\_vector, heine\_borel\}$ 
  assumes countable  $S$  and  $M: \bigwedge n :: nat. \bigwedge x. x \in S \implies norm(f\ n\ x) \leq M$ 
  obtains  $k$  where strict_mono  $(k :: nat \Rightarrow nat)$   $\bigwedge x. x \in S \implies \exists l. (\lambda n. f\ (k\ n)\ x)$ 
   $\longrightarrow l$ 
proof (cases  $S = \{\}$ )
  case True
  then show ?thesis
    using strict_mono_id that by fastforce
next
  case False
  with  $\langle countable\ S \rangle$  obtain  $\sigma :: nat \Rightarrow 'a$  where  $\sigma: S = range\ \sigma$ 
    using uncountable_def by blast
  obtain  $k$  where strict_mono  $k$  and  $k: \bigwedge i. \exists l. (\lambda n. (f \circ k)\ n\ (\sigma\ i)) \longrightarrow l$ 
  proof (rule subsequence_diagonalization_lemma
    [of  $\lambda i r. \exists l. ((\lambda n. (f \circ r)\ n\ (\sigma\ i)) \longrightarrow l)$  sequentially id])
    show  $\exists k :: nat \Rightarrow nat. strict\_mono\ k \wedge (\exists l. (\lambda n. (f \circ (r \circ k))\ n\ (\sigma\ i)) \longrightarrow l)$ 
  for  $i\ r$ 
  proof -
    have  $f\ (r\ n)\ (\sigma\ i) \in cball\ 0\ M$  for  $n$ 
      by (simp add:  $\sigma\ M$ )
    then show ?thesis
      using compact_def [of  $cball\ (0 :: 'b)\ M$ ] by (force simp: o_def)
  qed
  show  $\exists l. (\lambda n. (f \circ (r \circ k2))\ n\ (\sigma\ i)) \longrightarrow l$ 
  if  $\exists l. (\lambda n. (f \circ (r \circ k1))\ n\ (\sigma\ i)) \longrightarrow l \wedge j. N \leq j \implies \exists j' \geq j. k2\ j = k1\ j'$ 
  for  $i\ N$  and  $r\ k1\ k2 :: nat \Rightarrow nat$ 
  using that
  by (simp add: lim_sequentially) (metis (no_types, opaque_lifting) le_cases
order_trans)
qed auto
with  $\sigma$  that show ?thesis
  by force
qed

```

theorem *Arzela_Ascoli*:

```

fixes  $\mathcal{F} :: [nat, 'a :: euclidean\_space] \Rightarrow 'b :: \{real\_normed\_vector, heine\_borel\}$ 
assumes compact  $S$ 
  and  $M: \bigwedge n x. x \in S \implies norm(\mathcal{F}\ n\ x) \leq M$ 
  and equicont:

```

```

       $\bigwedge x e. \llbracket x \in S; 0 < e \rrbracket$ 
       $\implies \exists d. 0 < d \wedge (\forall n y. y \in S \wedge \text{norm}(x - y) < d \longrightarrow \text{norm}(\mathcal{F} n x - \mathcal{F} n y) < e)$ 
    obtains  $g k$  where continuous_on S g strict_mono ( $k :: \text{nat} \Rightarrow \text{nat}$ )
       $\bigwedge e. 0 < e \implies \exists N. \forall n x. n \geq N \wedge x \in S \longrightarrow \text{norm}(\mathcal{F}(k n) x - g x) < e$ 
  proof -
    have UEQ:  $\bigwedge e. 0 < e \implies \exists d. 0 < d \wedge (\forall n. \forall x \in S. \forall x' \in S. \text{dist } x' x < d \longrightarrow \text{dist } (\mathcal{F} n x') (\mathcal{F} n x) < e)$ 
    apply (rule compact_uniformly_equicontinuous [OF  $\langle \text{compact } S \rangle$ , of range  $\mathcal{F}$ ])
    using equicont by (force simp: dist_commute dist_norm)+
    have continuous_on S g
      if  $\bigwedge e. 0 < e \implies \exists N. \forall n x. n \geq N \wedge x \in S \longrightarrow \text{norm}(\mathcal{F}(r n) x - g x) < e$ 
      for  $g:: 'a \Rightarrow 'b$  and  $r :: \text{nat} \Rightarrow \text{nat}$ 
    proof (rule uniform_limit_theorem [of  $\_ \mathcal{F} \circ r$ ])
      have continuous_on S ( $\mathcal{F} (r n)$ ) for  $n$ 
        using UEQ by (force simp: continuous_on_iff)
      then show  $\forall \_ n$  in sequentially. continuous_on S ( $(\mathcal{F} \circ r) n$ )
        by (simp add: eventually_sequentially)
      show uniform_limit S ( $\mathcal{F} \circ r$ )  $g$  sequentially
        using that by (metis (mono_tags, opaque_lifting) comp_apply dist_norm uniform_limit_sequentially_iff)
    qed auto
  moreover
    obtain  $R$  where countable  $R R \subseteq S$  and  $SR: S \subseteq \text{closure } R$ 
      by (metis separable that)
    obtain  $k$  where strict_mono k and  $k: \bigwedge x. x \in R \implies \exists l. (\lambda n. \mathcal{F} (k n) x) \longrightarrow l$ 
      using  $\langle R \subseteq S \rangle$  by (force intro: function_convergent_subsequence [OF  $\langle \text{countable } R \rangle M$ ])
    then have Cauchy: Cauchy  $((\lambda n. \mathcal{F} (k n) x))$  if  $x \in R$  for  $x$ 
      using convergent_eq_Cauchy that by blast
    have  $\exists N. \forall m n x. N \leq m \wedge N \leq n \wedge x \in S \longrightarrow \text{dist } ((\mathcal{F} \circ k) m x) ((\mathcal{F} \circ k) n x) < e$ 
      if  $0 < e$  for  $e$ 
    proof -
      obtain  $d$  where  $0 < d$ 
        and  $d: \bigwedge n. \forall x \in S. \forall x' \in S. \text{dist } x' x < d \longrightarrow \text{dist } (\mathcal{F} n x') (\mathcal{F} n x) < e/3$ 
        by (metis UEQ  $\langle 0 < e \rangle$  divide_pos_pos zero_less_numeral)
      obtain  $T$  where  $T \subseteq R$  and finite T and  $T: S \subseteq (\bigcup c \in T. \text{ball } c d)$ 
    proof (rule compactE_image [OF  $\langle \text{compact } S \rangle$ , of  $R (\lambda x. \text{ball } x d)$ ])
      have closure  $R \subseteq (\bigcup c \in R. \text{ball } c d)$ 
        using  $\langle 0 < d \rangle$  by (auto simp: closure_approachable)
      with  $SR$  show  $S \subseteq (\bigcup c \in R. \text{ball } c d)$ 
        by auto
    qed auto
    have  $\exists M. \forall m \geq M. \forall n \geq M. \text{dist } (\mathcal{F} (k m) x) (\mathcal{F} (k n) x) < e/3$  if  $x \in R$  for  $x$ 
      using Cauchy  $\langle 0 < e \rangle$  that unfolding Cauchy_def
      by (metis less_divide_eq_numeral1(1) mult_zero_left)

```



```

then obtain MF where MF:  $\bigwedge x m n. [x \in R; m \geq MF x; n \geq MF x] \implies$ 
norm (F (k m) x - F (k n) x) < e/3
using dist_norm by metis
have dist ((F o k) m x) ((F o k) n x) < e
if m: Max (MF ' T)  $\leq m$  and n: Max (MF ' T)  $\leq n$  x  $\in S$  for m n x
proof -
obtain t where t  $\in T$  and t: x  $\in$  ball t d
using x  $\in S$  T by auto
have norm(F (k m) t - F (k m) x) < e / 3
by (metis x R  $\subseteq$  S) x T  $\subseteq$  R) t  $\in T$ ) d dist_norm mem_ball subset_iff t x
 $\in S$ )
moreover
have norm(F (k n) t - F (k n) x) < e / 3
by (metis x R  $\subseteq$  S) x T  $\subseteq$  R) t  $\in T$ ) subsetD d dist_norm mem_ball t x
 $\in S$ )
moreover
have norm(F (k m) t - F (k n) t) < e / 3
proof (rule MF)
show t  $\in R$ 
using x T  $\subseteq$  R) t  $\in T$ ) by blast
show MF t  $\leq m$  MF t  $\leq n$ 
by (meson Max_ge x finite T) t  $\in T$ ) finite_imageI imageI le_trans m
n)+
qed
ultimately
show ?thesis
unfolding dist_norm [symmetric] o_def
by (metis dist_triangle_third dist_commute)
qed
then show ?thesis
by force
qed
then obtain g where  $\forall e > 0. \exists N. \forall n x. N \leq n \wedge x \in S \implies$  norm ((F o k) n
x - g x) < e
using uniformly_convergent_eq_cauchy [of  $\lambda x. x \in S$  F o k] by (auto simp
add: dist_norm)
ultimately show thesis
by (metis x strict_mono k) comp_apply that)
qed

```

6.3.1 Montel's theorem

a sequence of holomorphic functions uniformly bounded on compact subsets of an open set S has a subsequence that converges to a holomorphic function, and converges *uniformly* on compact subsets of S .

theorem Montel:

fixes $\mathcal{F} :: [\text{nat}, \text{complex}] \Rightarrow \text{complex}$

assumes open S

and $\mathcal{H}: \bigwedge h. h \in \mathcal{H} \implies h$ holomorphic_on S

and *bounded*: $\bigwedge K. \llbracket \text{compact } K; K \subseteq S \rrbracket \implies \exists B. \forall h \in \mathcal{H}. \forall z \in K. \text{norm}(h z) \leq B$
and *rng_f*: $\text{range } \mathcal{F} \subseteq \mathcal{H}$
obtains *g r*
where *g holomorphic_on S strict_mono* ($r :: \text{nat} \Rightarrow \text{nat}$)
 $\bigwedge x. x \in S \implies ((\lambda n. \mathcal{F} (r n) x) \longrightarrow g x)$ *sequentially*
 $\bigwedge K. \llbracket \text{compact } K; K \subseteq S \rrbracket \implies \text{uniform_limit } K (\mathcal{F} \circ r)$ *g sequentially*

proof –

obtain *K* **where** *comK*: $\bigwedge n. \text{compact}(K n)$ **and** *KS*: $\bigwedge n::\text{nat}. K n \subseteq S$
and *subK*: $\bigwedge X. \llbracket \text{compact } X; X \subseteq S \rrbracket \implies \exists N. \forall n \geq N. X \subseteq K n$
using *open_Union_compact_subsets* [*OF* $\langle \text{open } S \rangle$] **by** *metis*
then have $\bigwedge i. \exists B. \forall h \in \mathcal{H}. \forall z \in K i. \text{norm}(h z) \leq B$
by (*simp add: bounded*)
then obtain *B* **where** *B*: $\bigwedge i h z. \llbracket h \in \mathcal{H}; z \in K i \rrbracket \implies \text{norm}(h z) \leq B i$
by *metis*
have $*$: $\exists r g. \text{strict_mono} (r::\text{nat} \Rightarrow \text{nat}) \wedge (\forall e > 0. \exists N. \forall n \geq N. \forall x \in K i. \text{norm}((\mathcal{F} \circ r) n x - g x) < e)$
if $\bigwedge n. \mathcal{F} n \in \mathcal{H}$ **for** $\mathcal{F} i$

proof –

obtain *g k* **where** *continuous_on* ($K i$) *g strict_mono* ($k::\text{nat} \Rightarrow \text{nat}$)
 $\bigwedge e. 0 < e \implies \exists N. \forall n \geq N. \forall x \in K i. \text{norm}(\mathcal{F}(k n) x - g x) < e$

proof (*rule Arzela_Ascoli* [*of* $K i \mathcal{F} B i$])

show $\exists d > 0. \forall n y. y \in K i \wedge \text{cmod} (z - y) < d \longrightarrow \text{cmod} (\mathcal{F} n z - \mathcal{F} n y) < e$

if $z: z \in K i$ **and** $0 < e$ **for** $z e$

proof –

obtain *r* **where** $0 < r$ **and** *r*: $\text{cball } z r \subseteq S$

using *z KS* [*of* i] $\langle \text{open } S \rangle$ **by** (*force simp: open_contains_cball*)

have $\text{cball } z (2/3 * r) \subseteq \text{cball } z r$

using $\langle 0 < r \rangle$ **by** (*simp add: cball_subset_cball_iff*)

then have *z23S*: $\text{cball } z (2/3 * r) \subseteq S$

using *r* **by** *blast*

obtain *M* **where** $0 < M$ **and** *M*: $\bigwedge n w. \text{dist } z w \leq 2/3 * r \implies \text{norm}(\mathcal{F} n w) \leq M$

proof –

obtain *N* **where** $N: \forall n \geq N. \text{cball } z (2/3 * r) \subseteq K n$

using *subK compact_cball* [*of* $z (2/3 * r)$] *z23S* **by** *force*

have $\text{cmod} (\mathcal{F} n w) \leq |B N| + 1$ **if** $\text{dist } z w \leq 2/3 * r$ **for** $n w$

proof –

have $w \in K N$

using *N mem_cball that* **by** *blast*

then have $\text{cmod} (\mathcal{F} n w) \leq B N$

using *B* $\langle \bigwedge n. \mathcal{F} n \in \mathcal{H} \rangle$ **by** *blast*

also have $\dots \leq |B N| + 1$

by *simp*

finally show *?thesis* .

qed

then show *?thesis*

```

    by (rule_tac M=|B N| + 1 in that) auto
  qed
  have cmod (F n z - F n y) < e
    if y ∈ K i and y_near_z: cmod (z - y) < r/3 cmod (z - y) < e * r
  / (6 * M)
    for n y
  proof -
    have ((λw. F n w / (w - ξ)) has_contour_integral
      (2 * pi) * i * winding_number (circlepath z (2/3 * r)) ξ * F n ξ)
      (circlepath z (2/3 * r))
    if dist ξ z < (2/3 * r) for ξ
  proof (rule Cauchy_integral_formula_convex_simple)
    have F n holomorphic_on S
      by (simp add: H ⟨∧n. F n ∈ H⟩)
    with z23S show F n holomorphic_on cball z (2/3 * r)
      using holomorphic_on_subset by blast
  qed (use that ⟨0 < r⟩ in ⟨auto simp: dist_commute⟩)
  then have *: ((λw. F n w / (w - ξ)) has_contour_integral (2 * pi) * i
  * F n ξ)
    (circlepath z (2/3 * r))
    if dist ξ z < (2/3 * r) for ξ
    using that by (simp add: winding_number_circlepath dist_norm)
  have y: ((λw. F n w / (w - y)) has_contour_integral (2 * pi) * i * F
  n y)
    (circlepath z (2/3 * r))
  proof (rule *)
    show dist y z < 2/3 * r
      using that ⟨0 < r⟩ by (simp only: dist_norm norm_minus_commute)
  qed
  have z: ((λw. F n w / (w - z)) has_contour_integral (2 * pi) * i * F n
  z)
    (circlepath z (2/3 * r))
    using ⟨0 < r⟩ by (force intro!: *)
  have le_er: cmod (F n x / (x - y) - F n x / (x - z)) ≤ e / r
    if cmod (x - z) = r/3 + r/3 for x
  proof -
    have ¬ (cmod (x - y) < r/3)
      using y_near_z(1) that ⟨M > 0⟩ ⟨r > 0⟩
    by (metis (full_types) norm_diff_triangle_less norm_minus_commute
  order_less_irrefl)
    then have r4_le_xy: r/4 ≤ cmod (x - y)
      using ⟨r > 0⟩ by simp
    then have neq: x ≠ y x ≠ z
      using that ⟨r > 0⟩ by (auto simp: field_split_simps norm_minus_commute)
    have leM: cmod (F n x) ≤ M
      by (simp add: M dist_commute dist_norm that)
    have cmod (F n x / (x - y) - F n x / (x - z)) = cmod (F n x) *
  cmod (1 / (x - y) - 1 / (x - z))
      by (metis (no_types, lifting) divide_inverse mult.left_neutral norm_mult

```

```

right_diff_distrib')
  also have ... = cmod (F n x) * cmod ((y - z) / ((x - y) * (x - z)))
    using neq by (simp add: field_split_simps)
  also have ... = cmod (F n x) * (cmod (y - z) / (cmod(x - y) * (2/3
* r)))
    by (simp add: norm_mult norm_divide that)
  also have ... ≤ M * (cmod (y - z) / (cmod(x - y) * (2/3 * r)))
    using ⟨r > 0⟩ ⟨M > 0⟩ by (intro mult_mono [OF leM]) auto
  also have ... < M * ((e * r / (6 * M)) / (cmod(x - y) * (2/3 * r)))
    unfolding mult_less_cancel_left
    using y_near_z(2) ⟨M > 0⟩ ⟨r > 0⟩ neq
    by (simp add: field_simps mult_less_0_iff norm_minus_commute)
  also have ... ≤ e/r
    using ⟨e > 0⟩ ⟨r > 0⟩ r4_le_xy by (simp add: field_split_simps)
  finally show ?thesis by simp
qed
have (2 * pi) * cmod (F n y - F n z) = cmod ((2 * pi) * i * F n y -
(2 * pi) * i * F n z)
  by (simp add: right_diff_distrib [symmetric] norm_mult)
  also have cmod ((2 * pi) * i * F n y - (2 * pi) * i * F n z) ≤ e / r *
(2 * pi * (2/3 * r))

proof (rule has_contour_integral_bound_circlepath [OF has_contour_integral_diff
[OF y z]])
  show ∧x. cmod (x - z) = 2/3 * r ⇒ cmod (F n x / (x - y) - F n
x / (x - z)) ≤ e / r
    using le_er by auto
  qed (use ⟨e > 0⟩ ⟨r > 0⟩ in auto)
  also have ... = (2 * pi) * e * (2/3)
    using ⟨r > 0⟩ by (simp add: field_split_simps)
  finally have cmod (F n y - F n z) ≤ e * (2/3)
    by simp
  also have ... < e
    using ⟨e > 0⟩ by simp
  finally show ?thesis by (simp add: norm_minus_commute)
qed
then show ?thesis
  apply (rule_tac x=min (r/3) ((e * r)/(6 * M)) in exI)
  using ⟨0 < e⟩ ⟨0 < r⟩ ⟨0 < M⟩ by simp
qed
show ∧n x. x ∈ K i ⇒ cmod (F n x) ≤ B i
  using B ⟨∧n. F n ∈ H⟩ by blast
next
fix g :: complex ⇒ complex and k :: nat ⇒ nat
assume *: ∧(g::complex⇒complex) (k::nat⇒nat). continuous_on (K i) g ⇒
strict_mono k ⇒
(∧e. 0 < e ⇒ ∃ N. ∀ n ≥ N. ∀ x ∈ K i. cmod (F (k n) x - g x) <
e) ⇒ thesis
  continuous_on (K i) g

```

```

      strict_mono k
       $\bigwedge e. 0 < e \implies \exists N. \forall n x. N \leq n \wedge x \in K i \longrightarrow \text{cmod } (\mathcal{F} (k n) x - g$ 
x) < e
    show ?thesis
      by (rule *(1)[OF *(2,3)], drule *(4)) auto
    qed (use comK in simp_all)
  then show ?thesis
    by auto
  qed
  define  $\Phi$  where  $\Phi \equiv \lambda g i r. \lambda k::\text{nat} \Rightarrow \text{nat}. \forall e > 0. \exists N. \forall n \geq N. \forall x \in K i. \text{cmod}$ 
(( $\mathcal{F} \circ (r \circ k)$ ) n x - g x) < e
  obtain k :: nat  $\Rightarrow$  nat where strict_mono k and k:  $\bigwedge i. \exists g. \Phi g i id k$ 
  proof (rule subsequence_diagonalization_lemma [where r=id])
    show  $\exists g. \Phi g i id (r \circ k2)$ 
      if ex:  $\exists g. \Phi g i id (r \circ k1)$  and  $\bigwedge j. N \leq j \implies \exists j' \geq j. k2 j = k1 j'$ 
      for i k1 k2 N and r::nat $\Rightarrow$ nat
    proof -
      obtain g where  $\Phi g i id (r \circ k1)$ 
        using ex by blast
      then have  $\Phi g i id (r \circ k2)$ 
        using that
        by (simp add:  $\Phi\_def$ ) (metis (no_types, opaque_lifting) le_trans linear)
      then show ?thesis
        by metis
    qed
  have  $\exists k g. \text{strict\_mono } (k::\text{nat} \Rightarrow \text{nat}) \wedge \Phi g i id (r \circ k)$  for i r
    unfolding  $\Phi\_def$  o_assoc using rng_f by (force intro!: *)
  then show  $\bigwedge i r. \exists k. \text{strict\_mono } (k::\text{nat} \Rightarrow \text{nat}) \wedge (\exists g. \Phi g i id (r \circ k))$ 
    by force
  qed fastforce
  have  $\exists l. \forall e > 0. \exists N. \forall n \geq N. \text{norm}(\mathcal{F} (k n) z - l) < e$  if z  $\in$  S for z
  proof -
    obtain G where G:  $\bigwedge i e. e > 0 \implies \exists M. \forall n \geq M. \forall x \in K i. \text{cmod } ((\mathcal{F} \circ k) n$ 
x - G i x) < e
      using k unfolding  $\Phi\_def$  by (metis id_comp)
    obtain N where  $\bigwedge n. n \geq N \implies z \in K n$ 
      using subK [of {z}] that  $\langle z \in S \rangle$  by auto
    moreover have  $\bigwedge e. e > 0 \implies \exists M. \forall n \geq M. \forall x \in K N. \text{cmod } ((\mathcal{F} \circ k) n x -$ 
G N x) < e
      using G by auto
    ultimately show ?thesis
      by (metis comp_apply order_refl)
  qed
  then obtain g where g:  $\bigwedge z e. \llbracket z \in S; e > 0 \rrbracket \implies \exists N. \forall n \geq N. \text{norm}(\mathcal{F} (k n)$ 
z - g z) < e
    by metis
  show ?thesis
  proof
    show g_lim:  $\bigwedge x. x \in S \implies (\lambda n. \mathcal{F} (k n) x) \longrightarrow g x$ 

```

```

    by (simp add: lim_sequentially g dist_norm)
  have dg_le_e:  $\exists N. \forall n \geq N. \forall x \in T. cmod (\mathcal{F} (k n) x - g x) < e$ 
  if T: compact T  $T \subseteq S$  and  $0 < e$  for T e
  proof -
    obtain N where N:  $\bigwedge n. n \geq N \implies T \subseteq K n$ 
      using subK [OF T] by blast
    obtain h where h:  $\bigwedge e. e > 0 \implies \exists M. \forall n \geq M. \forall x \in K N. cmod ((\mathcal{F} \circ k) n x - h x) < e$ 
      using k unfolding  $\Phi\_def$  by (metis id_comp)
    have geq:  $g w = h w$  if  $w \in T$  for w
    proof (rule LIMSEQ_unique)
      show  $(\lambda n. \mathcal{F} (k n) w) \longrightarrow g w$ 
        using  $\langle T \subseteq S \rangle g\_lim$  that by blast
      show  $(\lambda n. \mathcal{F} (k n) w) \longrightarrow h w$ 
        using h N that by (force simp: lim_sequentially dist_norm)
    qed
    show ?thesis
      using T h N  $\langle 0 < e \rangle$  by (fastforce simp add: geq)
  qed
  then show  $\bigwedge K. \llbracket compact K; K \subseteq S \rrbracket \implies uniform\_limit K (\mathcal{F} \circ k) g$  sequentially
    by (simp add: uniform_limit_iff dist_norm eventually_sequentially)
  show g holomorphic_on S
  proof (rule holomorphic_uniform_sequence [OF  $\langle open S \rangle \mathcal{H}$ ])
    show  $\bigwedge n. (\mathcal{F} \circ k) n \in \mathcal{H}$ 
      by (simp add: range_subsetD rng_f)
    show  $\exists d > 0. cball z d \subseteq S \wedge uniform\_limit (cball z d) (\lambda n. (\mathcal{F} \circ k) n) g$  sequentially
      if  $z \in S$  for z
    proof -
      obtain d where d:  $d > 0$   $cball z d \subseteq S$ 
        using  $\langle open S \rangle \langle z \in S \rangle open\_contains\_cball$  by blast
      then have uniform_limit (cball z d)  $(\mathcal{F} \circ k) g$  sequentially
        using dg_le_e compact_cball by (auto simp: uniform_limit_iff eventually_sequentially dist_norm)
      with d show ?thesis by blast
    qed
  qed
  qed (auto simp:  $\langle strict\_mono k \rangle$ )
  qed

```

6.4 Some simple but useful cases of Hurwitz's theorem

proposition Hurwitz_no_zeros:

assumes S : open S connected S

and holf: $\bigwedge n::nat. \mathcal{F} n$ holomorphic_on S

and holg: g holomorphic_on S

and ul_g: $\bigwedge K. \llbracket compact K; K \subseteq S \rrbracket \implies uniform_limit K \mathcal{F} g$ sequentially

and nonconst: \neg g constant_on S

```

    and nz:  $\bigwedge n z. z \in S \implies \mathcal{F} n z \neq 0$ 
    and z0  $\in S$ 
    shows  $g z0 \neq 0$ 
  proof
    assume g0:  $g z0 = 0$ 
    obtain h r m
      where  $0 < m 0 < r$  and subS:  $\text{ball } z0 \ r \subseteq S$ 
      and holh:  $h \text{ holomorphic\_on } \text{ball } z0 \ r$ 
      and geq:  $\bigwedge w. w \in \text{ball } z0 \ r \implies g w = (w - z0)^{\wedge m} * h w$ 
      and hnz:  $\bigwedge w. w \in \text{ball } z0 \ r \implies h w \neq 0$ 
    by (blast intro: holomorphic_factor_zero_nonconstant [OF holg S  $\langle z0 \in S \rangle$  g0
    nonconst])
    then have holf0:  $\mathcal{F} n \text{ holomorphic\_on } \text{ball } z0 \ r$  for  $n$ 
      by (meson holf holomorphic_on_subset)
    have *: (( $\lambda z. \text{deriv } (\mathcal{F} n) z / \mathcal{F} n z$ ) has_contour_integral 0) (circlepath z0
    (r/2)) for  $n$ 
      proof (rule Cauchy_theorem_disc_simple)
        show ( $\lambda z. \text{deriv } (\mathcal{F} n) z / \mathcal{F} n z$ ) holomorphic_on ball z0 r
          by (metis (no_types)  $\langle \text{open } S \rangle$  holf holomorphic_deriv holomorphic_on_divide
          holomorphic_on_subset nz subS)
        qed (use  $\langle 0 < r \rangle$  in auto)
      have hol_dg:  $\text{deriv } g \text{ holomorphic\_on } S$ 
        by (simp add:  $\langle \text{open } S \rangle$  holg holomorphic_deriv)
      have continuous_on (sphere z0 (r/2)) (deriv g)
        using  $\langle 0 < r \rangle$  subS
      by (intro holomorphic_on_imp_continuous_on holomorphic_on_subset [OF
      hol_dg]) auto
      then have compact (deriv g ' (sphere z0 (r/2)))
        by (rule compact_continuous_image [OF _ compact_sphere])
      then have bo_dg: bounded (deriv g ' (sphere z0 (r/2)))
        using compact_imp_bounded by blast
      have continuous_on (sphere z0 (r/2)) (cmod  $\circ$  g)
        using  $\langle 0 < r \rangle$  subS
      by (intro continuous_intros holomorphic_on_imp_continuous_on holomor-
      phic_on_subset [OF holg]) auto
      then have compact ((cmod  $\circ$  g) ' sphere z0 (r/2))
        by (rule compact_continuous_image [OF _ compact_sphere])
      moreover have (cmod  $\circ$  g) ' sphere z0 (r/2)  $\neq \{\}$ 
        using  $\langle 0 < r \rangle$  by auto
      ultimately obtain b where  $b \in (\text{cmod } \circ \ g) \text{ ' sphere } z0 \ (r/2)$ 
         $\bigwedge t. t \in (\text{cmod } \circ \ g) \text{ ' sphere } z0 \ (r/2) \implies b \leq t$ 
      using compact_attains_inf [of (norm  $\circ$  g) ' (sphere z0 (r/2))] by blast
      have ( $\lambda n. \text{contour\_integral } (\text{circlepath } z0 \ (r/2)) (\lambda z. \text{deriv } (\mathcal{F} n) z / \mathcal{F} n z)$ )
       $\longrightarrow$ 
        contour_integral (circlepath z0 (r/2)) ( $\lambda z. \text{deriv } g z / g z$ )
      proof (rule contour_integral_uniform_limit_circlepath)
        show  $\forall_F n$  in sequentially. ( $\lambda z. \text{deriv } (\mathcal{F} n) z / \mathcal{F} n z$ ) contour_integrable_on
        circlepath z0 (r/2)
          using * contour_integrable_on_def eventually_sequentiallyI by meson

```

```

show uniform_limit (sphere z0 (r/2)) ( $\lambda n z. \text{deriv } (\mathcal{F} n) z / \mathcal{F} n z$ ) ( $\lambda z. \text{deriv } g z / g z$ ) sequentially
proof (rule uniform_lim_divide [OF __ bo_dg])
show uniform_limit (sphere z0 (r/2)) ( $\lambda a. \text{deriv } (\mathcal{F} a)$ ) (deriv g) sequentially
proof (rule uniform_limitI)
  fix e::real
  assume 0 < e

  show  $\forall_F n$  in sequentially.  $\forall x \in \text{sphere } z0 (r/2). \text{dist } (\text{deriv } (\mathcal{F} n) x)$ 
  (deriv g x) < e
  proof -
    have dist (deriv ( $\mathcal{F} n$ ) w) (deriv g w) < e
      if e8:  $\bigwedge x. \text{dist } z0 x \leq 3 * r / 4 \implies \text{dist } (\mathcal{F} n x) (g x) * 8 < r * e$ 
      and w: w  $\in$  sphere z0 (r/2) for n w
    proof -
      have ball w (r/4)  $\subseteq$  ball z0 r cball w (r/4)  $\subseteq$  ball z0 r
      using <0 < r> w by (simp_all add: ball_subset_ball_iff cball_subset_ball_iff
  dist_commute)
      with subS have wr4_sub: ball w (r/4)  $\subseteq$  S cball w (r/4)  $\subseteq$  S by force+
      moreover
      have ( $\lambda z. \mathcal{F} n z - g z$ ) holomorphic_on S
        by (intro holomorphic_intros holf holg)
      ultimately have hol: ( $\lambda z. \mathcal{F} n z - g z$ ) holomorphic_on ball w (r/4)
        and cont: continuous_on (cball w (r / 4)) ( $\lambda z. \mathcal{F} n z - g z$ )
      using holomorphic_on_subset by (blast intro: holomorphic_on_imp_continuous_on)+
      have w  $\in$  S
        using <0 < r> wr4_sub by auto
      have dist z0 y  $\leq 3 * r / 4$  if dist w y < r/4 for y
      proof (rule dist_triangle_le [where z=w])
        show dist z0 w + dist y w  $\leq 3 * r / 4$ 
        using w that by (simp add: dist_commute)
      qed
      with e8 have in_ball:  $\bigwedge y. y \in \text{ball } w (r/4) \implies \mathcal{F} n y - g y \in \text{ball } 0$ 
  (r/4 * e/2)
        by (simp add: dist_norm [symmetric])
      have  $\mathcal{F} n$  field_differentiable at w
        by (metis holomorphic_on_imp_differentiable_at <w  $\in$  S> holf <open
  S>)
      moreover
      have g field_differentiable at w
        using <w  $\in$  S> <open S> holg holomorphic_on_imp_differentiable_at
by auto
      moreover
      have cmod (deriv ( $\lambda w. \mathcal{F} n w - g w$ ) w) * 2  $\leq e$ 
        using Cauchy_higher_deriv_bound [OF hol cont in_ball, of 1] <r >
  0> by auto
      ultimately have dist (deriv ( $\mathcal{F} n$ ) w) (deriv g w)  $\leq e/2$ 
        by (simp add: dist_norm)
      then show ?thesis

```



```

      using ‹e > 0› by auto
    qed
  moreover
  have cball z0 (3 * r / 4) ⊆ ball z0 r
    by (simp add: cball_subset_ball_iff ‹0 < r›)
  with subS have uniform_limit (cball z0 (3 * r/4))  $\mathcal{F}$  g sequentially
    by (force intro: ul_g)
  then have  $\forall_F n$  in sequentially.  $\forall x \in \text{cball } z0 (3 * r / 4)$ .  $\text{dist } (\mathcal{F} n x) (g$ 
 $x) < r / 4 * e / 2$ 
    using ‹0 < e› ‹0 < r› by (force simp: intro!: uniform_limitD)
  ultimately show ?thesis
    by (force simp add: eventually_sequentially)
  qed
  qed
  show uniform_limit (sphere z0 (r/2))  $\mathcal{F}$  g sequentially
  proof (rule uniform_limitI)
    fix e::real
    assume 0 < e
    have sphere z0 (r/2) ⊆ ball z0 r
      using ‹0 < r› by auto
    with subS have uniform_limit (sphere z0 (r/2))  $\mathcal{F}$  g sequentially
      by (force intro: ul_g)
    then show  $\forall_F n$  in sequentially.  $\forall x \in \text{sphere } z0 (r/2)$ .  $\text{dist } (\mathcal{F} n x) (g x)$ 
    < e
      using ‹0 < e› uniform_limit_iff by blast
    qed
    show  $b > 0 \wedge x. x \in \text{sphere } z0 (r/2) \implies b \leq \text{cmod } (g x)$ 
      using  $b < 0 < r$  by (fastforce simp: geq hnz)+
    qed
  qed (use ‹0 < r› in auto)
  then have  $(\lambda n. 0) \longrightarrow \text{contour\_integral } (\text{circlepath } z0 (r/2)) (\lambda z. \text{deriv } g z$ 
 $/ g z)$ 
    by (simp add: contour_integral_unique [OF *])
  then have  $\text{contour\_integral } (\text{circlepath } z0 (r/2)) (\lambda z. \text{deriv } g z / g z) = 0$ 
    by (simp add: LIMSEQ_const_iff)
  moreover
  have  $\text{contour\_integral } (\text{circlepath } z0 (r/2)) (\lambda z. \text{deriv } g z / g z) =$ 
 $\text{contour\_integral } (\text{circlepath } z0 (r/2)) (\lambda z. m / (z - z0) + \text{deriv } h z / h z)$ 
  proof (rule contour_integral_eq, use ‹0 < r› in simp)
    fix w
    assume w:  $\text{dist } z0 w * 2 = r$ 
    then have w_inb:  $w \in \text{ball } z0 r$ 
      using ‹0 < r› by auto
    have h_der: (h has_field_derivative deriv h w) (at w)
      using holh holomorphic_derivI w_inb by blast
    have  $\text{deriv } g w = ((\text{of\_nat } m * h w + \text{deriv } h w * (w - z0)) * (w - z0) ^ m)$ 
 $/ (w - z0)$ 
      if  $r = \text{dist } z0 w * 2$   $w \neq z0$ 
    proof -

```

```

have (( $\lambda w. (w - z0) ^ m * h w$ ) has_field_derivative
  ( $m * h w + \text{deriv } h w * (w - z0)$ ) * ( $w - z0$ ) ^  $m / (w - z0)$ ) (at w)
apply (rule derivative_eq_intros h_der refl)+
using that  $\langle m > 0 \rangle \langle 0 < r \rangle$  apply (simp add: divide_simps distrib_right)
by (metis Suc_pred mult.commute power_Suc)
then show ?thesis
proof (rule DERIV_imp_deriv [OF has_field_derivative_transform_within_open])
  show  $\bigwedge x. x \in \text{ball } z0 r \implies (x - z0) ^ m * h x = g x$ 
    by (simp add: hnz geq)
  qed (use that  $\langle m > 0 \rangle \langle 0 < r \rangle$  in auto)
qed
with  $\langle 0 < r \rangle \langle 0 < m \rangle w w\_inb$  show  $\text{deriv } g w / g w = \text{of\_nat } m / (w - z0)$ 
+  $\text{deriv } h w / h w$ 
  by (auto simp: geq field_split_simps hnz)
qed
moreover
have contour_integral (circlepath z0 (r/2)) ( $\lambda z. m / (z - z0) + \text{deriv } h z / h$ 
 $z$ ) =
   $2 * \text{of\_real } \pi * i * m + 0$ 
proof (rule contour_integral_unique [OF has_contour_integral_add])
  show ( $\lambda x. m / (x - z0)$ ) has_contour_integral  $2 * \text{of\_real } \pi * i * m$ 
(circlepath z0 (r/2))
  by (force simp:  $\langle 0 < r \rangle$  intro: Cauchy_integral_circlepath_simple)
  show ( $\lambda x. \text{deriv } h x / h x$ ) has_contour_integral 0 (circlepath z0 (r/2))
  using hnz holh holomorphic_deriv holomorphic_on_divide  $\langle 0 < r \rangle$ 
  by (fastforce intro!: Cauchy_theorem_disc_simple [of _ z0 r])
qed
ultimately show False using  $\langle 0 < m \rangle$  by auto
qed

corollary Hurwitz_injective:
assumes S: open S connected S
and holf:  $\bigwedge n::\text{nat}. \mathcal{F} n$  holomorphic_on S
and holg: g holomorphic_on S
and ul_g:  $\bigwedge K. \llbracket \text{compact } K; K \subseteq S \rrbracket \implies \text{uniform\_limit } K \mathcal{F} g$  sequentially
and nonconst:  $\neg g$  constant_on S
and inj:  $\bigwedge n. \text{inj\_on } (\mathcal{F} n) S$ 
shows inj_on g S
proof -
have False if z12: z1 ∈ S z2 ∈ S z1 ≠ z2 g z2 = g z1 for z1 z2
proof -
  obtain z0 where z0 ∈ S and z0: g z0 ≠ g z2
  using constant_on_def nonconst by blast
  have ( $\lambda z. g z - g z1$ ) holomorphic_on S
  by (intro holomorphic_intros holg)
  then obtain r where 0 < r ball z2 r ⊆ S  $\bigwedge z. \text{dist } z2 z < r \wedge z \neq z2 \implies g$ 
 $z \neq g z1$ 
  using isolated_zeros [of  $\lambda z. g z - g z1 S z2 z0$ ] S  $\langle z0 \in S \rangle z0 z12$  by auto
  have  $g z2 - g z1 \neq 0$ 

```

```

proof (rule Hurwitz_no_zeros [of  $S - \{z1\}$   $\lambda n z. \mathcal{F} n z - \mathcal{F} n z1$   $\lambda z. g z - g z1$ ])
  show open  $(S - \{z1\})$ 
    by (simp add: S open_delete)
  show connected  $(S - \{z1\})$ 
    by (simp add: connected_open_delete [OF S])
  show  $\bigwedge n. (\lambda z. \mathcal{F} n z - \mathcal{F} n z1)$  holomorphic_on  $S - \{z1\}$ 
    by (intro holomorphic_intros holomorphic_on_subset [OF holg]) blast
  show  $(\lambda z. g z - g z1)$  holomorphic_on  $S - \{z1\}$ 
    by (intro holomorphic_intros holomorphic_on_subset [OF holg]) blast
  show uniform_limit  $K$   $(\lambda n z. \mathcal{F} n z - \mathcal{F} n z1)$   $(\lambda z. g z - g z1)$  sequentially
    if compact  $K$   $K \subseteq S - \{z1\}$  for  $K$ 
  proof (rule uniform_limitI)
    fix  $e::\text{real}$ 
    assume  $e > 0$ 
    have uniform_limit  $K$   $\mathcal{F}$   $g$  sequentially
      using that ul_g by fastforce
    then have  $K: \forall_F n$  in sequentially.  $\forall x \in K. \text{dist} (\mathcal{F} n x) (g x) < e/2$ 
      using  $\langle 0 < e \rangle$  by (force simp: intro!: uniform_limitD)
    have uniform_limit  $\{z1\}$   $\mathcal{F}$   $g$  sequentially
      by (intro ul_g) (auto simp: z12)
    then have  $\forall_F n$  in sequentially.  $\forall x \in \{z1\}. \text{dist} (\mathcal{F} n x) (g x) < e/2$ 
      using  $\langle 0 < e \rangle$  by (force simp: intro!: uniform_limitD)
    then have  $z1: \forall_F n$  in sequentially.  $\text{dist} (\mathcal{F} n z1) (g z1) < e/2$ 
      by simp
    show  $\forall_F n$  in sequentially.  $\forall x \in K. \text{dist} (\mathcal{F} n x - \mathcal{F} n z1) (g x - g z1) < e$ 
      apply (intro eventually_mono [OF eventually_conj [OF K z1]])
    by (metis (no_types, opaque_lifting) diff_add_eq diff_diff_eq2 dist_commute
      dist_norm dist_triangle_add_half)
    qed
    show  $\neg (\lambda z. g z - g z1)$  constant_on  $S - \{z1\}$ 
      unfolding constant_on_def
      by (metis Diff_iff  $\langle z0 \in S \rangle$  empty_iff insert_iff right_minus_eq  $z0 z12$ )
    show  $\bigwedge n z. z \in S - \{z1\} \implies \mathcal{F} n z - \mathcal{F} n z1 \neq 0$ 
      by (metis DiffD1 DiffD2 eq_iff_diff_eq_0 inj inj_onD insertI1  $\langle z1 \in S \rangle$ )
    show  $z2 \in S - \{z1\}$ 
      using  $\langle z2 \in S \rangle$   $\langle z1 \neq z2 \rangle$  by auto
    qed
  with z12 show False by auto
  qed
  then show ?thesis by (auto simp: inj_on_def)
qed

```

6.5 The Great Picard theorem

lemma GPicard1:

```

assumes  $S$ : open  $S$  connected  $S$  and  $w \in S$   $0 < r$   $Y \subseteq X$ 
and holX:  $\bigwedge h. h \in X \implies h$  holomorphic_on  $S$ 
and X01:  $\bigwedge h z. \llbracket h \in X; z \in S \rrbracket \implies h z \neq 0 \wedge h z \neq 1$ 

```

```

    and r:  $\bigwedge h. h \in Y \implies \text{norm}(h w) \leq r$ 
  obtains B Z where  $0 < B$  open Z  $w \in Z$   $Z \subseteq S$   $\bigwedge h z. [h \in Y; z \in Z] \implies$ 
 $\text{norm}(h z) \leq B$ 
  proof -
    obtain e where  $e > 0$  and e:  $\text{cball } w \ e \subseteq S$ 
      using assms open_contains_cball_eq by blast
    show ?thesis
  proof
    show  $0 < \exp(\pi * \exp(\pi * (2 + 2 * r + 12)))$ 
      by simp
    show  $\text{ball } w \ (e / 2) \subseteq S$ 
      using e ball_divide_subset_numeral ball_subset_cball by blast
    show  $\text{cmod } (h z) \leq \exp(\pi * \exp(\pi * (2 + 2 * r + 12)))$ 
      if  $h \in Y$   $z \in \text{ball } w \ (e / 2)$  for h z
  proof -
    have  $h \in X$ 
      using  $\langle Y \subseteq X \rangle \langle h \in Y \rangle$  by blast
    have hol_h_o:  $(h \circ (\lambda z. (w + \text{of\_real } e * z)))$  holomorphic_on cball 0 1
  proof (intro holomorphic_intros holomorphic_on_compose)
    have h holomorphic_on S
      using holX  $\langle h \in X \rangle$  by auto
    then have h holomorphic_on cball w e
      by (metis e holomorphic_on_subset)
    moreover have  $(\lambda z. w + \text{complex\_of\_real } e * z) \text{ ` cball } 0 \ 1 \subseteq \text{cball } w \ e$ 
      using that  $\langle e > 0 \rangle$  by (auto simp: dist_norm norm_mult)
    ultimately show h holomorphic_on  $(\lambda z. w + \text{complex\_of\_real } e * z) \text{ ` cball } 0 \ 1$ 
      by (rule holomorphic_on_subset)
  qed
  have norm_le_r:  $\text{cmod } ((h \circ (\lambda z. w + \text{complex\_of\_real } e * z)) 0) \leq r$ 
    by (auto simp: r  $\langle h \in Y \rangle$ )
  have le12:  $\text{norm } (\text{of\_real } (\text{inverse } e) * (z - w)) \leq 1/2$ 
  using that  $\langle e > 0 \rangle$  by (simp add: inverse_eq_divide dist_norm norm_minus_commute
norm_divide)
  have non01:  $h(w + e * z) \neq 0 \wedge h(w + e * z) \neq 1$  if  $z \in \text{cball } 0 \ 1$  for
 $z::\text{complex}$ 
  proof (rule X01 [OF  $\langle h \in X \rangle$ ])
    have  $w + \text{complex\_of\_real } e * z \in \text{cball } w \ e$ 
      using  $\langle 0 < e \rangle$  that by (auto simp: dist_norm norm_mult)
    then show  $w + \text{complex\_of\_real } e * z \in S$ 
      by (rule subsetD [OF e])
  qed
  have  $\text{cmod } (h z) \leq \text{cmod } (h(w + \text{of\_real } e * (\text{inverse } e * (z - w))))$ 
    using  $\langle 0 < e \rangle$  by (simp add: field_split_simps)
  also have  $\dots \leq \exp(\pi * \exp(\pi * (14 + 2 * r)))$ 
  using r [OF  $\langle h \in Y \rangle$ ] Schottky [OF hol_h_o norm_le_r le12] non01
  by auto
  finally
  show ?thesis by simp

```

qed
qed (use $\langle e > 0 \rangle$ in auto)
qed

lemma *GPicard2*:

assumes $S \subseteq T$ *connected* $T \ S \neq \{\}$ *open* $S \wedge x. \llbracket x \text{ islimpt } S; x \in T \rrbracket \implies x \in S$
shows $S = T$
by (*metis* *assms* *open_subset* *connected_clopen* *closedin_limpt*)

lemma *GPicard3*:

assumes S : *open* S *connected* $S \ w \in S$ **and** $Y \subseteq X$
and *holX*: $\bigwedge h. h \in X \implies h \text{ holomorphic_on } S$
and *X01*: $\bigwedge h \ z. \llbracket h \in X; z \in S \rrbracket \implies h \ z \neq 0 \wedge h \ z \neq 1$
and *no_hw_le1*: $\bigwedge h. h \in Y \implies \text{norm}(h \ w) \leq 1$
and *compact* $K \ K \subseteq S$
obtains B **where** $\bigwedge h \ z. \llbracket h \in Y; z \in K \rrbracket \implies \text{norm}(h \ z) \leq B$

proof –

define U **where** $U \equiv \{z \in S. \exists B \ Z. 0 < B \wedge \text{open } Z \wedge z \in Z \wedge Z \subseteq S \wedge$
 $(\forall h \ z'. h \in Y \wedge z' \in Z \longrightarrow \text{norm}(h \ z') \leq B)\}$

then have $U \subseteq S$ **by** *blast*

have $U = S$

proof (*rule* *GPicard2* [*OF* $\langle U \subseteq S \rangle$ $\langle \text{connected } S \rangle$])

show $U \neq \{\}$

proof –

obtain $B \ Z$ **where** $0 < B$ *open* $Z \ w \in Z \ Z \subseteq S$

and $\bigwedge h \ z. \llbracket h \in Y; z \in Z \rrbracket \implies \text{norm}(h \ z) \leq B$

using *GPicard1* [*OF* S *zero_less_one* $\langle Y \subseteq X \rangle$ *holX*] *X01* *no_hw_le1* **by**

blast

then show *?thesis*

unfolding U_def **using** $\langle w \in S \rangle$ **by** *blast*

qed

show *open* U

unfolding *open_subopen* [*of* U] **by** (*auto* *simp*: U_def)

fix v

assume v : $v \text{ islimpt } U \ v \in S$

have $\neg (\forall r > 0. \exists h \in Y. r < \text{cmod}(h \ v))$

proof

assume $\forall r > 0. \exists h \in Y. r < \text{cmod}(h \ v)$

then have $\forall n. \exists h \in Y. \text{Suc } n < \text{cmod}(h \ v)$

by *simp*

then obtain \mathcal{F} **where** FY : $\bigwedge n. \mathcal{F} \ n \in Y$ **and** ltF : $\bigwedge n. \text{Suc } n < \text{cmod}(\mathcal{F}$

$n \ v)$

by *metis*

define \mathcal{G} **where** $\mathcal{G} \equiv \lambda n \ z. \text{inverse}(\mathcal{F} \ n \ z)$

have *holG*: $\mathcal{G} \ n \text{ holomorphic_on } S$ **for** n

proof (*simp* *add*: \mathcal{G}_def)

show $(\lambda z. \text{inverse}(\mathcal{F} \ n \ z)) \text{ holomorphic_on } S$

using $FY \ X01 \ \langle Y \subseteq X \rangle \text{ holX}$ **by** (*blast* *intro*: *holomorphic_on_inverse*)

```

qed
have  $\mathcal{G}_{not0}$ :  $\mathcal{G} \ n \ z \neq 0$  and  $\mathcal{G}_{not1}$ :  $\mathcal{G} \ n \ z \neq 1$  if  $z \in S$  for  $n \ z$ 
  using  $FY \ X01 \ \langle Y \subseteq X \rangle$  that by (force simp:  $\mathcal{G}_{def}$ )+
have  $\mathcal{G}_{le1}$ :  $cmod \ (\mathcal{G} \ n \ v) \leq 1$  for  $n$ 
  using less_le_trans linear ltF
  by (fastforce simp add:  $\mathcal{G}_{def}$  norm_inverse inverse_le_1_iff)
define  $W$  where  $W \equiv \{h. h \text{ holomorphic\_on } S \wedge (\forall z \in S. h \ z \neq 0 \wedge h \ z \neq 1)\}$ 
obtain  $B \ Z$  where  $0 < B$  open  $Z \ v \in Z \ Z \subseteq S$ 
  and  $B$ :  $\bigwedge h \ z. \llbracket h \in \text{range } \mathcal{G}; z \in Z \rrbracket \implies \text{norm}(h \ z) \leq B$ 
  apply (rule  $GPicard1$  [ $OF \ \langle \text{open } S \rangle \ \langle \text{connected } S \rangle \ \langle v \in S \rangle \ \text{zero\_less\_one,}$ 
of range  $\mathcal{G} \ W$ ])
  using  $hol\mathcal{G} \ \mathcal{G}_{not0} \ \mathcal{G}_{not1} \ \mathcal{G}_{le1}$  by (force simp:  $W_{def}$ )+
then obtain  $e$  where  $e > 0$  and  $e$ :  $\text{ball } v \ e \subseteq Z$ 
  by (meson open_contains_ball)
obtain  $h \ j$  where  $holh$ :  $h$  holomorphic_on ball  $v \ e$  and strict_mono  $j$ 
  and  $lim$ :  $\bigwedge x. x \in \text{ball } v \ e \implies (\lambda n. \mathcal{G} \ (j \ n) \ x) \longrightarrow h \ x$ 
  and  $ulim$ :  $\bigwedge K. \llbracket \text{compact } K; K \subseteq \text{ball } v \ e \rrbracket \implies \text{uniform\_limit } K \ (\mathcal{G} \circ j) \ h$  sequentially
proof (rule Montel)
  show  $\bigwedge h. h \in \text{range } \mathcal{G} \implies h$  holomorphic_on ball  $v \ e$ 
    by (metis  $\langle Z \subseteq S \rangle \ e \ hol\mathcal{G} \ \text{holomorphic\_on\_subset\_image}E$ )
  show  $\bigwedge K. \llbracket \text{compact } K; K \subseteq \text{ball } v \ e \rrbracket \implies \exists B. \forall h \in \text{range } \mathcal{G}. \forall z \in K. cmod$ 
( $h \ z$ )  $\leq B$ 
    using  $B \ e$  by blast
qed auto
have  $h \ v = 0$ 
proof (rule LIMSEQ_unique)
  show  $(\lambda n. \mathcal{G} \ (j \ n) \ v) \longrightarrow h \ v$ 
    using  $\langle e > 0 \rangle \ lim$  by simp
  have  $lt\_Fj$ :  $\text{real } x \leq cmod \ (\mathcal{F} \ (j \ x) \ v)$  for  $x$ 
    by (metis of_nat_Suc ltF  $\langle \text{strict\_mono } j \rangle \ \text{add.commute less\_eq\_real\_def}$ 
less_le_trans nat_le_real_less seq_suble)
  show  $(\lambda n. \mathcal{G} \ (j \ n) \ v) \longrightarrow 0$ 
proof (rule Lim_null_comparison [ $OF \ \text{eventually\_sequentiallyI} \ lim\_inverse\_n$ ])
  show  $cmod \ (\mathcal{G} \ (j \ x) \ v) \leq \text{inverse} \ (\text{real } x)$  if  $1 \leq x$  for  $x$ 
    using that by (simp add:  $\mathcal{G}_{def}$  norm_inverse_le_norm [ $OF \ lt\_Fj$ ])
qed
qed
have  $h \ v \neq 0$ 
proof (rule Hurwitz_no_zeros [of ball  $v \ e \ \mathcal{G} \circ j \ h$ ])
  show  $\bigwedge n. (\mathcal{G} \circ j) \ n$  holomorphic_on ball  $v \ e$ 
    using  $\langle Z \subseteq S \rangle \ e \ hol\mathcal{G}$  by force
  show  $\bigwedge n \ z. z \in \text{ball } v \ e \implies (\mathcal{G} \circ j) \ n \ z \neq 0$ 
    using  $\mathcal{G}_{not0} \ \langle Z \subseteq S \rangle \ e$  by fastforce
  show  $\neg h$  constant_on ball  $v \ e$ 
proof (clarsimp simp: constant_on_def)
  fix  $c$ 
  have False if  $\bigwedge z. \text{dist } v \ z < e \implies h \ z = c$ 

```

```

proof –
  have  $h\ v = c$ 
    by (simp add: <0 < e> that)
  obtain  $y$  where  $y \in U\ y \neq v$  and  $y: dist\ y\ v < e$ 
    using  $v\ \langle e > 0$  by (auto simp: islimpt_approachable)
  then obtain  $C\ T$  where  $y \in S\ C > 0$  open  $T\ y \in T\ T \subseteq S$ 
    and  $\bigwedge h\ z'. \llbracket h \in Y; z' \in T \rrbracket \implies cmod\ (h\ z') \leq C$ 
    using  $\langle y \in U \rangle$  by (auto simp: U_def)
  then have  $le\_C: \bigwedge n. cmod\ (\mathcal{F}\ n\ y) \leq C$ 
    using FY by blast

  have  $\forall_F\ n$  in sequentially. dist  $(\mathcal{G}\ (j\ n)\ y)\ (h\ y) < inverse\ C$ 
    using uniform_limitD [OF ulim [of  $\{y\}$ ], of inverse C]  $\langle C > 0 \rangle\ y$ 
    by (simp add: dist_commute)
  then obtain  $n$  where  $dist\ (\mathcal{G}\ (j\ n)\ y)\ (h\ y) < inverse\ C$ 
    by (meson eventually_at_top_linorder order_refl)
  moreover
    have  $h\ y = h\ v$ 
      by (metis <h v = c> dist_commute that y)
    ultimately have  $cmod\ (inverse\ (\mathcal{F}\ (j\ n)\ y)) < inverse\ C$ 
      by (simp add: <h v = 0> G_def)
    then have  $C < norm\ (\mathcal{F}\ (j\ n)\ y)$ 
      by (metis G_def Gnot0 <y \in S> inverse_less_imp_less inverse_zero
norm_inverse_zero_less_norm_iff)
    show False
      using  $\langle C < cmod\ (\mathcal{F}\ (j\ n)\ y) \rangle\ le\_C\ not\_less$  by blast
    qed
  then show  $\exists x \in ball\ v\ e. h\ x \neq c$  by force
  qed
  show  $h$  holomorphic_on ball v e
    by (simp add: holh)
  show  $\bigwedge K. \llbracket compact\ K; K \subseteq ball\ v\ e \rrbracket \implies uniform\_limit\ K\ (\mathcal{G} \circ j)\ h$ 
sequentially
    by (simp add: ulim)
  qed (use <e > 0> in auto)
  with  $\langle h\ v = 0 \rangle$  show False by blast
  qed
  then obtain  $r$  where  $0 < r$  and  $r: \bigwedge h. h \in Y \implies cmod\ (h\ v) \leq r$ 
    by (metis not_le)
  moreover
    obtain  $B\ Z$  where  $0 < B$  open  $Z\ v \in Z\ Z \subseteq S\ \bigwedge h\ z. \llbracket h \in Y; z \in Z \rrbracket \implies$ 
norm $(h\ z) \leq B$ 
      using X01
      by (auto simp: r intro: GPicard1 [OF  $\langle open\ S \rangle\ \langle connected\ S \rangle\ \langle v \in S \rangle\ \langle r > 0 \rangle$ 
 $\langle Y \subseteq X \rangle\ holX$ ] X01)
    ultimately show  $v \in U$ 
      using  $v$  by (simp add: U_def) meson
  qed
  have  $\bigwedge x. x \in K \longrightarrow x \in U$ 

```

```

  using ‹U = S› ‹K ⊆ S› by blast
  then have  $\bigwedge x. x \in K \longrightarrow (\exists B Z. 0 < B \wedge \text{open } Z \wedge x \in Z \wedge$ 
     $(\forall h z'. h \in Y \wedge z' \in Z \longrightarrow \text{norm}(h z') \leq B))$ 
  unfolding U_def by blast
  then obtain F Z where F:  $\bigwedge x. x \in K \implies \text{open } (Z x) \wedge x \in Z x \wedge$ 
     $(\forall h z'. h \in Y \wedge z' \in Z x \longrightarrow \text{norm}(h z') \leq F x)$ 
  by metis
  then obtain L where L ⊆ K finite L and L: K ⊆ (⋃ c ∈ L. Z c)
  by (auto intro: compactE_image [OF ‹compact K›, of K Z])
  then have *:  $\bigwedge x h z'. [x \in L; h \in Y \wedge z' \in Z x] \implies \text{cmod } (h z') \leq F x$ 
  using F by blast
  have  $\exists B. \forall h z. h \in Y \wedge z \in K \longrightarrow \text{norm}(h z) \leq B$ 
  proof (cases L = {})
    case True with L show ?thesis by simp
  next
    case False
    then have  $\forall h z. h \in Y \wedge z \in K \longrightarrow (\exists x \in L. \text{cmod } (h z) \leq F x)$ 
    by (metis * L UN_E subset_iff)
    with False ‹finite L› show ?thesis
    by (rule_tac x = Max (F ` L) in exI) (simp add: linorder_class.Max_ge_iff)
  qed
  with that show ?thesis by metis
qed

```

lemma GPicard4:

```

  assumes 0 < k and holf: f holomorphic_on (ball 0 k - {0})
  and AE:  $\bigwedge e. [0 < e; e < k] \implies \exists d. 0 < d \wedge d < e \wedge (\forall z \in \text{sphere } 0 d. \text{norm}(f z) \leq B)$ 
  obtains  $\varepsilon$  where 0 <  $\varepsilon$   $\varepsilon < k \wedge z. z \in \text{ball } 0 \varepsilon - \{0\} \implies \text{norm}(f z) \leq B$ 
  proof -
    obtain  $\varepsilon$  where 0 <  $\varepsilon$   $\varepsilon < k/2$  and  $\varepsilon$ :  $\bigwedge z. \text{norm } z = \varepsilon \implies \text{norm}(f z) \leq B$ 
    using AE [of k/2] ‹0 < k› by auto
    show ?thesis
  proof
    show  $\varepsilon < k$ 
    using ‹0 < k› ‹ $\varepsilon < k/2$ › by auto
    show  $\text{cmod } (f \xi) \leq B$  if  $\xi: \xi \in \text{ball } 0 \varepsilon - \{0\}$  for  $\xi$ 
  proof -
    obtain d where 0 < d d < norm  $\xi$  and d:  $\bigwedge z. \text{norm } z = d \implies \text{norm}(f z) \leq B$ 
    using AE [of norm  $\xi$ ] ‹ $\varepsilon < k$ ›  $\xi$  by auto
    have [simp]:  $\text{closure } (\text{cball } 0 \varepsilon - \text{ball } 0 d) = \text{cball } 0 \varepsilon - \text{ball } 0 d$ 
    by (blast intro!: closure_closed)
    have [simp]:  $\text{interior } (\text{cball } 0 \varepsilon - \text{ball } 0 d) = \text{ball } 0 \varepsilon - \text{cball } (0::\text{complex}) d$ 
    using ‹0 <  $\varepsilon$ › ‹0 < d› by (simp add: interior_diff)
    have *:  $\text{norm}(f w) \leq B$  if w ∈ cball 0  $\varepsilon$  - ball 0 d for w
    proof (rule maximum_modulus_frontier [of f cball 0  $\varepsilon$  - ball 0 d])
      show f holomorphic_on interior (cball 0  $\varepsilon$  - ball 0 d)
    qed
  qed

```



```

    using ‹ $\varepsilon < k$ › ‹ $0 < d$ › that by (auto intro: holomorphic_on_subset [OF
hol])
    show continuous_on (closure (cball 0  $\varepsilon$  - ball 0  $d$ )) f
    proof (intro holomorphic_on_imp_continuous_on holomorphic_on_subset
[OF hol])
      show closure (cball 0  $\varepsilon$  - ball 0  $d$ )  $\subseteq$  ball 0  $k$  - {0}
        using ‹ $0 < d$ › ‹ $\varepsilon < k$ › by auto
      qed
      show  $\bigwedge z. z \in \text{frontier (cball 0 } \varepsilon - \text{ball 0 } d) \implies cmod (f z) \leq B$ 
        unfolding frontier_def
        using  $\varepsilon d$  less_eq_real_def by force
      qed (use that in auto)
    show ?thesis
      using * ‹ $d < cmod \xi$ › that by auto
    qed
  qed (use ‹ $0 < \varepsilon$ › in auto)
qed

```

lemma GPicard5:

```

  assumes hol: f holomorphic_on (ball 0 1 - {0})
    and f01:  $\bigwedge z. z \in \text{ball 0 1} - \{0\} \implies f z \neq 0 \wedge f z \neq 1$ 
  obtains e B where 0 < e e < 1 0 < B
    ( $\forall z \in \text{ball 0 } e - \{0\}. \text{norm}(f z) \leq B$ )  $\vee$ 
    ( $\forall z \in \text{ball 0 } e - \{0\}. \text{norm}(f z) \geq B$ )
proof -
  have [simp]: 1 + of_nat n  $\neq$  (0::complex) for n
    using of_nat_eq_0_iff by fastforce
  have [simp]: cmod (1 + of_nat n) = 1 + of_nat n for n
    by (metis norm_of_nat of_nat_Suc)
  have *: ( $\lambda x::\text{complex}. x / \text{of\_nat (Suc n)}$ ) ‘ (ball 0 1 - {0})  $\subseteq$  ball 0 1 - {0}
for n
  by (auto simp: norm_divide field_split_simps split: if_split_asm)
  define h where h  $\equiv \lambda n z::\text{complex}. f (z / (\text{Suc } n))$ 
  have holh: (h n) holomorphic_on ball 0 1 - {0} for n
    unfolding h_def
  proof (rule holomorphic_on_compose_gen [unfolded o_def, OF _ hol *])
    show ( $\lambda x. x / \text{of\_nat (Suc n)}$ ) holomorphic_on ball 0 1 - {0}
      by (intro holomorphic_intros) auto
    qed
  qed
  have h01:  $\bigwedge n z. z \in \text{ball 0 1} - \{0\} \implies h n z \neq 0 \wedge h n z \neq 1$ 
    unfolding h_def
    using * by (force intro!: f01)
  obtain w where w: w  $\in$  ball 0 1 - {0::complex}
    by (rule_tac w = 1/2 in that) auto
  consider infinite {n. norm(h n w)  $\leq$  1} | infinite {n. 1  $\leq$  norm(h n w)}
  by (metis (mono_tags, lifting) infinite_nat_iff_unbounded_le le_cases mem_Collect_eq)
  then show ?thesis
  proof cases

```

```

case 1
with infinite_enumerate obtain r :: nat ⇒ nat
  where strict_mono r and r:  $\bigwedge n. r\ n \in \{n. \text{norm}(h\ n\ w) \leq 1\}$ 
  by blast
obtain B where B:  $\bigwedge j\ z. [\text{norm}\ z = 1/2; j \in \text{range}(h \circ r)] \implies \text{norm}(j\ z) \leq B$ 
proof (rule GPicard3 [OF _ _ w, where K = sphere 0 (1/2)])
  show range (h ∘ r) ⊆
    {g. g holomorphic_on ball 0 1 - {0} ∧ (∀ z ∈ ball 0 1 - {0}. g z ≠ 0 ∧ g z ≠ 1)}
  using h01 by (auto intro: holomorphic_intros holomorphic_on_compose holh)
  show connected (ball 0 1 - {0::complex})
    by (simp add: connected_open_delete)
qed (use r in auto)
have normf_le_B: cmod(f z) ≤ B if norm z = 1 / (2 * (1 + of_nat (r n)))
for z n
proof -
  have *:  $\bigwedge w. \text{norm}\ w = 1/2 \implies \text{cmod}((f\ (w / (1 + \text{of\_nat}\ (r\ n)))))) \leq B$ 
  using B by (auto simp: h_def o_def)
  have half:  $\text{norm}\ (z * (1 + \text{of\_nat}\ (r\ n))) = 1/2$ 
  by (simp add: norm_mult divide_simps that)
  show ?thesis
  using * [OF half] by simp
qed
obtain ε where 0 < ε ε < 1  $\bigwedge z. z \in \text{ball}\ 0\ \varepsilon - \{0\} \implies \text{cmod}(f\ z) \leq B$ 
proof (rule GPicard4 [OF zero_less_one half, of B])
  fix e::real
  assume 0 < e e < 1
  obtain n where (1/e - 2) / 2 < real n
  using reals_Archimedean2 by blast
  also have ... ≤ r n
  using <strict_mono r> by (simp add: seq_suble)
  finally have (1/e - 2) / 2 < real (r n) .
  with <0 < e> have e: e > 1 / (2 + 2 * real (r n))
  by (simp add: field_simps)
  show ∃ d>0. d < e ∧ (∀ z ∈ sphere 0 d. cmod (f z) ≤ B)
  apply (rule_tac x=1 / (2 * (1 + of_nat (r n))) in exI)
  using normf_le_B by (simp add: e)
qed blast
then have ε: cmod (f z) ≤ |B| + 1 if cmod z < ε z ≠ 0 for z
  using that by fastforce
have 0 < |B| + 1
  by simp
then show ?thesis
  using ε by (force intro!: that [OF <0 < ε> <ε < 1>])
next
case 2
with infinite_enumerate obtain r :: nat ⇒ nat

```

```

    where strict_mono r and r:  $\bigwedge n. r n \in \{n. \text{norm}(h n w) \geq 1\}$ 
  by blast
  obtain B where B:  $\bigwedge j z. [\text{norm } z = 1/2; j \in \text{range } (\lambda n. \text{inverse} \circ h (r n))]$ 
 $\implies \text{norm}(j z) \leq B$ 
  proof (rule GPicard3 [OF _ _ w, where K = sphere 0 (1/2)])
    show range  $(\lambda n. \text{inverse} \circ h (r n)) \subseteq$ 
       $\{g. g \text{ holomorphic\_on ball } 0 \ 1 - \{0\} \wedge (\forall z \in \text{ball } 0 \ 1 - \{0\}. g z \neq 0 \wedge$ 
 $g z \neq 1)\}$ 
    using h01 by (auto intro!: holomorphic_intros holomorphic_on_compose_gen
[unfolded o_def, OF _ holh] holomorphic_on_compose)
    show connected (ball 0 1 - {0::complex})
      by (simp add: connected_open_delete)
    show  $\bigwedge j. j \in \text{range } (\lambda n. \text{inverse} \circ h (r n)) \implies \text{cmod } (j w) \leq 1$ 
      using r norm_inverse_le_norm by fastforce
  qed (use r in auto)
  have norm_if_le_B:  $\text{cmod}(\text{inverse } (f z)) \leq B$  if  $\text{norm } z = 1 / (2 * (1 +$ 
of_nat (r n))) for z n
  proof -
    have *:  $\text{inverse } (\text{cmod}((f (z / (1 + \text{of\_nat } (r n)))))) \leq B$  if  $\text{norm } z = 1/2$ 
for z
      using B [OF that] by (force simp: norm_inverse h_def)
    have half:  $\text{norm } (z * (1 + \text{of\_nat } (r n))) = 1/2$ 
      by (simp add: norm_mult divide_simps that)
    show ?thesis
      using * [OF half] by (simp add: norm_inverse)
  qed
  have hol_if:  $(\text{inverse} \circ f)$  holomorphic_on (ball 0 1 - {0})
    by (metis (no_types, lifting) holf comp_apply f01 holomorphic_on_inverse
holomorphic_transform)
  obtain  $\varepsilon$  where  $0 < \varepsilon \varepsilon < 1$  and leB:  $\bigwedge z. z \in \text{ball } 0 \ \varepsilon - \{0\} \implies \text{cmod}((\text{inverse}$ 
 $\circ f) z) \leq B$ 
  proof (rule GPicard4 [OF zero_less_one hol_if, of B])
    fix e::real
    assume  $0 < e \ e < 1$ 
    obtain n where  $(1/e - 2) / 2 < \text{real } n$ 
      using reals_Archimedean2 by blast
    also have  $\dots \leq r n$ 
      using <strict_mono r> by (simp add: seq_suble)
    finally have  $(1/e - 2) / 2 < \text{real } (r n)$  .
    with  $\langle 0 < e \rangle$  have e:  $e > 1 / (2 + 2 * \text{real } (r n))$ 
      by (simp add: field_simps)
    show  $\exists d > 0. d < e \wedge (\forall z \in \text{sphere } 0 \ d. \text{cmod } ((\text{inverse} \circ f) z) \leq B)$ 
      apply (rule_tac x=1 / (2 * (1 + of_nat (r n))) in exI)
      using norm_if_le_B by (simp add: e)
  qed blast
  have  $\varepsilon: \text{cmod } (f z) \geq \text{inverse } B$  and  $B > 0$  if  $\text{cmod } z < \varepsilon \ z \neq 0$  for z
  proof -
    have inverse (cmod (f z))  $\leq B$ 
      using leB that by (simp add: norm_inverse)

```

```

moreover
have  $f z \neq 0$ 
  using  $\langle \varepsilon < 1 \rangle$  f01 that by auto
ultimately show  $cmod (f z) \geq inverse B$ 
  by (simp add: norm_inverse inverse_le_imp_le)
show  $B > 0$ 
  using  $\langle f z \neq 0 \rangle \langle inverse (cmod (f z)) \leq B \rangle$  not_le order.trans by fastforce
qed
then have  $B > 0$ 
  by (metis  $\langle 0 < \varepsilon \rangle$  dense leI order.asym vector_choose_size)
then have  $inverse B > 0$ 
  by (simp add: field_split_simps)
then show ?thesis
  using  $\varepsilon$  that [OF  $\langle 0 < \varepsilon \rangle \langle \varepsilon < 1 \rangle$ ]
  by (metis Diff_iff dist_0_norm insert_iff mem_ball)
qed
qed

```

lemma *GPicard6*:

```

assumes open  $M z \in M a \neq 0$  and hol $f$ :  $f$  holomorphic_on  $(M - \{z\})$ 
and f0a:  $\bigwedge w. w \in M - \{z\} \implies f w \neq 0 \wedge f w \neq a$ 
obtains  $r$  where  $0 < r$   $ball\ z\ r \subseteq M$ 
   $bounded(f \text{ ` } (ball\ z\ r - \{z\})) \vee$ 
   $bounded((inverse \circ f) \text{ ` } (ball\ z\ r - \{z\}))$ 
proof  $-$ 
  obtain  $r$  where  $0 < r$  and  $r$ :  $ball\ z\ r \subseteq M$ 
    using assms openE by blast
  let  $?g = \lambda w. f (z + of\_real\ r * w) / a$ 
  obtain  $e B$  where  $0 < e\ e < 1\ 0 < B$ 
    and  $B$ :  $(\forall z \in ball\ 0\ e - \{0\}. norm(?g\ z) \leq B) \vee (\forall z \in ball\ 0\ e - \{0\}. norm(?g\ z) \geq B)$ 
  proof (rule GPicard5)
    show  $?g$  holomorphic_on  $ball\ 0\ 1 - \{0\}$ 
  proof (intro holomorphic_intros holomorphic_on_compose_gen [unfolded o_def, OF_holf])
    show  $(\lambda x. z + complex\_of\_real\ r * x) \text{ ` } (ball\ 0\ 1 - \{0\}) \subseteq M - \{z\}$ 
      using  $\langle 0 < r \rangle$   $r$ 
      by (auto simp: dist_norm norm_mult subset_eq)
    qed (use  $\langle a \neq 0 \rangle$  in auto)
    show  $\bigwedge w. w \in ball\ 0\ 1 - \{0\} \implies f (z + of\_real\ r * w) / a \neq 0 \wedge f (z + of\_real\ r * w) / a \neq 1$ 
      using f0a  $\langle 0 < r \rangle \langle a \neq 0 \rangle$   $r$ 
      by (auto simp: field_split_simps dist_norm norm_mult subset_eq)
    qed
  show ?thesis
proof
  show  $0 < e * r$ 
    by (simp add:  $\langle 0 < e \rangle \langle 0 < r \rangle$ )

```

```

have ball z (e * r) ⊆ ball z r
  by (simp add: ‹0 < r› ‹e < 1› order.strict_implies_order subset_ball)
then show ball z (e * r) ⊆ M
  using r by blast
consider ⋀z. z ∈ ball 0 e - {0} ⇒ norm(?g z) ≤ B | ⋀z. z ∈ ball 0 e - {0}
⇒ norm(?g z) ≥ B
  using B by blast
then show bounded (f ‘ (ball z (e * r) - {z})) ∨
  bounded ((inverse ∘ f) ‘ (ball z (e * r) - {z}))
proof cases
  case 1
  have ‹‹dist z w < e * r; w ≠ z› ⇒ cmod (f w) ≤ B * norm a for w
    using ‹a ≠ 0› ‹0 < r› 1 [of (w - z) / r]
    by (simp add: norm_divide dist_norm field_split_simps)
  then show ?thesis
    by (force simp: intro!: boundedI)
  next
  case 2
  have ‹‹dist z w < e * r; w ≠ z› ⇒ cmod (f w) ≥ B * norm a for w
    using ‹a ≠ 0› ‹0 < r› 2 [of (w - z) / r]
    by (simp add: norm_divide dist_norm field_split_simps)
  then have ‹‹dist z w < e * r; w ≠ z› ⇒ inverse (cmod (f w)) ≤ inverse (B
* norm a) for w
    by (metis ‹0 < B› ‹a ≠ 0› mult_pos_pos norm_inverse norm_inverse_le_norm
zero_less_norm_iff)
  then show ?thesis
    by (force simp: norm_inverse intro!: boundedI)
qed
qed
qed

```

theorem great_Picard:

assumes open M z ∈ M a ≠ b **and** holf: f holomorphic_on (M - {z})

and fab: ⋀w. w ∈ M - {z} ⇒ f w ≠ a ∧ f w ≠ b

obtains l **where** (f ⟶ l) (at z) ∨ ((inverse ∘ f) ⟶ l) (at z)

proof -

obtain r **where** 0 < r **and** zrM: ball z r ⊆ M

and r: bounded((λz. f z - a) ‘ (ball z r - {z})) ∨

bounded((inverse ∘ (λz. f z - a)) ‘ (ball z r - {z}))

proof (rule GPicard6 [OF ‹open M› ‹z ∈ M›])

show b - a ≠ 0

using assms **by** auto

show (λz. f z - a) holomorphic_on M - {z}

by (intro holomorphic_intros holf)

qed (use fab in auto)

have holfb: f holomorphic_on ball z r - {z}

using zrM **by** (auto intro: holomorphic_on_subset [OF holf])

have holfb_i: (λz. inverse(f z - a)) holomorphic_on ball z r - {z}

```

    using fab zrM by (fastforce intro!: holomorphic_intros holfb)
  show ?thesis
    using r
  proof
    assume bounded ((λz. f z - a) ‘ (ball z r - {z}))
    then obtain B where B: ∧w. w ∈ (λz. f z - a) ‘ (ball z r - {z}) ⇒ norm
w ≤ B
    by (force simp: bounded_iff)
    then have ∀x. x ≠ z ∧ dist x z < r → cmod (f x - a) ≤ B
    by (simp add: dist_commute)
    with ‹0 < r› have ∀F w in at z. cmod (f w - a) ≤ B
    by (force simp add: eventually_at)
    moreover have ∧x. cmod (f x - a) ≤ B ⇒ cmod (f x) ≤ B + cmod a
    by (metis add_commute add_le_cancel_right norm_triangle_sub order.trans)
    ultimately have ∃B. ∀F w in at z. cmod (f w) ≤ B
    by (metis (mono_tags, lifting) eventually_at)
    then obtain g where holg: g holomorphic_on ball z r and gf: ∧w. w ∈ ball
z r - {z} ⇒ g w = f w
    using ‹0 < r› holomorphic_on_extend_bounded [OF holfb] by auto
    then have g -z → g z
    unfolding continuous_at [symmetric]
    using ‹0 < r› centre_in_ball field_differentiable_imp_continuous_at
holomorphic_on_imp_differentiable_at by blast
    then have (f → g z) (at z)
    using Lim_transform_within_open [of g g z z]
    using ‹0 < r› centre_in_ball gf by blast
    then show ?thesis
    using that by blast
  next
    assume bounded((inverse ∘ (λz. f z - a)) ‘ (ball z r - {z}))
    then obtain B where B: ∧w. w ∈ (inverse ∘ (λz. f z - a)) ‘ (ball z r - {z})
⇒ norm w ≤ B
    by (force simp: bounded_iff)
    then have ∀x. x ≠ z ∧ dist x z < r → cmod (inverse (f x - a)) ≤ B
    by (simp add: dist_commute)
    with ‹0 < r› have ∀F w in at z. cmod (inverse (f w - a)) ≤ B
    by (auto simp add: eventually_at)
    then have ∃B. ∀F z in at z. cmod (inverse (f z - a)) ≤ B
    by blast
    then obtain g where holg: g holomorphic_on ball z r and gf: ∧w. w ∈ ball
z r - {z} ⇒ g w = inverse (f w - a)
    using ‹0 < r› holomorphic_on_extend_bounded [OF holfb_i] by auto
    then have gz: g -z → g z
    unfolding continuous_at [symmetric]
    using ‹0 < r› centre_in_ball field_differentiable_imp_continuous_at
holomorphic_on_imp_differentiable_at by blast
    have gnz: ∧w. w ∈ ball z r - {z} ⇒ g w ≠ 0
    using gf fab zrM by fastforce
    show ?thesis

```

```

proof (cases g z = 0)
  case True
    have *:  $\llbracket g \neq 0; \text{inverse } g = f - a \rrbracket \implies g / (1 + a * g) = \text{inverse } f$  for f
  g::complex
    by (auto simp: field_simps)
    have (inverse o f) -z → 0
    proof (rule Lim_transform_within_open [of  $\lambda w. g w / (1 + a * g w)$  _ _
  UNIV ball z r])
      show ( $\lambda w. g w / (1 + a * g w)$ ) -z → 0
        using True by (auto simp: intro!: tendsto_eq_intros gz)
      show  $\bigwedge x. \llbracket x \in \text{ball } z r; x \neq z \rrbracket \implies g x / (1 + a * g x) = (\text{inverse } o f) x$ 
        using * gf gnz by simp
    qed (use <0 < r> in auto)
    with that show ?thesis by blast
  next
  case False
  show ?thesis
  proof (cases 1 + a * g z = 0)
    case True
      have (f → 0) (at z)
      proof (rule Lim_transform_within_open [of  $\lambda w. (1 + a * g w) / g w$  _ _
  _ ball z r])
        show ( $\lambda w. (1 + a * g w) / g w$ ) -z → 0
          by (rule tendsto_eq_intros refl gz <g z ≠ 0> | simp add: True)+
        show  $\bigwedge x. \llbracket x \in \text{ball } z r; x \neq z \rrbracket \implies (1 + a * g x) / g x = f x$ 
          using fab fab zrM by (fastforce simp add: gf field_split_simps)
      qed (use <0 < r> in auto)
      then show ?thesis
        using that by blast
    next
    case False
      have *:  $\llbracket g \neq 0; \text{inverse } g = f - a \rrbracket \implies g / (1 + a * g) = \text{inverse } f$  for f
  g::complex
      by (auto simp: field_simps)
      have (inverse o f) -z → g z / (1 + a * g z)
      proof (rule Lim_transform_within_open [of  $\lambda w. g w / (1 + a * g w)$  _ _
  UNIV ball z r])
        show ( $\lambda w. g w / (1 + a * g w)$ ) -z → g z / (1 + a * g z)
          using False by (auto simp: False intro!: tendsto_eq_intros gz)
        show  $\bigwedge x. \llbracket x \in \text{ball } z r; x \neq z \rrbracket \implies g x / (1 + a * g x) = (\text{inverse } o f) x$ 
          using * gf gnz by simp
      qed (use <0 < r> in auto)
      with that show ?thesis by blast
  qed
qed
qed
qed

```

corollary *great_Picard_alt*:

assumes M : open M $z \in M$ **and** *hol* f : *f* holomorphic_on $(M - \{z\})$
and *non*: $\bigwedge l. \neg (f \longrightarrow l) \text{ (at } z) \bigwedge l. \neg ((\text{inverse} \circ f) \longrightarrow l) \text{ (at } z)$
obtains a **where** $-\{a\} \subseteq f^{-1}(M - \{z\})$
unfolding *subset_iff_image_iff*
by (*metis* *great_Picard* [*OF* M hol] *non* *Compl_iff_insertI1*)

corollary *great_Picard_infinite*:

assumes M : open M $z \in M$ **and** *hol* f : *f* holomorphic_on $(M - \{z\})$
and *non*: $\bigwedge l. \neg (f \longrightarrow l) \text{ (at } z) \bigwedge l. \neg ((\text{inverse} \circ f) \longrightarrow l) \text{ (at } z)$
obtains a **where** $\bigwedge w. w \neq a \implies \text{infinite } \{x. x \in M - \{z\} \wedge f x = w\}$
proof –
have *False* **if** $a \neq b$ **and** *ab*: *finite* $\{x. x \in M - \{z\} \wedge f x = a\}$ *finite* $\{x. x \in M - \{z\} \wedge f x = b\}$ **for** a b
proof –
have *finab*: *finite* $\{x. x \in M - \{z\} \wedge f x \in \{a, b\}\}$
using *finite_UnI* [*OF* *ab*] **unfolding** *mem_Collect_eq* *insert_iff_empty_iff*
by (*simp* *add*: *conj_disj_distribL*)
obtain r **where** $0 < r$ **and** *zrM*: *ball* z $r \subseteq M$ **and** r : $\bigwedge x. [x \in M - \{z\}; f x \in \{a, b\}] \implies x \notin \text{ball } z$ r
proof –
obtain e **where** $e > 0$ **and** e : *ball* z $e \subseteq M$
using *assms* *openE* **by** *blast*
show *?thesis*
proof (*cases* $\{x \in M - \{z\}. f x \in \{a, b\}\} = \{\}$)
case *True*
then show *?thesis*
using e $\langle e > 0 \rangle$ **that** **by** *fastforce*
next
case *False*
let $?r = \min e$ (*Min* (*dist* z^{-1} $\{x \in M - \{z\}. f x \in \{a, b\}\}$))
show *?thesis*
proof
show $0 < ?r$
using *min_less_iff_conj* *Min_gr_iff* *finab* *False* $\langle 0 < e \rangle$ **by** *auto*
have *ball* z $?r \subseteq \text{ball } z$ e
by (*simp* *add*: *subset_ball*)
with e **show** *ball* z $?r \subseteq M$ **by** *blast*
show $\bigwedge x. [x \in M - \{z\}; f x \in \{a, b\}] \implies x \notin \text{ball } z$ $?r$
using *min_less_iff_conj* *Min_gr_iff* *finab* *False* $\langle 0 < e \rangle$ **by** *auto*
qed
qed
qed
have *holfb*: *f* holomorphic_on (*ball* z $r - \{z\}$)
apply (*rule* *holomorphic_on_subset* [*OF* *hol*])
using *zrM* **by** *auto*
show *?thesis*
apply (*rule* *great_Picard* [*OF* *open_ball* $\langle a \neq b \rangle$ *holfb*])


```

    using non ⟨0 < r⟩ r zrM by auto
  qed
  with that show thesis
    by meson
  qed

```

theorem *Casorati_Weierstrass*:

```

  assumes open M z ∈ M f holomorphic_on (M - {z})
    and  $\bigwedge l. \neg (f \longrightarrow l) (at z) \wedge l. \neg ((inverse \circ f) \longrightarrow l) (at z)$ 
  shows  $closure(f^{-1}(M - \{z\})) = UNIV$ 
  proof -
    obtain a where  $a: - \{a\} \subseteq f^{-1}(M - \{z\})$ 
      using great_Picard_alt [OF assms].
    have  $UNIV = closure(- \{a\})$ 
      by (simp add: closure_interior)
    also have  $\dots \subseteq closure(f^{-1}(M - \{z\}))$ 
      by (simp add: a_closure_mono)
    finally show ?thesis
      by blast
  qed
end

```

7 Moebius functions, Equivalents of Simply Connected Sets, Riemann Mapping Theorem

```

theory Riemann_Mapping
imports Great_Picard
begin

```

7.1 Moebius functions are biholomorphisms of the unit disc

definition *Moebius_function* :: $[real, complex, complex] \Rightarrow complex$ **where**
 $Moebius_function \equiv \lambda t w z. \exp(i * of_real t) * (z - w) / (1 - cnj w * z)$

lemma *Moebius_function_simple*:

```

  Moebius_function 0 w z = (z - w) / (1 - cnj w * z)
  by (simp add: Moebius_function_def)

```

lemma *Moebius_function_eq_zero*:

```

  Moebius_function t w w = 0
  by (simp add: Moebius_function_def)

```

lemma *Moebius_function_of_zero*:

```

  Moebius_function t w 0 = - exp(i * of_real t) * w
  by (simp add: Moebius_function_def)

```

lemma *Moebius_function_norm_lt_1*:

```

assumes w1: norm w < 1 and z1: norm z < 1
shows norm (Moebius_function t w z) < 1
proof -
  have 1 - cnj w * z ≠ 0
    by (metis complex_cnj_cnj complex_mod_sqrt_Re_mult_cnj mult.commute
mult_less_cancel_right1 norm_ge_zero norm_mult norm_one order.asym right_minus_eq
w1 z1)
  then have VV: 1 - w * cnj z ≠ 0
    by (metis complex_cnj_cnj complex_cnj_mult complex_cnj_one right_minus_eq)
  then have 1 - norm (Moebius_function t w z) ^ 2 =
    ((1 - norm w ^ 2) / (norm (1 - cnj w * z) ^ 2)) * (1 - norm z ^ 2)
    apply (cases w)
    apply (cases z)
  apply (simp add: Moebius_function_def divide_simps norm_divide norm_mult)
  apply (simp add: complex_norm complex_diff complex_mult one_complex.code
complex_cnj)
  apply (auto simp: algebra_simps power2_eq_square)
  done
  then have 1 - (cmod (Moebius_function t w z))^2 = (1 - cmod (w * w)) /
(cmod (1 - cnj w * z))^2 * (1 - cmod (z * z))
    by (simp add: norm_mult power2_eq_square)
  moreover have 0 < 1 - cmod (z * z)
    by (metis (no_types) z1 diff_gt_0_iff_gt mult.left_neutral norm_mult_less)
  ultimately have 0 < 1 - norm (Moebius_function t w z) ^ 2
    using ⟨1 - cnj w * z ≠ 0⟩ w1 norm_mult_less by fastforce
  then show ?thesis
    using linorder_not_less by fastforce
qed

```

lemma Moebius_function_holomorphic:

```

assumes norm w < 1
shows Moebius_function t w holomorphic_on ball 0 1
proof -
  have *: 1 - z * w ≠ 0 if norm z < 1 for z
  proof -
    have norm (1::complex) ≠ norm (z * w)
      using assms that norm_mult_less by fastforce
    then show ?thesis by auto
  qed
  show ?thesis
    unfolding Moebius_function_def
  proof (intro holomorphic_intros)
    show  $\bigwedge z. z \in \text{ball } 0 \ 1 \implies 1 - \text{cnj } w * z \neq 0$ 
      by (metis * complex_cnj_cnj complex_cnj_mult complex_mod_cnj mem_ball_0
mult.commute mult_1 right_minus_eq)
  qed
qed

```

lemma Moebius_function_compose:

```

assumes meq:  $-w1 = w2$  and  $norm\ w1 < 1$   $norm\ z < 1$ 
shows  $Moebius\_function\ 0\ w1\ (Moebius\_function\ 0\ w2\ z) = z$ 
proof -
  have  $norm\ w2 < 1$ 
    using assms by auto
  then have  $-w1 = z$  if  $cnj\ w2 * z = 1$ 
    by (metis assms(3) complex_mod_cnj_less_irrefl mult.right_neutral norm_mult_less_norm_one that)
  moreover have  $z=0$  if  $1 - cnj\ w2 * z = cnj\ w1 * (z - w2)$ 
  proof -
    have  $w2 * cnj\ w2 = 1$ 
      using that meq by (auto simp: algebra_simps)
    then show  $z = 0$ 
      using  $\langle cmod\ w2 < 1 \rangle$  complex_mod_sqrt_Re_mult_cnj by force
  qed
  moreover have  $z - w2 - w1 * (1 - cnj\ w2 * z) = z * (1 - cnj\ w2 * z - cnj\ w1 * (z - w2))$ 
    using meq by (fastforce simp: algebra_simps)
  ultimately
  show ?thesis
    by (simp add: Moebius_function_def divide_simps norm_divide norm_mult)
qed

```

lemma *ball_biholomorphism_exists*:

```

assumes  $a \in ball\ 0\ 1$ 
obtains  $f\ g$  where  $f\ a = 0$ 
   $f$  holomorphic_on  $ball\ 0\ 1$   $f^{-1}\ ball\ 0\ 1 \subseteq ball\ 0\ 1$ 
   $g$  holomorphic_on  $ball\ 0\ 1$   $g^{-1}\ ball\ 0\ 1 \subseteq ball\ 0\ 1$ 
   $\bigwedge z. z \in ball\ 0\ 1 \implies f\ (g\ z) = z$ 
   $\bigwedge z. z \in ball\ 0\ 1 \implies g\ (f\ z) = z$ 

```

proof

```

show  $Moebius\_function\ 0\ a$  holomorphic_on  $ball\ 0\ 1$   $Moebius\_function\ 0\ (-a)$  holomorphic_on  $ball\ 0\ 1$ 
  using Moebius_function_holomorphic assms mem_ball_0 by auto
show  $Moebius\_function\ 0\ a\ a = 0$ 
  by (simp add: Moebius_function_eq_zero)
show  $Moebius\_function\ 0\ a^{-1}\ ball\ 0\ 1 \subseteq ball\ 0\ 1$ 
   $Moebius\_function\ 0\ (-a)^{-1}\ ball\ 0\ 1 \subseteq ball\ 0\ 1$ 
  using Moebius_function_norm_lt_1 assms by auto
show  $Moebius\_function\ 0\ a\ (Moebius\_function\ 0\ (-a)\ z) = z$ 
   $Moebius\_function\ 0\ (-a)\ (Moebius\_function\ 0\ a\ z) = z$  if  $z \in ball\ 0\ 1$  for
 $z$ 
  using Moebius_function_compose assms that by auto
qed

```

7.2 A big chain of equivalents of simple connectedness for an open set

lemma *biholomorphic_to_disc_aux*:

```

assumes open S connected S 0 ∈ S and S01: S ⊆ ball 0 1
and prev: ⋀f. [f holomorphic_on S; ⋀z. z ∈ S ⇒ f z ≠ 0; inj_on f S]
    ⇒ ∃g. g holomorphic_on S ∧ (∀z ∈ S. f z = (g z)2)
shows ∃f g. f holomorphic_on S ∧ g holomorphic_on ball 0 1 ∧
    (∀z ∈ S. f z ∈ ball 0 1 ∧ g(f z) = z) ∧
    (∀z ∈ ball 0 1. g z ∈ S ∧ f(g z) = z)

proof –
define F where F ≡ {h. h holomorphic_on S ∧ h ‘ S ⊆ ball 0 1 ∧ h 0 = 0 ∧
inj_on h S}
have idF: id ∈ F
using S01 by (auto simp: F_def)
then have F ≠ {}
by blast
have imF_ne: ((λh. norm(deriv h 0)) ‘ F) ≠ {}
using idF by auto
have holF: ⋀h. h ∈ F ⇒ h holomorphic_on S
by (auto simp: F_def)
obtain f where f ∈ F and normf: ⋀h. h ∈ F ⇒ norm(deriv h 0) ≤ norm(deriv
f 0)
proof –
obtain r where r > 0 and r: ball 0 r ⊆ S
using ⟨open S⟩ ⟨0 ∈ S⟩ openE by auto
have bdd: bdd_above ((λh. norm(deriv h 0)) ‘ F)
proof (intro bdd_aboveI exI ballI, clarify)
show norm (deriv f 0) ≤ 1 / r if f ∈ F for f
proof –
have r01: (*) (complex_of_real r) ‘ ball 0 1 ⊆ S
using that ⟨r > 0⟩ by (auto simp: norm_mult r [THEN subsetD])
then have f holomorphic_on (*) (complex_of_real r) ‘ ball 0 1
using holomorphic_on_subset [OF holF] by (simp add: that)
then have holf: f ∘ (λz. (r * z)) holomorphic_on (ball 0 1)
by (intro holomorphic_intros holomorphic_on_compose)
have f0: (f ∘ (*) (complex_of_real r)) 0 = 0
using F_def that by auto
have f ‘ S ⊆ ball 0 1
using F_def that by blast
with r01 have fr1: ⋀z. norm z < 1 ⇒ norm ((f ∘ (*) (of_real r))z) < 1
by force
have *: ((λw. f (r * w)) has_field_derivative deriv f (r * z) * r) (at z)
if z ∈ ball 0 1 for z::complex
using DERIV_chain' [where g=f] ⟨open S⟩
by (meson DERIV_cmult_Id ⟨f ∈ F⟩ holF holomorphic_derivI im-
age_subset_iff
r01 that)
have df0: ((λw. f (r * w)) has_field_derivative deriv f 0 * r) (at 0)
using * [of 0] by simp
have deq: deriv (λx. f (complex_of_real r * x)) 0 = deriv f 0 * com-
plex_of_real r
using DERIV_imp_deriv df0 by blast

```

```

    have norm (deriv (f o (*)) (complex_of_real r)) 0 ≤ 1
      by (auto intro: Schwarz_Lemma [OF holf f0 fr1, of 0])
    with ⟨r > 0⟩ show ?thesis
      by (simp add: deq norm_mult divide_simps o_def)
  qed
  qed
  define l where l ≡ SUP h∈F. norm (deriv h 0)
  have eql: norm (deriv f 0) = l if le: l ≤ norm (deriv f 0) and f ∈ F for f
  proof (rule order_antisym [OF _ le])
    show cmod (deriv f 0) ≤ l
      using ⟨f ∈ F⟩ bdd cSUP_upper by (fastforce simp: l_def)
  qed
  obtain  $\mathcal{F}$  where  $\mathcal{F}$ in:  $\bigwedge n. \mathcal{F} n \in F$  and  $\mathcal{F}$ lim:  $(\lambda n. \text{norm} (\text{deriv} (\mathcal{F} n) 0))$ 
     $\longrightarrow l$ 
  proof -
    have  $\exists f. f \in F \wedge |\text{norm} (\text{deriv} f 0) - l| < 1 / (\text{Suc } n)$  for n
    proof -
      obtain f where f ∈ F and f:  $l < \text{norm} (\text{deriv} f 0) + 1/(\text{Suc } n)$ 
        using cSup_least [OF imF_ne, of l - 1/(Suc n)] by (fastforce simp:
      l_def)
      then have  $|\text{norm} (\text{deriv} f 0) - l| < 1 / (\text{Suc } n)$ 
        by (fastforce simp: abs_if_not_less eql)
      with ⟨f ∈ F⟩ show ?thesis
        by blast
    qed
    then obtain  $\mathcal{F}$  where fF:  $\bigwedge n. (\mathcal{F} n) \in F$ 
      and fless:  $\bigwedge n. |\text{norm} (\text{deriv} (\mathcal{F} n) 0) - l| < 1 / (\text{Suc } n)$ 
      by metis
    have  $(\lambda n. \text{norm} (\text{deriv} (\mathcal{F} n) 0)) \longrightarrow l$ 
    proof (rule metric_LIMSEQ_I)
      fix e::real
      assume e > 0
      then obtain N::nat where N: e > 1/(Suc N)
        using nat_approx_posE by blast
      show  $\exists N. \forall n \geq N. \text{dist} (\text{norm} (\text{deriv} (\mathcal{F} n) 0)) l < e$ 
      proof (intro exI allI impI)
        fix n assume N ≤ n
        have  $\text{dist} (\text{norm} (\text{deriv} (\mathcal{F} n) 0)) l < 1 / (\text{Suc } n)$ 
          using fless by (simp add: dist_norm)
        also have ... < e
          using N ⟨N ≤ n⟩ inverse_of_nat_le le_less_trans by blast
        finally show  $\text{dist} (\text{norm} (\text{deriv} (\mathcal{F} n) 0)) l < e$  .
      qed
    qed
  qed
  with fF show ?thesis
    using that by blast
  qed
  have  $\bigwedge K. [\text{compact } K; K \subseteq S] \implies \exists B. \forall h \in F. \forall z \in K. \text{norm} (h z) \leq B$ 
  by (rule_tac x=1 in exI) (force simp: F_def)

```

```

moreover have range  $\mathcal{F} \subseteq F$ 
using  $\langle \bigwedge n. \mathcal{F} n \in F \rangle$  by blast
ultimately obtain  $f$  and  $r :: \text{nat} \Rightarrow \text{nat}$ 
where  $\text{hol}f: f \text{ holomorphic\_on } S$  and  $r: \text{strict\_mono } r$ 
and  $\text{lim}f: \bigwedge x. x \in S \Longrightarrow (\lambda n. \mathcal{F} (r n) x) \longrightarrow f x$ 
and  $\text{ulim}f: \bigwedge K. [\text{compact } K; K \subseteq S] \Longrightarrow \text{uniform\_limit } K (\mathcal{F} \circ r) f$ 
sequentially
using Montel [of  $S F \mathcal{F}$ , OF  $\langle \text{open } S \rangle \text{hol}f$ ] by auto+
have  $\text{der}: \bigwedge n x. x \in S \Longrightarrow ((\mathcal{F} \circ r) n \text{ has\_field\_derivative } ((\lambda n. \text{deriv } (\mathcal{F} n)) \circ r) n x) \text{ (at } x)$ 
using  $\langle \bigwedge n. \mathcal{F} n \in F \rangle \langle \text{open } S \rangle \text{hol}f \text{ holomorphic\_derivI}$  by fastforce
have  $\text{ulim}: \bigwedge x. x \in S \Longrightarrow \exists d > 0. \text{cball } x d \subseteq S \wedge \text{uniform\_limit } (\text{cball } x d) (\mathcal{F} \circ r) f \text{ sequentially}$ 
by (meson  $\text{ulim}f \langle \text{open } S \rangle \text{compact\_cball\_open\_contains\_cball}$ )
obtain  $f' :: \text{complex} \Rightarrow \text{complex}$  where  $f': (f \text{ has\_field\_derivative } f' 0) \text{ (at } 0)$ 
and  $\text{tof}'0: (\lambda n. ((\lambda n. \text{deriv } (\mathcal{F} n)) \circ r) n 0) \longrightarrow f' 0$ 
using  $\text{has\_complex\_derivative\_uniform\_sequence}$  [OF  $\langle \text{open } S \rangle \text{der ulim}$ ]  $\langle 0 \in S \rangle$  by metis
then have  $\text{der}f0: \text{deriv } f 0 = f' 0$ 
by (simp add: DERIV_imp_deriv)
have  $f \text{ field\_differentiable (at } 0)$ 
using  $\text{field\_differentiable\_def } f'$  by blast
have  $(\lambda x. (\text{norm } (\text{deriv } (\mathcal{F} (r x)) 0))) \longrightarrow \text{norm } (\text{deriv } f 0)$ 
using  $\text{isCont\_tendsto\_compose}$  [OF  $\text{continuous\_norm}$  [OF  $\text{continuous\_ident}$ ]  $\text{tof}'0$ ]  $\text{der}f0$  by auto
with LIMSEQ_subseq_LIMSEQ [OF  $\mathcal{F} \text{lim } r$ ] have  $\text{no\_df}0: \text{norm}(\text{deriv } f 0) = l$ 
by (force simp:  $\text{o\_def}$  intro:  $\text{tendsto\_unique}$ )
have  $\text{nonconst}f: \neg f \text{ constant\_on } S$ 
using  $\langle \text{open } S \rangle \langle 0 \in S \rangle \text{no\_df}0 \text{ holomorphic\_nonconstant}$  [OF  $\text{hol}f$ ]  $\text{eql}$  [OF  $\_ \text{id}F$ ]
by force
show ?thesis
proof
show  $f \in F$ 
unfolding  $F\_def$ 
proof (intro CollectI conjI  $\text{hol}f$ )
have  $\text{norm}(f z) \leq 1$  if  $z \in S$  for  $z$ 
proof (intro Lim_norm_ubound [OF  $\_ \text{lim}f$ ]  $\text{always\_eventually allI}$  that)
fix  $n$ 
have  $\mathcal{F} (r n) \in F$ 
by (simp add:  $\mathcal{F} \text{in}$ )
then show  $\text{norm } (\mathcal{F} (r n) z) \leq 1$ 
using that by (auto simp:  $F\_def$ )
qed simp
then have  $\text{fless}1: \text{norm}(f z) < 1$  if  $z \in S$  for  $z$ 
using  $\text{maximum\_modulus\_principle}$  [OF  $\text{hol}f \langle \text{open } S \rangle \langle \text{connected } S \rangle \langle \text{open } S \rangle$ ]  $\text{nonconst}f$  that
by fastforce

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    then show  $f^{-1} S \subseteq \text{ball } 0 \ 1$ 
      by auto
    have  $(\lambda n. \mathcal{F} (r \ n) \ 0) \longrightarrow 0$ 
      using  $\mathcal{F}in$  by (auto simp:  $F\_def$ )
    then show  $f \ 0 = 0$ 
      using  $tendsto\_unique$  [ $OF \_ limf$ ]  $\langle 0 \in S \rangle$   $trivial\_limit\_sequentially$  by
blast
    show  $inj\_on \ f \ S$ 
    proof (rule  $Hurwitz\_injective$  [ $OF \langle open \ S \rangle \langle connected \ S \rangle \_ \ hol$ ])
      show  $\bigwedge n. (\mathcal{F} \circ r) \ n \ holomorphic\_on \ S$ 
        by (simp add:  $\mathcal{F}in \ holF$ )
      show  $\bigwedge K. [\![compact \ K; K \subseteq S]\!] \implies uniform\_limit \ K (\mathcal{F} \circ r) \ f \ sequentially$ 
        by (metis  $ulimf$ )
      show  $\neg f \ constant\_on \ S$ 
        using  $nonconstf$  by auto
      show  $\bigwedge n. inj\_on \ ((\mathcal{F} \circ r) \ n) \ S$ 
        using  $\mathcal{F}in$  by (auto simp:  $F\_def$ )
    qed
  qed
  show  $\bigwedge h. h \in F \implies norm (deriv \ h \ 0) \leq norm (deriv \ f \ 0)$ 
    by (metis  $eql \ le\_cases \ no\_df0$ )
  qed
  have  $hol$ :  $f \ holomorphic\_on \ S$  and  $inj$ :  $inj\_on \ f \ S$  and  $f01$ :  $f^{-1} S \subseteq \text{ball } 0 \ 1$ 
    using  $\langle f \in F \rangle$  by (auto simp:  $F\_def$ )
  obtain  $g$  where  $holg$ :  $g \ holomorphic\_on \ (f^{-1} S)$ 
    and  $derg$ :  $\bigwedge z. z \in S \implies deriv \ f \ z * deriv \ g \ (f \ z) = 1$ 
    and  $gf$ :  $\bigwedge z. z \in S \implies g(f \ z) = z$ 
    using  $holomorphic\_has\_inverse$  [ $OF \ hol \langle open \ S \rangle \ inj$ ] by metis
  have  $\text{ball } 0 \ 1 \subseteq f^{-1} S$ 
  proof
    fix  $a::complex$ 
    assume  $a$ :  $a \in \text{ball } 0 \ 1$ 
    have  $False$  if  $\bigwedge x. x \in S \implies f \ x \neq a$ 
    proof -
      obtain  $h \ k$  where  $h \ a = 0$ 
        and  $holh$ :  $h \ holomorphic\_on \ \text{ball } 0 \ 1$  and  $h01$ :  $h^{-1} \ \text{ball } 0 \ 1 \subseteq \text{ball } 0 \ 1$ 
        and  $holk$ :  $k \ holomorphic\_on \ \text{ball } 0 \ 1$  and  $k01$ :  $k^{-1} \ \text{ball } 0 \ 1 \subseteq \text{ball } 0 \ 1$ 
        and  $hk$ :  $\bigwedge z. z \in \text{ball } 0 \ 1 \implies h (k \ z) = z$ 
        and  $kh$ :  $\bigwedge z. z \in \text{ball } 0 \ 1 \implies k (h \ z) = z$ 
        using  $ball\_biholomorphism\_exists$  [ $OF \ a$ ] by blast
      have  $nf1$ :  $\bigwedge z. z \in S \implies norm(f \ z) < 1$ 
        using  $\langle f \in F \rangle$  by (auto simp:  $F\_def$ )
      have  $1$ :  $h \circ f \ holomorphic\_on \ S$ 
        using  $F\_def \langle f \in F \rangle \ holh \ holomorphic\_on\_compose \ holomorphic\_on\_subset$ 
        by blast
      have  $2$ :  $\bigwedge z. z \in S \implies (h \circ f) \ z \neq 0$ 
        by (metis  $\langle h \ a = 0 \rangle \ a \ comp\_eq\_dest\_lhs \ nf1 \ kh \ mem\_ball\_0 \ that$ )
      have  $3$ :  $inj\_on \ (h \circ f) \ S$ 

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    by (metis (no_types, lifting) F_def ‹f ∈ F› comp_inj_on inj_on_inverseI
    injf kh mem_Collect_eq inj_on_subset)
  obtain  $\psi$  where hol $\psi$ :  $\psi$  holomorphic_on ((h ∘ f) ‘ S)
  and  $\psi$ 2:  $\bigwedge z. z \in S \implies \psi(h(fz)) \wedge^2 = h(fz)$ 
  proof (rule exE [OF prev [OF 1 2 3]], safe)
    fix  $\vartheta$ 
    assume hol $\vartheta$ :  $\vartheta$  holomorphic_on S and  $\vartheta$ 2:  $(\forall z \in S. (h \circ f) z = (\vartheta z)^2)$ 
    show thesis
  proof
    show  $(\vartheta \circ g \circ k)$  holomorphic_on (h ∘ f) ‘ S
  proof (intro holomorphic_on_compose)
    show k holomorphic_on (h ∘ f) ‘ S
      using holomorphic_on_subset [OF holk] f01 h01 by force
    show g holomorphic_on k ‘ (h ∘ f) ‘ S
      using holomorphic_on_subset [OF holg] by (force simp: kh nf1)
    show  $\vartheta$  holomorphic_on g ‘ k ‘ (h ∘ f) ‘ S
      using holomorphic_on_subset [OF hol $\vartheta$ ] by (force simp: gf kh nf1)
  qed
  show  $((\vartheta \circ g \circ k)(h(fz)))^2 = h(fz)$  if  $z \in S$  for  $z$ 
    using  $\vartheta$ 2 gf kh nf1 that by fastforce
  qed
  qed
  have norm $\psi$ 1:  $\text{norm}(\psi(h(fz))) < 1$  if  $z \in S$  for  $z$ 
  by (metis  $\psi$ 2 h01 image_subset_iff mem_ball_0 nf1 norm_power power_less1_D
  that)
  then have  $\psi$ 01:  $\psi(h(f0)) \in \text{ball } 0 \ 1$ 
  by (simp add: ‹0 ∈ S›)
  obtain p q where p0:  $p(\psi(h(f0))) = 0$ 
  and holp: p holomorphic_on ball 0 1 and p01:  $p \text{ ‘ ball } 0 \ 1 \subseteq \text{ball } 0 \ 1$ 
  and holq: q holomorphic_on ball 0 1 and q01:  $q \text{ ‘ ball } 0 \ 1 \subseteq \text{ball } 0 \ 1$ 
  and pq:  $\bigwedge z. z \in \text{ball } 0 \ 1 \implies p(qz) = z$ 
  and qp:  $\bigwedge z. z \in \text{ball } 0 \ 1 \implies q(pz) = z$ 
  using ball_biholomorphism_exists [OF  $\psi$ 01] by metis
  have p ∘  $\psi$  ∘ h ∘ f ∈ F
  unfolding F_def
  proof (intro CollectI conjI holp)
    show p ∘  $\psi$  ∘ h ∘ f holomorphic_on S
  proof (intro holomorphic_on_compose holp)
    show h holomorphic_on f ‘ S
      using holomorphic_on_subset [OF holh] f01 by fastforce
    show  $\psi$  holomorphic_on h ‘ f ‘ S
      using holomorphic_on_subset [OF hol $\psi$ ] by fastforce
    show p holomorphic_on  $\psi \text{ ‘ h ‘ f ‘ S}$ 
      using holomorphic_on_subset [OF holp] by (simp add: image_subset_iff
  norm $\psi$ 1)
  qed
  show (p ∘  $\psi$  ∘ h ∘ f) ‘ S ⊆ ball 0 1
    using norm $\psi$ 1 p01 by fastforce
  show (p ∘  $\psi$  ∘ h ∘ f) 0 = 0

```



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    by (simp add: ⟨p (ψ (h (f 0))) = 0⟩)
  show inj_on (p ∘ ψ ∘ h ∘ f) S
    unfolding inj_on_def o_def
    by (metis ψ2 dist_0_norm gf kh mem_ball nf1 normψ1 qp)
qed
then have le_norm_df0: norm (deriv (p ∘ ψ ∘ h ∘ f) 0) ≤ norm (deriv f 0)
  by (rule normf)
have 1: k ∘ power2 ∘ q holomorphic_on ball 0 1
proof (intro holomorphic_on_compose holq)
  show power2 holomorphic_on q ‘ ball 0 1
    using holomorphic_on_subset holomorphic_on_power
    by (blast intro: holomorphic_on_ident)
  show k holomorphic_on power2 ‘ q ‘ ball 0 1
    using q01 holomorphic_on_subset [OF holk]
    by (force simp: norm_power abs_square_less_1)
qed
have 2: (k ∘ power2 ∘ q) 0 = 0
  using p0 F_def ⟨f ∈ F⟩ ψ01 ψ2 ⟨0 ∈ S⟩ kh qp by force
have 3: norm ((k ∘ power2 ∘ q) z) < 1 if norm z < 1 for z
proof -
  have norm ((power2 ∘ q) z) < 1
    using that q01 by (force simp: norm_power abs_square_less_1)
  with k01 show ?thesis
    by fastforce
qed
have False if c: ∀z. norm z < 1 ⟶ (k ∘ power2 ∘ q) z = c * z and norm
c = 1 for c
proof -
  have c ≠ 0 using that by auto
  have norm (p(1/2)) < 1 norm (p(-1/2)) < 1
    using p01 by force+
  then have (k ∘ power2 ∘ q) (p(1/2)) = c * p(1/2) (k ∘ power2 ∘ q)
(p(-1/2)) = c * p(-1/2)
    using c by force+
  then have p (1/2) = p (- (1/2))
    by (auto simp: ⟨c ≠ 0⟩ qp o_def)
  then have q (p (1/2)) = q (p (- (1/2)))
    by simp
  then have 1/2 = - (1/2::complex)
    by (auto simp: qp)
  then show False
    by simp
qed
moreover
have False if norm (deriv (k ∘ power2 ∘ q) 0) ≠ 1 norm (deriv (k ∘ power2
∘ q) 0) ≤ 1
  and le: ⋀ξ. norm ξ < 1 ⟶ norm ((k ∘ power2 ∘ q) ξ) ≤ norm ξ
proof -
  have norm (deriv (k ∘ power2 ∘ q) 0) < 1

```

```

    using that by simp
    moreover have eq: deriv f 0 = deriv (k ∘ (λz. z ^ 2) ∘ q) 0 * deriv (p ∘
ψ ∘ h ∘ f) 0
    proof (intro DERIV_imp_deriv_has_field_derivative_transform_within_open
[OF DERIV_chain])
      show (k ∘ power2 ∘ q has_field_derivative deriv (k ∘ power2 ∘ q) 0) (at
((p ∘ ψ ∘ h ∘ f) 0))
        using 1 holomorphic_derivI p0 by auto
      show (p ∘ ψ ∘ h ∘ f has_field_derivative deriv (p ∘ ψ ∘ h ∘ f) 0) (at 0)
        using ⟨p ∘ ψ ∘ h ∘ f ∈ F⟩ ⟨open S⟩ ⟨0 ∈ S⟩ holF holomorphic_derivI
by blast
      show ∧x. x ∈ S ⇒ (k ∘ power2 ∘ q ∘ (p ∘ ψ ∘ h ∘ f)) x = f x
        using ψ2 f01 kh normψ1 qp by auto
      qed (use assms in simp_all)
      ultimately have cmod (deriv (p ∘ ψ ∘ h ∘ f) 0) ≤ 0
        using le_norm_df0
      by (metis linorder_not_le mult.commute mult_less_cancel_left2 norm_mult)
      moreover have 1 ≤ norm (deriv f 0)
        using normf [of id] by (simp add: idF)
      ultimately show False
        by (simp add: eq)
      qed
      ultimately show ?thesis
        using Schwarz_Lemma [OF 1 2 3] norm_one by blast
      qed
      then show a ∈ f ' S
        by blast
      qed
      then have fS: f ' S = ball 0 1
        using F_def ⟨f ∈ F⟩ by blast
      then have ∀z∈ball 0 1. g z ∈ S ∧ f (g z) = z
        by (metis gf_imageE)
      with fS show ?thesis
        by (metis gf_holf_holg_image_eqI)
      qed

locale SC_Chain =
  fixes S :: complex set
  assumes openS: open S
begin

lemma winding_number_zero:
  assumes simply_connected S
  shows connected S ∧
    (∀γ z. path γ ∧ path_image γ ⊆ S ∧
      pathfinish γ = pathstart γ ∧ z ∉ S → winding_number γ z = 0)
  using assms
  by (auto simp: simply_connected_imp_connected simply_connected_imp_winding_number_zero)

```

lemma *contour_integral_zero*:

assumes *valid_path* *g* *path_image* $g \subseteq S$ *pathfinish* $g = \text{pathstart } g$ *f* *holomorphic_on* *S*

$\bigwedge \gamma z. \llbracket \text{path } \gamma; \text{path_image } \gamma \subseteq S; \text{pathfinish } \gamma = \text{pathstart } \gamma; z \notin S \rrbracket \implies \text{winding_number } \gamma z = 0$

shows $(f \text{ has_contour_integral } 0) g$

using *assms* **by** (*meson* *Cauchy_theorem_global* *openS* *valid_path_imp_path*)

lemma *global_primitive*:

assumes *connected* *S* **and** *holf*: *f* *holomorphic_on* *S*

and *prev*: $\bigwedge \gamma f. \llbracket \text{valid_path } \gamma; \text{path_image } \gamma \subseteq S; \text{pathfinish } \gamma = \text{pathstart } \gamma; f \text{ holomorphic_on } S \rrbracket \implies (f \text{ has_contour_integral } 0) \gamma$

shows $\exists h. \forall z \in S. (h \text{ has_field_derivative } f z) (\text{at } z)$

proof (*cases* $S = \{\}$)

case *True* **then show** *?thesis*

by *simp*

next

case *False*

then obtain *a* **where** $a \in S$

by *blast*

show *?thesis*

proof (*intro* *exI* *ballI*)

fix *x* **assume** $x \in S$

then obtain *d* **where** $d > 0$ **and** *d*: $\text{cball } x d \subseteq S$

using *openS* *open_contains_cball_eq* **by** *blast*

let $?g = \lambda z. (\text{SOME } g. \text{polynomial_function } g \wedge \text{path_image } g \subseteq S \wedge \text{pathstart } g = a \wedge \text{pathfinish } g = z)$

show $(\lambda z. \text{contour_integral } (?g z) f) \text{ has_field_derivative } f x$
(*at* *x*)

proof (*simp* *add*: *has_field_derivative_def* *has_derivative_at2* *bounded_linear_mult_right*, *rule* *Lim_transform*)

show $(\lambda y. \text{inverse}(\text{norm}(y - x)) *_{\mathbb{R}} (\text{contour_integral}(\text{linepath } x y) f - f x * (y - x))) -x \rightarrow 0$

proof (*clarsimp* *simp* *add*: *Lim_at*)

fix *e*:*real* **assume** $e > 0$

moreover have *continuous* (*at* *x*) *f*

using *openS* $\langle x \in S \rangle$ *holf* *continuous_on_eq_continuous_at* *holomorphic_on_imp_continuous_on* **by** *auto*

ultimately obtain *d1* **where** $d1 > 0$

and *d1*: $\bigwedge x'. \text{dist } x' x < d1 \implies \text{dist } (f x') (f x) < e/2$

unfolding *continuous_at_eps_delta*

by (*metis* *less_divide_eq_numeral1*(1) *mult_zero_left*)

obtain *d2* **where** $d2 > 0$ **and** *d2*: $\text{ball } x d2 \subseteq S$

using *openS* $\langle x \in S \rangle$ *open_contains_ball_eq* **by** *blast*

have $\text{inverse}(\text{norm}(y - x)) * \text{norm}(\text{contour_integral}(\text{linepath } x y) f - f x * (y - x)) < e$

if $0 < d1$ $0 < d2$ $y \neq x$ $\text{dist } y x < d1$ $\text{dist } y x < d2$ **for** *y*

proof -

```

have f contour_integrable_on linepath x y
proof (rule contour_integrable_continuous_linepath [OF continuous_on_subset])
  show continuous_on S f
    by (simp add: holf_holomorphic_on_imp_continuous_on)
  have closed_segment x y  $\subseteq$  ball x d2
    by (meson dist_commute_lessI dist_in_closed_segment le_less_trans
mem_ball_subsetI that(5))
  with d2 show closed_segment x y  $\subseteq$  S
    by blast
qed
then obtain z where z: (f has_contour_integral z) (linepath x y)
  by (force simp: contour_integrable_on_def)
have con: (( $\lambda$ w. f x) has_contour_integral f x * (y - x)) (linepath x y)
  using has_contour_integral_const_linepath [of f x y x] by metis
have norm (z - f x * (y - x))  $\leq$  (e/2) * norm (y - x)
proof (rule has_contour_integral_bound_linepath)
  show (( $\lambda$ w. f w - f x) has_contour_integral z - f x * (y - x)) (linepath
x y)
    by (rule has_contour_integral_diff [OF z con])
  show  $\bigwedge$ w. w  $\in$  closed_segment x y  $\implies$  norm (f w - f x)  $\leq$  e/2
    by (metis d1 dist_norm_less_le_trans not_less not_less_iff_gr_or_eq
segment_bound1 that(4))
qed (use  $\langle e > 0 \rangle$  in auto)
with  $\langle e > 0 \rangle$  have inverse (norm (y - x)) * norm (z - f x * (y - x))
 $\leq$  e/2
  by (simp add: field_split_simps)
also have ...  $<$  e
  using  $\langle e > 0 \rangle$  by simp
finally show ?thesis
  by (simp add: contour_integral_unique [OF z])
qed
with  $\langle d1 > 0 \rangle$   $\langle d2 > 0 \rangle$ 
show  $\exists d > 0. \forall z. z \neq x \wedge \text{dist } z \ x < d \implies$ 
inverse (norm (z - x)) * norm (contour_integral (linepath x z) f -
f x * (z - x))  $<$  e
  by (rule_tac x=min d1 d2 in exI) auto
qed
next
have *: (1 / norm (y - x)) *R (contour_integral (?g y) f -
(contour_integral (?g x) f + f x * (y - x))) =
(contour_integral (linepath x y) f - f x * (y - x)) /R norm (y - x)
if 0  $<$  d y  $\neq$  x and yx: dist y x  $<$  d for y
proof -
  have y  $\in$  S
    by (metis subsetD d dist_commute_less_eq_real_def mem_cball yx)
  have gxy: polynomial_function (?g x)  $\wedge$  path_image (?g x)  $\subseteq$  S  $\wedge$  pathstart
(?g x) = a  $\wedge$  pathfinish (?g x) = x
    polynomial_function (?g y)  $\wedge$  path_image (?g y)  $\subseteq$  S  $\wedge$  pathstart
(?g y) = a  $\wedge$  pathfinish (?g y) = y

```

```

      using someI_ex [OF connected_open_polynomial_connected [OF openS
⟨connected S⟩ ⟨a ∈ S⟩]] ⟨x ∈ S⟩ ⟨y ∈ S⟩
      by meson+
    then have vp: valid_path (?g x) valid_path (?g y)
      by (simp_all add: valid_path_polynomial_function)
      have f0: (f has_contour_integral 0) ((?g x) +++ linepath x y +++
reversepath (?g y))
    proof (rule prev)
      show valid_path ((?g x) +++ linepath x y +++ reversepath (?g y))
        using gxy vp by (auto simp: valid_path_join)
      have closed_segment x y ⊆ cball x d
        using yx by (auto simp: dist_commute dest!: dist_in_closed_segment)
      then have closed_segment x y ⊆ S
        using d by blast
      then show path_image ((?g x) +++ linepath x y +++ reversepath (?g
y)) ⊆ S
        using gxy by (auto simp: path_image_join)
    qed (use gxy holf in auto)
    then have fintxy: f contour_integrable_on linepath x y
      using gxy(2) has_contour_integral_integrable vp by fastforce
    have fintgx: f contour_integrable_on (?g x) f contour_integrable_on (?g y)
      using openS contour_integrable_holomorphic_simple gxy holf vp by blast+
    show ?thesis
      apply (clarsimp simp add: divide_simps)
      using contour_integral_unique [OF f0]
    apply (simp add: fintxy gxy contour_integrable_reversepath contour_integral_reversepath
fintgx vp)
      apply (simp add: algebra_simps)
    done
  qed
  show (λz. (1 / norm (z - x)) *R
(contour_integral (?g z) f - (contour_integral (?g x) f + f x * (z
- x))) -
(contour_integral (linepath x z) f - f x * (z - x)) /R norm (z - x))
-x → 0
    apply (rule tendsto_eventually)
    apply (simp add: eventually_at)
    apply (rule_tac x=d in exI)
    using ⟨d > 0⟩ * by simp
  qed
qed
qed

```

lemma *holomorphic_log*:

```

  assumes connected S and holf: f holomorphic_on S and nz:  $\bigwedge z. z \in S \implies f z \neq 0$ 
  and prev:  $\bigwedge f. f \text{ holomorphic\_on } S \implies \exists h. \forall z \in S. (h \text{ has\_field\_derivative } f z)$ 
  (at z)
  shows  $\exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S. f z = \exp(g z))$ 

```

```

proof –
  have ( $\lambda z. \text{deriv } f z / f z$ ) holomorphic_on S
    by (simp add: openS holf holomorphic_deriv holomorphic_on_divide nz)
  then obtain g where  $g: \bigwedge z. z \in S \implies (g \text{ has\_field\_derivative } \text{deriv } f z / f z)$ 
    (at z)
    using prev [of  $\lambda z. \text{deriv } f z / f z$ ] by metis
  have Df:  $\bigwedge x. x \in S \implies \text{DERIV } f x :> \text{deriv } f x$ 
    using holf holomorphic_derivI openS by force
  have hfd:  $\bigwedge x. x \in S \implies ((\lambda z. \text{exp } (g z) / f z) \text{ has\_field\_derivative } 0)$  (at x)
    by (rule derivative_eq_intros Df g nz | simp) +
  obtain c where  $c: \bigwedge x. x \in S \implies \text{exp } (g x) / f x = c$ 
  proof (rule DERIV_zero_connected_constant [OF <connected S> openS finite.emptyI])
    show continuous_on S ( $\lambda z. \text{exp } (g z) / f z$ )
      by (metis (full_types) openS g continuous_on_divide continuous_on_exp holf
holomorphic_on_imp_continuous_on holomorphic_on_open nz)
    then show  $\forall x \in S - \{c\}. ((\lambda z. \text{exp } (g z) / f z) \text{ has\_field\_derivative } 0)$  (at x)
      using hfd by (blast intro: DERIV_zero_connected_constant [OF <connected
S> openS finite.emptyI, of  $\lambda z. \text{exp } (g z) / f z$ ])
    qed auto
  show ?thesis
  proof (intro exI ballI conjI)
    have g holomorphic_on S
      using openS g holomorphic_on_open by blast
    then show ( $\lambda z. \text{Ln}(\text{inverse } c) + g z$ ) holomorphic_on S
      by (intro holomorphic_intros)
    fix z :: complex
    assume  $z \in S$ 
    then have  $\text{exp } (g z) / c = f z$ 
      by (metis c divide_divide_eq_right exp_not_eq_zero nonzero_mult_div_cancel_left)
    moreover have  $1 / c \neq 0$ 
      using  $\langle z \in S \rangle c \text{ nz}$  by fastforce
    ultimately show  $f z = \text{exp } (\text{Ln } (\text{inverse } c) + g z)$ 
      by (simp add: exp_add_inverse_eq_divide)
    qed
  qed

```

lemma *holomorphic_sqrt*:

```

  assumes holf: f holomorphic_on S and nz:  $\bigwedge z. z \in S \implies f z \neq 0$ 
  and prev:  $\bigwedge f. \llbracket f \text{ holomorphic\_on } S; \forall z \in S. f z \neq 0 \rrbracket \implies \exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S. f z = \text{exp}(g z))$ 
  shows  $\exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S. f z = (g z)^2)$ 
proof –
  obtain g where holg: g holomorphic_on S and  $g: \bigwedge z. z \in S \implies f z = \text{exp } (g z)$ 
    (at z)
    using prev [of f] holf nz by metis
  show ?thesis
  proof (intro exI ballI conjI)
    show ( $\lambda z. \text{exp}(g z/2)$ ) holomorphic_on S
      by (intro holomorphic_intros) (auto simp: holg)

```

```

  show  $\bigwedge z. z \in S \implies f z = (\exp (g z/2))^2$ 
  by (metis (no_types) g exp_double nonzero_mult_div_cancel_left times_divide_eq_right
zero_neq_numeral)
qed
qed

```

lemma *biholomorphic_to_disc*:

```

  assumes connected S and S:  $S \neq \{\}$   $S \neq UNIV$ 
  and prev:  $\bigwedge f. [f \text{ holomorphic\_on } S; \forall z \in S. f z \neq 0] \implies \exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S. f z = (g z)^2)$ 
  shows  $\exists f g. f \text{ holomorphic\_on } S \wedge g \text{ holomorphic\_on ball } 0 \ 1 \wedge$ 
     $(\forall z \in S. f z \in \text{ball } 0 \ 1 \wedge g(f z) = z) \wedge$ 
     $(\forall z \in \text{ball } 0 \ 1. g z \in S \wedge f(g z) = z)$ 

```

proof –

```

  obtain a b where  $a \in S$   $b \notin S$ 
  using S by blast
  then obtain  $\delta$  where  $\delta > 0$  and  $\delta: \text{ball } a \ \delta \subseteq S$ 
  using openS openE by blast
  obtain g where holg:  $g \text{ holomorphic\_on } S$  and eqg:  $\bigwedge z. z \in S \implies z - b = (g z)^2$ 

```

```

  proof (rule exE [OF prev [of  $\lambda z. z - b$ ]])

```

```

    show  $(\lambda z. z - b) \text{ holomorphic\_on } S$ 

```

```

    by (intro holomorphic_intros)

```

```

  qed (use  $\langle b \notin S \rangle$  in auto)

```

```

  have  $\neg g \text{ constant\_on } S$ 

```

proof –

```

  have  $(a + \delta/2) \in \text{ball } a \ \delta$   $a + (\delta/2) \neq a$ 

```

```

  using  $\langle \delta > 0 \rangle$  by (simp_all add: dist_norm)

```

```

  then show ?thesis

```

```

  unfolding constant_on_def

```

```

  using eqg [of a] eqg [of  $a + \delta/2$ ]  $\langle a \in S \rangle \ \delta$ 

```

```

  by (metis diff_add_cancel subset_eq)

```

qed

```

  then have open  $(g \text{ ` ball } a \ \delta)$ 

```

```

  using open_mapping_thm [of  $g$   $S$   $\text{ball } a \ \delta$ , OF holg openS  $\langle \text{connected } S \rangle$ ]  $\delta$  by
blast

```

```

  then obtain r where  $r > 0$  and  $r: \text{ball } (g a) \ r \subseteq (g \text{ ` ball } a \ \delta)$ 

```

```

  by (metis  $\langle 0 < \delta \rangle$  centre_in_ball imageI openE)

```

```

  have  $g\_not\_r: g z \notin \text{ball } (-(g a)) \ r$  if  $z \in S$  for  $z$ 

```

proof

```

  assume  $g z \in \text{ball } (-(g a)) \ r$ 

```

```

  then have  $-g z \in \text{ball } (g a) \ r$ 

```

```

  by (metis add.inverse_inverse dist_minus mem_ball)

```

```

  with r have  $-g z \in (g \text{ ` ball } a \ \delta)$ 

```

```

  by blast

```

```

  then obtain w where  $w: -g z = g w$   $\text{dist } a \ w < \delta$ 

```

```

  by auto

```

```

  with  $\delta$  have  $w \in S$ 

```

```

  by force

```

```

then have w = z
  by (metis diff_add_cancel eqg power_minus_Bit0 that w(1))
then have g z = 0
  using ⟨- g z = g w⟩ by auto
with eqg that ⟨b ∉ S⟩ show False
  by force
qed
then have nz:  $\bigwedge z. z \in S \implies g z + g a \neq 0$ 
  by (metis ⟨0 < r⟩ add commute add_diff_cancel_left' centre_in_ball diff_0)
let ?f =  $\lambda z. (r/3) / (g z + g a) - (r/3) / (g a + g a)$ 
obtain h where holh: h holomorphic_on S and h a = 0 and h01:  $h \text{ ' } S \subseteq \text{ball } 0 \ 1$  and inj_on h S
proof
  show ?f holomorphic_on S
    by (intro holomorphic_intros holg nz)
  have  $\exists: \llbracket \text{norm } x \leq 1/3; \text{norm } y \leq 1/3 \rrbracket \implies \text{norm}(x - y) < 1$  for  $x y::\text{complex}$ 
    using norm_triangle_ineq4 [of x y] by simp
  have  $\text{norm} ((r/3) / (g z + g a) - (r/3) / (g a + g a)) < 1$  if  $z \in S$  for  $z$ 
    apply (rule  $\exists$ )
    unfolding norm_divide
    using ⟨r > 0⟩ g_not_r [OF ⟨z ∈ S⟩] g_not_r [OF ⟨a ∈ S⟩]
    by (simp_all add: field_split_simps dist_commute dist_norm)
  then show ?f ' S  $\subseteq \text{ball } 0 \ 1$ 
    by auto
  show inj_on ?f S
    using ⟨r > 0⟩ eqg apply (clarsimp simp: inj_on_def)
    by (metis diff_add_cancel)
qed auto
obtain k where holk: k holomorphic_on (h ' S)
  and derk:  $\bigwedge z. z \in S \implies \text{deriv } h z * \text{deriv } k (h z) = 1$ 
  and kh:  $\bigwedge z. z \in S \implies k(h z) = z$ 
  using holomorphic_has_inverse [OF holh openS ⟨inj_on h S⟩] by metis

have 1: open (h ' S)
  by (simp add: ⟨inj_on h S⟩ holh openS open_mapping_thm3)
have 2: connected (h ' S)
  by (simp add: connected_continuous_image ⟨connected S⟩ holh holomorphic_on_imp_continuous_on)
have  $\exists: 0 \in h \text{ ' } S$ 
  using ⟨a ∈ S⟩ ⟨h a = 0⟩ by auto
have  $\exists g. g \text{ holomorphic\_on } h \text{ ' } S \wedge (\forall z \in h \text{ ' } S. f z = (g z)^2)$ 
  if holf: f holomorphic_on h ' S and nz:  $\bigwedge z. z \in h \text{ ' } S \implies f z \neq 0$  inj_on f (h ' S) for f
proof -
  obtain g where holg: g holomorphic_on S and eqg:  $\bigwedge z. z \in S \implies (f \circ h) z = (g z)^2$ 
    by (smt (verit) comp_def holf holh holomorphic_on_compose image_eqI nz(1) prev)
  show ?thesis
  proof (intro exI conjI)

```



```

    show  $g \circ k$  holomorphic_on  $h \text{ ' } S$ 
    by (smt (verit) holg holk holomorphic_on_compose holomorphic_on_subset
imageE image_subset_iff kh)
    show  $\forall z \in h \text{ ' } S. f z = ((g \circ k) z)^2$ 
    using eqg kh by auto
  qed
  qed
  obtain  $f g$  where  $f: f$  holomorphic_on  $h \text{ ' } S$  and  $g: g$  holomorphic_on ball 0 1
    and  $gf: \forall z \in h \text{ ' } S. f z \in \text{ball } 0 \ 1 \wedge g (f z) = z$  and  $fg: \forall z \in \text{ball } 0 \ 1. g z \in h$ 
    '  $S \wedge f (g z) = z$ 
    using biholomorphic_to_disc_aux [OF 1 2 3 h01 4] by blast
  show ?thesis
  proof (intro exI conjI)
    show  $f \circ h$  holomorphic_on  $S$ 
    by (simp add: f holh holomorphic_on_compose)
    show  $k \circ g$  holomorphic_on ball 0 1
    by (metis holomorphic_on_subset image_subset_iff fg holk g holomor-
phic_on_compose)
  qed (use fg gf kh in auto)
  qed

```

lemma *homeomorphic_to_disc:*

assumes $S = \text{UNIV} \vee$

$(\exists f g. f$ holomorphic_on $S \wedge g$ holomorphic_on ball 0 1 \wedge
 $(\forall z \in S. f z \in \text{ball } 0 \ 1 \wedge g(f z) = z) \wedge$
 $(\forall z \in \text{ball } 0 \ 1. g z \in S \wedge f(g z) = z))$ (**is** $_ \vee ?P$)

shows S homeomorphic ball (0::complex) 1

by (smt (verit, ccfv_SIG) holomorphic_on_imp_continuous_on homeomorphic_ball01_UNIV
homeomorphic_minimal assms)

lemma *homeomorphic_to_disc_imp_simply_connected:*

assumes $S = \{\}$ $\vee S$ homeomorphic ball (0::complex) 1

shows simply_connected S

using assms homeomorphic_simply_connected_eq convex_imp_simply_connected
by auto

end

proposition

assumes open S

shows simply_connected_eq_winding_number_zero:

simply_connected $S \iff$

connected $S \wedge$

$(\forall g z. \text{path } g \wedge \text{path_image } g \subseteq S \wedge$

$\text{pathfinish } g = \text{pathstart } g \wedge (z \notin S)$

$\implies \text{winding_number } g z = 0)$ (**is** ?wn0)

and simply_connected_eq_contour_integral_zero:

simply_connected $S \iff$

connected $S \wedge$

```

      (∀ g f. valid_path g ∧ path_image g ⊆ S ∧
        pathfinish g = pathstart g ∧ f holomorphic_on S
        → (f has_contour_integral 0) g) (is ?ci0)
and simply_connected_eq_global_primitive:
  simply_connected S ↔
  connected S ∧
  (∀ f. f holomorphic_on S →
    (∃ h. ∀ z. z ∈ S → (h has_field_derivative f z) (at z))) (is ?gp)
and simply_connected_eq_holomorphic_log:
  simply_connected S ↔
  connected S ∧
  (∀ f. f holomorphic_on S ∧ (∀ z ∈ S. f z ≠ 0)
    → (∃ g. g holomorphic_on S ∧ (∀ z ∈ S. f z = exp(g z)))) (is ?log)
and simply_connected_eq_holomorphic_sqrt:
  simply_connected S ↔
  connected S ∧
  (∀ f. f holomorphic_on S ∧ (∀ z ∈ S. f z ≠ 0)
    → (∃ g. g holomorphic_on S ∧ (∀ z ∈ S. f z = (g z)2))) (is ?sqrt)
and simply_connected_eq_biholomorphic_to_disc:
  simply_connected S ↔
  S = {} ∨ S = UNIV ∨
  (∃ f g. f holomorphic_on S ∧ g holomorphic_on ball 0 1 ∧
    (∀ z ∈ S. f z ∈ ball 0 1 ∧ g(f z) = z) ∧
    (∀ z ∈ ball 0 1. g z ∈ S ∧ f(g z) = z)) (is ?bih)
and simply_connected_eq_homeomorphic_to_disc:
  simply_connected S ↔ S = {} ∨ S homeomorphic ball (0::complex) 1
(is ?disc)
proof -
  interpret SC_Chain
  using assms by (simp add: SC_Chain_def)
  have ?wn0 ∧ ?ci0 ∧ ?gp ∧ ?log ∧ ?sqrt ∧ ?bih ∧ ?disc
proof -
  have *: [[α ⇒ β; β ⇒ γ; γ ⇒ δ; δ ⇒ ζ; ζ ⇒ η; η ⇒ ϑ; ϑ ⇒ ξ; ξ ⇒
α]]
    ⇒ (α ↔ β) ∧ (α ↔ γ) ∧ (α ↔ δ) ∧ (α ↔ ζ) ∧
      (α ↔ η) ∧ (α ↔ ϑ) ∧ (α ↔ ξ) for α β γ δ ζ η ϑ ξ
  by blast
show ?thesis
  apply (rule *)
  using winding_number_zero apply metis
  using contour_integral_zero apply metis
  using global_primitive apply metis
  using holomorphic_log apply metis
  using holomorphic_sqrt apply simp
  using biholomorphic_to_disc apply blast
  using homeomorphic_to_disc apply blast
  using homeomorphic_to_disc_imp_simply_connected apply blast
done
qed

```

```

then show ?wn0 ?ci0 ?gp ?log ?sqrt ?bih ?disc
  by safe
qed

```

```

corollary contractible_eqSimplyConnected2d:
  fixes S :: complex set
  assumes open S
  shows contractible S  $\longleftrightarrow$  simply_connected S
proof
  show contractible S  $\implies$  simply_connected S
    by (simp add: contractible_impSimplyConnected)
  show simply_connected S  $\implies$  contractible S
    using assms convex_impContractible homeomorphicContractibleEq
      simply_connectedEqHomeomorphicToDisc by auto
qed

```

7.3 A further chain of equivalences about components of the complement of a simply connected set

(following 1.35 in Burckel'S book)

```

context SC_Chain
begin

```

```

lemma frontier_properties:
  assumes simply_connected S
  shows if bounded S then connected(frontier S)
    else  $\forall C \in \text{components}(\text{frontier } S). \neg \text{bounded } C$ 
proof -
  have S = {}  $\vee$  S homeomorphic ball (0::complex) 1
    using simply_connectedEqHomeomorphicToDisc assms openS by blast
  then show ?thesis
  proof
    assume S = {}
    then show ?thesis
      by simp
  next
    assume S01: S homeomorphic ball (0::complex) 1
    then obtain g f
      where gim: g ' S = ball 0 1 and fg:  $\bigwedge x. x \in S \implies f(g x) = x$ 
        and fim: f ' ball 0 1 = S and gf:  $\bigwedge y. \text{cmod } y < 1 \implies g(f y) = y$ 
        and contg: continuous_on S g and contf: continuous_on (ball 0 1) f
        by (fastforce simp: homeomorphism_def homeomorphic_def)
    define D where D  $\equiv \lambda n. \text{ball } (0::\text{complex}) (1 - 1/(\text{of\_nat } n + 2))$ 
    define A where A  $\equiv \lambda n. \{z::\text{complex}. 1 - 1/(\text{of\_nat } n + 2) < \text{norm } z \wedge$ 
norm z < 1}
    define X where X  $\equiv \lambda n::\text{nat}. \text{closure}(f ' A n)$ 
    have D01: D n  $\subseteq$  ball 0 1 for n
      by (simp add: D_def ball_subset_ball_iff)
    have A01: A n  $\subseteq$  ball 0 1 for n

```

```

    by (auto simp: A_def)
  have cloX: closed(X n) for n
    by (simp add: X_def)
  have Xsubclo: X n  $\subseteq$  closure S for n
    unfolding X_def by (metis A01 closure_mono fim image_mono)
  have connected (A n) for n
    using connected_annulus [of _ 0::complex] by (simp add: A_def)
  then have connX: connected(X n) for n
    unfolding X_def
  by (metis A01 connected_continuous_image connected_imp_connected_closure
contf continuous_on_subset)
  have nestX: X n  $\subseteq$  X m if m  $\leq$  n for m n
  proof -
    have 1 - 1 / (real m + 2)  $\leq$  1 - 1 / (real n + 2)
      using that by (auto simp: field_simps)
    then show ?thesis
      by (auto simp: X_def A_def intro!: closure_mono)
  qed
  have closure S - S  $\subseteq$  ( $\bigcap$  n. X n)
  proof
    fix x
    assume x  $\in$  closure S - S
    then have x  $\in$  closure S x  $\notin$  S by auto
    show x  $\in$  ( $\bigcap$  n. X n)
  proof
    fix n
    have ball 0 1 = closure (D n)  $\cup$  A n
      by (auto simp: D_def A_def le_less_trans)
    with fim have Seq: S = f ' (closure (D n))  $\cup$  f ' (A n)
      by (simp add: image_Un)
    have continuous_on (closure (D n)) f
      by (simp add: D_def cball_subset_ball_iff continuous_on_subset [OF
contf])
    moreover have compact (closure (D n))
      by (simp add: D_def)
    ultimately have clo_fim: closed (f ' closure (D n))
      using compact_continuous_image compact_imp_closed by blast
    have *: (f ' cball 0 (1 - 1 / (real n + 2)))  $\subseteq$  S
      by (force simp: D_def Seq)
    show x  $\in$  X n
      using Seq X_def  $\langle$ x  $\in$  closure S $\rangle$   $\langle$ x  $\notin$  S $\rangle$  clo_fim by fastforce
  qed
  qed
  moreover have ( $\bigcap$  n. X n)  $\subseteq$  closure S - S
  proof -
    have ( $\bigcap$  n. X n)  $\subseteq$  closure S
      using Xsubclo by blast
    moreover have ( $\bigcap$  n. X n)  $\cap$  S  $\subseteq$  {}
    proof (clarify, clarsimp simp: X_def fim [symmetric])

```

```

fix  $x$  assume  $x$  [rule_format]:  $\forall n. f\ x \in \text{closure } (f\ ' A\ n)$  and  $\text{cmod } x < 1$ 
then obtain  $n$  where  $n: 1 / (1 - \text{norm } x) < \text{of\_nat } n$ 
  using reals_Archimedean2 by blast
with  $\langle \text{cmod } x < 1 \rangle$  gr0I have  $XX: 1 / \text{of\_nat } n < 1 - \text{norm } x$  and  $n > 0$ 
  by (fastforce simp: field_split_simps algebra_simps)+
have  $f\ x \in f\ ' (D\ n)$ 
  using  $n$   $\langle \text{cmod } x < 1 \rangle$  by (auto simp: field_split_simps algebra_simps
D_def)
moreover have  $f\ ' D\ n \cap \text{closure } (f\ ' A\ n) = \{\}$ 
proof  $-$ 
  have inj_on  $f$   $(D\ n)$ 
    unfolding inj_on_def using D01 by (metis gf mem_ball_0 subsetCE)
  then have op_fDn: open( $f\ ' (D\ n)$ )
  by (metis invariance_of_domain D_def Elementary_Metric_Spaces.open_ball
continuous_on_subset [OF contf D01])
have inj: inj_on  $f$   $(\text{ball } 0\ 1)$ 
  by (metis mem_ball_0 inj_on_def gf)
have  $D\ n \cup A\ n \subseteq \text{ball } 0\ 1$ 
  using D01 A01 by simp
moreover have  $D\ n \cap A\ n = \{\}$ 
  by (auto simp: D_def A_def)
ultimately have  $f\ ' D\ n \cap f\ ' A\ n = \{\}$ 
  by (metis A01 D01 image_is_empty inj_on_image_Int injf)
then show ?thesis
  by (simp add: open_Int_closure_eq_empty [OF op_fDn])
qed
ultimately show False
  using  $x$  [of n] by blast
qed
ultimately
show  $(\bigcap n. X\ n) \subseteq \text{closure } S - S$ 
  using closure_subset disjoint_iff_not_equal by blast
qed
ultimately have  $\text{closure } S - S = (\bigcap n. X\ n)$  by blast
then have frontierS:  $\text{frontier } S = (\bigcap n. X\ n)$ 
  by (simp add: frontier_def openS interior_open)
show ?thesis
proof (cases bounded S)
  case True
  have bouX: bounded  $(X\ n)$  for  $n$ 
    by (meson True Xsubclo bounded_closure bounded_subset)
  have compaX: compact  $(X\ n)$  for  $n$ 
    by (simp add: bouX cloX compact_eq_bounded_closed)
  have connected  $(\bigcap n. X\ n)$ 
    by (metis nestX compaX connX connected_nest)
  then show ?thesis
    by (simp add: True frontierS = (\bigcap n. X n))
  next

```

```

case False
have unboundedX:  $\neg$  bounded(X n) for n
proof
  assume bXn: bounded(X n)
  have continuous_on (cball 0 ( $1 - 1 / (2 + \text{real } n)$ )) f
    by (simp add: cball_subset_ball_iff continuous_on_subset [OF contf])
  then have bounded (f ' cball 0 ( $1 - 1 / (2 + \text{real } n)$ )))
    by (simp add: compact_imp_bounded [OF compact_continuous_image])
  moreover have bounded (f ' A n)
    by (auto simp: X_def closure_subset image_subset_iff bounded_subset
      [OF bXn])
  ultimately have bounded (f ' (cball 0 ( $1 - 1/(2 + \text{real } n)$ )  $\cup$  A n))
    by (simp add: image_Un)
  then have bounded (f ' ball 0 1)
    apply (rule bounded_subset)
    apply (auto simp: A_def algebra_simps)
  done
then show False
  using False by (simp add: fim [symmetric])
qed
have clo_INTX: closed( $\bigcap$ (range X))
  by (metis cloX closed_INT)
then have lcX: locally compact ( $\bigcap$ (range X))
  by (metis closed_imp_locally_compact)
have False if C: C  $\in$  components (frontier S) and boC: bounded C for C
proof –
  have closed C
    by (metis C closed_components frontier_closed)
  then have compact C
    by (metis boC compact_eq_bounded_closed)
  have Cco: C  $\in$  components ( $\bigcap$ (range X))
    by (metis frontierS C)
  obtain K where C  $\subseteq$  K compact K
    and Ksub: K  $\subseteq$   $\bigcap$ (range X) and clo: closed( $\bigcap$ (range X) – K)
  proof (cases {k. C  $\subseteq$  k  $\wedge$  compact k  $\wedge$  openin (top_of_set ( $\bigcap$ (range X)))
k} = {})
    case True
      then show ?thesis
        using Sura_Bura [OF lcX Cco  $\langle$ compact C $\rangle$ ] boC
          by (simp add: True)
    next
      case False
        then obtain L where compact L C  $\subseteq$  L and K: openin (top_of_set
          ( $\bigcap$ x. X x)) L
          by blast
        show ?thesis
          proof
            show L  $\subseteq$   $\bigcap$ (range X)
              by (metis K openin_imp_subset)

```

```

      show closed ( $\bigcap(\text{range } X) - L$ )
        by (metis closedin_diff closedin_self closedin_closed_trans [OF _
clo_INTX] K)
      qed (use ⟨compact L⟩ ⟨ $C \subseteq L$ ⟩ in auto)
    qed
  obtain U V where open U open V and compact (closure U)
    and V:  $\bigcap(\text{range } X) - K \subseteq V$  and U:  $K \subseteq U$   $U \cap V = \{\}$ 
    by (metis Diff_disjoint separation_normal_compact [OF ⟨compact K⟩
clo])
  then have  $U \cap (\bigcap(\text{range } X) - K) = \{\}$ 
    by blast
  have (closure U - U)  $\cap (\bigcap n. X n \cap \text{closure } U) \neq \{\}$ 
  proof (rule compact_imp_fip)
    show compact (closure U - U)
      by (metis ⟨compact (closure U)⟩ ⟨open U⟩ compact_diff)
    show  $\bigwedge T. T \in \text{range } (\lambda n. X n \cap \text{closure } U) \implies \text{closed } T$ 
      by clarify (metis cloX closed_Int closed_closure)
    show (closure U - U)  $\cap \bigcap \mathcal{F} \neq \{\}$ 
      if finite  $\mathcal{F}$  and  $\mathcal{F}: \mathcal{F} \subseteq \text{range } (\lambda n. X n \cap \text{closure } U)$  for  $\mathcal{F}$ 
    proof
      assume empty: (closure U - U)  $\cap \bigcap \mathcal{F} = \{\}$ 
      obtain J where finite J and J:  $\mathcal{F} = (\lambda n. X n \cap \text{closure } U) \text{ ` } J$ 
        using finite_subset_image [OF ⟨finite  $\mathcal{F}$ ⟩  $\mathcal{F}$ ] by auto
      show False
    proof (cases J =  $\{\}$ )
      case True
        with J_empty have closed U
          by (simp add: closure_subset_eq)
        have  $C \neq \{\}$ 
          using C_in_components_nonempty by blast
        then have  $U \neq \{\}$ 
          using ⟨ $K \subseteq U$ ⟩ ⟨ $C \subseteq K$ ⟩ by blast
        moreover have  $U \neq \text{UNIV}$ 
          using ⟨compact (closure U)⟩ by auto
        ultimately show False
          using ⟨open U⟩ ⟨closed U⟩ clopen by blast
      next
        case False
          define j where j  $\equiv \text{Max } J$ 
          have j  $\in J$ 
            by (simp add: False ⟨finite J⟩ j_def)
          have jmax:  $\bigwedge m. m \in J \implies m \leq j$ 
            by (simp add: j_def ⟨finite J⟩)
          have  $\bigcap ((\lambda n. X n \cap \text{closure } U) \text{ ` } J) = X j \cap \text{closure } U$ 
            using False jmax nestX ⟨j  $\in J$ ⟩ by auto
          then have XU:  $X j \cap \text{closure } U = X j \cap U$ 
            using J_closure_subset_empty by fastforce
          then have openin (top_of_set (X j)) (X j  $\cap \text{closure } U$ )
            by (simp add: openin_open_Int ⟨open U⟩)

```

```

moreover have closedin (top_of_set (X j)) (X j  $\cap$  closure U)
  by (simp add: closedin_closed_Int)
moreover have X j  $\cap$  closure U  $\neq$  X j
  by (metis unboundedX  $\langle$ compact (closure U) $\rangle$  bounded_subset
compact_eq_bounded_closed_inf.order_iff)
moreover have X j  $\cap$  closure U  $\neq$  {}
  by (metis Cco Ksub UNIV_I  $\langle$ C  $\subseteq$  K $\rangle$   $\langle$ K  $\subseteq$  U $\rangle$  XU bot.extremum_uniqueI
in_components_maximal le_INF_iff le_inf_iff)
ultimately show False
  using connX [of j] by (force simp: connected_clopen)
qed
qed
qed
moreover have ( $\bigcap$  n. X n  $\cap$  closure U) = ( $\bigcap$  n. X n)  $\cap$  closure U
  by blast
moreover have x  $\in$  U if  $\bigwedge$  n. x  $\in$  X n x  $\in$  closure U for x
  by (metis Diff_iff INT_I U V  $\langle$ open V $\rangle$  closure_iff_nhds_not_empty
order.refl subsetD that)
ultimately show False
  by (auto simp: open_Int_closure_eq_empty [OF  $\langle$ open V $\rangle$ , of U])
qed
then show ?thesis
  by (auto simp: False)
qed
qed
qed

```

lemma unbounded_complement_components:

```

assumes C: C  $\in$  components ( $-$  S) and S: connected S
and prev: if bounded S then connected(frontier S)
  else  $\forall$  C  $\in$  components(frontier S).  $\neg$  bounded C
shows  $\neg$  bounded C
proof (cases bounded S)
case True
with prev have S  $\neq$  UNIV and confr: connected(frontier S)
  by auto
obtain w where C_ccsw: C = connected_component_set ( $-$  S) w and w  $\notin$  S
  using C by (auto simp: components_def)
show ?thesis
proof (cases S = {})
  case True with C show ?thesis by auto
next
  case False
  show ?thesis
proof
  assume bounded C
  then have outside C  $\neq$  {}
  using outside_bounded_nonempty by metis

```



```

    then obtain z where z:  $\neg$  bounded (connected_component_set (- C) z)
and z  $\notin$  C
  by (auto simp: outside_def)
  have clo_ccs: closed (connected_component_set (- S) x) for x
  by (simp add: closed_Compl closed_connected_component openS)
  have connected_component_set (- S) w = connected_component_set (- S)
z
  proof (rule joinable_connected_component_eq [OF confr])
  show frontier S  $\subseteq$  - S
    using openS by (auto simp: frontier_def interior_open)
  have False if connected_component_set (- S) w  $\cap$  frontier (- S) = {}
  proof -
  have C  $\cap$  frontier S = {}
    using that by (simp add: C_ccsw)
  moreover have closed C
    using C_ccsw clo_ccs by blast
  ultimately show False
    by (metis C  $\langle$ S  $\neq$  {} $\rangle$   $\langle$ S  $\neq$  UNIV $\rangle$  C_ccsw bot_eq_sup_iff connected_component_eq_UNIV frontier_Int_closed
        frontier_closed frontier_complement frontier_eq_empty frontier_of_components_subset in_components_maximal inf.orderE)
  qed
  then show connected_component_set (- S) w  $\cap$  frontier S  $\neq$  {}
  by auto
  have *:  $\llbracket$ frontier C  $\subseteq$  C; frontier C  $\subseteq$  F; frontier C  $\neq$  {} $\rrbracket \implies C \cap F \neq$ 
{} for C F::complex set
  by blast
  have connected_component_set (- S) z  $\cap$  frontier (- S)  $\neq$  {}
  proof (rule *)
  show frontier (connected_component_set (- S) z)  $\subseteq$  connected_component_set
(- S) z
    by (auto simp: closed_Compl closed_connected_component frontier_def
openS)
  show frontier (connected_component_set (- S) z)  $\subseteq$  frontier (- S)
    using frontier_of_connected_component_subset by fastforce
  have connected (closure S - S)
    by (metis confr frontier_def interior_open openS)
  moreover have  $\neg$  bounded (-S)
    by (simp add: True cobounded_imp_unbounded)
  moreover have bounded (connected_component_set (- S) w)
    using C_ccsw  $\langle$ bounded C $\rangle$  by auto
  ultimately have z  $\notin$  S
    using  $\langle$ w  $\notin$  S $\rangle$  openS
  by (metis ComplI Compl_eq_Diff_UNIV connected_UNIV closed_closure
closure_subset
        connected_component_eq_self connected_diff_open_from_closed
subset_UNIV)
  then have connected_component_set (- S) z  $\neq$  {}
    by (metis ComplI connected_component_eq_empty)

```

```

      then show frontier (connected_component_set (- S) z) ≠ {}
        by (metis False ⟨S ≠ UNIV⟩ connected_component_eq_UNIV frontier_complement frontier_eq_empty)
    qed
    then show connected_component_set (- S) z ∩ frontier S ≠ {}
      by auto
    qed
    then show False
      by (metis C_ccsw Compl_iff ⟨w ∉ S⟩ ⟨z ∉ C⟩ connected_component_eq_empty connected_component_idemp)
    qed
  qed
next
  case False
  obtain w where C_ccsw: C = connected_component_set (- S) w and w ∉ S
    using C by (auto simp: components_def)
  have frontier (connected_component_set (- S) w) ⊆ connected_component_set (- S) w
    by (simp add: closed_Compl closed_connected_component frontier_subset_eq openS)
  moreover have frontier (connected_component_set (- S) w) ⊆ frontier S
    using frontier_complement frontier_of_connected_component_subset by blast
  moreover have frontier (connected_component_set (- S) w) ≠ {}
    by (metis C C_ccsw False bounded_empty_compl_top_eq connected_component_eq_UNIV double_compl frontier_not_empty_in_components_nonempty)
  ultimately obtain z where zin: z ∈ frontier S and z: z ∈ connected_component_set (- S) w
    by blast
  have connected_component_set (frontier S) z ∈ components(frontier S)
    by (simp add: ⟨z ∈ frontier S⟩ componentsI)
  with prev False have ¬ bounded (connected_component_set (frontier S) z)
    by simp
  moreover have connected_component (- S) w = connected_component (- S) z
    using connected_component_eq [OF z] by force
  ultimately show ?thesis
    by (metis C_ccsw SC_Chain.openS SC_Chain_axioms bounded_subset closed_Compl connected_component_mono frontier_complement frontier_subset_eq)
  qed

```

lemma empty_inside:

```

  assumes connected S ∧ C. C ∈ components (- S) ⇒ ¬ bounded C
  shows inside S = {}
  using assms by (auto simp: components_def inside_def)

```

lemma empty_inside_imp_simply_connected:

```

  [[connected S; inside S = {}]] ⇒ simply_connected S
  by (metis ComplI inside_Un_outside openS outside_mono simply_connected_eq_winding_number_zero subsetCE sup_bot.left_neutral winding_number_zero_in_outside)

```

end

proposition

fixes $S :: \text{complex set}$

assumes $\text{open } S$

shows $\text{simply_connected_eq_frontier_properties}$:

$\text{simply_connected } S \longleftrightarrow$

$\text{connected } S \wedge$

$(\text{if bounded } S \text{ then connected}(\text{frontier } S)$

$\text{else } (\forall C \in \text{components}(\text{frontier } S). \neg \text{bounded } C)) \text{ (is ?fp)}$

and $\text{simply_connected_eq_unbounded_complement_components}$:

$\text{simply_connected } S \longleftrightarrow$

$\text{connected } S \wedge (\forall C \in \text{components}(- S). \neg \text{bounded } C) \text{ (is ?ucc)}$

and $\text{simply_connected_eq_empty_inside}$:

$\text{simply_connected } S \longleftrightarrow$

$\text{connected } S \wedge \text{inside } S = \{\} \text{ (is ?ei)}$

proof –

interpret SC_Chain

using assms by $(\text{simp add: } SC_Chain_def)$

have $?fp \wedge ?ucc \wedge ?ei$

using $\text{empty_inside empty_inside_imp_simply_connected_frontier_properties}$

$\text{unbounded_complement_components winding_number_zero}$ by blast

then show $?fp ?ucc ?ei$

by blast+

qed

lemma $\text{simply_connected_iff_simple}$:

fixes $S :: \text{complex set}$

assumes $\text{open } S \text{ bounded } S$

shows $\text{simply_connected } S \longleftrightarrow \text{connected } S \wedge \text{connected}(- S) \text{ (is ?lhs = ?rhs)}$

proof

show $?lhs \implies ?rhs$

by $(\text{metis DIM_complex assms cobounded_has_bounded_component double_complement dual_order.order_iff_strict}$

$\text{simply_connected_eq_unbounded_complement_components})$

show $?rhs \implies ?lhs$

by $(\text{simp add: assms connected_frontier_simple simply_connected_eq_frontier_properties})$

qed

lemma $\text{subset_simply_connected_imp_inside_subset}$:

fixes $A :: \text{complex set}$

assumes $\text{simply_connected } A \text{ open } A \text{ } B \subseteq A$

shows $\text{inside } B \subseteq A$

by $(\text{metis assms Diff_eq_empty_iff_inside_mono subset_empty simply_connected_eq_empty_inside})$

7.4 Further equivalences based on continuous logs and sqrts

context *SC_Chain*

begin

lemma *continuous_log*:

fixes $f :: \text{complex} \Rightarrow \text{complex}$

assumes $S: \text{simply_connected } S$

and $\text{contf}: \text{continuous_on } S f$ and $\text{nz}: \bigwedge z. z \in S \implies f z \neq 0$

shows $\exists g. \text{continuous_on } S g \wedge (\forall z \in S. f z = \exp(g z))$

proof –

consider $S = \{ \} \mid S \text{ homeomorphic ball } (0::\text{complex}) 1$

using *simply_connected_eq_homeomorphic_to_disc* [*OF openS*] S by *metis*

then show *?thesis*

proof *cases*

case 1 then show *?thesis*

by *simp*

next

case 2

then obtain $h k :: \text{complex} \Rightarrow \text{complex}$

where $kh: \bigwedge x. x \in S \implies k(h x) = x$ and $him: h ' S = \text{ball } 0 1$

and $\text{conth}: \text{continuous_on } S h$

and $hk: \bigwedge y. y \in \text{ball } 0 1 \implies h(k y) = y$ and $kim: k ' \text{ball } 0 1 = S$

and $\text{contk}: \text{continuous_on } (\text{ball } 0 1) k$

unfolding *homeomorphism_def* *homeomorphic_def* by *metis*

obtain g where $\text{contg}: \text{continuous_on } (\text{ball } 0 1) g$

and $\text{expg}: \bigwedge z. z \in \text{ball } 0 1 \implies (f \circ k) z = \exp(g z)$

proof (*rule continuous_logarithm_on_ball*)

show $\text{continuous_on } (\text{ball } 0 1) (f \circ k)$

using *contf* *continuous_on_compose* *contk* *kim* by *blast*

show $\bigwedge z. z \in \text{ball } 0 1 \implies (f \circ k) z \neq 0$

using *kim* *nz* by *auto*

qed *auto*

then show *?thesis*

by (*metis comp_apply* *conth* *continuous_on_compose* *him* *imageI* *kh*)

qed

qed

lemma *continuous_sqrt*:

fixes $f :: \text{complex} \Rightarrow \text{complex}$

assumes $\text{contf}: \text{continuous_on } S f$ and $\text{nz}: \bigwedge z. z \in S \implies f z \neq 0$

and $\text{prev}: \bigwedge f::\text{complex} \Rightarrow \text{complex}.$

$\llbracket \text{continuous_on } S f; \bigwedge z. z \in S \implies f z \neq 0 \rrbracket$

$\implies \exists g. \text{continuous_on } S g \wedge (\forall z \in S. f z = \exp(g z))$

shows $\exists g. \text{continuous_on } S g \wedge (\forall z \in S. f z = (g z)^2)$

proof –

obtain g where $\text{contg}: \text{continuous_on } S g$ and $\text{geq}: \bigwedge z. z \in S \implies f z = \exp(g z)$

z)

using *contf* *nz* *prev* by *metis*

show *?thesis*

```

proof (intro exI ballI conjI)
  show continuous_on S (λz. exp(g z/2))
    by (intro continuous_intros) (auto simp: contg)
  show  $\bigwedge z. z \in S \implies f z = (\exp (g z/2))^2$ 
    by (metis (no_types, lifting) divide_inverse exp_double geq mult.left_commute
mult.right_neutral right_inverse zero_neq numeral)
qed
qed

lemma continuous_sqrt_imp_simply_connected:
  assumes connected S
    and prev:  $\bigwedge f::\text{complex} \implies \text{complex}. \llbracket \text{continuous\_on } S f; \forall z \in S. f z \neq 0 \rrbracket$ 
       $\implies \exists g. \text{continuous\_on } S g \wedge (\forall z \in S. f z = (g z)^2)$ 
  shows simply_connected S
proof (clarsimp simp add: simply_connected_eq_holomorphic_sqrt [OF openS]
⟨connected S⟩)
  fix f
  assume f holomorphic_on S and nz:  $\forall z \in S. f z \neq 0$ 
  then obtain g where contg: continuous_on S g and geq:  $\bigwedge z. z \in S \implies f z =$ 
   $(g z)^2$ 
    by (metis holomorphic_on_imp_continuous_on prev)
  show  $\exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S. f z = (g z)^2)$ 
proof (intro exI ballI conjI)
  show g holomorphic_on S
    proof (clarsimp simp add: holomorphic_on_open [OF openS])
      fix z
      assume z ∈ S
      with nz geq have g z ≠ 0
        by auto
      obtain δ where 0 < δ  $\bigwedge w. \llbracket w \in S; \text{dist } w z < \delta \rrbracket \implies \text{dist } (g w) (g z) <$ 
       $cmod (g z)$ 
        using contg [unfolded continuous_on_iff] by (metis ⟨g z ≠ 0⟩ ⟨z ∈ S⟩
zero_less_norm_iff)
      then have δ:  $\bigwedge w. \llbracket w \in S; w \in \text{ball } z \delta \rrbracket \implies g w + g z \neq 0$ 
        by (metis add.commute add_cancel_right_left dist_commute dist_complex_def
mem_ball
norm_increases_online norm_not_less_zero norm_zero order_less_asym)
      have *:  $(\lambda x. (f x - f z) / (x - z) / (g x + g z)) -z \rightarrow \text{deriv } f z / (g z + g z)$ 
        proof (intro tendsto_intros)
          show  $(\lambda x. (f x - f z) / (x - z)) -z \rightarrow \text{deriv } f z$ 
            using ⟨f holomorphic_on S⟩ ⟨z ∈ S⟩ has_field_derivative_iff holomor-
phic_derivI openS by blast
          show g -z → g z
            using ⟨z ∈ S⟩ contg continuous_on_eq_continuous_at isCont_def openS
by blast
          qed (simp add: ⟨g z ≠ 0⟩)
      then have (g has_field_derivative deriv f z / (g z + g z)) (at z)
        unfolding has_field_derivative_iff
        proof (rule Lim_transform_within_open)

```

```

show open (ball z  $\delta \cap S$ )
  by (simp add: openS open_Int)
show z  $\in$  ball z  $\delta \cap S$ 
  using ⟨z  $\in S$ ⟩ ⟨0 <  $\delta$ ⟩ by simp
show  $\bigwedge x. [x \in \text{ball } z \ \delta \cap S; x \neq z]$ 
   $\implies (f x - f z) / (x - z) / (g x + g z) = (g x - g z) / (x - z)$ 
  using  $\delta$  ⟨z  $\in S$ ⟩
  apply (simp add: geq_field_split_simps power2_eq_square)
  by (metis distrib_left mult_cancel_right)
qed
then show  $\exists f'. (g \text{ has\_field\_derivative } f') (at z) ..$ 
qed
qed (use geq in auto)
qed
end

```

proposition**fixes** $S :: \text{complex set}$ **assumes** open S **shows** simply_connected_eq_continuous_log:simply_connected $S \longleftrightarrow$ connected $S \wedge$ $(\forall f :: \text{complex} \Rightarrow \text{complex}. \text{continuous_on } S f \wedge (\forall z \in S. f z \neq 0)$ $\longrightarrow (\exists g. \text{continuous_on } S g \wedge (\forall z \in S. f z = \exp (g z))))$ (**is** ?log)**and** simply_connected_eq_continuous_sqrt:simply_connected $S \longleftrightarrow$ connected $S \wedge$ $(\forall f :: \text{complex} \Rightarrow \text{complex}. \text{continuous_on } S f \wedge (\forall z \in S. f z \neq 0)$ $\longrightarrow (\exists g. \text{continuous_on } S g \wedge (\forall z \in S. f z = (g z)^2)))$ (**is** ?sqrt)**proof** –**interpret** SC_Chain**using** asms **by** (simp add: SC_Chain_def)**show** ?log ?sqrt**using** local.continuous_log winding_number_zero continuous_sqrt continuous_sqrt_imp_simply_connected**by** auto**qed****7.5 More Borsukian results****lemma** Borsukian_componentwise_eq:**fixes** $S :: 'a :: \text{euclidean_space set}$ **assumes** S : locally_connected $S \vee$ compact S **shows** Borsukian $S \longleftrightarrow (\forall C \in \text{components } S. \text{Borsukian } C)$ **proof** –**have** *: ANR($-\{0 :: \text{complex}\}$)**by** (simp add: ANR_delete open_Comp open_imp_ANR)**show** ?thesis

using *cohomotopically_trivial_on_components* [*OF assms **] **by** (*auto simp: Borsukian_alt*)
qed

lemma *Borsukian_componentwise*:
fixes $S :: 'a::\text{euclidean_space set}$
assumes *locally connected S* \vee *compact S* \wedge $C \in \text{components } S \implies \text{Borsukian } C$
shows *Borsukian S*
by (*metis Borsukian_componentwise_eq assms*)

lemma *simply_connected_eq_Borsukian*:
fixes $S :: \text{complex set}$
shows *open S* \implies (*simply_connected S* \longleftrightarrow *connected S* \wedge *Borsukian S*)
by (*auto simp: simply_connected_eq_continuous_log Borsukian_continuous_logarithm*)

lemma *Borsukian_eq_simply_connected*:
fixes $S :: \text{complex set}$
shows *open S* $\implies \text{Borsukian } S \longleftrightarrow (\forall C \in \text{components } S. \text{simply_connected } C)$
by (*meson Borsukian_componentwise_eq_in_components_connected_open_components open_imp_locally_connected_simply_connected_eq_Borsukian*)

lemma *Borsukian_separation_open_closed*:
fixes $S :: \text{complex set}$
assumes $S: \text{open } S \vee \text{closed } S$ **and** *bounded S*
shows *Borsukian S* \longleftrightarrow *connected(- S)*
using S
proof
assume *open S*
show *?thesis*
unfolding *Borsukian_eq_simply_connected* [*OF* $\langle \text{open } S \rangle$]
by (*metis* $\langle \text{open } S \rangle$ $\langle \text{bounded } S \rangle$ *bounded_subset_in_components_maximal_non-separation_by_component_eq_open_components_simply_connected_iff_simple*)
next
assume *closed S*
with $\langle \text{bounded } S \rangle$ **show** *?thesis*
by (*simp add: Borsukian_separation_compact_compact_eq_bounded_closed*)
qed

7.6 Finally, the Riemann Mapping Theorem

theorem *Riemann_mapping_theorem*:
 $\text{open } S \wedge \text{simply_connected } S \longleftrightarrow$
 $S = \{\}$ $\vee S = \text{UNIV}$ \vee
 $(\exists f g. f \text{ holomorphic_on } S \wedge g \text{ holomorphic_on ball } 0 \ 1 \wedge$
 $(\forall z \in S. f z \in \text{ball } 0 \ 1 \wedge g(f z) = z) \wedge$
 $(\forall z \in \text{ball } 0 \ 1. g z \in S \wedge f(g z) = z))$
(is $_ = ?rhs$ **)**
proof –

```

have simply_connected S  $\longleftrightarrow$  ?rhs if open S
  by (simp add: simply_connected_eq_biholomorphic_to_disc that)
moreover have open S if ?rhs
proof -
  { fix f g
    assume g: g holomorphic_on ball 0 1  $\forall z \in \text{ball } 0 \ 1. g \ z \in S \wedge f (g \ z) = z$ 
      and  $\forall z \in S. \text{cmod } (f \ z) < 1 \wedge g (f \ z) = z$ 
    then have S = g ' (ball 0 1)
      by force
    then have open S
      by (metis open_ball g inj_on_def open_mapping_thm3)
  }
  with that show open S by auto
qed
ultimately show ?thesis by metis
qed

```

7.7 Applications to Winding Numbers

```

lemma simply_connected_inside_simple_path:
  fixes p :: real  $\Rightarrow$  complex
  shows simple_path p  $\implies$  simply_connected(inside(path_image p))
  using Jordan_inside_outside_connected_simple_path_image_inside_simple_curve_imp_closed
  simply_connected_eq_frontier_properties
  by fastforce

```

```

lemma simply_connected_Int:
  fixes S :: complex set
  assumes open S open T simply_connected S simply_connected T connected (S
 $\cap$  T)
  shows simply_connected (S  $\cap$  T)
  using assms by (force simp: simply_connected_eq_winding_number_zero open_Int)

```

7.8 The winding number defines a continuous logarithm for the path itself

```

lemma winding_number_as_continuous_log:
  assumes path p and  $\zeta: \zeta \notin \text{path\_image } p$ 
  obtains q where path q
    pathfinish q - pathstart q =  $2 * \text{of\_real } \pi * i * \text{winding\_number } p \ \zeta$ 
     $\wedge t. t \in \{0..1\} \implies p \ t = \zeta + \exp(q \ t)$ 
proof -
  let ?q =  $\lambda t. 2 * \text{of\_real } \pi * i * \text{winding\_number}(\text{subpath } 0 \ t \ p) \ \zeta + \text{Ln}(\text{pathstart } p - \zeta)$ 
  show ?thesis
  proof
    have *: continuous (at t within {0..1}) ( $\lambda x. \text{winding\_number}(\text{subpath } 0 \ x \ p) \ \zeta$ )
    if t:  $t \in \{0..1\}$  for t

```



```

proof –
  let ?B = ball (p t) (norm(p t - ζ))
  have p t ≠ ζ
    using path_image_def that ζ by blast
  then have simply_connected ?B
    by (simp add: convex_imp_simply_connected)
  then have  $\bigwedge f::\text{complex} \Rightarrow \text{complex. continuous\_on } ?B f \wedge (\forall \zeta \in ?B. f \zeta \neq$ 
0)
     $\longrightarrow (\exists g. \text{continuous\_on } ?B g \wedge (\forall \zeta \in ?B. f \zeta = \text{exp } (g \zeta)))$ 
    by (simp add: simply_connected_eq_continuous_log)
  moreover have continuous_on ?B ( $\lambda w. w - \zeta$ )
    by (intro continuous_intros)
  moreover have ( $\forall z \in ?B. z - \zeta \neq 0$ )
    by (auto simp: dist_norm)
  ultimately obtain g where contg: continuous_on ?B g
    and geg:  $\bigwedge z. z \in ?B \Longrightarrow z - \zeta = \text{exp } (g z)$  by blast
  obtain d where 0 < d and d:
     $\bigwedge x. \llbracket x \in \{0..1\}; \text{dist } x t < d \rrbracket \Longrightarrow \text{dist } (p x) (p t) < \text{cmod } (p t - \zeta)$ 
    using  $\langle \text{path } p \rangle t$  unfolding path_def continuous_on_iff
    by (metis  $\langle p t \neq \zeta \rangle$  right_minus_eq zero_less_norm_iff)
  have ( $(\lambda x. \text{winding\_number } (\lambda w. \text{subpath } 0 x p w - \zeta) 0 -$ 
     $\text{winding\_number } (\lambda w. \text{subpath } 0 t p w - \zeta) 0) \longrightarrow 0$ )
    (at t within {0..1})
  proof (rule Lim_transform_within [OF _  $\langle d > 0 \rangle$ ])
    have continuous (at t within {0..1}) (g o p)
    proof (rule continuous_within_compose)
      show continuous (at t within {0..1}) p
        using  $\langle \text{path } p \rangle$  continuous_on_eq_continuous_within path_def that by
blast
      show continuous (at (p t) within p ‘ {0..1}) g
        by (metis (no_types, lifting) open_ball UNIV_I  $\langle p t \neq \zeta \rangle$  centre_in_ball
contg continuous_on_eq_continuous_at continuous_within_topological_right_minus_eq
zero_less_norm_iff)
    qed
    with LIM_zero have ( $(\lambda u. (g (\text{subpath } t u p 1) - g (\text{subpath } t u p 0)))$ 
 $\longrightarrow 0$ ) (at t within {0..1})
      by (auto simp: subpath_def continuous_within_o_def)
    then show ( $(\lambda u. (g (\text{subpath } t u p 1) - g (\text{subpath } t u p 0)) / (2 * \text{of\_real}$ 
pi * i))  $\longrightarrow 0$ )
      (at t within {0..1})
      by (simp add: tendsto_divide_zero)
    show  $(g (\text{subpath } t u p 1) - g (\text{subpath } t u p 0)) / (2 * \text{of\_real } \text{pi} * i) =$ 
     $\text{winding\_number } (\lambda w. \text{subpath } 0 u p w - \zeta) 0 - \text{winding\_number } (\lambda w.$ 
subpath 0 t p w - ζ) 0
      if u ∈ {0..1} 0 < dist u t dist u t < d for u
    proof –
      have closed_segment t u ⊆ {0..1}
      using closed_segment_eq_real_ivl t that by auto
      then have  $\bigwedge r. \llbracket r \in \text{closed\_segment } t u \rrbracket \Longrightarrow \text{dist } (p t) (p r) < \text{cmod } (p$ 

```

```

t - ζ)
  by (smt (verit, best) d dist_commute dist_in_closed_segment subsetD
⟨dist u t < d⟩)
  then have piB: path_image(subpath t u p) ⊆ ?B
  by (auto simp: path_image_subpath_gen)
  have *: path (g ∘ subpath t u p)
  proof (rule path_continuous_image)
    show path (subpath t u p)
    using ⟨path p⟩ t that by auto
    show continuous_on (path_image (subpath t u p)) g
    using piB contg continuous_on_subset by blast
  qed
  have (g (subpath t u p 1) - g (subpath t u p 0)) / (2 * of_real pi * i)
    = winding_number (exp ∘ g ∘ subpath t u p) 0
  using winding_number_compose_exp [OF *]
  by (simp add: pathfinish_def pathstart_def o_assoc)
  also have ... = winding_number (λw. subpath t u p w - ζ) 0
  proof (rule winding_number_cong)
    have exp(g y) = y - ζ if y ∈ (subpath t u p) ‘{0..1} for y
    by (metis that geq_path_image_def piB subset_eq)
    then show ∧x. [0 ≤ x; x ≤ 1] ⇒ (exp ∘ g ∘ subpath t u p) x = subpath
t u p x - ζ
    by auto
  qed
  also have ... = winding_number (λw. subpath 0 u p w - ζ) 0 -
winding_number (λw. subpath 0 t p w - ζ) 0
  apply (simp add: winding_number_offset [symmetric])
  using winding_number_subpath_combine [OF ⟨path p⟩ ζ, of 0 t u] ⟨t ∈
{0..1}⟩ ⟨u ∈ {0..1}⟩
  by (simp add: add commute eq_diff_eq)
  finally show ?thesis .
  qed
  qed
  then show ?thesis
  by (subst winding_number_offset) (simp add: continuous_within LIM_zero_iff)
  qed
  show path ?q
  unfolding path_def
  by (intro continuous_intros) (simp add: continuous_on_eq_continuous_within
*)

have ζ ≠ p 0
  by (metis ζ pathstart_def pathstart_in_path_image)
  then show pathfinish ?q - pathstart ?q = 2 * of_real pi * i * winding_number
p ζ
  by (simp add: pathfinish_def pathstart_def)
  show p t = ζ + exp (?q t) if t ∈ {0..1} for t
  proof -
    have path (subpath 0 t p)

```

```

    using ‹path p› that by auto
  moreover
  have  $\zeta \notin \text{path\_image } (\text{subpath } 0 \ t \ p)$ 
  using  $\zeta$  [unfolded path_image_def] that by (auto simp: path_image_subpath)
  ultimately show ?thesis
    using winding_number_exp_2pi [of subpath 0 t p  $\zeta$ ] ‹ $\zeta \neq p \ 0$ ›
    by (auto simp: exp_add algebra_simps pathfinish_def pathstart_def sub-
path_def)
  qed
  qed
  qed

```

7.9 Winding number equality is the same as path/loop homotopy in $\mathbb{C} - 0$

lemma *winding_number_homotopic_loops_null_eq*:

```

  assumes path p and  $\zeta: \zeta \notin \text{path\_image } p$ 
  shows  $\text{winding\_number } p \ \zeta = 0 \iff (\exists a. \text{homotopic\_loops } (-\{\zeta\}) \ p \ (\lambda t. a))$ 
    (is ?lhs = ?rhs)

```

proof

```

  assume [simp]: ?lhs
  obtain q where path q
    and qeq:  $\text{pathfinish } q - \text{pathstart } q = 2 * \text{of\_real } \pi * i * \text{winding\_number } p \ \zeta$ 
    and peq:  $\bigwedge t. t \in \{0..1\} \implies p \ t = \zeta + \text{exp}(q \ t)$ 
  using winding_number_as_continuous_log [OF assms] by blast
  have *:  $\text{homotopic\_with\_canon } (\lambda r. \text{pathfinish } r = \text{pathstart } r) \ \{0..1\} \ (-\{\zeta\}) \ ((\lambda w. \zeta + \text{exp } w) \circ q) \ ((\lambda w. \zeta + \text{exp } w) \circ (\lambda t. 0))$ 
  proof (rule homotopic_with_compose_continuous_left)
    show  $\text{homotopic\_with\_canon } (\lambda f. \text{pathfinish } ((\lambda w. \zeta + \text{exp } w) \circ f) = \text{pathstart } ((\lambda w. \zeta + \text{exp } w) \circ f)) \ \{0..1\} \ \text{UNIV } q \ (\lambda t. 0)$ 
  proof (rule homotopic_with_mono, simp_all add: pathfinish_def pathstart_def)
    have  $\text{homotopic\_loops } \text{UNIV } q \ (\lambda t. 0)$ 
    by (rule homotopic_loops_linear) (use qeq ‹path q› in ‹auto simp: path_defs›)
    then have  $\text{homotopic\_with } (\lambda r. r \ 1 = r \ 0) \ (\text{top\_of\_set } \{0..1\}) \ \text{euclidean } q \ (\lambda t. 0)$ 
    by (simp add: homotopic_loops_def pathfinish_def pathstart_def)
    then show  $\text{homotopic\_with } (\lambda h. \text{exp } (h \ 1) = \text{exp } (h \ 0)) \ (\text{top\_of\_set } \{0..1\}) \ \text{euclidean } q \ (\lambda t. 0)$ 
    by (rule homotopic_with_mono) simp
  qed
  show  $\text{continuous\_on } \text{UNIV } (\lambda w. \zeta + \text{exp } w)$ 
  by (rule continuous_intros)+
  qed auto
  then have  $\text{homotopic\_with\_canon } (\lambda r. \text{pathfinish } r = \text{pathstart } r) \ \{0..1\} \ (-\{\zeta\}) \ p \ (\lambda x. \zeta + 1)$ 
  by (rule homotopic_with_eq) (auto simp: o_def peq pathfinish_def pathstart_def)
  then have  $\text{homotopic\_loops } (-\{\zeta\}) \ p \ (\lambda t. \zeta + 1)$ 

```

```

    by (simp add: homotopic_loops_def)
  then show ?rhs ..
next
  assume ?rhs
  then obtain a where homotopic_loops  $(-\{\zeta\})$  p  $(\lambda t. a)$  ..
  then have winding_number p  $\zeta =$  winding_number  $(\lambda t. a)$   $\zeta$   $a \neq \zeta$ 
    using winding_number_homotopic_loops homotopic_loops_imp_subset by
(force simp:)+
  then show ?lhs
    using winding_number_zero_const by auto
qed

```

lemma *winding_number_homotopic_paths_null_explicit_eq*:

```

  assumes path p and  $\zeta: \zeta \notin$  path_image p
  shows winding_number p  $\zeta = 0 \iff$  homotopic_paths  $(-\{\zeta\})$  p (linepath
(pathstart p) (pathstart p))
    (is ?lhs = ?rhs)

```

proof

```

  assume ?lhs
  then show ?rhs
    using homotopic_loops_imp_homotopic_paths_null
    by (force simp: linepath_refl winding_number_homotopic_loops_null_eq [OF
assms])
next
  assume ?rhs
  then show ?lhs
    by (metis  $\zeta$  pathstart_in_path_image winding_number_homotopic_paths wind-
ing_number_trivial)
qed

```

lemma *winding_number_homotopic_paths_null_eq*:

```

  assumes path p and  $\zeta: \zeta \notin$  path_image p
  shows winding_number p  $\zeta = 0 \iff$   $(\exists a. \text{homotopic\_paths } (-\{\zeta\})$  p  $(\lambda t. a))$ 
    (is ?lhs = ?rhs)

```

proof

```

  assume ?lhs
  then show ?rhs
    by (auto simp: winding_number_homotopic_paths_null_explicit_eq [OF assms]
linepath_refl)
next
  assume ?rhs
  then show ?lhs
    by (metis  $\zeta$  homotopic_paths_imp_pathfinish pathfinish_def pathfinish_in_path_image
winding_number_homotopic_paths winding_number_zero_const)
qed

```

proposition *winding_number_homotopic_paths_eq*:

```

  assumes path p and  $\zeta p: \zeta \notin$  path_image p
    and path q and  $\zeta q: \zeta \notin$  path_image q

```

```

    and qp: pathstart q = pathstart p pathfinish q = pathfinish p
    shows winding_number p  $\zeta$  = winding_number q  $\zeta$   $\longleftrightarrow$  homotopic_paths
(-{ $\zeta$ }) p q
    (is ?lhs = ?rhs)
proof
  assume ?lhs
  then have winding_number (p +++ reversepath q)  $\zeta$  = 0
    using assms by (simp add: winding_number_join winding_number_reversepath)
  moreover
  have path (p +++ reversepath q)  $\zeta$   $\notin$  path_image (p +++ reversepath q)
    using assms by (auto simp: not_in_path_image_join)
  ultimately obtain a where homotopic_paths (- { $\zeta$ }) (p +++ reversepath q)
(linepath a a)
    using winding_number_homotopic_paths_null_explicit_eq by blast
  then show ?rhs
    using homotopic_paths_imp_pathstart assms
    by (fastforce simp: dest: homotopic_paths_imp_homotopic_loops homotopic_paths_loop_parts)
qed (simp add: winding_number_homotopic_paths)

```

lemma winding_number_homotopic_loops_eq:

```

  assumes path p and  $\zeta$ p:  $\zeta$   $\notin$  path_image p
    and path q and  $\zeta$ q:  $\zeta$   $\notin$  path_image q
    and loops: pathfinish p = pathstart p pathfinish q = pathstart q
  shows winding_number p  $\zeta$  = winding_number q  $\zeta$   $\longleftrightarrow$  homotopic_loops
(-{ $\zeta$ }) p q
  (is ?lhs = ?rhs)

```

proof

```

  assume L: ?lhs
  have pathstart p  $\neq$   $\zeta$  pathstart q  $\neq$   $\zeta$ 
    using  $\zeta$ p  $\zeta$ q by blast+
  moreover have path_connected (-{ $\zeta$ })
    by (simp add: path_connected_punctured_universe)
  ultimately obtain r where path r and rim: path_image r  $\subseteq$  -{ $\zeta$ }
    and pas: pathstart r = pathstart p and paf: pathfinish r =
pathstart q
    by (auto simp: path_connected_def)
  then have pathstart r  $\neq$   $\zeta$  by blast
  have homotopic_loops (- { $\zeta$ }) p (r +++ q +++ reversepath r)
proof (rule homotopic_paths_imp_homotopic_loops)
  have path (r +++ q +++ reversepath r)
    by (simp add:  $\langle$ path r $\rangle$   $\langle$ path q $\rangle$  loops paf)
  moreover have  $\zeta$   $\notin$  path_image (r +++ q +++ reversepath r)
    by (metis  $\zeta$ q not_in_path_image_join path_image_reversepath rim subset_Cmpl_singleton)
  moreover have homotopic_loops (- { $\zeta$ }) (r +++ q +++ reversepath r) q
    using  $\langle$ path q $\rangle$   $\langle$ path r $\rangle$   $\zeta$ q homotopic_loops_conjugate loops(2) paf rim by
blast
  ultimately show homotopic_paths (- { $\zeta$ }) p (r +++ q +++ reversepath r)
    using loops pathfinish_join pathfinish_reversepath pathstart_join

```

```

    by (metis L ζ p ⟨path p⟩ pas winding_number_homotopic_loops winding_number_homotopic_paths_eq)
  qed (use loops pas in auto)
  moreover have homotopic_loops (- {ζ}) (r +++ q +++ reversepath r) q
    using rim ζ q by (auto simp: homotopic_loops_conjugate paf ⟨path q⟩ ⟨path r⟩
loops)
  ultimately show ?rhs
    using homotopic_loops_trans by metis
  qed (simp add: winding_number_homotopic_loops)

end
theory Complex_Singularities
  imports Conformal_Mappings
begin

```

7.10 Non-essential singular points

```

lemma at_to_0': NO_MATCH 0 z ⇒ at z = filtermap (λx. x + z) (at 0)
  for z :: 'a::real_normed_vector
  by (rule at_to_0)

```

```

lemma nhds_to_0: nhds (x :: 'a :: real_normed_vector) = filtermap ((+) x) (nhds
0)

```

```

proof -
  have (λxa. xa - - x) = (+) x
    by auto
  thus ?thesis
    using filtermap_nhds_shift[of -x 0] by simp
qed

```

```

lemma nhds_to_0': NO_MATCH 0 x ⇒ nhds (x :: 'a :: real_normed_vector)
= filtermap ((+) x) (nhds 0)
  by (rule nhds_to_0)

```

definition

```

is_pole :: ('a::topological_space ⇒ 'b::real_normed_vector) ⇒ 'a ⇒ bool
where is_pole f a = (LIM x (at a). f x :> at_infinity)

```

```

lemma is_pole_cong:
  assumes eventually (λx. f x = g x) (at a) a=b
  shows is_pole f a ⟷ is_pole g b
  unfolding is_pole_def using assms by (intro filterlim_cong, auto)

```

```

lemma is_pole_transform:
  assumes is_pole f a eventually (λx. f x = g x) (at a) a=b
  shows is_pole g b
  using is_pole_cong assms by auto

```

```

lemma is_pole_shift_iff:

```

```

fixes f :: ('a::real_normed_vector  $\Rightarrow$  'b::real_normed_vector)
shows is_pole (f  $\circ$  (+) d) a = is_pole f (a + d)
by (metis add_diff_cancel_right' filterlim_shift_iff is_pole_def)

lemma is_pole_tendsto:
fixes f :: ('a::topological_space  $\Rightarrow$  'b::real_normed_div_algebra)
shows is_pole f x  $\Longrightarrow$  ((inverse o f)  $\longrightarrow$  0) (at x)
unfolding is_pole_def
by (auto simp add: filterlim_inverse_at_iff[symmetric] comp_def filterlim_at)

lemma is_pole_shift_0:
fixes f :: ('a::real_normed_vector  $\Rightarrow$  'b::real_normed_vector)
shows is_pole f z  $\longleftrightarrow$  is_pole ( $\lambda x. f (z + x)$ ) 0
unfolding is_pole_def by (subst at_to_0) (auto simp: filterlim_filtermap_add_ac)

lemma is_pole_shift_0':
fixes f :: ('a::real_normed_vector  $\Rightarrow$  'b::real_normed_vector)
shows NO_MATCH 0 z  $\Longrightarrow$  is_pole f z  $\longleftrightarrow$  is_pole ( $\lambda x. f (z + x)$ ) 0
by (metis is_pole_shift_0)

lemma is_pole_compose_iff:
assumes filtermap g (at x) = (at y)
shows is_pole (f  $\circ$  g) x  $\longleftrightarrow$  is_pole f y
unfolding is_pole_def filterlim_def filtermap_compose assms ..

lemma is_pole_inverse_holomorphic:
assumes open s
  and f_holo: f holomorphic_on (s - {z})
  and pole: is_pole f z
  and non_z:  $\forall x \in s - \{z\}. f x \neq 0$ 
shows ( $\lambda x. \text{if } x=z \text{ then } 0 \text{ else inverse } (f x)$ ) holomorphic_on s
proof -
  define g where g  $\equiv \lambda x. \text{if } x=z \text{ then } 0 \text{ else inverse } (f x)$ 
  have isCont g z unfolding isCont_def using is_pole_tendsto[OF pole]
    by (simp add: g_def cong: LIM_cong)
  moreover have continuous_on (s - {z}) f using f_holo holomorphic_on_imp_continuous_on
by auto
  hence continuous_on (s - {z}) (inverse o f) unfolding comp_def
    by (auto elim!: continuous_on_inverse simp add: non_z)
  hence continuous_on (s - {z}) g unfolding g_def
    using continuous_on_eq by fastforce
  ultimately have continuous_on s g using open_delete[OF <open s>] <open s>
    by (auto simp add: continuous_on_eq_continuous_at)
  moreover have (inverse o f) holomorphic_on (s - {z})
    unfolding comp_def using f_holo
    by (auto elim!: holomorphic_on_inverse simp add: non_z)
  hence g holomorphic_on (s - {z})
    using g_def holomorphic_transform by force
  ultimately show ?thesis unfolding g_def using <open s>

```

by (auto elim!: no_isolated_singularity)
qed

lemma not_is_pole_holomorphic:

assumes open A x ∈ A f holomorphic_on A

shows ¬is_pole f x

proof –

have continuous_on A f

by (intro holomorphic_on_imp_continuous_on) fact

with assms have isCont f x

by (simp add: continuous_on_eq_continuous_at)

hence f -x → f x

by (simp add: isCont_def)

thus ¬is_pole f x

unfolding is_pole_def

using not_tendsto_and_filterlim_at_infinity[of at x f f x] by auto

qed

lemma is_pole_inverse_power: $n > 0 \implies \text{is_pole } (\lambda z::\text{complex}. 1 / (z - a) ^ n) a$

unfolding is_pole_def inverse_eq_divide [symmetric]

by (intro filterlim_compose[OF filterlim_inverse_at_infinity] tendsto_intros)

(auto simp: filterlim_at_eventually_at intro!: exI[of _ 1] tendsto_eq_intros)

lemma is_pole_cmult_iff [simp]:

assumes $c \neq 0$

shows $\text{is_pole } (\lambda z. c * f z :: 'a :: \text{real_normed_field}) z \longleftrightarrow \text{is_pole } f z$

proof

assume is_pole $(\lambda z. c * f z) z$

with $\langle c \neq 0 \rangle$ have is_pole $(\lambda z. \text{inverse } c * (c * f z)) z$

unfolding is_pole_def

by (force intro: tendsto_mult_filterlim_at_infinity)

thus is_pole f z

using $\langle c \neq 0 \rangle$ by (simp add: field_simps)

next

assume is_pole f z

with $\langle c \neq 0 \rangle$ show is_pole $(\lambda z. c * f z) z$

by (auto intro!: tendsto_mult_filterlim_at_infinity simp: is_pole_def)

qed

lemma is_pole_uminus_iff [simp]: $\text{is_pole } (\lambda z. -f z :: 'a :: \text{real_normed_field}) z \longleftrightarrow \text{is_pole } f z$

using is_pole_cmult_iff[of -1 f] by (simp del: is_pole_cmult_iff)

lemma is_pole_inverse: $\text{is_pole } (\lambda z::\text{complex}. 1 / (z - a)) a$

using is_pole_inverse_power[of 1 a] by simp

lemma is_pole_divide:

fixes $f :: 'a :: t2_space \Rightarrow 'b :: \text{real_normed_field}$


```

assumes isCont f z filterlim g (at 0) (at z) f z ≠ 0
shows is_pole (λz. f z / g z) z
proof -
  have filterlim (λz. f z * inverse (g z)) at_infinity (at z)
    using assms filterlim_compose filterlim_inverse_at_infinity isCont_def
      tendsto_mult_filterlim_at_infinity by blast
  thus ?thesis by (simp add: field_split_simps is_pole_def)
qed

lemma is_pole_basic:
  assumes f holomorphic_on A open A z ∈ A f z ≠ 0 n > 0
  shows is_pole (λw. f w / (w-z) ^ n) z
proof (rule is_pole_divide)
  have continuous_on A f by (rule holomorphic_on_imp_continuous_on) fact
  with assms show isCont f z by (auto simp: continuous_on_eq_continuous_at)
  have filterlim (λw. (w-z) ^ n) (nhds 0) (at z)
    using assms by (auto intro!: tendsto_eq_intros)
  thus filterlim (λw. (w-z) ^ n) (at 0) (at z)
    by (intro filterlim_atI tendsto_eq_intros)
    (use assms in ⟨auto simp: eventually_at_filter⟩)
qed fact+

lemma is_pole_basic':
  assumes f holomorphic_on A open A 0 ∈ A f 0 ≠ 0 n > 0
  shows is_pole (λw. f w / w ^ n) 0
  using is_pole_basic[of f A 0] assms by simp

lemma is_pole_compose:
  assumes is_pole f w g -z → w eventually (λz. g z ≠ w) (at z)
  shows is_pole (λx. f (g x)) z
  using assms(1) unfolding is_pole_def
  by (rule filterlim_compose) (use assms in ⟨auto simp: filterlim_at⟩)

lemma is_pole_plus_const_iff:
  is_pole f z ↔ is_pole (λx. f x + c) z
proof
  assume is_pole f z
  then have filterlim f at_infinity (at z) unfolding is_pole_def .
  moreover have ((λ_. c) → c) (at z) by auto
  ultimately have LIM x (at z). f x + c :> at_infinity
    using tendsto_add_filterlim_at_infinity'[of f at z] by auto
  then show is_pole (λx. f x + c) z unfolding is_pole_def .
next
  assume is_pole (λx. f x + c) z
  then have filterlim (λx. f x + c) at_infinity (at z)
    unfolding is_pole_def .
  moreover have ((λ_. -c) → -c) (at z) by auto
  ultimately have LIM x (at z). f x :> at_infinity
    using tendsto_add_filterlim_at_infinity'[of (λx. f x + c)

```

```

      at z ( $\lambda \_. - c$ ) - c]
    by auto
  then show is_pole f z unfolding is_pole_def .
qed

```

```

lemma is_pole_minus_const_iff:
  is_pole ( $\lambda x. f x - c$ ) z  $\longleftrightarrow$  is_pole f z
  using is_pole_plus_const_iff [of f z -c] by simp

```

```

lemma is_pole_alt:
  is_pole f x = ( $\forall B > 0. \exists U. \text{open } U \wedge x \in U \wedge (\forall y \in U. y \neq x \longrightarrow \text{norm } (f y) \geq B)$ )
  unfolding is_pole_def
  unfolding filterlim_at_infinity[of 0,simplified] eventually_at_topological
  by auto

```

```

lemma is_pole_mult_analytic_nonzero1:
  assumes is_pole g x f analytic_on {x}  $f x \neq 0$ 
  shows is_pole ( $\lambda x. f x * g x$ ) x
  unfolding is_pole_def
proof (rule tendsto_mult_filterlim_at_infinity)
  show  $f -x \rightarrow f x$ 
  using assms by (simp add: analytic_at_imp_isCont isContD)
qed (use assms in  $\langle$ auto simp: is_pole_def $\rangle$ )

```

```

lemma is_pole_mult_analytic_nonzero2:
  assumes is_pole f x g analytic_on {x}  $g x \neq 0$ 
  shows is_pole ( $\lambda x. f x * g x$ ) x
proof -
  have g: g analytic_on {x}
  using assms by auto
  show ?thesis
  using is_pole_mult_analytic_nonzero1 [OF  $\langle$ is_pole f x $\rangle$  g]  $\langle$ g x  $\neq 0$  $\rangle$ 
  by (simp add: mult.commute)
qed

```

```

lemma is_pole_mult_analytic_nonzero1_iff:
  assumes f analytic_on {x}  $f x \neq 0$ 
  shows is_pole ( $\lambda x. f x * g x$ ) x  $\longleftrightarrow$  is_pole g x
proof
  assume is_pole g x
  thus is_pole ( $\lambda x. f x * g x$ ) x
  by (intro is_pole_mult_analytic_nonzero1 assms)
next
  assume is_pole ( $\lambda x. f x * g x$ ) x
  hence is_pole ( $\lambda x. \text{inverse } (f x) * (f x * g x)$ ) x
  by (rule is_pole_mult_analytic_nonzero1)
  (use assms in  $\langle$ auto intro!: analytic_intros $\rangle$ )
  also have ?this  $\longleftrightarrow$  is_pole g x
  proof (rule is_pole_cong)

```

```

have eventually ( $\lambda x. f x \neq 0$ ) (at x)
  using assms by (simp add: analytic_at_neq_imp_eventually_neq)
thus eventually ( $\lambda x. \text{inverse } (f x) * (f x * g x) = g x$ ) (at x)
  by eventually_elim auto
qed auto
finally show is_pole g x .
qed

```

```

lemma is_pole_mult_analytic_nonzero2_iff:
  assumes g analytic_on {x} g x  $\neq 0$ 
  shows is_pole ( $\lambda x. f x * g x$ ) x  $\longleftrightarrow$  is_pole f x
  by (subst mult.commute, rule is_pole_mult_analytic_nonzero1_iff) (fact assms)+

```

```

lemma frequently_const_imp_not_is_pole:

```

```

  fixes z :: 'a::first_countable_topology
  assumes frequently ( $\lambda w. f w = c$ ) (at z)
  shows  $\neg$  is_pole f z

```

```

proof

```

```

  assume is_pole f z
  from assms have z islimpt {w. f w = c}
    by (simp add: islimpt_conv_frequently_at)
  then obtain g where g:  $\bigwedge n. g n \in \{w. f w = c\} - \{z\}$  g  $\longrightarrow$  z
    unfolding islimpt_sequential by blast
  then have (f  $\circ$  g)  $\longrightarrow$  c
    by (simp add: tendsto_eventually)
  moreover have filterlim g (at z) sequentially
    using g by (auto simp: filterlim_at)
  then have filterlim (f  $\circ$  g) at_infinity sequentially
    unfolding o_def
    using  $\langle$ is_pole f z $\rangle$  filterlim_compose is_pole_def by blast
  ultimately show False
    using not_tendsto_and_filterlim_at_infinity trivial_limit_sequentially by

```

```

blast

```

```

qed

```

7.11 Isolated singularities

The proposition $\exists x. f -z \rightarrow x \vee \text{is_pole } f z$ can be interpreted as the complex function f has a non-essential singularity at z (i.e. the singularity is either removable or a pole).

```

definition not_essential:: [complex  $\Rightarrow$  complex, complex]  $\Rightarrow$  bool where
  not_essential f z = ( $\exists x. f -z \rightarrow x \vee \text{is\_pole } f z$ )

```

```

definition isolated_singularity_at:: [complex  $\Rightarrow$  complex, complex]  $\Rightarrow$  bool where
  isolated_singularity_at f z = ( $\exists r > 0. f$  analytic_on ball z r - {z})

```

```

lemma not_essential_cong:

```

```

  assumes eventually ( $\lambda x. f x = g x$ ) (at z) z = z'
  shows not_essential f z  $\longleftrightarrow$  not_essential g z'

```

unfolding *not_essential_def* **using** *assms filterlim_cong is_pole_cong* **by** *fast-force*

lemma *not_essential_compose_iff*:
assumes *filtermap g (at z) = at z'*
shows *not_essential (f ∘ g) z = not_essential f z'*
unfolding *not_essential_def filterlim_def filtermap_compose assms is_pole_compose_iff [OF assms]*
by *blast*

lemma *isolated_singularity_at_cong*:
assumes *eventually (λx. f x = g x) (at z) z = z'*
shows *isolated_singularity_at f z ↔ isolated_singularity_at g z'*
proof –
have *isolated_singularity_at g z*
if *isolated_singularity_at f z eventually (λx. f x = g x) (at z)* **for** *f g*
proof –
from *that(1)* **obtain** *r* **where** *r: r > 0 f analytic_on ball z r - {z}*
by *(auto simp: isolated_singularity_at_def)*
from *that(2)* **obtain** *r'* **where** *r': r' > 0 ∀x∈ball z r' - {z}. f x = g x*
unfolding *eventually_at_filter eventually_nhds_metric* **by** *(auto simp: dist_commute)*

have *f holomorphic_on ball z r - {z}*
using *r(2)* **by** *(subst (asm) analytic_on_open) auto*
hence *f holomorphic_on ball z (min r r') - {z}*
by *(rule holomorphic_on_subset) auto*
also have *?this ↔ g holomorphic_on ball z (min r r') - {z}*
using *r'* **by** *(intro holomorphic_cong) auto*
also have *... ↔ g analytic_on ball z (min r r') - {z}*
by *(subst analytic_on_open) auto*
finally show *?thesis*
unfolding *isolated_singularity_at_def*
by *(intro exI[of _ min r r']) (use ⟨r > 0⟩ ⟨r' > 0⟩ in auto)*

qed
from *this[of f g] this[of g f] assms* **show** *?thesis*
by *(auto simp: eq_commute)*

qed

lemma *removable_singularity*:
assumes *f holomorphic_on A - {x} open A*
assumes *f -x→ c*
shows *(λy. if y = x then c else f y) holomorphic_on A (is ?g holomorphic_on _)*
proof –
have *continuous_on A ?g*
unfolding *continuous_on_def*
proof
fix *y* **assume** *y: y ∈ A*

```

show (?g  $\longrightarrow$  ?g y) (at y within A)
proof (cases y = x)
  case False
  have continuous_on (A - {x}) f
    using assms(1) by (meson holomorphic_on_imp_continuous_on)
  hence (f  $\longrightarrow$  ?g y) (at y within A - {x})
    using False y by (auto simp: continuous_on_def)
  also have ?this  $\longleftrightarrow$  (?g  $\longrightarrow$  ?g y) (at y within A - {x})
    by (intro filterlim_cong refl) (auto simp: eventually_at_filter)
  also have at y within A - {x} = at y within A
    using y assms False
    by (intro at_within_nhd[of _ A - {x}]) auto
  finally show ?thesis .
next
  case [simp]: True
  have f -x  $\rightarrow$  c
    by fact
  also have ?this  $\longleftrightarrow$  (?g  $\longrightarrow$  ?g y) (at y)
    by (simp add: LIM_equal)
  finally show ?thesis
    by (meson Lim_at_imp_Lim_at_within)
qed
qed
moreover {
  have ?g holomorphic_on A - {x}
    using assms(1) holomorphic_transform by fastforce
}
ultimately show ?thesis
  using assms by (simp add: no_isolated_singularity)
qed

lemma removable_singularity':
  assumes isolated_singularity_at f z
  assumes f -z  $\rightarrow$  c
  shows ( $\lambda y$ . if y = z then c else f y) analytic_on {z}
proof -
  from assms obtain r where r: r > 0 f analytic_on ball z r - {z}
    by (auto simp: isolated_singularity_at_def)
  have ( $\lambda y$ . if y = z then c else f y) holomorphic_on ball z r
  proof (rule removable_singularity)
    show f holomorphic_on ball z r - {z}
      using r(2) by (subst (asm) analytic_on_open) auto
  qed (use assms in auto)
  thus ?thesis
    using r(1) unfolding analytic_at by (intro exI[of _ ball z r]) auto
qed

named_theorems singularity_intros introduction rules for singularities

```

```

lemma holomorphic_factor_unique:
  fixes f:: complex  $\Rightarrow$  complex and z::complex and r::real and m n::int
  assumes r>0 g z $\neq$ 0 h z $\neq$ 0
    and asm:  $\forall w \in \text{ball } z \ r - \{z\}. f \ w = g \ w * (w-z)^{\text{powi } n} \wedge g \ w \neq 0 \wedge f \ w = h \ w$ 
    *  $(w-z)^{\text{powi } m} \wedge h \ w \neq 0$ 
    and g_holo: g holomorphic_on ball z r and h_holo: h holomorphic_on ball z r
  shows n=m
proof -
  have [simp]: at z within ball z r  $\neq$  bot using <r>0>
    by (auto simp add: at_within_ball_bot_iff)
  have False when n>m
proof -
  have (h  $\longrightarrow$  0) (at z within ball z r)
    proof (rule Lim_transform_within[OF _ <r>0>, where f= $\lambda w. (w-z)^{\text{powi } (n-m)} * g \ w$ ])
    have  $\forall w \in \text{ball } z \ r - \{z\}. h \ w = (w-z)^{\text{powi } (n-m)} * g \ w$ 
      using <n>m> asm <r>0> by (simp add: field_simps power_int_diff) force
    then show  $\llbracket x' \in \text{ball } z \ r; 0 < \text{dist } x' \ z; \text{dist } x' \ z < r \rrbracket$ 
       $\implies (x' - z)^{\text{powi } (n-m)} * g \ x' = h \ x'$  for x' by auto
  next
  define F where F  $\equiv$  at z within ball z r
  define f' where f'  $\equiv \lambda x. (x - z)^{\text{powi } (n-m)}$ 
  have f' z=0 using <n>m> unfolding f'_def by auto
  moreover have continuous F f' unfolding f'_def F_def continuous_def
    using <n>m>
    by (auto simp add: Lim_ident_at intro!: tendsto_powr_complex_0 tendsto_eq_intros)
  ultimately have (f'  $\longrightarrow$  0) F unfolding F_def
    by (simp add: continuous_within)
  moreover have (g  $\longrightarrow$  g z) F
    unfolding F_def
    using <r>0> centre_in_ball continuous_on_def g_holo holomorphic_on_imp_continuous_on
by blast
  ultimately show (( $\lambda w. f' \ w * g \ w$ )  $\longrightarrow$  0) F using tendsto_mult by
fastforce
qed
  moreover have (h  $\longrightarrow$  h z) (at z within ball z r)
    using holomorphic_on_imp_continuous_on[OF h_holo]
    by (auto simp add: continuous_on_def <r>0>)
  ultimately have h z=0 by (auto intro!: tendsto_unique)
  thus False using <h z $\neq$ 0> by auto
qed
  moreover have False when m>n
proof -
  have (g  $\longrightarrow$  0) (at z within ball z r)
    proof (rule Lim_transform_within[OF _ <r>0>, where f= $\lambda w. (w-z)^{\text{powi } (m-n)} * h \ w$ ])
    have  $\forall w \in \text{ball } z \ r - \{z\}. g \ w = (w-z)^{\text{powi } (m-n)} * h \ w$  using <m>n> asm
      by (simp add: field_simps power_int_diff) force

```

```

then show  $\llbracket x' \in \text{ball } z \ r; \ 0 < \text{dist } x' \ z; \text{dist } x' \ z < r \rrbracket$ 
   $\implies (x' - z) \text{ powi } (m - n) * h \ x' = g \ x'$  for  $x'$  by auto
next
define  $F$  where  $F \equiv \text{at } z \ \text{within } \text{ball } z \ r$ 
define  $f'$  where  $f' \equiv \lambda x. (x - z) \text{ powi } (m - n)$ 
have  $f' \ z = 0$  using  $\langle m > n \rangle$  unfolding  $f'_\text{def}$  by auto
moreover have continuous  $F \ f'$  unfolding  $f'_\text{def}$   $F_\text{def}$  continuous_def
  using  $\langle m > n \rangle$ 
by (auto simp: Lim_ident_at intro!: tendsto_powr_complex_0 tendsto_eq_intros)
ultimately have  $(f' \longrightarrow 0) \ F$  unfolding  $F_\text{def}$ 
  by (simp add: continuous_within)
moreover have  $(h \longrightarrow h \ z) \ F$ 
  using holomorphic_on_imp_continuous_on [ $OF \ h\_holo, \text{unfolded } \text{continuous\_on\_on\_def}$ ]  $\langle r > 0 \rangle$ 
  unfolding  $F_\text{def}$  by auto
ultimately show  $((\lambda w. f' \ w * h \ w) \longrightarrow 0) \ F$  using tendsto_mult by
fastforce
qed
moreover have  $(g \longrightarrow g \ z)$  (at  $z$  within  $\text{ball } z \ r$ )
  using holomorphic_on_imp_continuous_on [ $OF \ g\_holo$ ]
  by (auto simp add: continuous_on_def  $\langle r > 0 \rangle$ )
ultimately have  $g \ z = 0$  by (auto intro!: tendsto_unique)
thus  $False$  using  $\langle g \ z \neq 0 \rangle$  by auto
qed
ultimately show  $n = m$  by fastforce
qed

lemma holomorphic_factor_puncture:
assumes  $f\_iso$ : isolated_singularity_at  $f \ z$ 
  and  $not\_essential$   $f \ z$  —  $f$  has either a removable singularity or a pole at  $z$ 
  and  $non\_zero$ :  $\exists F \ w$  in (at  $z$ ).  $f \ w \neq 0$  —  $f$  will not be constantly zero in a
  neighbour of  $z$ 
shows  $\exists ! n :: \text{int}. \exists g \ r. 0 < r \wedge g \ \text{holomorphic\_on } \text{cball } z \ r \wedge g \ z \neq 0$ 
   $\wedge (\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w * (w - z) \text{ powi } n \wedge g \ w \neq 0)$ 
proof —
define  $P$  where  $P = (\lambda f \ n \ g \ r. 0 < r \wedge g \ \text{holomorphic\_on } \text{cball } z \ r \wedge g \ z \neq 0$ 
   $\wedge (\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w * (w - z) \text{ powi } n \wedge g \ w \neq 0))$ 
have  $imp\_unique$ :  $\exists ! n :: \text{int}. \exists g \ r. P \ f \ n \ g \ r$  when  $\exists n \ g \ r. P \ f \ n \ g \ r$ 
proof (rule ex_ex1I [ $OF \ that$ ])
  fix  $n1 \ n2 :: \text{int}$ 
  assume  $g1\_asm$ :  $\exists g1 \ r1. P \ f \ n1 \ g1 \ r1$  and  $g2\_asm$ :  $\exists g2 \ r2. P \ f \ n2 \ g2 \ r2$ 
define  $fac$  where  $fac \equiv \lambda n \ g \ r. \forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w * (w - z) \text{ powi } n$ 
 $\wedge g \ w \neq 0$ 
obtain  $g1 \ r1$  where  $0 < r1$  and  $g1\_holo$ :  $g1 \ \text{holomorphic\_on } \text{cball } z \ r1$  and
 $g1 \ z \neq 0$ 
and  $fac \ n1 \ g1 \ r1$  using  $g1\_asm$  unfolding  $P_\text{def}$   $fac_\text{def}$  by auto
obtain  $g2 \ r2$  where  $0 < r2$  and  $g2\_holo$ :  $g2 \ \text{holomorphic\_on } \text{cball } z \ r2$  and
 $g2 \ z \neq 0$ 
and  $fac \ n2 \ g2 \ r2$  using  $g2\_asm$  unfolding  $P_\text{def}$   $fac_\text{def}$  by auto

```

```

define r where r ≡ min r1 r2
have r>0 using ⟨r1>0⟩ ⟨r2>0⟩ unfolding r_def by auto
moreover have ∀ w∈ball z r-⟨z⟩. f w = g1 w * (w-z) powi n1 ∧ g1 w≠0
  ∧ f w = g2 w * (w-z) powi n2 ∧ g2 w≠0
  using ⟨fac n1 g1 r1⟩ ⟨fac n2 g2 r2⟩ unfolding fac_def r_def
  by fastforce
ultimately show n1=n2
  using g1_holo g2_holo ⟨g1 z≠0⟩ ⟨g2 z≠0⟩
  apply (elim holomorphic_factor_unique)
  by (auto simp add: r_def)
qed

have P_exist: ∃ n g r. P h n g r when
  ∃ z'. (h ⟶ z') (at z) isolated_singularity_at h z ∃_F w in (at z). h w≠0
for h
proof -
  from that(2) obtain r where r>0 and r: h analytic_on ball z r - {z}
  unfolding isolated_singularity_at_def by auto
  obtain z' where (h ⟶ z') (at z) using ⟨∃ z'. (h ⟶ z') (at z)⟩ by auto
  define h' where h'=(λx. if x=z then z' else h x)
  have h' holomorphic_on ball z r
  proof (rule no_isolated_singularity'[of {z}])
    show ∧ w. w ∈ {z} ⟹ (h' ⟶ h' w) (at w within ball z r)
    by (simp add: LIM_cong Lim_at_imp_Lim_at_within ⟨h -z→ z'⟩ h'_def)
    show h' holomorphic_on ball z r - {z}
    using r analytic_imp_holomorphic h'_def holomorphic_transform by
fastforce
  qed auto
  have ?thesis when z'=0
  proof -
    have h' z=0 using that unfolding h'_def by auto
    moreover have ¬ h' constant_on ball z r
    using ⟨∃_F w in (at z). h w≠0⟩ unfolding constant_on_def frequently_def
eventually_at h'_def
    by (metis ⟨0 < r⟩ centre_in_ball dist_commute mem_ball that)
    moreover note ⟨h' holomorphic_on ball z r⟩
    ultimately obtain g r1 n where 0 < n 0 < r1 ball z r1 ⊆ ball z r and
      g: g holomorphic_on ball z r1
      ∧ w. w ∈ ball z r1 ⟹ h' w = (w-z) ^ n * g w
      ∧ w. w ∈ ball z r1 ⟹ g w ≠ 0
    using holomorphic_factor_zero_nonconstant[of _ ball z r z thesis,simplified,
      OF ⟨h' holomorphic_on ball z r⟩ ⟨r>0⟩ ⟨h' z=0⟩ ⟨¬ h' constant_on
ball z r⟩]
    by (auto simp add: dist_commute)
    define rr where rr=r1/2
    have P h' n g rr
    unfolding P_def rr_def
    using ⟨n>0⟩ ⟨r1>0⟩ g by (auto simp add: powr_nat)
    then have P h n g rr

```



```

    unfolding h'_def P_def by auto
  then show ?thesis unfolding P_def by blast
qed
moreover have ?thesis when z'≠0
proof -
  have h' z≠0 using that unfolding h'_def by auto
  obtain r1 where r1>0 cball z r1 ⊆ ball z r ∀ x∈cball z r1. h' x≠0
  proof -
    have isCont h' z h' z≠0
      by (auto simp add: Lim_cong_within ⟨h -z→ z'⟩ ⟨z'≠0⟩ continuous_at
h'_def)
    then obtain r2 where r2>0 ∀ x∈ball z r2. h' x≠0
      using continuous_at_avoid[of z h' 0 ] unfolding ball_def by auto
    define r1 where r1=min r2 r / 2
    have 0 < r1 cball z r1 ⊆ ball z r
      using ⟨r2>0⟩ ⟨r>0⟩ unfolding r1_def by auto
    moreover have ∀ x∈cball z r1. h' x ≠ 0
      using r2 unfolding r1_def by simp
    ultimately show ?thesis using that by auto
  qed
  then have P h' 0 h' r1 using ⟨h' holomorphic_on ball z r⟩ unfolding P_def
by auto
  then have P h 0 h' r1 unfolding P_def h'_def by auto
  then show ?thesis unfolding P_def by blast
qed
ultimately show ?thesis by auto
qed

have ?thesis when ∃ x. (f ⟶ x) (at z)
  apply (rule_tac imp_unique[unfolded P_def])
  using P_exist[OF that(1) f_iso non_zero] unfolding P_def .
moreover have ?thesis when is_pole f z
proof (rule imp_unique[unfolded P_def])
  obtain e where [simp]: e>0 and e_holo: f holomorphic_on ball z e - {z}
and e_nz: ∀ x∈ball z e-{z}. f x≠0
  proof -
    have ∀_F z in at z. f z ≠ 0
      using ⟨is_pole f z⟩ filterlim_at_infinity_imp_eventually_ne unfolding
is_pole_def
      by auto
    then obtain e1 where e1: e1>0 ∀ x∈ball z e1-{z}. f x≠0
      using that eventually_at[of λx. f x≠0 z UNIV,simplified] by (auto simp
add: dist_commute)
    obtain e2 where e2: e2>0 f holomorphic_on ball z e2 - {z}
      using f_iso analytic_imp_holomorphic unfolding isolated_singularity_at_def
by auto
    show ?thesis
      using e1 e2 by (force intro: that[of min e1 e2])
  qed
qed

```

```

define h where  $h \equiv \lambda x. \text{inverse } (f x)$ 
have  $\exists n g r. P h n g r$ 
proof -
  have  $(\lambda x. \text{inverse } (f x)) \text{analytic\_on\_ball } z e - \{z\}$ 
  by (metis e_holo_e_nz_open_ball_analytic_on_open_holomorphic_on_inverse
open_delete)
  moreover have  $h -z \rightarrow 0$ 
  using Lim_transform_within_open assms(2) h_def_is_pole_tendsto that
by fastforce
  moreover have  $\exists_F w \text{ in } (at z). h w \neq 0$ 
  using non_zero by (simp add: h_def)
  ultimately show ?thesis
  using P_exist[of h] <e > 0>
  unfolding isolated_singularity_at_def h_def
  by blast
qed
then obtain n g r
where  $0 < r$  and
  g_holo: g holomorphic_on cball z r and g z ≠ 0 and
  g_fac:  $(\forall w \in \text{cball } z r - \{z\}. h w = g w * (w - z)^{\text{powi } n} \wedge g w \neq 0)$ 
  unfolding P_def by auto
have  $P f (-n) (\text{inverse } o g) r$ 
proof -
  have  $f w = \text{inverse } (g w) * (w - z)^{\text{powi } (-n)}$  when  $w \in \text{cball } z r - \{z\}$  for w
  by (metis g_fac h_def inverse_inverse_eq inverse_mult_distrib power_int_minus
that)
  then show ?thesis
  unfolding P_def comp_def
  using  $\langle r > 0 \rangle$  g_holo g_fac <g z ≠ 0> by (auto intro: holomorphic_intros)
qed
then show  $\exists x g r. 0 < r \wedge g \text{holomorphic\_on } \text{cball } z r \wedge g z \neq 0$ 
   $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w - z)^{\text{powi } x} \wedge g w \neq 0)$ 
  unfolding P_def by blast
qed
ultimately show ?thesis
using  $\langle \text{not\_essential } f z \rangle$  unfolding not_essential_def by presburger
qed

lemma not_essential_transform:
assumes not_essential g z
assumes  $\forall_F w \text{ in } (at z). g w = f w$ 
shows not_essential f z
using assms unfolding not_essential_def
by (simp add: filterlim_cong_is_pole_cong)

lemma isolated_singularity_at_transform:
assumes isolated_singularity_at g z
assumes  $\forall_F w \text{ in } (at z). g w = f w$ 

```

shows *isolated_singularity_at* $f z$
 using *assms isolated_singularity_at_cong* by blast

lemma *not_essential_powr*[*singularity_intros*]:
 assumes *LIM* w (at z). $f w :>$ (at x)
 shows *not_essential* $(\lambda w. (f w) \text{ powi } n) z$
proof –
 define *fp* where $fp = (\lambda w. (f w) \text{ powi } n)$
 have *?thesis* when $n > 0$
proof –
 have $(\lambda w. (f w) \wedge (\text{nat } n)) -z \rightarrow x \wedge \text{nat } n$
 using that *assms unfolding filterlim_at* by (auto intro!: *tendsto_eq_intros*)
 then have $fp -z \rightarrow x \wedge \text{nat } n$ **unfolding** *fp_def*
 by (smt (verit) *LIM_cong power_int_def* that)
 then show *?thesis unfolding not_essential_def fp_def* by auto
qed
 moreover have *?thesis* when $n = 0$
proof –
 have $\forall_F x$ in at z . $fp x = 1$
 using that *filterlim_at_within_not_equal[OF assms]* by (auto simp: *fp_def*)
 then have $fp -z \rightarrow 1$
 by (simp add: *tendsto_eventually*)
 then show *?thesis unfolding fp_def not_essential_def* by auto
qed
 moreover have *?thesis* when $n < 0$
proof (cases $x=0$)
 case True
 have $(\lambda x. f x \wedge \text{nat } (-n)) -z \rightarrow 0$
 using *assms True that unfolding filterlim_at* by (auto intro!: *tendsto_eq_intros*)
 moreover have $\forall_F x$ in at z . $f x \wedge \text{nat } (-n) \neq 0$
 by (smt (verit) *True assms eventually_at_topological filterlim_at_power_eq_0_iff*)
 ultimately have *LIM* w (at z). *inverse* $((f w) \wedge (\text{nat } (-n))) :>$ *at_infinity*
 by (metis *filterlim_atI filterlim_compose filterlim_inverse_at_infinity*)
 then have *LIM* w (at z). $fp w :>$ *at_infinity*
proof (elim *filterlim_mono_eventually*)
 show $\forall_F x$ in at z . *inverse* $(f x \wedge \text{nat } (-n)) = fp x$
 using *filterlim_at_within_not_equal[OF assms, of 0]* **unfolding** *fp_def*
 by (smt (verit) *eventuallyI power_int_def power_inverse* that)
qed auto
 then show *?thesis unfolding fp_def not_essential_def is_pole_def* by auto
next
 case False
 let *?xx* = *inverse* $(x \wedge (\text{nat } (-n)))$
 have $(\lambda w. \text{inverse} ((f w) \wedge (\text{nat } (-n)))) -z \rightarrow ?xx$
 using *assms False unfolding filterlim_at* by (auto intro!: *tendsto_eq_intros*)
 then have $fp -z \rightarrow ?xx$
 by (smt (verit, best) *LIM_cong fp_def power_int_def power_inverse* that)
 then show *?thesis unfolding fp_def not_essential_def* by auto
qed

ultimately show ?thesis by linarith
qed

lemma isolated_singularity_at_powr[singularity_intros]:
 assumes isolated_singularity_at f z $\forall_F w$ in (at z). f w $\neq 0$
 shows isolated_singularity_at $(\lambda w. (f w) \text{ powi } n)$ z
 proof –
 obtain r1 where r1 > 0 f analytic_on ball z r1 – {z}
 using assms(1) unfolding isolated_singularity_at_def by auto
 then have r1: f holomorphic_on ball z r1 – {z}
 using analytic_on_open[of ball z r1 – {z} f] by blast
 obtain r2 where r2 > 0 and r2: $\forall w. w \neq z \wedge \text{dist } w \ z < r2 \longrightarrow f w \neq 0$
 using assms(2) unfolding eventually_at by auto
 define r3 where r3 = min r1 r2
 have $(\lambda w. (f w) \text{ powi } n)$ holomorphic_on ball z r3 – {z}
 by (intro holomorphic_on_power_int) (use r1 r2 in ‹auto simp: dist_commute r3_def›)
 moreover have r3 > 0 unfolding r3_def using ‹0 < r1› ‹0 < r2› by linarith
 ultimately show ?thesis
 by (meson open_ball analytic_on_open isolated_singularity_at_def open_delete)
 qed

lemma non_zero_neighbour:
 assumes f_iso: isolated_singularity_at f z
 and f_ness: not_essential f z
 and f_nconst: $\exists_F w$ in (at z). f w $\neq 0$
 shows $\forall_F w$ in (at z). f w $\neq 0$
 proof –
 obtain fn fp fr
 where [simp]: fp z $\neq 0$ and fr > 0
 and fr: fp holomorphic_on cball z fr
 $\wedge \lambda w. w \in \text{cball } z \text{ fr} - \{z\} \implies f w = fp w * (w - z) \text{ powi } fn \wedge fp w \neq 0$
 using holomorphic_factor_puncture[OF f_iso f_ness f_nconst] by auto
 have f w $\neq 0$ when w $\neq z$ dist w z < fr for w
 proof –
 have f w = fp w * (w - z) powi fn fp w $\neq 0$
 using fr that by (auto simp add: dist_commute)
 moreover have (w - z) powi fn $\neq 0$
 unfolding powr_eq_0_iff using ‹w $\neq z$ › by auto
 ultimately show ?thesis by auto
 qed
 then show ?thesis using ‹fr > 0› unfolding eventually_at by auto
 qed

lemma non_zero_neighbour_pole:
 assumes is_pole f z
 shows $\forall_F w$ in (at z). f w $\neq 0$
 using assms filterlim_at_infinity_imp_eventually_ne[of f at z 0]
 unfolding is_pole_def by auto

```

lemma non_zero_neighbour_alt:
  assumes holo: f holomorphic_on S
    and open S connected S  $z \in S$   $\beta \in S$   $f \beta \neq 0$ 
    shows  $\forall_F w$  in (at z).  $f w \neq 0 \wedge w \in S$ 
proof (cases  $f z = 0$ )
  case True
    from isolated_zeros[OF holo ⟨open S⟩ ⟨connected S⟩ ⟨ $z \in S$ ⟩ True ⟨ $\beta \in S$ ⟩ ⟨ $f \beta \neq 0$ ⟩]
    obtain r where  $0 < r$  ball z r  $\subseteq S$   $\forall w \in$  ball z r  $- \{z\}. f w \neq 0$  by metis
    then show ?thesis
      by (smt (verit) open_ball centre_in_ball eventually_at_topological_insertE
insert_Diff subsetD)
  next
    case False
    obtain r1 where  $r1 > 0$   $\forall y. \text{dist } z \ y < r1 \longrightarrow f y \neq 0$ 
    using continuous_at_avoid[OF f, OF _ False] assms continuous_on_eq_continuous_at
      holo holomorphic_on_imp_continuous_on by blast
    obtain r2 where  $r2 > 0$  ball z r2  $\subseteq S$ 
    using assms openE by blast
    show ?thesis unfolding eventually_at
      by (metis (no_types) dist_commute order.strict_trans linorder_less_linear
mem_ball r1 r2 subsetD)
qed

```

```

lemma not_essential_times[singularity_intros]:
  assumes f_ness: not_essential f z and g_ness: not_essential g z
  assumes f_iso: isolated_singularity_at f z and g_iso: isolated_singularity_at g
z
  shows not_essential ( $\lambda w. f w * g w$ ) z
proof -
  define fg where  $fg = (\lambda w. f w * g w)$ 
  have ?thesis when  $\neg ((\exists_F w$  in (at z).  $f w \neq 0) \wedge (\exists_F w$  in (at z).  $g w \neq 0))$ 
  proof -
    have  $\forall_F w$  in (at z).  $fg w = 0$ 
    using fg_def frequently_elim1 not_eventually_that by fastforce
    from tendsto_cong[OF this] have  $fg -z \rightarrow 0$  by auto
    then show ?thesis unfolding not_essential_def fg_def by auto
  qed
  moreover have ?thesis when f_nconst:  $\exists_F w$  in (at z).  $f w \neq 0$  and g_nconst:
 $\exists_F w$  in (at z).  $g w \neq 0$ 
  proof -
    obtain fn fp fr where [simp]:  $fp \ z \neq 0$  and  $fr > 0$ 
    and fr: fp holomorphic_on cball z fr
       $\forall w \in$  cball z fr  $- \{z\}. f w = fp \ w * (w - z)^{powi \ fn} \wedge fp \ w \neq 0$ 
    using holomorphic_factor_puncture[OF f_iso f_ness f_nconst] by auto
    obtain gn gp gr where [simp]:  $gp \ z \neq 0$  and  $gr > 0$ 
    and gr: gp holomorphic_on cball z gr
       $\forall w \in$  cball z gr  $- \{z\}. g w = gp \ w * (w - z)^{powi \ gn} \wedge gp \ w \neq 0$ 

```

```

using holomorphic_factor_puncture[OF g_iso g_ness g_nconst] by auto

define r1 where r1=(min fr gr)
have r1>0 unfolding r1_def using ‹fr>0› ‹gr>0› by auto
have fg_times: fg w = (fp w * gp w) * (w-z) powi (fn+gn) and fgp_nz: fp
w*gp w≠0
  when w∈ball z r1 - {z} for w
proof -
  have f w = fp w * (w-z) powi fn fp w≠0
    using fr that unfolding r1_def by auto
  moreover have g w = gp w * (w-z) powi gn gp w ≠ 0
    using gr that unfolding r1_def by auto
  ultimately show fg w = (fp w * gp w) * (w-z) powi (fn+gn) fp w*gp w≠0
    using that by (auto simp add: fg_def power_int_add)
qed

obtain [intro]: fp -z→fp z gp -z→gp z
  using fr(1) ‹fr>0› gr(1) ‹gr>0›
  by (metis centre_in_ball continuous_at continuous_on_interior
    holomorphic_on_imp_continuous_on_interior_cball)
have ?thesis when fn+gn>0
proof -
  have (λw. (fp w * gp w) * (w-z) ^ (nat (fn+gn))) -z→0
    using that by (auto intro!:tendsto_eq_intros)
  then have fg -z→ 0
    using Lim_transform_within[OF _ ‹r1>0›]
  by (smt (verit, best) Diff_iff dist_commute fg_times mem_ball power_int_def
    singletonD that zero_less_dist_iff)
  then show ?thesis unfolding not_essential_def fg_def by auto
qed
moreover have ?thesis when fn+gn=0
proof -
  have (λw. fp w * gp w) -z→fp z*gp z
    using that by (auto intro!:tendsto_eq_intros)
  then have fg -z→ fp z*gp z
    apply (elim Lim_transform_within[OF _ ‹r1>0›])
    apply (subst fg_times)
    by (auto simp add: dist_commute that)
  then show ?thesis unfolding not_essential_def fg_def by auto
qed
moreover have ?thesis when fn+gn<0
proof -
  have LIM x at z. (x - z) ^ nat (- (fn + gn)) :=> at 0
    using eventually_at_topological that
    by (force intro!: tendsto_eq_intros filterlim_atI)
  moreover have ∃ c. (λc. fp c * gp c) -z→ c ∧ 0 ≠ c
    using ‹fp -z→ fp z› ‹gp -z→ gp z› tendsto_mult by fastforce
  ultimately have LIM w (at z). fp w * gp w / (w-z) ^ nat (- (fn+gn)) :=>
at_infinity

```

```

    using filterlim_divide_at_infinity by blast
  then have is_pole fg z unfolding is_pole_def
    apply (elim filterlim_transform_within[OF _ _ <r1>0])
    using that
      by (simp_all add: dist_commute fg_times of_int_of_nat divide_simps
power_int_def del: minus_add_distrib)
    then show ?thesis unfolding not_essential_def fg_def by auto
  qed
  ultimately show ?thesis unfolding not_essential_def fg_def by fastforce
  qed
  ultimately show ?thesis by auto
  qed

lemma not_essential_inverse[singularity_intros]:
  assumes f_ness: not_essential f z
  assumes f_iso: isolated_singularity_at f z
  shows not_essential ( $\lambda w. \text{inverse } (f w)$ ) z
  proof -
    define vf where vf = ( $\lambda w. \text{inverse } (f w)$ )
    have ?thesis when  $\neg(\exists_F w \text{ in } (at z). f w \neq 0)$ 
    proof -
      have  $\forall_F w \text{ in } (at z). f w = 0$ 
        using not_eventually_that by fastforce
      then have vf  $-z \rightarrow 0$ 
        unfolding vf_def by (simp add: tendsto_eventually)
      then show ?thesis
        unfolding not_essential_def vf_def by auto
    qed
    moreover have ?thesis when is_pole f z
      by (metis (mono_tags, lifting) filterlim_at filterlim_inverse_at_iff is_pole_def
not_essential_def that)
    moreover have ?thesis when  $\exists x. f -z \rightarrow x$  and f_nconst:  $\exists_F w \text{ in } (at z). f w \neq 0$ 
    proof -
      from that obtain fz where fz:  $f -z \rightarrow fz$  by auto
      have ?thesis when  $fz = 0$ 

      proof -
        have ( $\lambda w. \text{inverse } (vf w)$ )  $-z \rightarrow 0$ 
          using fz that unfolding vf_def by auto
        moreover have  $\forall_F w \text{ in } at z. \text{inverse } (vf w) \neq 0$ 
          using non_zero_neighbour[OF f_iso f_ness f_nconst]
          unfolding vf_def by auto
        ultimately show ?thesis unfolding not_essential_def vf_def
          using filterlim_atI filterlim_inverse_at_iff is_pole_def by blast
      qed
    moreover have ?thesis when  $fz \neq 0$ 
      using fz not_essential_def tendsto_inverse that by blast
    ultimately show ?thesis by auto
  qed
  qed

```

ultimately show *?thesis* using *f_ness unfolding not_essential_def* by auto
qed

lemma *isolated_singularity_at_inverse*[*singularity_intros*]:

assumes *f_iso*: *isolated_singularity_at* *f* *z*

and *f_ness*: *not_essential* *f* *z*

shows *isolated_singularity_at* ($\lambda w. \text{inverse } (f w)$) *z*

proof –

define *vf* where *vf* = ($\lambda w. \text{inverse } (f w)$)

have *?thesis* when $\neg(\exists_{Fw} \text{in } (at z). f w \neq 0)$

proof –

have $\forall_{Fw} \text{in } (at z). f w = 0$

using *that*[*unfolded_frequently_def, simplified*] by (*auto elim: eventually_rev_mp*)

then have $\forall_{Fw} \text{in } (at z). vf w = 0$

unfolding *vf_def* by auto

then obtain *d1* where *d1* > 0 and *d1*: $\forall x. x \neq z \wedge \text{dist } x z < d1 \longrightarrow vf x =$

0

unfolding *eventually_at* by auto

then have *vf* holomorphic_on ball *z* *d1* - {*z*}

using *holomorphic_transform*[of $\lambda_. 0$]

by (*metis Diff_iff dist_commute holomorphic_on_const insert_iff mem_ball*)

then have *vf* analytic_on ball *z* *d1* - {*z*}

by (*simp add: analytic_on_open open_delete*)

then show *?thesis*

using $\langle d1 > 0 \rangle$ unfolding *isolated_singularity_at_def* *vf_def* by auto

qed

moreover have *?thesis* when *f_nconst*: $\exists_{Fw} \text{in } (at z). f w \neq 0$

proof –

have $\forall_{Fw} \text{in } at z. f w \neq 0$ using *non_zero_neighbour*[OF *f_iso f_ness f_nconst*].

then obtain *d1* where *d1*: *d1* > 0 $\forall x. x \neq z \wedge \text{dist } x z < d1 \longrightarrow f x \neq 0$

unfolding *eventually_at* by auto

obtain *d2* where *d2* > 0 and *d2*: *f* analytic_on ball *z* *d2* - {*z*}

using *f_iso* unfolding *isolated_singularity_at_def* by auto

define *d3* where *d3* = min *d1* *d2*

have *d3* > 0 unfolding *d3_def* using $\langle d1 > 0 \rangle \langle d2 > 0 \rangle$ by auto

moreover

have *f* analytic_on ball *z* *d3* - {*z*}

by (*smt* (*verit, best*) *Diff_iff analytic_on_analytic_at* *d2 d3_def mem_ball*)

then have *vf* analytic_on ball *z* *d3* - {*z*}

unfolding *vf_def*

by (*intro analytic_on_inverse; simp add: d1(2) d3_def dist_commute*)

ultimately show *?thesis*

unfolding *isolated_singularity_at_def* *vf_def* by auto

qed

ultimately show *?thesis* by auto

qed

lemma *not_essential_divide*[*singularity_intros*]:


```

assumes f_ness: not_essential f z and g_ness: not_essential g z
assumes f_iso: isolated_singularity_at f z and g_iso: isolated_singularity_at g
z
shows not_essential ( $\lambda w. f w / g w$ ) z
proof –
  have not_essential ( $\lambda w. f w * \text{inverse } (g w)$ ) z
    by (simp add: f_iso f_ness g_iso g_ness isolated_singularity_at_inverse
      not_essential_inverse not_essential_times)
  then show ?thesis by (simp add: field_simps)
qed

```

lemma

```

assumes f_iso: isolated_singularity_at f z
and g_iso: isolated_singularity_at g z
shows isolated_singularity_at_times[singularity_intros]:
  isolated_singularity_at ( $\lambda w. f w * g w$ ) z
and isolated_singularity_at_add[singularity_intros]:
  isolated_singularity_at ( $\lambda w. f w + g w$ ) z
proof –
  obtain d1 d2 where d1>0 d2>0
    and d1: f analytic_on ball z d1 – {z} and d2: g analytic_on ball z d2 – {z}
    using f_iso g_iso unfolding isolated_singularity_at_def by auto
  define d3 where d3=min d1 d2
  have d3>0 unfolding d3_def using  $\langle d1 > 0 \rangle \langle d2 > 0 \rangle$  by auto

  have fan: f analytic_on ball z d3 – {z}
    by (smt (verit, best) Diff_iff analytic_on_analytic_at d1 d3_def mem_ball)
  have gan: g analytic_on ball z d3 – {z}
    by (smt (verit, best) Diff_iff analytic_on_analytic_at d2 d3_def mem_ball)
  have ( $\lambda w. f w * g w$ ) analytic_on ball z d3 – {z}
    using analytic_on_mult fan gan by blast
  then show isolated_singularity_at ( $\lambda w. f w * g w$ ) z
    using  $\langle d3 > 0 \rangle$  unfolding isolated_singularity_at_def by auto
  have ( $\lambda w. f w + g w$ ) analytic_on ball z d3 – {z}
    using analytic_on_add fan gan by blast
  then show isolated_singularity_at ( $\lambda w. f w + g w$ ) z
    using  $\langle d3 > 0 \rangle$  unfolding isolated_singularity_at_def by auto
qed

```

lemma *isolated_singularity_at_uminus*[*singularity_intros*]:

```

assumes f_iso: isolated_singularity_at f z
shows isolated_singularity_at ( $\lambda w. - f w$ ) z
using assms unfolding isolated_singularity_at_def using analytic_on_neg by
blast

```

lemma *isolated_singularity_at_id*[*singularity_intros*]:

```

isolated_singularity_at ( $\lambda w. w$ ) z
unfolding isolated_singularity_at_def by (simp add: gt_ex)

```

lemma *isolated_singularity_at_minus*[singularity_intros]:
assumes *isolated_singularity_at* f z **and** *isolated_singularity_at* g z
shows *isolated_singularity_at* $(\lambda w. f\ w - g\ w)$ z
unfolding *diff_conv_add_uminus*
using *assms isolated_singularity_at_add isolated_singularity_at_uminus* **by**
blast

lemma *isolated_singularity_at_divide*[singularity_intros]:
assumes *isolated_singularity_at* f z
and *isolated_singularity_at* g z
and *not_essential* g z
shows *isolated_singularity_at* $(\lambda w. f\ w / g\ w)$ z
unfolding *divide_inverse*
by (*simp add: assms isolated_singularity_at_inverse isolated_singularity_at_times*)

lemma *isolated_singularity_at_const*[singularity_intros]:
isolated_singularity_at $(\lambda w. c)$ z
unfolding *isolated_singularity_at_def* **by** (*simp add: gt_ex*)

lemma *isolated_singularity_at_holomorphic*:
assumes f *holomorphic_on* $s - \{z\}$ *open* $s\ z \in s$
shows *isolated_singularity_at* f z
using *assms* **unfolding** *isolated_singularity_at_def*
by (*metis analytic_on_holomorphic centre_in_ball insert_Diff openE open_delete subset_insert_iff*)

lemma *isolated_singularity_at_altdef*:
isolated_singularity_at f $z \longleftrightarrow$ *eventually* $(\lambda z. f$ *analytic_on* $\{z\})$ $(at\ z)$
proof
assume *isolated_singularity_at* f z
then obtain r **where** $r: r > 0$ f *analytic_on* $ball\ z\ r - \{z\}$
unfolding *isolated_singularity_at_def* **by** *blast*
have *eventually* $(\lambda w. w \in ball\ z\ r - \{z\})$ $(at\ z)$
using $r(1)$ **by** (*intro eventually_at_in_open*) *auto*
thus *eventually* $(\lambda z. f$ *analytic_on* $\{z\})$ $(at\ z)$
by *eventually_elim* (*use r analytic_on_subset in auto*)
next
assume *eventually* $(\lambda z. f$ *analytic_on* $\{z\})$ $(at\ z)$
then obtain A **where** $A: open\ A\ z \in A \wedge w. w \in A - \{z\} \implies f$ *analytic_on*
 $\{w\}$
unfolding *eventually_at_topological* **by** *blast*
then show *isolated_singularity_at* f z
by (*meson analytic_imp_holomorphic analytic_on_analytic_at isolated_singularity_at_holomorphic*)
qed

lemma *isolated_singularity_at_shift*:
assumes *isolated_singularity_at* $(\lambda x. f\ (x + w))$ z
shows *isolated_singularity_at* f $(z + w)$
proof –

```

from assms obtain r where r:  $r > 0$  and ana:  $(\lambda x. f (x + w))$  analytic_on
ball z r - {z}
  unfolding isolated_singularity_at_def by blast
  have  $((\lambda x. f (x + w)) \circ (\lambda x. x - w))$  analytic_on (ball (z + w) r - {z + w})
  by (rule analytic_on_compose_gen[OF _ ana])
    (auto simp: dist_norm algebra_simps intro!: analytic_intros)
  hence f analytic_on (ball (z + w) r - {z + w})
  by (simp add: o_def)
  thus ?thesis using r
  unfolding isolated_singularity_at_def by blast
qed

```

```

lemma isolated_singularity_at_shift_iff:
isolated_singularity_at f (z + w)  $\longleftrightarrow$  isolated_singularity_at  $(\lambda x. f (x + w))$  z
using isolated_singularity_at_shift[of f w z]
  isolated_singularity_at_shift[of  $\lambda x. f (x + w) - w$  w + z]
by (auto simp: algebra_simps)

```

```

lemma isolated_singularity_at_shift_0:
  NO_MATCH 0 z  $\implies$  isolated_singularity_at f z  $\longleftrightarrow$  isolated_singularity_at
 $(\lambda x. f (z + x))$  0
using isolated_singularity_at_shift_iff[of f 0 z] by (simp add: add_ac)

```

```

lemma not_essential_shift:
  assumes not_essential  $(\lambda x. f (x + w))$  z
  shows not_essential f (z + w)
proof -
  from assms consider c where  $(\lambda x. f (x + w)) -z \rightarrow c$  | is_pole  $(\lambda x. f (x + w))$  z
  unfolding not_essential_def by blast
  thus ?thesis
proof cases
  case (1 c)
  hence  $f -z + w \rightarrow c$ 
  by (smt (verit, ccfv_SIG) LIM_cong add.assoc filterlim_at_to_0)
  thus ?thesis
  by (auto simp: not_essential_def)
next
  case 2
  hence is_pole f (z + w)
  by (subst is_pole_shift_iff [symmetric]) (auto simp: o_def add_ac)
  thus ?thesis
  by (auto simp: not_essential_def)
qed
qed

```

```

lemma not_essential_shift_iff: not_essential f (z + w)  $\longleftrightarrow$  not_essential  $(\lambda x. f (x + w))$  z
using not_essential_shift[of f w z]

```

$\text{not_essential_shift}$ [of $\lambda x. f(x + w) - w w + z$]
by (auto simp: algebra_simps)

lemma *not_essential_shift_0*:

NO_MATCH $0 z \implies \text{not_essential } f z \longleftrightarrow \text{not_essential } (\lambda x. f(z + x)) 0$
using *not_essential_shift_iff*[of $f 0 z$] **by** (simp add: add_ac)

lemma *not_essential_holomorphic*:

assumes f holomorphic_on A $x \in A$ open A
shows $\text{not_essential } f x$
by (metis assms at_within_open continuous_on holomorphic_on_imp_continuous_on
not_essential_def)

lemma *not_essential_analytic*:

assumes f analytic_on $\{z\}$
shows $\text{not_essential } f z$
using *analytic_at_assms_not_essential_holomorphic* **by** blast

lemma *not_essential_id* [*singularity_intros*]: $\text{not_essential } (\lambda w. w) z$
by (simp add: not_essential_analytic)

lemma *is_pole_imp_not_essential* [*intro*]: $\text{is_pole } f z \implies \text{not_essential } f z$
by (auto simp: not_essential_def)

lemma *tendsto_imp_not_essential* [*intro*]: $f -z \rightarrow c \implies \text{not_essential } f z$
by (auto simp: not_essential_def)

lemma *eventually_not_pole*:

assumes *isolated_singularity_at* $f z$
shows *eventually* $(\lambda w. \neg \text{is_pole } f w)$ (at z)
proof –
from assms **obtain** r **where** $r > 0$ **and** r : f analytic_on ball $z r - \{z\}$
by (auto simp: isolated_singularity_at_def)
then have *eventually* $(\lambda w. w \in \text{ball } z r - \{z\})$ (at z)
by (intro eventually_at_in_open) auto
thus *eventually* $(\lambda w. \neg \text{is_pole } f w)$ (at z)
by (metis (no_types, lifting) analytic_at analytic_on_analytic_at eventually_mono not_is_pole_holomorphic r)
qed

lemma *not_islimpt_poles*:

assumes *isolated_singularity_at* $f z$
shows $\neg z$ islimpt $\{w. \text{is_pole } f w\}$
using *eventually_not_pole* [*OF* assms]
by (auto simp: islimpt_conv_frequently_at_frequently_def)

lemma *analytic_at_imp_no_pole*: f analytic_on $\{z\} \implies \neg \text{is_pole } f z$
using *analytic_at_not_is_pole_holomorphic* **by** blast

lemma *not_essential_const* [*singularity_intros*]: *not_essential* ($\lambda_. c$) *z*
 by *blast*

lemma *not_essential_uminus* [*singularity_intros*]:

assumes *f_ness*: *not_essential* *f z*

assumes *f_iso*: *isolated_singularity_at* *f z*

shows *not_essential* ($\lambda w. -f w$) *z*

proof –

have *not_essential* ($\lambda w. -1 * f w$) *z*

by (*intro* *assms* *singularity_intros*)

thus ?*thesis* by *simp*

qed

lemma *isolated_singularity_at_analytic*:

assumes *f* *analytic_on* {*z*}

shows *isolated_singularity_at* *f z*

by (*meson* *Diff_subset_analytic_at* *assms* *holomorphic_on_subset* *isolated_singularity_at_holomorphic*)

lemma *isolated_singularity_sum* [*singularity_intros*]:

assumes $\bigwedge x. x \in A \implies \text{isolated_singularity_at } (f x) z$

shows *isolated_singularity_at* ($\lambda w. \sum_{x \in A}. f x w$) *z*

using *assms* by (*induction* *A* *rule*: *infinite_finite_induct*) (*auto* *intro!*: *singularity_intros*)

lemma *isolated_singularity_prod* [*singularity_intros*]:

assumes $\bigwedge x. x \in A \implies \text{isolated_singularity_at } (f x) z$

shows *isolated_singularity_at* ($\lambda w. \prod_{x \in A}. f x w$) *z*

using *assms* by (*induction* *A* *rule*: *infinite_finite_induct*) (*auto* *intro!*: *singularity_intros*)

lemma *isolated_singularity_sum_list* [*singularity_intros*]:

assumes $\bigwedge f. f \in \text{set } fs \implies \text{isolated_singularity_at } f z$

shows *isolated_singularity_at* ($\lambda w. \sum_{f \leftarrow fs}. f w$) *z*

using *assms* by (*induction* *fs*) (*auto* *intro!*: *singularity_intros*)

lemma *isolated_singularity_prod_list* [*singularity_intros*]:

assumes $\bigwedge f. f \in \text{set } fs \implies \text{isolated_singularity_at } f z$

shows *isolated_singularity_at* ($\lambda w. \prod_{f \leftarrow fs}. f w$) *z*

using *assms* by (*induction* *fs*) (*auto* *intro!*: *singularity_intros*)

lemma *isolated_singularity_sum_mset* [*singularity_intros*]:

assumes $\bigwedge f. f \in \# F \implies \text{isolated_singularity_at } f z$

shows *isolated_singularity_at* ($\lambda w. \sum_{f \in \# F}. f w$) *z*

using *assms* by (*induction* *F*) (*auto* *intro!*: *singularity_intros*)

lemma *isolated_singularity_prod_mset* [*singularity_intros*]:

assumes $\bigwedge f. f \in \# F \implies \text{isolated_singularity_at } f z$

shows *isolated_singularity_at* ($\lambda w. \prod_{f \in \# F}. f w$) *z*

using *assms* by (*induction* *F*) (*auto* *intro!*: *singularity_intros*)

lemma *analytic_nhd_imp_isolated_singularity*:
assumes f *analytic_on* $A - \{x\}$ $x \in A$ *open* A
shows *isolated_singularity_at* f x
unfolding *isolated_singularity_at_def* **using** *assms*
using *analytic_imp_holomorphic* *isolated_singularity_at_def* *isolated_singularity_at_holomorphic*
by *blast*

lemma *isolated_singularity_at_iff_analytic_nhd*:
isolated_singularity_at f $x \longleftrightarrow (\exists A. x \in A \wedge \text{open } A \wedge f \text{ analytic_on } A - \{x\})$
by (*meson open_ball analytic_nhd_imp_isolated_singularity*
centre_in_ball isolated_singularity_at_def)

7.12 The order of non-essential singularities (i.e. removable singularities or poles)

definition *zorder* :: $(\text{complex} \Rightarrow \text{complex}) \Rightarrow \text{complex} \Rightarrow \text{int}$ **where**
 $\text{zorder } f z = (\text{THE } n. (\exists h r. r > 0 \wedge h \text{ holomorphic_on } \text{cball } z r \wedge h z \neq 0$
 $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = h w * (w - z)^{\text{powi } n}$
 $\wedge h w \neq 0)))$

definition *zor_poly*
:: $[\text{complex} \Rightarrow \text{complex}, \text{complex}] \Rightarrow \text{complex} \Rightarrow \text{complex}$ **where**
 $\text{zor_poly } f z = (\text{SOME } h. \exists r. r > 0 \wedge h \text{ holomorphic_on } \text{cball } z r \wedge h z \neq 0$
 $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = h w * (w - z)^{\text{powi } (\text{zorder } f z)}$
 $\wedge h w \neq 0))$

lemma *zorder_exist*:
fixes $f :: \text{complex} \Rightarrow \text{complex}$ **and** $z :: \text{complex}$
defines $n \equiv \text{zorder } f z$ **and** $g \equiv \text{zor_poly } f z$
assumes f *iso*: *isolated_singularity_at* $f z$
and f *ness*: *not_essential* $f z$
and f *nconst*: $\exists_F w$ *in* (*at* z). $f w \neq 0$
shows $g z \neq 0 \wedge (\exists r. r > 0 \wedge g \text{ holomorphic_on } \text{cball } z r$
 $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w - z)^{\text{powi } n} \wedge g w \neq 0))$
proof –
define P **where** $P = (\lambda n g r. 0 < r \wedge g \text{ holomorphic_on } \text{cball } z r \wedge g z \neq 0$
 $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w - z)^{\text{powi } n} \wedge g w \neq 0))$
have $\exists! k. \exists g r. P k g r$
using *holomorphic_factor_puncture* [*OF assms*(\exists -)] **unfolding** P *def* **by** *auto*
then **have** $\exists g r. P n g r$
unfolding n *def* P *def* *zorder_def* **by** (*rule theI'*)
then **have** $\exists r. P n g r$
unfolding P *def* *zor_poly_def* g *def* n *def* **by** (*rule someI_ex*)
then **obtain** $r1$ **where** $P n g r1$
by *auto*
then **show** *?thesis*
unfolding P *def* **by** *auto*
qed

```

lemma zorder_shift:
  shows zorder f z = zorder ( $\lambda u. f (u + z)$ ) 0
  unfolding zorder_def
  apply (rule arg_cong [of concl: The])
  apply (auto simp: fun_eq_iff Ball_def dist_norm)
  subgoal for x h r
    apply (rule_tac x=h o (+)z in exI)
    apply (rule_tac x=r in exI)
    apply (intro conjI holomorphic_on_compose holomorphic_intros)
      apply (simp_all flip: cball_translation)
    apply (simp add: add commute)
  done
  subgoal for x h r
    apply (rule_tac x=h o ( $\lambda u. u - z$ ) in exI)
    apply (rule_tac x=r in exI)
    apply (intro conjI holomorphic_on_compose holomorphic_intros)
      apply (simp_all flip: cball_translation_subtract)
    by (metis diff_add_cancel eq_iff_diff_eq_0 norm_minus_commute)
  done

lemma zorder_shift': NO_MATCH 0 z  $\implies$  zorder f z = zorder ( $\lambda u. f (u + z)$ ) 0
  by (rule zorder_shift)

lemma
  fixes f:: complex  $\Rightarrow$  complex and z::complex
  assumes f_iso: isolated_singularity_at f z
    and f_ness: not_essential f z
    and f_nconst:  $\exists_F w$  in (at z). f w  $\neq$  0
  shows zorder_inverse: zorder ( $\lambda w. inverse (f w)$ ) z = - zorder f z
    and zor_poly_inverse:  $\forall_F w$  in (at z). zor_poly ( $\lambda w. inverse (f w)$ ) z w
      = inverse (zor_poly f z w)

proof -
  define vf where vf = ( $\lambda w. inverse (f w)$ )
  define fn vfn where
    fn = zorder f z and vfn = zorder vf z
  define fp vfp where
    fp = zor_poly f z and vfp = zor_poly vf z

  obtain fr where [simp]: fp z  $\neq$  0 and fr > 0
    and fr: fp holomorphic_on cball z fr
       $\forall w \in cball z fr - \{z\}. f w = fp w * (w - z)^{powi fn} \wedge fp w \neq 0$ 
  using zorder_exist[OF f_iso f_ness f_nconst, folded fn_def fp_def]
  by auto
  have fr_inverse: vf w = (inverse (fp w)) * (w - z)^{powi (-fn)}
    and fr_nz: inverse (fp w)  $\neq$  0
    when w  $\in$  ball z fr - {z} for w
  proof -
    have f w = fp w * (w - z)^{powi fn} fp w  $\neq$  0

```

```

    using fr(2) that by auto
  then show  $vf\ w = (\text{inverse } (fp\ w)) * (w-z)\ \text{powi } (-fn)\ \text{inverse } (fp\ w) \neq 0$ 
    by (simp_all add: power_int_minus vf_def)
qed
obtain vfr where [simp]:  $vfp\ z \neq 0$  and  $vfr > 0$  and vfr:  $vfp\ \text{holomorphic\_on } cball\ z\ vfr$ 
  ( $\forall w \in cball\ z\ vfr - \{z\}. vfw = vfp\ w * (w-z)\ \text{powi } vfn \wedge vfp\ w \neq 0$ )
proof -
  have isolated_singularity_at  $vf\ z$ 
  using isolated_singularity_at_inverse[OF f_iso f_ness] unfolding vf_def .
  moreover have not_essential  $vf\ z$ 
  using not_essential_inverse[OF f_ness f_iso] unfolding vf_def .
  moreover have  $\exists_F\ w\ \text{in } at\ z. vfw \neq 0$ 
  using f_nconst unfolding vf_def by (auto elim: frequently_elim1)
  ultimately show ?thesis using zorder_exist[of  $vf\ z$ , folded vfn_def vfp_def]
that by auto
qed

define r1 where  $r1 = \min\ fr\ vfr$ 
have  $r1 > 0$  using  $\langle fr > 0 \rangle\ \langle vfr > 0 \rangle$  unfolding r1_def by simp
show  $vfn = -fn$ 
proof (rule holomorphic_factor_unique)
  have  $\S: \bigwedge w. \llbracket fp\ w = 0; dist\ z\ w < fr \rrbracket \implies False$ 
  using fr_nz by force
  then show  $\forall w \in ball\ z\ r1 - \{z\}. vfw = vfp\ w * (w-z)\ \text{powi } vfn \wedge vfw \neq 0 \wedge vfw = \text{inverse } (fp\ w) * (w-z)\ \text{powi } (-fn) \wedge \text{inverse } (fp\ w) \neq 0$ 
  using fr_inverse r1_def vfr(2)
  by (smt (verit) Diff_iff inverse_nonzero_iff_nonzero mem_ball mem_cball)
  show  $vfp\ \text{holomorphic\_on } ball\ z\ r1$ 
  using r1_def vfr(1) by auto
  show  $(\lambda w. \text{inverse } (fp\ w))\ \text{holomorphic\_on } ball\ z\ r1$ 
  by (metis  $\S\ ball\_subset\_cball\ fr(1)\ \text{holomorphic\_on\_inverse}\ \text{holomorphic\_on\_subset}\ mem\_ball\ min.cobounded2\ min.commute\ r1\_def\ subset\_ball$ )
qed (use  $\langle r1 > 0 \rangle$  in auto)
have  $vfp\ w = \text{inverse } (fp\ w)$  when  $w \in ball\ z\ r1 - \{z\}$  for  $w$ 
proof -
  have  $w \in ball\ z\ fr - \{z\} \wedge w \in cball\ z\ vfr - \{z\} \wedge w \neq z$ 
  using that unfolding r1_def by auto
  then show ?thesis
  by (metis  $\langle vfn = -fn \rangle\ \text{power\_int\_not\_zero}\ \text{right\_minus\_eq}\ fr\_inverse\ vfr(2)\ \text{vector\_space\_over\_itself.scale\_right\_imp\_eq}$ )
qed
then show  $\forall_F\ w\ \text{in } (at\ z). vfw = \text{inverse } (fp\ w)$ 
  unfolding eventually_at by (metis DiffI dist_commute mem_ball singletonD  $\langle r1 > 0 \rangle$ )
qed

```



```

lemma zor_poly_shift:
  assumes iso1: isolated_singularity_at f z
    and ness1: not_essential f z
    and nzero1:  $\exists_F w$  in at z. f w  $\neq 0$ 
  shows  $\forall_F w$  in nhds z. zor_poly f z w = zor_poly ( $\lambda u. f (z + u)$ ) 0 (w-z)
proof -
  obtain r1 where r1>0 zor_poly f z z  $\neq 0$  and
    holo1: zor_poly f z holomorphic_on cball z r1 and
    rball1:  $\forall w \in$  cball z r1 - {z}.
      f w = zor_poly f z w * (w-z) powi (zorder f z)  $\wedge$ 
      zor_poly f z w  $\neq 0$ 
    using zorder_exist[OF iso1 ness1 nzero1] by blast

  define ff where ff=( $\lambda u. f (z + u)$ )
  have isolated_singularity_at ff 0
    unfolding ff_def
    using iso1 isolated_singularity_at_shift_iff[of f 0 z]
    by (simp add: algebra_simps)
  moreover have not_essential ff 0
    unfolding ff_def
    using ness1 not_essential_shift_iff[of f 0 z]
    by (simp add: algebra_simps)
  moreover have  $\exists_F w$  in at 0. ff w  $\neq 0$ 
    unfolding ff_def using nzero1
    by (smt (verit, ccfv_SIG) add.commute eventually_at_to_0
      eventually_mono not_frequently)
  ultimately
  obtain r2 where r2>0 zor_poly ff 0 0  $\neq 0$ 
    and holo2: zor_poly ff 0 holomorphic_on cball 0 r2
    and rball2:  $\forall w \in$  cball 0 r2 - {0}.
      ff w = zor_poly ff 0 w * w powi (zorder ff 0)  $\wedge$  zor_poly ff 0 w  $\neq 0$ 
    using zorder_exist[of ff 0] by auto

  define r where r=min r1 r2
  have r>0 using <r1>0> <r2>0> unfolding r_def by auto

  have zor_poly f z w = zor_poly ff 0 (w-z)
    if w $\in$ ball z r - {z} for w
  proof -
    define n where n  $\equiv$  zorder f z

    have f w = zor_poly f z w * (w-z) powi n
      using n_def r_def rball1 that by auto
    moreover have f w = zor_poly ff 0 (w-z) * (w-z) powi n
  proof -
    have w-z $\in$ cball 0 r2 - {0}
      using r_def that by (auto simp: dist_complex_def)
    then have ff (w-z) = zor_poly ff 0 (w-z) * (w-z) powi (zorder ff 0)

```

```

    using rball2 by blast
  moreover have of_int (zorder ff 0) = n
    unfolding n_def ff_def by (simp add:zorder_shift' add.commute)
  ultimately show ?thesis unfolding ff_def by auto
qed
  ultimately have zor_poly f z w * (w-z) powi n = zor_poly ff 0 (w-z) *
(w-z) powi n
  by auto
  moreover have (w-z) powi n ≠ 0
    using that by auto
  ultimately show ?thesis
    using mult_cancel_right by blast
qed
then have  $\forall_F w \text{ in } \text{at } z. \text{zor\_poly } f z w = \text{zor\_poly } ff 0 (w-z)$ 
  unfolding eventually_at
  by (metis DiffI ⟨0 < r⟩ dist_commute mem_ball singletonD)
moreover have isCont (zor_poly f z) z
  using holo1[THEN holomorphic_on_imp_continuous_on]
  by (simp add: ⟨0 < r1⟩ continuous_on_interior)
moreover
have isCont (zor_poly ff 0) 0
  using ⟨0 < r2⟩ continuous_on_interior holo2 holomorphic_on_imp_continuous_on
  by fastforce
then have isCont ( $\lambda w. \text{zor\_poly } ff 0 (w-z)$ ) z
  unfolding isCont_iff by simp
ultimately show  $\forall_F w \text{ in } \text{nhds } z. \text{zor\_poly } f z w = \text{zor\_poly } ff 0 (w-z)$ 
  by (elim at_within_isCont_imp_nhds; auto)
qed

lemma
  fixes f g:: complex  $\Rightarrow$  complex and z::complex
  assumes f_iso: isolated_singularity_at f z and g_iso: isolated_singularity_at g
z
  and f_ess: not_essential f z and g_ess: not_essential g z
  and fg_nconst:  $\exists_F w \text{ in } (\text{at } z). f w * g w \neq 0$ 
  shows zorder_times:  $\text{zorder } (\lambda w. f w * g w) z = \text{zorder } f z + \text{zorder } g z$  and
  zor_poly_times:  $\forall_F w \text{ in } (\text{at } z). \text{zor\_poly } (\lambda w. f w * g w) z w$ 
= zor_poly f z w * zor_poly g z w

proof -
  define fg where fg = ( $\lambda w. f w * g w$ )
  define fn gn fgn where
    fn = zorder f z and gn = zorder g z and fgn = zorder fg z
  define fp gp fgp where
    fp = zor_poly f z and gp = zor_poly g z and fgp = zor_poly fg z
  have f_nconst:  $\exists_F w \text{ in } (\text{at } z). f w \neq 0$  and g_nconst:  $\exists_F w \text{ in } (\text{at } z). g w \neq 0$ 
    using fg_nconst by (auto elim!:frequently_elim1)
  obtain fr where [simp]: fp z ≠ 0 and fr > 0
    and fr: fp holomorphic_on cball z fr
     $\forall w \in \text{cball } z \text{ fr} - \{z\}. f w = fp w * (w-z) \text{ powi } fn \wedge fp w \neq 0$ 

```

```

    using zorder_exist[OF f_iso f_ness f_nconst, folded fp_def fn_def] by auto
  obtain gr where [simp]: gp z ≠ 0 and gr > 0
    and gr: gp holomorphic_on cball z gr
      ∀ w ∈ cball z gr - {z}. g w = gp w * (w - z) powi gn ∧ gp w ≠ 0
    using zorder_exist[OF g_iso g_ness g_nconst, folded gn_def gp_def] by auto
  define r1 where r1 = min fr gr
  have r1 > 0 unfolding r1_def using ⟨fr > 0⟩ ⟨gr > 0⟩ by auto
  have fg_times: fg w = (fp w * gp w) * (w - z) powi (fn + gn) and fgp_nz: fp w * gp
w ≠ 0
    when w ∈ ball z r1 - {z} for w
  proof -
    have f w = fp w * (w - z) powi fn fp w ≠ 0
      using fr(2) r1_def that by auto
    moreover have g w = gp w * (w - z) powi gn gp w ≠ 0
      using gr(2) that unfolding r1_def by auto
    ultimately show fg w = (fp w * gp w) * (w - z) powi (fn + gn) fp w * gp w ≠ 0
      using that unfolding fg_def by (auto simp add: power_int_add)
  qed

  obtain fgr where [simp]: fgp z ≠ 0 and fgr > 0
    and fgr: fgp holomorphic_on cball z fgr
      ∀ w ∈ cball z fgr - {z}. fg w = fgp w * (w - z) powi fgn ∧ fgp w ≠ 0
  proof -
    have isolated_singularity_at fg z
      unfolding fg_def using isolated_singularity_at_times[OF f_iso g_iso] .
    moreover have not_essential fg z
      by (simp add: f_iso f_ness fg_def g_iso g_ness not_essential_times)
    moreover have ∃_F w in at z. fg w ≠ 0
      using fg_def fg_nconst by blast
    ultimately show ?thesis
      using that zorder_exist[of fg z] fgn_def fgp_def by fastforce
  qed
  define r2 where r2 = min fgr r1
  have r2 > 0 using ⟨r1 > 0⟩ ⟨fgr > 0⟩ unfolding r2_def by simp
  show fgn = fn + gn
  proof (rule holomorphic_factor_unique)
    show ∀ w ∈ ball z r2 - {z}. fg w = fgp w * (w - z) powi fgn ∧ fgp w ≠ 0 ∧
fg w = fp w * gp w * (w - z) powi (fn + gn) ∧ fp w * gp w ≠ 0
      using fg_times fgp_nz fgr(2) r2_def by fastforce
    next
      show fgp holomorphic_on ball z r2
        using fgr(1) r2_def by auto
    next
      show (λw. fp w * gp w) holomorphic_on ball z r2
        by (metis ball_subset_cball fr(1) gr(1) holomorphic_on_mult holomor-
phic_on_subset
          min.cobounded1 min.cobounded2 r1_def r2_def subset_ball)
  qed (auto simp add: ⟨0 < r2⟩)

```

```

have fgp w = fp w * gp w when w: w ∈ ball z r2 - {z} for w
proof -
  have w ∈ ball z r1 - {z} w ∈ cball z fgr - {z} w ≠ z
  using w unfolding r2_def by auto
  then show ?thesis
  by (metis ⟨fgn = fn + gn⟩ eq_iff_diff_eq_0 fg_times fgr(2) power_int_eq_0_iff
    mult_right_cancel)
qed
then show ∀_F w in (at z). fgp w = fp w * gp w
  using ⟨r2>0 unfolding eventually_at by (auto simp add: dist_commute)
qed

```

lemma

```

fixes f g:: complex ⇒ complex and z::complex
assumes f_iso: isolated_singularity_at f z and g_iso: isolated_singularity_at g
z
  and f_ness: not_essential f z and g_ness: not_essential g z
  and fg_nconst: ∃_F w in (at z). f w * g w ≠ 0
shows zorder_divide: zorder (λw. f w / g w) z = zorder f z - zorder g z and
zorder_poly_divide: ∀_F w in (at z). zorder_poly (λw. f w / g w) z w
= zorder_poly f z w / zorder_poly g z w

```

proof -

```

have f_nconst: ∃_F w in (at z). f w ≠ 0 and g_nconst: ∃_F w in (at z). g w ≠ 0
  using fg_nconst by (auto elim!: frequently_elim1)
define vg where vg = (λw. inverse (g w))
have 1: isolated_singularity_at vg z
  by (simp add: g_iso g_ness isolated_singularity_at_inverse vg_def)
moreover have 2: not_essential vg z
  by (simp add: g_iso g_ness not_essential_inverse vg_def)
moreover have 3: ∃_F w in at z. f w * vg w ≠ 0
  using fg_nconst vg_def by auto
ultimately have zorder (λw. f w * vg w) z = zorder f z + zorder vg z
  using zorder_times[OF f_iso _ f_ness] by blast
then show zorder (λw. f w / g w) z = zorder f z - zorder g z
  using zorder_inverse[OF g_iso g_ness g_nconst, folded vg_def] unfolding
vg_def
  by (auto simp add: field_simps)
have ∀_F w in at z. zorder_poly (λw. f w * vg w) z w = zorder_poly f z w * zorder_poly
vg z w
  using zorder_poly_times[OF f_iso _ f_ness, of vg] 1 2 3 by blast
then show ∀_F w in (at z). zorder_poly (λw. f w / g w) z w = zorder_poly f z w /
zorder_poly g z w
  using zorder_poly_inverse[OF g_iso g_ness g_nconst, folded vg_def] unfolding
vg_def
  by eventually_elim (auto simp add: field_simps)
qed

```

lemma zorder_exist_zero:

```

fixes f:: complex ⇒ complex and z::complex

```

```

defines  $n \equiv zorder\ f\ z$  and  $g \equiv zor\_poly\ f\ z$ 
assumes holo:  $f$  holomorphic_on  $S$  and open  $S$  connected  $S\ z \in S$ 
and non_const:  $\exists w \in S. f\ w \neq 0$ 
shows (if  $f\ z = 0$  then  $n > 0$  else  $n = 0$ )  $\wedge$  ( $\exists r. r > 0 \wedge cball\ z\ r \subseteq S \wedge g$  holomorphic_on  $cball\ z\ r$ 
 $\wedge$  ( $\forall w \in cball\ z\ r. f\ w = g\ w * (w - z) ^ nat\ n \wedge g\ w \neq 0$ ))
proof -
obtain  $r$  where  $g\ z \neq 0$  and  $r: r > 0\ cball\ z\ r \subseteq S\ g$  holomorphic_on  $cball\ z\ r$ 
( $\forall w \in cball\ z\ r - \{z\}. f\ w = g\ w * (w - z) ^ powi\ n \wedge g\ w \neq 0$ )
proof -
have  $g\ z \neq 0 \wedge (\exists r > 0. g$  holomorphic_on  $cball\ z\ r$ 
 $\wedge$  ( $\forall w \in cball\ z\ r - \{z\}. f\ w = g\ w * (w - z) ^ powi\ n \wedge g\ w \neq 0$ ))
proof (rule zorder_exist[of  $f\ z$ , folded g_def n_def])
show isolated_singularity_at  $f\ z$ 
using  $\langle open\ S \rangle\ \langle z \in S \rangle$  holo holomorphic_on_imp_analytic_at isolated_singularity_at_analytic

by force
show not_essential  $f\ z$  unfolding not_essential_def
using  $\langle open\ S \rangle\ \langle z \in S \rangle$  at_within_open continuous_on_holo holomorphic_on_imp_continuous_on
by fastforce
have  $\forall_F\ w$  in at  $z. f\ w \neq 0 \wedge w \in S$ 
using assms(4,5,6) holo non_const non_zero_neighbour_alt by blast
then show  $\exists_F\ w$  in at  $z. f\ w \neq 0$ 
by (auto elim: eventually_frequentlyE)
qed
then obtain  $r1$  where  $g\ z \neq 0\ r1 > 0$  and  $r1: g$  holomorphic_on  $cball\ z\ r1$ 
( $\forall w \in cball\ z\ r1 - \{z\}. f\ w = g\ w * (w - z) ^ powi\ n \wedge g\ w \neq 0$ )
by auto
obtain  $r2$  where  $r2: r2 > 0\ cball\ z\ r2 \subseteq S$ 
using assms(4,6) open_contains_cball_eq by blast
define  $r3$  where  $r3 \equiv \min\ r1\ r2$ 
have  $r3 > 0\ cball\ z\ r3 \subseteq S$  using  $\langle r1 > 0 \rangle\ r2$  unfolding r3_def by auto
moreover have  $g$  holomorphic_on  $cball\ z\ r3$ 
using r1(1) unfolding r3_def by auto
moreover have ( $\forall w \in cball\ z\ r3 - \{z\}. f\ w = g\ w * (w - z) ^ powi\ n \wedge g\ w \neq 0$ )
using r1(2) unfolding r3_def by auto
ultimately show ?thesis using that[of r3]  $\langle g\ z \neq 0 \rangle$  by auto
qed

have fz_lim:  $f - z \rightarrow f\ z$ 
by (metis assms(4,6) at_within_open continuous_on_holo holomorphic_on_imp_continuous_on)
have gz_lim:  $g - z \rightarrow g\ z$ 
using r
by (meson Elementary_Metric_Spaces.open_ball analytic_at_analytic_at_imp_isCont

ball_subset_cball centre_in_ball holomorphic_on_subset isContD)
have if_0: if  $f\ z = 0$  then  $n > 0$  else  $n = 0$ 
proof -

```

```

have ( $\lambda w. g w * (w-z)^{\text{powi } n} - z \rightarrow f z$ )
  using fz_lim Lim_transform_within_open[where  $s=\text{ball } z \ r$ ] by fastforce
then have ( $\lambda w. (g w * (w-z)^{\text{powi } n} / g w) - z \rightarrow f z/g z$ )
  using gz_lim  $\langle g z \neq 0 \rangle$  tendsto_divide by blast
then have powi_tendsto: ( $\lambda w. (w-z)^{\text{powi } n} - z \rightarrow f z/g z$ )
  using Lim_transform_within_open[where  $s=\text{ball } z \ r$ ] by fastforce

have ?thesis when  $n \geq 0 \ f z = 0$ 
proof -
  have ( $\lambda w. (w-z)^{\wedge \text{nat } n} - z \rightarrow f z/g z$ )
    using Lim_transform_within[OF powi_tendsto, where  $d=r$ ]
    by (meson power_int_def r(1) that(1))
  then have *: ( $\lambda w. (w-z)^{\wedge \text{nat } n} - z \rightarrow 0$ ) using  $\langle f z = 0 \rangle$  by simp
  moreover have False when  $n = 0$ 
  proof -
    have ( $\lambda w. (w-z)^{\wedge \text{nat } n} - z \rightarrow 1$ )
      using  $\langle n = 0 \rangle$  by auto
    then show False using * using LIM_unique zero_neq_one by blast
  qed
  ultimately show ?thesis using that by fastforce
qed
moreover have ?thesis when  $n \geq 0 \ f z \neq 0$ 
proof -
  have False when  $n > 0$ 
  proof -
    have ( $\lambda w. (w-z)^{\wedge \text{nat } n} - z \rightarrow f z/g z$ )
      using Lim_transform_within[OF powi_tendsto, where  $d=r$ ]
      by (meson  $\langle 0 \leq n \rangle$  power_int_def r(1))
    moreover have ( $\lambda w. (w-z)^{\wedge \text{nat } n} - z \rightarrow 0$ )
      using  $\langle n > 0 \rangle$  by (auto intro!: tendsto_eq_intros)
    ultimately show False
      using  $\langle f z \neq 0 \rangle \langle g z \neq 0 \rangle$  LIM_unique divide_eq_0_iff by blast
  qed
  then show ?thesis using that by force
qed
moreover have False when  $n < 0$ 
proof -
  have ( $\lambda w. \text{inverse } ((w-z)^{\wedge \text{nat } (-n)}) - z \rightarrow f z/g z$ )
    by (smt (verit) LIM_cong power_int_def power_inverse powi_tendsto that)
  moreover
  have ( $\lambda w. (w-z)^{\wedge \text{nat } (-n)} - z \rightarrow 0$ )
    using that by (auto intro!: tendsto_eq_intros)
  ultimately
  have ( $\lambda x. \text{inverse } ((x-z)^{\wedge \text{nat } (-n)}) * (x-z)^{\wedge \text{nat } (-n)} - z \rightarrow 0$ )
    using tendsto_mult by fastforce
  then have ( $\lambda x. 1::\text{complex}$ )  $- z \rightarrow 0$ 
    using Lim_transform_within_open by fastforce
  then show False
    using LIM_const_eq by fastforce

```

```

qed
ultimately show ?thesis by fastforce
qed
moreover have  $f w = g w * (w-z) ^ \wedge \text{nat } n \wedge g w \neq 0$  when  $w \in \text{cball } z r$  for  $w$ 
proof (cases  $w=z$ )
case True
then have  $f -z \rightarrow f w$ 
using  $fz\_lim$  by blast
then have  $(\lambda w. g w * (w-z) ^ \wedge \text{nat } n) -z \rightarrow f w$ 
proof (elim  $\text{Lim\_transform\_within}[OF\_ \langle r>0 \rangle]$ )
fix  $x$  assume  $0 < \text{dist } x z \text{ dist } x z < r$ 
then have  $x \in \text{cball } z r - \{z\} \ x \neq z$ 
unfolding  $\text{cball\_def}$  by (auto simp add:  $\text{dist\_commute}$ )
then have  $f x = g x * (x - z) \text{ powi } n$ 
using  $r(4)[\text{rule\_format, of } x]$  by simp
also have  $\dots = g x * (x - z) ^ \wedge \text{nat } n$ 
by (smt (verit, best)  $\text{if\_0 int\_nat\_eq power\_int\_of\_nat}$ )
finally show  $f x = g x * (x - z) ^ \wedge \text{nat } n$  .
qed
moreover have  $(\lambda w. g w * (w-z) ^ \wedge \text{nat } n) -z \rightarrow g w * (w-z) ^ \wedge \text{nat } n$ 
using True by (auto intro!:  $\text{tendsto\_eq\_intros gz\_lim}$ )
ultimately have  $f w = g w * (w-z) ^ \wedge \text{nat } n$  using  $\text{LIM\_unique}$  by blast
then show ?thesis using  $\langle g z \neq 0 \rangle$  True by auto
next
case False
then have  $f w = g w * (w-z) \text{ powi } n \ g w \neq 0$ 
using  $r(4)$  that by auto
then show ?thesis
by (smt (verit, best)  $\text{False if\_0 int\_nat\_eq power\_int\_of\_nat}$ )
qed
ultimately show ?thesis using  $r$  by auto
qed

```

lemma zorder_exist_pole :

```

fixes  $f :: \text{complex} \Rightarrow \text{complex}$  and  $z :: \text{complex}$ 
defines  $n \equiv \text{zorder } f z$  and  $g \equiv \text{zor\_poly } f z$ 
assumes  $\text{holo}: f \text{ holomorphic\_on } S - \{z\}$  and  $\text{open } S \ z \in S$  and  $\text{is\_pole } f z$ 
shows  $n < 0 \wedge g z \neq 0 \wedge (\exists r. r > 0 \wedge \text{cball } z r \subseteq S \wedge g \text{ holomorphic\_on } \text{cball } z r$ 
 $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w / (w-z) ^ \wedge \text{nat } (-n) \wedge g w \neq 0))$ 
proof -
obtain  $r$  where  $g z \neq 0$  and  $r: r > 0 \ \text{cball } z r \subseteq S \ g \text{ holomorphic\_on } \text{cball } z r$ 
 $(\forall w \in \text{cball } z r - \{z\}. f w = g w * (w-z) \text{ powi } n \wedge g w \neq 0)$ 
proof -
have  $g z \neq 0 \wedge (\exists r > 0. g \text{ holomorphic\_on } \text{cball } z r$ 
 $\wedge (\forall w \in \text{cball } z r - \{z\}. f w = g w * (w-z) \text{ powi } n \wedge g w \neq 0))$ 
proof (rule  $\text{zorder\_exist}$ [of  $f z, \text{folded } g\_def \ n\_def$ ])
show  $\text{isolated\_singularity\_at } f z$  unfolding  $\text{isolated\_singularity\_at\_def}$ 
using  $\text{holo\_assms}(4,5)$ 
by (metis  $\text{analytic\_on\_holomorphic centre\_in\_ball insert\_Diff openE}$ )

```

```

open_delete subset_insert_iff)
  show not_essential f z unfolding not_essential_def
  using assms(4,6) at_within_open continuous_on_holo holomorphic_on_imp_continuous_on
  by fastforce
  from non_zero_neighbour_pole[OF ‹is_pole f z›] show  $\exists_F w$  in at z.  $f w \neq$ 
0
  by (auto elim: eventually_frequentlyE)
qed
then obtain r1 where  $g z \neq 0$   $r1 > 0$  and  $r1: g$  holomorphic_on cball z r1
  ( $\forall w \in \text{cball } z \text{ } r1 - \{z\}. f w = g w * (w - z)^{\text{powi } n} \wedge g w \neq 0$ )
  by auto
obtain r2 where  $r2: r2 > 0$  cball z r2  $\subseteq S$ 
  using assms(4,5) open_contains_cball_eq by metis
define r3 where  $r3 = \min r1 r2$ 
have  $r3 > 0$  cball z r3  $\subseteq S$  using ‹ $r1 > 0$ ›  $r2$  unfolding r3_def by auto
moreover have  $g$  holomorphic_on cball z r3
  using r1(1) unfolding r3_def by auto
moreover have ( $\forall w \in \text{cball } z \text{ } r3 - \{z\}. f w = g w * (w - z)^{\text{powi } n} \wedge g w \neq 0$ )
  using r1(2) unfolding r3_def by auto
ultimately show ?thesis
  using that[of r3] ‹ $g z \neq 0$ › by auto
qed

have  $n < 0$ 
proof (rule ccontr)
  assume  $\neg n < 0$ 
  define c where  $c = (\text{if } n = 0 \text{ then } g z \text{ else } 0)$ 
  have [simp]:  $g -z \rightarrow g z$ 
  using r
  by (metis centre_in_ball continuous_on_interior holomorphic_on_imp_continuous_on
interior_cball isContD)
  have  $\forall x \in \text{ball } z \text{ } r. x \neq z \rightarrow f x = g x * (x - z)^{\wedge \text{nat } n}$ 
  by (simp add: ‹ $\neg n < 0$ › linorder_not_le power_int_def r)
  then have  $\forall_F x$  in at z.  $f x = g x * (x - z)^{\wedge \text{nat } n}$ 
  using centre_in_ball eventually_at_topological r(1) by blast
  moreover have ( $\lambda x. g x * (x - z)^{\wedge \text{nat } n} -z \rightarrow c$ )
  proof (cases  $n = 0$ )
    case True
    then show ?thesis unfolding c_def by simp
  next
    case False
    then have ( $\lambda x. (x - z)^{\wedge \text{nat } n} -z \rightarrow 0$ ) using ‹ $\neg n < 0$ ›
    by (auto intro!: tendsto_eq_intros)
    from tendsto_mult[OF _ this, of  $g z$ , simplified]
    show ?thesis unfolding c_def using False by simp
  qed
  ultimately have  $f -z \rightarrow c$  using tendsto_cong by fast
then show False using ‹is_pole f z› at_neq_bot not_tendsto_and_filterlim_at_infinity
  unfolding is_pole_def by blast

```



```

qed
moreover have  $\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w / (w-z) \wedge \text{nat } (-n) \wedge g \ w \neq 0$ 
  using  $r(4) \langle n < 0 \rangle$ 
  by (smt (verit) inverse_eq_divide mult.right_neutral power_int_def power_inverse
times_divide_eq_right)
ultimately show ?thesis
  using  $r \langle g \ z \neq 0 \rangle$  by auto
qed

```

lemma *zorder_eqI*:

```

assumes open  $S \ z \in S \ g \ \text{holomorphic\_on } S \ g \ z \neq 0$ 
assumes fg_eq:  $\bigwedge w. \llbracket w \in S; w \neq z \rrbracket \implies f \ w = g \ w * (w-z) \text{ powi } n$ 
shows  $zorder \ f \ z = n$ 

```

proof –

```

have continuous_on  $S \ g$  by (rule holomorphic_on_imp_continuous_on) fact
moreover have open  $(-\{0::\text{complex}\})$  by auto
ultimately have open  $((g -' (-\{0\})) \cap S)$ 
  unfolding continuous_on_open_vimage[OF  $\langle \text{open } S \rangle$ ] by blast
moreover from assms have  $z \in (g -' (-\{0\})) \cap S$  by auto
ultimately obtain  $r$  where  $r: r > 0 \ \text{cball } z \ r \subseteq S \cap (g -' (-\{0\}))$ 
  unfolding open_contains_cball by blast

```

```

let ?gg =  $(\lambda w. g \ w * (w-z) \text{ powi } n)$ 

```

```

define  $P$  where  $P = (\lambda n \ g \ r. 0 < r \wedge g \ \text{holomorphic\_on } \text{cball } z \ r \wedge g \ z \neq 0$ 
 $\wedge (\forall w \in \text{cball } z \ r - \{z\}. f \ w = g \ w * (w-z) \text{ powi } n \wedge g \ w \neq 0))$ 

```

```

have  $P \ n \ g \ r$ 

```

```

  unfolding  $P\_def$  using  $r \ \text{assms}(3,4,5)$  by auto

```

```

then have  $\exists g \ r. P \ n \ g \ r$  by auto

```

```

moreover have unique:  $\exists! n. \exists g \ r. P \ n \ g \ r$  unfolding  $P\_def$ 

```

```

proof (rule holomorphic_factor_puncture)

```

```

  have  $\text{ball } z \ r - \{z\} \subseteq S$  using  $r$  using ball_subset_cball by blast

```

```

  then have ?gg holomorphic_on  $\text{ball } z \ r - \{z\}$ 

```

```

    using  $\langle g \ \text{holomorphic\_on } S \rangle \ r$  by (auto intro!: holomorphic_intros)

```

```

  then have  $f$  holomorphic_on  $\text{ball } z \ r - \{z\}$ 

```

```

    by (smt (verit, best) DiffD2  $\langle \text{ball } z \ r - \{z\} \subseteq S \rangle \ fg\_eq \ \text{holomorphic\_cong}$ 
singleton_iff subset_iff)

```

```

  then show isolated_singularity_at  $f \ z$  unfolding isolated_singularity_at_def

```

```

    using analytic_on_open open_delete  $r(1)$  by blast

```

next

```

  have not_essential ?gg  $z$ 

```

```

  proof (intro singularity_intros)

```

```

    show not_essential  $g \ z$ 

```

```

      by (meson  $\langle \text{continuous\_on } S \rangle \ \text{assms} \ \text{continuous\_on\_eq\_continuous\_at}$ 
isCont_def not_essential_def)

```

```

  show  $\forall_F w \ \text{in at } z. w - z \neq 0$  by (simp add: eventually_at_filter)

```

```

  then show LIM  $w \ \text{at } z. w - z :> \text{at } 0$ 

```

```

    unfolding filterlim_at by (auto intro: tendsto_eq_intros)

```

```

  show isolated_singularity_at  $g \ z$ 

```

```

    by (meson Diff_subset open_ball analytic_on_holomorphic)

```

```

      assms holomorphic_on_subset isolated_singularity_at_def openE)
qed
moreover
have  $\forall_F w \text{ in at } z. g w * (w-z) \text{ powi } n = f w$ 
  unfolding eventually_at_topological using assms fg_eq by force
ultimately show not_essential f z
  using not_essential_transform by blast
show  $\exists_F w \text{ in at } z. f w \neq 0$  unfolding frequently_at
proof (intro strip)
  fix d::real assume  $0 < d$ 
  define z' where  $z' \equiv z + \min d r / 2$ 
  have  $z' \neq z \text{ dist } z' z < d$ 
    unfolding z'_def using  $\langle d > 0 \rangle \langle r > 0 \rangle$  by (auto simp add: dist_norm)
  moreover have  $f z' \neq 0$ 
  proof (subst fg_eq[OF _  $\langle z' \neq z \rangle$ ])
    have  $z' \in \text{cball } z r$ 
      unfolding z'_def using  $\langle r > 0 \rangle \langle d > 0 \rangle$  by (auto simp add: dist_norm)
    then show  $z' \in S$  using  $r(2)$  by blast
  show  $g z' * (z' - z) \text{ powi } n \neq 0$ 
    using P_def  $\langle P n g r \rangle \langle z' \in \text{cball } z r \rangle \langle z' \neq z \rangle$  by auto
  qed
  ultimately show  $\exists x \in UNIV. x \neq z \wedge \text{dist } x z < d \wedge f x \neq 0$  by auto
qed
qed
ultimately have  $(THE n. \exists g r. P n g r) = n$ 
  by (rule_tac the1_equality)
then show ?thesis unfolding zorder_def P_def by blast
qed

lemma simple_zeroI:
  assumes open S z  $\in S$  g holomorphic_on S g z  $\neq 0$ 
  assumes  $\bigwedge w. w \in S \implies f w = g w * (w-z)$ 
  shows zorder f z = 1
  using assms zorder_eqI by force

lemma higher_deriv_power:
  shows  $(\text{deriv } \sim j) (\lambda w. (w-z) \wedge n) w =$ 
    pochhammer (of_nat (Suc n - j)) j *  $(w-z) \wedge (n - j)$ 
proof (induction j arbitrary: w)
  case 0
  thus ?case by auto
next
  case (Suc j w)
  have  $(\text{deriv } \sim \text{Suc } j) (\lambda w. (w-z) \wedge n) w = \text{deriv } ((\text{deriv } \sim j) (\lambda w. (w-z) \wedge$ 
 $n)) w$ 
  by simp
  also have  $(\text{deriv } \sim j) (\lambda w. (w-z) \wedge n) =$ 
     $(\lambda w. \text{pochhammer } (\text{of\_nat } (\text{Suc } n - j)) j * (w-z) \wedge (n - j))$ 
  using Suc by (intro Suc.IH ext)

```

```

also {
  have (... has_field_derivative of_nat (n - j) *
        pochhammer (of_nat (Suc n - j)) j * (w - z) ^ (n - Suc j)) (at w)
    using Suc.premis by (auto intro!: derivative_eq_intros)
  also have of_nat (n - j) * pochhammer (of_nat (Suc n - j)) j =
        pochhammer (of_nat (Suc n - Suc j)) (Suc j)
    by (cases Suc j ≤ n, subst pochhammer_rec)
        (use Suc.premis in ‹simp_all add: algebra_simps Suc_diff_le pochhammer_0_left›)
  finally have deriv (λw. pochhammer (of_nat (Suc n - j)) j * (w - z) ^ (n -
j)) w =
        ... * (w - z) ^ (n - Suc j)
    by (rule DERIV_imp_deriv)
}
finally show ?case .
qed

```

lemma zorder_zero_eqI:

```

assumes f_holo: f holomorphic_on S and open S z ∈ S
assumes zero: ∧i. i < nat n ⇒ (deriv ^^ i) f z = 0
assumes nz: (deriv ^^ nat n) f z ≠ 0 and n ≥ 0
shows zorder f z = n

```

proof -

```

obtain r where [simp]: r > 0 and ball z r ⊆ S
  using ‹open S› ‹z ∈ S› openE by blast
have nz': ∃ w ∈ ball z r. f w ≠ 0
proof (rule ccontr)
  assume ¬ (∃ w ∈ ball z r. f w ≠ 0)
  then have eventually (λu. f u = 0) (nhds z)
    using open_ball ‹0 < r› centre_in_ball eventually_nhds by blast
  then have (deriv ^^ nat n) f z = (deriv ^^ nat n) (λ_. 0) z
    by (intro higher_deriv_cong_ev) auto
  also have (deriv ^^ nat n) (λ_. 0) z = 0
    by (induction n) simp_all
  finally show False using nz by contradiction
qed

```

define zn g where zn = zorder f z and g = zor_poly f z

```

obtain e where e_if: if f z = 0 then 0 < zn else zn = 0 and
  [simp]: e > 0 and cball z e ⊆ ball z r and
  g_holo: g holomorphic_on cball z e and
  e_fac: (∀ w ∈ cball z e. f w = g w * (w - z) ^ nat zn ∧ g w ≠ 0)

```

proof -

```

  have f_holomorphic_on ball z r
    using f_holo ‹ball z r ⊆ S› by auto
  from that zorder_exist_zero[of f ball z r z,simplified,OF this nz',folded zn_def
g_def]
  show thesis by blast
qed

```

```

then obtain zn ≥ 0 g z ≠ 0
  by (metis centre_in_cball less_le_not_le order_refl)

define A where A ≡ (λi. of_nat (i choose (nat zn)) * fact (nat zn) * (deriv ^^
(i - nat zn)) g z)
have deriv_A: (deriv ^^ i) f z = (if zn ≤ int i then A i else 0) for i
proof -
  have eventually (λw. w ∈ ball z e) (nhds z)
    using ⟨cball z e ⊆ ball z r⟩ ⟨e>0⟩ by (intro eventually_nhds_in_open) auto
  hence eventually (λw. f w = (w-z) ^ (nat zn) * g w) (nhds z)
    using e_fac eventually_mono by fastforce
  hence (deriv ^^ i) f z = (deriv ^^ i) (λw. (w-z) ^ nat zn * g w) z
    by (intro higher_deriv_cong_ev) auto
  also have ... = (∑ j=0..i. of_nat (i choose j) *
    (deriv ^^ j) (λw. (w-z) ^ nat zn) z * (deriv ^^ (i - j)) g z)
    using g_holo ⟨e>0⟩
    by (intro higher_deriv_mult[of_ball z e]) (auto intro!: holomorphic_intros)
  also have ... = (∑ j=0..i. if j = nat zn then
    of_nat (i choose nat zn) * fact (nat zn) * (deriv ^^ (i - nat zn))
g z else 0)
proof (intro sum.cong refl, goal_cases)
  case (1 j)
  have (deriv ^^ j) (λw. (w-z) ^ nat zn) z =
    pochhammer (of_nat (Suc (nat zn) - j)) j * 0 ^ (nat zn - j)
    by (subst higher_deriv_power) auto
  also have ... = (if j = nat zn then fact j else 0)
    by (auto simp: not_less pochhammer_0_left pochhammer_fact)
  also have of_nat (i choose j) * ... * (deriv ^^ (i - j)) g z =
    (if j = nat zn then of_nat (i choose (nat zn)) * fact (nat zn)
    * (deriv ^^ (i - nat zn)) g z else 0)
    by simp
  finally show ?case .
qed
also have ... = (if i ≥ zn then A i else 0)
  by (auto simp: A_def)
finally show (deriv ^^ i) f z = ... .
qed

have False when n < zn
  using deriv_A[of nat n] that ⟨n ≥ 0⟩ by (simp add: nz)
moreover have n ≤ zn
proof -
  have g z ≠ 0
    by (simp add: ⟨g z ≠ 0⟩)
  then have (deriv ^^ nat zn) f z ≠ 0
    using deriv_A[of nat zn] by (auto simp add: A_def)
  then have nat zn ≥ nat n using zero[of nat zn] by linarith
  moreover have zn ≥ 0 using e_if by (auto split: if_splits)
  ultimately show ?thesis using nat_le_eq_zle by blast

```

qed
ultimately show *?thesis* unfolding *zn_def* by fastforce
qed

lemma
assumes eventually $(\lambda z. f z = g z)$ (at *z*) $z = z'$
shows *zorder_cong*: *zorder* $f z = \text{zorder } g z'$ and *zor_poly_cong*: *zor_poly* $f z = \text{zor_poly } g z'$
proof -
define *P* where $P = (\lambda f n h r. 0 < r \wedge h \text{ holomorphic_on } \text{cball } z r \wedge h z \neq 0 \wedge (\forall w \in \text{cball } z r - \{z\}. \text{ff } w = h w * (w-z) \text{ powi } n \wedge h w \neq 0))$
have $(\exists r. P f n h r) = (\exists r. P g n h r)$ for *n h*
proof -
have *: $\exists r. P g n h r$ if $\exists r. P f n h r$ and eventually $(\lambda x. f x = g x)$ (at *z*)
for *f g*
proof -
from *that(1)* obtain *r1* where *r1_P*: *P f n h r1* by auto
from *that(2)* obtain *r2* where *r2>0* and *r2_dist*: $\forall x. x \neq z \wedge \text{dist } x z \leq r2 \rightarrow f x = g x$
unfolding *eventually_at_le* by auto
define *r* where $r = \min r1 r2$
have $r > 0$ *h z* $\neq 0$ using *r1_P* $\langle r2 > 0 \rangle$ unfolding *r_def P_def* by auto
moreover have *h holomorphic_on cball z r*
using *r1_P* unfolding *P_def r_def* by auto
moreover have $g w = h w * (w-z) \text{ powi } n \wedge h w \neq 0$ when $w \in \text{cball } z r - \{z\}$ for *w*
proof -
have $f w = h w * (w-z) \text{ powi } n \wedge h w \neq 0$
using *r1_P* that unfolding *P_def r_def* by auto
moreover have $f w = g w$
using *r2_dist* that by (*simp add: dist_commute r_def*)
ultimately show *?thesis* by *simp*
qed
ultimately show *?thesis* unfolding *P_def* by auto
qed
from *assms* have *eq'*: eventually $(\lambda z. g z = f z)$ (at *z*)
by (*simp add: eq_commute*)
show *?thesis*
using * *assms(1)* *eq'* by *blast*
qed
then show *zorder* $f z = \text{zorder } g z'$ *zor_poly* $f z = \text{zor_poly } g z'$
using $\langle z = z' \rangle$ unfolding *P_def zorder_def zor_poly_def* by auto
qed

lemma *zorder_times_analytic'*:
assumes *isolated_singularity_at f z not_essential f z*
assumes *g analytic_on {z}* frequently $(\lambda z. f z * g z \neq 0)$ (at *z*)
shows *zorder* $(\lambda x. f x * g x) z = \text{zorder } f z + \text{zorder } g z$
using *assms* *isolated_singularity_at_analytic not_essential_analytic zorder_times*

by *blast*

lemma *zorder_cmult*:

assumes $c \neq 0$

shows $\text{zorder } (\lambda z. c * f z) z = \text{zorder } f z$

proof –

define P **where**

$$P = (\lambda f n h r. 0 < r \wedge h \text{ holomorphic_on } \text{cball } z r \wedge \\ h z \neq 0 \wedge (\forall w \in \text{cball } z r - \{z\}. f w = h w * (w - z)^{\text{powi } n} \wedge h w \neq 0))$$

have $*$: $P (\lambda x. c * f x) n (\lambda x. c * h x) r$

if $P f n h r c \neq 0$ **for** $f n h r c$

using *that unfolding P_def by (auto intro!: holomorphic_intros)*

have $(\exists h r. P (\lambda x. c * f x) n h r) \longleftrightarrow (\exists h r. P f n h r)$ **for** n

using $*$ [*of f n _ _ c*] $*$ [*of $\lambda x. c * f x n _ _$ inverse c*] $\langle c \neq 0 \rangle$

by (*fastforce simp: field_simps*)

hence (*THE n. $\exists h r. P (\lambda x. c * f x) n h r$*) = (*THE n. $\exists h r. P f n h r$*)

by *simp*

thus *?thesis*

by (*simp add: zorder_def P_def*)

qed

lemma *zorder_uminus [simp]*: $\text{zorder } (\lambda z. -f z) z = \text{zorder } f z$

using *zorder_cmult[of -1 f]* **by** *simp*

lemma *zorder_nonzero_div_power*:

assumes *sz: open S z $\in S$ f holomorphic_on S f z $\neq 0$ and $n > 0$*

shows $\text{zorder } (\lambda w. f w / (w - z)^{\wedge n}) z = - n$

by (*intro zorder_eqI [OF sz]*) (*simp add: inverse_eq_divide power_int_minus*)

lemma *zor_poly_eq*:

assumes *isolated_singularity_at f z not_essential f z $\exists_F w$ in at z. f w $\neq 0$*

shows *eventually* $(\lambda w. \text{zor_poly } f z w = f w * (w - z)^{\text{powi } - \text{zorder } f z})$ (*at z*)

proof –

obtain r **where** $r: r > 0$

$$(\forall w \in \text{cball } z r - \{z\}. f w = \text{zor_poly } f z w * (w - z)^{\text{powi } (\text{zorder } f z)})$$

using *zorder_exist[OF assms]* **by** *blast*

then have $*$: $\forall w \in \text{ball } z r - \{z\}. \text{zor_poly } f z w = f w * (w - z)^{\text{powi } - \text{zorder } f z}$

by (*auto simp: field_simps power_int_minus*)

have *eventually* $(\lambda w. w \in \text{ball } z r - \{z\})$ (*at z*)

using *r eventually_at_ball'[of r z UNIV]* **by** *auto*

thus *?thesis* **by** *eventually_elim (insert *, auto)*

qed

lemma *zor_poly_zero_eq*:

assumes *f holomorphic_on S open S connected S z $\in S$ $\exists w \in S. f w \neq 0$*

shows *eventually* $(\lambda w. \text{zor_poly } f z w = f w / (w - z)^{\wedge \text{nat } (\text{zorder } f z)})$ (*at z*)

proof –

```

obtain  $r$  where  $r: r > 0$ 
  ( $\forall w \in \text{cball } z \ r - \{z\}. f \ w = \text{zor\_poly } f \ z \ w * (w - z) ^{\text{nat } (\text{zorder } f \ z)}$ )
  using zorder_exist_zero[OF assms] by auto
then have  $*$ :  $\forall w \in \text{ball } z \ r - \{z\}. \text{zor\_poly } f \ z \ w = f \ w / (w - z) ^{\text{nat } (\text{zorder } f \ z)}$ 
by (auto simp: field_simps powr_minus)
have eventually ( $\lambda w. w \in \text{ball } z \ r - \{z\}$ ) (at  $z$ )
  using r eventually_at_ball'[of  $r \ z \ UNIV$ ] by auto
thus ?thesis by eventually_elim (insert *, auto)
qed

```

```

lemma zor_poly_pole_eq:
  assumes f_iso: isolated_singularity_at f z is_pole f z
  shows eventually ( $\lambda w. \text{zor\_poly } f \ z \ w = f \ w * (w - z) ^{\text{nat } (- \text{zorder } f \ z)}$ ) (at  $z$ )
proof -
  obtain  $e$  where [simp]:  $e > 0$  and f_holo: f holomorphic_on ball z e - {z}
  using f_iso analytic_imp_holomorphic unfolding isolated_singularity_at_def
by blast
  obtain  $r$  where  $r: r > 0$ 
    ( $\forall w \in \text{cball } z \ r - \{z\}. f \ w = \text{zor\_poly } f \ z \ w / (w - z) ^{\text{nat } (- \text{zorder } f \ z)}$ )
    using zorder_exist_pole[OF f_holo,simplified,OF <is_pole f z>] by auto
  then have  $*$ :  $\forall w \in \text{ball } z \ r - \{z\}. \text{zor\_poly } f \ z \ w = f \ w * (w - z) ^{\text{nat } (- \text{zorder } f \ z)}$ 
  by (auto simp: field_simps)
  have eventually ( $\lambda w. w \in \text{ball } z \ r - \{z\}$ ) (at  $z$ )
    using r eventually_at_ball'[of  $r \ z \ UNIV$ ] by auto
  thus ?thesis by eventually_elim (insert *, auto)
qed

```

```

lemma zor_poly_eqI:
  fixes  $f :: \text{complex} \Rightarrow \text{complex}$  and  $z0 :: \text{complex}$ 
  defines  $n \equiv \text{zorder } f \ z0$ 
  assumes isolated_singularity_at f z0 not_essential f z0  $\exists_F w \text{ in at } z0. f \ w \neq 0$ 
  assumes lim: ((\lambda x. f (g x) * (g x - z0) powi - n) \longrightarrow c) F
  assumes g: filterlim g (at z0) F and  $F \neq \text{bot}$ 
  shows zor_poly f z0 z0 = c
proof -
  from zorder_exist[OF assms(2-4)] obtain  $r$  where
     $r: r > 0$  zor_poly f z0 holomorphic_on cball z0 r
     $\bigwedge w. w \in \text{cball } z0 \ r - \{z0\} \Longrightarrow f \ w = \text{zor\_poly } f \ z0 \ w * (w - z0) \text{ powi } n$ 
  unfolding n_def by blast
  from  $r(1)$  have eventually ( $\lambda w. w \in \text{ball } z0 \ r \wedge w \neq z0$ ) (at  $z0$ )
    using eventually_at_ball'[of  $r \ z0 \ UNIV$ ] by auto
  hence eventually ( $\lambda w. \text{zor\_poly } f \ z0 \ w = f \ w * (w - z0) \text{ powi } - n$ ) (at  $z0$ )
    by eventually_elim (insert r, auto simp: field_simps power_int_minus)
  moreover have continuous_on (ball z0 r) (zor_poly f z0)
    using  $r$  by (intro holomorphic_on_imp_continuous_on) auto
  with  $r$  have isCont (zor_poly f z0) z0
    by (auto simp: continuous_on_eq_continuous_at)

```

hence $(\text{zor_poly } f \ z0 \longrightarrow \text{zor_poly } f \ z0 \ z0) \text{ (at } z0)$
 unfolding *isCont_def* .
 ultimately have $((\lambda w. f \ w * (w - z0) \text{ powi } - n) \longrightarrow \text{zor_poly } f \ z0 \ z0) \text{ (at } z0)$
 by (*blast intro: Lim_transform_eventually*)
 hence $((\lambda x. f \ (g \ x) * (g \ x - z0) \text{ powi } - n) \longrightarrow \text{zor_poly } f \ z0 \ z0) \ F$
 by (*rule filterlim_compose[OF _ g]*)
 from *tendsto_unique[OF <F ≠ bot> this lim]* show ?thesis .
 qed

lemma *zor_poly_zero_eqI*:

fixes $f :: \text{complex} \Rightarrow \text{complex}$ and $z0 :: \text{complex}$
 defines $n \equiv \text{zorder } f \ z0$
 assumes f *holomorphic_on* A *open* A *connected* A $z0 \in A \exists z \in A. f \ z \neq 0$
 assumes *lim*: $((\lambda x. f \ (g \ x) / (g \ x - z0) \wedge \text{nat } n) \longrightarrow c) \ F$
 assumes g : *filterlim* g (at $z0$) F and $F \neq \text{bot}$
 shows $\text{zor_poly } f \ z0 \ z0 = c$

proof -

from *zorder_exist_zero[OF assms(2-6)]* obtain r where

$r: r > 0 \text{ cball } z0 \ r \subseteq A \text{ zor_poly } f \ z0 \text{ holomorphic_on } \text{cball } z0 \ r$

$\bigwedge w. w \in \text{cball } z0 \ r \implies f \ w = \text{zor_poly } f \ z0 \ w * (w - z0) \wedge \text{nat } n$

unfolding *n_def* by *blast*

from $r(1)$ have *eventually* $(\lambda w. w \in \text{ball } z0 \ r \wedge w \neq z0) \text{ (at } z0)$

using *eventually_at_ball'*[of $r \ z0 \ \text{UNIV}$] by *auto*

hence *eventually* $(\lambda w. \text{zor_poly } f \ z0 \ w = f \ w / (w - z0) \wedge \text{nat } n) \text{ (at } z0)$

by *eventually_elim* (*insert* r , *auto simp: field_simps*)

moreover have *continuous_on* $(\text{ball } z0 \ r) (\text{zor_poly } f \ z0)$

using r by (*intro holomorphic_on_imp_continuous_on*) *auto*

with $r(1,2)$ have *isCont* $(\text{zor_poly } f \ z0) \ z0$

by (*auto simp: continuous_on_eq_continuous_at*)

hence $(\text{zor_poly } f \ z0 \longrightarrow \text{zor_poly } f \ z0 \ z0) \text{ (at } z0)$

unfolding *isCont_def* .

ultimately have $((\lambda w. f \ w / (w - z0) \wedge \text{nat } n) \longrightarrow \text{zor_poly } f \ z0 \ z0) \text{ (at } z0)$

by (*blast intro: Lim_transform_eventually*)

hence $((\lambda x. f \ (g \ x) / (g \ x - z0) \wedge \text{nat } n) \longrightarrow \text{zor_poly } f \ z0 \ z0) \ F$

by (*rule filterlim_compose[OF _ g]*)

from *tendsto_unique[OF <F ≠ bot> this lim]* show ?thesis .

qed

lemma *zor_poly_pole_eqI*:

fixes $f :: \text{complex} \Rightarrow \text{complex}$ and $z0 :: \text{complex}$

defines $n \equiv \text{zorder } f \ z0$

assumes f *iso: isolated_singularity_at* $f \ z0$ and *is_pole* $f \ z0$

assumes *lim*: $((\lambda x. f \ (g \ x) * (g \ x - z0) \wedge \text{nat } (-n)) \longrightarrow c) \ F$

assumes g : *filterlim* g (at $z0$) F and $F \neq \text{bot}$

shows $\text{zor_poly } f \ z0 \ z0 = c$

proof -

obtain r where $r: r > 0 \text{ zor_poly } f \ z0 \text{ holomorphic_on } \text{cball } z0 \ r$

proof -

have $\exists_F w$ in at $z0$. $f w \neq 0$
using *non_zero_neighbour_pole*[*OF* $\langle is_pole f z0 \rangle$]
by (*auto elim: eventually_frequentlyE*)
moreover have *not_essential* $f z0$
unfolding *not_essential_def* **using** $\langle is_pole f z0 \rangle$ **by** *simp*
ultimately show *?thesis*
using *that zorder_exist*[*OF* $f_iso, folded n_def$] **by** *auto*
qed
from $r(1)$ **have** *eventually* $(\lambda w. w \in ball\ z0\ r \wedge w \neq z0)$ (at $z0$)
using *eventually_at_ball'*[*of* $r\ z0\ UNIV$] **by** *auto*
have *eventually* $(\lambda w. zor_poly\ f\ z0\ w = f w * (w - z0) \wedge nat\ (-n))$ (at $z0$)
using *zor_poly_pole_eq*[*OF* $f_iso\ \langle is_pole f z0 \rangle$] **unfolding** n_def .
moreover have *continuous_on* $(ball\ z0\ r)$ $(zor_poly\ f\ z0)$
using r **by** (*intro holomorphic_on_imp_continuous_on*) *auto*
with $r(1,2)$ **have** *isCont* $(zor_poly\ f\ z0)$ $z0$
by (*auto simp: continuous_on_eq_continuous_at*)
hence $(zor_poly\ f\ z0 \longrightarrow zor_poly\ f\ z0\ z0)$ (at $z0$)
unfolding *isCont_def* .
ultimately have $((\lambda w. f w * (w - z0) \wedge nat\ (-n)) \longrightarrow zor_poly\ f\ z0\ z0)$ (at $z0$)
by (*blast intro: Lim_transform_eventually*)
hence $((\lambda x. f (g x) * (g x - z0) \wedge nat\ (-n)) \longrightarrow zor_poly\ f\ z0\ z0)$ F
by (*rule filterlim_compose*[*OF* g])
from *tendsto_unique*[*OF* $\langle F \neq bot \rangle$ *this lim*] **show** *?thesis* .
qed

lemma

assumes *is_pole* f $(x :: complex)$ *open* A $x \in A$
assumes $\bigwedge y. y \in A - \{x\} \implies (f\ has_field_derivative\ f'\ y)$ (at y)
shows *is_pole_deriv'*: *is_pole* $f' x$
and *zorder_deriv'*: *zorder* $f' x = zorder\ f\ x - 1$
proof -
have *holo*: *f holomorphic_on* $A - \{x\}$
using *assms* **by** (*subst holomorphic_on_open*) *auto*
obtain r **where** $r: r > 0\ ball\ x\ r \subseteq A$
using *assms*($2,3$) *openE* **by** *blast*
moreover have *open* $(ball\ x\ r - \{x\})$
by *auto*
ultimately have *isolated_singularity_at* $f x$
by (*auto simp: isolated_singularity_at_def analytic_on_open*
intro!: exI[*of* r] *holomorphic_on_subset*[*OF* *holo*])
hence *ev*: $\forall_F w$ in at x . $zor_poly\ f\ x\ w = f w * (w - x) \wedge nat\ (-zorder\ f\ x)$
using $\langle is_pole\ f\ x \rangle$ *zor_poly_pole_eq* **by** *blast*

define P **where** $P = zor_poly\ f\ x$

define n **where** $n = nat\ (-zorder\ f\ x)$

obtain r **where** $r: r > 0\ cball\ x\ r \subseteq A\ P\ holomorphic_on\ cball\ x\ r\ zorder\ f\ x < 0\ P\ x \neq 0$

```

   $\forall w \in \text{cball } x \ r - \{x\}. f \ w = P \ w / (w-x)^\wedge n \wedge P \ w \neq 0$ 
  using P_def assms holo_n_def zorder_exist_pole by blast
  have n:  $n > 0$ 
  using r(4) by (auto simp: n_def)

  have [derivative_intros]: (P has_field_derivative deriv P w) (at w)
  if  $w \in \text{ball } x \ r$  for w
  using that by (intro holomorphic_derivI[OF holomorphic_on_subset[OF r(3),
of ball x r]]) auto

  define D where  $D = (\lambda w. (\text{deriv } P \ w * (w-x) - \text{of\_nat } n * P \ w) / (w-x)^\wedge (n+1))$ 
  define n' where  $n' = n - 1$ 
  have n':  $n = \text{Suc } n'$ 
  using n by (simp add: n'_def)

  have eventually ( $\lambda w. w \in \text{ball } x \ r$ ) (nhds x)
  using  $\langle r > 0 \rangle$  by (intro eventually_nhds_in_open) auto
  hence ev'': eventually ( $\lambda w. w \in \text{ball } x \ r - \{x\}$ ) (at x)
  by (auto simp: eventually_at_filter elim: eventually_mono)

  {
    fix w assume w:  $w \in \text{ball } x \ r - \{x\}$ 
    have ev': eventually ( $\lambda w. w \in \text{ball } x \ r - \{x\}$ ) (nhds w)
    using w by (intro eventually_nhds_in_open) auto

    have §: ( $\text{deriv } P \ w * (w-x)^\wedge n - P \ w * (n * (w-x)^\wedge (n-1))$ ) / ( $(w-x)^\wedge n * (w-x)^\wedge n$ ) = D w
    using w n' by (simp add: divide_simps D_def) (simp add: algebra_simps)
    have (( $\lambda w. P \ w / (w-x)^\wedge n$ ) has_field_derivative D w) (at w)
    by (rule derivative_eq_intros refl | use w § in force)+
    also have ?this  $\longleftrightarrow$  (f has_field_derivative D w) (at w)
    using r by (intro has_field_derivative_cong_ev refl eventually_mono[OF ev'']) auto
    finally have (f has_field_derivative D w) (at w) .
    moreover have (f has_field_derivative f' w) (at w)
    using w r by (intro assms) auto
    ultimately have  $D \ w = f' \ w$ 
    using DERIV_unique by blast
  } note D_eq = this

  have is_pole D x
  unfolding D_def using n  $\langle r > 0 \rangle$   $\langle P \ x \neq 0 \rangle$ 
  by (intro is_pole_basic[where A = ball x r] holomorphic_intros holomorphic_on_subset[OF r(3)]) auto
  also have ?this  $\longleftrightarrow$  is_pole f' x
  by (intro is_pole_cong eventually_mono[OF ev'']) D_eq auto
  finally show is_pole f' x .

```

```

have zorder  $f' x = -int (Suc n)$ 
proof (rule zorder_eqI)
  show open (ball x r)  $x \in ball x r$ 
    using  $\langle r > 0 \rangle$  by auto
  show  $f' w = (deriv P w * (w-x) - of\_nat n * P w) * (w-x) powi (- int (Suc n))$ 
    if  $w \in ball x r$   $w \neq x$  for  $w$ 
    using that  $D\_eq[of w] n$  by (auto simp: D_def power_int_diff power_int_minus
powr_nat' divide_simps)
  qed (use r n in  $\langle auto intro! : holomorphic\_intros \rangle$ )
  thus zorder  $f' x = zorder f x - 1$ 
    using n by (simp add: n_def)
qed

```

lemma

```

assumes is_pole f (x :: complex) isolated_singularity_at f x
shows is_pole_deriv: is_pole (deriv f) x
  and zorder_deriv: zorder (deriv f) x = zorder f x - 1
proof -
  from assms(2) obtain r where  $r > 0$  f analytic_on ball x r - {x}
    by (auto simp: isolated_singularity_at_def)
  hence holo: f holomorphic_on ball x r - {x}
    by (subst (asm) analytic_on_open) auto
  have *:  $x \in ball x r$  open (ball x r) open (ball x r - {x})
    using  $\langle r > 0 \rangle$  by auto
  show is_pole (deriv f) x zorder (deriv f) x = zorder f x - 1
    by (meson * assms(1) holo holomorphic_derivI is_pole_deriv' zorder_deriv')+
qed

```

lemma removable_singularity_deriv':

```

assumes  $f -x \rightarrow c$   $x \in A$  open (A :: complex set)
assumes  $\bigwedge y. y \in A - \{x\} \implies (f \text{ has\_field\_derivative } f' y) (at y)$ 
shows  $\exists c. f' -x \rightarrow c$ 

```

proof -

```

have holo: f holomorphic_on A - {x}
  using assms by (subst holomorphic_on_open) auto

```

define g **where** $g = (\lambda y. \text{if } y = x \text{ then } c \text{ else } f y)$

have deriv_g_eq: $deriv g y = f' y$ **if** $y \in A - \{x\}$ **for** y

proof -

```

have ev: eventually  $(\lambda y. y \in A - \{x\})$  (nhds y)
  using that assms by (intro eventually_nhds_in_open) auto

```

have (f has_field_derivative f' y) (at y)

using assms that **by** auto

also have ?this \longleftrightarrow (g has_field_derivative f' y) (at y)

by (intro has_field_derivative_cong_ev refl eventually_mono[OF ev]) (auto simp: g_def)

finally show ?thesis

by (intro DERIV_imp_deriv assms)

qed

```

have  $g$  holomorphic_on  $A$ 
  unfolding  $g\_def$  using  $assms$   $assms(1)$  holo
  by (intro removable_singularity) auto
hence  $deriv\ g$  holomorphic_on  $A$ 
  by (intro holomorphic_deriv  $assms$ )
hence continuous_on  $A$  ( $deriv\ g$ )
  by (meson holomorphic_on_imp_continuous_on)
hence ( $deriv\ g \longrightarrow deriv\ g\ x$ ) (at  $x$  within  $A$ )
  using  $assms$  by (auto simp: continuous_on_def)
also have  $?thesis \iff (f' \longrightarrow deriv\ g\ x)$  (at  $x$  within  $A$ )
  by (intro filterlim_cong refl) (auto simp: eventually_at_filter deriv_g_eq)
finally have  $f' -x \rightarrow deriv\ g\ x$ 
  using  $\langle open\ A \rangle \langle x \in A \rangle$  by (meson tendsto_within_open)
thus  $?thesis$ 
  by blast

```

qed

lemma removable_singularity_deriv:

```

assumes  $f -x \rightarrow c$  isolated_singularity_at  $f\ x$ 
shows  $\exists c. deriv\ f -x \rightarrow c$ 

```

proof -

```

from  $assms(2)$  obtain  $r$  where  $r: r > 0$   $f$  analytic_on ball  $x\ r - \{x\}$ 
  by (auto simp: isolated_singularity_at_def)
hence holo:  $f$  holomorphic_on ball  $x\ r - \{x\}$ 
  using analytic_imp_holomorphic by blast
show  $?thesis$ 
  using  $assms(1)$ 
proof (rule removable_singularity_deriv')
  show  $x \in ball\ x\ r$  open (ball  $x\ r$ )
    using  $\langle r > 0 \rangle$  by auto
  qed (auto intro!: holomorphic_derivI[OF holo])

```

qed

lemma not_essential_deriv':

```

assumes not_essential  $f\ x$   $x \in A$  open  $A$ 
assumes  $\bigwedge y. y \in A - \{x\} \implies (f\ has\_field\_derivative\ f'\ y)$  (at  $y$ )
shows not_essential  $f'\ x$ 

```

proof -

```

have holo:  $f$  holomorphic_on  $A - \{x\}$ 
  using  $assms$  by (subst holomorphic_on_open) auto
from  $assms$  consider is_pole  $f\ x \mid c$  where  $f -x \rightarrow c$ 
  by (auto simp: not_essential_def)
thus  $?thesis$ 
proof cases
  case 1
  thus  $?thesis$ 
    using  $assms$  is_pole_deriv' by blast

```

```

next
  case (2 c)
  thus ?thesis
    by (meson assms removable_singularity_deriv' tendsto_imp_not_essential)
qed
qed

```

```

lemma not_essential_deriv[singularity_intros]:
  assumes not_essential f x isolated_singularity_at f x
  shows not_essential (deriv f) x
proof -
  from assms(2) obtain r where r: r > 0 f analytic_on ball x r - {x}
  by (auto simp: isolated_singularity_at_def)
  hence holo: f holomorphic_on ball x r - {x}
  by (subst (asm) analytic_on_open) auto
  show ?thesis
  using assms(1)
  proof (rule not_essential_deriv')
    show x ∈ ball x r open (ball x r)
    using ‹r > 0› by auto
  qed (auto intro!: holomorphic_derivI[OF holo])
qed

```

```

lemma not_essential_frequently_0_imp_tendsto_0:
  fixes f :: complex ⇒ complex
  assumes sing: isolated_singularity_at f z not_essential f z
  assumes freq: frequently (λz. f z = 0) (at z)
  shows f -z→ 0
proof -
  from freq obtain g :: nat ⇒ complex where g: filterlim g (at z) at_top ∧ n. f
  (g n) = 0
  using frequently_atE by blast
  have eventually (λx. f (g x) = 0) sequentially
  using g by auto
  hence fg: (λx. f (g x)) → 0
  by (simp add: tendsto_eventually)

```

```

from assms(2) consider c where f -z→ c | is_pole f z
  unfolding not_essential_def by blast
thus ?thesis
proof cases
  case (1 c)
  have (λx. f (g x)) → c
  by (rule filterlim_compose[OF 1 g(1)])
  with fg have c = 0
  using LIMSEQ_unique by blast
  with 1 show ?thesis by simp
next
  case 2

```

```

    have filterlim ( $\lambda x. f (g x)$ ) at_infinity sequentially
      using 2 filterlim_compose g(1) is_pole_def by blast
    with fg have False
      by (meson not_tendsto_and_filterlim_at_infinity sequentially_bot)
    thus ?thesis ..
  qed
qed

lemma not_essential_frequently_0_imp_eventually_0:
  fixes f :: complex  $\Rightarrow$  complex
  assumes sing: isolated_singularity_at f z not_essential f z
  assumes freq: frequently ( $\lambda z. f z = 0$ ) (at z)
  shows eventually ( $\lambda z. f z = 0$ ) (at z)
proof -
  from sing obtain r where r:  $r > 0$  and f_analytic_on_ball z r - {z}
    by (auto simp: isolated_singularity_at_def)
  hence holo: f_holomorphic_on_ball z r - {z}
    by (subst (asm) analytic_on_open) auto
  have eventually ( $\lambda w. w \in \text{ball } z \ r - \{z\}$ ) (at z)
    using r by (intro eventually_at_in_open) auto
  from freq and this have frequently ( $\lambda w. f w = 0 \wedge w \in \text{ball } z \ r - \{z\}$ ) (at z)
    using frequently_eventually_frequently by blast
  hence frequently ( $\lambda w. w \in \{w \in \text{ball } z \ r - \{z\}. f w = 0\}$ ) (at z)
    by (simp add: conj_commute)
  hence limpt: z_islimpt { $w \in \text{ball } z \ r - \{z\}. f w = 0$ }
    using islimpt_conv_frequently_at by blast

  define g where g = ( $\lambda w. \text{if } w = z \text{ then } 0 \text{ else } f w$ )
  have f_minus_z_tendsto_0
    by (intro not_essential_frequently_0_imp_tendsto_0 assms)
  hence g_holo: g_holomorphic_on_ball z r
    unfolding g_def by (intro removable_singularity holo) auto

  have g_eq_0: g w = 0 if w  $\in$  ball z r for w
  proof (rule analytic_continuation[where f = g])
    show open (ball z r) connected (ball z r)
      using r by auto
    show z_islimpt { $w \in \text{ball } z \ r - \{z\}. f w = 0$ }
      by fact
    show g w = 0 if w  $\in$  { $w \in \text{ball } z \ r - \{z\}. f w = 0$ } for w
      using that by (auto simp: g_def)
  qed (use r that g_holo in auto)

  have eventually ( $\lambda w. w \in \text{ball } z \ r - \{z\}$ ) (at z)
    using r by (intro eventually_at_in_open) auto
  thus eventually ( $\lambda w. f w = 0$ ) (at z)
    by (metis freq non_zero_neighbour not_eventually_not_frequently sing)
qed

```

```

lemma pole_imp_not_constant:
  fixes f :: 'a :: {perfect_space}  $\Rightarrow$  _
  assumes is_pole f x open A x  $\in$  A A  $\subseteq$  insert x B
  shows  $\neg$ f constant_on B
proof
  assume *: f constant_on B
  then obtain c where c:  $\forall x \in B. f x = c$ 
    by (auto simp: constant_on_def)
  have eventually  $(\lambda y. y \in A - \{x\})$  (at x)
    using assms by (intro eventually_at_in_open) auto
  hence eventually  $(\lambda y. f y = c)$  (at x)
    by eventually_elim (use c assms in auto)
  hence **: f  $\rightarrow$  c
    by (simp add: tendsto_eventually)
  show False
    using **  $\langle$ is_pole f x $\rangle$  at_neq_bot is_pole_def
      not_tendsto_and_filterlim_at_infinity by blast
qed

lemma neg_zorder_imp_is_pole:
  assumes iso: isolated_singularity_at f z and f_ess: not_essential f z
    and zorder f z < 0 and fre_nz:  $\exists_F w$  in at z. f w  $\neq$  0
  shows is_pole f z
proof -
  define P where P = zor_poly f z
  define n where n = zorder f z
  have n<0 unfolding n_def by (simp add: assms(3))
  define nn where nn = nat (-n)

  obtain r where r: P z  $\neq$  0 r>0 and r_holo: P holomorphic_on cball z r and
    w_Pn:  $(\forall w \in \text{cball } z r - \{z\}. f w = P w * (w-z)^{\text{powi } n} \wedge P w \neq 0)$ 
    using zorder_exist[OF iso f_ess fre_nz, folded P_def n_def] by auto

  have is_pole  $(\lambda w. P w * (w-z)^{\text{powi } n}) z$ 
    unfolding is_pole_def
  proof (rule tendsto_mult_filterlim_at_infinity)
  show P  $\rightarrow$  z by
    by (metis  $\langle$ r>0 $\rangle$  r_holo centre_in_ball continuous_on_interior
      holomorphic_on_imp_continuous_on interior_cball isContD)
  show P z  $\neq$  0 by (simp add:  $\langle$ P z  $\neq$  0 $\rangle$ )

  have LIM x at z. inverse  $((x - z) \wedge \text{nat } (-n))$  :> at_infinity
    apply (subst filterlim_inverse_at_iff[symmetric])
    using  $\langle$ n<0 $\rangle$ 
    by (auto intro!: tendsto_eq_intros filterlim_atI
      simp add: eventually_at_filter)
  then show LIM x at z.  $(x - z)^{\text{powi } n}$  :> at_infinity
    proof (elim filterlim_mono_eventually)

```

```

have inverse ((x - z) ^ nat (-n)) = (x - z) powi n
if x ≠ z for x
by (metis ‹n < 0› linorder_not_le power_int_def power_inverse)
then show ∀F x in at z. inverse ((x - z) ^ nat (-n))
      = (x - z) powi n
by (simp add: eventually_at_filter)
qed auto
qed
moreover have ∀F w in at z. f w = P w * (w - z) powi n
unfolding eventually_at_le
using w_Pn ‹r > 0› by (force simp add: dist_commute)
ultimately show ?thesis using is_pole_cong by fast
qed

lemma is_pole_divide_zorder:
  fixes f g:: complex ⇒ complex and z::complex
  assumes f_iso: isolated_singularity_at f z and g_iso: isolated_singularity_at g
  z
    and f_ess: not_essential f z and g_ess: not_essential g z
    and fg_nconst: ∃F w in (at z). f w * g w ≠ 0
    and z_less: zorder f z < zorder g z
  shows is_pole (λz. f z / g z) z
proof -
  define fn gn fg where fn=zorder f z and gn=zorder g z
    and fg=(λw. f w / g w)

  have isolated_singularity_at fg z
    unfolding fg_def using f_iso g_iso g_ess
    by (auto intro: singularity_intros)
  moreover have not_essential fg z
    unfolding fg_def using f_iso g_iso g_ess f_ess
    by (auto intro: singularity_intros)
  moreover have zorder fg z < 0
proof -
  have zorder fg z = fn - gn
    using zorder_divide[OF f_iso g_iso f_ess g_ess fg_nconst]
    by (simp add: fg_def fn_def gn_def)
  then show ?thesis
    using z_less by (simp add: fn_def gn_def)
qed
moreover have ∃F w in at z. fg w ≠ 0
using fg_nconst unfolding fg_def by force
ultimately show is_pole fg z
using neg_zorder_imp_is_pole by auto
qed

lemma isolated_pole_imp_nzero_times:
  assumes f_iso: isolated_singularity_at f z
    and is_pole f z

```



```

shows  $\exists_F w \text{ in } (at\ z). \text{deriv } f\ w * f\ w \neq 0$ 
proof (rule ccontr)
  assume  $\neg (\exists_F w \text{ in } at\ z. \text{deriv } f\ w * f\ w \neq 0)$ 
  then have  $\forall_F x \text{ in } at\ z. \text{deriv } f\ x * f\ x = 0$ 
    unfolding not_frequently_by simp
  moreover have  $\forall_F w \text{ in } at\ z. f\ w \neq 0$ 
    using non_zero_neighbour_pole[OF ‹is_pole f z›] .
  moreover have  $\forall_F w \text{ in } at\ z. \text{deriv } f\ w \neq 0$ 
    using is_pole_deriv[OF ‹is_pole f z› f_iso, THEN non_zero_neighbour_pole]
  .
  ultimately have  $\forall_F w \text{ in } at\ z. \text{False}$ 
    by eventually_elim auto
  then show False by auto
qed

lemma isolated_pole_imp_neg_zorder:
  assumes isolated_singularity_at f z and is_pole f z
  shows zorder f z < 0
  using analytic_imp_holomorphic_assms centre_in_ball isolated_singularity_at_def
  zorder_exist_pole by blast

lemma isolated_singularity_at_deriv[singularity_intros]:
  assumes isolated_singularity_at f x
  shows isolated_singularity_at (deriv f) x
  by (meson analytic_deriv_assms isolated_singularity_at_def)

lemma zorder_deriv_minus_1:
  fixes f g:: complex  $\Rightarrow$  complex and z::complex
  assumes f_iso: isolated_singularity_at f z
    and f_ess: not_essential f z
    and f_nconst:  $\exists_F w \text{ in } at\ z. f\ w \neq 0$ 
    and f_ord: zorder f z  $\neq 0$ 
  shows zorder (deriv f) z = zorder f z - 1
proof -
  define P where P = zor_poly f z
  define n where n = zorder f z
  have n $\neq 0$  unfolding n_def using f_ord by auto

  obtain r where P z  $\neq 0$  r>0 and P_holo: P holomorphic_on cball z r
    and ( $\forall w \in cball\ z\ r - \{z\}. f\ w$ 
      = P w * (w-z) powi n  $\wedge$  P w  $\neq 0$ )
  using zorder_exist[OF f_iso f_ess f_nconst, folded P_def n_def] by auto
from this(4)
  have f_eq: ( $\forall w \in cball\ z\ r - \{z\}. f\ w$ 
    = P w * (w-z) powi n  $\wedge$  P w  $\neq 0$ )
  using complex_powr_of_int f_ord n_def by presburger

  define D where D = ( $\lambda w. (\text{deriv } P\ w * (w-z) + \text{of\_int } n * P\ w)$ )

```

$*(w-z) \text{ powi } (n - 1)$

have *deriv_f_eq*: $\text{deriv } f \ w = D \ w$ **if** $w \in \text{ball } z \ r - \{z\}$ **for** w

proof -

have *ev'*: *eventually* $(\lambda w. w \in \text{ball } z \ r - \{z\})$ (*nhds* w)

using *that* **by** (*intro eventually_nhds_in_open*) *auto*

define *wz* **where** $wz = w - z$

have $wz \neq 0$ **unfolding** *wz_def* **using** *that* **by** *auto*

moreover **have** (*P* *has_field_derivative* *deriv P w*) (*at w*)

by (*meson DiffD1 Elementary_Metric_Spaces.open_ball P_holo*

ball_subset_cball holomorphic_derivI holomorphic_on_subset that)

ultimately **have** $((\lambda w. P \ w * (w-z) \ \text{powi } n)$ *has_field_derivative* $D \ w$) (*at w*)

unfolding *D_def* **using** *that*

apply (*auto intro!*: *derivative_eq_intros*)

by (*auto simp*: *algebra_simps simp flip:power_int_add_1' wz_def*)

also **have** $?this \longleftrightarrow$ (*f* *has_field_derivative* $D \ w$) (*at w*)

using *f_eq*

by (*intro has_field_derivative_cong_ev refl eventually_mono[OF ev']*) *auto*

ultimately **have** (*f* *has_field_derivative* $D \ w$) (*at w*) **by** *simp*

moreover **have** (*f* *has_field_derivative* *deriv f w*) (*at w*)

by (*metis DERIV_imp_deriv calculation*)

ultimately **show** *?thesis* **using** *DERIV_imp_deriv* **by** *blast*

qed

show *zorder* (*deriv f*) $z = n - 1$

proof (*rule zorder_eqI*)

show *open* ($\text{ball } z \ r$) $z \in \text{ball } z \ r$

using $\langle r > 0 \rangle$ **by** *auto*

define *g* **where** $g = (\lambda w. (\text{deriv } P \ w * (w-z) + \text{of_int } n * P \ w))$

show *g* *holomorphic_on* $\text{ball } z \ r$

unfolding *g_def* **using** *P_holo*

by (*auto intro!*:*holomorphic_intros*)

show $g \ z \neq 0$

unfolding *g_def* **using** $\langle P \ z \neq 0 \rangle \langle n \neq 0 \rangle$ **by** *auto*

show $\text{deriv } f \ w = (\text{deriv } P \ w * (w-z) + \text{of_int } n * P \ w) * (w-z) \ \text{powi } (n - 1)$

if $w \in \text{ball } z \ r \ w \neq z$ **for** w

using *D_def deriv_f_eq* *that* **by** *blast*

qed

qed

lemma *deriv_divide_is_pole*: — Generalises $\llbracket \text{is_pole } ?f \ ?x; \text{isolated_singularity_at } ?f \ ?x \rrbracket \implies \text{zorder } (\text{deriv } ?f) \ ?x = \text{zorder } ?f \ ?x - 1$

fixes $f \ g :: \text{complex} \Rightarrow \text{complex}$ **and** $z :: \text{complex}$

assumes *f_iso*: *isolated_singularity_at* $f \ z$

and *f_ness*: *not_essential* $f \ z$

```

    and fg_nconst:  $\exists_F w \text{ in } (at\ z). \text{deriv } f\ w * f\ w \neq 0$ 
    and f_ord:  $\text{zorder } f\ z \neq 0$ 
    shows is_pole  $(\lambda z. \text{deriv } f\ z / f\ z)\ z$ 
  proof (rule neg_zorder_imp_is_pole)
    define ff where ff =  $(\lambda w. \text{deriv } f\ w / f\ w)$ 
    show isolated_singularity_at ff z
      using f_iso f_ness unfolding ff_def
      by (auto intro: singularity_intros)
    show not_essential ff z
      unfolding ff_def using f_ness f_iso by (auto intro: singularity_intros)

    have zorder ff z =  $\text{zorder } (\text{deriv } f)\ z - \text{zorder } f\ z$ 
      unfolding ff_def using f_iso f_ness fg_nconst
      using isolated_singularity_at_deriv not_essential_deriv zorder_divide by blast
    moreover have  $\text{zorder } (\text{deriv } f)\ z = \text{zorder } f\ z - 1$ 
      using f_iso f_ness f_ord fg_nconst frequently_elim1 zorder_deriv_minus_1
    by fastforce
    ultimately show  $\text{zorder } ff\ z < 0$  by auto

    show  $\exists_F w \text{ in } at\ z. ff\ w \neq 0$ 
      unfolding ff_def using fg_nconst by auto
  qed

```

```

lemma is_pole_deriv_divide_is_pole:
  fixes f g :: complex  $\Rightarrow$  complex and z :: complex
  assumes f_iso: isolated_singularity_at f z
    and is_pole f z
    shows is_pole  $(\lambda z. \text{deriv } f\ z / f\ z)\ z$ 
  proof (rule deriv_divide_is_pole[OF f_iso])
    show not_essential f z
      using <is_pole f z> unfolding not_essential_def by auto
    show  $\exists_F w \text{ in } at\ z. \text{deriv } f\ w * f\ w \neq 0$ 
      using assms f_iso isolated_pole_imp_nzero_times by blast
    show  $\text{zorder } f\ z \neq 0$ 
      using isolated_pole_imp_neg_zorder assms by fastforce
  qed

```

7.13 Isolated points

```

definition isolated_points_of :: complex set  $\Rightarrow$  complex set where
  isolated_points_of A =  $\{z \in A. \text{eventually } (\lambda w. w \notin A)\ (at\ z)\}$ 

```

```

lemma isolated_points_of_altdef: isolated_points_of A =  $\{z \in A. \neg z \text{ islimpt } A\}$ 
  unfolding isolated_points_of_def islimpt_def eventually_at_filter eventually_nhds
  by blast

```

```

lemma isolated_points_of_empty [simp]: isolated_points_of {} = {}
  and isolated_points_of_UNIV [simp]: isolated_points_of UNIV = {}
  by (auto simp: isolated_points_of_def)

```

lemma *isolated_points_of_open_is_empty* [simp]: $open\ A \implies isolated_points_of\ A = \{\}$

unfolding *isolated_points_of_altdef*

by (*simp add: interior_limit_point interior_open*)

lemma *isolated_points_of_subset*: $isolated_points_of\ A \subseteq A$

by (*auto simp: isolated_points_of_def*)

lemma *isolated_points_of_discrete*:

assumes *discrete A*

shows $isolated_points_of\ A = A$

using *assms* **by** (*auto simp: isolated_points_of_def discrete_altdef*)

lemmas *uniform_discreteI1 = uniformI1*

lemmas *uniform_discreteI2 = uniformI2*

lemma *zorder_zero_eqI'*:

assumes *f analytic_on {z}*

assumes $\bigwedge i. i < nat\ n \implies (deriv\ \sim i)\ f\ z = 0$

assumes $(deriv\ \sim nat\ n)\ f\ z \neq 0$ **and** $n \geq 0$

shows $zorder\ f\ z = n$

proof –

from *assms(1)* **obtain** *A* **where** *open A z ∈ A f holomorphic_on A*

using *analytic_at by blast*

thus *?thesis*

using *zorder_zero_eqI[of f A z n] assms* **by** *blast*

qed

7.14 Isolated zeros

definition *isolated_zero* :: $('a::topological_space \Rightarrow 'b::real_normed_div_algebra)$

$\Rightarrow 'a \Rightarrow bool$ **where**

$isolated_zero\ f\ a \iff f\ -a \rightarrow 0 \wedge eventually\ (\lambda x. f\ x \neq 0)\ (at\ a)$

lemma *isolated_zero_shift*:

fixes $z :: 'a :: real_normed_vector$

shows $isolated_zero\ f\ z \iff isolated_zero\ (\lambda w. f\ (z + w))\ 0$

unfolding *isolated_zero_def*

by (*simp add: at_to_0' eventually_filtermap filterlim_filtermap add_ac*)

lemma *isolated_zero_shift'*:

fixes $z :: 'a :: real_normed_vector$

assumes *NO_MATCH 0 z*

shows $isolated_zero\ f\ z \iff isolated_zero\ (\lambda w. f\ (z + w))\ 0$

by (*rule isolated_zero_shift*)

lemma *isolated_zero_imp_not_essential* [*intro*]:

$isolated_zero\ f\ z \implies not_essential\ f\ z$

unfolding *isolated_zero_def not_essential_def*
using *tendsto_nhds_iff* **by** *blast*

lemma *pole_is_not_zero*:
fixes *f* :: 'a::perfect_space \Rightarrow 'b::real_normed_field
assumes *is_pole* *f* *z*
shows \neg *isolated_zero* *f* *z*
proof
assume *isolated_zero* *f* *z*
then have *filterlim* *f* (*nhds* 0) (at *z*)
unfolding *isolated_zero_def* **using** *tendsto_nhds_iff* **by** *blast*
moreover have *filterlim* *f* *at_infinity* (at *z*)
using \langle *is_pole* *f* *z* \rangle **unfolding** *is_pole_def* .
ultimately show *False*
using *not_tendsto_and_filterlim_at_infinity[OF at_neq_bot]*
by *auto*
qed

lemma *isolated_zero_imp_pole_inverse*:
fixes *f* :: $_ \Rightarrow$ 'b::{real_normed_div_algebra, division_ring}
assumes *isolated_zero* *f* *z*
shows *is_pole* (λz . *inverse* (*f* *z*)) *z*
proof –
from *assms* **have** *ev*: *eventually* (λz . *f* *z* \neq 0) (at *z*)
by (*auto simp: isolated_zero_def*)
have *filterlim* *f* (*nhds* 0) (at *z*)
using *assms* **by** (*simp add: isolated_zero_def*)
with *ev* **have** *filterlim* *f* (at 0) (at *z*)
using *filterlim_atI* **by** *blast*
also have $?this \iff$ *filterlim* (λz . *inverse* (*inverse* (*f* *z*))) (at 0) (at *z*)
by (*rule filterlim_cong*) (use *ev* **in** \langle *auto elim!: eventually_mono* \rangle)
finally have *filterlim* (λz . *inverse* (*f* *z*)) *at_infinity* (at *z*)
by (*subst filterlim_inverse_at_iff [symmetric]*)
thus $?thesis$
by (*simp add: is_pole_def*)
qed

lemma *is_pole_imp_isolated_zero_inverse*:
fixes *f* :: $_ \Rightarrow$ 'b::{real_normed_div_algebra, division_ring}
assumes *is_pole* *f* *z*
shows *isolated_zero* (λz . *inverse* (*f* *z*)) *z*
proof –
from *assms* **have** *ev*: *eventually* (λz . *f* *z* \neq 0) (at *z*)
by (*simp add: non_zero_neighbour_pole*)
have *filterlim* *f* *at_infinity* (at *z*)
using *assms* **by** (*simp add: is_pole_def*)
also have $?this \iff$ *filterlim* (λz . *inverse* (*inverse* (*f* *z*))) *at_infinity* (at *z*)
by (*rule filterlim_cong*) (use *ev* **in** \langle *auto elim!: eventually_mono* \rangle)
finally have *filterlim* (λz . *inverse* (*f* *z*)) (at 0) (at *z*)

```

  by (subst (asm) filterlim_inverse_at_iff [symmetric]) auto
  hence filterlim ( $\lambda z. \text{inverse } (f z)$ ) (nhds 0) (at z)
  using filterlim_at by blast
  moreover have eventually ( $\lambda z. \text{inverse } (f z) \neq 0$ ) (at z)
  using ev by eventually_elim auto
  ultimately show ?thesis
  by (simp add: isolated_zero_def)
qed

```

```

lemma is_pole_inverse_iff: is_pole ( $\lambda z. \text{inverse } (f z)$ ) z  $\longleftrightarrow$  isolated_zero f z
  using is_pole_imp_isolated_zero_inverse isolated_zero_imp_pole_inverse by
  fastforce

```

```

lemma isolated_zero_inverse_iff: isolated_zero ( $\lambda z. \text{inverse } (f z)$ ) z  $\longleftrightarrow$  is_pole
  f z
  using is_pole_imp_isolated_zero_inverse isolated_zero_imp_pole_inverse by
  fastforce

```

```

lemma zero_isolated_zero:
  fixes f :: 'a :: {t2_space, perfect_space}  $\Rightarrow$  _
  assumes isolated_zero f z isCont f z
  shows f z = 0
proof (rule tendsto_unique)
  show f  $-z \rightarrow$  f z
  using assms(2) by (rule isContD)
  show f  $-z \rightarrow$  0
  using assms(1) by (simp add: isolated_zero_def)
qed auto

```

```

lemma zero_isolated_zero_analytic:
  assumes isolated_zero f z f analytic_on {z}
  shows f z = 0
  using assms(1) analytic_at_imp_isCont[OF assms(2)] by (rule zero_isolated_zero)

```

```

lemma isolated_zero_analytic_iff:
  assumes f analytic_on {z}
  shows isolated_zero f z  $\longleftrightarrow$  f z = 0  $\wedge$  eventually ( $\lambda z. f z \neq 0$ ) (at z)
proof safe
  assume f z = 0 eventually ( $\lambda z. f z \neq 0$ ) (at z)
  with assms show isolated_zero f z
  unfolding isolated_zero_def by (metis analytic_at_imp_isCont isCont_def)
qed (use zero_isolated_zero_analytic[OF _ assms] in  $\langle$  auto simp: isolated_zero_def  $\rangle$ )

```

```

lemma non_isolated_zero_imp_eventually_zero:
  assumes f analytic_on {z} f z = 0  $\neg$  isolated_zero f z
  shows eventually ( $\lambda z. f z = 0$ ) (at z)
proof (rule not_essential_frequently_0_imp_eventually_0)
  from assms(1) show isolated_singularity_at f z not_essential f z
  by (simp_all add: isolated_singularity_at_analytic not_essential_analytic)

```

```

from assms(1,2) have  $f \rightarrow z \rightarrow 0$ 
  by (metis analytic_at_imp_isCont continuous_within)
thus frequently ( $\lambda z. f z = 0$ ) (at z)
  using assms(2,3) by (auto simp: isolated_zero_def frequently_def)
qed

```

```

lemma non_isolated_zero_imp_eventually_zero':
  assumes  $f$  analytic_on { $z$ }  $f z = 0$   $\neg$ isolated_zero  $f z$ 
  shows eventually ( $\lambda z. f z = 0$ ) (nhds z)
  using non_isolated_zero_imp_eventually_zero[OF assms] assms(2)
  using eventually_nhds_conv_at by blast

```

```

end
theory Complex_Residues
  imports Complex_Singularities
begin

```

7.15 Definition of residues

Wenda Li and LC Paulson (2016). A Formal Proof of Cauchy's Residue Theorem. Interactive Theorem Proving

```

definition residue :: (complex  $\Rightarrow$  complex)  $\Rightarrow$  complex  $\Rightarrow$  complex where
  residue  $f z =$  (SOME int.  $\exists e > 0. \forall \varepsilon > 0. \varepsilon < e \rightarrow$ 
     $\rightarrow$  (f has_contour_integral  $2 * \pi * i * \text{int}$ ) (circlepath z  $\varepsilon$ ))

```

```

lemma residue_cong:

```

```

  assumes eq: eventually ( $\lambda z. f z = g z$ ) (at z) and  $z = z'$ 
  shows residue  $f z = \text{residue } g z'$ 

```

```

proof –

```

```

  from assms have eq': eventually ( $\lambda z. g z = f z$ ) (at z)
    by (simp add: eq_commute)

```

```

  let  $?P = \lambda f c e. (\forall \varepsilon > 0. \varepsilon < e \rightarrow$ 
    (f has_contour_integral_of_real  $(2 * \pi) * i * c$ ) (circlepath z  $\varepsilon$ ))

```

```

  have residue  $f z = \text{residue } g z$  unfolding residue_def

```

```

  proof (rule Eps_cong)

```

```

    fix  $c :: \text{complex}$ 

```

```

    have  $\exists e > 0. ?P g c e$ 

```

```

    if  $\exists e > 0. ?P f c e$  and eventually ( $\lambda z. f z = g z$ ) (at z) for  $f g$ 

```

```

    proof –

```

```

      from that(1) obtain  $e$  where  $e > 0$   $?P f c e$ 

```

```

      by blast

```

```

      from that(2) obtain  $e'$  where  $e' > 0 \wedge z'. z' \neq z \implies \text{dist } z' z < e' \implies$ 
 $f z' = g z'$ 

```

```

      unfolding eventually_at by blast

```

```

      have  $?P g c (\min e e')$ 

```

```

      proof (intro allI exI impI, goal_cases)

```

```

        case (1  $\varepsilon$ )

```

```

        hence (f has_contour_integral_of_real  $(2 * \pi) * i * c$ ) (circlepath z  $\varepsilon$ )

```

```

        using  $e(2)$  by auto

```

```

thus ?case
proof (rule has_contour_integral_eq)
  fix z' assume z' ∈ path_image (circlepath z ε)
  hence dist z' z < e' and z' ≠ z
  using 1 by (auto simp: dist_commute)
  with e'(2)[of z'] show f z' = g z' by simp
qed
qed
moreover from e and e' have min e e' > 0 by auto
ultimately show ?thesis by blast
qed
from this[OF _ eq] and this[OF _ eq]
  show (∃ e>0. ?P f c e) ↔ (∃ e>0. ?P g c e)
  by blast
qed
with assms show ?thesis by simp
qed

lemma residue_shift_0: residue f z = residue (λx. f (z + x)) 0
proof -
  define Q where
    Q = (λr f z ε. (f has_contour_integral complex_of_real (2 * pi) * i * r)
      (circlepath z ε))
  define P where
    P = (λr f z. ∃ e>0. ∀ ε>0. ε < e → Q r f z ε)
  have path_eq: circlepath (z - w) ε = (+) (-w) o circlepath z ε for z w ε
  by (simp add: circlepath_def o_def part_circlepath_def algebra_simps)
  have *: P r f z if P r (λx. f (x + w)) (z - w) for r w f z
  using that by (auto simp: P_def Q_def path_eq has_contour_integral_translate)
  have (SOME r. P r f z) = (SOME r. P r (λx. f (z + x)) 0)
  using *[of _ f z z] *[of _ λx. f (z + x) -z]
  by (intro arg_cong[where f = Eps] ext iffI) (simp_all add: add_ac)
  thus ?thesis
  by (simp add: residue_def P_def Q_def)
qed

lemma residue_shift_0': NO_MATCH 0 z ⇒ residue f z = residue (λx. f (z +
x)) 0
  by (rule residue_shift_0)

lemma contour_integral_circlepath_eq:
  assumes open s and f_holo: f holomorphic_on (s - {z}) and 0 < e1 e1 ≤ e2
  and e2_cball: cball z e2 ⊆ s
  shows
    f contour_integrable_on circlepath z e1
    f contour_integrable_on circlepath z e2
    contour_integral (circlepath z e2) f = contour_integral (circlepath z e1) f
proof -
  define l where l ≡ linepath (z + e2) (z + e1)

```



```

  have [simp]: valid_path l pathstart l=z+e2 pathfinish l=z+e1 unfolding l_def
  by auto
  have e2>0 using ⟨e1>0⟩ ⟨e1≤e2⟩ by auto
  have zℓ_img:z∉path_image l
  proof
    assume z ∈ path_image l
    then have e2 ≤ cmod (e2 - e1)
      using segment_furthest_le[of z z+e2 z+e1 z+e2,simplified] ⟨e1>0⟩ ⟨e2>0⟩
  unfolding l_def
    by (auto simp add:closed_segment_commute)
  thus False using ⟨e2>0⟩ ⟨e1>0⟩ ⟨e1≤e2⟩
    apply (subst (asm) norm_of_real)
    by auto
  qed
  define g where g ≡ circlepath z e2 +++ l +++ reversepath (circlepath z e1)
  +++ reversepath l
  show [simp]: f contour_integrable_on circlepath z e2 f contour_integrable_on
  (circlepath z e1)
  proof -
    show f contour_integrable_on circlepath z e2
      apply (intro contour_integrable_continuous_circlepath[OF
        continuous_on_subset[OF holomorphic_on_imp_continuous_on[OF
f_holo]]])
      using ⟨e2>0⟩ e2_cball by auto
    show f contour_integrable_on (circlepath z e1)
      apply (intro contour_integrable_continuous_circlepath[OF
        continuous_on_subset[OF holomorphic_on_imp_continuous_on[OF
f_holo]]])
      using ⟨e1>0⟩ ⟨e1≤e2⟩ e2_cball by auto
  qed
  have [simp]:f contour_integrable_on l
  proof -
    have closed_segment (z + e2) (z + e1) ⊆ cball z e2 using ⟨e2>0⟩ ⟨e1>0⟩
    ⟨e1≤e2⟩
      by (intro closed_segment_subset,auto simp add:dist_norm)
    hence closed_segment (z + e2) (z + e1) ⊆ s - {z} using zℓ_img e2_cball
  unfolding l_def
    by auto
  then show f contour_integrable_on l unfolding l_def
    apply (intro contour_integrable_continuous_linepath[OF
        continuous_on_subset[OF holomorphic_on_imp_continuous_on[OF
f_holo]]])
    by auto
  qed
  let ?ig=λg. contour_integral g f
  have (f has_contour_integral 0) g
  proof (rule Cauchy_theorem_global[OF _ f_holo])
    show open (s - {z}) using ⟨open s⟩ by auto
    show valid_path g unfolding g_def l_def by auto
  end

```

```

    show pathfinish g = pathstart g unfolding g_def l_def by auto
  next
    have path_img:path_image g  $\subseteq$  cball z e2
      proof -
        have closed_segment (z + e2) (z + e1)  $\subseteq$  cball z e2 using  $\langle e2 > 0 \rangle$   $\langle e1 > 0 \rangle$ 
         $\langle e1 \leq e2 \rangle$ 
          by (intro closed_segment_subset, auto simp add:dist_norm)
        moreover have sphere z |e1|  $\subseteq$  cball z e2 using  $\langle e2 > 0 \rangle$   $\langle e1 \leq e2 \rangle$   $\langle e1 > 0 \rangle$ 
      by auto
      ultimately show ?thesis unfolding g_def l_def using  $\langle e2 > 0 \rangle$ 
        by (simp add: path_image_join closed_segment_commute)
    qed
  show path_image g  $\subseteq$  s - {z}
    proof -
      have  $z \notin$  path_image g using z_l_img
      unfolding g_def l_def by (auto simp add: path_image_join closed_segment_commute)
      moreover note  $\langle$  cball z e2  $\subseteq$  s  $\rangle$  and path_img
      ultimately show ?thesis by auto
    qed
  show winding_number g w = 0 when  $w \notin$  s - {z} for w
    proof -
      have winding_number g w = 0 when  $w \notin$  s using that e2_cball
      apply (intro winding_number_zero_outside[OF _____ path_img])
      by (auto simp add:g_def l_def)
      moreover have winding_number g z = 0
      proof -
        let ?Wz =  $\lambda$ g. winding_number g z
        have ?Wz g = ?Wz (circlepath z e2) + ?Wz l + ?Wz (reversepath
        (circlepath z e1))
          + ?Wz (reversepath l)
          using  $\langle e2 > 0 \rangle$   $\langle e1 > 0 \rangle$  z_l_img unfolding g_def l_def
          by (subst winding_number_join, auto simp add:path_image_join
          closed_segment_commute)+
        also have ... = ?Wz (circlepath z e2) + ?Wz (reversepath (circlepath
        z e1))
          using z_l_img
          apply (subst (2) winding_number_reversepath)
          by (auto simp add:l_def closed_segment_commute)
        also have ... = 0
      proof -
        have ?Wz (circlepath z e2) = 1 using  $\langle e2 > 0 \rangle$ 
          by (auto intro: winding_number_circlepath_centre)
        moreover have ?Wz (reversepath (circlepath z e1)) = -1 using
         $\langle e1 > 0 \rangle$ 
          apply (subst winding_number_reversepath)
          by (auto intro: winding_number_circlepath_centre)
        ultimately show ?thesis by auto
      qed
    qed
  finally show ?thesis .

```

```

      qed
      ultimately show ?thesis using that by auto
    qed
  qed
  then have 0 = ?ig g using contour_integral_unique by simp
  also have ... = ?ig (circlepath z e2) + ?ig l + ?ig (reversepath (circlepath z e1))
    + ?ig (reversepath l)
    unfolding g_def
    by (auto simp add:contour_integrable_reversepath_eq)
  also have ... = ?ig (circlepath z e2) - ?ig (circlepath z e1)
    by (auto simp add:contour_integral_reversepath)
  finally show contour_integral (circlepath z e2) f = contour_integral (circlepath
z e1) f
    by simp
  qed

lemma base_residue:
  assumes open s z∈s r>0 and f_holo:f holomorphic_on (s - {z})
    and r_cball:cball z r ⊆ s
  shows (f has_contour_integral 2 * pi * i * (residue f z)) (circlepath z r)
proof -
  obtain e where e>0 and e_cball:cball z e ⊆ s
    using open_contains_cball[of s] ⟨open s⟩ ⟨z∈s⟩ by auto
  define c where c ≡ 2 * pi * i
  define i where i ≡ contour_integral (circlepath z e) f / c
  have (f has_contour_integral c*i) (circlepath z e) when ε>0 ε<e for ε
  proof -
    have contour_integral (circlepath z e) f = contour_integral (circlepath z ε) f
      f contour_integrable_on circlepath z ε
      f contour_integrable_on circlepath z e
    using ⟨ε<e⟩
      by (intro contour_integral_circlepath_eq[OF ⟨open s⟩ f_holo ⟨ε>0⟩ _
e_cball],auto)+
    then show ?thesis unfolding i_def c_def
      by (auto intro:has_contour_integral_integral)
  qed
  then have ∃ e>0. ∀ ε>0. ε<e → (f has_contour_integral c * (residue f z))
(circlepath z ε)
    unfolding residue_def c_def
    apply (rule_tac someI[of _ i],intro exI[where x=e])
    by (auto simp add:⟨e>0⟩ c_def)
  then obtain e' where e'>0
    and e'_def:∀ ε>0. ε<e' → (f has_contour_integral c * (residue f z))
(circlepath z ε)
    by auto
  let ?int=λe. contour_integral (circlepath z e) f
  define ε where ε ≡ Min {r,e'} / 2
  have ε>0 ε≤r ε<e' using ⟨r>0⟩ ⟨e'>0⟩ unfolding ε_def by auto
  have (f has_contour_integral c * (residue f z)) (circlepath z ε)

```

```

    using e'_def[rule_format,OF ‹ $\epsilon > 0$ › ‹ $\epsilon < e'$ ›] .
  then show ?thesis unfolding c_def
    using contour_integral_circlepath_eq[OF ‹open s› f_holo ‹ $\epsilon > 0$ › ‹ $\epsilon \leq r$ › r_cball]
    by (auto elim: has_contour_integral_eqpath[of _ _ circlepath z  $\epsilon$  circlepath z
r])
qed

```

lemma residue_holo:

```

  assumes open s z  $\in$  s and f_holo: f holomorphic_on s
  shows residue f z = 0
proof -
  define c where c  $\equiv$  2 * pi * i
  obtain e where e > 0 and e_cball: cball z e  $\subseteq$  s using ‹open s› ‹z  $\in$  s›
  using open_contains_cball_eq by blast
  have (f has_contour_integral c * residue f z) (circlepath z e)
  using f_holo
  by (auto intro: base_residue[OF ‹open s› ‹z  $\in$  s› ‹e > 0› _ e_cball, folded c_def])
  moreover have (f has_contour_integral 0) (circlepath z e)
  using f_holo e_cball ‹e > 0›
  by (auto intro: Cauchy_theorem_convex_simple[of _ cball z e])
  ultimately have c * residue f z = 0
  using has_contour_integral_unique by blast
  thus ?thesis unfolding c_def by auto
qed

```

lemma residue_const: residue ($\lambda _.$ c) z = 0

```

  by (intro residue_holo[of UNIV::complex set], auto intro: holomorphic_intros)

```

lemma residue_add:

```

  assumes open s z  $\in$  s and f_holo: f holomorphic_on s - {z}
  and g_holo: g holomorphic_on s - {z}
  shows residue ( $\lambda z.$  f z + g z) z = residue f z + residue g z
proof -
  define c where c  $\equiv$  2 * pi * i
  define fg where fg  $\equiv$  ( $\lambda z.$  f z + g z)
  obtain e where e > 0 and e_cball: cball z e  $\subseteq$  s using ‹open s› ‹z  $\in$  s›
  using open_contains_cball_eq by blast
  have (fg has_contour_integral c * residue fg z) (circlepath z e)
  unfolding fg_def using f_holo g_holo
  apply (intro base_residue[OF ‹open s› ‹z  $\in$  s› ‹e > 0› _ e_cball, folded c_def])
  by (auto intro: holomorphic_intros)
  moreover have (fg has_contour_integral c * residue f z + c * residue g z) (circlepath
z e)
  unfolding fg_def using f_holo g_holo
  by (auto intro: has_contour_integral_add base_residue[OF ‹open s› ‹z  $\in$  s›
‹e > 0› _ e_cball, folded c_def])
  ultimately have c * (residue f z + residue g z) = c * residue fg z
  using has_contour_integral_unique by (auto simp add: distrib_left)
  thus ?thesis unfolding fg_def

```

by (auto simp add:c_def)
qed

lemma *residue_lmul*:

assumes *open s z ∈ s* and *f_holo: f holomorphic_on s - {z}*
shows *residue (λz. c * (f z)) z = c * residue f z*
proof (cases *c=0*)
case *True*
thus ?thesis using *residue_const* by auto
next
case *False*
define *c'* where *c' ≡ 2 * pi * i*
define *f'* where *f' ≡ (λz. c * (f z))*
obtain *e* where *e > 0* and *e_cball: cball z e ⊆ s* using ⟨*open s*⟩ ⟨*z ∈ s*⟩
using *open_contains_cball_eq* by blast
have (*f'* has_contour_integral *c' * residue f' z*) (*circlepath z e*)
unfolding *f'_def* using *f_holo*
apply (intro *base_residue[OF ⟨open s⟩ ⟨z ∈ s⟩ ⟨e > 0⟩ _ e_cball, folded c'_def]*)
by (auto intro: *holomorphic_intros*)
moreover have (*f'* has_contour_integral *c * (c' * residue f z)*) (*circlepath z e*)
unfolding *f'_def* using *f_holo*
by (auto intro: *has_contour_integral_lmul*
base_residue[OF ⟨open s⟩ ⟨z ∈ s⟩ ⟨e > 0⟩ _ e_cball, folded c'_def])
ultimately have *c' * residue f' z = c * (c' * residue f z)*
using *has_contour_integral_unique* by auto
thus ?thesis unfolding *f'_def c'_def* using *False*
by (auto simp add: *field_simps*)
qed

lemma *residue_rmul*:

assumes *open s z ∈ s* and *f_holo: f holomorphic_on s - {z}*
shows *residue (λz. (f z) * c) z = residue f z * c*
using *residue_lmul[OF assms, of c]* by (auto simp add: *algebra_simps*)

lemma *residue_div*:

assumes *open s z ∈ s* and *f_holo: f holomorphic_on s - {z}*
shows *residue (λz. (f z) / c) z = residue f z / c*
using *residue_lmul[OF assms, of 1/c]* by (auto simp add: *algebra_simps*)

lemma *residue_neg*:

assumes *open s z ∈ s* and *f_holo: f holomorphic_on s - {z}*
shows *residue (λz. - (f z)) z = - residue f z*
using *residue_lmul[OF assms, of -1]* by auto

lemma *residue_diff*:

assumes *open s z ∈ s* and *f_holo: f holomorphic_on s - {z}*
and *g_holo: g holomorphic_on s - {z}*
shows *residue (λz. f z - g z) z = residue f z - residue g z*
using *residue_add[OF assms(1,2,3), of λz. - g z]* *residue_neg[OF assms(1,2,4)]*

by (auto intro:holomorphic_intros g_holo)

lemma *residue_simple*:

assumes *open s z ∈ s* and *f_holo: f holomorphic_on s*

shows *residue* $(\lambda w. f w / (w - z)) z = f z$

proof –

define *c* where $c \equiv 2 * \pi * i$

define *f'* where $f' \equiv \lambda w. f w / (w - z)$

obtain *e* where $e > 0$ and *e_cball: cball z e ⊆ s* using $\langle open s \rangle \langle z \in s \rangle$

using *open_contains_cball_eq* by blast

have (*f' has_contour_integral c * f z*) (*circlepath z e*)

unfolding *f'_def c_def* using $\langle e > 0 \rangle$ *f_holo e_cball*

by (auto intro!: *Cauchy_integral_circlepath_simple holomorphic_intros*)

moreover have (*f' has_contour_integral c * residue f' z*) (*circlepath z e*)

unfolding *f'_def* using *f_holo*

apply (intro *base_residue[OF \langle open s \rangle \langle z \in s \rangle \langle e > 0 \rangle _ e_cball, folded c_def]*)

by (auto intro!:holomorphic_intros)

ultimately have $c * f z = c * \text{residue } f' z$

using *has_contour_integral_unique* by blast

thus *?thesis* unfolding *c_def f'_def* by auto

qed

lemma *residue_simple'*:

assumes *s: open s z ∈ s* and *holo: f holomorphic_on (s - {z})*

and *lim: ((λw. f w * (w - z)) → c) (at z)*

shows *residue* $f z = c$

proof –

define *g* where $g = (\lambda w. \text{if } w = z \text{ then } c \text{ else } f w * (w - z))$

from *holo* have $(\lambda w. f w * (w - z)) \text{ holomorphic_on } (s - \{z\})$ (is *?P*)

by (*force intro: holomorphic_intros*)

also have $?P \longleftrightarrow g \text{ holomorphic_on } (s - \{z\})$

by (*intro holomorphic_cong refl*) (*simp_all add: g_def*)

finally have $*$: $g \text{ holomorphic_on } (s - \{z\})$.

note *lim*

also have $(\lambda w. f w * (w - z)) -z \rightarrow c \longleftrightarrow g -z \rightarrow g z$

by (*intro filterlim_cong refl*) (*simp_all add: g_def [abs_def] eventually_at_filter*)

finally have $**$: $g -z \rightarrow g z$.

have *g_holo: g holomorphic_on s*

by (*rule no_isolated_singularity'[where K = {z}]*)

(*insert assms * **, simp_all add: at_within_open_NO_MATCH*)

from *s* and *this* have *residue* $(\lambda w. g w / (w - z)) z = g z$

by (*rule residue_simple*)

also have $\forall_F za \text{ in } at z. g za / (za - z) = f za$

unfolding *eventually_at* by (auto intro!: *exI[of _ 1] simp: field_simps g_def*)

hence *residue* $(\lambda w. g w / (w - z)) z = \text{residue } f z$

by (*intro residue_cong refl*)

finally show *?thesis*

by (simp add: g_def)
qed

lemma *residue_holomorphic_over_power*:

assumes *open A z0 ∈ A f holomorphic_on A*
shows $\text{residue } (\lambda z. f z / (z - z0) ^ \text{Suc } n) z0 = (\text{deriv } ^ n) f z0 / \text{fact } n$
proof –
let $?f = \lambda z. f z / (z - z0) ^ \text{Suc } n$
from *assms(1,2)* **obtain** *r* **where** $r: r > 0 \text{ cball } z0 r \subseteq A$
by (auto simp: *open_contains_cball*)
have (*?f has_contour_integral* $2 * \pi * i * \text{residue } ?f z0$) (*circlepath* *z0 r*)
using *r assms* by (*intro base_residue[of A]*) (*auto intro!: holomorphic_intros*)
moreover have (*?f has_contour_integral* $2 * \pi * i / \text{fact } n * (\text{deriv } ^ n) f z0$) (*circlepath* *z0 r*)
using *assms r*
by (*intro Cauchy_has_contour_integral_higher_derivative_circlepath*)
(*auto intro!: holomorphic_on_subset[OF assms(3)] holomorphic_on_imp_continuous_on*)
ultimately have $2 * \pi * i * \text{residue } ?f z0 = 2 * \pi * i / \text{fact } n * (\text{deriv } ^ n) f z0$
by (*rule has_contour_integral_unique*)
thus *?thesis* by (*simp add: field_simps*)
qed

lemma *residue_holomorphic_over_power'*:

assumes *open A 0 ∈ A f holomorphic_on A*
shows $\text{residue } (\lambda z. f z / z ^ \text{Suc } n) 0 = (\text{deriv } ^ n) f 0 / \text{fact } n$
using *residue_holomorphic_over_power[OF assms]* by *simp*

theorem *residue_fps_expansion_over_power_at_0*:

assumes *f has_fps_expansion F*
shows $\text{residue } (\lambda z. f z / z ^ \text{Suc } n) 0 = \text{fps_nth } F n$
proof –
from *has_fps_expansion_imp_holomorphic[OF assms]* **obtain** *s*
where *open s 0 ∈ s f holomorphic_on s* $\bigwedge z. z \in s \implies f z = \text{eval_fps } F z$
by *auto*
with *assms* **have** $\text{residue } (\lambda z. f z / (z - 0) ^ \text{Suc } n) 0 = (\text{deriv } ^ n) f 0 / \text{fact } n$
unfolding *has_fps_expansion_def*
by (*intro residue_holomorphic_over_power[of s]*) (*auto simp: zero_ereal_def*)
also from *assms* **have** $\dots = \text{fps_nth } F n$
by (*subst fps_nth_fps_expansion*) *auto*
finally show *?thesis* by *simp*
qed

lemma *residue_pole_order*:

fixes *f::complex ⇒ complex* **and** *z::complex*
defines $n \equiv \text{nat } (- \text{zorder } f z)$ **and** $h \equiv \text{zor_poly } f z$
assumes *f_iso:isolated_singularity_at f z*
and *pole:is_pole f z*

```

shows residue f z = ((deriv  $\hat{\sim}$  (n - 1)) h z / fact (n-1))
proof -
  define g where g  $\equiv$   $\lambda x$ . if x=z then 0 else inverse (f x)
  obtain e where [simp]:e>0 and f_holo:f holomorphic_on ball z e - {z}
  using f_iso analytic_imp_holomorphic unfolding isolated_singularity_at_def
  by blast
  obtain r where 0 < n 0 < r and r_cball:cball z r  $\subseteq$  ball z e and h_holo: h
  holomorphic_on cball z r
  and h_divide:( $\forall w \in$  cball z r. ( $w \neq z \longrightarrow f w = h w / (w - z) ^ n \wedge h w \neq 0$ ))
  proof -
    obtain r where r:zorder f z < 0 h z  $\neq$  0 r>0 cball z r  $\subseteq$  ball z e h holomor-
    phic_on cball z r
    ( $\forall w \in$  cball z r - {z}.  $f w = h w / (w - z) ^ n \wedge h w \neq 0$ )
    using zorder_exist_pole[OF f_holo,simplified,OF  $\langle$ is_pole f z $\rangle$ ,folded n_def
    h_def] by auto
    have n>0 using  $\langle$ zorder f z < 0 $\rangle$  unfolding n_def by simp
    moreover have ( $\forall w \in$  cball z r. ( $w \neq z \longrightarrow f w = h w / (w - z) ^ n \wedge h w \neq$ 
    0))
      using  $\langle$ h z  $\neq$  0 $\rangle$  r(6) by blast
    ultimately show ?thesis using r(3,4,5) that by blast
  qed
  have r_nonzero: $\bigwedge w$ .  $w \in$  ball z r - {z}  $\implies f w \neq 0$ 
  using h_divide by simp
  define c where c  $\equiv$  2 * pi * i
  define der_f where der_f  $\equiv$  ((deriv  $\hat{\sim}$  (n - 1)) h z / fact (n-1))
  define h' where h'  $\equiv$   $\lambda u$ . h u / (u - z) ^ n
  have (h' has_contour_integral c / fact (n - 1) * (deriv  $\hat{\sim}$  (n - 1)) h z)
  (circlepath z r)
  unfolding h'_def
  proof (rule Cauchy_has_contour_integral_higher_derivative_circlepath[of z r
  h z n-1,
  folded c_def Suc_pred'[OF  $\langle$ n>0 $\rangle$ ]])
    show continuous_on (cball z r) h using holomorphic_on_imp_continuous_on
    h_holo by simp
    show h holomorphic_on ball z r using h_holo by auto
    show z  $\in$  ball z r using  $\langle$ r>0 $\rangle$  by auto
  qed
  then have (h' has_contour_integral c * der_f) (circlepath z r) unfolding
  der_f_def by auto
  then have (f has_contour_integral c * der_f) (circlepath z r)
  proof (elim has_contour_integral_eq)
    fix x assume x  $\in$  path_image (circlepath z r)
    hence x  $\in$  cball z r - {z} using  $\langle$ r>0 $\rangle$  by auto
    then show h' x = f x using h_divide unfolding h'_def by auto
  qed
  moreover have (f has_contour_integral c * residue f z) (circlepath z r)
  using base_residue[of  $\langle$ ball z e $\rangle$  z,simplified,OF  $\langle$ r>0 $\rangle$  f_holo r_cball,folded
  c_def]
  unfolding c_def by simp

```


ultimately have $c * \text{der}_f = c * \text{residue } f z$ using *has_contour_integral_unique*
 by *blast*
 hence $\text{der}_f = \text{residue } f z$ unfolding *c_def* by *auto*
 thus *?thesis* unfolding *der_f_def* by *auto*
 qed

lemma *residue_simple_pole*:
 assumes *isolated_singularity_at* $f z0$
 assumes *is_pole* $f z0$ $\text{zorder } f z0 = - 1$
 shows $\text{residue } f z0 = \text{zor_poly } f z0 z0$
 using *assms* by (*subst residue_pole_order*) *simp_all*

lemma *residue_simple_pole_limit*:
 assumes *isolated_singularity_at* $f z0$
 assumes *is_pole* $f z0$ $\text{zorder } f z0 = - 1$
 assumes $((\lambda x. f (g x) * (g x - z0)) \longrightarrow c) F$
 assumes *filterlim* g (at $z0$) $F F \neq \text{bot}$
 shows $\text{residue } f z0 = c$

proof –
 have $\text{residue } f z0 = \text{zor_poly } f z0 z0$
 by (*rule residue_simple_pole assms*)
 also have $\dots = c$
 apply (*rule zor_poly_pole_eqI*)
 using *assms* by *auto*
 finally show *?thesis* .
 qed

lemma
 assumes *f_holo*: f *holomorphic_on* s and *g_holo*: g *holomorphic_on* s
 and *open* s *connected* $s z \in s$
 assumes *g_deriv*:(g *has_field_derivative* g') (at z)
 assumes $f z \neq 0$ $g z = 0$ $g' \neq 0$
 shows *porder_simple_pole_deriv*: $\text{zorder } (\lambda w. f w / g w) z = - 1$
 and *residue_simple_pole_deriv*: $\text{residue } (\lambda w. f w / g w) z = f z / g'$
 proof –
 have [*simp*]:*isolated_singularity_at* $f z$ *isolated_singularity_at* $g z$
 using *isolated_singularity_at_holomorphic*[*OF* $\langle \text{open } s \rangle \langle z \in s \rangle$] *f_holo g_holo*
 by (*meson Diff_subset holomorphic_on_subset*)
 have [*simp*]:*not_essential* $f z$ *not_essential* $g z$
 unfolding *not_essential_def* using *f_holo g_holo assms*(3,5)
 by (*meson continuous_on_eq_continuous_at continuous_within holomorphic_on_imp_continuous_on*)
 have *g_nconst*: $\exists_F w$ in at $z. g w \neq 0$
 proof (*rule ccontr*)
 assume $\neg (\exists_F w$ in at $z. g w \neq 0)$
 then have $\forall_F w$ in *nhds* $z. g w = 0$
 unfolding *eventually_at eventually_nhds frequently_at* using $\langle g z = 0 \rangle$
 by (*metis open_ball UNIV_I centre_in_ball dist_commute mem_ball*)
 then have *deriv* $g z = \text{deriv } (\lambda_. 0) z$
 by (*intro deriv_cong_ev*) *auto*

```

    then have deriv g z = 0 by auto
    then have g' = 0 using g_deriv DERIV_imp_deriv by blast
    then show False using ⟨g'≠0⟩ by auto
  qed

  have zorder (λw. f w / g w) z = zorder f z - zorder g z
  proof -
    have ∀F w in at z. f w ≠ 0 ∧ w ∈ s
    apply (rule non_zero_neighbour_alt)
    using assms by auto
    with g_nconst have ∃F w in at z. f w * g w ≠ 0
    by (elim frequently_rev_mp eventually_rev_mp, auto)
    then show ?thesis using zorder_divide[of f z g] by auto
  qed

  moreover have zorder f z = 0
  apply (rule zorder_zero_eqI[OF f_holo ⟨open s⟩ ⟨z ∈ s⟩])
  using ⟨f z ≠ 0⟩ by auto
  moreover have zorder g z = 1
  apply (rule zorder_zero_eqI[OF g_holo ⟨open s⟩ ⟨z ∈ s⟩])
  subgoal using assms(8) by auto
  subgoal using DERIV_imp_deriv assms(9) g_deriv by auto
  subgoal by simp
  done
  ultimately show zorder (λw. f w / g w) z = - 1 by auto

  show residue (λw. f w / g w) z = f z / g'
  proof (rule residue_simple_pole_limit[where g=id and F=at z, simplified])
    show zorder (λw. f w / g w) z = - 1 by fact
    show isolated_singularity_at (λw. f w / g w) z
    by (auto intro: singularity_intros)
    show is_pole (λw. f w / g w) z
    proof (rule is_pole_divide)
      have ∀F x in at z. g x ≠ 0
      apply (rule non_zero_neighbour)
      using g_nconst by auto
      moreover have g -z → 0
      using DERIV_isCont assms(8) continuous_at g_deriv by force
      ultimately show filterlim g (at 0) (at z) unfolding filterlim_at by simp
      show isCont f z
      using assms(3,5) continuous_on_eq_continuous_at f_holo holomorphic_on_imp_continuous_on
      by auto
      show f z ≠ 0 by fact
    qed
    show filterlim id (at z) (at z) by (simp add: filterlim_iff)
    have ((λw. (f w * (w - z)) / g w) → f z / g') (at z)
    proof (rule lhopital_complex_simple)
      show ((λw. f w * (w - z)) has_field_derivative f z) (at z)
      using assms by (auto intro!: derivative_eq_intros holomorphic_derivI[OF

```

```

f_holo])
  show (g has_field_derivative g') (at z) by fact
qed (insert assms, auto)
then show (( $\lambda w. (f w / g w) * (w - z)$ )  $\longrightarrow$  f z / g') (at z)
  by (simp add: field_split_simps)
qed
qed

```

7.16 Poles and residues of some well-known functions

```

lemma is_pole_Gamma: is_pole Gamma (-of_nat n)
  unfolding is_pole_def using Gamma_poles .

```

```

lemma Gamma_residue:
  residue Gamma (-of_nat n) = (-1) ^ n / fact n
proof (rule residue_simple')
  show open (- (Z<0 - {-of_nat n}) :: complex set)
    by (intro open_Compl closed_subset_Ints) auto
  show Gamma_holomorphic_on (- (Z<0 - {-of_nat n}) - {- of_nat n})
    by (rule holomorphic_Gamma) auto
  show ( $\lambda w. Gamma w * (w - (-of\_nat\ n))$ )  $\rightarrow$  (-of_nat n)  $\rightarrow$  (-1) ^ n / fact n
    using Gamma_residues[of n] by simp
qed auto

```

end

8 The Residue Theorem, the Argument Principle and Rouché's Theorem

```

theory Residue_Theorem
  imports Complex_Residues HOL-Library.Landau_Symbols
begin

```

Several theorems that could be moved up, IF there were a previous theory importing both Landau Symbols and Elementary Metric Spaces

```

lemma continuous_bounded_at_infinity_imp_bounded:
  fixes f :: real  $\Rightarrow$  'a :: real_normed_field
  assumes f  $\in$  O[at_bot]( $\lambda_. 1$ )
  assumes f  $\in$  O[at_top]( $\lambda_. 1$ )
  assumes cf: continuous_on UNIV f
  shows bounded (range f)
proof -
  obtain c1 c2
    where eventually ( $\lambda x. norm (f x) \leq c1$ ) at_bot eventually ( $\lambda x. norm (f x) \leq c2$ ) at_top
  using assms by (auto elim!: landau_o.bigE)
  then obtain x1 x2 where x1:  $\bigwedge x. x \leq x1 \implies norm (f x) \leq c1$  and x2:  $\bigwedge x. x \geq x2 \implies norm (f x) \leq c2$ 

```

```

  by (auto simp: eventually_at_bot_linorder eventually_at_top_linorder)
  have compact (f ' {x1..x2})
  by (intro compact_continuous_image continuous_on_subset[OF cf]) auto
  hence bounded (f ' {x1..x2})
  by (rule compact_imp_bounded)
  then obtain c3 where c3:  $\bigwedge x. x \in \{x1..x2\} \implies \text{norm } (f x) \leq c3$ 
  unfolding bounded_iff by fast
  have norm (f x)  $\leq \text{Max } \{c1, c2, c3\}$  for x
  by (cases x  $\leq x1$ ; cases x  $\geq x2$ ) (use x1 x2 c3 in (auto simp: le_max_iff_disj))
  thus ?thesis
  unfolding bounded_iff by blast
qed

```

lemma holomorphic_on_extend:

```

  assumes f holomorphic_on S - {ξ} ξ ∈ interior S f ∈ O[at ξ](λ_. 1)
  shows (∃ g. g holomorphic_on S ∧ (∀ z ∈ S - {ξ}. g z = f z))
  by (subst holomorphic_on_extend_bounded) (insert assms, auto elim!: landau_o.bigE)

```

lemma removable_singularities:

```

  assumes finite X X ⊆ interior S f holomorphic_on (S - X)
  assumes  $\bigwedge z. z \in X \implies f \in O[at z](\lambda_. 1)$ 
  shows  $\exists g. g \text{ holomorphic\_on } S \wedge (\forall z \in S - X. g z = f z)$ 
  using assms

```

proof (induction arbitrary: f rule: finite_induct)

```

  case empty
  thus ?case by auto

```

next

```

  case (insert z0 X f)
  from insert.prem1 and insert.hyps have z0: z0 ∈ interior (S - X)
  by (auto simp: interior_diff finite_imp_closed)
  hence  $\exists g. g \text{ holomorphic\_on } (S - X) \wedge (\forall z \in S - X - \{z0\}. g z = f z)$ 
  using insert.prem1 insert.hyps by (intro holomorphic_on_extend) auto
  then obtain g where g: g holomorphic_on (S - X)  $\forall z \in S - X - \{z0\}. g z =$ 
  fz by blast

```

```

  have  $\exists h. h \text{ holomorphic\_on } S \wedge (\forall z \in S - X. h z = g z)$ 

```

proof (rule insert.IH)

```

  fix z0' assume z0': z0' ∈ X

```

```

  hence eventually (λz. z ∈ interior S - (X - {z0'}) - {z0}) (nhds z0')

```

```

  using insert.prem1 insert.hyps

```

```

  by (intro eventually_nhds_in_open open_Diff finite_imp_closed) auto

```

```

  hence ev: eventually (λz. z ∈ S - X - {z0'}) (at z0')

```

```

  unfolding eventually_at_filter

```

```

  by eventually_elim (insert z0' insert.hyps interior_subset[of S], auto)

```

```

  have g ∈ Θ[at z0'](f)

```

```

  by (intro bigthetaI_cong eventually_mono[OF ev]) (insert g, auto)

```

```

  also have f ∈ O[at z0'](λ_. 1)

```

```

  using z0' by (intro insert.prem1) auto

```

```

  finally show g ∈ ...

```

qed (use insert.prem1 g in auto)

then obtain h where h holomorphic_on $S \forall z \in S - X. h z = g z$ by blast
with g have h holomorphic_on $S \forall z \in S - \text{insert } z0 X. h z = f z$ by auto
thus ?case by blast
qed

lemma continuous_imp_bigo_1:
assumes continuous (at x within A) f
shows $f \in O[\text{at } x \text{ within } A](\lambda_. 1)$
proof (rule bigoI_tendsto)
from **assms show** $((\lambda x. f x / 1) \longrightarrow f x)$ (at x within A)
by (auto simp: continuous_within)
qed auto

lemma taylor_bigo_linear:
assumes f field_differentiable at $x0$ within A
shows $(\lambda x. f x - f x0) \in O[\text{at } x0 \text{ within } A](\lambda x. x - x0)$
proof –
from **assms obtain** f' **where** (f has_field_derivative f') (at $x0$ within A)
by (auto simp: field_differentiable_def)
hence $((\lambda x. (f x - f x0) / (x - x0)) \longrightarrow f')$ (at $x0$ within A)
by (auto simp: has_field_derivative_iff)
thus ?thesis by (intro bigoI_tendsto[**where** $c = f'$]) (auto simp: eventually_at_filter)
qed

8.1 Cauchy's residue theorem

lemma get_integrable_path:
assumes open S connected (S -pts) finite pts f holomorphic_on (S -pts) $a \in S$ -pts
 $b \in S$ -pts
obtains g **where** valid_path g pathstart $g = a$ pathfinish $g = b$
 $\text{path_image } g \subseteq S$ -pts f contour_integrable_on g **using** **assms**
proof (induct arbitrary: S thesis a rule: finite_induct[OF <finite pts>])
case 1
obtain g **where** valid_path g path_image $g \subseteq S$ pathstart $g = a$ pathfinish $g = b$
using connected_open_polynomial_connected[OF <open S >, of $a b$] <connected
 $(S - \{\})$ >
 $\text{valid_path_polynomial_function } 1.\text{prems}(6) 1.\text{prems}(7)$ **by** auto
moreover have f contour_integrable_on g
using contour_integrable_holomorphic_simple[OF __ <open S > <valid_path g >
<path_image $g \subseteq S$ >, of f]
< f holomorphic_on $S - \{\}$ >
by auto
ultimately show ?case **using** 1(1)[of g] **by** auto
next
case idt:($2 p$ pts)
obtain e **where** $e > 0$ **and** $e: \forall w \in \text{ball } a e. w \in S \wedge (w \neq a \longrightarrow w \notin \text{insert } p \text{ pts})$
using finite_ball_avoid[OF <open S > <finite (insert p pts)>, of a]
< $a \in S - \text{insert } p \text{ pts}$ >

```

    by auto
  define a' where a'  $\equiv$  a+e/2
  have a'  $\in$  S - {p} - pts using e[rule_format, of a+e/2]  $\langle e > 0 \rangle$ 
    by (auto simp add: dist_complex_def a'_def)
  then obtain g' where g'[simp]: valid_path g' pathstart g' = a' pathfinish g' = b
    path_image g'  $\subseteq$  S - {p} - pts f contour_integrable_on g'
    using idt.hyps(3)[of a' S - {p}] idt.prem1 idt.hyps(1)
    by (metis Diff_insert2 open_delete)
  define g where g  $\equiv$  linepath a a' +++ g'
  have valid_path g unfolding g_def by (auto intro: valid_path_join)
  moreover have pathstart g = a and pathfinish g = b unfolding g_def by auto
  moreover have path_image g  $\subseteq$  S - insert p pts
    unfolding g_def
  proof (rule subset_path_image_join)
    have closed_segment a a'  $\subseteq$  ball a e using  $\langle e > 0 \rangle$ 
    by (auto dest!: segment_bound1 simp: a'_def dist_complex_def norm_minus_commute)
    then show path_image (linepath a a')  $\subseteq$  S - insert p pts using e idt(9)
    by auto
  next
    show path_image g'  $\subseteq$  S - insert p pts using g'(4) by blast
  qed
  moreover have f contour_integrable_on g
  proof -
    have closed_segment a a'  $\subseteq$  ball a e using  $\langle e > 0 \rangle$ 
    by (auto dest!: segment_bound1 simp: a'_def dist_complex_def norm_minus_commute)
    then have closed_segment a a'  $\subseteq$  S - insert p pts
      using e idt.prem1(6) by auto
    then have continuous_on (closed_segment a a') f
      using holomorphic_on_imp_continuous_on holomorphic_on_subset idt.prem1(5)
  by presburger
    then show ?thesis
      using contour_integrable_continuous_linepath by (simp add: g_def)
  qed
  ultimately show ?case using idt.prem1(1)[of g] by auto
qed

lemma Cauchy_theorem_aux:
  assumes open S connected (S-pts) finite pts pts  $\subseteq$  S f holomorphic_on S-pts
    valid_path g pathfinish g = pathstart g path_image g  $\subseteq$  S-pts
     $\forall z. (z \notin S) \longrightarrow$  winding_number g z = 0
     $\forall p \in S. h p > 0 \wedge (\forall w \in cball p (h p). w \in S \wedge (w \neq p \longrightarrow w \notin pts))$ 
  shows contour_integral g f =  $(\sum p \in pts. winding\_number\ g\ p * contour\_integral$ 
    (circlepath p (h p)) f)
    using assms
  proof (induct arbitrary: S g rule: finite_induct[OF  $\langle$ finite pts $\rangle$ ])
    case 1
    then show ?case by (simp add: Cauchy_theorem_global contour_integral_unique)
  next
    case (2 p pts)

```

```

note  $fin[simp] = \langle finite (insert\ p\ pts) \rangle$ 
and  $connected = \langle connected\ (S - insert\ p\ pts) \rangle$ 
and  $valid[simp] = \langle valid\_path\ g \rangle$ 
and  $g\_loop[simp] = \langle pathfinish\ g = pathstart\ g \rangle$ 
and  $holo[simp] = \langle f\ holomorphic\_on\ S - insert\ p\ pts \rangle$ 
and  $path\_img = \langle path\_image\ g \subseteq S - insert\ p\ pts \rangle$ 
and  $winding = \langle \forall z. z \notin S \longrightarrow winding\_number\ g\ z = 0 \rangle$ 
and  $h = \langle \forall pa \in S. 0 < h\ pa \wedge (\forall w \in cball\ pa\ (h\ pa). w \in S \wedge (w \neq pa \longrightarrow w \notin insert\ p\ pts)) \rangle$ 
have  $h\ p > 0$  and  $p \in S$ 
and  $h\_p: \forall w \in cball\ p\ (h\ p). w \in S \wedge (w \neq p \longrightarrow w \notin insert\ p\ pts)$ 
using  $h\ \langle insert\ p\ pts \subseteq S \rangle$  by auto
obtain  $pg$  where  $pg[simp]: valid\_path\ pg\ pathstart\ pg = pathstart\ g\ pathfinish\ g$ 
 $pg = p + h\ p$ 
 $path\_image\ pg \subseteq S - insert\ p\ pts$  f  $contour\_integrable\_on\ pg$ 
proof -
have  $p + h\ p \in cball\ p\ (h\ p)$  using  $h[rule\_format, of\ p]$ 
by  $(simp\ add: \langle p \in S \rangle\ dist\_norm)$ 
then have  $p + h\ p \in S - insert\ p\ pts$  using  $h[rule\_format, of\ p]\ \langle insert\ p\ pts \subseteq S \rangle$ 
by fastforce
moreover have  $pathstart\ g \in S - insert\ p\ pts$  using  $path\_img$  by auto
ultimately show ?thesis
using  $get\_integrable\_path[OF\ \langle open\ S \rangle\ connected\ fin\ holo, of\ pathstart\ g\ p + h\ p]$  that
by blast
qed
obtain  $n::int$  where  $n = winding\_number\ g\ p$ 
using  $integer\_winding\_number[OF\ \_g\_loop, of\ p]\ valid\ path\_img$ 
by  $(metis\ DiffD2\ Ints\_cases\ insertI1\ subset\_eq\ valid\_path\_imp\_path)$ 
define  $p\_circ$  where  $p\_circ \equiv circlepath\ p\ (h\ p)$ 
define  $p\_circ\_pt$  where  $p\_circ\_pt \equiv linepath\ (p + h\ p)\ (p + h\ p)$ 
define  $n\_circ$  where  $n\_circ \equiv \lambda n. ((+++)\ p\_circ\ \overset{\sim}{n}\ p\_circ\_pt)$ 
define  $cp$  where  $cp \equiv if\ n \geq 0\ then\ reversepath\ (n\_circ\ (nat\ n))\ else\ n\_circ\ (nat\ (-\ n))$ 

have  $n\_circ: valid\_path\ (n\_circ\ k)$ 
 $winding\_number\ (n\_circ\ k)\ p = k$ 
 $pathstart\ (n\_circ\ k) = p + h\ p$ 
 $pathfinish\ (n\_circ\ k) = p + h\ p$ 
 $path\_image\ (n\_circ\ k) = (if\ k = 0\ then\ \{p + h\ p\}\ else\ sphere\ p\ (h\ p))$ 
 $p \notin path\_image\ (n\_circ\ k)$ 
 $\wedge p'. p' \notin S - pts \implies winding\_number\ (n\_circ\ k)\ p' = 0 \wedge p' \notin path\_image\ (n\_circ\ k)$ 
 $f\ contour\_integrable\_on\ (n\_circ\ k)$ 
 $contour\_integral\ (n\_circ\ k)\ f = k * contour\_integral\ p\_circ\ f$ 
for  $k$ 
proof (induct k)
case 0
show  $valid\_path\ (n\_circ\ 0)$ 

```

```

and path_image (n_circ 0) = (if 0=0 then {p + h p} else sphere p (h p))
and winding_number (n_circ 0) p = of_nat 0
and pathstart (n_circ 0) = p + h p
and pathfinish (n_circ 0) = p + h p
and p ∉ path_image (n_circ 0)
unfolding n_circ_def p_circ_pt_def using ⟨h p > 0⟩
by (auto simp add: dist_norm)
show winding_number (n_circ 0) p'=0 ∧ p'∉path_image (n_circ 0) when
p'∉S- pts for p'
  unfolding n_circ_def p_circ_pt_def
  apply (auto intro!:winding_number_trivial)
  by (metis Diff_iff pathfinish_in_path_image pg(3) pg(4) subsetCE subset_insertI that)+
  show f contour_integrable_on (n_circ 0)
  unfolding n_circ_def p_circ_pt_def
  by (auto intro!:contour_integrable_continuous_linepath simp add:continuous_on_sing)
  show contour_integral (n_circ 0) f = of_nat 0 * contour_integral p_circ f
  unfolding n_circ_def p_circ_pt_def by auto
next
  case (Suc k)
  have n_Suc:n_circ (Suc k) = p_circ +++ n_circ k unfolding n_circ_def
by auto
  have pcirc:p ∉ path_image p_circ valid_path p_circ pathfinish p_circ =
pathstart (n_circ k)
  using Suc(3) unfolding p_circ_def using ⟨h p > 0⟩ by (auto simp add:
p_circ_def)
  have pcirc_image:path_image p_circ ⊆ S - insert p pts
  proof -
  have path_image p_circ ⊆ cball p (h p) using ⟨0 < h p⟩ p_circ_def by
auto
  then show ?thesis using h_p pcirc(1) by auto
  qed
  have pcirc_integrable:f contour_integrable_on p_circ
  by (auto simp add:p_circ_def intro!: pcirc_image[unfolded p_circ_def]
contour_integrable_continuous_circlepath holomorphic_on_imp_continuous_on
holomorphic_on_subset[OF holo])
  show valid_path (n_circ (Suc k))
  using valid_path_join[OF pcirc(2) Suc(1) pcirc(3)] unfolding n_circ_def
by auto
  show path_image (n_circ (Suc k))
  = (if Suc k = 0 then {p + complex_of_real (h p)} else sphere p (h p))
  proof -
  have path_image p_circ = sphere p (h p)
  unfolding p_circ_def using ⟨0 < h p⟩ by auto
  then show ?thesis unfolding n_Suc using Suc.hyps(5) ⟨h p>0⟩
  by (auto simp add: path_image_join[OF pcirc(3)] dist_norm)
  qed
  then show p ∉ path_image (n_circ (Suc k)) using ⟨h p>0⟩ by auto
  show winding_number (n_circ (Suc k)) p = of_nat (Suc k)

```



```

proof –
  have winding_number p_circ p = 1
  by (simp add: ⟨h p > 0⟩ p_circ_def winding_number_circlepath_centre)
  moreover have p ∉ path_image (n_circ k) using Suc(5) ⟨h p > 0⟩ by
auto
  then have winding_number (p_circ +++ n_circ k) p
    = winding_number p_circ p + winding_number (n_circ k) p
  using valid_path_imp_path Suc.hyps(1) Suc.hyps(2) pcirc
  apply (intro winding_number_join)
  by auto
  ultimately show ?thesis using Suc(2) unfolding n_circ_def
  by auto
qed
show pathstart (n_circ (Suc k)) = p + h p
  by (simp add: n_circ_def p_circ_def)
show pathfinish (n_circ (Suc k)) = p + h p
  using Suc(4) unfolding n_circ_def by auto
show winding_number (n_circ (Suc k)) p'=0 ∧ p'∉path_image (n_circ
(Suc k)) when p'∉S-pts for p'
  proof –
    have p' ∉ path_image p_circ using ⟨p ∈ S⟩ h p_circ_def that using
pcirc_image by blast
    moreover have p' ∉ path_image (n_circ k)
    using Suc.hyps(7) that by blast
    moreover have winding_number p_circ p' = 0
    proof –
      have path_image p_circ ⊆ cball p (h p)
      using h unfolding p_circ_def using ⟨p ∈ S⟩ by fastforce
      moreover have p'∉cball p (h p) using ⟨p ∈ S⟩ h that 2.hyps(2) by
fastforce
    ultimately show ?thesis
    unfolding p_circ_def
    by (intro winding_number_zero_outside) auto
    qed
    ultimately show ?thesis
    unfolding n_Suc using Suc.hyps pcirc
    by (metis add.right_neutral not_in_path_image_join that valid_path_imp_path
winding_number_join)
    qed
    show f contour_integrable_on (n_circ (Suc k))
    unfolding n_Suc
    by (rule contour_integrable_joinI[OF pcirc_integrable Suc(8) pcirc(2)
Suc(1)])
    show contour_integral (n_circ (Suc k)) f = (Suc k) * contour_integral
p_circ f
    by (simp add: Rings.ring_distrib(2) Suc.hyps n_Suc pcirc pcirc_integrable)
    qed
have cp[simp]:pathstart cp = p + h p pathfinish cp = p + h p
  valid_path cp path_image cp ⊆ S – insert p pts

```

```

winding_number cp p = - n
 $\wedge p'. p' \notin S - pts \implies \text{winding\_number } cp \ p' = 0 \wedge p' \notin \text{path\_image } cp$ 
f contour_integrable_on cp
contour_integral cp f = - n * contour_integral p_circ f
proof -
  show pathstart cp = p + h p and pathfinish cp = p + h p and valid_path cp
  using n_circ unfolding cp_def by auto
next
  have sphere p (h p)  $\subseteq S - \text{insert } p \ pts$ 
  using h[rule_format,of p]  $\langle \text{insert } p \ pts \subseteq S \rangle$  by force
  moreover have p + complex_of_real (h p)  $\in S - \text{insert } p \ pts$ 
  using pg(3) pg(4) by (metis pathfinish_in_path_image subsetCE)
  ultimately show path_image cp  $\subseteq S - \text{insert } p \ pts$  unfolding cp_def
  using n_circ(5) by auto
next
  show winding_number cp p = - n
  unfolding cp_def using winding_number_reversepath n_circ  $\langle h \ p > 0 \rangle$ 
  by (auto simp: valid_path_imp_path)
next
  show winding_number cp  $p' = 0 \wedge p' \notin \text{path\_image } cp$  when  $p' \notin S - pts$  for
  p'
  proof -
    have winding_number (reversepath (n_circ (nat n)))  $p' = 0$ 
    using n_circ that
    by (metis add.inverse_neutral valid_path_imp_path winding_number_reversepath)
    then show ?thesis
    using cp_def n_circ(7) that by force
  qed
next
  show f contour_integrable_on cp unfolding cp_def
  using contour_integrable_reversepath_eq n_circ(1,8) by auto
next
  show contour_integral cp f = - n * contour_integral p_circ f
  unfolding cp_def using contour_integral_reversepath[OF n_circ(1)]
  n_circ(9)
  by auto
  qed
  define g' where  $g' \equiv g +++ pg +++ cp +++ (\text{reversepath } pg)$ 
  have contour_integral g' f =  $(\sum p \in pts. \text{winding\_number } g' \ p * \text{contour\_integral } (\text{circlepath } p \ (h \ p)) \ f)$ 
  proof (rule 2.hyps(3)[of  $S - \{p\}$  g', OF __  $\langle \text{finite } pts \rangle$  ])
    show connected (S - {p} - pts) using connected by (metis Diff_insert2)
    show open (S - {p}) using  $\langle \text{open } S \rangle$  by auto
    show pts  $\subseteq S - \{p\}$  using  $\langle \text{insert } p \ pts \subseteq S \rangle \langle p \notin pts \rangle$  by blast
    show f holomorphic_on S - {p} - pts using holo  $\langle p \notin pts \rangle$  by (metis
  Diff_insert2)
  show valid_path g'
  unfolding g'_def cp_def using n_circ valid_pg g_loop
  by (auto intro!: valid_path_join)

```

```

show pathfinish g' = pathstart g'
  unfolding g'_def cp_def using pg(2) by simp
show path_image g'  $\subseteq$  S - {p} - pts
  proof -
    define s' where s'  $\equiv$  S - {p} - pts
    have s':s' = S - insert p pts unfolding s'_def by auto
    then show ?thesis using path_img pg(4) cp(4)
      by (simp add: g'_def s'_def subset_path_image_join)
  qed
note path_join_imp[simp]
show  $\forall z. z \notin S - \{p\} \longrightarrow$  winding_number g' z = 0
  proof clarify
    fix z assume z:z  $\notin$  S - {p}
    have z_notin_cp: z  $\notin$  path_image cp
      using cp(6) cp_def n_circ(6) z by auto
    have z_notin_pg: z  $\notin$  path_image pg
      by (metis Diff_iff Diff_insert2 pg(4) subsetD z)
    have winding_number (g +++ pg +++ cp +++ reversepath pg) z =
winding_number g z
      + winding_number (pg +++ cp +++ (reversepath pg)) z
    proof (rule winding_number_join)
      show path g using <valid_path g> by (simp add: valid_path_imp_path)
      show z  $\notin$  path_image g using z path_img by auto
      show path (pg +++ cp +++ reversepath pg) using pg(3) cp
        by (simp add: valid_path_imp_path)
    next
      have path_image (pg +++ cp +++ reversepath pg)  $\subseteq$  S - insert p pts
        using pg(4) cp(4) by (auto simp: subset_path_image_join)
      then show z  $\notin$  path_image (pg +++ cp +++ reversepath pg) using
z by auto
    next
      show pathfinish g = pathstart (pg +++ cp +++ reversepath pg) using
g_loop by auto
    qed
    also have ... = winding_number g z + (winding_number pg z
      + winding_number (cp +++ (reversepath pg)) z)
    proof (subst add_left_cancel, rule winding_number_join)
      show path pg and path (cp +++ reversepath pg)
        and pathfinish pg = pathstart (cp +++ reversepath pg)
        by (auto simp add: valid_path_imp_path)
      show z  $\notin$  path_image pg using pg(4) z by blast
      show z  $\notin$  path_image (cp +++ reversepath pg) using z
        by (metis Diff_iff <z  $\notin$  path_image pg> contra_subsetD cp(4) insertI1
          not_in_path_image_join path_image_reversepath singletonD)
    qed
    also have ... = winding_number g z + (winding_number pg z
      + (winding_number cp z + winding_number (reversepath pg) z))
    by (simp add: valid_path_imp_path winding_number_join z_notin_cp
z_notin_pg)

```

```

    also have ... = winding_number g z + winding_number cp z
      by (simp add: valid_path_imp_path winding_number_reversepath
z_notin_pg)
    finally have winding_number g' z = winding_number g z + winding_number cp z
      unfolding g'_def .
    moreover have winding_number g z + winding_number cp z = 0
      using winding z ⟨n=winding_number g p⟩ by auto
    ultimately show winding_number g' z = 0 unfolding g'_def by auto
  qed
  show ∀ pa ∈ S - {p}. 0 < h pa ∧ (∀ w ∈ cball pa (h pa). w ∈ S - {p} ∧ (w
≠ pa → w ∉ pts))
    using h by fastforce
  qed
  moreover have contour_integral g' f = contour_integral g f
    - winding_number g p * contour_integral p_circ f
  proof -
    have *: f contour_integrable_on g f contour_integrable_on pg f contour_integrable_on
cp
      by (auto simp add: open_Diff[OF ⟨open S⟩, OF finite_imp_closed[OF fin]]
intro!: contour_integrable_holomorphic_simple[OF holo__path_img])
    have contour_integral g' f = contour_integral g f + contour_integral pg f
      + contour_integral cp f + contour_integral (reversepath pg) f
      using * by (simp add: g'_def contour_integrable_reversepath_eq)
    also have ... = contour_integral g f + contour_integral cp f
      using contour_integral_reversepath
      by (auto simp add: contour_integrable_reversepath)
    also have ... = contour_integral g f - winding_number g p * contour_integral
p_circ f
      using ⟨n=winding_number g p⟩ by auto
    finally show ?thesis .
  qed
  moreover have winding_number g' p' = winding_number g p' when p' ∈ pts for
p'
  proof -
    obtain [simp]: p' ∉ path_image g p' ∉ path_image pg p' ∉ path_image cp
      using 2.prem8 that by (metis Diff_iff Diff_insert2 ⟨p' ∈ pts⟩ cp(4)
pg(4) subsetD)
    have winding_number g' p' = winding_number g p' + winding_number pg p'
      + winding_number (cp +++ reversepath pg) p'
      by (simp add: g'_def not_in_path_image_join valid_path_imp_path winding_number_join)
    also have ... = winding_number g p' using that
      by (simp add: valid_path_imp_path winding_number_join winding_number_reversepath)
    finally show ?thesis .
  qed
  ultimately show ?case unfolding p_circ_def
  apply (subst (asm) sum.cong[OF refl,
of pts _ λp. winding_number g p * contour_integral (circlepath p (h p)) f])

```

by (auto simp: sum.insert[OF ‹finite pts› ‹p∉pts›] algebra_simps)
qed

lemma *Cauchy_theorem_singularities:*

assumes *open S connected S finite pts and*

holo: f holomorphic_on S-pts and

valid_path g and

loop:pathfinish g = pathstart g and

path_image g ⊆ S-pts and

homo:∀ z. (z ∉ S) ⟶ winding_number g z = 0 and

avoid:∀ p∈S. h p>0 ∧ (∀ w∈cball p (h p). w∈S ∧ (w≠p ⟶ w ∉ pts))

shows *contour_integral g f = (∑ p∈pts. winding_number g p * contour_integral (circlepath p (h p)) f)*

(is ?L=?R)

proof –

define *circ* **where** *circ ≡ λp. winding_number g p * contour_integral (circlepath p (h p)) f*

define *pts1* **where** *pts1 ≡ pts ∩ S*

define *pts2* **where** *pts2 ≡ pts - pts1*

have *pts=pts1 ∪ pts2 pts1 ∩ pts2 = {} pts2 ∩ S={}* *pts1⊆S*

unfolding *pts1_def pts2_def* **by** *auto*

have *contour_integral g f = (∑ p∈pts1. circ p)* **unfolding** *circ_def*

proof (*rule Cauchy_theorem_aux[OF ‹open S› _ _ ‹pts1⊆S› _ ‹valid_path g› loop _ homo]*)

have *finite pts1* **unfolding** *pts1_def* **using** *‹finite pts›* **by** *auto*

then show *connected (S - pts1)*

using *‹open S› ‹connected S› connected_open_delete_finite[of S]* **by** *auto*

next

show *finite pts1* **using** *‹pts = pts1 ∪ pts2› assms(3)* **by** *auto*

show *f holomorphic_on S - pts1* **by** (*metis Diff_Int2 Int_absorb holo pts1_def*)

show *path_image g ⊆ S - pts1* **using** *assms(7) pts1_def* **by** *auto*

show *∀ p∈S. 0 < h p ∧ (∀ w∈cball p (h p). w ∈ S ∧ (w ≠ p ⟶ w ∉ pts1))*

by (*simp add: avoid pts1_def*)

qed

moreover have *sum circ pts2 = 0*

by (*metis ‹pts2 ∩ S = {}› circ_def disjoint_iff_not_equal homo mult_zero_left sum.neutral*)

moreover have *?R=sum circ pts1 + sum circ pts2*

unfolding *circ_def*

using *sum.union_disjoint[OF _ _ ‹pts1 ∩ pts2 = {}›] ‹finite pts› ‹pts=pts1 ∪ pts2›*

by *blast*

ultimately show *?thesis*

by *simp*

qed

theorem *Residue_theorem:*

fixes *S pts::complex set and f::complex ⇒ complex*

```

  and  $g::\text{real} \Rightarrow \text{complex}$ 
  assumes open  $S$  connected  $S$  finite pts and
    holo: $f$  holomorphic_on  $S$ -pts and
    valid_path  $g$  and
    loop:pathfinish  $g = \text{pathstart } g$  and
    path_image  $g \subseteq S$ -pts and
    homo: $\forall z. (z \notin S) \longrightarrow \text{winding\_number } g z = 0$ 
  shows contour_integral  $g f = 2 * \pi * i * (\sum_{p \in \text{pts.}} \text{winding\_number } g p * \text{residue } f p)$ 
  proof -
    define  $c$  where  $c \equiv 2 * \pi * i$ 
    obtain  $h$  where avoid: $\forall p \in S. h p > 0 \wedge (\forall w \in \text{cball } p (h p). w \in S \wedge (w \neq p \longrightarrow w \notin \text{pts}))$ 
    using finite_cball_avoid[OF <open  $S$ > <finite pts>] by metis
    have contour_integral  $g f$ 
      =  $(\sum_{p \in \text{pts.}} \text{winding\_number } g p * \text{contour\_integral } (\text{circlepath } p (h p)) f)$ 
    using Cauchy_theorem_singularities[OF assms avoid] .
    also have ... =  $(\sum_{p \in \text{pts.}} c * \text{winding\_number } g p * \text{residue } f p)$ 
    proof (intro sum.cong)
      show pts = pts by simp
    next
      fix  $x$  assume  $x \in \text{pts}$ 
      show winding_number  $g x * \text{contour\_integral } (\text{circlepath } x (h x)) f$ 
        =  $c * \text{winding\_number } g x * \text{residue } f x$ 
      proof (cases  $x \in S$ )
        case False
          then have winding_number  $g x = 0$  using homo by auto
          thus ?thesis by auto
        next
          case True
            have contour_integral  $(\text{circlepath } x (h x)) f = c * \text{residue } f x$ 
              using < $x \in \text{pts}$ > <finite pts> avoid[rule_format, OF True]
            apply (intro base_residue[of  $S - (\text{pts} - \{x\})$ , THEN contour_integral_unique, folded
              c_def])
              by (auto intro:holomorphic_on_subset[OF holo] open_Diff[OF <open  $S$ >
                finite_imp_closed])
            then show ?thesis by auto
          qed
      qed
    also have ... =  $c * (\sum_{p \in \text{pts.}} \text{winding\_number } g p * \text{residue } f p)$ 
      by (simp add: sum_distrib_left algebra_simps)
    finally show ?thesis unfolding c_def .
  qed

```

8.2 The argument principle

theorem *argument_principle*:

fixes $f::\text{complex} \Rightarrow \text{complex}$ and poles $S::\text{complex set}$

defines $pz \equiv \{w \in S. f w = 0 \vee w \in \text{poles}\}$ — pz is the set of poles and zeros

```

assumes open S connected S and
  f_holo:f holomorphic_on S - poles and
  h_holo:h holomorphic_on S and
  valid_path g and
  loop:pathfinish g = pathstart g and
  path_img:path_image g  $\subseteq$  S - pz and
  homo: $\forall z. (z \notin S) \longrightarrow$  winding_number g z = 0 and
  finite:finite pz and
  poles: $\forall p \in S \cap$  poles. is_pole f p
shows contour_integral g ( $\lambda x. \text{deriv } f x * h x / f x$ ) = 2 * pi * i *
  ( $\sum p \in$  pz. winding_number g p * h p * zorder f p)
(is ?L=?R)
proof -
  define c where c  $\equiv$  2 * complex_of_real pi * i
  define ff where ff  $\equiv$  ( $\lambda x. \text{deriv } f x * h x / f x$ )
  define cont where cont  $\equiv$   $\lambda$ ff p e. (ff has_contour_integral c * zorder f p * h p
) (circlepath p e)
  define avoid where avoid  $\equiv$   $\lambda$ p e.  $\forall w \in$  cball p e.  $w \in S \wedge (w \neq p \longrightarrow w \notin$  pz)

  have  $\exists e > 0. \text{avoid } p e \wedge (p \in$  pz  $\longrightarrow \text{cont ff } p e)$  when  $p \in S$  for p
  proof -
    obtain e1 where e1 > 0 and e1_avoid:avoid p e1
    using finite_cball_avoid[OF <open S> finite] <p ∈ S> unfolding avoid_def by
  auto
  have  $\exists e2 > 0. \text{cball } p e2 \subseteq \text{ball } p e1 \wedge \text{cont ff } p e2$  when  $p \in$  pz
  proof -
    define po where po  $\equiv$  zorder f p
    define pp where pp  $\equiv$  zor_poly f p
    define f' where f'  $\equiv$   $\lambda w. pp w * (w - p) \text{powi } po$ 
    define ff' where ff'  $\equiv$  ( $\lambda x. \text{deriv } f' x * h x / f' x$ )
    obtain r where pp p  $\neq 0$  r > 0 and
      r < e1 and
      pp_holo:pp holomorphic_on cball p r and
      pp_po:( $\forall w \in$  cball p r - {p}. f w = pp w * (w - p) powi po  $\wedge$  pp w  $\neq 0$ )
    proof -
      have isolated_singularity_at f p
      proof -
        have ball p e1 - {p}  $\subseteq$  S - poles
        using avoid_def e1_avoid pz_def by fastforce
        then have f holomorphic_on ball p e1 - {p}
        by (intro holomorphic_on_subset[OF f_holo])
        then show ?thesis unfolding isolated_singularity_at_def
        using <e1 > 0> analytic_on_open open_delete by blast
      qed
    moreover have not_essential f p
    proof (cases is_pole f p)
      case True
      then show ?thesis unfolding not_essential_def by auto
    next

```

```

case False
then have  $p \in S$ -poles using  $\langle p \in S \rangle$  poles unfolding pz_def by auto
moreover have open ( $S$ -poles)
proof -
  have closed ( $S \cap$  poles)
    using finite by (simp add: pz_def finite_imp_closed rev_finite_subset
subset_eq)
  then show ?thesis
    by (metis Diff_Cmpl Diff_Diff_Int Diff_eq  $\langle$ open  $S \rangle$  open_Diff)
  qed
ultimately have isCont  $f$   $p$ 
  using holomorphic_on_imp_continuous_on[OF f_holo] continuous_on_eq_continuous_at
  by auto
  then show ?thesis unfolding isCont_def not_essential_def by auto
qed
moreover have  $\exists_F w$  in at  $p$ .  $f w \neq 0$ 
proof (rule ccontr)
  assume  $\neg (\exists_F w$  in at  $p$ .  $f w \neq 0$ )
  then have  $\forall_F w$  in at  $p$ .  $f w = 0$  unfolding frequently_def by auto
  then obtain  $r1$  where  $r1 > 0$  and  $r1: \forall w \in$ ball  $p$   $r1 - \{p\}$ .  $f w = 0$ 
    unfolding eventually_at by (auto simp add: dist_commute)
  obtain  $r2$  where  $r2 > 0$  and  $r2: \text{ball } p \ r2 \subseteq S$ 
    using  $\langle p \in S \rangle$   $\langle$ open  $S \rangle$  openE by blast
  define  $rr$  where  $rr = \min \ r1 \ r2$ 

from  $r1 \ r2$ 
have ball  $p$   $rr - \{p\} \subseteq \{w \in S \cap \text{ball } p \ rr - \{p\}. f w = 0\}$ 
  unfolding rr_def by auto
moreover have infinite (ball  $p$   $rr - \{p\}$ )
  using  $\langle r1 > 0 \rangle$   $\langle r2 > 0 \rangle$  finite_imp_not_open
  unfolding rr_def by fastforce
  ultimately have infinite  $\{w \in S \cap \text{ball } p \ rr - \{p\}. f w = 0\}$  using infinite_super by blast
  then have infinite pz
  unfolding pz_def by (smt (verit) infinite_super Collect_mono_iff DiffE
Int_iff)
  then show False using  $\langle$ finite pz $\rangle$  by auto
qed
ultimately obtain  $r$  where  $pp \ p \neq 0$  and  $r: r > 0$   $pp$  holomorphic_on cball
 $p \ r$ 
  ( $\forall w \in$ ball  $p \ r - \{p\}$ .  $f w = pp \ w * (w - p) \text{ pow } i \ po \wedge pp \ w \neq 0$ )
  using zorder_exist[of f p, folded po_def pp_def] by auto
define  $r1$  where  $r1 = \min \ r \ e1 / 2$ 
have  $r1 < e1$  unfolding r1_def using  $\langle e1 > 0 \rangle$   $\langle r > 0 \rangle$  by auto
moreover have  $r1 > 0$   $pp$  holomorphic_on cball  $p \ r1$ 
  ( $\forall w \in$ ball  $p \ r1 - \{p\}$ .  $f w = pp \ w * (w - p) \text{ pow } i \ po \wedge pp \ w \neq 0$ )
  unfolding r1_def using  $\langle e1 > 0 \rangle$   $r$  by auto
ultimately show ?thesis using that  $\langle pp \ p \neq 0 \rangle$  by auto

```



```

qed

define e2 where e2  $\equiv$  r/2
have e2>0 using <r>0> unfolding e2_def by auto
define anal where anal  $\equiv$   $\lambda w. \text{deriv } pp \ w * h \ w / pp \ w$ 
define prin where prin  $\equiv$   $\lambda w. po * h \ w / (w - p)$ 
have (( $\lambda w. \text{prin } w + anal \ w$ ) has_contour_integral c * po * h p) (circlepath
p e2)
proof (rule has_contour_integral_add[of _ _ _ _ 0,simplified])
  have ball_p_r  $\subseteq$  S
  using <r<e1> avoid_def ball_subset_cball e1_avoid by (simp add:
subset_eq)
  then have cball_p_e2  $\subseteq$  S
  using <r>0> unfolding e2_def by auto
  then have ( $\lambda w. po * h \ w$ ) holomorphic_on cball_p_e2
  using h_holo by (auto intro!: holomorphic_intros)
  then show (prin has_contour_integral c * po * h p) (circlepath p e2)
  using Cauchy_integral_circlepath_simple[folded c_def, of  $\lambda w. po * h \ w$ ]
  <e2>0>
  unfolding prin_def by (auto simp add: mult.assoc)
  have anal holomorphic_on ball_p_r unfolding anal_def
  using pp_holo h_holo pp_po <ball_p_r  $\subseteq$  S> <pp  $\neq$  0>
  by (auto intro!: holomorphic_intros)
  then show (anal has_contour_integral 0) (circlepath p e2)
  using e2_def <r>0>
  by (auto elim!: Cauchy_theorem_disc_simple)
qed
then have cont_ff' p e2 unfolding cont_def po_def
proof (elim has_contour_integral_eq)
  fix w assume w  $\in$  path_image (circlepath p e2)
  then have w $\in$ ball_p_r and w $\neq$ p unfolding e2_def using <r>0> by auto
  define wp where wp  $\equiv$  w-p
  have wp $\neq$ 0 and pp w  $\neq$  0
  unfolding wp_def using <w $\neq$ p> <w $\in$ ball_p_r> pp_po by auto
  moreover have der_f':deriv f' w = po * pp w * (w-p) powi (po - 1) +
deriv pp w * (w-p) powi po
  proof (rule DERIV_imp_deriv)
    have (pp has_field_derivative (deriv pp w)) (at w)
    using DERIV_deriv_iff_has_field_derivative pp_holo <w $\neq$ p>
    by (meson open_ball <w  $\in$  ball_p_r> ball_subset_cball holomorphic_derivI
holomorphic_on_subset)
    then show (f' has_field_derivative of_int po * pp w * (w - p) powi (po
- 1)
+ deriv pp w * (w - p) powi po) (at w)
    unfolding f'_def using <w $\neq$ p>
  by (auto intro!: derivative_eq_intros DERIV_cong[OF has_field_derivative_powr_of_int])
qed
ultimately show prin w + anal w = ff' w
  unfolding f'_def ff'_def prin_def anal_def

```

```

    apply (simp add: field_simps flip: wp_def)
    by (metis (no_types, lifting) mult.commute power_int_minus_mult)
qed
then have cont ff p e2 unfolding cont_def
proof (elim has_contour_integral_eq)
  fix w assume w ∈ path_image (circlepath p e2)
  then have w ∈ ball p r and w ≠ p unfolding e2_def using ‹r > 0› by auto
  have deriv f' w = deriv f w
  proof (rule complex_derivative_transform_within_open[where s = ball p r
- {p}])
    show f' holomorphic_on ball p r - {p} unfolding f'_def using pp_holo
      by (auto intro!: holomorphic_intros)
  next
    have ball p e1 - {p} ⊆ S - poles
      using ball_subset_cball e1_avoid[unfolded avoid_def] unfolding pz_def
      by auto
    then have ball p r - {p} ⊆ S - poles
      using ‹r < e1› by force
    then show f holomorphic_on ball p r - {p} using f_holo
      by auto
  next
    show open (ball p r - {p}) by auto
    show w ∈ ball p r - {p} using ‹w ∈ ball p r› ‹w ≠ p› by auto
  next
    fix x assume x ∈ ball p r - {p}
    then show f' x = f x
      using pp_po unfolding f'_def by auto
  qed
  moreover have f' w = f w
    using ‹w ∈ ball p r› ball_subset_cball subset_iff pp_po ‹w ≠ p›
    unfolding f'_def by auto
  ultimately show ff' w = ff w
    unfolding ff'_def ff_def by simp
qed
moreover have cball p e2 ⊆ ball p e1
  using ‹0 < r› ‹r < e1› e2_def by auto
ultimately show ?thesis using ‹e2 > 0› by auto
qed
then obtain e2 where e2: p ∈ pz → e2 > 0 ∧ cball p e2 ⊆ ball p e1 ∧ cont ff
p e2
  by auto
define e4 where e4 ≡ if p ∈ pz then e2 else e1
have e4 > 0 using e2 ‹e1 > 0› unfolding e4_def by auto
moreover have avoid p e4 using e2 ‹e1 > 0› e1_avoid unfolding e4_def
avoid_def by auto
moreover have p ∈ pz → cont ff p e4
  by (auto simp add: e2 e4_def)
ultimately show ?thesis by auto
qed

```

```

then obtain get_e where get_e:  $\forall p \in S. \text{get\_e } p > 0 \wedge \text{avoid } p (\text{get\_e } p)$ 
   $\wedge (p \in pz \longrightarrow \text{cont\_ff } p (\text{get\_e } p))$ 
by metis
define ci where ci  $\equiv \lambda p. \text{contour\_integral } (\text{circlepath } p (\text{get\_e } p)) \text{ ff}$ 
define w where w  $\equiv \lambda p. \text{winding\_number } g p$ 
have contour_integral g ff =  $(\sum p \in pz. w p * ci p)$  unfolding ci_def w_def
proof (rule Cauchy_theorem_singularities[OF ‹open S› ‹connected S› finite _
  ‹valid_path g› loop
    path_img homo])
  have open (S - pz) using open_Diff[OF _ finite_imp_closed[OF finite]] ‹open
  S› by auto
  then show ff holomorphic_on S - pz unfolding ff_def using f_holo h_holo
    by (auto intro!: holomorphic_intros simp add: pz_def)
  next
    show  $\forall p \in S. 0 < \text{get\_e } p \wedge (\forall w \in \text{cball } p (\text{get\_e } p). w \in S \wedge (w \neq p \longrightarrow w \notin pz))$ 
      using get_e using avoid_def by blast
  qed
also have ... =  $(\sum p \in pz. c * w p * h p * \text{zorder } f p)$ 
proof (rule sum.cong[of pz pz, simplified])
  fix p assume p  $\in$  pz
  show  $w p * ci p = c * w p * h p * (\text{zorder } f p)$ 
proof (cases p  $\in$  S)
  assume p  $\in$  S
  have  $ci p = c * h p * (\text{zorder } f p)$ 
    unfolding ci_def
  using ‹p  $\in$  S› ‹p  $\in$  pz› cont_def contour_integral_unique get_e by fastforce
  thus ?thesis by auto
next
  assume p  $\notin$  S
  then have w p = 0 using homo unfolding w_def by auto
  then show ?thesis by auto
qed
qed
also have ... =  $c * (\sum p \in pz. w p * h p * \text{zorder } f p)$ 
  unfolding sum_distrib_left by (simp add: algebra_simps)
finally have contour_integral g ff =  $c * (\sum p \in pz. w p * h p * \text{of\_int } (\text{zorder } f p))$ .
  then show ?thesis unfolding ff_def c_def w_def by simp
qed

```

8.3 Coefficient asymptotics for generating functions

For a formal power series that has a meromorphic continuation on some disc in the context plane, we can use the Residue Theorem to extract precise asymptotic information from the residues at the poles. This can be used to derive the asymptotic behaviour of the coefficients ($a_n \sim \dots$). If the function is meromorphic on the entire complex plane, one can even derive a

full asymptotic expansion.

We will first show the relationship between the coefficients and the sum over the residues with an explicit remainder term (the contour integral along the circle used in the Residue theorem).

theorem

fixes $f :: \text{complex} \Rightarrow \text{complex}$ **and** $n :: \text{nat}$ **and** $r :: \text{real}$
defines $g \equiv (\lambda w. f w / w \wedge \text{Suc } n)$ **and** $\gamma \equiv \text{circlepath } 0 r$
assumes $\text{open } A$ $\text{connected } A$ $\text{cball } 0 r \subseteq A$ $r > 0$
assumes $f \text{ holomorphic_on } A - S$ $S \subseteq \text{ball } 0 r$ $\text{finite } S$ $0 \notin S$
shows $\text{fps_coeff_conv_residues}$:
 $(\text{deriv } \wedge n) f 0 / \text{fact } n =$
 $\text{contour_integral } \gamma g / (2 * \text{pi} * i) - (\sum z \in S. \text{residue } g z)$ **(is ?thesis1)**
and $\text{fps_coeff_residues_bound}$:
 $(\bigwedge z. \text{norm } z = r \implies z \notin k \implies \text{norm } (f z) \leq C) \implies C \geq 0 \implies \text{finite}$
 $k \implies$
 $\text{norm } ((\text{deriv } \wedge n) f 0 / \text{fact } n + (\sum z \in S. \text{residue } g z)) \leq C / r \wedge n$

proof –

have $\text{holo}: g \text{ holomorphic_on } A - \text{insert } 0 S$
unfolding g_def **using** assms **by** $(\text{auto intro!}: \text{holomorphic_intros})$
have $\text{contour_integral } \gamma g = 2 * \text{pi} * i * (\sum z \in \text{insert } 0 S. \text{winding_number } \gamma z * \text{residue } g z)$
proof $(\text{rule Residue_theorem})$
show $g \text{ holomorphic_on } A - \text{insert } 0 S$ **by** fact
from assms **show** $\forall z. z \notin A \longrightarrow \text{winding_number } \gamma z = 0$
unfolding γ_def **by** $(\text{intro allI impI winding_number_zero_outside}[of _ \text{cball } 0 r]) \text{ auto}$
qed $(\text{insert } \text{assms}, \text{auto simp}: \gamma_def)$
also have $\text{winding_number } \gamma z = 1$ **if** $z \in \text{insert } 0 S$ **for** z
unfolding γ_def **using** assms **that** **by** $(\text{intro winding_number_circlepath}) \text{ auto}$
hence $(\sum z \in \text{insert } 0 S. \text{winding_number } \gamma z * \text{residue } g z) = (\sum z \in \text{insert } 0 S. \text{residue } g z)$
by $(\text{intro sum.cong}) \text{ simp_all}$
also have $\dots = \text{residue } g 0 + (\sum z \in S. \text{residue } g z)$
using $\langle 0 \notin S \rangle$ **and** $\langle \text{finite } S \rangle$ **by** $(\text{subst sum.insert}) \text{ auto}$
also from $\langle r > 0 \rangle$ **have** $0 \in \text{cball } 0 r$ **by** simp
with assms **have** $0 \in A - S$ **by** blast
with assms **have** $\text{residue } g 0 = (\text{deriv } \wedge n) f 0 / \text{fact } n$
unfolding g_def **by** $(\text{subst residue_holomorphic_over_power}[of } A - S])$
 $(\text{auto simp}: \text{finite_imp_closed})$
finally show $?thesis1$
by $(\text{simp add}: \text{field_simps})$

assume $C: \bigwedge z. \text{norm } z = r \implies z \notin k \implies \text{norm } (f z) \leq C$ $C \geq 0$ **and** $k: \text{finite } k$

have $(\text{deriv } \wedge n) f 0 / \text{fact } n + (\sum z \in S. \text{residue } g z) = \text{contour_integral } \gamma g / (2 * \text{pi} * i)$
using $\langle ?thesis1 \rangle$ **by** $(\text{simp add}: \text{algebra_simps})$
also have $\text{norm } \dots = \text{norm } (\text{contour_integral } \gamma g) / (2 * \text{pi})$

```

  by (simp add: norm_divide norm_mult)
  also have norm (contour_integral  $\gamma$  g)  $\leq C / r^{\text{Suc } n} * (2 * \pi * r)$ 
  proof (rule has_contour_integral_bound_circlepath_strong)
    from ‹open A› and ‹finite S› have open (A - insert 0 S)
      by (blast intro: finite_imp_closed)
    with assms show (g has_contour_integral contour_integral  $\gamma$  g) (circlepath 0
  r)
    unfolding  $\gamma\_def$ 
    by (intro has_contour_integral_integral contour_integrable_holomorphic_simple
  [OF holo]) auto
  next
    fix z assume z: norm (z - 0) = r z  $\notin k$ 
    hence norm (g z) = norm (f z) / rSuc n
      by (simp add: norm_divide g_def norm_mult norm_power)
    also have ...  $\leq C / r^{\text{Suc } n}$ 
      using k and ‹r > 0› and z by (intro divide_right_mono C zero_le_power)
  auto
    finally show norm (g z)  $\leq C / r^{\text{Suc } n}$  .
  qed (insert C(2) k ‹r > 0›, auto)
  also from ‹r > 0› have C / rSuc n * (2 * pi * r) / (2 * pi) = C / rn
    by simp
  finally show norm ((deriv  $\hat{\sim}$  n) f 0 / fact n + ( $\sum z \in S$ . residue g z))  $\leq \dots$ 
    by - (simp_all add: divide_right_mono)
  qed

```

Since the circle is fixed, we can get an upper bound on the values of the generating function on the circle and therefore show that the integral is $O(r^{-n})$.

corollary *fps_coeff_residues_bigo*:

```

  fixes f :: complex  $\Rightarrow$  complex and r :: real
  assumes open A connected A cball 0 r  $\subseteq$  A r > 0
  assumes f holomorphic_on A - S S  $\subseteq$  ball 0 r finite S 0  $\notin$  S
  assumes g: eventually ( $\lambda n$ . g n = -( $\sum z \in S$ . residue ( $\lambda z$ . f z / zSuc n) z))
  sequentially
    (is eventually ( $\lambda n$ .  $\_ = -?g' n$ )  $\_$ )
  shows ( $\lambda n$ . (deriv  $\hat{\sim}$  n) f 0 / fact n - g n)  $\in O(\lambda n$ . 1 / rn) (is ( $\lambda n$ . ?c n
  -  $\_$ )  $\in O(\_)$ )
  proof -
    from assms have compact (fˆ sphere 0 r)
      by (intro compact_continuous_image holomorphic_on_imp_continuous_on
    holomorphic_on_subset[OF ‹f holomorphic_on A - S›]) auto
    hence bounded (fˆ sphere 0 r) by (rule compact_imp_bounded)
    then obtain C where C:  $\bigwedge z$ . z  $\in$  sphere 0 r  $\implies$  norm (f z)  $\leq C$ 
      by (auto simp: bounded_iff sphere_def)
    have 0  $\leq$  norm (f (of_real r)) by simp
    also from C[of of_real r] and ‹r > 0› have ...  $\leq C$  by simp
    finally have C_nonneg: C  $\geq 0$  .

    have ( $\lambda n$ . ?c n + ?g' n)  $\in O(\lambda n$ . of_real (1 / rn))

```

```

proof (intro bigoI[of _ C] always_eventually allI )
  fix n :: nat
  from assms and C and C_nonneg have norm (?c n + ?g' n) ≤ C / r ^ n
  by (intro fps_coeff_residues_bound[where A = A and k = {}]) auto
  also have ... = C * norm (complex_of_real (1 / r ^ n))
  using ⟨r > 0⟩ by (simp add: norm_divide norm_power)
  finally show norm (?c n + ?g' n) ≤ ... .
qed
also have ?this ↔ (λn. ?c n - g n) ∈ O(λn. of_real (1 / r ^ n))
  by (intro landau_o.big.in_cong eventually_mono[OF g]) simp_all
finally show ?thesis .
qed

```

corollary *fps_coeff_residues_bigo'*:

```

fixes f :: complex ⇒ complex and r :: real
assumes exp: f has_fps_expansion F
assumes open A connected A cball 0 r ⊆ A r > 0
assumes f_holomorphic_on A - S S ⊆ ball 0 r finite S 0 ∉ S
assumes eventually (λn. g n = -(∑ z∈S. residue (λz. f z / z ^ Suc n) z))
  sequentially
  (is eventually (λn. _ = -?g' n) _)
shows (λn. fps_nth F n - g n) ∈ O(λn. 1 / r ^ n) (is (λn. ?c n - _) ∈
O(_))
proof -
  have fps_nth F = (λn. (deriv ~ n) f 0 / fact n)
  using fps_nth_fps_expansion[OF exp] by (intro ext) simp_all
  with fps_coeff_residues_bigo[OF assms(2-)] show ?thesis by simp
qed

```

8.4 Rouché's theorem

theorem *Rouche_theorem*:

```

fixes f g :: complex ⇒ complex and s :: complex set
defines fg ≡ (λp. f p + g p)
defines zeros_fg ≡ {p ∈ s. fg p = 0} and zeros_f ≡ {p ∈ s. f p = 0}
assumes
  open s and connected s and
  finite zeros_fg and
  finite zeros_f and
  f_holo: f holomorphic_on s and
  g_holo: g holomorphic_on s and
  valid_path γ and
  loop: pathfinish γ = pathstart γ and
  path_img: path_image γ ⊆ s and
  path_less: ∀ z ∈ path_image γ. cmod(f z) > cmod(g z) and
  homo: ∀ z. (z ∉ s) → winding_number γ z = 0
shows (∑ p ∈ zeros_fg. winding_number γ p * zorder fg p)
  = (∑ p ∈ zeros_f. winding_number γ p * zorder f p)
proof -

```

```

have path_fg:path_image  $\gamma \subseteq s - \text{zeros\_fg}$ 
proof -
  have False when  $z \in \text{path\_image } \gamma$  and  $f z + g z = 0$  for  $z$ 
  proof -
    have  $\text{cmod } (f z) > \text{cmod } (g z)$  using  $\langle z \in \text{path\_image } \gamma \rangle$  path_less by auto
    moreover have  $f z = -g z$  using  $\langle f z + g z = 0 \rangle$  by (simp add:
eq_neg_iff_add_eq_0)
    then have  $\text{cmod } (f z) = \text{cmod } (g z)$  by auto
    ultimately show False by auto
  qed
  then show ?thesis unfolding zeros_fg_def fg_def using path_img by auto
qed
have path_f:path_image  $\gamma \subseteq s - \text{zeros\_f}$ 
proof -
  have False when  $z \in \text{path\_image } \gamma$  and  $f z = 0$  for  $z$ 
  proof -
    have  $\text{cmod } (g z) < \text{cmod } (f z)$  using  $\langle z \in \text{path\_image } \gamma \rangle$  path_less by auto
    then have  $\text{cmod } (g z) < 0$  using  $\langle f z = 0 \rangle$  by auto
    then show False by auto
  qed
  then show ?thesis unfolding zeros_f_def using path_img by auto
qed
define w where  $w \equiv \lambda p. \text{winding\_number } \gamma p$ 
define c where  $c \equiv 2 * \text{complex\_of\_real } \pi * i$ 
define h where  $h \equiv \lambda p. g p / f p + 1$ 
obtain spikes
  where finite_spikes_and_spikes:  $\forall x \in \{0..1\} - \text{spikes}. \gamma$  differentiable at  $x$ 
  using  $\langle \text{valid\_path } \gamma \rangle$ 
by (auto simp: valid_path_def piecewise_C1_differentiable_on_def C1_differentiable_on_eq)
have h_contour:  $(\lambda x. \text{deriv } h x / h x)$  has_contour_integral 0
proof -
  have outside_img:  $0 \in \text{outside } (\text{path\_image } (h \circ \gamma))$ 
  proof -
    have  $h p \in \text{ball } 1 1$  when  $p \in \text{path\_image } \gamma$  for  $p$ 
    proof -
      have  $\text{cmod } (g p / f p) < 1$ 
      by (smt (verit) divide_less_eq_1_pos norm_divide norm_ge_zero
path_less that)
    then show ?thesis
      unfolding h_def by (auto simp add: dist_complex_def)
    qed
  then have path_image  $(h \circ \gamma) \subseteq \text{ball } 1 1$ 
  by (simp add: image_subset_iff path_image_compose)
  moreover have  $(0::\text{complex}) \notin \text{ball } 1 1$  by (simp add: dist_norm)
  ultimately show ?thesis
    using convex_in_outside[of ball 1 1 0] outside_mono by blast
  qed
  have valid_h: valid_path  $(h \circ \gamma)$ 
  proof (rule valid_path_compose_holomorphic[OF  $\langle \text{valid\_path } \gamma \rangle$  __ path_f])

```

```

    show  $h$  holomorphic_on  $s - \text{zeros}_f$ 
      unfolding  $h\_def$  using  $f\_holo$   $g\_holo$ 
      by (auto intro!: holomorphic_intros simp add:zeros_f_def)
  next
    show  $open$  ( $s - \text{zeros}_f$ ) using ⟨finite  $\text{zeros}_f$ ⟩ ⟨open  $s$ ⟩ finite_imp_closed
      by auto
  qed
  have (( $\lambda z. 1/z$ ) has_contour_integral 0) ( $h \circ \gamma$ )
  proof -
    have  $0 \notin \text{path\_image}$  ( $h \circ \gamma$ ) using outside_img by (simp add: outside_def)
    then have (( $\lambda z. 1/z$ ) has_contour_integral  $c * \text{winding\_number}$  ( $h \circ \gamma$ ) 0)
      ( $h \circ \gamma$ )
      using has_contour_integral_winding_number[of  $h \circ \gamma$  0,simplified] valid_h
      unfolding  $c\_def$  by auto
    moreover have  $\text{winding\_number}$  ( $h \circ \gamma$ ) 0 = 0
  proof -
    have  $0 \in \text{outside}$  ( $\text{path\_image}$  ( $h \circ \gamma$ )) using outside_img .
    moreover have  $\text{path}$  ( $h \circ \gamma$ )
      using valid_h by (simp add: valid_path_imp_path)
    moreover have  $\text{pathfinish}$  ( $h \circ \gamma$ ) =  $\text{pathstart}$  ( $h \circ \gamma$ )
      by (simp add: loop_pathfinish_compose pathstart_compose)
    ultimately show ?thesis using winding_number_zero_in_outside by auto
  qed
  ultimately show ?thesis by auto
  qed
  moreover have  $\text{vector\_derivative}$  ( $h \circ \gamma$ ) (at  $x$ ) =  $\text{vector\_derivative}$   $\gamma$  (at  $x$ )
  *  $\text{deriv}$   $h$  ( $\gamma$   $x$ )
    when  $x \in \{0..1\} - \text{spikes}$  for  $x$ 
  proof (rule vector_derivative_chain_at_general)
    show  $\gamma$  differentiable at  $x$  using that ⟨valid_path  $\gamma$ ⟩ spikes by auto
  next
    define  $der$  where  $der \equiv \lambda p. (\text{deriv } g \ p * f \ p - g \ p * \text{deriv } f \ p) / (f \ p * f \ p)$ 
    define  $t$  where  $t \equiv \gamma \ x$ 
    have  $f \ t \neq 0$  unfolding zeros_f_def t_def
      by (metis DiffD1 image_eqI norm_not_less_zero norm_zero path_defs(4)
      path_less that)
    moreover have  $t \in s$ 
      using contra_subsetD path_image_def path_fg t_def that by fastforce
    ultimately have ( $h$  has_field_derivative  $der \ t$ ) (at  $t$ )
      unfolding  $h\_def$   $der\_def$  using  $g\_holo$   $f\_holo$  ⟨open  $s$ ⟩
      by (auto intro!: holomorphic_derivI derivative_eq_intros)
    then show  $h$  field_differentiable at ( $\gamma \ x$ )
      unfolding  $t\_def$  field_differentiable_def by blast
  qed
  then have ((/) 1 has_contour_integral 0) ( $h \circ \gamma$ )
    = (( $\lambda x. \text{deriv } h \ x / h \ x$ ) has_contour_integral 0)  $\gamma$ 
    unfolding has_contour_integral
  by (force intro!: has_integral_spike_eq[OF negligible_finite, OF ⟨finite spikes⟩])
  ultimately show ?thesis by auto

```



```

qed
then have contour_integral  $\gamma$  ( $\lambda x. \text{deriv } h \ x / h \ x$ ) = 0
  using contour_integral_unique by simp
moreover have contour_integral  $\gamma$  ( $\lambda x. \text{deriv } fg \ x / fg \ x$ ) = contour_integral  $\gamma$ 
( $\lambda x. \text{deriv } f \ x / f \ x$ )
  + contour_integral  $\gamma$  ( $\lambda p. \text{deriv } h \ p / h \ p$ )
proof -
  have ( $\lambda p. \text{deriv } f \ p / f \ p$ ) contour_integrable_on  $\gamma$ 
  proof (rule contour_integrable_holomorphic_simple[OF _ _ <valid_path  $\gamma$ >
path_f])
    show open (s - zeros_f)
      using finite_imp_closed[OF <finite zeros_f>] <open s> by auto
    then show ( $\lambda p. \text{deriv } f \ p / f \ p$ ) holomorphic_on s - zeros_f
      using f_holo
    by (auto intro!: holomorphic_intros simp add:zeros_f_def)
  qed
moreover have ( $\lambda p. \text{deriv } h \ p / h \ p$ ) contour_integrable_on  $\gamma$ 
  using h_contour
  by (simp add: has_contour_integral_integrable)
ultimately have contour_integral  $\gamma$  ( $\lambda x. \text{deriv } f \ x / f \ x + \text{deriv } h \ x / h \ x$ ) =
  contour_integral  $\gamma$  ( $\lambda p. \text{deriv } f \ p / f \ p$ ) + contour_integral  $\gamma$ 
( $\lambda p. \text{deriv } h \ p / h \ p$ )
  using contour_integral_add[of ( $\lambda p. \text{deriv } f \ p / f \ p$ )  $\gamma$  ( $\lambda p. \text{deriv } h \ p / h \ p$ )]
  by auto
moreover have  $\text{deriv } fg \ p / fg \ p = \text{deriv } f \ p / f \ p + \text{deriv } h \ p / h \ p$ 
  when  $p \in \text{path\_image } \gamma$  for  $p$ 
proof -
  have  $fg \ p \neq 0$  and  $f \ p \neq 0$ 
    using path_f path_fg that unfolding zeros_f_def zeros_fg_def by auto
  have  $h \ p \neq 0$ 
  proof (rule ccontr)
    assume  $\neg h \ p \neq 0$ 
    then have  $cmod \ (g \ p / f \ p) = 1$ 
      by (simp add: add_eq_0_iff2 h_def)
    then show False
      by (smt (verit) divide_eq_1_iff norm_divide path_less that)
  qed
  have der_fg:  $\text{deriv } fg \ p = \text{deriv } f \ p + \text{deriv } g \ p$  unfolding fg_def
    using f_holo g_holo holomorphic_on_imp_differentiable_at[OF _ <open
s>] path_img that
    by auto
  have der_h:  $\text{deriv } h \ p = (\text{deriv } g \ p * f \ p - g \ p * \text{deriv } f \ p) / (f \ p * f \ p)$ 
  proof -
    define der where  $der \equiv \lambda p. (\text{deriv } g \ p * f \ p - g \ p * \text{deriv } f \ p) / (f \ p * f \ p)$ 
    have  $p \in s$  using path_img that by auto
    then have (h has_field_derivative der p) (at p)
      unfolding h_def der_def using g_holo f_holo <open s> <f p  $\neq 0$ >
      by (auto intro!: derivative_eq_intros holomorphic_derivI)
    then show ?thesis unfolding der_def using DERIV_imp_deriv by auto
  
```

```

qed
show ?thesis
  using ‹h p≠0› ‹f p≠0› ‹fg p≠0›
  unfolding der_fg der_h
  apply (simp add: divide_simps h_def fg_def)
  by (simp add: mult.commute mult.left_commute ring_class.ring_distrib(1))
qed
then have contour_integral  $\gamma$  ( $\lambda p. \text{deriv } fg \ p / fg \ p$ )
      = contour_integral  $\gamma$  ( $\lambda p. \text{deriv } f \ p / f \ p + \text{deriv } h \ p / h \ p$ )
  by (elim contour_integral_eq)
ultimately show ?thesis by auto
qed
moreover have contour_integral  $\gamma$  ( $\lambda x. \text{deriv } fg \ x / fg \ x$ ) =  $c * (\sum_{p \in \text{zeros\_fg}} w \ p * \text{zorder } fg \ p)$ 
proof -
  have fg_holomorphic_on s
    unfolding fg_def using f_holo g_holo holomorphic_on_add by auto
  moreover
  have path_image  $\gamma \subseteq s - \{p \in s. fg \ p = 0\}$ 
    using path_fg unfolding zeros_fg_def .
  moreover
  have finite { $p \in s. fg \ p = 0$ }
    using ‹finite zeros_fg› unfolding zeros_fg_def .
  ultimately show ?thesis
    unfolding c_def zeros_fg_def w_def
    using argument_principle[OF ‹open s› ‹connected s› __ ‹valid_path  $\gamma$ › loop
    _ homo, of_ {}  $\lambda_. 1$ ]
    by simp
qed
moreover have contour_integral  $\gamma$  ( $\lambda x. \text{deriv } f \ x / f \ x$ ) =  $c * (\sum_{p \in \text{zeros\_f}} w \ p * \text{zorder } f \ p)$ 
  unfolding c_def zeros_f_def w_def
proof (rule argument_principle[OF ‹open s› ‹connected s› __ ‹valid_path  $\gamma$ ›
loop _ homo
, of_ {}  $\lambda_. 1, \text{simplified}$ ])
  show f_holomorphic_on s
    using f_holo g_holo holomorphic_on_add by auto
  show path_image  $\gamma \subseteq s - \{p \in s. f \ p = 0\}$ 
    using path_f unfolding zeros_f_def .
  show finite { $p \in s. f \ p = 0$ }
    using ‹finite zeros_f› unfolding zeros_f_def .
qed
ultimately have  $c * (\sum_{p \in \text{zeros\_fg}} w \ p * (\text{zorder } fg \ p)) = c * (\sum_{p \in \text{zeros\_f}} w \ p * (\text{zorder } f \ p))$ 
  by auto
then show ?thesis unfolding c_def using w_def by auto
qed
end

```

theory *Laurent_Convergence*

imports *HOL-Computational_Algebra.Formal_Laurent_Series HOL-Library.Landau_Symbols Residue_Theorem*

begin

definition *fls_conv_radius* :: *complex fls* \Rightarrow *ereal* **where**
fls_conv_radius *f* = *fps_conv_radius* (*fls_regp*art *f*)

definition *eval_fls* :: *complex fls* \Rightarrow *complex* \Rightarrow *complex* **where**
eval_fls *F* *z* = *eval_fps* (*fls_base_factor_to_fps* *F*) *z* * *z* *powi fls_subdegree* *F*

definition
has_laurent_expansion :: (*complex* \Rightarrow *complex*) \Rightarrow *complex fls* \Rightarrow *bool*
(infixl \langle *has'_laurent'_expansion* \rangle 60)
where (*f* *has_laurent_expansion* *F*) \longleftrightarrow
fls_conv_radius *F* > 0 \wedge *eventually* ($\lambda z. \text{eval_fls } F z = f z$) (at 0)

lemma *has_laurent_expansion_schematicI*:
f *has_laurent_expansion* *F* \Longrightarrow *F* = *G* \Longrightarrow *f* *has_laurent_expansion* *G*
by *simp*

lemma *has_laurent_expansion_cong*:
assumes *eventually* ($\lambda x. f x = g x$) (at 0) *F* = *G*
shows (*f* *has_laurent_expansion* *F*) \longleftrightarrow (*g* *has_laurent_expansion* *G*)
proof –
have *eventually* ($\lambda z. \text{eval_fls } F z = g z$) (at 0)
if *eventually* ($\lambda z. \text{eval_fls } F z = f z$) (at 0) *eventually* ($\lambda x. f x = g x$) (at 0)
for *f g*
using *that* **by** *eventually_elim auto*
from *this*[of *f g*] *this*[of *g f*] **show** *?thesis*
using *assms* **by** (*auto simp: eq_commute has_laurent_expansion_def*)
qed

lemma *has_laurent_expansion_cong'*:
assumes *eventually* ($\lambda x. f x = g x$) (at *z*) *F* = *G* *z* = *z'*
shows ($(\lambda x. f (z + x))$ *has_laurent_expansion* *F*) \longleftrightarrow ($(\lambda x. g (z' + x))$ *has_laurent_expansion* *G*)
by (*intro has_laurent_expansion_cong*)
(use assms in \langle *auto simp: at_to_0' eventually_filtermap add_ac* \rangle)

lemma *fls_conv_radius_altdef*:
fls_conv_radius *F* = *fps_conv_radius* (*fls_base_factor_to_fps* *F*)
proof –
have *conv_radius* ($\lambda n. \text{fls_nth } F (\text{int } n)$) = *conv_radius* ($\lambda n. \text{fls_nth } F (\text{int } n + \text{fls_subdegree } F)$)
proof (*cases fls_subdegree* *F* \geq 0)
case *True*

```

hence conv_radius ( $\lambda n. \text{fls\_nth } F \text{ (int } n + \text{fls\_subdegree } F)$ ) =
  conv_radius ( $\lambda n. \text{fls\_nth } F \text{ (int (} n + \text{nat (fls\_subdegree } F))$ ))
by auto
thus ?thesis
by (subst (asm) conv_radius_shift) auto
next
case False
hence conv_radius ( $\lambda n. \text{fls\_nth } F \text{ (int } n)$ ) =
  conv_radius ( $\lambda n. \text{fls\_nth } F \text{ (fls\_subdegree } F + \text{int (} n + \text{nat (-fls\_subdegree } F))$ ))
by auto
thus ?thesis
by (subst (asm) conv_radius_shift) (auto simp: add_ac)
qed
thus ?thesis
by (simp add: fls_conv_radius_def fps_conv_radius_def)
qed

```

```

lemma eval_fps_of_nat [simp]: eval_fps (of_nat n) z = of_nat n
and eval_fps_of_int [simp]: eval_fps (of_int m) z = of_int m
by (simp_all flip: fps_of_nat fps_of_int)

```

```

lemma fps_conv_radius_of_nat [simp]: fps_conv_radius (of_nat n) =  $\infty$ 
and fps_conv_radius_of_int [simp]: fps_conv_radius (of_int m) =  $\infty$ 
by (simp_all flip: fps_of_nat fps_of_int)

```

```

lemma fps_conv_radius_fls_regpart: fps_conv_radius (fls_regpart F) = fls_conv_radius
F
by (simp add: fls_conv_radius_def)

```

```

lemma fls_conv_radius_0 [simp]: fls_conv_radius 0 =  $\infty$ 
and fls_conv_radius_1 [simp]: fls_conv_radius 1 =  $\infty$ 
and fls_conv_radius_const [simp]: fls_conv_radius (fls_const c) =  $\infty$ 
and fls_conv_radius_numeral [simp]: fls_conv_radius (numeral num) =  $\infty$ 
and fls_conv_radius_of_nat [simp]: fls_conv_radius (of_nat n) =  $\infty$ 
and fls_conv_radius_of_int [simp]: fls_conv_radius (of_int m) =  $\infty$ 
and fls_conv_radius_X [simp]: fls_conv_radius fls_X =  $\infty$ 
and fls_conv_radius_X_inv [simp]: fls_conv_radius fls_X_inv =  $\infty$ 
and fls_conv_radius_X_intpow [simp]: fls_conv_radius (fls_X_intpow m) =
 $\infty$ 
by (simp_all add: fls_conv_radius_def fls_X_intpow_regpart)

```

```

lemma fls_conv_radius_shift [simp]: fls_conv_radius (fls_shift n F) = fls_conv_radius
F
unfolding fls_conv_radius_altdef by (subst fls_base_factor_to_fps_shift) (rule
refl)

```

```

lemma fls_conv_radius_fps_to_fls [simp]: fls_conv_radius (fps_to_fls F) = fps_conv_radius
F

```

by (simp add: fls_conv_radius_def)

lemma fls_conv_radius_deriv [simp]: fls_conv_radius (fls_deriv F) ≥ fls_conv_radius F

proof –

have fls_conv_radius (fls_deriv F) = fps_conv_radius (fls_repart (fls_deriv F))

by (simp add: fls_conv_radius_def)

also have fls_repart (fls_deriv F) = fps_deriv (fls_repart F)

by (intro fps_ext) (auto simp: add_ac)

also have fps_conv_radius ... ≥ fls_conv_radius F

by (simp add: fls_conv_radius_def fps_conv_radius_deriv)

finally show ?thesis .

qed

lemma fls_conv_radius_uminus [simp]: fls_conv_radius (-F) = fls_conv_radius F

by (simp add: fls_conv_radius_def)

lemma fls_conv_radius_add: fls_conv_radius (F + G) ≥ min (fls_conv_radius F) (fls_conv_radius G)

by (simp add: fls_conv_radius_def fps_conv_radius_add)

lemma fls_conv_radius_diff: fls_conv_radius (F - G) ≥ min (fls_conv_radius F) (fls_conv_radius G)

by (simp add: fls_conv_radius_def fps_conv_radius_diff)

lemma fls_conv_radius_mult: fls_conv_radius (F * G) ≥ min (fls_conv_radius F) (fls_conv_radius G)

proof (cases F = 0 ∨ G = 0)

case False

hence [simp]: F ≠ 0 G ≠ 0

by auto

have fls_conv_radius (F * G) = fps_conv_radius (fls_repart (fls_shift (fls_subdegree F + fls_subdegree G) (F * G)))

by (simp add: fls_conv_radius_altdef)

also have fls_repart (fls_shift (fls_subdegree F + fls_subdegree G) (F * G))

=

fls_base_factor_to_fps F * fls_base_factor_to_fps G

by (simp add: fls_times_def)

also have fps_conv_radius ... ≥ min (fls_conv_radius F) (fls_conv_radius G)

unfolding fls_conv_radius_altdef by (rule fps_conv_radius_mult)

finally show ?thesis .

qed auto

lemma fps_conv_radius_add_ge:

fps_conv_radius F ≥ r ⇒ fps_conv_radius G ≥ r ⇒ fps_conv_radius (F + G) ≥ r

using fps_conv_radius_add[of F G] by (simp add: min_def split: if_splits)

lemma *fps_conv_radius_diff_ge*:

$\text{fps_conv_radius } F \geq r \implies \text{fps_conv_radius } G \geq r \implies \text{fps_conv_radius } (F - G) \geq r$
using *fps_conv_radius_diff*[of *F G*] **by** (*simp add: min_def split: if_splits*)

lemma *fps_conv_radius_mult_ge*:

$\text{fps_conv_radius } F \geq r \implies \text{fps_conv_radius } G \geq r \implies \text{fps_conv_radius } (F * G) \geq r$
using *fps_conv_radius_mult*[of *F G*] **by** (*simp add: min_def split: if_splits*)

lemma *fls_conv_radius_add_ge*:

$\text{fls_conv_radius } F \geq r \implies \text{fls_conv_radius } G \geq r \implies \text{fls_conv_radius } (F + G) \geq r$
using *fls_conv_radius_add*[of *F G*] **by** (*simp add: min_def split: if_splits*)

lemma *fls_conv_radius_diff_ge*:

$\text{fls_conv_radius } F \geq r \implies \text{fls_conv_radius } G \geq r \implies \text{fls_conv_radius } (F - G) \geq r$
using *fls_conv_radius_diff*[of *F G*] **by** (*simp add: min_def split: if_splits*)

lemma *fls_conv_radius_mult_ge*:

$\text{fls_conv_radius } F \geq r \implies \text{fls_conv_radius } G \geq r \implies \text{fls_conv_radius } (F * G) \geq r$
using *fls_conv_radius_mult*[of *F G*] **by** (*simp add: min_def split: if_splits*)

lemma *fls_conv_radius_power*: $\text{fls_conv_radius } (F \wedge n) \geq \text{fls_conv_radius } F$
by (*induction n*) (*auto intro!: fls_conv_radius_mult_ge*)

lemma *eval_fls_0* [*simp*]: $\text{eval_fls } 0 z = 0$

and *eval_fls_1* [*simp*]: $\text{eval_fls } 1 z = 1$

and *eval_fls_const* [*simp*]: $\text{eval_fls } (\text{fls_const } c) z = c$

and *eval_fls_numeral* [*simp*]: $\text{eval_fls } (\text{numeral } \text{num}) z = \text{numeral } \text{num}$

and *eval_fls_of_nat* [*simp*]: $\text{eval_fls } (\text{of_nat } n) z = \text{of_nat } n$

and *eval_fls_of_int* [*simp*]: $\text{eval_fls } (\text{of_int } m) z = \text{of_int } m$

and *eval_fls_X* [*simp*]: $\text{eval_fls } \text{fls_X } z = z$

and *eval_fls_X_intpow* [*simp*]: $\text{eval_fls } (\text{fls_X_intpow } m) z = z \text{ powi } m$

by (*simp_all add: eval_fls_def*)

lemma *eval_fls_at_0*: $\text{eval_fls } F 0 = (\text{if } \text{fls_subdegree } F \geq 0 \text{ then } \text{fls_nth } F 0 \text{ else } 0)$

by (*cases fls_subdegree F = 0*)

(*simp_all add: eval_fls_def fls_regpart_def eval_fps_at_0*)

lemma *eval_fps_to_fls*:

assumes $\text{norm } z < \text{fps_conv_radius } F$

shows $\text{eval_fls } (\text{fps_to_fls } F) z = \text{eval_fps } F z$

proof (*cases F = 0*)

case [*simp*]: *False*

```

have eval_fps F z = eval_fps (unit_factor F * normalize F) z
  by (metis unit_factor_mult_normalize)
also have ... = eval_fps (unit_factor F * fps_X ^ subdegree F) z
  by simp
also have ... = eval_fps (unit_factor F) z * z ^ subdegree F
  using assms by (subst eval_fps_mult) auto
also have ... = eval_fls (fps_to_fls F) z
  unfolding eval_fls_def fls_base_factor_to_fps_to_fls fls_subdegree_fls_to_fps
    power_int_of_nat ..
finally show ?thesis ..
qed auto

```

```

lemma eval_fls_shift:
  assumes [simp]: z ≠ 0
  shows eval_fls (fls_shift n F) z = eval_fls F z * z powi -n
proof (cases F = 0)
  case [simp]: False
  show ?thesis
  unfolding eval_fls_def
  by (subst fls_base_factor_to_fps_shift, subst fls_shift_subdegree[OF ‹F ≠ 0›],
    subst power_int_diff)
    (auto simp: power_int_minus divide_simps)
qed auto

```

```

lemma eval_fls_add:
  assumes ereal (norm z) < fls_conv_radius F ereal (norm z) < fls_conv_radius
    G z ≠ 0
  shows eval_fls (F + G) z = eval_fls F z + eval_fls G z
  using assms
proof (induction fls_subdegree F fls_subdegree G arbitrary: F G rule: linorder_wlog)
  case (sym F G)
  show ?case
    using sym(1)[of G F] sym(2-) by (simp add: add_ac)
next
  case (le F G)
  show ?case
  proof (cases F = 0 ∨ G = 0)
  case False
  hence [simp]: F ≠ 0 G ≠ 0
    by auto
  note [simp] = ‹z ≠ 0›
  define F' G' where F' = fls_base_factor_to_fps F G' = fls_base_factor_to_fps
    G
  define m n where m = fls_subdegree F n = fls_subdegree G
  have m ≤ n
    using le by (auto simp: m_n_def)
  have conv1: ereal (cmod z) < fps_conv_radius F' ereal (cmod z) < fps_conv_radius
    G'
    using assms le by (simp_all add: F'_G'_def fls_conv_radius_altdef)

```

```

have conv2: ereal (cmod z) < fps_conv_radius (G' * fps_X ^ nat (n - m))
using conv1 by (intro less_le_trans[OF _ fps_conv_radius_mult]) auto
have conv3: ereal (cmod z) < fps_conv_radius (F' + G' * fps_X ^ nat (n -
m))
using conv1 conv2 by (intro less_le_trans[OF _ fps_conv_radius_add])
auto

have eval_fls F z + eval_fls G z = eval_fps F' z * z powi m + eval_fps G' z
* z powi n
unfolding eval_fls_def m_n_def[symmetric] F'_G'_def[symmetric]
by (simp add: power_int_add algebra_simps)
also have ... = (eval_fps F' z + eval_fps G' z * z powi (n - m)) * z powi m
by (simp add: algebra_simps power_int_diff)
also have eval_fps G' z * z powi (n - m) = eval_fps (G' * fps_X ^ nat (n
- m)) z
using assms ⟨m ≤ n⟩ conv1 by (subst eval_fps_mult) (auto simp: power_int_def)
also have eval_fps F' z + ... = eval_fps (F' + G' * fps_X ^ nat (n - m)) z
using conv1 conv2 by (subst eval_fps_add) auto
also have ... = eval_fls (fps_to_fls (F' + G' * fps_X ^ nat (n - m))) z
using conv3 by (subst eval_fps_to_fls) auto
also have ... * z powi m = eval_fls (fls_shift (-m) (fps_to_fls (F' + G' *
fps_X ^ nat (n - m)))) z
by (subst eval_fls_shift) auto
also have fls_shift (-m) (fps_to_fls (F' + G' * fps_X ^ nat (n - m))) = F
+ G
using ⟨m ≤ n⟩
by (simp add: fls_times_fps_to_fls fps_to_fls_power fls_X_power_conv_shift_1
fls_shifted_times_simps F'_G'_def m_n_def)
finally show ?thesis ..
qed auto
qed

```

lemma eval_fls_minus:

```

assumes ereal (norm z) < fls_conv_radius F
shows eval_fls (-F) z = -eval_fls F z
using assms by (simp add: eval_fls_def eval_fps_minus fls_conv_radius_altdef)

```

lemma eval_fls_diff:

```

assumes ereal (norm z) < fls_conv_radius F ereal (norm z) < fls_conv_radius
G
and [simp]: z ≠ 0
shows eval_fls (F - G) z = eval_fls F z - eval_fls G z
proof -
have eval_fls (F + (-G)) z = eval_fls F z - eval_fls G z
using assms by (subst eval_fls_add) (auto simp: eval_fls_minus)
thus ?thesis
by simp
qed

```


lemma *eval_fls_mult*:
assumes $\text{ereal}(\text{norm } z) < \text{fls_conv_radius } F$ $\text{ereal}(\text{norm } z) < \text{fls_conv_radius } G$ $z \neq 0$
shows $\text{eval_fls } (F * G) z = \text{eval_fls } F z * \text{eval_fls } G z$
proof (*cases* $F = 0 \vee G = 0$)
case *False*
hence [*simp*]: $F \neq 0 \ G \neq 0$
by *auto*
note [*simp*] = $\langle z \neq 0 \rangle$
define $F' \ G'$ **where** $F' = \text{fls_base_factor_to_fps } F$ $G' = \text{fls_base_factor_to_fps } G$
define $m \ n$ **where** $m = \text{fls_subdegree } F$ $n = \text{fls_subdegree } G$
have $\text{eval_fls } F z * \text{eval_fls } G z = (\text{eval_fps } F' z * \text{eval_fps } G' z) * z^{\text{powi } (m + n)}$
unfolding *eval_fls_def m_n_def[symmetric] F'_G'_def[symmetric]*
by (*simp add: power_int_add algebra_simps*)
also have $\dots = \text{eval_fps } (F' * G') z * z^{\text{powi } (m + n)}$
using *assms by (subst eval_fps_mult) (auto simp: F'_G'_def fls_conv_radius_altdef)*
also have $\dots = \text{eval_fls } (F * G) z$
by (*simp add: eval_fls_def F'_G'_def m_n_def*) (*simp add: fls_times_def*)
finally show *?thesis ..*
qed *auto*

lemma *eval_fls_power*:
assumes $\text{ereal}(\text{norm } z) < \text{fls_conv_radius } F$ $z \neq 0$
shows $\text{eval_fls } (F \wedge n) z = \text{eval_fls } F z \wedge n$
proof (*induction n*)
case (*Suc n*)
have $\text{eval_fls } (F \wedge \text{Suc } n) z = \text{eval_fls } (F * F \wedge n) z$
by *simp*
also have $\dots = \text{eval_fls } F z * \text{eval_fls } (F \wedge n) z$
using *assms by (subst eval_fls_mult) (auto intro!: less_le_trans[OF fls_conv_radius_power])*
finally show *?case*
using *Suc by simp*
qed *auto*

lemma *eval_fls_eq*:
assumes $N \leq \text{fls_subdegree } F$ $\text{fls_subdegree } F \geq 0 \vee z \neq 0$
assumes $(\lambda n. \text{fls_nth } F (\text{int } n + N) * z^{\text{powi } (\text{int } n + N)}) \text{ sums } S$
shows $\text{eval_fls } F z = S$
proof (*cases z = 0*)
case [*simp*]: *True*
have $(\lambda n. \text{fls_nth } F (\text{int } n + N) * z^{\text{powi } (\text{int } n + N)}) =$
 $(\lambda n. \text{if } n \in (\text{if } N \leq 0 \text{ then } \{\text{nat } (-N)\} \text{ else } \{\}) \text{ then } \text{fls_nth } F (\text{int } n + N)$
 $\text{else } 0)$
by (*auto simp: fun_eq_iff split: if_splits*)
also have $\dots \text{ sums } (\sum n \in (\text{if } N \leq 0 \text{ then } \{\text{nat } (-N)\} \text{ else } \{\}). \text{fls_nth } F (\text{int } n + N))$
by (*rule sums_If_finite_set*) *auto*

```

also have ... = fls_nth F 0
  using assms by auto
also have ... = eval_fls F z
  using assms by (auto simp: eval_fls_def eval_fps_at_0 power_int_0_left_if)
finally show ?thesis
  using assms by (simp add: sums_iff)
next
case [simp]: False
define N' where N' = fls_subdegree F
define d where d = nat (N' - N)

have ( $\lambda n. \text{fls\_nth } F \text{ (int } n + N) * z^{\text{powi } (int } n + N)$ ) sums S
  by fact
also have ?this  $\longleftrightarrow$  ( $\lambda n. \text{fls\_nth } F \text{ (int } (n+d) + N) * z^{\text{powi } (int } (n+d) + N)$ ) sums S
  by (rule sums_zero_iff_shift [symmetric]) (use assms in ‹auto simp: d_def N'_def›)
also have ( $\lambda n. \text{int } (n+d) + N = (\lambda n. \text{int } n + N')$ )
  using assms by (auto simp: N'_def d_def)
finally have ( $\lambda n. \text{fls\_nth } F \text{ (int } n + N') * z^{\text{powi } (int } n + N')$ ) sums S .
  hence ( $\lambda n. z^{\text{powi } (-N')} * (\text{fls\_nth } F \text{ (int } n + N') * z^{\text{powi } (int } n + N'))$ )
sums (z powi (-N') * S)
  by (intro sums_mult)
hence ( $\lambda n. \text{fls\_nth } F \text{ (int } n + N') * z^{\wedge n}$ ) sums (z powi (-N') * S)
  by (simp add: power_int_add power_int_minus field_simps)
thus ?thesis
  by (simp add: eval_fls_def eval_fps_def sums_iff power_int_minus N'_def)
qed

```

```

lemma norm_summable_fls:
  norm z < fls_conv_radius f  $\implies$  summable ( $\lambda n. \text{norm } (\text{fls\_nth } f \ n * z^{\wedge n})$ )
  using norm_summable_fps[of z fls_regpart f] by (simp add: fls_conv_radius_def)

```

```

lemma norm_summable_fls':
  norm z < fls_conv_radius f  $\implies$  summable ( $\lambda n. \text{norm } (\text{fls\_nth } f \ (n + \text{fls\_subdegree } f) * z^{\wedge n})$ )
  using norm_summable_fps[of z fls_base_factor_to_fps f] by (simp add: fls_conv_radius_altdef)

```

```

lemma summable_fls:
  norm z < fls_conv_radius f  $\implies$  summable ( $\lambda n. \text{fls\_nth } f \ n * z^{\wedge n}$ )
  by (rule summable_norm_cancel[OF norm_summable_fls])

```

```

theorem sums_eval_fls:
  fixes f
  defines n  $\equiv$  fls_subdegree f
  assumes norm z < fls_conv_radius f and z  $\neq$  0  $\vee$  n  $\geq$  0
  shows ( $\lambda k. \text{fls\_nth } f \ (int \ k + n) * z^{\text{powi } (int \ k + n)}$ ) sums eval_fls f z
proof (cases z = 0)
  case [simp]: False

```

```

have ( $\lambda k. \text{fps\_nth } (\text{fls\_base\_factor\_to\_fps } f) k * z ^ k * z \text{ powi } n$ ) sums
  ( $\text{eval\_fps } (\text{fls\_base\_factor\_to\_fps } f) z * z \text{ powi } n$ )
using assms(2) by (intro sums_eval_fps sums_mult2) (auto simp: fls_conv_radius_altdef)
thus ?thesis
  by (simp add: power_int_add n_def eval_fls_def mult_ac)
next
case [simp]: True
with assms have  $n \geq 0$ 
  by auto
have ( $\lambda k. \text{fls\_nth } f (int k + n) * z \text{ powi } (int k + n)$ ) sums
  ( $\sum_{k \in (\text{if } n \leq 0 \text{ then } \{\text{nat } (-n)\} \text{ else } \{\})}. \text{fls\_nth } f (int k + n) * z \text{ powi } (int k + n)$ )
  by (intro sums_finite) (auto split: if_splits)
also have  $\dots = \text{eval\_fls } f z$ 
  using  $\langle n \geq 0 \rangle$  by (auto simp: eval_fls_at_0 n_def not_le)
finally show ?thesis .
qed

```

```

lemma holomorphic_on_eval_fls:
  fixes f
  defines  $n \equiv \text{fls\_subdegree } f$ 
  assumes  $A \subseteq \text{eball } 0 (\text{fls\_conv\_radius } f) - (\text{if } n \geq 0 \text{ then } \{\} \text{ else } \{0\})$ 
  shows  $\text{eval\_fls } f \text{ holomorphic\_on } A$ 
proof (cases  $n \geq 0$ )
  case True
  have  $\text{eval\_fls } f = (\lambda z. \text{eval\_fps } (\text{fls\_base\_factor\_to\_fps } f) z * z ^ \text{nat } n)$ 
    using True by (simp add: fun_eq_iff eval_fls_def power_int_def n_def)
  moreover have  $\dots \text{ holomorphic\_on } A$ 
    using True assms(2) by (intro holomorphic_intros) (auto simp: fls_conv_radius_altdef)
  ultimately show ?thesis
    by simp
  next
  case False
  show ?thesis using assms
    unfolding  $\text{eval\_fls\_def}$  by (intro holomorphic_intros) (auto simp: fls_conv_radius_altdef)
qed

```

```

lemma holomorphic_on_eval_fls' [holomorphic_intros]:
  assumes  $g \text{ holomorphic\_on } A$ 
  assumes  $g ' A \subseteq \text{eball } 0 (\text{fls\_conv\_radius } f) - (\text{if } \text{fls\_subdegree } f \geq 0 \text{ then } \{\} \text{ else } \{0\})$ 
  shows  $(\lambda x. \text{eval\_fls } f (g x)) \text{ holomorphic\_on } A$ 
  by (meson assms holomorphic_on_compose holomorphic_on_eval_fls holomorphic_transform_o_def)

```

```

lemma continuous_on_eval_fls:
  fixes f
  defines  $n \equiv \text{fls\_subdegree } f$ 
  assumes  $A \subseteq \text{eball } 0 (\text{fls\_conv\_radius } f) - (\text{if } n \geq 0 \text{ then } \{\} \text{ else } \{0\})$ 

```

shows *continuous_on* A (*eval_fls* f)
using *assms* *holomorphic_on_eval_fls* *holomorphic_on_imp_continuous_on*
by *blast*

lemma *continuous_on_eval_fls'* [*continuous_intros*]:
fixes f
defines $n \equiv \text{fls_subdegree } f$
assumes $g \cdot A \subseteq \text{eball } 0 (\text{fls_conv_radius } f) - (\text{if } n \geq 0 \text{ then } \{\} \text{ else } \{0\})$
assumes *continuous_on* A g
shows *continuous_on* A ($\lambda x. \text{eval_fls } f (g x)$)
by (*metis* *assms* *continuous_on_compose2* *continuous_on_eval_fls* *order.refl*)

lemmas *has_field_derivative_eval_fps'* [*derivative_intros*] =
DERIV_chain2[*OF* *has_field_derivative_eval_fps*]

lemma *has_field_derivative_eval_fls*:
assumes $z \in \text{eball } 0 (\text{fls_conv_radius } f) - \{0\}$
shows (*eval_fls* f *has_field_derivative* *eval_fls* (*fls_deriv* f) z) (*at* z *within* A)
proof –
define g **where** $g = \text{fls_base_factor_to_fps } f$
define n **where** $n = \text{fls_subdegree } f$
have [*simp*]: *fps_conv_radius* $g = \text{fls_conv_radius } f$
by (*simp* *add*: *fls_conv_radius_altdef* *g_def*)
have *conv1*: *fps_conv_radius* (*fps_deriv* $g * \text{fps}_X$) $\geq \text{fls_conv_radius } f$
by (*intro* *fps_conv_radius_mult_ge* *order.trans*[*OF* *fps_conv_radius_deriv*])
auto
have *conv2*: *fps_conv_radius* (*of_int* $n * g$) $\geq \text{fls_conv_radius } f$
by (*intro* *fps_conv_radius_mult_ge*) *auto*
have *conv3*: *fps_conv_radius* (*fps_deriv* $g * \text{fps}_X + \text{of_int } n * g$) $\geq \text{fls_conv_radius } f$
by (*intro* *fps_conv_radius_add_ge* *conv1* *conv2*)

have [*simp*]: *fps_conv_radius* $g = \text{fls_conv_radius } f$
by (*simp* *add*: *g_def* *fls_conv_radius_altdef*)
have ($\lambda z. \text{eval_fps } g z * z^{\text{powi } \text{fls_subdegree } f}$) *has_field_derivative*
(*eval_fps* (*fps_deriv* g) $z * z^{\text{powi } n} + \text{of_int } n * z^{\text{powi } (n - 1)} * \text{eval_fps } g z$)
(*at* z *within* A)
using *assms* **by** (*auto* *intro!*: *derivative_eq_intros* *simp*: *n_def*)
also **have** ($\lambda z. \text{eval_fps } g z * z^{\text{powi } \text{fls_subdegree } f}$) = *eval_fls* f
by (*simp* *add*: *eval_fls_def* *g_def* *fun_eq_iff*)
also **have** *eval_fps* (*fps_deriv* g) $z * z^{\text{powi } n} + \text{of_int } n * z^{\text{powi } (n - 1)} * \text{eval_fps } g z =$
($z * \text{eval_fps } (\text{fps_deriv } g) z + \text{of_int } n * \text{eval_fps } g z$) $* z^{\text{powi } (n - 1)}$
using *assms* **by** (*auto* *simp*: *power_int_diff* *field_simps*)
also **have** ($z * \text{eval_fps } (\text{fps_deriv } g) z + \text{of_int } n * \text{eval_fps } g z$) =
eval_fps (*fps_deriv* $g * \text{fps}_X + \text{of_int } n * g$) z

```

    using conv1 conv2 assms fps_conv_radius_deriv[of g]
    by (subst eval_fps_add) (auto simp: eval_fps_mult)
  also have ... = eval_fls (fps_to_fls (fps_deriv g * fps_X + of_int n * g)) z
    using conv3 assms by (subst eval_fps_to_fls) auto
  also have ... * z powi (n - 1) = eval_fls (fls_shift (1 - n) (fps_to_fls
(fps_deriv g * fps_X + of_int n * g))) z
    using assms by (subst eval_fls_shift) auto
  also have fls_shift (1 - n) (fps_to_fls (fps_deriv g * fps_X + of_int n * g))
= fls_deriv f
    by (intro fls_eqI) (auto simp: g_def n_def algebra_simps eq_commute[of _
fls_subdegree f])
  finally show ?thesis .
qed

```

lemma *eval_fls_deriv*:

```

  assumes z ∈ eball 0 (fls_conv_radius F) - {0}
  shows eval_fls (fls_deriv F) z = deriv (eval_fls F) z
  by (metis DERIV_imp_deriv assms has_field_derivative_eval_fls)

```

lemma *analytic_on_eval_fls*:

```

  assumes A ⊆ eball 0 (fls_conv_radius f) - (if fls_subdegree f ≥ 0 then {} else
{0})
  shows eval_fls f analytic_on A
proof (rule analytic_on_subset [OF __ assms])
  show eval_fls f analytic_on eball 0 (fls_conv_radius f) - (if fls_subdegree f ≥
0 then {} else {0})
    using holomorphic_on_eval_fls[OF order.refl]
    by (subst analytic_on_open) auto
qed

```

lemma *analytic_on_eval_fls'* [analytic_intros]:

```

  assumes g analytic_on A
  assumes g ' A ⊆ eball 0 (fls_conv_radius f) - (if fls_subdegree f ≥ 0 then {}
else {0})
  shows (λx. eval_fls f (g x)) analytic_on A
proof -
  have eval_fls f ∘ g analytic_on A
    by (intro analytic_on_compose[OF assms(1) analytic_on_eval_fls]) (use assms
in auto)
  thus ?thesis
    by (simp add: o_def)
qed

```

lemma *continuous_eval_fls* [continuous_intros]:

```

  assumes z ∈ eball 0 (fls_conv_radius F) - (if fls_subdegree F ≥ 0 then {} else
{0})
  shows continuous (at z within A) (eval_fls F)
proof -
  have isCont (eval_fls F) z

```

```

    using continuous_on_eval_fls[OF order.refl] assms
  by (subst (asm) continuous_on_eq_continuous_at) auto
thus ?thesis
  using continuous_at_imp_continuous_at_within by blast
qed

```

named_theorems laurent_expansion_intros

```

lemma has_laurent_expansion_imp_asymp_equiv_0:
  assumes F: f has_laurent_expansion F
  defines n  $\equiv$  fls_subdegree F
  shows f  $\sim$ [at 0] ( $\lambda z. fls\_nth F n * z^{powi n}$ )
proof (cases F = 0)
  case True
  thus ?thesis using assms
    by (auto simp: has_laurent_expansion_def)
next
  case [simp]: False
  define G where G = fls_base_factor_to_fps F
  have fls_conv_radius F > 0
    using F by (auto simp: has_laurent_expansion_def)
  hence isCont (eval_fps G) 0
    by (intro continuous_intros) (auto simp: G_def fps_conv_radius_fls_regpart
zero_ereal_def)
  hence lim: eval_fps G  $\rightarrow$  eval_fps G 0
    by (meson isContD)
  have [simp]: fps_nth G 0  $\neq$  0
    by (auto simp: G_def)

  have f  $\sim$ [at 0] eval_fls F
    using F by (intro asymp_equiv_refl_ev) (auto simp: has_laurent_expansion_def
eq_commute)
  also have ... = ( $\lambda z. eval\_fps G z * z^{powi n}$ )
    by (intro ext) (simp_all add: eval_fls_def G_def n_def)
  also have ...  $\sim$ [at 0] ( $\lambda z. fps\_nth G 0 * z^{powi n}$ ) using lim
    by (intro asymp_equiv_intros tendsto_imp_asymp_equiv_const) (auto simp:
eval_fps_at_0)
  also have fps_nth G 0 = fls_nth F n
    by (simp add: G_def n_def)
  finally show ?thesis
    by simp
qed

```

```

lemma has_laurent_expansion_imp_asymp_equiv:
  assumes F: ( $\lambda w. f (z + w)$ ) has_laurent_expansion F
  defines n  $\equiv$  fls_subdegree F

```

shows $f \sim[at\ z] (\lambda w. fls_nth\ F\ n * (w - z)\ powi\ n)$
using $has_laurent_expansion_imp_asympt_equiv_0[OF\ assms(1)]$ **unfolding**
 n_def
by $(simp\ add: at_to_0[of\ z]\ asympt_equiv_filtermap_iff\ add_ac)$

lemmas $[tendsto_intros\ del] = tendsto_power_int$

lemma $has_laurent_expansion_imp_tendsto_0$:
assumes $F: f\ has_laurent_expansion\ F$ **and** $fls_subdegree\ F \geq 0$
shows $f \rightarrow 0$
proof $(rule\ asympt_equiv_tendsto_transfer)$
show $(\lambda z. fls_nth\ F\ (fls_subdegree\ F) * z\ powi\ fls_subdegree\ F) \sim[at\ 0] f$
by $(rule\ asympt_equiv_symI, rule\ has_laurent_expansion_imp_asympt_equiv_0)$
 $fact$
show $(\lambda z. fls_nth\ F\ (fls_subdegree\ F) * z\ powi\ fls_subdegree\ F) \rightarrow 0$
by $(rule\ tendsto_eq_intros\ refl\ |\ use\ assms(2)\ in\ simp)+$
 $(use\ assms(2)\ in\ \langle auto\ simp: power_int_0_left_if \rangle)$
qed

lemma $has_laurent_expansion_imp_tendsto$:
assumes $F: (\lambda w. f\ (z + w))\ has_laurent_expansion\ F$ **and** $fls_subdegree\ F \geq 0$
shows $f \rightarrow z$
using $has_laurent_expansion_imp_tendsto_0[OF\ assms]$
by $(simp\ add: at_to_0[of\ z]\ filterlim_filtermap\ add_ac)$

lemma $has_laurent_expansion_imp_filterlim_infinity_0$:
assumes $F: f\ has_laurent_expansion\ F$ **and** $fls_subdegree\ F < 0$
shows $filterlim\ f\ at_infinity\ (at\ 0)$
proof $(rule\ asympt_equiv_at_infinity_transfer)$
have $[simp]: F \neq 0$
using $assms(2)$ **by** $auto$
show $(\lambda z. fls_nth\ F\ (fls_subdegree\ F) * z\ powi\ fls_subdegree\ F) \sim[at\ 0] f$
by $(rule\ asympt_equiv_symI, rule\ has_laurent_expansion_imp_asympt_equiv_0)$
 $fact$
show $filterlim\ (\lambda z. fls_nth\ F\ (fls_subdegree\ F) * z\ powi\ fls_subdegree\ F)\ at_infinity\ (at\ 0)$
by $(rule\ tendsto_mult_filterlim_at_infinity\ tendsto_const$
 $filterlim_power_int_neg_at_infinity\ |\ use\ assms(2)\ in\ simp)+$
 $(auto\ simp: eventually_at_filter)$
qed

lemma $has_laurent_expansion_imp_neg_fls_subdegree$:
assumes $F: f\ has_laurent_expansion\ F$
and $infy: filterlim\ f\ at_infinity\ (at\ 0)$
shows $fls_subdegree\ F < 0$
proof $(rule\ ccontr)$

```

assume asm:  $\neg$  fls_subdegree F < 0
define ff where ff = ( $\lambda z$ . fls_nth F (fls_subdegree F)
  * z powi fls_subdegree F)

have (ff  $\longrightarrow$  (if fls_subdegree F = 0 then fls_nth F 0 else 0)) (at 0)
  using asm unfolding ff_def
  by (auto intro!: tendsto_eq_intros)
moreover have filterlim ff at_infinity (at 0)
proof (rule asympt_equiv_at_infinity_transfer)
  show f  $\sim$ [at 0] ff unfolding ff_def
  using has_laurent_expansion_imp_asympt_equiv_0[OF F] unfolding ff_def
  .
  show filterlim f at_infinity (at 0) by fact
qed
ultimately show False
  using not_tendsto_and_filterlim_at_infinity[of at (0::complex)] by auto
qed

lemma has_laurent_expansion_imp_filterlim_infinity:
  assumes F: ( $\lambda w$ . f (z + w)) has_laurent_expansion F and fls_subdegree F < 0
  shows filterlim f at_infinity (at z)
  using has_laurent_expansion_imp_filterlim_infinity_0[OF assms]
  by (simp add: at_to_0[of z] filterlim_filtermap add_ac)

lemma has_laurent_expansion_imp_is_pole_0:
  assumes F: f has_laurent_expansion F and fls_subdegree F < 0
  shows is_pole f 0
  using has_laurent_expansion_imp_filterlim_infinity_0[OF assms]
  by (simp add: is_pole_def)

lemma is_pole_0_imp_neg_fls_subdegree:
  assumes F: f has_laurent_expansion F and is_pole f 0
  shows fls_subdegree F < 0
  using F assms(2) has_laurent_expansion_imp_neg_fls_subdegree is_pole_def
  by blast

lemma has_laurent_expansion_imp_is_pole:
  assumes F: ( $\lambda x$ . f (z + x)) has_laurent_expansion F and fls_subdegree F < 0
  shows is_pole f z
  using has_laurent_expansion_imp_is_pole_0[OF assms]
  by (simp add: is_pole_shift_0')

lemma is_pole_imp_neg_fls_subdegree:
  assumes F: ( $\lambda x$ . f (z + x)) has_laurent_expansion F and is_pole f z
  shows fls_subdegree F < 0
proof -
  have is_pole ( $\lambda x$ . f (z + x)) 0
    using assms(2) is_pole_shift_0 by blast

```



```

then show ?thesis
  using F is_pole_0_imp_neg_fls_subdegree by blast
qed

lemma is_pole_fls_subdegree_iff:
  assumes  $(\lambda x. f(z + x))$  has_laurent_expansion F
  shows is_pole f z  $\longleftrightarrow$  fls_subdegree F < 0
  using assms is_pole_imp_neg_fls_subdegree has_laurent_expansion_imp_is_pole
  by auto

lemma
  assumes f has_laurent_expansion F
  shows has_laurent_expansion_isolated_0: isolated_singularity_at f 0
  and has_laurent_expansion_not_essential_0: not_essential f 0
proof -
  from assms have eventually  $(\lambda z. \text{eval\_fls } F z = f z)$  (at 0)
  by (auto simp: has_laurent_expansion_def)
  then obtain r where  $r: r > 0 \wedge z. z \in \text{ball } 0 r - \{0\} \implies \text{eval\_fls } F z = f z$ 
  by (auto simp: eventually_at_filter_ball_def eventually_nhds_metric)

  have fls_conv_radius F > 0
  using assms by (auto simp: has_laurent_expansion_def)
  then obtain R :: real where  $R: R > 0 \wedge R \leq \min r (\text{fls\_conv\_radius } F)$ 
  using  $\langle r > 0 \rangle$  by (metis dual_order.strict_implies_order ereal_dense2 ereal_less(2)
min_def)

  have eval_fls F holomorphic_on ball 0 R - {0}
  using r R by (intro holomorphic_intros ball_eball_mono Diff_mono) (auto
simp: ereal_le_less)
  also have  $?this \longleftrightarrow f \text{ holomorphic\_on ball } 0 R - \{0\}$ 
  using r R by (intro holomorphic_cong) auto
  also have  $\dots \longleftrightarrow f \text{ analytic\_on ball } 0 R - \{0\}$ 
  by (subst analytic_on_open) auto
  finally show isolated_singularity_at f 0
  unfolding isolated_singularity_at_def using  $\langle R > 0 \rangle$  by blast

show not_essential f 0
proof (cases fls_subdegree F  $\geq$  0)
  case True
  hence  $f -0 \rightarrow \text{fls\_nth } F 0$ 
  by (intro has_laurent_expansion_imp_tendsto_0[OF assms])
  thus ?thesis
  by (auto simp: not_essential_def)
next
  case False
  hence is_pole f 0
  by (intro has_laurent_expansion_imp_is_pole_0[OF assms]) auto
  thus ?thesis
  by (auto simp: not_essential_def)

```

qed

qed

lemma

assumes $(\lambda w. f(z + w))$ has_laurent_expansion F
 shows has_laurent_expansion_isolated: isolated_singularity_at f z
 and has_laurent_expansion_not_essential: not_essential f z
 using has_laurent_expansion_isolated_0[OF assms] has_laurent_expansion_not_essential_0[OF assms]
 by (simp_all add: isolated_singularity_at_shift_0 not_essential_shift_0)

lemma has_laurent_expansion_fps:

assumes f has_fps_expansion F
 shows f has_laurent_expansion fps_to_fls F
 proof -
 from assms have radius: $0 < \text{fps_conv_radius } F$ and eval: $\forall_F z$ in nhds 0.
 eval_fps F z = f z
 by (auto simp: has_fps_expansion_def)
 from eval have eval': $\forall_F z$ in at 0. eval_fps F z = f z
 using eventually_at_filter eventually_mono by fastforce
 moreover have eventually $(\lambda z. z \in \text{eball } 0 (\text{fps_conv_radius } F) - \{0\})$ (at 0)
 using radius by (intro eventually_at_in_open) (auto simp: zero_ereal_def)
 ultimately have eventually $(\lambda z. \text{eval_fls } (\text{fps_to_fls } F) z = f z)$ (at 0)
 by eventually_elim (auto simp: eval_fps_to_fls)
 thus ?thesis using radius
 by (auto simp: has_laurent_expansion_def)

qed

lemma has_laurent_expansion_const [simp, intro, laurent_expansion_intros]:

$(\lambda_. c)$ has_laurent_expansion fls_const c
 by (auto simp: has_laurent_expansion_def)

lemma has_laurent_expansion_0 [simp, intro, laurent_expansion_intros]:

$(\lambda_. 0)$ has_laurent_expansion 0
 by (auto simp: has_laurent_expansion_def)

lemma has_fps_expansion_0_iff: f has_fps_expansion 0 \longleftrightarrow eventually $(\lambda z. f z = 0)$ (nhds 0)

by (auto simp: has_fps_expansion_def)

lemma has_laurent_expansion_1 [simp, intro, laurent_expansion_intros]:

$(\lambda_. 1)$ has_laurent_expansion 1
 by (auto simp: has_laurent_expansion_def)

lemma has_laurent_expansion_numeral [simp, intro, laurent_expansion_intros]:

$(\lambda_. \text{numeral } n)$ has_laurent_expansion numeral n
 by (auto simp: has_laurent_expansion_def)

lemma has_laurent_expansion_fps_X_power [laurent_expansion_intros]:

($\lambda x. x \hat{=} n$) *has_laurent_expansion* (*fls_X_intpow* n)
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_fps_X_power_int* [*laurent_expansion_intros*]:
($\lambda x. x \text{ powi } n$) *has_laurent_expansion* (*fls_X_intpow* n)
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_fps_X* [*laurent_expansion_intros*]:
($\lambda x. x$) *has_laurent_expansion* *fls_X*
by (*auto simp: has_laurent_expansion_def*)

lemma *has_laurent_expansion_cmult_left* [*laurent_expansion_intros*]:
assumes *f has_laurent_expansion F*
shows ($\lambda x. c * f x$) *has_laurent_expansion* *fls_const c * F*
proof –
from *assms have eventually* ($\lambda z. z \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\}$) (*at 0*)
by (*intro eventually_at_in_open*) (*auto simp: has_laurent_expansion_def zero_ereal_def*)
moreover from *assms have eventually* ($\lambda z. \text{eval_fls } F z = f z$) (*at 0*)
by (*auto simp: has_laurent_expansion_def*)
**ultimately have eventually ($\lambda z. \text{eval_fls } (\text{fls_const } c * F) z = c * f z$) (*at 0*)
by *eventually_elim (simp_all add: eval_fls_mult)*
with *assms show ?thesis*
by (*auto simp: has_laurent_expansion_def intro!: less_le_trans[OF fls_conv_radius_mult]*)
qed**

lemma *has_laurent_expansion_cmult_right* [*laurent_expansion_intros*]:
assumes *f has_laurent_expansion F*
shows ($\lambda x. f x * c$) *has_laurent_expansion* *F * fls_const c*
proof –
have *F * fls_const c = fls_const c * F*
by (*intro fls_eqI*) (*auto simp: mult.commute*)
with *has_laurent_expansion_cmult_left [OF assms]* **show** *?thesis*
by (*simp add: mult.commute*)
qed

lemma *has_fps_expansion_scaleR* [*fps_expansion_intros*]:
fixes *F :: 'a :: {banach, real_normed_div_algebra, comm_ring_1}* *fps*
shows *f has_fps_expansion F* \implies ($\lambda x. c *_R f x$) *has_fps_expansion* *fps_const*
(*of_real c*) * *F*
unfolding *scaleR_conv_of_real* **by** (*intro fps_expansion_intros*)

lemma *has_laurent_expansion_scaleR* [*laurent_expansion_intros*]:
f has_laurent_expansion F \implies ($\lambda x. c *_R f x$) *has_laurent_expansion* *fls_const*
(*of_real c*) * *F*
unfolding *scaleR_conv_of_real* **by** (*intro laurent_expansion_intros*)

lemma *has_laurent_expansion_minus* [*laurent_expansion_intros*]:
assumes *f has_laurent_expansion F*

shows $(\lambda x. -f x)$ has_laurent_expansion $-F$
proof –
from *assms* **have** eventually $(\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\})$ (at 0)
by (intro eventually_at_in_open) (auto simp: has_laurent_expansion_def zero_ereal_def)
moreover from *assms* **have** eventually $(\lambda x. \text{eval_fls } F x = f x)$ (at 0)
by (auto simp: has_laurent_expansion_def)
ultimately have eventually $(\lambda x. \text{eval_fls } (-F) x = -f x)$ (at 0)
by eventually_elim (auto simp: eval_fls_minus)
thus ?thesis using *assms* **by** (auto simp: has_laurent_expansion_def)
qed

lemma has_laurent_expansion_add [laurent_expansion_intros]:
assumes *f* has_laurent_expansion *F* *g* has_laurent_expansion *G*
shows $(\lambda x. f x + g x)$ has_laurent_expansion $F + G$
proof –
from *assms* **have** $0 < \min (\text{fls_conv_radius } F) (\text{fls_conv_radius } G)$
by (auto simp: has_laurent_expansion_def)
also have $\dots \leq \text{fls_conv_radius } (F + G)$
by (rule fls_conv_radius_add)
finally have radius: $\dots > 0$.

from *assms* **have** eventually $(\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\})$ (at 0)
by (intro eventually_at_in_open; force simp: has_laurent_expansion_def zero_ereal_def)+
moreover have eventually $(\lambda x. \text{eval_fls } F x = f x)$ (at 0)
and eventually $(\lambda x. \text{eval_fls } G x = g x)$ (at 0)
using *assms* **by** (auto simp: has_laurent_expansion_def)
ultimately have eventually $(\lambda x. \text{eval_fls } (F + G) x = f x + g x)$ (at 0)
by eventually_elim (auto simp: eval_fls_add)
with radius **show** ?thesis **by** (auto simp: has_laurent_expansion_def)
qed

lemma has_laurent_expansion_diff [laurent_expansion_intros]:
assumes *f* has_laurent_expansion *F* *g* has_laurent_expansion *G*
shows $(\lambda x. f x - g x)$ has_laurent_expansion $F - G$
using has_laurent_expansion_add[of *f* *F* $\lambda x. -g x - G$] *assms*
by (simp add: has_laurent_expansion_minus)

lemma has_laurent_expansion_mult [laurent_expansion_intros]:
assumes *f* has_laurent_expansion *F* *g* has_laurent_expansion *G*
shows $(\lambda x. f x * g x)$ has_laurent_expansion $F * G$
proof –
from *assms* **have** $0 < \min (\text{fls_conv_radius } F) (\text{fls_conv_radius } G)$
by (auto simp: has_laurent_expansion_def)
also have $\dots \leq \text{fls_conv_radius } (F * G)$
by (rule fls_conv_radius_mult)

finally have $\text{radius: } \dots > 0$.

from assms **have** $\text{eventually } (\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\})$ (at 0)
 $\text{eventually } (\lambda x. x \in \text{eball } 0 (\text{fls_conv_radius } G) - \{0\})$ (at 0)
by (intro $\text{eventually_at_in_open}$; force $\text{simp: has_laurent_expansion_def zero_ereal_def}$) +
moreover have $\text{eventually } (\lambda x. \text{eval_fls } F x = f x)$ (at 0)
 and $\text{eventually } (\lambda x. \text{eval_fls } G x = g x)$ (at 0)
 using assms **by** (auto $\text{simp: has_laurent_expansion_def}$)
ultimately have $\text{eventually } (\lambda x. \text{eval_fls } (F * G) x = f x * g x)$ (at 0)
 by eventually_elim (auto $\text{simp: eval_fls_mult}$)
with radius **show** $?thesis$ **by** (auto $\text{simp: has_laurent_expansion_def}$)
qed

lemma $\text{has_fps_expansion_power}$ [$\text{fps_expansion_intros}$]:
 fixes $F :: 'a :: \{\text{banach, real_normed_div_algebra, comm_ring_1}\}$ fps
 shows $f \text{ has_fps_expansion } F \implies (\lambda x. f x \hat{=} m) \text{ has_fps_expansion } F \hat{=} m$
by (induction m) (auto intro!: $\text{fps_expansion_intros}$)

lemma $\text{has_laurent_expansion_power}$ [$\text{laurent_expansion_intros}$]:
 assumes $f \text{ has_laurent_expansion } F$
 shows $(\lambda x. f x \hat{=} n) \text{ has_laurent_expansion } F \hat{=} n$
by (induction n) (auto intro!: $\text{laurent_expansion_intros assms}$)

lemma $\text{has_laurent_expansion_sum}$ [$\text{laurent_expansion_intros}$]:
 assumes $\bigwedge x. x \in I \implies f x \text{ has_laurent_expansion } F x$
 shows $(\lambda y. \sum x \in I. f x y) \text{ has_laurent_expansion } (\sum x \in I. F x)$
using assms **by** (induction I rule: $\text{infinite_finite_induct}$) (auto intro!: $\text{laurent_expansion_intros}$)

lemma $\text{has_laurent_expansion_prod}$ [$\text{laurent_expansion_intros}$]:
 assumes $\bigwedge x. x \in I \implies f x \text{ has_laurent_expansion } F x$
 shows $(\lambda y. \prod x \in I. f x y) \text{ has_laurent_expansion } (\prod x \in I. F x)$
using assms **by** (induction I rule: $\text{infinite_finite_induct}$) (auto intro!: $\text{laurent_expansion_intros}$)

lemma $\text{has_laurent_expansion_deriv}$ [$\text{laurent_expansion_intros}$]:
 assumes $f \text{ has_laurent_expansion } F$
 shows $\text{deriv } f \text{ has_laurent_expansion } \text{fls_deriv } F$

proof –

have $\text{eventually } (\lambda z. z \in \text{eball } 0 (\text{fls_conv_radius } F) - \{0\})$ (at 0)
 using assms **by** (intro $\text{eventually_at_in_open}$)
 (auto $\text{simp: has_laurent_expansion_def zero_ereal_def}$)
moreover from assms **have** $\text{eventually } (\lambda z. \text{eval_fls } F z = f z)$ (at 0)
 by (auto $\text{simp: has_laurent_expansion_def}$)
then obtain s **where** $\text{open } s$ $0 \in s$ **and** $s: \bigwedge w. w \in s - \{0\} \implies \text{eval_fls } F w = f w$
 by (auto $\text{simp: eventually_nhds eventually_at_filter}$)
hence $\text{eventually } (\lambda w. w \in s - \{0\})$ (at 0)
 by (intro $\text{eventually_at_in_open}$) auto
ultimately have $\text{eventually } (\lambda z. \text{eval_fls } (\text{fls_deriv } F) z = \text{deriv } f z)$ (at 0)

```

proof eventually_elim
  case (elim z)
  hence eval_fls (fls_deriv F) z = deriv (eval_fls F) z
    by (simp add: eval_fls_deriv)
  also have eventually ( $\lambda w. w \in s - \{0\}$ ) (nhds z)
    using elim and ‹open s› by (intro eventually_nhds_in_open) auto
  hence eventually ( $\lambda w. \text{eval\_fls } F w = f w$ ) (nhds z)
    by eventually_elim (use s in auto)
  hence deriv (eval_fls F) z = deriv f z
    by (intro deriv_cong_ev refl)
  finally show ?case .
qed
with assms show ?thesis
  by (auto simp: has_laurent_expansion_def intro!: less_le_trans[OF fls_conv_radius_deriv])
qed

```

```

lemma has_laurent_expansion_shift [laurent_expansion_intros]:
  assumes f has_laurent_expansion F
  shows ( $\lambda x. f x * x \text{ powi } n$ ) has_laurent_expansion (fls_shift (-n) F)
proof -
  have eventually ( $\lambda x. x \in \text{eball } 0 (\text{fls\_conv\_radius } F) - \{0\}$ ) (at 0)
    using assms by (intro eventually_at_in_open) (auto simp: has_laurent_expansion_def
zero_ereal_def)
  moreover have eventually ( $\lambda x. \text{eval\_fls } F x = f x$ ) (at 0)
    using assms by (auto simp: has_laurent_expansion_def)
  ultimately have eventually ( $\lambda x. \text{eval\_fls } (\text{fls\_shift } (-n) F) x = f x * x \text{ powi } n$ ) (at 0)
    by eventually_elim (auto simp: eval_fls_shift assms)
  with assms show ?thesis by (auto simp: has_laurent_expansion_def)
qed

```

```

lemma has_laurent_expansion_shift' [laurent_expansion_intros]:
  assumes f has_laurent_expansion F
  shows ( $\lambda x. f x * x \text{ powi } (-n)$ ) has_laurent_expansion (fls_shift n F)
  using has_laurent_expansion_shift[OF assms, of -n] by simp

```

```

lemma has_laurent_expansion_deriv':
  assumes f has_laurent_expansion F
  assumes open A 0  $\in A \wedge x. x \in A - \{0\} \implies (f \text{ has\_field\_derivative } f' x)$  (at x)
  shows f' has_laurent_expansion fls_deriv F
proof -
  have deriv f has_laurent_expansion fls_deriv F
    by (intro laurent_expansion_intros assms)
  also have ?this  $\longleftrightarrow$  ?thesis
proof (intro has_laurent_expansion_cong refl)
  have eventually ( $\lambda z. z \in A - \{0\}$ ) (at 0)
    by (intro eventually_at_in_open assms)

```

```

thus eventually ( $\lambda z. \text{deriv } f z = f' z$ ) (at 0)
  by eventually_elim (auto intro!: DERIV_imp_deriv assms)
qed
finally show ?thesis .
qed

```

definition laurent_expansion :: (complex \Rightarrow complex) \Rightarrow complex \Rightarrow complex fls
where

```

laurent_expansion f z =
  (if eventually ( $\lambda z. f z = 0$ ) (at z) then 0
   else fls_shift (-zorder f z) (fps_to_fls (fps_expansion (zor_poly f z) z)))

```

lemma laurent_expansion_cong:

```

assumes eventually ( $\lambda w. f w = g w$ ) (at z) z = z'
shows laurent_expansion f z = laurent_expansion g z'
unfolding laurent_expansion_def
using zor_poly_cong[OF assms(1,2)] zorder_cong[OF assms] assms
by (intro if_cong refl) (auto elim: eventually_elim2)

```

theorem not_essential_has_laurent_expansion_0:

```

assumes isolated_singularity_at f 0 not_essential f 0
shows f has_laurent_expansion laurent_expansion f 0
proof (cases  $\exists_F w$  in at 0.  $f w \neq 0$ )
  case False
  have ( $\lambda_. 0$ ) has_laurent_expansion 0
    by simp
  also have ?this  $\longleftrightarrow$  f has_laurent_expansion 0
    using False by (intro has_laurent_expansion_cong) (auto simp: frequently_def)
  finally show ?thesis
    using False by (simp add: laurent_expansion_def frequently_def)

```

next

```

case True
define n where n = zorder f 0
obtain r where r: zor_poly f 0 0  $\neq 0$  zor_poly f 0 holomorphic_on cball 0 r r
  > 0

```

$$\forall w \in \text{cball } 0 \ r - \{0\}. f w = \text{zor_poly } f \ 0 \ w * w^{\text{powi } n} \wedge \text{zor_poly } f \ 0 \ w \neq 0$$

```

using zorder_exist[OF assms True] unfolding n_def by auto
have holo: zor_poly f 0 holomorphic_on ball 0 r
by (rule holomorphic_on_subset[OF r(2)]) auto

```

define F **where** F = fps_expansion (zor_poly f 0) 0

```

have F: zor_poly f 0 has_fps_expansion F
  unfolding F_def by (rule has_fps_expansion_fps_expansion[OF _ _ holo])
  (use <r > 0) in auto)
have ( $\lambda z. \text{zor\_poly } f \ 0 \ z * z^{\text{powi } n}$ ) has_laurent_expansion fls_shift (-n)
  (fps_to_fls F)
by (intro laurent_expansion_intros has_laurent_expansion_fps[OF F])
also have ?this  $\longleftrightarrow$  f has_laurent_expansion fls_shift (-n) (fps_to_fls F)

```

```

  by (intro has_laurent_expansion_cong refl eventually_mono[OF eventually_at_in_open[of
ball 0 r]])
    (use r in ⟨auto simp: complex_powr_of_int⟩)
  finally show ?thesis using True
    by (simp add: laurent_expansion_def F_def n_def frequently_def)
qed

```

lemma *not_essential_has_laurent_expansion:*

```

  assumes isolated_singularity_at f z not_essential f z
  shows (λx. f (z + x)) has_laurent_expansion laurent_expansion f z
proof -
  from assms(1) have iso:isolated_singularity_at (λx. f (z + x)) 0
    by (simp add: isolated_singularity_at_shift_0)
  moreover from assms(2) have ness:not_essential (λx. f (z + x)) 0
    by (simp add: not_essential_shift_0)
  ultimately have (λx. f (z + x)) has_laurent_expansion laurent_expansion (λx.
f (z + x)) 0
    by (rule not_essential_has_laurent_expansion_0)

```

also have ... = laurent_expansion f z

proof (cases $\exists_F w$ in at z. f w \neq 0)

case False

then have $\forall_F w$ in at z. f w = 0 **using** not_frequently **by** force

then have laurent_expansion (λx. f (z + x)) 0 = 0

by (smt (verit, best) add commute eventually_at_to_0 eventually_mono laurent_expansion_def)

moreover have laurent_expansion f z = 0

using $\langle \forall_F w$ in at z. f w = 0 **unfolding** laurent_expansion_def **by** auto

ultimately show ?thesis **by** auto

next

case True

define df **where** df=zor_poly (λx. f (z + x)) 0

define g **where** g=(λu. u-z)

have fps_expansion df 0

= fps_expansion (df o g) z

proof -

have $\exists_F w$ in at 0. f (z + w) \neq 0 **using** True

by (smt (verit, best) add commute eventually_at_to_0 eventually_mono not_frequently)

from zorder_exist[OF iso ness this,folded df_def]

obtain r **where** r>0 **and** df_holo:df holomorphic_on cball 0 r **and** df 0 \neq

0

$\forall w \in \text{cball } 0 \text{ } r - \{0\}$.

f (z + w) = df w * w powi (zorder (λw. f (z + w)) 0) \wedge

df w \neq 0

by auto

then have df_nz: $\forall w \in \text{ball } 0 \text{ } r$. df w \neq 0 **by** auto


```

have (deriv  $\widehat{\sim}$  n) df 0 = (deriv  $\widehat{\sim}$  n) (df  $\circ$  g) z for n
unfolding comp_def g_def
proof (subst higher_deriv_compose_linear'[where u=1 and c=-z,simplified])
show df holomorphic_on ball 0 r
using df_holo by auto
show open (ball z r) open (ball 0 r) z  $\in$  ball z r
using <r>0 by auto
show  $\bigwedge w. w \in \text{ball } z \ r \implies w - z \in \text{ball } 0 \ r$ 
by (simp add: dist_norm)
qed auto
then show ?thesis
unfolding fps_expansion_def by auto
qed
also have ... = fps_expansion (zor_poly f z) z
proof (rule fps_expansion_cong)
have  $\forall_F w \text{ in nhds } z. \text{zor\_poly } f \ z \ w$ 
= zor_poly ( $\lambda u. f (z + u)$ ) 0 (w - z)
apply (rule zor_poly_shift)
using True assms by auto
then show  $\forall_F w \text{ in nhds } z. (df \circ g) \ w = \text{zor\_poly } f \ z \ w$ 
unfolding df_def g_def comp_def
by (auto elim:eventually_mono)
qed
finally show ?thesis unfolding df_def
by (auto simp: laurent_expansion_def at_to_0[of z]
eventually_filtermap_add_ac zorder_shift')
qed
finally show ?thesis .
qed

lemma has_fps_expansion_to_laurent:
f has_fps_expansion F  $\longleftrightarrow$  f has_laurent_expansion fps_to_fls F  $\wedge$  f 0 =
fps_nth F 0
proof safe
assume *: f has_laurent_expansion fps_to_fls F f 0 = fps_nth F 0
have eventually ( $\lambda z. z \in \text{eball } 0 \ (\text{fps\_conv\_radius } F)$ ) (nhds 0)
using * by (intro eventually_nhds_in_open) (auto simp: has_laurent_expansion_def
zero_ereal_def)
moreover have eventually ( $\lambda z. z \neq 0 \longrightarrow \text{eval\_fls } (\text{fps\_to\_fls } F) \ z = f \ z$ ) (nhds
0)
using * by (auto simp: has_laurent_expansion_def eventually_at_filter)
ultimately have eventually ( $\lambda z. f \ z = \text{eval\_fps } F \ z$ ) (nhds 0)
by eventually_elim
(auto simp: has_laurent_expansion_def eventually_at_filter eval_fps_at_0
eval_fps_to_fls *(2))
thus f has_fps_expansion F
using * by (auto simp: has_fps_expansion_def has_laurent_expansion_def
eq_commute)
next

```

```

assume f has_fps_expansion F
thus f 0 = fps_nth F 0
  by (metis eval_fps_at_0 has_fps_expansion_imp_holomorphic)
qed (auto intro: has_laurent_expansion_fps)

```

```

lemma eval_fps_fls_base_factor [simp]:
  assumes z ≠ 0
  shows eval_fps (fls_base_factor_to_fps F) z = eval_fls F z * z powi -fls_subdegree
  F
  using assms unfolding eval_fls_def by (simp add: power_int_minus_field_simps)

```

```

lemma has_fps_expansion_imp_analytic_0:
  assumes f has_fps_expansion F
  shows f analytic_on {0}
  by (meson analytic_at_two assms has_fps_expansion_imp_holomorphic)

```

```

lemma has_fps_expansion_imp_analytic:
  assumes (λx. f (z + x)) has_fps_expansion F
  shows f analytic_on {z}
proof -
  have (λx. f (z + x)) analytic_on {0}
    by (rule has_fps_expansion_imp_analytic_0) fact
  hence (λx. f (z + x)) ∘ (λx. x - z) analytic_on {z}
    by (intro analytic_on_compose_gen analytic_intros) auto
  thus ?thesis
    by (simp add: o_def)
qed

```

```

lemma is_pole_cong_asymp_equiv:
  assumes f ~[at z] g z = z'
  shows is_pole f z = is_pole g z'
  using asymp_equiv_at_infinity_transfer[OF assms(1)]
    asymp_equiv_at_infinity_transfer[OF asymp_equiv_symI[OF assms(1)]]
  assms(2)
  unfolding is_pole_def by auto

```

```

lemma not_is_pole_const [simp]: ¬is_pole (λ_::'a::perfect_space. c :: complex)
  z
  using not_tendsto_and_filterlim_at_infinity[of at z λ_::'a. c c] by (auto simp:
is_pole_def)

```

```

lemma has_laurent_expansion_imp_is_pole_iff:
  assumes F: (λx. f (z + x)) has_laurent_expansion F
  shows is_pole f z ↔ fls_subdegree F < 0
proof
  assume pole: is_pole f z
  have [simp]: F ≠ 0
  proof
    assume F = 0

```

```

hence is_pole f z  $\longleftrightarrow$  is_pole ( $\lambda \_. 0 :: \text{complex}$ ) z using assms
  by (intro is_pole_cong)
    (auto simp: has_laurent_expansion_def at_to_0[of z] eventually_filtermap
add_ac)
  with pole show False
    by simp
qed

note pole
also have is_pole f z  $\longleftrightarrow$ 
  is_pole ( $\lambda w. \text{fls\_nth } F (\text{fls\_subdegree } F) * (w - z) \text{ powi fls\_subdegree}$ 
F) z
  using has_laurent_expansion_imp_asymp_equiv[OF F] by (intro is_pole_cong_asymp_equiv
refl)
also have  $\dots \longleftrightarrow$  is_pole ( $\lambda w. (w - z) \text{ powi fls\_subdegree } F$ ) z
  by simp
finally have pole':  $\dots$  .

have False if fls_subdegree F  $\geq 0$ 
proof -
  have ( $\lambda w. (w - z) \text{ powi fls\_subdegree } F$ ) holomorphic_on UNIV
    using that by (intro holomorphic_intros) auto
  hence  $\neg \text{is\_pole } (\lambda w. (w - z) \text{ powi fls\_subdegree } F)$  z
    by (meson UNIV_I not_is_pole_holomorphic_open_UNIV)
  with pole' show False
    by simp
qed
thus fls_subdegree F  $< 0$ 
  by force
qed (use has_laurent_expansion_imp_is_pole[OF assms] in auto)

lemma analytic_at_imp_has_fps_expansion_0:
  assumes f analytic_on {0}
  shows f has_fps_expansion fps_expansion f 0
  using assms has_fps_expansion_fps_expansion analytic_at by fast

lemma analytic_at_imp_has_fps_expansion:
  assumes f analytic_on {z}
  shows ( $\lambda x. f (z + x)$ ) has_fps_expansion fps_expansion f z
proof -
  have f  $\circ (\lambda x. z + x)$  analytic_on {0}
    by (intro analytic_on_compose_gen[OF assms] analytic_intros) auto
  hence (f  $\circ (\lambda x. z + x)$ ) has_fps_expansion fps_expansion (f  $\circ (\lambda x. z + x)$ ) 0
    unfolding o_def by (intro analytic_at_imp_has_fps_expansion_0) auto
  also have  $\dots = \text{fps\_expansion } f$  z
    by (simp add: fps_expansion_def higher_deriv_shift_0')
  finally show ?thesis by (simp add: add_ac)
qed

```

```

lemma has_laurent_expansion_zorder_0:
  assumes f has_laurent_expansion F F ≠ 0
  shows zorder f 0 = fls_subdegree F
proof -
  define G where G = fls_base_factor_to_fps F
  from assms obtain A where A: 0 ∈ A open A ∧ x. x ∈ A - {0} ⇒ eval_fl
F x = f x
  unfolding has_laurent_expansion_def eventually_at_filter eventually_nhds
  by blast

  show ?thesis
  proof (rule zorder_eqI)
  show open (A ∩ eball 0 (fls_conv_radius F)) 0 ∈ A ∩ eball 0 (fls_conv_radius
F)
    using assms A by (auto simp: has_laurent_expansion_def zero_ereal_def)
  show eval_fps G holomorphic_on A ∩ eball 0 (fls_conv_radius F)
    by (intro holomorphic_intros) (auto simp: fls_conv_radius_altdef G_def)
  show eval_fps G 0 ≠ 0 using ⟨F ≠ 0⟩
    by (auto simp: eval_fps_at_0 G_def)
  next
  fix w :: complex assume w ∈ A ∩ eball 0 (fls_conv_radius F) w ≠ 0
  thus f w = eval_fps G w * (w - 0) powi (fls_subdegree F)
    using A unfolding G_def
    by (subst eval_fps_fls_base_factor)
    (auto simp: complex_powr_of_int power_int_minus field_simps)
  qed
qed

lemma has_laurent_expansion_zorder:
  assumes (λw. f (z + w)) has_laurent_expansion F F ≠ 0
  shows zorder f z = fls_subdegree F
  using has_laurent_expansion_zorder_0[OF assms] by (simp add: zorder_shift'
add_ac)

lemma has_fps_expansion_zorder_0:
  assumes f has_fps_expansion F F ≠ 0
  shows zorder f 0 = int (subdegree F)
  using assms has_laurent_expansion_zorder_0[of f fps_to_fl F]
  by (auto simp: has_fps_expansion_to_laurent fls_subdegree_fls_to_fps)

lemma has_fps_expansion_zorder:
  assumes (λw. f (z + w)) has_fps_expansion F F ≠ 0
  shows zorder f z = int (subdegree F)
  using has_fps_expansion_zorder_0[OF assms]
  by (simp add: zorder_shift' add_ac)

lemma has_fps_expansion_fls_base_factor_to_fps:
  assumes f has_laurent_expansion F
  defines n ≡ fls_subdegree F

```

defines $c \equiv \text{fps_nth } (\text{fls_base_factor_to_fps } F) 0$
shows $(\lambda z. \text{if } z = 0 \text{ then } c \text{ else } f z * z^{\text{powi } -n}) \text{ has_fps_expansion } \text{fls_base_factor_to_fps } F$
proof –
have $(\lambda z. f z * z^{\text{powi } -n}) \text{ has_laurent_expansion } \text{fls_shift } (-(-n)) F$
by $(\text{intro laurent_expansion_intros } \text{assms})$
also have $\text{fls_shift } (-(-n)) F = \text{fps_to_fls } (\text{fls_base_factor_to_fps } F)$
by $(\text{simp add: } n_def \text{ fls_shift_nonneg_subdegree})$
also have $(\lambda z. f z * z^{\text{powi } -n}) \text{ has_laurent_expansion } \text{fps_to_fls } (\text{fls_base_factor_to_fps } F) \iff$
 $(\lambda z. \text{if } z = 0 \text{ then } c \text{ else } f z * z^{\text{powi } -n}) \text{ has_laurent_expansion } \text{fps_to_fls } (\text{fls_base_factor_to_fps } F)$
by $(\text{intro has_laurent_expansion_cong}) (\text{auto simp: eventually_at_filter})$
also have $\dots \iff (\lambda z. \text{if } z = 0 \text{ then } c \text{ else } f z * z^{\text{powi } -n}) \text{ has_fps_expansion } \text{fls_base_factor_to_fps } F$
by $(\text{subst has_fps_expansion_to_laurent}) (\text{auto simp: } c_def)$
finally show $?thesis$.
qed

lemma zero_has_laurent_expansion_imp_eq_0:
assumes $(\lambda_. 0) \text{ has_laurent_expansion } F$
shows $F = 0$
proof –
have $at (0 :: \text{complex}) \neq \text{bot}$
by auto
moreover have $(\lambda z. \text{if } z = 0 \text{ then } \text{fls_nth } F (\text{fls_subdegree } F) \text{ else } 0) \text{ has_fps_expansion } \text{fls_base_factor_to_fps } F$ **(is** $?f \text{ has_fps_expansion } _)$
using $\text{has_fps_expansion_fls_base_factor_to_fps}[OF \text{assms}]$ **by** $(\text{simp cong: if_cong})$
hence $\text{isCont } ?f 0$
using $\text{has_fps_expansion_imp_continuous}$ **by** blast
hence $?f -0 \rightarrow \text{fls_nth } F (\text{fls_subdegree } F)$
by $(\text{auto simp: isCont_def})$
moreover have $?f -0 \rightarrow 0 \iff (\lambda_::\text{complex. } 0 :: \text{complex}) -0 \rightarrow 0$
by $(\text{intro filterlim_cong}) (\text{auto simp: eventually_at_filter})$
hence $?f -0 \rightarrow 0$
by simp
ultimately have $\text{fls_nth } F (\text{fls_subdegree } F) = 0$
by $(\text{rule tendsto_unique})$
thus $?thesis$
by $(\text{meson nth_fls_subdegree_nonzero})$
qed

lemma has_laurent_expansion_unique:
assumes $f \text{ has_laurent_expansion } F$ $f \text{ has_laurent_expansion } G$
shows $F = G$
proof –
from assms **have** $(\lambda x. f x - f x) \text{ has_laurent_expansion } F - G$
by $(\text{intro laurent_expansion_intros})$

```

hence  $(\lambda_. 0)$  has_laurent_expansion  $F - G$ 
  by simp
hence  $F - G = 0$ 
  by (rule zero_has_laurent_expansion_imp_eq_0)
thus ?thesis
  by simp
qed

```

```

lemma laurent_expansion_eqI:
  assumes  $(\lambda x. f (z + x))$  has_laurent_expansion  $F$ 
  shows laurent_expansion  $f z = F$ 
  using assms has_laurent_expansion_isolated has_laurent_expansion_not_essential
    has_laurent_expansion_unique not_essential_has_laurent_expansion by
blast

```

```

lemma laurent_expansion_0_eqI:
  assumes  $f$  has_laurent_expansion  $F$ 
  shows laurent_expansion  $f 0 = F$ 
  using assms laurent_expansion_eqI[of f 0] by simp

```

```

lemma has_laurent_expansion_nonzero_imp_eventually_nonzero:
  assumes  $f$  has_laurent_expansion  $F$   $F \neq 0$ 
  shows eventually  $(\lambda x. f x \neq 0)$  (at 0)
proof (rule ccontr)
  assume  $\neg$ eventually  $(\lambda x. f x \neq 0)$  (at 0)
  with assms have eventually  $(\lambda x. f x = 0)$  (at 0)
  by (intro not_essential_frequently_0_imp_eventually_0 has_laurent_expansion_isolated
    has_laurent_expansion_not_essential)
    (auto simp: frequently_def)
  hence  $(f \text{ has\_laurent\_expansion } 0) \longleftrightarrow ((\lambda_. 0) \text{ has\_laurent\_expansion } 0)$ 
  by (intro has_laurent_expansion_cong) auto
  hence  $f$  has_laurent_expansion 0
  by simp
  with assms(1) have  $F = 0$ 
  using has_laurent_expansion_unique by blast
  with  $\langle F \neq 0 \rangle$  show False
  by contradiction
qed

```

```

lemma has_laurent_expansion_eventually_nonzero_iff':
  assumes  $f$  has_laurent_expansion  $F$ 
  shows eventually  $(\lambda x. f x \neq 0)$  (at 0)  $\longleftrightarrow F \neq 0$ 
proof
  assume  $\forall_F x$  in at 0.  $f x \neq 0$ 
  moreover have  $\neg (\forall_F x$  in at 0.  $f x \neq 0)$  if  $F=0$ 
  proof  $-$ 
    have  $\forall_F x$  in at 0.  $f x = 0$ 
    using assms that unfolding has_laurent_expansion_def by simp
    then show ?thesis unfolding not_eventually

```

```

    by (auto elim:eventually_frequentlyE)
  qed
  ultimately show  $F \neq 0$  by auto
qed (simp add:has_laurent_expansion_nonzero_imp_eventually_nonzero[OF assms])

lemma has_laurent_expansion_eventually_nonzero_iff:
  assumes  $(\lambda w. f (z+w))$  has_laurent_expansion  $F$ 
  shows eventually  $(\lambda x. f x \neq 0)$  (at  $z$ )  $\longleftrightarrow F \neq 0$ 
  apply (subst eventually_at_to_0)
  apply (rule has_laurent_expansion_eventually_nonzero_iff')
  using assms by (simp add:add commute)

lemma has_laurent_expansion_inverse [laurent_expansion_intros]:
  assumes  $f$  has_laurent_expansion  $F$ 
  shows  $(\lambda x. \text{inverse } (f x))$  has_laurent_expansion  $\text{inverse } F$ 
proof (cases  $F = 0$ )
  case True
  thus ?thesis using assms
    by (auto simp: has_laurent_expansion_def)
next
  case False
  define  $G$  where  $G = \text{laurent\_expansion } (\lambda x. \text{inverse } (f x)) 0$ 
  from False have  $ev$ : eventually  $(\lambda z. f z \neq 0)$  (at 0)
  by (intro has_laurent_expansion_nonzero_imp_eventually_nonzero[OF assms])

  have *:  $(\lambda x. \text{inverse } (f x))$  has_laurent_expansion  $G$  unfolding  $G$ _def
  by (intro not_essential_has_laurent_expansion_0 isolated_singularity_at_inverse
not_essential_inverse
  has_laurent_expansion_isolated_0[OF assms] has_laurent_expansion_not_essential_0[OF
assms])
  have  $(\lambda x. f x * \text{inverse } (f x))$  has_laurent_expansion  $F * G$ 
  by (intro laurent_expansion_intros assms *)
  also have ?this  $\longleftrightarrow (\lambda x. 1)$  has_laurent_expansion  $F * G$ 
  by (intro has_laurent_expansion_cong_refl eventually_mono[OF ev]) auto
  finally have  $(\lambda x. 1)$  has_laurent_expansion  $F * G$  .
  moreover have  $(\lambda x. 1)$  has_laurent_expansion 1
  by simp
  ultimately have  $F * G = 1$ 
  using has_laurent_expansion_unique by blast
  hence  $G = \text{inverse } F$ 
  using inverse_unique by blast
  with * show ?thesis
  by simp
qed

lemma has_laurent_expansion_power_int [laurent_expansion_intros]:
   $f$  has_laurent_expansion  $F \implies (\lambda x. f x \text{ powi } n)$  has_laurent_expansion  $(F \text{ powi } n)$ 
  by (auto simp: power_int_def intro!: laurent_expansion_intros)

```

```

lemma has_fps_expansion_0_analytic_continuation:
  assumes f has_fps_expansion 0 f holomorphic_on A
  assumes open A connected A 0 ∈ A x ∈ A
  shows  $f x = 0$ 
proof –
  have eventually ( $\lambda z. z \in A \wedge f z = 0$ ) (nhds 0) using assms
  by (intro eventually_conj eventually_nhds_in_open) (auto simp: has_fps_expansion_def)
  then obtain B where B: open B 0 ∈ B  $\forall z \in B. z \in A \wedge f z = 0$ 
  unfolding eventually_nhds by blast
  show ?thesis
  proof (rule analytic_continuation_open[where f = f and g =  $\lambda_. 0$ ])
  show  $B \neq \{\}$ 
  using  $\langle \text{open } B \rangle B$  by auto
  show connected A
  using assms by auto
  qed (use assms B in auto)
qed

```

```

lemma has_laurent_expansion_0_analytic_continuation:
  assumes f has_laurent_expansion 0 f holomorphic_on A - {0}
  assumes open A connected A 0 ∈ A x ∈ A - {0}
  shows  $f x = 0$ 
proof –
  have eventually ( $\lambda z. z \in A - \{0\} \wedge f z = 0$ ) (at 0) using assms
  by (intro eventually_conj eventually_at_in_open) (auto simp: has_laurent_expansion_def)
  then obtain B where B: open B 0 ∈ B  $\forall z \in B - \{0\}. z \in A - \{0\} \wedge f z = 0$ 
  unfolding eventually_at_filter eventually_nhds by blast
  show ?thesis
  proof (rule analytic_continuation_open[where f = f and g =  $\lambda_. 0$ ])
  show  $B - \{0\} \neq \{\}$ 
  using  $\langle \text{open } B \rangle \langle 0 \in B \rangle$  by (metis insert_Diff not_open_singleton)
  show connected (A - {0})
  using assms by (intro connected_open_delete) auto
  qed (use assms B in auto)
qed

```

```

lemma has_fps_expansion_cong:
  assumes eventually ( $\lambda x. f x = g x$ ) (nhds 0)  $F = G$ 
  shows f has_fps_expansion F  $\longleftrightarrow$  g has_fps_expansion G
  using assms(2) by (auto simp: has_fps_expansion_def elim!: eventually_elim2[OF assms(1)])

```

```

lemma zor_poly_has_fps_expansion:
  assumes f has_laurent_expansion F F  $\neq 0$ 
  shows zor_poly f 0 has_fps_expansion fls_base_factor_to_fps F
proof –
  note [simp] =  $\langle F \neq 0 \rangle$ 

```



```

have eventually ( $\lambda z. f z \neq 0$ ) (at 0)
by (rule has_laurent_expansion_nonzero_imp_eventually_nonzero[OF assms])
hence freq: frequently ( $\lambda z. f z \neq 0$ ) (at 0)
by (rule eventually_frequently[rotated]) auto

have *: isolated_singularity_at f 0 not_essential f 0
using has_laurent_expansion_isolated_0[OF assms(1)] has_laurent_expansion_not_essential_0[OF
assms(1)]
by auto

define G where G = fls_base_factor_to_fps F
define n where n = zorder f 0
have n_altdef: n = fls_subdegree F
using has_laurent_expansion_zorder_0 [OF assms(1)] by (simp add: n_def)
obtain r where r: zor_poly f 0 0  $\neq 0$  zor_poly f 0 holomorphic_on cball 0 r r
> 0

$$\forall w \in \text{cball } 0 \ r - \{0\}. f \ w = \text{zor\_poly } f \ 0 \ w * w^{\text{powi } n} \wedge$$


$$\text{zor\_poly } f \ 0 \ w \neq 0$$

using zorder_exist[OF * freq] unfolding n_def by auto
obtain r' where r': r' > 0  $\forall x \in \text{ball } 0 \ r' - \{0\}. \text{eval\_fls } F \ x = f \ x$ 
using assms(1) unfolding has_laurent_expansion_def eventually_at_filter
eventually_nhds_metric ball_def
by (auto simp: dist_commute)
have holo: zor_poly f 0 holomorphic_on ball 0 r
by (rule holomorphic_on_subset[OF r(2)]) auto

have ( $\lambda z. \text{if } z = 0 \text{ then fps\_nth } G \ 0 \text{ else } f \ z * z^{\text{powi } -n}$ ) has_fps_expansion G
unfolding G_def n_altdef by (intro has_fps_expansion_fls_base_factor_to_fps
assms)
also have ?this  $\longleftrightarrow$  zor_poly f 0 has_fps_expansion G
proof (intro has_fps_expansion_cong)
have eventually ( $\lambda z. z \in \text{ball } 0 \ (\min \ r \ r')$ ) (nhds 0)
using  $\langle r > 0 \rangle \langle r' > 0 \rangle$  by (intro eventually_nhds_in_open) auto
thus  $\forall_F \ x \text{ in nhds } 0. (\text{if } x = 0 \text{ then } G \ \$ \ 0 \text{ else } f \ x * x^{\text{powi } -n}) = \text{zor\_poly } f$ 
0 x
proof eventually_elim
case (elim w)
have w: w  $\in \text{ball } 0 \ r \ w \in \text{ball } 0 \ r'$ 
using elim by auto
show ?case
proof (cases w = 0)
case False
hence f w = zor_poly f 0 w * wpowi n
using r w by auto
thus ?thesis using False
by (simp add: powr_minus complex_powr_of_int power_int_minus)
next
case [simp]: True
obtain R where R: R > 0 R  $\leq r$  R  $\leq r'$  R  $\leq \text{fls\_conv\_radius } F$ 

```

```

    using ⟨r > 0⟩ ⟨r' > 0⟩ assms(1) unfolding has_laurent_expansion_def
    by (smt (verit, ccfv_SIG) ereal_dense2 ereal_less(2) less_ereal.simps(1)
order.strict_implies_order order_trans)
  have eval_fps G 0 = zor_poly f 0 0
  proof (rule analytic_continuation_open[where f = eval_fps G and g =
zor_poly f 0])
    show connected (ball 0 R :: complex set)
      by auto
    have of_real R / 2 ∈ ball 0 R - {0 :: complex}
      using R by auto
    thus ball 0 R - {0 :: complex} ≠ {}
      by blast
    show eval_fps G holomorphic_on ball 0 R
      using R less_le_trans[OF _ R(4)] unfolding G_def
      by (intro holomorphic_intros) (auto simp: fls_conv_radius_altdef)
    show zor_poly f 0 holomorphic_on ball 0 R
      by (rule holomorphic_on_subset[OF holo]) (use R in auto)
    show eval_fps G z = zor_poly f 0 z if z ∈ ball 0 R - {0} for z
      using that r r' R n_altdef unfolding G_def
      by (subst eval_fps_fls_base_factor)
        (auto simp: complex_pow_of_int field_simps power_int_minus
n_def)
    qed (use R in auto)
    hence zor_poly f 0 0 = fps_nth G 0
      by (simp add: eval_fps_at_0)
    thus ?thesis by simp
  qed
qed
qed
qed (use r' in auto)
finally show ?thesis
  by (simp add: G_def)
qed

lemma zorder_geI_0:
  assumes f_analytic_on {0} f_holomorphic_on A open A connected A 0 ∈ A z ∈
A f z ≠ 0
  assumes ∧k. k < n ⇒ (deriv ^^ k) f 0 = 0
  shows zorder f 0 ≥ n
proof -
  define F where F = fps_expansion f 0
  from assms have f_has_fps_expansion F
  unfolding F_def using analytic_at_imp_has_fps_expansion_0 by blast
  hence laurent: f_has_laurent_expansion fps_to_fls F and [simp]: f 0 = fps_nth
F 0
  by (simp_all add: has_fps_expansion_to_laurent)

  have [simp]: F ≠ 0
  proof
    assume [simp]: F = 0

```

```

hence  $f z = 0$ 
proof (cases  $z = 0$ )
  case False
    have  $f$  has_laurent_expansion 0
      using laurent by simp
    thus ?thesis
  proof (rule has_laurent_expansion_0_analytic_continuation)
    show  $f$  holomorphic_on  $A - \{0\}$ 
      using assms(2) by (rule holomorphic_on_subset) auto
    qed (use assms False in auto)
  qed auto
with  $\langle f z \neq 0 \rangle$  show False by contradiction
qed

have  $\text{zorder } f 0 = \text{int } (\text{subdegree } F)$ 
  using has_laurent_expansion_zorder_0[OF laurent] by (simp add: fls_subdegree_fls_to_fps)
also have  $\text{subdegree } F \geq n$ 
  using assms by (intro subdegree_geI  $\langle F \neq 0 \rangle$ ) (auto simp: F_def fps_expansion_def)
hence  $\text{int } (\text{subdegree } F) \geq \text{int } n$ 
  by simp
finally show ?thesis .
qed

lemma zorder_geI:
  assumes  $f$  analytic_on  $\{x\}$   $f$  holomorphic_on  $A$  open  $A$  connected  $A$   $x \in A$   $z \in A$ 
   $f z \neq 0$ 
  assumes  $\bigwedge k. k < n \implies (\text{deriv } \wedge k) f x = 0$ 
  shows  $\text{zorder } f x \geq n$ 
proof -
  have  $\text{zorder } f x = \text{zorder } (f \circ (\lambda u. u + x)) 0$ 
    by (subst zorder_shift) (auto simp: o_def)
  also have  $\dots \geq n$ 
  proof (rule zorder_geI_0)
    show  $f \circ (\lambda u. u + x)$  analytic_on  $\{0\}$ 
      by (intro analytic_on_compose_gen[OF assms(1)] analytic_intros) auto
    show  $f \circ (\lambda u. u + x)$  holomorphic_on  $((+) (-x)) 'A$ 
      by (intro holomorphic_on_compose_gen[OF assms(2)] holomorphic_intros)
  auto
  show connected  $((+) (-x)) 'A$ 
    by (intro connected_continuous_image continuous_intros assms)
  show open  $((+) (-x)) 'A$ 
    by (intro open_translation assms)
  show  $z - x \in (+) (-x)) 'A$ 
    using  $\langle z \in A \rangle$  by auto
  show  $0 \in (+) (-x)) 'A$ 
    using  $\langle x \in A \rangle$  by auto
  show  $f \circ (\lambda u. u + x)$   $(z - x) \neq 0$ 
    using  $\langle f z \neq 0 \rangle$  by auto
next

```

```

fix  $k :: \text{nat}$  assume  $k < n$ 
hence  $(\text{deriv } \overset{\sim}{\sim} k) f x = 0$ 
  using assms by simp
also have  $(\text{deriv } \overset{\sim}{\sim} k) f x = (\text{deriv } \overset{\sim}{\sim} k) (f \circ (+) x) 0$ 
  by (subst higher_deriv_shift_0) auto
finally show  $(\text{deriv } \overset{\sim}{\sim} k) (f \circ (\lambda u. u + x)) 0 = 0$ 
  by (subst add.commute) auto
qed
finally show ?thesis .
qed

```

```

lemma has_laurent_expansion_divide [laurent_expansion_intros]:
  assumes  $f$  has_laurent_expansion  $F$  and  $g$  has_laurent_expansion  $G$ 
  shows  $(\lambda x. f x / g x)$  has_laurent_expansion  $(F / G)$ 
proof -
  have  $(\lambda x. f x * \text{inverse } (g x))$  has_laurent_expansion  $(F * \text{inverse } G)$ 
    by (intro laurent_expansion_intros assms)
  thus ?thesis
    by (simp add: field_simps)
qed

```

```

lemma has_laurent_expansion_residue_0:
  assumes  $f$  has_laurent_expansion  $F$ 
  shows residue  $f$   $0 = \text{fls\_residue } F$ 
proof (cases fls_subdegree F ≥ 0)
  case True
  have residue  $f$   $0 = \text{residue } (\text{eval\_fls } F) 0$ 
    using assms by (intro residue_cong) (auto simp: has_laurent_expansion_def eq_commute)
  also have  $\dots = 0$ 
    by (rule residue_holo[OF  $\_\_\_ \text{holomorphic\_on\_eval\_fls}$ [OF order.refl]])
    (use True assms in <auto simp: has_laurent_expansion_def zero_ereal_def>)
  also have  $\dots = \text{fls\_residue } F$ 
    using True by simp
  finally show ?thesis .
next
  case False
  hence  $F \neq 0$ 
    by auto
  have  $*$ : zor_poly  $f$   $0$  has_fps_expansion fls_base_factor_to_fps  $F$ 
    by (intro zor_poly_has_fps_expansion False assms <F ≠ 0>)

  have residue  $f$   $0 = (\text{deriv } \overset{\sim}{\sim} (\text{nat } (-\text{zorder } f 0) - 1)) (\text{zor\_poly } f 0) 0 / \text{fact}$ 
     $(\text{nat } (-\text{zorder } f 0) - 1)$ 
    by (intro residue_pole_order has_laurent_expansion_isolated_0[OF assms]
      has_laurent_expansion_imp_is_pole_0[OF assms]) (use False in auto)
  also have  $\dots = \text{fls\_residue } F$ 
    using has_laurent_expansion_zorder_0[OF assms <F ≠ 0>] False
    by (subst fps_nth_fps_expansion [OF *, symmetric]) (auto simp: of_nat_diff)

```

finally show ?thesis .
qed

lemma has_laurent_expansion_residue:
assumes $(\lambda x. f(z + x))$ has_laurent_expansion F
shows residue $f z =$ fls_residue F
using has_laurent_expansion_residue_0[OF assms] by (simp add: residue_shift_0')

lemma eval_fls_has_laurent_expansion [laurent_expansion_intros]:
assumes fls_conv_radius $F > 0$
shows eval_fls F has_laurent_expansion F
using assms by (auto simp: has_laurent_expansion_def)

lemma fps_expansion_unique_complex:
fixes $F G ::$ complex fps
assumes f has_fps_expansion F f has_fps_expansion G
shows $F = G$
using assms unfolding fps_eq_iff by (auto simp: fps_eq_iff fps_nth_fps_expansion)

lemma fps_expansion_eqI:
assumes f has_fps_expansion F
shows fps_expansion $f 0 = F$
using assms unfolding fps_eq_iff
by (auto simp: fps_eq_iff fps_nth_fps_expansion fps_expansion_def)

lemma holomorphic_on_imp_fps_conv_radius_ge:
assumes f has_fps_expansion F f holomorphic_on eball $0 r$
shows fps_conv_radius $F \geq r$
proof -
define n where $n =$ subdegree F
have fps_conv_radius (fps_expansion $f 0) \geq r$
by (intro conv_radius_fps_expansion assms)
also have fps_expansion $f 0 = F$
using assms by (intro fps_expansion_eqI)
finally show ?thesis
by simp
qed

lemma has_fps_expansion_imp_eval_fps_eq:
assumes f has_fps_expansion F norm $z < r$
assumes f holomorphic_on ball $0 r$
shows eval_fps $F z = f z$
proof -
have [simp]: fps_expansion $f 0 = F$
by (rule fps_expansion_eqI) fact
have *: f holomorphic_on eball 0 (ereal r)
using assms by simp
from conv_radius_fps_expansion[OF *] have fps_conv_radius $F \geq$ ereal r
by simp

```

have eval_fps (fps_expansion f 0) z = f (0 + z)
  by (rule eval_fps_expansion'[OF *]) (use assms in auto)
thus ?thesis
  by simp
qed

```

lemma has_fps_expansion_imp_sums_complex:

```

fixes F :: complex fps
assumes f has_fps_expansion F f holomorphic_on eball 0 r ereal (norm z) < r
shows (λn. fps_nth F n * z ^ n) sums f z
proof -
  have r: fps_conv_radius F ≥ r
    using assms(1,2) by (rule holomorphic_on_imp_fps_conv_radius_ge)
  from assms obtain R where R: norm z < R ereal R < r
    using ereal_dense2 less_ereal.simps(1) by blast
  have z: norm z < fps_conv_radius F
    using r R assms(3) by order

  have summable (λn. fps_nth F n * z ^ n)
    by (rule summable_fps) (use z in auto)
  moreover have eval_fps F z = f z
  proof (rule has_fps_expansion_imp_eval_fps_eq[where r = R])
    have *: ereal (norm z) < r if norm z < R for z :: complex
      using that R ereal_le_less less_imp_le by blast
    show f holomorphic_on ball 0 R
      using assms(2) by (rule holomorphic_on_subset) (use * in auto)
  qed (use R assms(1) in auto)
  ultimately show ?thesis
  unfolding eval_fps_def sums_iff by simp
qed

```

lemma fls_conv_radius_ge:

```

assumes f has_laurent_expansion F
assumes f holomorphic_on eball 0 r - {0}
shows fls_conv_radius F ≥ r
proof -
  define n where n = fls_subdegree F
  define G where G = fls_base_factor_to_fps F
  define g where g = (λz. if z = 0 then fps_nth G 0 else f z * z powi -n)
  have G: g has_fps_expansion G
    unfolding G_def g_def n_def
    by (intro has_fps_expansion_fls_base_factor_to_fps assms)
  have (λz. f z * z powi -n) holomorphic_on eball 0 r - {0}
    by (intro holomorphic_intros assms) auto
  also have ?this ↔ g holomorphic_on eball 0 r - {0}
    by (intro holomorphic_cong) (auto simp: g_def)
  finally have g analytic_on eball 0 r - {0}
    by (subst analytic_on_open) auto
  moreover have g analytic_on {0}

```

```

    using G has_fps_expansion_imp_analytic_0 by auto
  ultimately have g analytic_on (eball 0 r - {0} ∪ {0})
    by (subst analytic_on_Un) auto
  hence g analytic_on eball 0 r
    by (rule analytic_on_subset) auto
  hence g holomorphic_on eball 0 r
    by (subst (asm) analytic_on_open) auto
  hence fps_conv_radius (fps_expansion g 0) ≥ r
    by (intro conv_radius_fps_expansion)
  also have fps_expansion g 0 = G
    using G by (intro fps_expansion_eqI)
  finally show ?thesis
    by (simp add: fls_conv_radius_altdef G_def)
qed

```

lemma eval_fls_eqI:

```

  assumes f has_laurent_expansion F f holomorphic_on eball 0 r - {0}
  assumes z ∈ eball 0 r - {0}
  shows eval_fls F z = f z
proof -
  have conv: fls_conv_radius F ≥ r
    by (intro fls_conv_radius_ge[OF assms(1,2)])
  have (λz. eval_fls F z - f z) has_laurent_expansion F - F
    using assms by (intro laurent_expansion_intros assms) (auto simp: has_laurent_expansion_def)
  hence (λz. eval_fls F z - f z) has_laurent_expansion 0
    by simp
  hence eval_fls F z - f z = 0
  proof (rule has_laurent_expansion_0_analytic_continuation)
    have ereal 0 ≤ ereal (norm z)
      by simp
    also have norm z < r
      using assms by auto
    finally have r > 0
      by (simp add: zero_ereal_def)
    thus open (eball 0 r :: complex set) connected (eball 0 r :: complex set)
      0 ∈ eball 0 r z ∈ eball 0 r - {0}
      using assms by (auto simp: zero_ereal_def)
  qed (auto intro!: holomorphic_intros assms less_le_trans[OF _ conv] split:
if_splits)
  thus ?thesis by simp
qed

```

lemma fls_nth_as_contour_integral:

```

  assumes F: f has_laurent_expansion F
  assumes holo: f holomorphic_on ball 0 r - {0}
  assumes R: 0 < R R < r
  shows ((λz. f z * z powi (-(n+1))) has_contour_integral
    complex_of_real (2 * pi) * i * fls_nth F n) (circlepath 0 R)
proof -

```

```

define I where  $I = (\lambda z. f z * z \text{ powi } (-(n+1)))$ 
have (I has_contour_integral complex_of_real ( $2 * \pi$ ) * i * residue I 0)
(circlepath 0 R)
proof (rule base_residue)
  show open (ball (0::complex) r) 0  $\in$  ball (0::complex) r
  using R F by (auto simp: has_laurent_expansion_def zero_ereal_def)
qed (use R in  $\langle$ auto intro!: holomorphic_intros holomorphic_on_subset[OF holo]
  simp: I_def split: if_splits $\rangle$ )
also have residue I 0 = fls_residue (fls_shift ( $n + 1$ ) F)
unfolding I_def by (intro has_laurent_expansion_residue_0 laurent_expansion_intros
F)
also have  $\dots = \text{fls\_nth } F \ n$ 
by simp
finally show ?thesis
by (simp add: I_def)
qed

```

lemma *tendsto_0_subdegree_iff_0*:

```

assumes F:f has_laurent_expansion F and  $F \neq 0$ 
shows ( $f \rightarrow 0$ )  $\longleftrightarrow$  fls_subdegree F > 0
proof -
  have ?thesis if is_pole f 0
  proof -
    have fls_subdegree F < 0
    using is_pole_0_imp_neg_fls_subdegree[OF F that] .
    moreover then have  $\neg f \rightarrow 0$ 
    using  $\langle$ is_pole f 0 $\rangle$  F at_neg_bot
      has_laurent_expansion_imp_filterlim_infinity_0
      not_tendsto_and_filterlim_at_infinity that
    by blast
    ultimately show ?thesis by auto
  qed
moreover have ?thesis if  $\neg \text{is\_pole } f \ 0 \ \exists x. f \rightarrow x$ 
proof -
  have fls_subdegree F  $\geq 0$ 
  using has_laurent_expansion_imp_is_pole_0[OF F] that(1)
  by linarith
  have  $f \rightarrow 0$  if fls_subdegree F > 0
  using fls_eq0_below_subdegree[OF that]
  by (metis F  $\langle 0 \leq \text{fls\_subdegree } F \rangle$  has_laurent_expansion_imp_tendsto_0)
  moreover have fls_subdegree F > 0 if  $f \rightarrow 0$ 
  proof -
  have False if fls_subdegree F = 0
  proof -
    have  $f \rightarrow 0$  if fls_nth F 0
    using has_laurent_expansion_imp_tendsto_0
      [OF F  $\langle \text{fls\_subdegree } F \geq 0 \rangle$ ] .
    then have fls_nth F 0 = 0 using  $\langle f \rightarrow 0 \rangle$ 
    using LIM_unique by blast
  qed

```



```

    then have  $F = 0$ 
      using nth_fls_subdegree_zero_iff  $\langle \text{fls\_subdegree } F = 0 \rangle$ 
      by metis
    with  $\langle F \neq 0 \rangle$  show False by auto
  qed
  with  $\langle \text{fls\_subdegree } F \geq 0 \rangle$ 
  show ?thesis by fastforce
  qed
  ultimately show ?thesis by auto
  qed
  moreover have  $\text{is\_pole } f \ 0 \vee (\exists x. f \ -0 \rightarrow x)$ 
  proof -
    have not_essential  $f \ 0$ 
      using F has laurent_expansion_not_essential_0 by auto
    then show ?thesis unfolding not_essential_def
      by auto
  qed
  ultimately show ?thesis by auto
  qed

lemma tendsto_0_subdegree_iff:
  assumes  $F: (\lambda w. f \ (z+w))$  has laurent_expansion  $F$  and  $F \neq 0$ 
  shows  $(f \ -z \rightarrow 0) \longleftrightarrow \text{fls\_subdegree } F > 0$ 
  apply (subst Lim_at_zero)
  apply (rule tendsto_0_subdegree_iff_0)
  using assms by auto

lemma is_pole_0_deriv_divide_iff:
  assumes  $F: f$  has laurent_expansion  $F$  and  $F \neq 0$ 
  shows  $\text{is\_pole } (\lambda x. \text{deriv } f \ x / f \ x) \ 0 \longleftrightarrow \text{is\_pole } f \ 0 \vee (f \ -0 \rightarrow 0)$ 
  proof -
    have  $(\lambda x. \text{deriv } f \ x / f \ x)$  has laurent_expansion  $\text{fls\_deriv } F / F$ 
      using  $F$  by (auto intro: laurent_expansion_intros)

    have  $\text{is\_pole } (\lambda x. \text{deriv } f \ x / f \ x) \ 0 \longleftrightarrow$ 
       $\text{fls\_subdegree } (\text{fls\_deriv } F / F) < 0$ 
      apply (rule is_pole_fls_subdegree_iff)
      using  $F$  by (auto intro: laurent_expansion_intros)
    also have  $\dots \longleftrightarrow \text{is\_pole } f \ 0 \vee (f \ -0 \rightarrow 0)$ 
    proof (cases fls_subdegree F = 0)
      case True
      then have  $\text{fls\_subdegree } (\text{fls\_deriv } F / F) \geq 0$ 
        by (metis diff_zero div_0  $\langle F \neq 0 \rangle$  fls_deriv_subdegree0
          fls_divide_subdegree)
      moreover then have  $\neg \text{is\_pole } f \ 0$ 
        by (metis F True is_pole_0_imp_neg_fls_subdegree less_le)
      moreover have  $\neg (f \ -0 \rightarrow 0)$ 
        using tendsto_0_subdegree_iff_0  $[OF \ F \ \langle F \neq 0 \rangle]$  True by auto
      ultimately show ?thesis by auto
    qed
  qed

```

```

next
  case False
  then have fls_deriv F  $\neq$  0
    by (metis fls_const_subdegree fls_deriv_eq_0_iff)
  then have fls_subdegree (fls_deriv F / F) =
    fls_subdegree (fls_deriv F) - fls_subdegree F
    by (rule fls_divide_subdegree[OF  $\langle F \neq 0 \rangle$ ])
  moreover have fls_subdegree (fls_deriv F) = fls_subdegree F - 1
    using fls_subdegree_deriv[OF False] .
  ultimately have fls_subdegree (fls_deriv F / F) < 0 by auto
  moreover have  $f - 0 \rightarrow 0 = (0 < \text{fls\_subdegree } F)$ 
    using tendsto_0_subdegree_iff_0[OF F  $\langle F \neq 0 \rangle$ ] .
  moreover have is_pole f 0 = (fls_subdegree F < 0)
    using is_pole_fls_subdegree_iff F by auto
  ultimately show ?thesis using False by auto
qed
finally show ?thesis .
qed

```

lemma is_pole_deriv_divide_iff:

```

  assumes F:( $\lambda w. f(z+w)$ ) has_laurent_expansion F and  $F \neq 0$ 
  shows is_pole ( $\lambda x. \text{deriv } f x / f x$ ) z  $\longleftrightarrow$  is_pole f z  $\vee$  (f - z  $\rightarrow$  0)
proof -
  define ff df where ff=( $\lambda w. f(z+w)$ ) and df=( $\lambda w. \text{deriv } f(z+w)$ )
  have is_pole ( $\lambda x. \text{deriv } f x / f x$ ) z
     $\longleftrightarrow$  is_pole ( $\lambda x. \text{deriv } ff x / ff x$ ) 0
  unfolding ff_def df_def
  by (simp add:deriv_shift_0' is_pole_shift_0' comp_def algebra_simps)
  moreover have is_pole f z  $\longleftrightarrow$  is_pole ff 0
  unfolding ff_def by (auto simp:is_pole_shift_0')
  moreover have (f - z  $\rightarrow$  0)  $\longleftrightarrow$  (ff - 0  $\rightarrow$  0)
  unfolding ff_def by (simp add:LIM_offset_zero_iff)
  moreover have is_pole ( $\lambda x. \text{deriv } ff x / ff x$ ) 0 = (is_pole ff 0  $\vee$  ff - 0  $\rightarrow$  0)
  apply (rule is_pole_0_deriv_divide_iff)
  using F ff_def  $\langle F \neq 0 \rangle$  by blast+
  ultimately show ?thesis by auto
qed

```

lemma subdegree_imp_eventually_deriv_nzero_0:

```

  assumes F:f has_laurent_expansion F and fls_subdegree F  $\neq$  0
  shows eventually ( $\lambda z. \text{deriv } f z \neq 0$ ) (at 0)
proof -
  have deriv f has_laurent_expansion fls_deriv F
    using has_laurent_expansion_deriv[OF F] .
  moreover have fls_deriv F  $\neq$  0
    using  $\langle \text{fls\_subdegree } F \neq 0 \rangle$ 
  by (metis fls_const_subdegree fls_deriv_eq_0_iff)
  ultimately show ?thesis
    using has_laurent_expansion_eventually_nonzero_iff' by blast

```

qed

lemma *subdegree_imp_eventually_deriv_nzero*:

assumes $F: (\lambda w. f(z+w))$ *has_laurent_expansion* F

and *fls_subdegree* $F \neq 0$

shows *eventually* $(\lambda w. \text{deriv } f \ w \neq 0)$ (at z)

proof –

have $\forall_F x$ in at 0. *deriv* $(\lambda w. f(z+w))$ $x \neq 0$

using *subdegree_imp_eventually_deriv_nzero_0* *assms* **by** *auto*

then show *?thesis*

apply (*subst eventually_at_to_0*)

by (*simp add: deriv_shift_0' comp_def algebra_simps*)

qed

lemma *has_fps_expansion_imp_asymp_equiv_0*:

fixes $f :: \text{complex} \Rightarrow \text{complex}$

assumes $F: f$ *has_fps_expansion* F

defines $n \equiv \text{subdegree } F$

shows $f \sim[\text{nhds } 0]$ $(\lambda z. \text{fps_nth } F \ n * z^n)$

proof –

have F' : f *has_laurent_expansion* *fps_to_fl* F

using F *has_laurent_expansion_fps* **by** *blast*

have $f \sim[\text{at } 0]$ $(\lambda z. \text{fps_nth } F \ n * z^n)$

using *has_laurent_expansion_imp_asymp_equiv_0[OF F']*

by (*simp add: fls_subdegree_fl_to_fps_n_def*)

moreover have $f \ 0 = \text{fps_nth } F \ n * 0^n$

using F **by** (*auto simp: n_def has_fps_expansion_to_laurent_power_0_left*)

ultimately show *?thesis*

by (*auto simp: asymp_equiv_nhds_iff*)

qed

lemma *has_fps_expansion_imp_tendsto_0*:

fixes $f :: \text{complex} \Rightarrow \text{complex}$

assumes f *has_fps_expansion* F

shows $(f \longrightarrow \text{fps_nth } F \ 0)$ (*nhds* 0)

proof (*rule asymp_equiv_tendsto_transfer*)

show $(\lambda z. \text{fps_nth } F \ (\text{subdegree } F) * z^{\text{subdegree } F}) \sim[\text{nhds } 0]$ f

by (*rule asymp_equiv_symI, rule has_fps_expansion_imp_asymp_equiv_0*)

fact

have $(\lambda z. F \ \$ \ \text{subdegree } F * z^{\text{subdegree } F}) \longrightarrow F \ \$ \ 0$ (at 0)

by (*rule tendsto_eq_intros refl | simp*) + (*auto simp: power_0_left*)

thus $(\lambda z. F \ \$ \ \text{subdegree } F * z^{\text{subdegree } F}) \longrightarrow F \ \$ \ 0$ (*nhds* 0)

by (*auto simp: tendsto_nhds_iff power_0_left*)

qed

lemma *has_fps_expansion_imp_0_eq_fps_nth_0*:

assumes f *has_fps_expansion* F

shows $f \ 0 = \text{fps_nth } F \ 0$

proof –

```

have eventually ( $\lambda x. f x = \text{eval\_fps } F x$ ) (nhds 0)
  using assms by (auto simp: has_fps_expansion_def eq_commute)
then obtain A where open A  $0 \in A \forall x \in A. f x = \text{eval\_fps } F x$ 
  unfolding eventually_nhds by blast
hence  $f 0 = \text{eval\_fps } F 0$ 
  by blast
thus ?thesis
  by (simp add: eval_fps_at_0)
qed

```

lemma fls_nth_compose_aux:

```

assumes f has_fps_expansion F
assumes G: g has_fps_expansion G fps_nth G 0 = 0 fps_deriv G  $\neq 0$ 
assumes (f  $\circ$  g) has_laurent_expansion H
shows fls_nth H (int n) = fps_nth (fps_compose F G) n
using assms(1,5)
proof (induction n arbitrary: f F H rule: less_induct)
case (less n f F H)
have [simp]:  $g 0 = 0$ 
  using has_fps_expansion_imp_0_eq_fps_nth_0[OF G(1)] G(2) by simp
have ana_f: f analytic_on {0}
  using less.prem1 by (meson has_fps_expansion_imp_analytic_0)
have ana_g: g analytic_on {0}
  using G by (meson has_fps_expansion_imp_analytic_0)
have (f  $\circ$  g) has_laurent_expansion fps_to_fls (fps_expansion (f  $\circ$  g) 0)
  by (rule analytic_at_imp_has_fps_expansion_0 analytic_intros has_laurent_expansion_fps
    analytic_on_compose_gen ana_f ana_g)+ auto
with less.prem1 have H = fps_to_fls (fps_expansion (f  $\circ$  g) 0)
  using has_laurent_expansion_unique by blast
also have fls_subdegree ...  $\geq 0$ 
  by (simp add: fls_subdegree_fls_to_fps)
finally have subdeg: fls_subdegree H  $\geq 0$  .

```

show ?case

```

proof (cases n = 0)
case [simp]: True
have lim_g:  $g -0 \rightarrow 0$ 
  using has_laurent_expansion_imp_tendsto_0[of g fps_to_fls G] G
  by (auto simp: fls_subdegree_fls_to_fps_gt0 has_fps_expansion_to_laurent)
have lim_f: (f  $\longrightarrow$  fps_nth F 0) (nhds 0)
  by (intro has_fps_expansion_imp_tendsto_0 less.prem1)
have ( $\lambda x. f (g x) -0 \rightarrow \text{fps\_nth } F 0$ )
  by (rule filterlim_compose[OF lim_f lim_g])
moreover have (f  $\circ$  g)  $-0 \rightarrow \text{fls\_nth } H 0$ 
  by (intro has_laurent_expansion_imp_tendsto_0 less.prem1 subdeg)
ultimately have fps_nth F 0 = fls_nth H 0
  using tendsto_unique by (force simp: o_def)
thus ?thesis

```

```

    by simp
  next
    case n: False
    define GH where GH = (fls_deriv H / fls_deriv (fps_to_fls G))
    define GH' where GH' = fls_regpart GH

    have ( $\lambda x. \text{deriv } (f \circ g) x / \text{deriv } g x$ ) has_laurent_expansion
      fls_deriv H / fls_deriv (fps_to_fls G)
    by (intro laurent_expansion_intros less.premis has_laurent_expansion_fps[of
    _ G] G)
    also have ?this  $\longleftrightarrow$  ( $\text{deriv } f \circ g$ ) has_laurent_expansion fls_deriv H /
    fls_deriv (fps_to_fls G)
    proof (rule has_laurent_expansion_cong)
      from ana_f obtain r1 where r1:  $r1 > 0$  f holomorphic_on ball 0 r1
      unfolding analytic_on_def by blast
      from ana_g obtain r2 where r2:  $r2 > 0$  g holomorphic_on ball 0 r2
      unfolding analytic_on_def by blast
      have lim_g:  $g \rightarrow 0$ 
      using has_laurent_expansion_imp_tendsto_0[of g fps_to_fls G] G
      by (auto simp: fls_subdegree_fls_to_fps_gt0 has_fps_expansion_to_laurent)
      moreover have open (ball 0 r1)  $0 \in \text{ball } 0 r1$ 
      using r1 by auto
      ultimately have eventually ( $\lambda x. g x \in \text{ball } 0 r1$ ) (at 0)
      unfolding tendsto_def by blast
      moreover have eventually ( $\lambda x. \text{deriv } g x \neq 0$ ) (at 0)
      using G fps_to_fls_eq_0_iff has_fps_expansion_deriv has_fps_expansion_to_laurent
      has_laurent_expansion_nonzero_imp_eventually_nonzero by blast
      moreover have eventually ( $\lambda x. x \in \text{ball } 0 (\min r1 r2) - \{0\}$ ) (at 0)
      by (intro eventually_at_in_open) (use r1 r2 in auto)
      ultimately show eventually ( $\lambda x. \text{deriv } (f \circ g) x / \text{deriv } g x = (\text{deriv } f \circ g)$ 
    x) (at 0)
      proof eventually_elim
        case (elim x)
        thus ?case using r1 r2
          by (subst deriv_chain)
          (auto simp: field_simps holomorphic_on_def at_within_open[of _ ball
    _])
      qed
    qed auto
    finally have GH: ( $\text{deriv } f \circ g$ ) has_laurent_expansion GH
      unfolding GH_def .

    have ( $\text{deriv } f \circ g$ ) has_laurent_expansion fps_to_fls (fps_expansion (deriv f
     $\circ g$ ) 0)
    by (rule analytic_at_imp_has_fps_expansion_0 analytic_intros has_laurent_expansion_fps
    analytic_on_compose_gen ana_f ana_g)+ auto
    with GH have GH = fps_to_fls (fps_expansion (deriv f  $\circ g$ ) 0)
    using has_laurent_expansion_unique by blast
    also have fls_subdegree ...  $\geq 0$ 

```

```

    by (simp add: fls_subdegree_fls_to_fps)
  finally have subdeg': fls_subdegree GH ≥ 0 .

  have deriv f has_fps_expansion fps_deriv F
    by (intro fps_expansion_intros less.prem)
  from this and GH have IH: fls_nth GH (int k) = fps_nth (fps_compose
(fps_deriv F) G) k
    if k < n for k
  by (intro less.IH that)

  have fps_nth (fps_compose (fps_deriv F) G) n = (∑ i=0..n. of_nat (Suc i)
* F $ Suc i * G ^ i $ n)
    by (simp add: fps_compose_nth)

  have fps_nth (fps_compose F G) n =
    fps_nth (fps_deriv (fps_compose F G)) (n - 1) / of_nat n
    using n by (cases n) (auto simp del: of_nat_Suc)
  also have fps_deriv (fps_compose F G) = fps_compose (fps_deriv F) G *
fps_deriv G
    using G by (subst fps_compose_deriv) auto
  also have fps_nth ... (n - 1) = (∑ i=0..n-1. (fps_deriv F oo G) $ i *
fps_deriv G $ (n - 1 - i))
    unfolding fps_mult_nth ..
  also have ... = (∑ i=0..n-1. fps_nth GH' i * of_nat (n - i) * G $ (n -
i))
    using n by (intro sum.cong) (auto simp: IH Suc_diff_Suc GH'_def)
  also have ... = (∑ i=0..n. fps_nth GH' i * of_nat (n - i) * G $ (n - i))
    by (intro sum.mono_neutral_left) auto
  also have ... = fps_nth (GH' * Abs_fps (λi. of_nat i * fps_nth G i)) n
    by (simp add: fps_mult_nth mult_ac)
  also have Abs_fps (λi. of_nat i * fps_nth G i) = fps_X * fps_deriv G
    by (simp add: fps_mult_fps_X_deriv_shift)
  also have fps_nth (GH' * (fps_X * fps_deriv G)) n =
    fls_nth (fps_to_fls (GH' * (fps_X * fps_deriv G))) (int n)
    by simp
  also have fps_to_fls (GH' * (fps_X * fps_deriv G)) =
    GH * fps_to_fls (fps_deriv G) * fls_X
    using subdeg' by (simp add: mult_ac fls_times_fps_to_fls GH'_def)
  also have GH * fps_to_fls (fps_deriv G) = fls_deriv H
    unfolding GH_def using G by (simp add: fls_deriv_fps_to_fls)
  also have fls_deriv H * fls_X = fls_shift (-1) (fls_deriv H)
    using fls_X_times_conv_shift(2) by blast
  finally show ?thesis
    using n by simp
qed
qed

```

```

lemma has_fps_expansion_compose [fps_expansion_intros]:
  fixes f g :: complex ⇒ complex

```

```

assumes  $F$ :  $f$  has_fps_expansion  $F$ 
assumes  $G$ :  $g$  has_fps_expansion  $G$  fps_nth  $G$   $0 = 0$ 
shows  $(f \circ g)$  has_fps_expansion fps_compose  $F$   $G$ 
proof (cases fps_deriv  $G = 0$ )
  case False
    have [simp]:  $g$   $0 = 0$ 
      using has_fps_expansion_imp_0_eq_fps_nth_0[OF  $G(1)$ ]  $G(2)$  False by
simp
    have  $ana\_f$ :  $f$  analytic_on  $\{0\}$ 
      using  $F$  by (meson has_fps_expansion_imp_analytic_0)
    have  $ana\_g$ :  $g$  analytic_on  $\{0\}$ 
      using  $G$  by (meson has_fps_expansion_imp_analytic_0)
    have  $fg$ :  $(f \circ g)$  has_fps_expansion fps_expansion  $(f \circ g)$   $0$ 
      by (rule analytic_at_imp_has_fps_expansion_0 analytic_intros
        analytic_on_compose_gen  $ana\_f$   $ana\_g$ ) + auto

    have  $fls\_nth$  (fps_to_fls (fps_expansion  $(f \circ g)$   $0$ )) (int  $n$ ) = fps_nth (fps_compose
 $F$   $G$ )  $n$  for  $n$ 
      by (rule fls_nth_compose_aux has_laurent_expansion_fps  $F$   $G$  False  $fg$ ) +
    hence fps_expansion  $(f \circ g)$   $0 =$  fps_compose  $F$   $G$ 
      by (simp add: fps_eq_iff)
    thus ?thesis using  $fg$ 
      by simp
  next
    case True
    have [simp]:  $f$   $0 =$  fps_nth  $F$   $0$ 
      using  $F$  by (auto dest: has_fps_expansion_imp_0_eq_fps_nth_0)
    from True have fps_nth  $G$   $n = 0$  for  $n$ 
      using  $G(2)$  by (cases  $n$ ) (auto simp del: of_nat_Suc)
    hence [simp]:  $G = 0$ 
      by (auto simp: fps_eq_iff)
    have  $(\lambda\_ . f$   $0)$  has_fps_expansion fps_const  $(f$   $0)$ 
      by (intro fps_expansion_intros)
    also have eventually  $(\lambda x. g$   $x = 0)$  (nhds  $0$ )
      using  $G$  by (auto simp: has_fps_expansion_def)
    hence  $(\lambda\_ . f$   $0)$  has_fps_expansion fps_const  $(f$   $0) \longleftrightarrow$   $(f \circ g)$  has_fps_expansion
fps_const  $(f$   $0)$ 
      by (intro has_fps_expansion_cong) (auto elim!: eventually_mono)
    thus ?thesis
      by simp
  qed

```

lemma has_fps_expansion_fps_to_fls:

assumes f has_laurent_expansion fps_to_fls F

shows $(\lambda z. \text{if } z = 0 \text{ then fps_nth } F$ 0 **else** f $z)$ has_fps_expansion F

(is ? f' has_fps_expansion $_$)

proof $-$

have f has_laurent_expansion fps_to_fls $F \longleftrightarrow$? f' has_laurent_expansion

```

fps_to_fls F
  by (intro has_laurent_expansion_cong) (auto simp: eventually_at_filter)
  with assms show ?thesis
  by (auto simp: has_fps_expansion_to_laurent)
qed

lemma has_laurent_expansion_compose [laurent_expansion_intros]:
  fixes f g :: complex  $\Rightarrow$  complex
  assumes F: f has_laurent_expansion F
  assumes G: g has_laurent_expansion fps_to_fls G fps_nth G 0 = 0 G  $\neq$  0
  shows (f  $\circ$  g) has_laurent_expansion fls_compose_fps F G
proof -
  from assms have lim_g: g  $\rightarrow$  0
  by (subst tendsto_0_subdegree_iff_0[OF G(1)])
  (auto simp: fls_subdegree_fls_to_fps_subdegree_pos_iff)
  have ev1: eventually ( $\lambda z. g z \neq 0$ ) (at 0)
  using  $\langle G \neq 0 \rangle$  G(1) fps_to_fls_eq_0_iff has_laurent_expansion_fps
  has_laurent_expansion_nonzero_imp_eventually_nonzero by blast
  moreover have eventually ( $\lambda z. z \neq 0$ ) (at (0 :: complex))
  by (auto simp: eventually_at_filter)
  ultimately have ev: eventually ( $\lambda z. z \neq 0 \wedge g z \neq 0$ ) (at 0)
  by eventually_elim blast
  from ev1 and lim_g have lim_g': filterlim g (at 0) (at 0)
  by (auto simp: filterlim_at)
  define g' where g' = ( $\lambda z. \text{if } z = 0 \text{ then } \text{fps\_nth } G \ 0 \text{ else } g z$ )

  show ?thesis
  proof (cases F = 0)
    assume [simp]: F = 0
    have eventually ( $\lambda z. f z = 0$ ) (at 0)
    using F by (auto simp: has_laurent_expansion_def)
    hence eventually ( $\lambda z. f (g z) = 0$ ) (at 0)
    using lim_g' by (rule eventually_compose_filterlim)
    thus ?thesis
    by (auto simp: has_laurent_expansion_def)
  next
    assume [simp]: F  $\neq$  0
    define n where n = fls_subdegree F
    define f' where
      f' = ( $\lambda z. \text{if } z = 0 \text{ then } \text{fps\_nth } (\text{fls\_base\_factor\_to\_fps } F) \ 0 \text{ else } f z * z^{\text{powi } -n}$ )
    have (( $\lambda z. (f' \circ g') z * g z^{\text{powi } n}$ ) has_laurent_expansion fls_compose_fps
  F G
  unfolding f'_def n_def fls_compose_fps_def g'_def
  by (intro fps_expansion_intros laurent_expansion_intros has_fps_expansion_fps_to_fls
  has_fps_expansion_fls_base_factor_to_fps assms has_laurent_expansion_fps)
  also have ?this  $\longleftrightarrow$  ?thesis
  by (intro has_laurent_expansion_cong eventually_mono[OF ev])
  (auto simp: f'_def power_int_minus g'_def)

```


finally show ?thesis .
qed
qed

lemma has_laurent_expansion_fls_X_inv [laurent_expansion_intros]:
inverse has_laurent_expansion_fls_X_inv
using has_laurent_expansion_inverse[OF has_laurent_expansion_fps_X]
by (simp add: fls_inverse_X)

lemma zorder_times_analytic:
assumes f_analytic_on {z} g_analytic_on {z}
assumes eventually ($\lambda z. f z * g z \neq 0$) (at z)
shows zorder ($\lambda z. f z * g z$) z = zorder f z + zorder g z
proof -
have *: ($\lambda w. f (z + w)$) has_fps_expansion fps_expansion f z
 ($\lambda w. g (z + w)$) has_fps_expansion fps_expansion g z
 ($\lambda w. f (z + w) * g (z + w)$) has_fps_expansion fps_expansion f z *
fps_expansion g z
by (intro fps_expansion_intros analytic_at_imp_has_fps_expansion assms)+
have [simp]: fps_expansion f z $\neq 0$
proof
assume fps_expansion f z = 0
hence eventually ($\lambda z. f z * g z = 0$) (at z) using *(1)
by (auto simp: has_fps_expansion_0_iff_nhds_to_0' eventually_filtermap
eventually_at_filter
elim: eventually_mono)
with assms(3) have eventually ($\lambda z. False$) (at z)
by eventually_elim auto
thus False by simp
qed
have [simp]: fps_expansion g z $\neq 0$
proof
assume fps_expansion g z = 0
hence eventually ($\lambda z. f z * g z = 0$) (at z) using *(2)
by (auto simp: has_fps_expansion_0_iff_nhds_to_0' eventually_filtermap
eventually_at_filter
elim: eventually_mono)
with assms(3) have eventually ($\lambda z. False$) (at z)
by eventually_elim auto
thus False by simp
qed
from *[THEN has_fps_expansion_zorder] show ?thesis
by auto
qed

lemma zorder_const [simp]: $c \neq 0 \implies zorder (\lambda_. c) z = 0$
by (intro zorder_eqI[where S = UNIV]) auto

lemma zorder_prod_analytic:

```

assumes  $\bigwedge x. x \in A \implies f\ x\ \text{analytic\_on}\ \{z\}$ 
assumes eventually  $(\lambda z. (\prod x \in A. f\ x\ z) \neq 0)$  (at z)
shows  $\text{zorder } (\lambda z. \prod x \in A. f\ x\ z)\ z = (\sum x \in A. \text{zorder } (f\ x)\ z)$ 
using assms
proof (induction A rule: infinite_finite_induct)
  case (insert x A)
  have  $\text{zorder } (\lambda z. f\ x\ z * (\prod x \in A. f\ x\ z))\ z = \text{zorder } (f\ x)\ z + \text{zorder } (\lambda z. \prod x \in A. f\ x\ z)\ z$ 
  using insert.prem1 insert.hyps by (intro zorder_times_analytic analytic_intros)
  auto
  also have  $\text{zorder } (\lambda z. \prod x \in A. f\ x\ z)\ z = (\sum x \in A. \text{zorder } (f\ x)\ z)$ 
  using insert.prem1 insert.hyps by (intro insert.IH) (auto elim!: eventually_mono)
  finally show ?case using insert
    by simp
qed auto

```

```

lemma zorder_eq_0I:
assumes g analytic_on {z} g z ≠ 0
shows  $\text{zorder } g\ z = 0$ 
using analytic_at assms zorder_eqI by fastforce

```

```

lemma zorder_pos_iff:
assumes f holomorphic_on A open A z ∈ A frequently (λz. f z ≠ 0) (at z)
shows  $\text{zorder } f\ z > 0 \iff f\ z = 0$ 
proof -
  have f analytic_on {z}
  using assms analytic_at by blast
  hence  $*$ :  $(\lambda w. f\ (z + w))\ \text{has\_fps\_expansion}\ \text{fps\_expansion}\ f\ z$ 
  using analytic_at_imp_has_fps_expansion by blast
  have nz: fps_expansion f z ≠ 0
  proof
    assume fps_expansion f z = 0
    hence eventually (λz. f z = 0) (nhds z)
    using  $*$  by (auto simp: has_fps_expansion_def nhds_to_0' eventually_filtermap add_ac)
    hence eventually (λz. f z = 0) (at z)
    by (auto simp: eventually_at_filter elim: eventually_mono)
    with assms show False
    by (auto simp: frequently_def)
  qed
  from has_fps_expansion_zorder[OF * this] have eq: zorder f z = int (subdegree (fps_expansion f z))
  by auto
  moreover have  $\text{subdegree } (\text{fps\_expansion } f\ z) = 0 \iff \text{fps\_nth } (\text{fps\_expansion } f\ z)\ 0 \neq 0$ 
  using nz by (auto simp: subdegree_eq_0_iff)
  moreover have  $\text{fps\_nth } (\text{fps\_expansion } f\ z)\ 0 = f\ z$ 
  by (auto simp: fps_expansion_def)
  ultimately show ?thesis

```

by auto
qed

lemma *zorder_pos_iff'*:
assumes *f analytic_on {z} frequently* ($\lambda z. f z \neq 0$) (*at z*)
shows $zorder f z > 0 \longleftrightarrow f z = 0$
using *analytic_at assms zorder_pos_iff* **by** blast

lemma *zorder_ge_0*:
assumes *f analytic_on {z} frequently* ($\lambda z. f z \neq 0$) (*at z*)
shows $zorder f z \geq 0$
proof –
have *: ($\lambda w. f (z + w)$) *has_laurent_expansion fps_to_fls* (*fps_expansion f z*)
using *assms by (simp add: analytic_at_imp_has_fps_expansion_has_laurent_expansion_fps)*
from * *assms(2)* **have** *fps_to_fls* (*fps_expansion f z*) $\neq 0$
by (*auto simp: has_laurent_expansion_def frequently_def at_to_0' eventually_filtermap add_ac*)
with *has_laurent_expansion_zorder[OF *]* **show** ?thesis
by (*simp add: fls_subdegree_fls_to_fps*)
qed

lemma *zorder_eq_0_iff*:
assumes *f analytic_on {z} frequently* ($\lambda w. f w \neq 0$) (*at z*)
shows $zorder f z = 0 \longleftrightarrow f z \neq 0$
using *assms zorder_eq_0I zorder_pos_iff'* **by** fastforce

lemma *zorder_scale*:
assumes *f analytic_on {a * z} eventually* ($\lambda w. f w \neq 0$) (*at (a * z)*) $a \neq 0$
shows $zorder (\lambda w. f (a * w)) z = zorder f (a * z)$
proof –
from *assms(1)* **obtain** *r* **where** $r > 0$ *f holomorphic_on ball (a * z) r*
by (*auto simp: analytic_on_def*)
have *: *open (ball (a * z) r) connected (ball (a * z) r) a * z ∈ ball (a * z) r*
using $r < a \neq 0$ **by** (*auto simp: dist_norm*)
from *assms(2)* **have** *eventually* ($\lambda w. f w \neq 0 \wedge w \in ball (a * z) r - \{a * z\}$)
(*at (a * z)*)
using $\langle r > 0 \rangle$ **by** (*intro eventually_conj eventually_at_in_open*) *auto*
then obtain *z0* **where** $f z0 \neq 0 \wedge z0 \in ball (a * z) r - \{a * z\}$
using *eventually_happens[of _ at (a * z)]* **by** force
hence **: $\exists w \in ball (a * z) r. f w \neq 0$
by blast

define *n* **where** $n = nat (zorder f (a * z))$
obtain *r'* **where** *r'*:
(*if* $f (a * z) = 0$ *then* $0 < zorder f (a * z)$ *else* $zorder f (a * z) = 0$)
 $r' > 0$ *cball (a * z) r' ⊆ ball (a * z) r zor_poly f (a * z) holomorphic_on*
*cball (a * z) r'*
 $\bigwedge w. w \in cball (a * z) r' \implies$
 $f w = zor_poly f (a * z) w * (w - a * z) \wedge^n \wedge zor_poly f (a * z) w \neq 0$

```

unfolding n_def using zorder_exist_zero[OF r(2) **] by blast

show ?thesis
proof (rule zorder_eqI)
  show open (ball z (r' / norm a)) z ∈ ball z (r' / norm a)
    using r ⟨r' > 0⟩ ⟨a ≠ 0⟩ by auto
  have (*) a ' ball z (r' / cmod a) ⊆ cball (a * z) r'
  proof safe
    fix w assume w ∈ ball z (r' / cmod a)
    thus a * w ∈ cball (a * z) r'
    using dist_mult_left[of a z w] ⟨a ≠ 0⟩ by (auto simp: divide_simps mult_ac)
  qed
  thus (λw. a ^ n * (zor_poly f (a * z) ∘ (λw. a * w)) w) holomorphic_on ball
z (r' / norm a)
    using ⟨a ≠ 0⟩ by (intro holomorphic_on_compose_gen[OF _ r'(4)] holo-
morphic_intros) auto
  show a ^ n * (zor_poly f (a * z) ∘ (λw. a * w)) z ≠ 0
    using r' ⟨a ≠ 0⟩ by auto
  show f (a * w) = a ^ n * (zor_poly f (a * z) ∘ (*) a) w * (w - z) powi (zorder
f (a * z))
    if w ∈ ball z (r' / norm a) w ≠ z for w
  proof -
    have f (a * w) = zor_poly f (a * z) (a * w) * (a * (w - z)) ^ n
      using that r'(5)[of a * w] dist_mult_left[of a z w] ⟨a ≠ 0⟩ unfolding
ring_distrib
    by (auto simp: divide_simps mult_ac)
    also have ... = a ^ n * zor_poly f (a * z) (a * w) * (w - z) ^ n
      by (subst power_mult_distrib) (auto simp: mult_ac)
    also have (w - z) ^ n = (w - z) powi of_nat n
      by simp
    also have of_nat n = zorder f (a * z)
      using r'(1) by (auto simp: n_def split: if_splits)
    finally show ?thesis
      unfolding o_def n_def .
  qed
qed
qed

lemma zorder_compose_aux:
  assumes isolated_singularity_at f 0 not_essential f 0
  assumes G: g has_fps_expansion G G ≠ 0 g 0 = 0
  assumes eventually (λw. f w ≠ 0) (at 0)
  shows zorder (f ∘ g) 0 = zorder f 0 * subdegree G
proof -
  obtain F where F: f has_laurent_expansion F
    using not_essential_has_laurent_expansion_0[OF assms(1,2)] by blast
  have [simp]: fps_nth G 0 = 0
    using G ⟨g 0 = 0⟩ by (simp add: has_fps_expansion_imp_0_eq_fps_nth_0)
  note [simp] = ⟨G ≠ 0⟩ ⟨g 0 = 0⟩

```

```

have [simp]:  $F \neq 0$ 
  using has_laurent_expansion_eventually_nonzero_iff[of  $f 0 F$ ]  $F$  assms by
simp
have  $FG$ :  $(f \circ g)$  has_laurent_expansion fls_compose_fps  $F G$ 
  by (intro has_laurent_expansion_compose has_laurent_expansion_fps  $F G$ )
auto

have  $zorder (f \circ g) 0 = fls\_subdegree (fls\_compose\_fps F G)$ 
  using has_laurent_expansion_zorder_0 [OF  $FG$ ] by (auto simp: fls_compose_fps_eq_0_iff)
also have  $\dots = fls\_subdegree F * int (subdegree G)$ 
  by simp
also have  $fls\_subdegree F = zorder f 0$ 
  using has_laurent_expansion_zorder_0 [OF  $F$ ] by auto
finally show ?thesis .
qed

```

lemma *zorder_compose*:

```

assumes isolated_singularity_at  $f (g z)$  not_essential  $f (g z)$ 
assumes  $G$ :  $(\lambda x. g (z + x) - g z)$  has_fps_expansion  $G G \neq 0$ 
assumes eventually  $(\lambda w. f w \neq 0)$  (at  $(g z)$ )
shows  $zorder (f \circ g) z = zorder f (g z) * subdegree G$ 
proof -
define  $f'$  where  $f' = (\lambda w. f (g z + w))$ 
define  $g'$  where  $g' = (\lambda w. g (z + w) - g z)$ 
have  $zorder f (g z) = zorder f' 0$ 
  by (simp add: f'_def zorder_shift' add_ac)
have  $zorder (\lambda x. g x - g z) z = zorder g' 0$ 
  by (simp add: g'_def zorder_shift' add_ac)
have  $zorder (f \circ g) z = zorder (f' \circ g') 0$ 
  by (simp add: zorder_shift' f'_def g'_def add_ac o_def)
also have  $\dots = zorder f' 0 * int (subdegree G)$ 
proof (rule zorder_compose_aux)
  show isolated_singularity_at  $f' 0$  unfolding f'_def
    using assms has_laurent_expansion_isolated_0 not_essential_has_laurent_expansion
by blast
  show not_essential  $f' 0$  unfolding f'_def
    using assms has_laurent_expansion_not_essential_0 not_essential_has_laurent_expansion
by blast
  qed (use assms in ⟨auto simp: f'_def g'_def at_to_0' eventually_filtermap
add_ac⟩)
also have  $zorder f' 0 = zorder f (g z)$ 
  by (simp add: f'_def zorder_shift' add_ac)
finally show ?thesis .
qed

```

```

lemma fps_to_fls_eq_fls_const_iff [simp]:  $fps\_to\_fls F = fls\_const c \iff F =$ 
 $fps\_const c$ 
  using fps_to_fls_eq_iff by fastforce

```

```

lemma zorder_compose':
  assumes isolated_singularity_at f (g z) not_essential f (g z)
  assumes g_analytic_on {z}
  assumes eventually ( $\lambda w. f w \neq 0$ ) (at (g z))
  assumes eventually ( $\lambda w. g w \neq g z$ ) (at z)
  shows zorder (f  $\circ$  g) z = zorder f (g z) * zorder ( $\lambda x. g x - g z$ ) z
proof -
  obtain G where G [fps_expansion_intros]: ( $\lambda x. g (z + x)$ ) has_fps_expansion
  G
  using assms_analytic_at_imp_has_fps_expansion by blast
  have G': ( $\lambda x. g (z + x) - g z$ ) has_fps_expansion G - fps_const (g z)
  by (intro fps_expansion_intros)
  hence G'': ( $\lambda x. g (z + x) - g z$ ) has_laurent_expansion_fps_to_fls (G -
  fps_const (g z))
  using has_laurent_expansion_fps by blast
  have nz: G - fps_const (g z)  $\neq 0$ 
  using has_laurent_expansion_eventually_nonzero_iff[OF G''] assms by auto
  have zorder (f  $\circ$  g) z = zorder f (g z) * subdegree (G - fps_const (g z))
  proof (rule zorder_compose)
  show ( $\lambda x. g (z + x) - g z$ ) has_fps_expansion G - fps_const (g z)
  by (intro fps_expansion_intros)
  qed (use assms nz in auto)
  also have int (subdegree (G - fps_const (g z))) = fls_subdegree (fps_to_fls G
  - fls_const (g z))
  by (simp flip: fls_subdegree_fls_to_fps)
  also have ... = zorder ( $\lambda x. g x - g z$ ) z
  using has_laurent_expansion_zorder [OF G''] nz by auto
  finally show ?thesis .
qed

lemma analytic_at_cong:
  assumes eventually ( $\lambda x. f x = g x$ ) (nhds x) x = y
  shows f_analytic_on {x}  $\longleftrightarrow$  g_analytic_on {y}
proof -
  have g_analytic_on {x} if f_analytic_on {x} eventually ( $\lambda x. f x = g x$ ) (nhds x)
for f g
  proof -
  have ( $\lambda y. f (x + y)$ ) has_fps_expansion_fps_expansion f x
  by (rule analytic_at_imp_has_fps_expansion) fact
  also have ?this  $\longleftrightarrow$  ( $\lambda y. g (x + y)$ ) has_fps_expansion_fps_expansion f x
  using that by (intro has_fps_expansion_cong_refl) (auto simp: nhds_to_0'
  eventually_filtermap)
  finally show ?thesis
  by (rule has_fps_expansion_imp_analytic)
qed
from this[of f g] this[of g f] show ?thesis using assms
  by (auto simp: eq_commute)
qed

```

lemma *has_laurent_expansion_sin'* [laurent_expansion_intros]:
sin has_laurent_expansion fps_to_fls (fps_sin 1)
using *has_fps_expansion_sin'* *has_fps_expansion_to_laurent* **by** *blast*

lemma *has_laurent_expansion_cos'* [laurent_expansion_intros]:
cos has_laurent_expansion fps_to_fls (fps_cos 1)
using *has_fps_expansion_cos'* *has_fps_expansion_to_laurent* **by** *blast*

lemma *has_laurent_expansion_sin* [laurent_expansion_intros]:
 $(\lambda z. \sin (c * z))$ has_laurent_expansion fps_to_fls (fps_sin c)
by (*intro* *has_laurent_expansion_fps* *has_fps_expansion_sin*)

lemma *has_laurent_expansion_cos* [laurent_expansion_intros]:
 $(\lambda z. \cos (c * z))$ has_laurent_expansion fps_to_fls (fps_cos c)
by (*intro* *has_laurent_expansion_fps* *has_fps_expansion_cos*)

lemma *has_laurent_expansion_tan'* [laurent_expansion_intros]:
tan has_laurent_expansion fps_to_fls (fps_tan 1)
using *has_fps_expansion_tan'* *has_fps_expansion_to_laurent* **by** *blast*

lemma *has_laurent_expansion_tan* [laurent_expansion_intros]:
 $(\lambda z. \tan (c * z))$ has_laurent_expansion fps_to_fls (fps_tan c)
by (*intro* *has_laurent_expansion_fps* *has_fps_expansion_tan*)

8.5 More Laurent expansions

lemma *has_laurent_expansion_frequently_zero_iff*:
assumes $(\lambda w. f (z + w))$ has_laurent_expansion *F*
shows *frequently* $(\lambda z. f z = 0)$ (at *z*) $\longleftrightarrow F = 0$
using *assms* **by** (*simp* *add*: *frequently_def* *has_laurent_expansion_eventually_nonzero_iff*)

lemma *has_laurent_expansion_eventually_zero_iff*:
assumes $(\lambda w. f (z + w))$ has_laurent_expansion *F*
shows *eventually* $(\lambda z. f z = 0)$ (at *z*) $\longleftrightarrow F = 0$
using *assms*
by (*metis* *has_laurent_expansion_frequently_zero_iff* *has_laurent_expansion_isolated*

has_laurent_expansion_not_essential *laurent_expansion_def*
not_essential_frequently_0_imp_eventually_0 *not_essential_has_laurent_expansion*)

lemma *has_laurent_expansion_frequently_nonzero_iff*:
assumes $(\lambda w. f (z + w))$ has_laurent_expansion *F*
shows *frequently* $(\lambda z. f z \neq 0)$ (at *z*) $\longleftrightarrow F \neq 0$
using *assms* **by** (*metis* *has_laurent_expansion_eventually_zero_iff* *not_eventually*)

lemma *has_laurent_expansion_sum_list* [laurent_expansion_intros]:
assumes $\bigwedge x. x \in \text{set } xs \implies f x$ has_laurent_expansion *F* *x*
shows $(\lambda y. \sum x \leftarrow xs. f x y)$ has_laurent_expansion $(\sum x \leftarrow xs. F x)$
using *assms* **by** (*induction* *xs*) (*auto* *intro!*: *laurent_expansion_intros*)

lemma *has_laurent_expansion_prod_list* [*laurent_expansion_intros*]:
assumes $\bigwedge x. x \in \text{set } xs \implies f\ x \text{ has_laurent_expansion } F\ x$
shows $(\lambda y. \prod x \leftarrow xs. f\ x\ y) \text{ has_laurent_expansion } (\prod x \leftarrow xs. F\ x)$
using *assms* **by** (*induction xs*) (*auto intro!*: *laurent_expansion_intros*)

lemma *has_laurent_expansion_sum_mset* [*laurent_expansion_intros*]:
assumes $\bigwedge x. x \in \# I \implies f\ x \text{ has_laurent_expansion } F\ x$
shows $(\lambda y. \sum x \in \# I. f\ x\ y) \text{ has_laurent_expansion } (\sum x \in \# I. F\ x)$
using *assms* **by** (*induction I*) (*auto intro!*: *laurent_expansion_intros*)

lemma *has_laurent_expansion_prod_mset* [*laurent_expansion_intros*]:
assumes $\bigwedge x. x \in \# I \implies f\ x \text{ has_laurent_expansion } F\ x$
shows $(\lambda y. \prod x \in \# I. f\ x\ y) \text{ has_laurent_expansion } (\prod x \in \# I. F\ x)$
using *assms* **by** (*induction I*) (*auto intro!*: *laurent_expansion_intros*)

8.6 Formal convergence versus analytic convergence

The convergence of a sequence of formal power series and the convergence of the functions in the complex plane do not imply each other:

- If we have the sequence of constant power series $(1/n)_{n \geq 0}$, this clearly converges to the zero function analytically, but as a series of formal power series it is divergent (since the 0-th coefficient never stabilises).
- Conversely, the sequence of series $(n!x^n)_{n \geq 0}$ converges formally to 0, but the corresponding sequence of functions diverges for every $x \neq 0$.

However, if the sequence of series converges to some limit series h and the corresponding series of functions converges uniformly to some limit function $g(x)$, then h is also a series expansion of $g(x)$, i.e. in that case, formal and analytic convergence agree.

proposition *uniform_limit_imp_fps_expansion_eq*:
fixes $f :: 'a \Rightarrow \text{complex fps}$
assumes *lim1*: $(f \longrightarrow h)\ F$
assumes *lim2*: *uniform_limit* $A\ (\lambda x\ z. f'\ x\ z)\ g'\ F$
assumes *expansions*: *eventually* $(\lambda x. f'\ x \text{ has_fps_expansion } f\ x)\ F\ g'\ \text{has_fps_expansion } g$
assumes *holo*: *eventually* $(\lambda x. f'\ x \text{ holomorphic_on } A)\ F$
assumes *A*: *open* $A\ 0 \in A$
assumes *nontriv* [*simp*]: $F \neq \text{bot}$
shows $g = h$
proof (*rule fps_ext*)
fix $n :: \text{nat}$
have *eventually* $(\lambda x. \text{fps_nth } (f\ x)\ n = \text{fps_nth } h\ n)\ F$
using *lim1* **unfolding** *tendsto_fps_iff* **by** *blast*
hence *eventually* $(\lambda x. (\text{deriv } \widehat{\sim} n)\ (f'\ x)\ 0 / \text{fact } n = \text{fps_nth } h\ n)\ F$


```

    using expansions(1)
  proof eventually_elim
    case (elim x)
    have fps_nth (f x) n = (deriv  $\hat{\sim}$  n) (f' x) 0 / fact n
      by (rule fps_nth_fps_expansion) (use elim in auto)
    with elim show ?case
      by simp
  qed
  hence (( $\lambda$ x. (deriv  $\hat{\sim}$  n) (f' x) 0 / fact n)  $\longrightarrow$  fps_nth h n) F
    by (simp add: tendsto_eventually)

  moreover have (( $\lambda$ x. (deriv  $\hat{\sim}$  n) (f' x) 0)  $\longrightarrow$  (deriv  $\hat{\sim}$  n) g' 0) F
    using lim2
  proof (rule higher_deriv_complex_uniform_limit)
    show eventually ( $\lambda$ x. f' x holomorphic_on A) F
      using holo by eventually_elim auto
  qed (use A in auto)
  hence (( $\lambda$ x. (deriv  $\hat{\sim}$  n) (f' x) 0 / fact n)  $\longrightarrow$  (deriv  $\hat{\sim}$  n) g' 0 / fact n) F
    by (intro tendsto_divide) auto

  ultimately have fps_nth h n = (deriv  $\hat{\sim}$  n) g' 0 / fact n
    using tendsto_unique[OF nontriv] by blast
  also have ... = fps_nth g n
    by (rule fps_nth_fps_expansion [symmetric]) fact
  finally show fps_nth g n = fps_nth h n ..
qed

end

theory Meromorphic imports
  Laurent_Convergence
  Cauchy_Integral_Formula
  HOL-Analysis.Sparse_In
begin

```

8.7 Remove singular points

This function takes a complex function and returns a version of it where all removable singularities have been removed and all other singularities (to be more precise, unremovable discontinuities) are set to 0.

This is very useful since it is sometimes difficult to avoid introducing removable singularities. For example, consider the meromorphic functions $f(z) = z$ and $g(z) = 1/z$. Then a mathematician would write $f(z)g(z) = 1$, but in Isabelle of course this is not so.

Using the `remove_sings` function, we indeed have `remove_sings (λ z. f z * g z) = (λ _. 1)`.

definition `remove_sings :: (complex \Rightarrow complex) \Rightarrow complex \Rightarrow complex` **where**

$remove_sings\ f\ z = (if\ \exists\ c.\ f\ -z \rightarrow c\ then\ Lim\ (at\ z)\ f\ else\ 0)$

lemma *remove_sings_eqI* [intro]:

assumes $f\ -z \rightarrow c$

shows $remove_sings\ f\ z = c$

using *assms* **unfolding** *remove_sings_def* **by** (*auto simp: tendsto_Lim*)

lemma *remove_sings_at_analytic* [simp]:

assumes $f\ analytic_on\ \{z\}$

shows $remove_sings\ f\ z = f\ z$

using *assms* **by** (*intro remove_sings_eqI*) (*simp add: analytic_at_imp_isCont isContD*)

lemma *remove_sings_at_pole* [simp]:

assumes $is_pole\ f\ z$

shows $remove_sings\ f\ z = 0$

using *assms* **unfolding** *remove_sings_def is_pole_def*

by (*meson at_neq_bot not_tendsto_and_filterlim_at_infinity*)

lemma *eventually_remove_sings_eq_at*:

assumes $isolated_singularity_at\ f\ z$

shows $eventually\ (\lambda w.\ remove_sings\ f\ w = f\ w)\ (at\ z)$

proof –

from *assms* **obtain** r **where** $r:\ r > 0\ f\ analytic_on\ ball\ z\ r - \{z\}$

by (*auto simp: isolated_singularity_at_def*)

hence $*$: $f\ analytic_on\ \{w\}$ **if** $w \in ball\ z\ r - \{z\}$ **for** w

using r **that** **by** (*auto intro: analytic_on_subset*)

have $eventually\ (\lambda w.\ w \in ball\ z\ r - \{z\})\ (at\ z)$

using r **by** (*intro eventually_at_in_open*) *auto*

thus *?thesis*

by *eventually_elim* (*auto simp: remove_sings_at_analytic **)

qed

lemma *eventually_remove_sings_eq_nhds*:

assumes $f\ analytic_on\ \{z\}$

shows $eventually\ (\lambda w.\ remove_sings\ f\ w = f\ w)\ (nhds\ z)$

proof –

from *assms* **obtain** A **where** $A:\ open\ A\ z \in A\ f\ holomorphic_on\ A$

by (*auto simp: analytic_at*)

have $eventually\ (\lambda z.\ z \in A)\ (nhds\ z)$

by (*intro eventually_nhds_in_open A*)

thus *?thesis*

proof *eventually_elim*

case (*elim w*)

from *elim* **have** $f\ analytic_on\ \{w\}$

using A *analytic_at* **by** *blast*

thus *?case* **by** *auto*

qed

qed

```

lemma remove_sings_compose:
  assumes filtermap g (at z) = at z'
  shows remove_sings (f ∘ g) z = remove_sings f z'
proof (cases  $\exists c. f -z' \rightarrow c$ )
  case True
  then obtain c where c: f -z' → c
    by auto
  from c have remove_sings f z' = c
    by blast
  moreover from c have remove_sings (f ∘ g) z = c
    using c by (intro remove_sings_eqI) (auto simp: filterlim_def filtermap_compose
assms)
  ultimately show ?thesis
    by simp
next
  case False
  hence  $\neg(\exists c. (f \circ g) -z \rightarrow c)$ 
    by (auto simp: filterlim_def filtermap_compose assms)
  with False show ?thesis
    by (auto simp: remove_sings_def)
qed

```

```

lemma remove_sings_cong:
  assumes eventually  $(\lambda x. f x = g x)$  (at z) z = z'
  shows remove_sings f z = remove_sings g z'
proof (cases  $\exists c. f -z \rightarrow c$ )
  case True
  then obtain c where c: f -z → c by blast
  hence remove_sings f z = c
    by blast
  moreover have  $f -z \rightarrow c \iff g -z' \rightarrow c$ 
    using assms by (intro filterlim_cong refl) auto
  with c have remove_sings g z' = c
    by (intro remove_sings_eqI) auto
  ultimately show ?thesis
    by simp
next
  case False
  have  $f -z \rightarrow c \iff g -z' \rightarrow c$  for c
    using assms by (intro filterlim_cong) auto
  with False show ?thesis
    by (auto simp: remove_sings_def)
qed

```

```

lemma deriv_remove_sings_at_analytic [simp]:
  assumes f analytic_on {z}
  shows deriv (remove_sings f) z = deriv f z

```

```

apply (rule deriv_cong_ev)
apply (rule eventually_remove_sings_eq_nhds)
using assms by auto

```

```

lemma isolated_singularity_at_remove_sings [simp, intro]:
  assumes isolated_singularity_at f z
  shows isolated_singularity_at (remove_sings f) z
  using isolated_singularity_at_cong[OF eventually_remove_sings_eq_at[OF assms]
refl] assms
  by simp

```

```

lemma not_essential_remove_sings_iff [simp]:
  assumes isolated_singularity_at f z
  shows not_essential (remove_sings f) z  $\longleftrightarrow$  not_essential f z
  using not_essential_cong[OF eventually_remove_sings_eq_at[OF assms(1)]
refl]
  by simp

```

```

lemma not_essential_remove_sings [intro]:
  assumes isolated_singularity_at f z not_essential f z
  shows not_essential (remove_sings f) z
  by (subst not_essential_remove_sings_iff) (use assms in auto)

```

```

lemma
  assumes isolated_singularity_at f z
  shows is_pole_remove_sings_iff [simp]:
    is_pole (remove_sings f) z  $\longleftrightarrow$  is_pole f z
  and zorder_remove_sings [simp]:
    zorder (remove_sings f) z = zorder f z
  and zor_poly_remove_sings [simp]:
    zor_poly (remove_sings f) z = zor_poly f z
  and has_laurent_expansion_remove_sings_iff [simp]:
    ( $\lambda w. \text{remove\_sings } f (z + w)$ ) has_laurent_expansion F  $\longleftrightarrow$ 
    ( $\lambda w. f (z + w)$ ) has_laurent_expansion F
  and tendsto_remove_sings_iff [simp]:
    remove_sings f  $-z \rightarrow c \longleftrightarrow$  f  $-z \rightarrow c$ 
  by (intro is_pole_cong eventually_remove_sings_eq_at refl zorder_cong
    zor_poly_cong has_laurent_expansion_cong' tendsto_cong assms)+

```

```

lemma remove_sings_has_laurent_expansion [laurent_expansion_intros]:
  assumes f has_laurent_expansion F
  shows remove_sings f has_laurent_expansion F
proof -
  have remove_sings f has_laurent_expansion F  $\longleftrightarrow$  f has_laurent_expansion F
  proof (rule has_laurent_expansion_cong)
    have isolated_singularity_at f 0
    using assms by (metis has_laurent_expansion_isolated_0)
  thus eventually ( $\lambda x. \text{remove\_sings } f x = f x$ ) (at 0)
  by (rule eventually_remove_sings_eq_at)

```

```

qed auto
with assms show ?thesis
  by simp
qed

lemma get_all_poles_from_remove_sings:
  fixes f:: complex  $\Rightarrow$  complex
  defines ff $\equiv$ remove_sings f
  assumes f_holo:f holomorphic_on s - pts and finite pts
    pts $\subseteq$ s open s and not_ess: $\forall x \in pts. \text{not\_essential } f x$ 
  obtains pts' where
    pts'  $\subseteq$  pts finite pts' ff holomorphic_on s - pts'  $\forall x \in pts'. \text{is\_pole } ff x$ 
proof -
  define pts' where pts' = {x  $\in$  pts. is_pole f x}

  have pts'  $\subseteq$  pts unfolding pts'_def by auto
  then have finite pts' using  $\langle$ finite pts $\rangle$ 
    using rev_finite_subset by blast
  then have open (s - pts') using  $\langle$ open s $\rangle$ 
    by (simp add: finite_imp_closed open_Diff)

  have isolated:isolated_singularity_at f z if z  $\in$  pts for z
proof (rule isolated_singularity_at_holomorphic)
  show f holomorphic_on (s - (pts - {z})) -  $\{z\}$ 
    by (metis Diff_insert f_holo insert_Diff that)
  show open (s - (pts - {z}))
    by (meson assms(3) assms(5) finite_Diff finite_imp_closed open_Diff)
  show z  $\in$  s - (pts - {z})
    using assms(4) that by auto
qed

  have ff holomorphic_on s - pts'
proof (rule no_isolated_singularity')
  show (ff  $\longrightarrow$  ff z) (at z within s - pts') if z  $\in$  pts - pts' for z
proof -
  have at z within s - pts' = at z
    apply (rule at_within_open)
    using  $\langle$ open (s - pts') $\rangle$  that  $\langle$ pts $\subseteq$ s $\rangle$  by auto
  moreover have ff -z  $\rightarrow$  ff z
    unfolding ff_def
proof (subst tendsto_remove_sings_iff)
  show isolated_singularity_at f z
    apply (rule isolated)
    using that by auto
  have not_essential f z
    using not_ess that by auto
  moreover have  $\neg \text{is\_pole } f z$ 
    using that unfolding pts'_def by auto
  ultimately have  $\exists c. f -z \rightarrow c$ 

```

```

      unfolding not_essential_def by auto
    then show  $f -z \rightarrow \text{remove\_sings } f z$ 
      using remove_sings_eqI by blast
  qed
  ultimately show ?thesis by auto
  qed
  have ff holomorphic_on  $s - pts$ 
    using f_holo
  proof (elim holomorphic_transform)
    fix  $x$  assume  $x \in s - pts$ 
    then have  $f$  analytic_on  $\{x\}$ 
      using assms(3) assms(5) f_holo
    by (meson finite_imp_closed
        holomorphic_on_imp_analytic_at_open_Diff)
    from remove_sings_at_analytic[OF this]
    show  $f x = ff x$  unfolding ff_def by auto
  qed
  then show ff holomorphic_on  $s - pts' - (pts - pts')$ 
    apply (elim holomorphic_on_subset)
    by blast
  show open  $(s - pts')$ 
    by (simp add: ⟨open  $(s - pts')$ ⟩)
  show finite  $(pts - pts')$ 
    by (simp add: assms(3))
  qed
  moreover have  $\forall x \in pts'. \text{is\_pole } ff x$ 
    unfolding pts'_def
    using ff_def is_pole_remove_sings_iff isolated by blast
  moreover note  $\langle pts' \subseteq pts \rangle \langle \text{finite } pts' \rangle$ 
  ultimately show ?thesis using that by auto
  qed

lemma remove_sings_eq_0_iff:
  assumes not_essential f w
  shows  $\text{remove\_sings } f w = 0 \iff \text{is\_pole } f w \vee f -w \rightarrow 0$ 
  proof (cases is_pole f w)
  case False
  then obtain  $c$  where  $c: f -w \rightarrow c$ 
    using ⟨not_essential f w⟩ unfolding not_essential_def by auto
  then show ?thesis
    using False remove_sings_eqI by auto
  qed simp

lemma remove_sings_analytic_at:
  assumes isolated_singularity_at f z f -z → c
  shows  $\text{remove\_sings } f$  analytic_on  $\{z\}$ 
  proof -
  from assms(1) obtain  $A$  where  $A: \text{open } A z \in A f$  holomorphic_on  $(A - \{z\})$ 
    using analytic_imp_holomorphic isolated_singularity_at_iff_analytic_nhd by

```

```

auto
  have ana: f analytic_on (A - {z})
    by (subst analytic_on_open) (use A in auto)

  have remove_sings f holomorphic_on A
  proof (rule no_isolated_singularity)
    have f holomorphic_on (A - {z})
      by fact
    moreover have remove_sings f holomorphic_on (A - {z})  $\longleftrightarrow$  f holomor-
    phic_on (A - {z})
      by (intro holomorphic_cong remove_sings_at_analytic) (auto intro!: ana-
    lytic_on_subset[OF ana])
    ultimately show remove_sings f holomorphic_on (A - {z})
      by blast
    hence continuous_on (A - {z}) (remove_sings f)
      by (intro holomorphic_on_imp_continuous_on)
    moreover have isCont (remove_sings f) z
      using assms isCont_def remove_sings_eqI tendsto_remove_sings_iff by
    blast
    ultimately show continuous_on A (remove_sings f)
      by (metis A(1) DiffI closed_singleton continuous_on_eq_continuous_at
    open_Diff_singletonD)
    qed (use A(1) in auto)
    thus ?thesis
      using A(1,2) analytic_at by blast
  qed

lemma remove_sings_analytic_on:
  assumes f analytic_on A
  shows remove_sings f analytic_on A
proof -
  from assms obtain B where B: open B A  $\subseteq$  B f holomorphic_on B
  by (metis analytic_on_holomorphic)
  have remove_sings f holomorphic_on B  $\longleftrightarrow$  f holomorphic_on B
  proof (rule holomorphic_cong)
    fix z assume z  $\in$  B
    have f analytic_on {z}
      using  $\langle z \in B \rangle$  B holomorphic_on_imp_analytic_at by blast
    thus remove_sings f z = f z
      by (rule remove_sings_at_analytic)
  qed auto
  thus ?thesis
    using B analytic_on_holomorphic by blast
qed

lemma residue_remove_sings [simp]:
  assumes isolated_singularity_at f z
  shows residue (remove_sings f) z = residue f z
proof -

```

from *assms* **have** *eventually* $(\lambda w. \text{remove_sings } f w = f w)$ *(at z)*
using *eventually_remove_sings_eq_at* **by** *blast*
then obtain *A* **where** *A*: *open A z* $\in A \wedge w. w \in A - \{z\} \implies \text{remove_sings } f w = f w$
by *(subst (asm) eventually_at_topological)* *blast*
from *A(1,2)* **obtain** ε **where** $\varepsilon: \varepsilon > 0$ *cball z* $\varepsilon \subseteq A$
using *open_contains_cball_eq* **by** *blast*
hence *eq*: *remove_sings f w = f w* **if** $w \in \text{cball } z \varepsilon - \{z\}$ **for** *w*
using *that A* ε **by** *blast*

define *P* **where** $P = (\lambda f c \varepsilon. (\text{f has_contour_integral of_real } (2 * \text{pi}) * i * c)$
(circlepath z $\varepsilon))$
have $P (\text{remove_sings } f) c \delta \longleftrightarrow P f c \delta$ **if** $0 < \delta \delta < \varepsilon$ **for** *c* δ
unfolding *P_def* **using** $\langle \varepsilon > 0 \rangle$ **that** **by** *(intro has_contour_integral_cong)*
(auto simp: eq)
hence $*$: $(\forall \varepsilon > 0. \varepsilon < e \longrightarrow P (\text{remove_sings } f) c \varepsilon) \longleftrightarrow (\forall \varepsilon > 0. \varepsilon < e \longrightarrow P f c \varepsilon)$ **if** $e \leq \varepsilon$ **for** *c* *e*
using *that* **by** *force*
have $**$: $(\exists e > 0. \forall \varepsilon > 0. \varepsilon < e \longrightarrow P (\text{remove_sings } f) c \varepsilon) \longleftrightarrow (\exists e > 0. \forall \varepsilon > 0. \varepsilon < e \longrightarrow P f c \varepsilon)$ **for** *c*
proof
assume $(\exists e > 0. \forall \varepsilon > 0. \varepsilon < e \longrightarrow P (\text{remove_sings } f) c \varepsilon)$
then obtain *e* **where** $e > 0 \forall \varepsilon > 0. \varepsilon < e \longrightarrow P (\text{remove_sings } f) c \varepsilon$
by *blast*
thus $(\exists e > 0. \forall \varepsilon > 0. \varepsilon < e \longrightarrow P f c \varepsilon)$
by *(intro exI[of _ min e* $\varepsilon])$ *(use* $*$ *[of min e* $\varepsilon c] \varepsilon(1)$ **in** *auto*)
next
assume $(\exists e > 0. \forall \varepsilon > 0. \varepsilon < e \longrightarrow P f c \varepsilon)$
then obtain *e* **where** $e > 0 \forall \varepsilon > 0. \varepsilon < e \longrightarrow P f c \varepsilon$
by *blast*
thus $(\exists e > 0. \forall \varepsilon > 0. \varepsilon < e \longrightarrow P (\text{remove_sings } f) c \varepsilon)$
by *(intro exI[of _ min e* $\varepsilon])$ *(use* $*$ *[of min e* $\varepsilon c] \varepsilon(1)$ **in** *auto*)
qed
show *?thesis*
unfolding *residue_def* **by** *(intro arg_cong[of _ _ Eps] ext **[unfolded P_def])*
qed

lemma *remove_sings_shift_0*:
 $\text{remove_sings } f z = \text{remove_sings } (\lambda w. f (z + w)) 0$
proof *(cases* $\exists c. f -z \rightarrow c)$
case *True*
then obtain *c* **where** $c: f -z \rightarrow c$
by *blast*
from *c* **have** $\text{remove_sings } f z = c$
by *(rule remove_sings_eqI)*
moreover **have** $\text{remove_sings } (\lambda w. f (z + w)) 0 = c$
by *(rule remove_sings_eqI)* *(use* c **in** $\langle \text{simp_all add: at_to_0' filterlim_filtermap add_ac} \rangle$)
ultimately show *?thesis*


```

  by simp
next
case False
hence  $\neg(\exists c. (\lambda w. f (z + w)) -0 \rightarrow c)$ 
  by (simp add: at_to_0' filterlim_filtermap add_ac)
with False show ?thesis
  by (simp add: remove_sings_def)
qed

```

```

lemma remove_sings_shift_0':
  NO_MATCH 0 z  $\implies$  remove_sings f z = remove_sings ( $\lambda w. f (z + w)$ ) 0
  by (rule remove_sings_shift_0)

```

8.8 Meromorphicity

We define meromorphicity in terms of Laurent series expansions. This has the advantage of giving us a particularly simple definition that makes many of the lemmas below trivial because they reduce to similar statements about Laurent series that are already in the library.

On open domains, this definition coincides with the usual one from the literature, namely that the function be holomorphic on its domain except for a set of non-essential singularities that is *sparse*, i.e. that has no limit point inside the domain.

However, unlike the definitions found in the literature, our definition also makes sense for non-open domains: similarly to (*analytic_on*), we consider a function meromorphic on a non-open domain if it is meromorphic on some open superset of that domain.

We will prove all of this below.

```

definition meromorphic_on :: (complex  $\Rightarrow$  complex)  $\Rightarrow$  complex set  $\Rightarrow$  bool
  (infixl  $\langle$ (meromorphic'_on) $\rangle$  50) where
  f meromorphic_on A  $\longleftrightarrow$  ( $\forall z \in A. \exists F. (\lambda w. f (z + w))$  has_laurent_expansion F)

```

```

lemma meromorphic_at_iff: f meromorphic_on {z}  $\longleftrightarrow$  isolated_singularity_at f z  $\wedge$  not_essential f z

```

```

  unfolding meromorphic_on_def
  by (metis has_laurent_expansion_isolated has_laurent_expansion_not_essential insertI1 singletonD not_essential_has_laurent_expansion)

```

```

named_theorems meromorphic_intros

```

```

lemma meromorphic_on_empty [simp, intro]: f meromorphic_on {}
  by (auto simp: meromorphic_on_def)

```

```

lemma meromorphic_on_def':
  f meromorphic_on A  $\longleftrightarrow$  ( $\forall z \in A. (\lambda w. f (z + w))$  has_laurent_expansion laurent_expansion f z)

```

unfolding meromorphic_on_def **using** laurent_expansion_eqI **by** blast

lemma meromorphic_on_meromorphic_at: f meromorphic_on $A \longleftrightarrow (\forall x \in A. f$
meromorphic_on $\{x\})$
by (auto simp: meromorphic_on_def)

lemma meromorphic_on_altdef:
 f meromorphic_on $A \longleftrightarrow (\forall z \in A. \text{isolated_singularity_at } f z \wedge \text{not_essential } f$
 $z)$
by (subst meromorphic_on_meromorphic_at) (auto simp: meromorphic_at_iff)

lemma meromorphic_on_cong:
assumes $\bigwedge z. z \in A \implies \text{eventually } (\lambda w. f w = g w) \text{ (at } z) A = B$
shows f meromorphic_on $A \longleftrightarrow g$ meromorphic_on B
unfolding meromorphic_on_def **using** assms
by (intro ball_cong refl arg_cong[of _ _ Ex] has_laurent_expansion_cong ext)
(simp_all add: at_to_0' eventually_filtermap_add_ac)

lemma meromorphic_on_subset: f meromorphic_on $A \implies B \subseteq A \implies f$ meromorphic_on B
by (auto simp: meromorphic_on_def)

lemma meromorphic_on_Un:
assumes f meromorphic_on A f meromorphic_on B
shows f meromorphic_on $(A \cup B)$
using assms **unfolding** meromorphic_on_def **by** blast

lemma meromorphic_on_Union:
assumes $\bigwedge A. A \in B \implies f$ meromorphic_on A
shows f meromorphic_on $(\bigcup B)$
using assms **unfolding** meromorphic_on_def **by** blast

lemma meromorphic_on_UN:
assumes $\bigwedge x. x \in X \implies f$ meromorphic_on $(A x)$
shows f meromorphic_on $(\bigcup x \in X. A x)$
using assms **unfolding** meromorphic_on_def **by** blast

lemma meromorphic_on_imp_has_laurent_expansion:
assumes f meromorphic_on A $z \in A$
shows $(\lambda w. f (z + w))$ has_laurent_expansion laurent_expansion $f z$
using assms laurent_expansion_eqI **unfolding** meromorphic_on_def **by** blast

lemma meromorphic_on_open_nhd:
assumes f meromorphic_on A
obtains B **where** open B $A \subseteq B$ f meromorphic_on B
proof –
obtain F **where** $F: \bigwedge z. z \in A \implies (\lambda w. f (z + w))$ has_laurent_expansion $F z$
using assms **unfolding** meromorphic_on_def **by** metis
have $\exists Z. \text{open } Z \wedge z \in Z \wedge (\forall w \in Z - \{z\}. \text{eval_fls } (F z) (w - z) = f w)$ **if** $z: z$

```

∈ A for z
  proof -
    obtain Z where Z: open Z 0 ∈ Z ∧ w. w ∈ Z - {0} ⇒ eval_fls (F z) w =
    f (z + w)
    using F[OF z] unfolding has_laurent_expansion_def eventually_at_topological
  by blast
    hence open ((+) z ' Z) and z ∈ (+) z ' Z
    using open_translation by auto
    moreover have eval_fls (F z) (w - z) = f w if w ∈ (+) z ' Z - {z} for w
    using Z(3)[of w - z] that by auto
    ultimately show ?thesis by blast
  qed
  then obtain Z where Z:
    ∧ z. z ∈ A ⇒ open (Z z) ∧ z ∈ Z z ∧ (∀ w ∈ Z z - {z}. eval_fls (F z) (w - z)
    = f w)
  by metis

define B where B = (⋃ z ∈ A. Z z ∩ eball z (fls_conv_radius (F z)))
show ?thesis
proof (rule that[of B])
  show open B
  using Z unfolding B_def by auto
  show A ⊆ B
  unfolding B_def using F Z by (auto simp: has_laurent_expansion_def
zero_ereal_def)
  show f meromorphic_on B
  unfolding meromorphic_on_def
proof
  fix z assume z: z ∈ B
  show ∃ F. (λ w. f (z + w)) has_laurent_expansion F
  proof (cases z ∈ A)
    case True
    thus ?thesis using F by blast
  next
    case False
    then obtain z0 where z0: z0 ∈ A z ∈ Z z0 - {z0} dist z0 z <
    fls_conv_radius (F z0)
    using z False Z unfolding B_def by auto
    hence (λ w. eval_fls (F z0) (w - z0)) analytic_on {z}
    by (intro analytic_on_eval_fls' analytic_intros) (auto simp: dist_norm)
    also have ?this ⟷ f analytic_on {z}
    proof (intro analytic_at_cong_refl)
      have eventually (λ w. w ∈ Z z0 - {z0}) (nhds z)
      using Z[of z0] z0 by (intro eventually_nhds_in_open) auto
      thus ∀ F x in nhds z. eval_fls (F z0) (x - z0) = f x
      by eventually_elim (use Z[of z0] z0 in auto)
    qed
  finally show ?thesis
  using analytic_at_imp_has_fps_expansion has_fps_expansion_to_laurent

```

```

by blast
  qed
  qed
  qed
  qed

```

```

lemma meromorphic_on_not_essential:
  assumes f meromorphic_on {z}
  shows not_essential f z
  using assms has_laurent_expansion_not_essential unfolding meromorphic_on_def
  by blast

```

```

lemma meromorphic_on_isolated_singularity:
  assumes f meromorphic_on {z}
  shows isolated_singularity_at f z
  using assms has_laurent_expansion_isolated unfolding meromorphic_on_def
  by blast

```

```

lemma meromorphic_on_imp_not_islimpt_singularities:
  assumes f meromorphic_on A z ∈ A
  shows ¬z islimpt {w. ¬f analytic_on {w}}
  proof -
    obtain B where B: open B A ⊆ B f meromorphic_on B
      using assms meromorphic_on_open_nhd by blast
    obtain F where F: (λw. f (z + w)) has_laurent_expansion F
      using B assms(2) unfolding meromorphic_on_def by blast
    from F have isolated_singularity_at f z
      using has_laurent_expansion_isolated assms(2) by blast
    then obtain r where r: r > 0 f analytic_on ball z r - {z}
      unfolding isolated_singularity_at_def by blast
    have f analytic_on {w} if w ∈ ball z r - {z} for w
      by (rule analytic_on_subset[OF r(2)]) (use that in auto)
    hence eventually (λw. f analytic_on {w}) (at z)
      using eventually_at_in_open[of ball z r z] ⟨r > 0⟩ by (auto elim!: eventu-
ally_mono)
    thus ¬z islimpt {w. ¬f analytic_on {w}}
      by (auto simp: islimpt_conv_frequently_at frequently_def)
  qed

```

```

lemma meromorphic_on_imp_sparse_singularities:
  assumes f meromorphic_on A
  shows {w. ¬f analytic_on {w}} sparse_in A
  by (metis assms meromorphic_on_imp_not_islimpt_singularities
meromorphic_on_open_nhd sparse_in_open sparse_in_subset)

```

```

lemma meromorphic_on_imp_sparse_singularities':
  assumes f meromorphic_on A
  shows {w ∈ A. ¬f analytic_on {w}} sparse_in A
  using meromorphic_on_imp_sparse_singularities[OF assms]

```

by (rule sparse_in_subset2) auto

lemma meromorphic_onE:

assumes f meromorphic_on A

obtains pts where $pts \subseteq A$ pts sparse_in A f analytic_on $A - pts$

$\bigwedge z. z \in A \implies \text{not_essential } f z \bigwedge z. z \in A \implies \text{isolated_singularity_at } f z$

proof (rule that)

show $\{z \in A. \neg f \text{ analytic_on } \{z\}\}$ sparse_in A

using *assms* by (rule meromorphic_on_imp_sparse_singularities¹)

show f analytic_on $A - \{z \in A. \neg f \text{ analytic_on } \{z\}\}$

by (subst analytic_on_analytic_at) auto

qed (use *assms* in \langle auto intro: meromorphic_on_isolated_singularity meromorphic_on_not_essential meromorphic_on_subset \rangle)

lemma meromorphic_onI_weak:

assumes f analytic_on $A - pts \bigwedge z. z \in pts \implies \text{not_essential } f z$ pts sparse_in A

$pts \cap \text{frontier } A = \{\}$

shows f meromorphic_on A

unfolding meromorphic_on_def

proof

fix z assume $z: z \in A$

show $(\exists F. (\lambda w. f (z + w))) \text{ has_laurent_expansion } F$

proof (cases $z \in pts$)

case *False*

have f analytic_on $\{z\}$

using *assms*(1) by (rule analytic_on_subset) (use *False z* in auto)

thus ?thesis

using isolated_singularity_at_analytic not_essential_analytic

not_essential_has_laurent_expansion by blast

next

case *True*

show ?thesis

proof (rule exI, rule not_essential_has_laurent_expansion)

show not_essential $f z$

using *assms*(2) *True* by blast

next

show isolated_singularity_at $f z$

proof (rule isolated_singularity_at_holomorphic)

show open (interior $A - (pts - \{z\})$)

proof (rule open_diff_sparse_pts)

show $pts - \{z\}$ sparse_in interior A

using sparse_in_subset sparse_in_subset2 *assms* interior_subset

Diff_subset by metis

qed auto

next

have f analytic_on interior $A - (pts - \{z\}) - \{z\}$

using *assms*(1) by (rule analytic_on_subset) (use interior_subset in

blast)

```

      thus  $f$  holomorphic_on interior  $A - (pts - \{z\}) - \{z\}$ 
        by (rule analytic_imp_holomorphic)
    next
      from  $assms(4)$  and True have  $z \in \text{interior } A$ 
        unfolding frontier_def using closure_subset  $z$  by blast
      thus  $z \in \text{interior } A - (pts - \{z\})$ 
        by blast
    qed
  qed
qed
qed

```

lemma meromorphic_onI_open:

```

  assumes open  $A$   $f$  analytic_on  $A - pts \wedge z. z \in pts \implies \text{not\_essential } f z$ 
  assumes  $\wedge z. z \in A \implies \neg z \text{ islimpt } pts \cap A$ 
  shows  $f$  meromorphic_on  $A$ 
proof (rule meromorphic_onI_weak)
  have *:  $A - pts \cap A = A - pts$ 
    by blast
  show  $f$  analytic_on  $A - pts \cap A$ 
    unfolding * by fact
  show  $pts \cap A$  sparse_in  $A$ 
    using  $assms(1,4)$  by (subst sparse_in_open) auto
  show not_essential  $f z$  if  $z \in pts \cap A$  for  $z$ 
    using  $assms(3)$  that by blast
  show  $pts \cap A \cap \text{frontier } A = \{\}$ 
    using  $\langle \text{open } A \rangle$  frontier_disjoint_eq by blast
qed

```

lemma meromorphic_at_isCont_imp_analytic:

```

  assumes  $f$  meromorphic_on  $\{z\}$  isCont  $f z$ 
  shows  $f$  analytic_on  $\{z\}$ 
proof -
  have *:  $(\lambda w. f (z + w))$  has_laurent_expansion laurent_expansion  $f z$ 
    using  $assms$  by (auto intro: meromorphic_on_imp_has_laurent_expansion)
  from  $assms$  have  $\neg \text{is\_pole } f z$ 
    using is_pole_def not_tendsto_and_filterlim_at_infinity trivial_limit_at by
    (metis isContD)
  with * have  $fls\_subdegree (laurent\_expansion f z) \geq 0$ 
    using has_laurent_expansion_imp_is_pole linorder_not_le by blast
  hence **:  $(\lambda w. \text{eval\_fls } (laurent\_expansion f z) (w - z))$  analytic_on  $\{z\}$ 
    by (intro analytic_intros)+ (use * in  $\langle \text{auto simp: has\_laurent\_expansion\_def}$ 
    zero_ereal_def  $\rangle$ )
  have  $(\lambda w. \text{eval\_fls } (laurent\_expansion f z) (w - z)) -z \rightarrow \text{eval\_fls } (laurent\_expansion$ 
 $f z) (z - z)$ 
    by (intro isContD analytic_at_imp_isCont **)
  also have  $?this \longleftrightarrow f -z \rightarrow \text{eval\_fls } (laurent\_expansion f z) (z - z)$ 
    by (intro filterlim_cong refl)
    (use * in  $\langle \text{auto simp: has\_laurent\_expansion\_def at\_to\_0' eventually\_filtermap}$ 

```

```

add_ac)
  finally have  $f - z \rightarrow \text{eval\_fls } (\text{laurent\_expansion } f z) 0$ 
    by simp
  moreover from assms have  $f - z \rightarrow f z$ 
    by (auto intro: isContD)
  ultimately have ***:  $\text{eval\_fls } (\text{laurent\_expansion } f z) 0 = f z$ 
    by (rule LIM_unique)

  have eventually  $(\lambda w. f w = \text{eval\_fls } (\text{laurent\_expansion } f z) (w - z))$  (at z)
    using * by (simp add: has_laurent_expansion_def at_to_0' eventually_filtermap
add_ac eq_commute)
  hence eventually  $(\lambda w. f w = \text{eval\_fls } (\text{laurent\_expansion } f z) (w - z))$  (nhds z)
    unfolding eventually_at_filter by eventually_elim (use *** in auto)
  hence  $f \text{ analytic\_on } \{z\} \longleftrightarrow (\lambda w. \text{eval\_fls } (\text{laurent\_expansion } f z) (w - z))$ 
analytic_on {z}
    by (intro analytic_at_cong refl)
  with ** show ?thesis
    by simp
qed

lemma analytic_on_imp_meromorphic_on:
  assumes  $f \text{ analytic\_on } A$ 
  shows  $f \text{ meromorphic\_on } A$ 
  by (rule meromorphic_onI_weak[of _ _ {}]) (use assms in auto)

lemma meromorphic_on_compose:
  assumes  $g \text{ meromorphic\_on } A$   $f \text{ analytic\_on } B$   $f ' B \subseteq A$ 
  shows  $(\lambda w. g (f w)) \text{ meromorphic\_on } B$ 
  unfolding meromorphic_on_def
proof safe
  fix z assume  $z: z \in B$ 
  have  $f \text{ analytic\_on } \{z\}$ 
    using assms(2) by (rule analytic_on_subset) (use assms(3) z in auto)
  hence  $(\lambda w. f w - f z) \text{ analytic\_on } \{z\}$ 
    by (intro analytic_intros)
  then obtain F where F:  $(\lambda w. f (z + w) - f z) \text{ has\_fps\_expansion } F$ 
    using analytic_at_imp_has_fps_expansion by blast

  from assms(3) and z have  $f z \in A$ 
    by auto
  with assms(1) obtain G where G:  $(\lambda w. g (f z + w)) \text{ has\_laurent\_expansion } G$ 
  using z by (auto simp: meromorphic_on_def)

  have  $\exists H. ((\lambda w. g (f z + w)) \circ (\lambda w. f (z + w) - f z)) \text{ has\_laurent\_expansion } H$ 
  proof (cases F = 0)
  case False
  show ?thesis
  proof (rule exI, rule has_laurent_expansion_compose)

```

```

show ( $\lambda w. f(z + w) - f z$ ) has_laurent_expansion fps_to_fls F
  using F by (rule has_laurent_expansion_fps)
show fps_nth F 0 = 0
  using has_fps_expansion_imp_0_eq_fps_nth_0[OF F] by simp
qed fact+
next
case True
have ( $\lambda w. g(f z)$ ) has_laurent_expansion fls_const (g (f z))
  by auto
also have  $?this \longleftrightarrow (\lambda w. ((\lambda w. g(f z + w)) \circ (\lambda w. f(z + w) - f z)) w)$ 
  has_laurent_expansion fls_const (g (f z))
proof (rule has_laurent_expansion_cong, goal_cases)
  case 1
  from F and True have eventually ( $\lambda w. f(z + w) - f z = 0$ ) (nhds 0)
    by (simp add: has_fps_expansion_0_iff)
  hence eventually ( $\lambda w. f(z + w) - f z = 0$ ) (at 0)
    by (simp add: eventually_nhds_conv_at)
  thus  $?case$ 
    by eventually_elim auto
qed auto
finally show  $?thesis$ 
  by blast
qed
thus  $\exists H. (\lambda w. g(f(z + w)))$  has_laurent_expansion H
  by (simp add: o_def)
qed

```

```

lemma constant_on_imp_meromorphic_on:
  assumes f constant_on A open A
  shows f meromorphic_on A
  using assms analytic_on_imp_meromorphic_on
  constant_on_imp_analytic_on
  by blast

```

8.9 Nice meromorphicity

This is probably very non-standard, but we call a function “nicely meromorphic” if it is meromorphic and has no removable singularities. That means that the only singularities are poles. We furthermore require that the function return 0 at any pole, for convenience.

```

definition nicely_meromorphic_on :: (complex  $\Rightarrow$  complex)  $\Rightarrow$  complex set  $\Rightarrow$  bool
  (infixl  $\langle$ nicely'_meromorphic'_on $\rangle$  50)
  where f nicely_meromorphic_on A  $\longleftrightarrow$  f meromorphic_on A
   $\wedge (\forall z \in A. (is\_pole f z \wedge f z = 0) \vee f -z \rightarrow f z)$ 

```

```

lemma nicely_meromorphic_on_subset:
  f nicely_meromorphic_on A  $\Longrightarrow$  B  $\subseteq$  A  $\Longrightarrow$  f nicely_meromorphic_on B
  using meromorphic_on_subset unfolding nicely_meromorphic_on_def by blast

```


lemma *constant_on_imp_nicely_meromorphic_on*:

assumes *f* *constant_on* *A* *open* *A*
shows *f* *nicely_meromorphic_on* *A*
by (*meson* *analytic_at_imp_isCont* *assms*
constant_on_imp_holomorphic_on
constant_on_imp_meromorphic_on
holomorphic_on_imp_analytic_at *isCont_def*
nicely_meromorphic_on_def)

lemma *nicely_meromorphic_on_imp_analytic_at*:

assumes *f* *nicely_meromorphic_on* *A* *z* \in *A* \neg *is_pole* *f* *z*
shows *f* *analytic_on* $\{z\}$
proof (*rule* *meromorphic_at_isCont_imp_analytic*)
show *f* *meromorphic_on* $\{z\}$
by (*rule* *meromorphic_on_subset[of _ A]*) (*use* *assms* **in** \langle *auto simp: nicely_meromorphic_on_def* \rangle)
next
from *assms* **have** *f* $-z \rightarrow$ *f* *z*
by (*auto simp: nicely_meromorphic_on_def*)
thus *isCont* *f* *z*
by (*auto simp: isCont_def*)
qed

lemma *analytic_on_imp_nicely_meromorphic_on*:

f *analytic_on* *A* \implies *f* *nicely_meromorphic_on* *A*
by (*meson* *analytic_at_imp_isCont* *analytic_on_analytic_at*
analytic_on_imp_meromorphic_on *isContD* *nicely_meromorphic_on_def*)

lemma *remove_sings_meromorphic* [*meromorphic_intros*]:

assumes *f* *meromorphic_on* *A*
shows *remove_sings* *f* *meromorphic_on* *A*
unfolding *meromorphic_on_def*
proof *safe*
fix *z* **assume** *z*: *z* \in *A*
show $\exists F$. $(\lambda w$. *remove_sings* *f* (*z* + *w*)) *has_laurent_expansion* *F*
using *assms* *meromorphic_on_isolated_singularity* *meromorphic_on_not_essential*
not_essential_has_laurent_expansion *z* *meromorphic_on_subset* **by** *blast*
qed

lemma *remove_sings_nicely_meromorphic*:

assumes *f* *meromorphic_on* *A*
shows *remove_sings* *f* *nicely_meromorphic_on* *A*
proof –
have *remove_sings* *f* *meromorphic_on* *A*
by (*simp* *add: assms* *remove_sings_meromorphic*)
moreover **have** *is_pole* (*remove_sings* *f*) *z*
 \wedge *remove_sings* *f* *z* = 0 \vee
remove_sings *f* $-z \rightarrow$ *remove_sings* *f* *z*
if *z* \in *A* **for** *z*

```

proof (cases  $\exists c. f -z \rightarrow c$ )
  case True
  then have remove_sings f  $-z \rightarrow$  remove_sings f z
    by (metis remove_sings_eqI tendsto_remove_sings_iff
      assms meromorphic_onE that)
  then show ?thesis by simp
next
  case False
  then have is_pole (remove_sings f) z
     $\wedge$  remove_sings f z = 0
  by (meson is_pole_remove_sings_iff remove_sings_def
    remove_sings_eq_0_iff assms meromorphic_onE that)
  then show ?thesis by simp
qed
ultimately show ?thesis
  unfolding nicely_meromorphic_on_def by simp
qed

```

A nicely meromorphic function that frequently takes the same value in the neighbourhood of some point is constant.

lemma frequently_eq_meromorphic_imp_constant:

```

assumes frequently ( $\lambda z. f z = c$ ) (at z)
assumes f nicely_meromorphic_on A open A connected A z  $\in$  A
shows  $\bigwedge w. w \in A \implies f w = c$ 
proof -
  from assms(2) have mero: f meromorphic_on A
    by (auto simp: nicely_meromorphic_on_def)
  have sparse: {z. is_pole f z} sparse_in A
    using assms(2) mero
  by (meson assms(3) meromorphic_on_isolated_singularity meromorphic_on_meromorphic_at
    not_islimpt_poles_sparse_in_open)

  have eq: f w = c if w: w  $\in$  A  $\neg$ is_pole f w for w
  proof -
    have f w - c = 0
  proof (rule analytic_continuation[of  $\lambda w. f w - c$ ])
    show ( $\lambda w. f w - c$ ) holomorphic_on {z  $\in$  A.  $\neg$ is_pole f z} using assms(2)
    by (intro holomorphic_intros)
    (metis (mono_tags, lifting) analytic_imp_holomorphic analytic_on_analytic_at

      mem_Collect_eq nicely_meromorphic_on_imp_analytic_at)
  next
  from sparse and assms(3) have open (A - {z. is_pole f z})
    by (simp add: open_diff_sparse_pts)
  also have A - {z. is_pole f z} = {z  $\in$  A.  $\neg$ is_pole f z}
    by blast
  finally show open ... .
  next
  from sparse have connected (A - {z. is_pole f z})

```

```

    using assms(3,4) by (intro sparse_imp_connected) auto
  also have  $A - \{z. \text{is\_pole } f z\} = \{z \in A. \neg \text{is\_pole } f z\}$ 
    by blast
  finally show connected ... .
next
have eventually  $(\lambda w. w \in A)$  (at z)
  using assms by (intro eventually_at_in_open') auto
moreover have eventually  $(\lambda w. \neg \text{is\_pole } f w)$  (at z) using mero
  by (metis assms(5) eventually_not_pole meromorphic_onE)
ultimately have ev: eventually  $(\lambda w. w \in A \wedge \neg \text{is\_pole } f w)$  (at z)
  by eventually_elim auto
show z islimpt  $\{z \in A. \neg \text{is\_pole } f z \wedge f z = c\}$ 
  using frequently_eventually_frequently[OF assms(1) ev]
  unfolding islimpt_conv_frequently_at by (rule frequently_elim1) auto
next
from assms(1) have  $\neg \text{is\_pole } f z$ 
  by (simp add: frequently_const_imp_not_is_pole)
with  $\langle z \in A \rangle$  show  $z \in \{z \in A. \neg \text{is\_pole } f z\}$ 
  by auto
qed (use w in auto)
thus  $f w = c$ 
  by simp
qed

have not_pole:  $\neg \text{is\_pole } f w$  if  $w: w \in A$  for  $w$ 
proof -
  have eventually  $(\lambda w. \neg \text{is\_pole } f w)$  (at w)
    using mero by (metis eventually_not_pole meromorphic_onE that)
  moreover have eventually  $(\lambda w. w \in A)$  (at w)
    using w  $\langle \text{open } A \rangle$  by (intro eventually_at_in_open')
  ultimately have eventually  $(\lambda w. f w = c)$  (at w)
    by eventually_elim (auto simp: eq)
  hence  $\text{is\_pole } f w \longleftrightarrow \text{is\_pole } (\lambda_. c) w$ 
    by (intro is_pole_cong refl)
  thus ?thesis
    by simp
qed

show  $f w = c$  if  $w: w \in A$  for  $w$ 
  using eq[OF w not_pole[OF w]] .
qed

lemma nicely_meromorphic_on_unop:
  assumes f nicely_meromorphic_on A
  assumes f meromorphic_on A  $\implies (\lambda z. h (f z))$  meromorphic_on A
  assumes  $\bigwedge z. z \in A \implies \text{is\_pole } f z \implies \text{is\_pole } (\lambda z. h (f z)) z$ 
  assumes  $\bigwedge z. z \in f^{-1} A \implies \text{isCont } h z$ 
  assumes  $h 0 = 0$ 
  shows  $(\lambda z. h (f z))$  nicely_meromorphic_on A

```

```

unfolding nicely_meromorphic_on_def
proof (intro conjI ball)
  show ( $\lambda z. h (f z)$ ) meromorphic_on A
    using assms(1,2) by (auto simp: nicely_meromorphic_on_def)
next
  fix z assume z:  $z \in A$ 
  hence is_pole f z  $\wedge f z = 0 \vee f -z \rightarrow f z$ 
    using assms by (auto simp: nicely_meromorphic_on_def)
  thus is_pole ( $\lambda z. h (f z)$ ) z  $\wedge h (f z) = 0 \vee (\lambda z. h (f z)) -z \rightarrow h (f z)$ 
proof (rule disj_forward)
  assume is_pole f z  $\wedge f z = 0$ 
  thus is_pole ( $\lambda z. h (f z)$ ) z  $\wedge h (f z) = 0$ 
    using assms z by auto
next
  assume *:  $f -z \rightarrow f z$ 
  from z assms have isCont h (f z)
    by auto
  with * show ( $\lambda z. h (f z)$ )  $-z \rightarrow h (f z)$ 
    using continuous_within continuous_within_compose3 by blast
qed
qed

```

8.10 Closure properties and proofs for individual functions

```

lemma meromorphic_on_const [intro, meromorphic_intros]: ( $\lambda_. c$ ) meromorphic_on A
  by (rule analytic_on_imp_meromorphic_on) auto

```

```

lemma meromorphic_on_id [intro, meromorphic_intros]: ( $\lambda w. w$ ) meromorphic_on A
  by (auto simp: meromorphic_on_def intro!: exI laurent_expansion_intros)

```

```

lemma meromorphic_on_id' [intro, meromorphic_intros]: id meromorphic_on A
  by (auto simp: meromorphic_on_def intro!: exI laurent_expansion_intros)

```

```

lemma meromorphic_on_add [meromorphic_intros]:
  assumes f meromorphic_on A g meromorphic_on A
  shows ( $\lambda w. f w + g w$ ) meromorphic_on A
  unfolding meromorphic_on_def
  by (rule laurent_expansion_intros exI ball
    assms[THEN meromorphic_on_imp_has_laurent_expansion] | assumption)+

```

```

lemma meromorphic_on_uminus [meromorphic_intros]:
  assumes f meromorphic_on A
  shows ( $\lambda w. -f w$ ) meromorphic_on A
  unfolding meromorphic_on_def
  by (rule laurent_expansion_intros exI ball
    assms[THEN meromorphic_on_imp_has_laurent_expansion] | assumption)+

```

tion)+

lemma *meromorphic_on_diff* [*meromorphic_intros*]:
assumes f *meromorphic_on* A g *meromorphic_on* A
shows $(\lambda w. f w - g w)$ *meromorphic_on* A
using *meromorphic_on_add*[*OF assms(1) meromorphic_on_uminus*[*OF assms(2)*]]
by *simp*

lemma *meromorphic_on_mult* [*meromorphic_intros*]:
assumes f *meromorphic_on* A g *meromorphic_on* A
shows $(\lambda w. f w * g w)$ *meromorphic_on* A
unfolding *meromorphic_on_def*
by (*rule laurent_expansion_intros exI ballI*
assms[*THEN meromorphic_on_imp_has_laurent_expansion*] | *assump-*
tion)+

lemma *meromorphic_on_power* [*meromorphic_intros*]:
assumes f *meromorphic_on* A
shows $(\lambda w. f w ^ n)$ *meromorphic_on* A
unfolding *meromorphic_on_def*
by (*rule laurent_expansion_intros exI ballI*
assms[*THEN meromorphic_on_imp_has_laurent_expansion*] | *assump-*
tion)+

lemma *meromorphic_on_powi* [*meromorphic_intros*]:
assumes f *meromorphic_on* A
shows $(\lambda w. f w \text{ powi } n)$ *meromorphic_on* A
unfolding *meromorphic_on_def*
by (*rule laurent_expansion_intros exI ballI*
assms[*THEN meromorphic_on_imp_has_laurent_expansion*] | *assump-*
tion)+

lemma *meromorphic_on_scaleR* [*meromorphic_intros*]:
assumes f *meromorphic_on* A
shows $(\lambda w. \text{scaleR } x (f w))$ *meromorphic_on* A
unfolding *meromorphic_on_def*
by (*rule laurent_expansion_intros exI ballI*
assms[*THEN meromorphic_on_imp_has_laurent_expansion*] | *assump-*
tion)+

lemma *meromorphic_on_inverse* [*meromorphic_intros*]:
assumes f *meromorphic_on* A
shows $(\lambda w. \text{inverse } (f w))$ *meromorphic_on* A
unfolding *meromorphic_on_def*
by (*rule laurent_expansion_intros exI ballI*
assms[*THEN meromorphic_on_imp_has_laurent_expansion*] | *assump-*
tion)+

lemma *meromorphic_on_divide* [*meromorphic_intros*]:

assumes f meromorphic_on A g meromorphic_on A
shows $(\lambda w. f w / g w)$ meromorphic_on A
using meromorphic_on_mult[OF assms(1) meromorphic_on_inverse[OF assms(2)]]
by (simp add: field_simps)

lemma meromorphic_on_sum [meromorphic_intros]:
assumes $\bigwedge i. i \in I \implies f i$ meromorphic_on A
shows $(\lambda w. \sum_{i \in I}. f i w)$ meromorphic_on A
unfolding meromorphic_on_def
by (rule laurent_expansion_intros exI ballI
 assms[THEN meromorphic_on_imp_has_laurent_expansion] | assumption)
 +

lemma meromorphic_on_sum_list [meromorphic_intros]:
assumes $\bigwedge i. i \in \text{set } fs \implies f i$ meromorphic_on A
shows $(\lambda w. \sum_{i \leftarrow fs}. f i w)$ meromorphic_on A
unfolding meromorphic_on_def
by (rule laurent_expansion_intros exI ballI
 assms[THEN meromorphic_on_imp_has_laurent_expansion] | assumption)
 +

lemma meromorphic_on_sum_mset [meromorphic_intros]:
assumes $\bigwedge i. i \in \# I \implies f i$ meromorphic_on A
shows $(\lambda w. \sum_{i \in \# I}. f i w)$ meromorphic_on A
unfolding meromorphic_on_def
by (rule laurent_expansion_intros exI ballI
 assms[THEN meromorphic_on_imp_has_laurent_expansion] | assumption)
 +

lemma meromorphic_on_prod [meromorphic_intros]:
assumes $\bigwedge i. i \in I \implies f i$ meromorphic_on A
shows $(\lambda w. \prod_{i \in I}. f i w)$ meromorphic_on A
unfolding meromorphic_on_def
by (rule laurent_expansion_intros exI ballI
 assms[THEN meromorphic_on_imp_has_laurent_expansion] | assumption)
 +

lemma meromorphic_on_prod_list [meromorphic_intros]:
assumes $\bigwedge i. i \in \text{set } fs \implies f i$ meromorphic_on A
shows $(\lambda w. \prod_{i \leftarrow fs}. f i w)$ meromorphic_on A
unfolding meromorphic_on_def
by (rule laurent_expansion_intros exI ballI
 assms[THEN meromorphic_on_imp_has_laurent_expansion] | assumption)
 +

lemma meromorphic_on_prod_mset [meromorphic_intros]:
assumes $\bigwedge i. i \in \# I \implies f i$ meromorphic_on A
shows $(\lambda w. \prod_{i \in \# I}. f i w)$ meromorphic_on A
unfolding meromorphic_on_def

by (rule laurent_expansion_intros exI ballI
 assms[THEN meromorphic_on_imp_has_laurent_expansion] | assumption)
 +

lemma nicely_meromorphic_on_const [intro]: $(\lambda_. c)$ nicely_meromorphic_on
 A
unfolding nicely_meromorphic_on_def **by** auto

lemma nicely_meromorphic_on_cmult_left [intro]:
assumes f nicely_meromorphic_on A
shows $(\lambda z. c * f z)$ nicely_meromorphic_on A
proof (cases $c = 0$)
 case [simp]: False
show ?thesis
using assms **by** (rule nicely_meromorphic_on_unop) (auto intro!: meromorphic_intros)
qed (auto intro!: nicely_meromorphic_on_const)

lemma nicely_meromorphic_on_cmult_right [intro]:
assumes f nicely_meromorphic_on A
shows $(\lambda z. f z * c)$ nicely_meromorphic_on A
using nicely_meromorphic_on_cmult_left[OF assms, of c] **by** (simp add: mult.commute)

lemma nicely_meromorphic_on_scaleR [intro]:
assumes f nicely_meromorphic_on A
shows $(\lambda z. c *_R f z)$ nicely_meromorphic_on A
using assms **by** (auto simp: scaleR_conv_of_real)

lemma nicely_meromorphic_on_uminus [intro]:
assumes f nicely_meromorphic_on A
shows $(\lambda z. -f z)$ nicely_meromorphic_on A
using nicely_meromorphic_on_cmult_left[OF assms, of -1] **by** simp

lemma meromorphic_on_If [meromorphic_intros]:
assumes f meromorphic_on A g meromorphic_on B
assumes $\bigwedge z. z \in A \implies z \in B \implies f z = g z$ open A open B $C \subseteq A \cup B$
shows $(\lambda z. \text{if } z \in A \text{ then } f z \text{ else } g z)$ meromorphic_on C
proof (rule meromorphic_on_subset)
show $(\lambda z. \text{if } z \in A \text{ then } f z \text{ else } g z)$ meromorphic_on $(A \cup B)$
proof (rule meromorphic_on_Un)
have $(\lambda z. \text{if } z \in A \text{ then } f z \text{ else } g z)$ meromorphic_on A \longleftrightarrow f meromorphic_on
 A
proof (rule meromorphic_on_cong)
fix z **assume** $z \in A$
hence eventually $(\lambda z. z \in A)$ (at z)
using $\langle \text{open } A \rangle$ **by** (intro eventually_at_in_open') auto
thus $\forall_F w$ in at z . (if $w \in A$ then $f w$ else $g w$) = $f w$
by eventually_elim auto
qed auto

```

with assms(1) show ( $\lambda z. \text{if } z \in A \text{ then } f z \text{ else } g z$ ) meromorphic_on A
  by blast
next
have ( $\lambda z. \text{if } z \in A \text{ then } f z \text{ else } g z$ ) meromorphic_on B  $\longleftrightarrow$  g meromorphic_on
B
proof (rule meromorphic_on_cong)
  fix z assume  $z \in B$ 
  hence eventually ( $\lambda z. z \in B$ ) (at z)
    using  $\langle \text{open } B \rangle$  by (intro eventually_at_in_open') auto
  thus  $\forall_F w \text{ in } \text{at } z. (\text{if } w \in A \text{ then } f w \text{ else } g w) = g w$ 
    by eventually_elim (use assms(3)) in auto
qed auto
with assms(2) show ( $\lambda z. \text{if } z \in A \text{ then } f z \text{ else } g z$ ) meromorphic_on B
  by blast
qed
qed fact

```

```

lemma meromorphic_on_deriv [meromorphic_intros]:
  f meromorphic_on A  $\implies$  deriv f meromorphic_on A
by (metis meromorphic_on_def isolated_singularity_at_deriv meromorphic_on_isolated_singularity

```

```

    meromorphic_on_meromorphic_at meromorphic_on_not_essential
not_essential_deriv
    not_essential_has_laurent_expansion)

```

```

lemma meromorphic_on_higher_deriv [meromorphic_intros]:
  f meromorphic_on A  $\implies$  (deriv  $\widehat{\sim}$  n) f meromorphic_on A
by (induction n) (auto intro!; meromorphic_intros)

```

```

lemma analytic_on_eval_fps [analytic_intros]:
  assumes f analytic_on A
  assumes  $\bigwedge z. z \in A \implies \text{norm } (f z) < \text{fps\_conv\_radius } F$ 
  shows ( $\lambda w. \text{eval\_fps } F (f w)$ ) analytic_on A
  by (rule analytic_on_compose[OF assms(1) analytic_on_eval_fps, unfolded
o_def])
    (use assms(2)) in auto)

```

```

lemma meromorphic_on_eval_fps [meromorphic_intros]:
  assumes f analytic_on A
  assumes  $\bigwedge z. z \in A \implies \text{norm } (f z) < \text{fps\_conv\_radius } F$ 
  shows ( $\lambda w. \text{eval\_fps } F (f w)$ ) meromorphic_on A
by (rule analytic_on_imp_meromorphic_on analytic_intros analytic_on_eval_fps
assms)+

```

```

lemma meromorphic_on_eval_fls [meromorphic_intros]:
  assumes f analytic_on A
  assumes  $\bigwedge z. z \in A \implies \text{norm } (f z) < \text{fls\_conv\_radius } F$ 
  shows ( $\lambda w. \text{eval\_fls } F (f w)$ ) meromorphic_on A
proof (cases fls_conv_radius F  $> 0$ )

```



```

case False
with assms(2) have  $A = \{\}$ 
  by (metis all_not_in_conv ereal_less(2) norm_eq_zero order.strict_trans
      zero_ereal_def zero_less_norm_iff)
thus ?thesis
  by auto
next
case True
have  $F: \text{eval\_fls } F \text{ has\_laurent\_expansion } F$ 
  using True by (rule eval_fls_has_laurent_expansion)
show ?thesis
proof (rule meromorphic_on_compose[OF __ assms(1)])
  show  $\text{eval\_fls } F \text{ meromorphic\_on eball } 0 \text{ (fls\_conv\_radius } F)$ 
  proof (rule meromorphic_onI_open)
    show  $\text{eval\_fls } F \text{ analytic\_on eball } 0 \text{ (fls\_conv\_radius } F) - \{0\}$ 
    by (rule analytic_on_eval_fls) auto
    show  $\text{not\_essential } (\text{eval\_fls } F) z \text{ if } z \in \{0\} \text{ for } z$ 
    using that F has_laurent_expansion_not_essential_0 by blast
    qed (auto simp: islimpt_finite)
  qed (use assms(2) in auto)
qed

lemma meromorphic_on_imp_analytic_cosparse:
  assumes  $f \text{ meromorphic\_on } A$ 
  shows  $\text{eventually } (\lambda z. f \text{ analytic\_on } \{z\}) \text{ (cosparse } A)$ 
  unfolding eventually_cosparse using assms meromorphic_on_imp_sparse_singularities
by auto

lemma meromorphic_on_imp_not_pole_cosparse:
  assumes  $f \text{ meromorphic\_on } A$ 
  shows  $\text{eventually } (\lambda z. \neg \text{is\_pole } f z) \text{ (cosparse } A)$ 
proof –
  have  $\text{eventually } (\lambda z. f \text{ analytic\_on } \{z\}) \text{ (cosparse } A)$ 
  by (rule meromorphic_on_imp_analytic_cosparse) fact
  thus ?thesis
  by eventually_elim (blast dest: analytic_at_imp_no_pole)
qed

lemma eventually_remove_sings_eq:
  assumes  $f \text{ meromorphic\_on } A$ 
  shows  $\text{eventually } (\lambda z. \text{remove\_sings } f z = f z) \text{ (cosparse } A)$ 
proof –
  have  $\text{eventually } (\lambda z. f \text{ analytic\_on } \{z\}) \text{ (cosparse } A)$ 
  using assms by (rule meromorphic_on_imp_analytic_cosparse)
  thus ?thesis
  by eventually_elim auto
qed

lemma remove_sings_constant_on_open_iff:

```

```

assumes  $f$  meromorphic_on  $A$  open  $A$ 
shows remove_sings  $f$  constant_on  $A \longleftrightarrow (\exists c. \forall_{\approx} x \in A. f x = c)$ 
proof
  assume remove_sings  $f$  constant_on  $A$ 
  then obtain  $c$  where  $c$ : remove_sings  $f z = c$  if  $z \in A$  for  $z$ 
    using that by (auto simp: constant_on_def)
  have  $\forall_{\approx} x \in A. x \in A$ 
    using  $\langle$ open  $A$  $\rangle$  by (simp add: eventually_in_cosparse)
  hence  $\forall_{\approx} x \in A. f x = c$ 
    using eventually_remove_sings_eq[OF assms(1)] by eventually_elim (use  $c$ 
in auto)
  thus  $\exists c. \forall_{\approx} x \in A. f x = c$ 
    by blast
next
  assume  $\exists c. \forall_{\approx} x \in A. f x = c$ 
  then obtain  $c$  where  $c$ :  $\forall_{\approx} x \in A. f x = c$ 
    by blast
  have  $\forall_{\approx} x \in A. \text{remove\_sings } f x = c$ 
    using eventually_remove_sings_eq[OF assms(1)]  $c$  by eventually_elim auto
  hence remove_sings  $f z = c$  if  $z \in A$  for  $z$  using that
    by (meson assms(2)  $c$  eventually_cosparse_open_eq remove_sings_eqI tendsto_eventually)
  thus remove_sings  $f$  constant_on  $A$ 
    unfolding constant_on_def by blast
qed

```

A meromorphic function on a connected domain takes any given value either almost everywhere or almost nowhere.

lemma meromorphic_imp_constant_or_avoid:

```

assumes mero:  $f$  meromorphic_on  $A$  and  $A$ : open  $A$  connected  $A$ 
shows eventually  $(\lambda z. f z = c)$  (cosparse  $A$ )  $\vee$  eventually  $(\lambda z. f z \neq c)$  (cosparse  $A$ )
proof –
  have eventually  $(\lambda z. f z = c)$  (cosparse  $A$ ) if freq: frequently  $(\lambda z. f z = c)$ 
(cosparse  $A$ )
  proof –
    let  $?f = \text{remove\_sings } f$ 
    have  $ev$ : eventually  $(\lambda z. ?f z = f z)$  (cosparse  $A$ )
      by (rule eventually_remove_sings_eq) fact
    have frequently  $(\lambda z. ?f z = c)$  (cosparse  $A$ )
      using frequently_eventually_frequently[OF freq  $ev$ ] by (rule frequently_elim1)
  auto
  then obtain  $z0$  where  $z0$ :  $z0 \in A$  frequently  $(\lambda z. ?f z = c)$  (at  $z0$ )
    using  $A$  by (auto simp: eventually_cosparse_open_eq frequently_def)
  have mero':  $?f$  nicely_meromorphic_on  $A$ 
    using mero remove_sings_nicely_meromorphic by blast
  have  $eq$ :  $?f w = c$  if  $w$ :  $w \in A$  for  $w$ 
    using frequently_eq_meromorphic_imp_constant[OF  $z0(2)$  mero']  $A$   $z0(1)$ 
  w by blast

```

```

have eventually ( $\lambda z. z \in A$ ) (cosparse A)
  by (rule eventually_in_cosparse) (use A in auto)
thus eventually ( $\lambda z. f z = c$ ) (cosparse A)
  using ev by eventually_elim (use eq in auto)
qed
thus ?thesis
  by (auto simp: frequently_def)
qed

lemma nicely_meromorphic_imp_constant_or_avoid:
  assumes f nicely_meromorphic_on A open A connected A
  shows ( $\forall x \in A. f x = c$ )  $\vee$  ( $\forall \approx x \in A. f x \neq c$ )
proof -
  have ( $\forall \approx x \in A. f x = c$ )  $\vee$  ( $\forall \approx x \in A. f x \neq c$ )
    by (intro meromorphic_imp_constant_or_avoid)
      (use assms in <auto simp: nicely_meromorphic_on_def>)
  thus ?thesis
proof
  assume ev:  $\forall \approx x \in A. f x = c$ 
  have  $\forall x \in A. f x = c$ 
proof
  fix x assume x:  $x \in A$ 
  have not_essential f x
    using assms x unfolding nicely_meromorphic_on_def by blast
  moreover have is_pole f x  $\longleftrightarrow$  is_pole ( $\lambda \_. c$ ) x
  by (intro is_pole_cong) (use ev x in <auto simp: eventually_cosparse_open_eq
  assms>)
  hence  $\neg$ is_pole f x
    by auto
  ultimately have f analytic_on {x}
    using assms(1) nicely_meromorphic_on_imp_analytic_at x by blast
  hence f  $\rightarrow$  f x
    by (intro isContD analytic_at_imp_isCont)
  also have ?this  $\longleftrightarrow$  ( $\lambda \_. c$ )  $\rightarrow$  f x
  by (intro tendsto_cong) (use ev x in <auto simp: eventually_cosparse_open_eq
  assms>)
  finally have ( $\lambda \_. c$ )  $\rightarrow$  f x .
  moreover have ( $\lambda \_. c$ )  $\rightarrow$  c
    by simp
  ultimately show f x = c
    using LIM_unique by blast
qed
thus ?thesis
  by blast
qed blast
qed

lemma nicely_meromorphic_onE:
  assumes f nicely_meromorphic_on A

```

obtains *pts* **where** $pts \subseteq A$ *pts sparse_in A*
 f *analytic_on A - pts*
 $\bigwedge z. z \in pts \implies is_pole\ f\ z \wedge f\ z = 0$
proof –
define *pts* **where** $pts = \{z \in A. \neg f\ analytic_on\ \{z\}\}$
have $pts \subseteq A$ *pts sparse_in A*
using *assms unfolding pts_def nicely_meromorphic_on_def*
by (*auto intro:meromorphic_on_imp_sparse_singularities'*)
moreover **have** f *analytic_on A - pts* **unfolding** *pts_def*
by (*subst analytic_on_analytic_at*) *auto*
moreover **have** $\bigwedge z. z \in pts \implies is_pole\ f\ z \wedge f\ z = 0$
by (*metis (no_types, lifting) remove_sings_eqI*
remove_sings_eq_0_iff assms is_pole_imp_not_essential
mem_Collect_eq nicely_meromorphic_on_def
nically_meromorphic_on_imp_analytic_at pts_def)
ultimately show *?thesis* **using** *that* **by** *auto*
qed

lemma *nically_meromorphic_onI_open*:
assumes *open A* **and**
 $analytic: f$ *analytic_on A - pts* **and**
 $pole: \bigwedge x. x \in pts \implies is_pole\ f\ x \wedge f\ x = 0$ **and**
 $isolated: \bigwedge x. x \in A \implies isolated_singularity_at\ f\ x$
shows f *nically_meromorphic_on A*
proof –
have f *meromorphic_on A*
proof (*rule meromorphic_onI_open*)
show $\bigwedge z. z \in pts \implies not_essential\ f\ z$
using *pole unfolding not_essential_def* **by** *auto*
show $\bigwedge z. z \in A \implies \neg z\ islimpt\ pts \cap A$
by (*metis assms(3) assms(4) inf_commute inf_le2*
islimpt_subset mem_Collect_eq not_islimpt_poles subsetI)
qed *fact+*
moreover **have** $(\forall z \in A. (is_pole\ f\ z \wedge f\ z = 0) \vee f\ -z \rightarrow f\ z)$
by (*meson DiffI analytic_analytic_at_imp_isCont*
analytic_on_analytic_at assms(3) isContD)
ultimately show *?thesis* **unfolding** *nically_meromorphic_on_def*
by *auto*
qed

lemma *nically_meromorphic_without_singularities*:
assumes f *nically_meromorphic_on A* $\forall z \in A. \neg is_pole\ f\ z$
shows f *analytic_on A*
by (*meson analytic_on_analytic_at assms*
nically_meromorphic_on_imp_analytic_at)

lemma *meromorphic_on_cong'*:
assumes *eventually* $(\lambda z. f\ z = g\ z)$ (*cosparses A*) $A = B$
shows f *meromorphic_on A* \longleftrightarrow g *meromorphic_on B*

unfolding *assms(2)[symmetric]*
by (*rule meromorphic_on_cong eventually_cospars_imp_eventually_at assms*)+
auto

8.11 Meromorphic functions and zorder

lemma *zorder_power_int*:
assumes *f meromorphic_on {z} frequently ($\lambda z. f z \neq 0$) (at z)*
shows *zorder ($\lambda z. f z \text{ powi } n$) z = n * zorder f z*
proof –
from *assms(1)* **obtain** *L* **where** *L: ($\lambda w. f (z + w)$) has_laurent_expansion L*
by (*auto simp: meromorphic_on_def*)
from *assms(2)* **and** *L* **have** [*simp*]: *L \neq 0*
by (*metis assms(1) has_laurent_expansion_eventually_nonzero_iff meromorphic_at_iff*
not_essential_frequently_0_imp_eventually_0 not_eventually_not_frequently)
from *L* **have** *L': ($\lambda w. f (z + w) \text{ powi } n$) has_laurent_expansion L powi n*
by (*intro laurent_expansion_intros*)
have *zorder f z = fls_subdegree L*
using *L assms(2) $\langle L \neq 0 \rangle$* **by** (*simp add: has_laurent_expansion_zorder*)
moreover **have** *zorder ($\lambda z. f z \text{ powi } n$) z = fls_subdegree (L powi n)*
using *L' assms(2) $\langle L \neq 0 \rangle$* **by** (*simp add: has_laurent_expansion_zorder*)
moreover **have** *fls_subdegree (L powi n) = n * fls_subdegree L*
by *simp*
ultimately show *?thesis*
by *simp*
qed

lemma *zorder_power*:
assumes *f meromorphic_on {z} frequently ($\lambda z. f z \neq 0$) (at z)*
shows *zorder ($\lambda z. f z ^ n$) z = n * zorder f z*
using *zorder_power_int[OF assms, of int n]* **by** *simp*

lemma *zorder_add1*:
assumes *f meromorphic_on {z} frequently ($\lambda z. f z \neq 0$) (at z)*
assumes *g meromorphic_on {z} frequently ($\lambda z. g z \neq 0$) (at z)*
assumes *zorder f z < zorder g z*
shows *zorder ($\lambda z. f z + g z$) z = zorder f z*
proof –
from *assms(1)* **obtain** *F* **where** *F: ($\lambda w. f (z + w)$) has_laurent_expansion F*
by (*auto simp: meromorphic_on_def*)
from *assms(3)* **obtain** *G* **where** *G: ($\lambda w. g (z + w)$) has_laurent_expansion G*
by (*auto simp: meromorphic_on_def*)
have [*simp*]: *F \neq 0 G \neq 0*
by (*metis assms has_laurent_expansion_eventually_nonzero_iff meromorphic_at_iff*
not_essential_frequently_0_imp_eventually_0 not_eventually_not_frequently
F G)+
have *: *zorder f z = fls_subdegree F zorder g z = fls_subdegree G*

```

    using F G assms by (simp_all add: has_laurent_expansion_zorder)
  from assms * have F ≠ -G
    by auto
  hence [simp]: F + G ≠ 0
    by (simp add: add_eq_0_iff2)
  moreover have zorder (λz. f z + g z) z = fls_subdegree (F + G)
    using has_laurent_expansion_zorder[OF has_laurent_expansion_add[OF F
G]] ⟨F ≠ -G⟩ by simp
  moreover have fls_subdegree (F + G) = fls_subdegree F
    using assms by (simp add: * fls_subdegree_add_eq1)
  ultimately show ?thesis
    by (simp add: *)
qed

```

lemma zorder_add2:

```

  assumes f meromorphic_on {z} frequently (λz. f z ≠ 0) (at z)
  assumes g meromorphic_on {z} frequently (λz. g z ≠ 0) (at z)
  assumes zorder f z > zorder g z
  shows zorder (λz. f z + g z) z = zorder g z
  using zorder_add1[OF assms(3,4) assms(1,2)] assms(5-) by (simp add: add commute)

```

lemma zorder_add_ge:

```

  fixes f g :: complex ⇒ complex
  assumes f meromorphic_on {z} frequently (λz. f z ≠ 0) (at z)
  assumes g meromorphic_on {z} frequently (λz. g z ≠ 0) (at z)
  assumes frequently (λz. f z + g z ≠ 0) (at z) zorder f z ≥ c zorder g z ≥ c
  shows zorder (λz. f z + g z) z ≥ c
proof -
  from assms(1) obtain F where F: (λw. f (z + w)) has_laurent_expansion F
    by (auto simp: meromorphic_on_def)
  from assms(3) obtain G where G: (λw. g (z + w)) has_laurent_expansion G
    by (auto simp: meromorphic_on_def)
  have [simp]: F ≠ 0 G ≠ 0
    using assms F G has_laurent_expansion_frequently_nonzero_iff by blast+
  have FG: (λw. f (z + w) + g (z + w)) has_laurent_expansion F + G
    by (intro laurent_expansion_intros F G)
  have [simp]: F + G ≠ 0
    using assms(5) has_laurent_expansion_frequently_nonzero_iff[OF FG] by
blast

```

```

  have *: zorder f z = fls_subdegree F zorder g z = fls_subdegree G
    zorder (λz. f z + g z) z = fls_subdegree (F + G)
  using F G FG has_laurent_expansion_zorder by simp_all
  moreover have zorder (λz. f z + g z) z = fls_subdegree (F + G)
    using has_laurent_expansion_zorder[OF has_laurent_expansion_add[OF F
G]] by simp
  moreover have fls_subdegree (F + G) ≥ min (fls_subdegree F) (fls_subdegree
G)

```

```

    by (intro fls_plus_subdegree) simp
  ultimately show ?thesis
    using assms(6,7) unfolding * by linarith
qed

```

```

lemma zorder_diff_ge:
  fixes f g :: complex  $\Rightarrow$  complex
  assumes f meromorphic_on {z} frequently ( $\lambda z. f z \neq 0$ ) (at z)
  assumes g meromorphic_on {z} frequently ( $\lambda z. g z \neq 0$ ) (at z)
  assumes frequently ( $\lambda z. f z \neq g z$ ) (at z) zorder f z  $\geq$  c zorder g z  $\geq$  c
  shows zorder ( $\lambda z. f z - g z$ ) z  $\geq$  c
proof -
  have ( $\lambda z. -g z$ ) meromorphic_on {z}
    by (auto intro: meromorphic_intros assms)
  thus ?thesis
    using zorder_add_ge[of f z  $\lambda z. -g z$  c] assms by simp
qed

```

```

lemma zorder_diff1:
  assumes f meromorphic_on {z} frequently ( $\lambda z. f z \neq 0$ ) (at z)
  assumes g meromorphic_on {z} frequently ( $\lambda z. g z \neq 0$ ) (at z)
  assumes zorder f z < zorder g z
  shows zorder ( $\lambda z. f z - g z$ ) z = zorder f z
proof -
  have zorder ( $\lambda z. f z + (-g z)$ ) z = zorder f z
    by (intro zorder_add1 meromorphic_intros assms) (use assms in auto)
  thus ?thesis
    by simp
qed

```

```

lemma zorder_diff2:
  assumes f meromorphic_on {z} frequently ( $\lambda z. f z \neq 0$ ) (at z)
  assumes g meromorphic_on {z} frequently ( $\lambda z. g z \neq 0$ ) (at z)
  assumes zorder f z > zorder g z
  shows zorder ( $\lambda z. f z - g z$ ) z = zorder g z
proof -
  have zorder ( $\lambda z. f z + (-g z)$ ) z = zorder ( $\lambda z. -g z$ ) z
    by (intro zorder_add2 meromorphic_intros assms) (use assms in auto)
  thus ?thesis
    by simp
qed

```

```

lemma zorder_mult:
  assumes f meromorphic_on {z} frequently ( $\lambda z. f z \neq 0$ ) (at z)
  assumes g meromorphic_on {z} frequently ( $\lambda z. g z \neq 0$ ) (at z)
  shows zorder ( $\lambda z. f z * g z$ ) z = zorder f z + zorder g z
proof -
  from assms(1) obtain F where F: ( $\lambda w. f (z + w)$ ) has_laurent_expansion F
    by (auto simp: meromorphic_on_def)

```

from *assms*(β) **obtain** G **where** $G: (\lambda w. g (z + w))$ *has_laurent_expansion* G
by (*auto simp: meromorphic_on_def*)
have [*simp*]: $F \neq 0 \ G \neq 0$
by (*metis assms has_laurent_expansion_eventually_nonzero_iff meromorphic_at_iff*
not_essential_frequently_0_imp_eventually_0 not_eventually_not_frequently
 $F \ G$)
have *: *zorder* $f \ z = \text{fls_subdegree } F$ *zorder* $g \ z = \text{fls_subdegree } G$
using $F \ G$ *assms* **by** (*simp_all add: has_laurent_expansion_zorder*)
moreover **have** *zorder* $(\lambda z. f \ z * g \ z) \ z = \text{fls_subdegree } (F * G)$
using *has_laurent_expansion_zorder*[*OF has_laurent_expansion_mult*[*OF F G*]] **by** *simp*
moreover **have** $\text{fls_subdegree } (F * G) = \text{fls_subdegree } F + \text{fls_subdegree } G$
using *assms* **by** *simp*
ultimately show *?thesis*
by (*simp add: **)
qed

lemma *zorder_divide*:

assumes f *meromorphic_on* $\{z\}$ *frequently* $(\lambda z. f \ z \neq 0)$ (*at z*)
assumes g *meromorphic_on* $\{z\}$ *frequently* $(\lambda z. g \ z \neq 0)$ (*at z*)
shows *zorder* $(\lambda z. f \ z / g \ z) \ z = \text{zorder } f \ z - \text{zorder } g \ z$
proof –
from *assms*(1) **obtain** F **where** $F: (\lambda w. f (z + w))$ *has_laurent_expansion* F
by (*auto simp: meromorphic_on_def*)
from *assms*(β) **obtain** G **where** $G: (\lambda w. g (z + w))$ *has_laurent_expansion* G
by (*auto simp: meromorphic_on_def*)
have [*simp*]: $F \neq 0 \ G \neq 0$
by (*metis assms has_laurent_expansion_eventually_nonzero_iff meromorphic_at_iff*
not_essential_frequently_0_imp_eventually_0 not_eventually_not_frequently
 $F \ G$)
have *: *zorder* $f \ z = \text{fls_subdegree } F$ *zorder* $g \ z = \text{fls_subdegree } G$
using $F \ G$ *assms* **by** (*simp_all add: has_laurent_expansion_zorder*)
moreover **have** *zorder* $(\lambda z. f \ z / g \ z) \ z = \text{fls_subdegree } (F / G)$
using *has_laurent_expansion_zorder*[*OF has_laurent_expansion_divide*[*OF F G*]] **by** *simp*
moreover **have** $\text{fls_subdegree } (F / G) = \text{fls_subdegree } F - \text{fls_subdegree } G$
using *assms* **by** (*simp add: fls_divide_subdegree*)
ultimately show *?thesis*
by (*simp add: **)
qed

lemma *constant_on_extend_nicely_meromorphic_on*:

assumes f *nicely_meromorphic_on* B f *constant_on* A
assumes *open* A *open* B *connected* $B \ A \neq \{\}$ $A \subseteq B$
shows f *constant_on* B
proof –
from *assms* **obtain** c **where** $c: \bigwedge z. z \in A \implies f \ z = c$


```

  by (auto simp: constant_on_def)
  have eventually ( $\lambda z. z \in A$ ) (cosparse A)
  by (intro eventually_in_cosparse assms order.refl)
  hence eventually ( $\lambda z. f z = c$ ) (cosparse A)
  by eventually_elim (use c in auto)
  hence freq: frequently ( $\lambda z. f z = c$ ) (cosparse A)
  by (intro eventually_frequently) (use assms in auto)
  then obtain z0 where z0:  $z0 \in A$  frequently ( $\lambda z. f z = c$ ) (at z0)
  using assms by (auto simp: frequently_def eventually_cosparse_open_eq)

  have  $f z = c$  if  $z \in B$  for  $z$ 
  proof (rule frequently_eq_meromorphic_imp_constant[OF _ assms(1)])
    show  $z0 \in B$  frequently ( $\lambda z. f z = c$ ) (at z0)
    using z0 assms by auto
  qed (use assms that in auto)
  thus f_constant_on B
  by (auto simp: constant_on_def)
qed

```

8.12 More on poles and zeros

lemma zorder_prod:

```

  assumes  $\bigwedge x. x \in A \implies f x$  meromorphic_on  $\{z\}$ 
  assumes eventually ( $\lambda z. (\prod_{x \in A}. f x z) \neq 0$ ) (at z)
  shows zorder ( $\lambda z. \prod_{x \in A}. f x z$ )  $z = (\sum_{x \in A}. zorder (f x) z)$ 
  using assms
proof (induction A rule: infinite_finite_induct)
  case (insert a A)
  have zorder ( $\lambda z. \prod_{x \in \text{insert } a A}. f x z$ )  $z = zorder (\lambda z. f a z * (\prod_{x \in A}. f x z)) z$ 
  using insert.hyps by simp
  also have  $\dots = zorder (f a) z + zorder (\lambda z. \prod_{x \in A}. f x z) z$ 
  proof (subst zorder_mult)
    have  $\forall_F z$  in at  $z. f a z \neq 0$ 
    using insert.premis(2) by eventually_elim (use insert.hyps in auto)
    thus  $\exists_F z$  in at  $z. f a z \neq 0$ 
    using eventually_frequently_at_neq_bot by blast
  next
    have  $\forall_F z$  in at  $z. (\prod_{x \in A}. f x z) \neq 0$ 
    using insert.premis(2) by eventually_elim (use insert.hyps in auto)
    thus  $\exists_F z$  in at  $z. (\prod_{x \in A}. f x z) \neq 0$ 
    using eventually_frequently_at_neq_bot by blast
  qed (use insert.premis in <auto intro!: meromorphic_intros>)
  also have zorder ( $\lambda z. \prod_{x \in A}. f x z$ )  $z = (\sum_{x \in A}. zorder (f x) z)$ 
  by (intro insert.IH) (use insert.premis insert.hyps in <auto elim!: eventually_mono>)
  also have zorder (f a)  $z + \dots = (\sum_{x \in \text{insert } a A}. zorder (f x) z)$ 
  using insert.hyps by simp
  finally show ?case .
qed auto

```

lemma *zorder_scale*:

assumes f meromorphic_on $\{a * z\}$ $a \neq 0$
shows $zorder (\lambda w. f (a * w)) z = zorder f (a * z)$
proof (cases eventually $(\lambda z. f z = 0)$ (at $(a * z)$))
case True
hence ev : eventually $(\lambda z. f (a * z) = 0)$ (at z)
proof (rule eventually_compose_filterlim)
show filterlim $((*) a)$ (at $(a * z)$) (at z)
proof (rule filterlim_atI)
show $\forall_F x$ in at $z. a * x \neq a * z$
using eventually_neq_at_within[$of z z$] **by** eventually_elim (use $\langle a \neq 0 \rangle$)
in auto)
qed (auto intro!: tendsto_intros)
qed

have $zorder (\lambda w. f (a * w)) z = zorder (\lambda_. 0) z$
by (rule zorder_cong) (use ev **in** auto)
also have $\dots = zorder (\lambda_. 0) (a * z)$
by (simp add: zorder_shift')
also have $\dots = zorder f (a * z)$
by (rule zorder_cong) (use True **in** auto)
finally show ?thesis .

next

case False

define G **where** $G = fps_const a * fps_X$

have $zorder (f \circ (\lambda z. a * z)) z = zorder f (a * z) * int (subdegree G)$

proof (rule zorder_compose)

show isolated_singularity_at $f (a * z)$ not_essential $f (a * z)$

using $assms(1)$ **by** (auto simp: meromorphic_on_altdef)

next

have $(\lambda x. a * x)$ has_fps_expansion G

unfolding G_def **by** (intro fps_expansion_intros)

thus $(\lambda x. a * (z + x) - a * z)$ has_fps_expansion G

by (simp add: algebra_simps)

next

show $\forall_F w$ in at $(a * z). f w \neq 0$ **using** False

by (metis $assms(1)$ has_laurent_expansion_isolated has_laurent_expansion_not_essential meromorphic_on_def non_zero_neighbour not_eventually_singletonI)

qed (use $\langle a \neq 0 \rangle$ **in** $\langle auto simp: G_def \rangle$)

also have subdegree $G = 1$

using $\langle a \neq 0 \rangle$ **by** (simp add: G_def)

finally show ?thesis

by (simp add: o_def)

qed

lemma *zorder_uminus*:

assumes f meromorphic_on $\{-z\}$

shows $zorder (\lambda w. f (-w)) z = zorder f (-z)$

using *assms zorder_scale*[of $f - 1 z$] by *simp*

lemma *is_pole_deriv_iff*:

assumes *f meromorphic_on* $\{z\}$

shows *is_pole* (*deriv f*) $z \longleftrightarrow$ *is_pole f z*

proof –

from *assms* obtain *F* where *F*: $(\lambda w. f (z + w))$ *has_laurent_expansion F*
by (*auto simp: meromorphic_on_def*)

have *deriv* $(\lambda w. f (z + w))$ *has_laurent_expansion fls_deriv F*

using *F* by (*rule has_laurent_expansion_deriv*)

also have *deriv* $(\lambda w. f (z + w)) = (\lambda w. \text{deriv } f (z + w))$

by (*simp add: deriv_shift_0' add_ac o_def fun_eq_iff*)

finally have *F'*: $(\lambda w. \text{deriv } f (z + w))$ *has_laurent_expansion fls_deriv F* .

have *is_pole* (*deriv f*) $z \longleftrightarrow$ *fls_subdegree* (*fls_deriv F*) < 0

using *is_pole_fls_subdegree_iff*[*OF F'*] by *simp*

also have $\dots \longleftrightarrow$ *fls_subdegree F* < 0

using *fls_deriv_subdegree0 fls_subdegree_deriv linorder_less_linear* by *fast-force*

also have $\dots \longleftrightarrow$ *is_pole f z*

using *F* by (*simp add: has_laurent_expansion_imp_is_pole_iff*)

finally show *?thesis* .

qed

lemma *isolated_zero_remove_sings_iff* [*simp*]:

assumes *isolated_singularity_at f z*

shows *isolated_zero* (*remove_sings f*) $z \longleftrightarrow$ *isolated_zero f z*

proof –

have *: $(\forall_F x \text{ in at } z. \text{remove_sings } f x \neq 0) \longleftrightarrow (\forall_F x \text{ in at } z. f x \neq 0)$

proof

assume $(\forall_F x \text{ in at } z. f x \neq 0)$

thus $(\forall_F x \text{ in at } z. \text{remove_sings } f x \neq 0)$

using *eventually_remove_sings_eq_at*[*OF assms*]

by *eventually_elim auto*

next

assume $(\forall_F x \text{ in at } z. \text{remove_sings } f x \neq 0)$

thus $(\forall_F x \text{ in at } z. f x \neq 0)$

using *eventually_remove_sings_eq_at*[*OF assms*]

by *eventually_elim auto*

qed

show *?thesis*

unfolding *isolated_zero_def* using *assms ** by *simp*

qed

lemma *zorder_isolated_zero_pos*:

assumes *isolated_zero f z f analytic_on* $\{z\}$

shows *zorder f z* > 0

proof (*subst zorder_pos_iff'* [*OF assms(2)*])

show $f z = 0$

using *assms* by (*simp add: zero_isolated_zero_analytic*)

next

have $\forall_F z \text{ in } at\ z. f\ z \neq 0$
using *assms* **by** (*auto simp: isolated_zero_def*)
thus $\exists_F z \text{ in } at\ z. f\ z \neq 0$
by (*simp add: eventually_frequently*)

qed

lemma *zorder_isolated_zero_pos'*:

assumes *isolated_zero f z isolated_singularity_at f z*
shows $zorder\ f\ z > 0$

proof –

from *assms(1)* **have** $f\ -z \rightarrow 0$
by (*simp add: isolated_zero_def*)
with *assms(2)* **have** *remove_sings f analytic_on {z}*
by (*intro remove_sings_analytic_at*)
hence $zorder\ (remove_sings\ f)\ z > 0$
using *assms* **by** (*intro zorder_isolated_zero_pos*) *auto*
thus *?thesis*
using *assms* **by** *simp*

qed

lemma *zero_isolated_zero_nicely_meromorphic*:

assumes *isolated_zero f z f nicely_meromorphic_on {z}*
shows $f\ z = 0$

proof –

have $\neg is_pole\ f\ z$
using *assms pole_is_not_zero* **by** *blast*
with *assms(2)* **have** *f analytic_on {z}*
by (*simp add: nicely_meromorphic_on_imp_analytic_at*)
with *zero_isolated_zero_analytic assms(1)* **show** *?thesis*
by *blast*

qed

lemma *meromorphic_on_imp_not_zero_cosparse*:

assumes *f meromorphic_on A*
shows *eventually* $(\lambda z. \neg is_pole\ f\ z)$ *(cosparse A)*

proof –

have *eventually* $(\lambda z. \neg is_pole\ (\lambda z. inverse\ (f\ z))\ z)$ *(cosparse A)*
by (*intro meromorphic_on_imp_not_pole_cosparse meromorphic_intros assms*)
thus *?thesis*
by (*simp add: is_pole_inverse_iff*)

qed

lemma *nicely_meromorphic_on_inverse [meromorphic_intros]*:

assumes *f nicely_meromorphic_on A*
shows $(\lambda x. inverse\ (f\ x))$ *nicely_meromorphic_on A*
unfolding *nicely_meromorphic_on_def*

proof (*intro conjI ballI*)

fix *z* **assume** $z: z \in A$

```

have is_pole  $f z \wedge f z = 0 \vee f -z \rightarrow f z$ 
  using assms z by (auto simp: nicely_meromorphic_on_def)
thus is_pole  $(\lambda x. \text{inverse } (f x)) z \wedge \text{inverse } (f z) = 0 \vee$ 
   $(\lambda x. \text{inverse } (f x)) -z \rightarrow \text{inverse } (f z)$ 
proof
  assume is_pole  $f z \wedge f z = 0$ 
  hence isolated_zero  $(\lambda z. \text{inverse } (f z)) z \wedge \text{inverse } (f z) = 0$ 
    by (auto simp: isolated_zero_inverse_iff)
  hence  $(\lambda x. \text{inverse } (f x)) -z \rightarrow \text{inverse } (f z)$ 
    by (simp add: isolated_zero_def)
  thus ?thesis ..
next
  assume lim:  $f -z \rightarrow f z$ 
  hence ana: f analytic_on {z}
    using assms is_pole_def nicely_meromorphic_on_imp_analytic_at
      not_tendsto_and_filterlim_at_infinity trivial_limit_at z by blast
  show ?thesis
  proof (cases isolated_zero f z)
    case True
      with lim have  $f z = 0$ 
        using continuous_within_zero_isolated_zero by blast
      with True have is_pole  $(\lambda z. \text{inverse } (f z)) z \wedge \text{inverse } (f z) = 0$ 
        by (auto simp: is_pole_inverse_iff)
      thus ?thesis ..
    next
      case False
      hence  $f z \neq 0 \vee (f z = 0 \wedge \text{eventually } (\lambda z. f z = 0) \text{ (at } z))$ 
        using non_isolated_zero_imp_eventually_zero[OF ana] by blast
      thus ?thesis
      proof (elim disjE conjE)
        assume  $f z \neq 0$ 
        hence  $(\lambda z. \text{inverse } (f z)) -z \rightarrow \text{inverse } (f z)$ 
          by (intro tendsto_intros lim)
        thus ?thesis ..
      next
        assume *:  $f z = 0 \text{ eventually } (\lambda z. f z = 0) \text{ (at } z)$ 
        have eventually  $(\lambda z. \text{inverse } (f z) = 0) \text{ (at } z)$ 
          using *(2) by eventually_elim auto
        hence  $(\lambda z. \text{inverse } (f z)) -z \rightarrow 0$ 
          by (simp add: tendsto_eventually)
        with *(1) show ?thesis
          by auto
      qed
    qed
  qed
qed (use assms in  $\langle \text{auto simp: nicely_meromorphic_on_def intro!: meromorphic_intros} \rangle$ )

lemma is_pole_zero_at_nicely_mero:
  assumes f nicely_meromorphic_on A is_pole f z z  $z \in A$ 

```

```

shows  $f z = 0$ 
by (meson assms at_neq_bot
      is_pole_def nicely_meromorphic_on_def
      not_tendsto_and_filterlim_at_infinity)

lemma zero_or_pole:
  assumes mero:  $f$  nicely_meromorphic_on  $A$ 
    and  $z \in A$   $f z = 0$  and event:  $\forall_F x$  in at  $z$ .  $f x \neq 0$ 
  shows isolated_zero  $f z \vee$  is_pole  $f z$ 
proof -
  from mero  $\langle z \in A \rangle$ 
  have (is_pole  $f z \wedge f z = 0$ )  $\vee f -z \rightarrow f z$ 
    unfolding nicely_meromorphic_on_def by simp
  moreover have isolated_zero  $f z$  if  $f -z \rightarrow f z$ 
    unfolding isolated_zero_def
    using  $\langle f z = 0 \rangle$  that event tendsto_nhds_iff by auto
  ultimately show ?thesis by auto
qed

lemma isolated_zero_fls_subdegree_iff:
  assumes  $(\lambda x. f (z + x))$  has_laurent_expansion  $F$ 
  shows isolated_zero  $f z \iff$  fls_subdegree  $F > 0$ 
  using assms unfolding isolated_zero_def
  by (metis Lim_at_zero_fls_zero_subdegree has_laurent_expansion_eventually_nonzero_iff
    not_le
    order.refl tendsto_0_subdegree_iff_0)

lemma zorder_pos_imp_isolated_zero:
  assumes  $f$  meromorphic_on  $\{z\}$  eventually  $(\lambda z. f z \neq 0)$  (at  $z$ ) zorder  $f z > 0$ 
  shows isolated_zero  $f z$ 
  using assms isolated_zero_fls_subdegree_iff
  by (metis has_laurent_expansion_eventually_nonzero_iff
    has_laurent_expansion_zorder insertI1
    meromorphic_on_def)

lemma zorder_neg_imp_is_pole:
  assumes  $f$  meromorphic_on  $\{z\}$  eventually  $(\lambda z. f z \neq 0)$  (at  $z$ ) zorder  $f z < 0$ 
  shows is_pole  $f z$ 
  using assms is_pole_fls_subdegree_iff at_neq_bot eventually_frequently_meromorphic_at_iff
    neg_zorder_imp_is_pole by blast

lemma not_pole_not_isolated_zero_imp_zorder_eq_0:
  assumes  $f$  meromorphic_on  $\{z\}$   $\neg$ is_pole  $f z$   $\neg$ isolated_zero  $f z$  frequently  $(\lambda z. f z \neq 0)$  (at  $z$ )
  shows zorder  $f z = 0$ 
proof -
  have remove_sings  $f$  analytic_on  $\{z\}$ 
    using assms meromorphic_at_iff not_essential_def remove_sings_analytic_at

```

```

by blast
  moreover from this and assms have remove_sings f z  $\neq$  0
  using isolated_zero_def meromorphic_at_iff non_zero_neighbour remove_sings_eq_0_iff
by blast
  moreover have frequently ( $\lambda z$ . remove_sings f z  $\neq$  0) (at z)
  using assms analytic_at_neq_imp_eventually_neq calculation(1,2)
  eventually_frequently trivial_limit_at by blast
  ultimately have zorder (remove_sings f) z = 0
  using zorder_eq_0_iff by blast
  thus ?thesis
  using assms(1) meromorphic_at_iff by auto
qed

lemma not_essential_compose:
  assumes not_essential f (g z) g analytic_on {z}
  shows not_essential ( $\lambda x$ . f (g x)) z
proof (cases isolated_zero ( $\lambda w$ . g w - g z) z)
  case False
  hence eventually ( $\lambda w$ . g w - g z = 0) (nhds z)
  by (intro non_isolated_zero_imp_eventually_zero' analytic_intros assms) auto
  hence not_essential ( $\lambda x$ . f (g x)) z  $\longleftrightarrow$  not_essential ( $\lambda \_$ . f (g z)) z
  by (intro not_essential_cong refl)
  (auto elim!: eventually_mono simp: eventually_at_filter)
  thus ?thesis
  by (simp add: not_essential_const)
next
  case True
  hence ev: eventually ( $\lambda w$ . g w  $\neq$  g z) (at z)
  by (auto simp: isolated_zero_def)
  from assms consider c where f -g z  $\rightarrow$  c | is_pole f (g z)
  by (auto simp: not_essential_def)
  have isCont g z
  by (rule analytic_at_imp_isCont) fact
  hence lim: g -z  $\rightarrow$  g z
  using isContD by blast

  from assms(1) consider c where f -g z  $\rightarrow$  c | is_pole f (g z)
  unfolding not_essential_def by blast
  thus ?thesis
proof cases
  fix c assume f -g z  $\rightarrow$  c
  hence ( $\lambda x$ . f (g x)) -z  $\rightarrow$  c
  by (rule filterlim_compose) (use lim ev in <auto simp: filterlim_at>)
  thus ?thesis
  by (auto simp: not_essential_def)
next
  assume is_pole f (g z)
  hence is_pole ( $\lambda x$ . f (g x)) z
  by (rule is_pole_compose) fact+

```

```

thus ?thesis
  by (auto simp: not_essential_def)
qed
qed

lemma isolated_singularity_at_compose:
  assumes isolated_singularity_at f (g z) g analytic_on {z}
  shows isolated_singularity_at ( $\lambda x. f (g x)$ ) z
proof (cases isolated_zero ( $\lambda w. g w - g z$ ) z)
  case False
  hence eventually ( $\lambda w. g w - g z = 0$ ) (nhds z)
    by (intro non_isolated_zero_imp_eventually_zero') (use assms in ⟨auto intro!:
analytic_intros⟩)
  hence isolated_singularity_at ( $\lambda x. f (g x)$ ) z  $\longleftrightarrow$  isolated_singularity_at ( $\lambda_.$ 
f (g z)) z
    by (intro isolated_singularity_at_cong_refl)
      (auto elim!: eventually_mono simp: eventually_at_filter)
  thus ?thesis
    by (simp add: isolated_singularity_at_const)
next
  case True
  from assms(1) obtain r where r: r > 0 f analytic_on ball (g z) r - {g z}
    by (auto simp: isolated_singularity_at_def)
  hence holo_f: f holomorphic_on ball (g z) r - {g z}
    by (subst (asm) analytic_on_open) auto
  from assms(2) obtain r' where r': r' > 0 g holomorphic_on ball z r'
    by (auto simp: analytic_on_def)

  have continuous_on (ball z r') g
    using holomorphic_on_imp_continuous_on r' by blast
  hence isCont g z
    using r' by (subst (asm) continuous_on_eq_continuous_at) auto
  hence g - z  $\rightarrow$  g z
    using isContD by blast
  hence eventually ( $\lambda w. g w \in \text{ball } (g z) r$ ) (at z)
    using ⟨r > 0⟩ unfolding tendsto_def by force
  moreover have eventually ( $\lambda w. g w \neq g z$ ) (at z) using True
    by (auto simp: isolated_zero_def elim!: eventually_mono)
  ultimately have eventually ( $\lambda w. g w \in \text{ball } (g z) r - \{g z\}$ ) (at z)
    by eventually_elim auto
  then obtain r'' where r'': r'' > 0  $\forall w \in \text{ball } z r'' - \{z\}. g w \in \text{ball } (g z) r - \{g$ 
z}
    unfolding eventually_at_filter eventually_nhds_metric ball_def
    by (auto simp: dist_commute)
  have f  $\circ$  g holomorphic_on ball z (min r' r'') - {z}
proof (rule holomorphic_on_compose_gen)
  show g holomorphic_on ball z (min r' r'') - {z}
    by (rule holomorphic_on_subset[OF r'(2)]) auto

```



```

show  $f$  holomorphic_on ball (g z) r - {g z}
  by fact
show  $g^{-1}(\text{ball } z (\min r' r'') - \{z\}) \subseteq \text{ball } (g z) r - \{g z\}$ 
  using  $r''$  by force
qed
hence  $f \circ g$  analytic_on ball z (min r' r'') - {z}
  by (subst analytic_on_open) auto
thus ?thesis using  $\langle r' > 0 \rangle \langle r'' > 0 \rangle$ 
  by (auto simp: isolated_singularity_at_def o_def intro!: exI[of _ min r' r''])
qed

```

```

lemma is_pole_power_int_0:
  assumes  $f$  analytic_on {x} isolated_zero f x n < 0
  shows is_pole ( $\lambda x. f x \text{ powi } n$ ) x
proof -
  have  $f -x \rightarrow f x$ 
    using  $\text{assms}(1)$  by (simp add: analytic_at_imp_isCont isContD)
  with  $\text{assms}$  show ?thesis
    unfolding is_pole_def
    by (intro filterlim_power_int_neg_at_infinity) (auto simp: isolated_zero_def)
qed

```

```

lemma isolated_zero_imp_not_constant_on:
  fixes  $f :: 'a :: \text{perfect\_space} \Rightarrow 'b :: \text{real\_normed\_div\_algebra}$ 
  assumes isolated_zero f x x  $\in A$  open A
  shows  $\neg f$  constant_on A
proof
  assume  $f$  constant_on A
  then obtain c where  $c: \bigwedge x. x \in A \implies f x = c$ 
    by (auto simp: constant_on_def)
  have eventually ( $\lambda z. z \in A - \{x\}$ ) (at x)
    by (intro eventually_at_in_open  $\text{assms}$ )
  hence eventually ( $\lambda z. f z = c$ ) (at x)
    by eventually_elim (use c in auto)
  hence  $f -x \rightarrow c$ 
    using tendsto_eventually by blast
  moreover from  $\text{assms}$  have  $f -x \rightarrow 0$ 
    by (simp add: isolated_zero_def)
  ultimately have [simp]:  $c = 0$ 
    using tendsto_unique[of at x f c 0] by (simp add: at_neq_bot)

  have eventually ( $\lambda x. f x \neq 0$ ) (at x)
    using  $\text{assms}$  by (auto simp: isolated_zero_def)
  moreover have eventually ( $\lambda x. x \in A$ ) (at x)
    using  $\text{assms}$  by (intro eventually_at_in_open') auto
  ultimately have eventually ( $\lambda x. \text{False}$ ) (at x)
    by eventually_elim (use c in auto)
  thus False
    by simp

```

qed

end

9 The Weierstraß Factorisation Theorem

```
theory Weierstrass_Factorization
  imports Meromorphic
begin
```

9.1 The elementary factors

The Weierstraß elementary factors are the family of entire functions

$$E_n(z) = (1 - z) \exp\left(z + \frac{z^2}{2} + \dots + \frac{z^n}{n}\right)$$

with the key properties that they have a single zero at $z = 1$ and satisfy $E_n(z) = 1 + O(z^{n+1})$ around the origin.

definition *weierstrass_factor* :: *nat* \Rightarrow *complex* \Rightarrow *complex* **where**
weierstrass_factor *n* *z* = $(1 - z) * \exp(\sum_{k=1..n}. z^k / \text{of_nat } k)$

lemma *weierstrass_factor_continuous_on* [*continuous_intros*]:
continuous_on *A* *f* \implies *continuous_on* *A* $(\lambda z. \text{weierstrass_factor } n (f z))$
by (*auto simp: weierstrass_factor_def intro!: continuous_intros*)

lemma *weierstrass_factor_holomorphic* [*holomorphic_intros*]:
f *holomorphic_on* *A* \implies $(\lambda z. \text{weierstrass_factor } n (f z))$ *holomorphic_on* *A*
by (*auto simp: weierstrass_factor_def intro!: holomorphic_intros*)

lemma *weierstrass_factor_analytic* [*analytic_intros*]:
f *analytic_on* *A* \implies $(\lambda z. \text{weierstrass_factor } n (f z))$ *analytic_on* *A*
by (*auto simp: weierstrass_factor_def intro!: analytic_intros*)

lemma *weierstrass_factor_0* [*simp*]: *weierstrass_factor* *n* 0 = 1
by (*auto simp: weierstrass_factor_def power_0_left*)

lemma *weierstrass_factor_1* [*simp*]: *weierstrass_factor* *n* 1 = 0
by (*simp add: weierstrass_factor_def*)

lemma *weierstrass_factor_eq_0_iff* [*simp*]: *weierstrass_factor* *n* *z* = 0 \iff *z* = 1
by (*simp add: weierstrass_factor_def*)

lemma *zorder_weierstrass_factor* [*simp*]: *zorder* (*weierstrass_factor* *n*) 1 = 1

proof (*rule zorder_eqI*)

show $(\lambda z. -\exp(\sum_{k=1..n}. z^k / \text{of_nat } k))$ *holomorphic_on* *UNIV*
by (*intro holomorphic_intros auto*)

qed (*auto simp: weierstrass_factor_def algebra_simps*)

lemma

fixes $z :: 'a :: \{\text{banach, real_normed_field}\}$
assumes $\text{norm } z \leq 1 / 2$
shows $\text{norm_exp_bounds_lemma: norm } (\text{exp } z - 1 - z) \leq \text{norm } z / 2$
and $\text{norm_exp_bounds: norm } (\text{exp } z - 1) \geq 1 / 2 * \text{norm } z$
 $\text{norm } (\text{exp } z - 1) \leq 3 / 2 * \text{norm } z$

proof –
show $*$: $\text{norm } (\text{exp } z - 1 - z) \leq \text{norm } z / 2$
proof (*cases* $z = 0$)
case *False*
have $\text{sums: } (\lambda k. z ^ (k + 2) /_R \text{fact } (k + 2)) \text{ sums } (\text{exp } z - (\sum_{k < 2}. z ^ k /_R \text{fact } k))$
by (*intro sums_split_initial_segment exp_converges*)

have $\text{summable } (\lambda k. \text{norm } z ^ (k + 2) /_R \text{fact } (k + 2))$
using $\text{summable_norm_exp[of } z]$
by (*intro summable_norm summable_ignore_initial_segment*)
(auto simp: norm_power norm_divide divide_simps)
also have $?this \longleftrightarrow \text{summable } (\lambda k. \text{norm } z ^ 2 * (\text{norm } z ^ k / \text{fact } (k + 2)))$
by (*simp add: power_add mult_ac divide_simps power2_eq_square del: of_nat_Suc of_nat_add*)
also have $\dots \longleftrightarrow \text{summable } (\lambda k. \text{norm } z ^ k / \text{fact } (k + 2))$
by (*subst summable_cmult_iff*) (*use* $\langle z \neq 0 \rangle$ **in** *auto*)
finally have $\text{summable: summable } (\lambda k. \text{norm } z ^ k / \text{fact } (k + 2))$.

have $\text{exp } z - 1 - z = (\sum k. z ^ (k + 2) / \text{fact } (k + 2))$
using $\text{sums by (simp add: sums_iff scaleR_conv_of_real divide_simps eval_nat_numeral)}$
also have $\text{norm } \dots \leq (\sum k. \text{norm } (z ^ (k + 2) / \text{fact } (k + 2)))$
using $\text{summable_norm_exp[of } z]$
by (*intro summable_norm summable_ignore_initial_segment*)
(auto simp: norm_power norm_divide divide_simps)
also have $\dots = (\sum k. \text{norm } z ^ 2 * (\text{norm } z ^ k / \text{fact } (k + 2)))$
by (*simp add: power_add norm_power norm_divide mult_ac norm_mult power2_eq_square del: of_nat_Suc*)
also have $\dots = \text{norm } z ^ 2 * (\sum k. \text{norm } z ^ k / \text{fact } (k + 2))$
using $\text{summable by (rule suminf_mult)}$
also have $\dots \leq \text{norm } z ^ 2 * (1 / (1 - \text{norm } z) / 2)$
proof (*intro mult_left_mono, rule sums_le*)
show $(\lambda k. \text{norm } z ^ k / \text{fact } (k + 2)) \text{ sums } (\sum k. \text{norm } z ^ k / \text{fact } (k + 2))$
using summable by blast
show $(\lambda k. \text{norm } z ^ k / 2) \text{ sums } (1 / (1 - \text{norm } z) / 2)$
using $\text{assms by (intro geometric_sums sums_divide) auto}$

next
fix $k :: \text{nat}$
have $\text{norm } z ^ k / \text{fact } (k + 2) \leq \text{norm } z ^ k / \text{fact } 2$
by (*intro divide_left_mono fact_mono*) *auto*

```

    thus norm z ^ k / fact (k + 2) ≤ norm z ^ k / 2
      by simp
  qed (auto simp: divide_simps)
  also have ... = norm z * (norm z / (1 - norm z)) / 2
    by (simp add: power2_eq_square)
  also have ... ≤ norm z * ((1 / 2) / (1 - (1 / 2))) / 2
    using assms by (intro mult_mono frac_le diff_mono) auto
  also have ... = norm z / 2
    by simp
  finally show norm (exp z - 1 - z) ≤ norm z / 2 .
qed auto

have norm (exp z - 1) ≤ norm z + norm (exp z - 1 - z)
  by (rule norm_triangle_sub)
with * show norm (exp z - 1) ≤ 3 / 2 * norm z
  by simp

have norm z - norm (exp z - 1 - z) ≤ norm (exp z - 1)
  using norm_triangle_ineq3[of exp z - 1 - z - z] by simp
with * show norm (exp z - 1) ≥ 1 / 2 * norm z
  by simp
qed

lemma weierstrass_factor_bound:
  assumes norm z ≤ 1 / 2
  shows norm (weierstrass_factor n z - 1) ≤ 3 * norm z ^ Suc n
proof (cases n = 0 ∨ z = 0)
  case True
  thus ?thesis
  proof
    assume n = 0
    thus ?thesis by (auto simp: weierstrass_factor_def)
  qed auto
next
  case False
  with assms have z ≠ 1 n > 0 z ≠ 0
    by auto

  have summable (λk. cmod z ^ (k + Suc n) / real (k + Suc n))
    using ln_series'[of -norm z] assms
    by (intro summable_norm summable_ignore_initial_segment)
      (simp_all add: sums_iff summable_minus_iff power_minus' norm_divide
norm_power)
  also have ?this ⟷ summable (λk. norm z ^ Suc n * (norm z ^ k / real (k +
Suc n)))
    by (simp add: power_add mult_ac)
  also have ... ⟷ summable (λk. norm z ^ k / real (k + Suc n))
    by (subst summable_cmult_iff) (use ⟨z ≠ 0⟩ in auto)
  finally have summable: summable (λk. norm z ^ k / real (k + Suc n)) .

```

```

have ( $\lambda k. z^k / \text{of\_nat } k$ ) sums -  $\text{Ln } (1 - z)$ 
  using sums_minus[OF Ln_series[of -z]] assms by (simp add: power_minus')
hence ( $\lambda k. z^{(k + \text{Suc } n)} / \text{of\_nat } (k + \text{Suc } n)$ ) sums (-  $\text{Ln } (1 - z)$  -
( $\sum_{k < \text{Suc } n}. z^k / \text{of\_nat } k$ ))
  by (intro sums_split_initial_segment)
also have ( $\sum_{k < \text{Suc } n}. z^k / \text{of\_nat } k$ ) = ( $\sum_{k=1..n}. z^k / \text{of\_nat } k$ )
  by (intro sum.mono_neutral_right) auto
finally have  $\text{norm } (\text{ln } (1 - z) + (\sum_{k=1..n}. z^k / \text{of\_nat } k)) =$ 
   $\text{norm } (\sum_{k. z^{(k + \text{Suc } n)} / \text{of\_nat } (k + \text{Suc } n)})$ 
  by (simp add: sums_iff norm_uminus_minus)

also have  $\dots \leq (\sum_{k. \text{norm } (z^{(k + \text{Suc } n)} / \text{of\_nat } (k + \text{Suc } n)))$ 
  using ln_series'[of -norm z] assms
  by (intro summable_norm summable_ignore_initial_segment)
  (simp_all add: sums_iff summable_minus_iff power_minus' norm_divide
norm_power)
also have  $\dots = (\sum_{k. \text{norm } z^{\text{Suc } n} * (\text{norm } z^k / \text{real } (k + \text{Suc } n)))$ 
  by (simp add: algebra_simps norm_mult norm_power norm_divide power_add
del: of_nat_add of_nat_Suc)
also have  $\dots = \text{norm } z^{\text{Suc } n} * (\sum_{k. \text{norm } z^k / \text{real } (k + \text{Suc } n)}$ 
  by (intro suminf_mult summable)
also have  $\dots \leq \text{norm } z^{\text{Suc } n} * (1 / (1 - \text{norm } z))$ 
proof (intro mult_left_mono[OF sums_le])
  show ( $\lambda k. \text{norm } z^k / \text{real } (k + \text{Suc } n)$ ) sums ( $\sum_{k. \text{norm } z^k / \text{real } (k +$ 
Suc } n))
  using summable by blast
  show ( $\lambda k. \text{norm } z^k$ ) sums ( $1 / (1 - \text{norm } z)$ )
  by (rule geometric_sums) (use assms in auto)
qed (auto simp: field_simps)
also have  $\text{norm } z^{\text{Suc } n} * (1 / (1 - \text{norm } z)) \leq \text{norm } z^{\text{Suc } n} * (1 / (1 -$ 
(1 / 2)))
  using assms by (intro mult_mono power_mono divide_left_mono diff_mono
mult_pos_pos) auto
also have  $\dots = 2 * \text{norm } z^{\text{Suc } n}$ 
  by simp
finally have norm_le:  $\text{norm } (\text{ln } (1 - z) + (\sum_{k=1..n}. z^k / \text{of\_nat } k)) \leq 2$ 
* norm } z^{\text{Suc } n} .

also have  $\dots \leq 2 * \text{norm } z^2$ 
  using  $\langle n > 0 \rangle$  assms by (intro mult_left_mono power_decreasing) auto
also have  $\dots \leq 2 * (1 / 2)^2$ 
  by (intro mult_left_mono assms power_mono) auto
finally have norm_le':  $\text{norm } (\text{ln } (1 - z) + (\sum_{k=1..n}. z^k / \text{of\_nat } k)) \leq 1$ 
/ 2
  by (simp add: power2_eq_square)

have weierstrass_factor  $n z = \text{exp } (\text{ln } (1 - z) + (\sum_{k=1..n}. z^k / \text{of\_nat } k))$ 
  using  $\langle z \neq 1 \rangle$  by (simp add: exp_add weierstrass_factor_def)

```

```

also have norm (... - 1) ≤ (3 / 2) * norm (ln (1 - z) + (∑ k=1..n. z ^ k /
of_nat k))
  by (intro norm_exp_bounds norm_le^')
also have ... ≤ (3 / 2) * (2 * norm z ^ Suc n)
  by (intro mult_left_mono norm_le) auto
finally show ?thesis
  by simp
qed

```

9.2 Infinite products of elementary factors

We now show that the elementary factors can be used to construct an entire function with its zeros at a certain set of points (given by a sequence of non-zero numbers a_n with no accumulation point).

```

locale weierstrass_product =
  fixes a :: nat ⇒ complex
  fixes p :: nat ⇒ nat
  assumes a_nonzero: ∧n. a n ≠ 0
  assumes filterlim_a: filterlim a at_infinity at_top
  assumes summable_a_p: ∧r. r > 0 ⇒ summable (λn. (r / norm (a n)) ^ Suc
(p n))
begin

```

```

definition f :: complex ⇒ complex where
  f z = (∏ n. weierstrass_factor (p n) (z / a n))

```

```

lemma abs_convergent: abs_convergent_prod (λn. weierstrass_factor (p n) (z /
a n))

```

```

  unfolding abs_convergent_prod_conv_summable
proof (rule summable_comparison_test_ev)
  have eventually (λn. norm (a n) > 2 * norm z) at_top
    using filterlim_a by (metis filterlim_at_infinity_imp_norm_at_top filter-
lim_at_top_dense)
  thus eventually (λn. norm (norm (weierstrass_factor (p n) (z / a n) - 1)) ≤
3 * norm (z / a n) ^ Suc (p n)) at_top
proof eventually_elim
  case (elim n)
  hence norm (z / a n) ≤ 1 / 2
    by (auto simp: norm_divide divide_simps)
  thus ?case using weierstrass_factor_bound[of z / a n p n]
    by simp
qed
next
  show summable (λn. 3 * norm (z / a n) ^ Suc (p n))
    using summable_mult[OF summable_a_p[of norm z], of 3]
    by (cases z = 0) (auto simp: norm_divide)
qed

```

```

lemma convergent: convergent_prod (λn. weierstrass_factor (p n) (z / a n))

```

using *abs_convergent*[of *z*] *abs_convergent_prod_imp_convergent_prod* **by** *blast*

lemma *has_prod*: $(\lambda n. \text{weierstrass_factor } (p \ n) \ (z / a \ n)) \text{ has_prod } f \ z$
using *convergent*[of *z*] **unfolding** *f_def* **by** *auto*

lemma *finite_occs_a*: *finite* $(a - \{z\})$

proof –

have *eventually* $(\lambda n. \text{norm } (a \ n) > \text{norm } z) \text{ at_top}$
using *filterlim_a* **by** $(\text{metis } \text{filterlim_at_infinity_imp_norm_at_top } \text{filterlim_at_top_dense})$
then obtain *N* **where** $N: \bigwedge n. n \geq N \implies \text{norm } (a \ n) > \text{norm } z$
by $(\text{auto } \text{simp: } \text{eventually_at_top_linorder})$
have $n < N$ **if** $n \in a - \{z\}$ **for** *n*
using $N[\text{of } n]$ **that** **by** $(\text{cases } n < N) \text{ auto}$
hence $a - \{z\} \subseteq \{..<N\}$ *finite* $\{..<N\}$
by *auto*
thus *?thesis*
using *finite_subset* **by** *blast*

qed

context

fixes *P*

defines $P \equiv (\lambda N \ z. \prod_{n < N. \text{weierstrass_factor } (p \ n) \ (z / a \ n)})$

begin

lemma *uniformly_convergent*:

assumes $R > 0$

shows *uniformly_convergent_on* $(\text{cball } 0 \ R) \ P$

unfolding *P_def*

proof $(\text{rule } \text{uniformly_convergent_on_prod'})$

show *uniformly_convergent_on* $(\text{cball } 0 \ R) \ (\lambda N \ z. \sum_{n < N. \text{norm } (\text{weierstrass_factor } (p \ n) \ (z / a \ n) - 1))$

proof $(\text{rule } \text{Weierstrass_m_test'_ev})$

have *eventually* $(\lambda n. \text{norm } (a \ n) \geq 2 * R) \text{ sequentially}$

using *filterlim_a* **by** $(\text{metis } \text{filterlim_at_infinity_imp_norm_at_top } \text{filterlim_at_top})$

thus $\forall_F n \text{ in } \text{sequentially}. \forall z \in \text{cball } 0 \ R. \text{norm } (\text{norm } (\text{weierstrass_factor } (p \ n) \ (z / a \ n) - 1)) \leq$

$3 * (R / \text{norm } (a \ n)) \wedge \text{Suc } (p \ n)$

proof *eventually_elim*

case $(\text{elim } n)$

show *?case*

proof *safe*

fix *z* :: *complex* **assume** $z: z \in \text{cball } 0 \ R$

have $2 * \text{norm } z \leq 2 * R$

using *z* **by** *auto*

also have $\dots \leq \text{norm } (a \ n)$

using *elim* **by** *simp*

finally have $\text{norm } (a \ n) \geq 2 * \text{norm } z .$

```

    hence norm (weierstrass_factor (p n) (z / a n) - 1) ≤ 3 * norm (z / a
n) ^ Suc (p n)
    by (intro weierstrass_factor_bound) (auto simp: norm_divide divide_simps)
    also have ... = 3 * (norm z / norm (a n)) ^ Suc (p n)
      by (simp add: norm_divide)
    also have ... ≤ 3 * (R / norm (a n)) ^ Suc (p n)
      by (intro mult_left_mono power_mono divide_right_mono) (use z in
auto)
    finally show norm (norm (weierstrass_factor (p n) (z / a n) - 1)) ≤
      3 * (R / norm (a n)) ^ Suc (p n) by simp
  qed
next
show summable (λn. 3 * (R / norm (a n)) ^ Suc (p n))
  by (intro summable_mult summable_a_p assms)
qed
qed (auto intro!: continuous_intros simp: a_nonzero)

```

lemma *uniform_limit*:

```

  assumes R > 0
  shows uniform_limit (cball 0 R) P f at_top
proof -
  obtain g where g: uniform_limit (cball 0 R) P g at_top
  using uniformly_convergent[OF assms] by (auto simp: uniformly_convergent_on_def)
  also have ?this ↔ uniform_limit (cball 0 R) P f at_top
  proof (intro uniform_limit_cong)
    fix z :: complex assume z ∈ cball 0 R
    with g have (λn. P (Suc n) z) ⟶ g z
      by (metis tendsto_uniform_limitI filterlim_sequentially_Suc)
    moreover have (λn. P (Suc n) z) ⟶ f z
      using convergent_prod_LIMSEQ[OF convergent[of z]] unfolding P_def
lessThan_Suc_atMost
      by (simp add: f_def)
    ultimately show g z = f z
      using tendsto_unique by force
  qed auto
  finally show ?thesis .
qed

```

lemma *holomorphic* [*holomorphic_intros*]: *f* holomorphic_on *A*

```

proof (rule holomorphic_on_subset)
  show f holomorphic_on UNIV
  proof (rule holomorphic_uniform_sequence)
    fix z :: complex
    have *: uniform_limit (cball 0 (norm z + 1)) P f sequentially
      by (rule uniform_limit) (auto intro: add_nonneg_pos)
    hence uniform_limit (cball z 1) P f sequentially
      by (rule uniform_limit_on_subset) (simp add: cball_subset_cball_iff)
    thus ∃ d > 0. cball z d ⊆ UNIV ∧ uniform_limit (cball z d) P f sequentially
  qed

```



```

    by (intro exI[of _ 1]) auto
  qed (auto intro!: holomorphic_intros simp: P_def)
qed auto

lemma analytic [analytic_intros]: f analytic_on A
  using holomorphic[of UNIV] analytic_on_holomorphic by blast

end

lemma zero: f z = 0  $\longleftrightarrow$  z  $\in$  range a
  using has_prod_eq_0_iff[OF has_prod, of z] by (auto simp: a_nonzero)

lemma not_islimpt_range_a:  $\neg$ z islimpt (range a)
proof
  assume z islimpt (range a)
  then obtain r :: nat  $\Rightarrow$  nat where r: strict_mono r (a  $\circ$  r)  $\longrightarrow$  z
    using islimpt_range_imp_convergent_subsequence by metis
  moreover have filterlim (a  $\circ$  r) at_infinity sequentially
    unfolding o_def by (rule filterlim_compose[OF filterlim_a filterlim_subseq[OF
r(1)]])
  ultimately show False
    by (meson not_tendsto_and_filterlim_at_infinity trivial_limit_sequentially)
qed

lemma isolated_zero:
  assumes z  $\in$  range a
  shows isolated_zero f z
proof -
  have eventually ( $\lambda$ z. f z  $\neq$  0) (at z)
    using not_islimpt_range_a[of z] by (auto simp: islimpt_iff_eventually zero)
  moreover have f  $-z \rightarrow$  f z
    by (intro isContD analytic_at_imp_isCont analytic)
  hence f  $-z \rightarrow$  0
    using assms zero[of z] by auto
  ultimately show ?thesis
    by (auto simp: isolated_zero_def)
qed

lemma zorder: zorder f z = card (a - $\{z\}$ )
proof -
  obtain N where N: a - $\{z\} \subseteq \{..N\}$ 
    using finite_occ_a[of z] by (meson finite_nat_iff_bounded_le)
  define g where g = ( $\lambda$ z n. weierstrass_factor (p n) (z / a n))
  define h1 where h1 = ( $\lambda$ w. ( $\prod$  n $\in\{..N\}$  - a- $\{z\}$ . g w n) * ( $\prod$  n. g w (n + Suc N)))
  define h2 where h2 = ( $\lambda$ w. ( $\prod$  n $\in\{..N\} \cap$  a- $\{z\}$ . g w n))

  have has_prod_h1': ( $\lambda$ n. g w (n + Suc N)) has_prod ( $\prod$  n. g w (n + Suc N))

```

```

for w
  unfolding g_def
  by (intro convergent_prod_has_prod convergent_prod_ignore_initial_segment
convergent)

  have f_eq: f w = h1 w * h2 w for w
  proof -
    have f w = ( $\prod n < \text{Suc } N. g w n$ ) * ( $\prod n. g w (n + \text{Suc } N)$ )
    proof (rule has_prod_unique2)
      show ( $\lambda n. g w n$ ) has_prod (( $\prod n < \text{Suc } N. g w n$ ) * ( $\prod n. g w (n + \text{Suc } N)$ ))
        unfolding g_def by (intro has_prod_ignore_initial_segment' convergent)
      show g w has_prod f w
        unfolding g_def by (rule has_prod)
    qed
    also have  $\{.. < \text{Suc } N\} = (\{..N\} - a-\{z\}) \cup (\{..N\} \cap a-\{z\})$ 
      by auto
    also have ( $\prod k \in \dots. g w k$ ) = ( $\prod k \in \{..N\} - a-\{z\}. g w k$ ) * ( $\prod k \in \{..N\} \cap$ 
 $a-\{z\}. g w k$ )
      by (intro prod.union_disjoint) auto
    finally show ?thesis
      by (simp add: h1_def h2_def mult_ac)
    qed

  have ana_h1: h1 analytic_on {z}
  proof -
    interpret h1: weierstrass_product  $\lambda n. a (n + \text{Suc } N) \lambda n. p (n + \text{Suc } N)$ 
    proof
      have filterlim ( $\lambda n. n + \text{Suc } N$ ) at_top at_top
        by (rule filterlim_add_const_nat_at_top)
      thus filterlim ( $\lambda n. a (n + \text{Suc } N)$ ) at_infinity at_top
        by (intro filterlim_compose[OF filterlim_a])
      show summable ( $\lambda n. (r / \text{cmod } (a (n + \text{Suc } N))) \wedge \text{Suc } (p (n + \text{Suc } N)))$ ) if
 $r > 0$  for r
        by (intro summable_ignore_initial_segment summable_a_p that)
      qed (auto simp: a_nonzero)

      show ?thesis using h1.analytic
        unfolding h1_def g_def h1.f_def by (intro analytic_intros) (auto simp:
 $a\_nonzero$ )
      qed

  have ana_h2: h2 analytic_on {z}
    unfolding h2_def g_def by (intro analytic_intros) (auto simp: a_nonzero)

  have zorder f z = zorder ( $\lambda w. h1 w * h2 w$ ) z
    by (simp add: f_eq [abs_def])
  also have  $\dots = \text{zorder } h1 z + \text{zorder } h2 z$ 
  proof (rule zorder_times_analytic)
    have eventually ( $\lambda w. f w \neq 0$ ) (at z)

```

```

    using not_islimpt_range_a[of z] by (auto simp: islimpt_conv_frequently_at
frequently_def zero)
    thus eventually ( $\lambda w. h1 w * h2 w \neq 0$ ) (at z)
      by (simp add: f_eq)
  qed fact+
  also have zorder h2 z = ( $\sum n \in \{..N\} \cap a^{-' \{z\}}. zorder (\lambda w. g w n) z$ )
    unfolding h2_def
    by (intro zorder_prod_analytic)
      (auto simp: a_nonzero g_def eventually_at_filter intro!: analytic_intros)
  also have h1 z  $\neq 0$  using N_has_prod_eq_0_iff[OF has_prod_h1 [of z]]
    by (auto simp: h1_def g_def)
  hence zorder h1 z = 0
    by (intro zorder_eq_0I ana_h1)
  also have ( $\sum n \in \{..N\} \cap a^{-' \{z\}}. zorder (\lambda w. g w n) z$ ) = ( $\sum n \in \{..N\} \cap a^{-' \{z\}}. 1$ )
  proof (intro sum.cong refl)
    fix n :: nat
    assume n:  $n \in \{..N\} \cap a^{-' \{z\}}$ 
    have zorder ( $\lambda w. weierstrass_factor (p n) (1 / a n * w)$ ) z =
      zorder (weierstrass_factor (p n)) (1 / a n * z)
      using a_nonzero[of n] eventually_neq_at_within[of 1 z / a n UNIV]
      by (intro zorder_scale analytic_intros analytic_on_imp_meromorphic_on)
  auto
  hence zorder ( $\lambda w. g w n$ ) z = zorder (weierstrass_factor (p n)) 1
    using n a_nonzero[of n] by (auto simp: g_def)
  thus zorder ( $\lambda w. g w n$ ) z = 1
    by simp
  qed
  also have ... = card ( $\{..N\} \cap a^{-' \{z\}}$ )
    by simp
  also have  $\{..N\} \cap a^{-' \{z\}} = a^{-' \{z\}}$ 
    using N by blast
  finally show ?thesis
    by simp
  qed
end

```

The following locale is the most common case of $p(n) = n$.

```

locale weierstrass_product' =
  fixes a :: nat  $\Rightarrow$  complex
  assumes a_nonzero:  $\bigwedge n. a n \neq 0$ 
  assumes filterlim_a: filterlim a at_infinity at_top
  assumes finite_occs_a':  $\bigwedge z. z \in \text{range } a \implies \text{finite } (a^{-' \{z\}})$ 
begin

lemma finite_occs_a: finite ( $a^{-' \{z\}}$ )
proof (cases  $z \in \text{range } a$ )
  case False

```

```

  hence a - ' {z} = {}
    by auto
  thus ?thesis by simp
qed (use finite_occs_a'[of z] in auto)

sublocale weierstrass_product a  $\lambda n. n$ 
proof
  fix r :: real assume r: r > 0
  show summable ( $\lambda n. (r / \text{norm } (a n)) \wedge \text{Suc } n$ )
  proof (rule summable_comparison_test_ev)
    have eventually ( $\lambda n. \text{norm } (a n) > 2 * r$ ) at_top
      using filterlim_a by (metis filterlim_at_infinity_imp_norm_at_top filterlim_at_top_dense)
    thus eventually ( $\lambda n. \text{norm } ((r / \text{norm } (a n)) \wedge \text{Suc } n) \leq (1 / 2) \wedge \text{Suc } n$ )
      at_top
    proof eventually_elim
      case (elim n)
      have norm ((r / norm (a n))  $\wedge$  Suc n) = (r / norm (a n))  $\wedge$  Suc n
        using <r > 0> by (simp add: abs_mult)
      also have ...  $\leq (1 / 2) \wedge \text{Suc } n$ 
        using <r > 0> elim by (intro power_mono) (auto simp: divide_simps)
      finally show ?case .
    qed
  next
    show summable ( $\lambda n. (1 / 2) \wedge \text{Suc } n :: \text{real}$ )
    unfolding summable_Suc_iff by (intro summable_geometric) auto
  qed
qed (use a_nonzero filterlim_a finite_occs_a in auto)

end

```

9.3 Writing a quotient as an exponential

If two holomorphic functions f and g on a simply connected domain have the same zeros with the same multiplicities, they can be written as $g(x) = e^{h(x)}f(x)$ for some holomorphic function $h(x)$.

lemma *holomorphic_zorder_factorization:*

assumes g holomorphic_on A open A connected A

assumes f holomorphic_on $A \wedge z. z \in A \implies f z = 0 \iff g z = 0$

$\wedge z. z \in A \implies \text{zorder } f z = \text{zorder } g z$

obtains h where h holomorphic_on $A \wedge z. z \in A \implies h z \neq 0 \wedge z. z \in A \implies g z = h z * f z$

proof (cases $\exists z \in A. g z \neq 0$)

case *False*

show ?thesis

by (rule that[of $\lambda_. 1$]) (use *False* assms in auto)

next

case *True*

define F where $F = \text{fps_expansion } f$

```

define G where  $G = \text{fps\_expansion } g$ 
define c where  $c = (\lambda z. \text{fps\_nth } (G \ z) \ (\text{subdegree } (G \ z)) \ / \ \text{fps\_nth } (F \ z) \ (\text{subdegree } (F \ z)))$ 
define h where  $h = (\lambda z. \text{if } f \ z = 0 \ \text{then } c \ z \ \text{else } g \ z \ / \ f \ z)$ 

have ev_nonzero: eventually  $(\lambda w. g \ w \neq 0 \wedge w \in A) \ (\text{at } z) \ \text{if } z \in A \ \text{for } z$ 
proof -
  from True obtain z0 where  $z0: z0 \in A \ g \ z0 \neq 0$ 
  by auto
  show ?thesis
  by  $(\text{rule } \text{non\_zero\_neighbour\_alt}[\text{where } ?\beta = z0])$ 
   $(\text{use } \text{assms } \text{that } z0 \ \text{in } \langle \text{auto } \text{intro: } \text{simply\_connected\_imp\_connected} \rangle)$ 
qed

have g_ana: g analytic_on  $\{z\} \ \text{if } z \in A \ \text{for } z$ 
  using assms  $\langle \text{open } A \rangle \ \text{analytic\_at } \text{that} \ \text{by } \text{blast}$ 
have f_ana: f analytic_on  $\{z\} \ \text{if } z \in A \ \text{for } z$ 
  using assms  $\langle \text{open } A \rangle \ \text{analytic\_at } \text{that} \ \text{by } \text{blast}$ 

have F:  $(\lambda w. f \ (z + w)) \ \text{has\_fps\_expansion } F \ z \ \text{if } z \in A \ \text{for } z$ 
  unfolding F_def by  $(\text{rule } \text{analytic\_at\_imp\_has\_fps\_expansion}[OF \ i_ana[OF \ \text{that}]])$ 
have G:  $(\lambda w. g \ (z + w)) \ \text{has\_fps\_expansion } G \ z \ \text{if } z \in A \ \text{for } z$ 
  unfolding G_def by  $(\text{rule } \text{analytic\_at\_imp\_has\_fps\_expansion}[OF \ i_g_ana[OF \ \text{that}]])$ 

have [simp]:  $G \ z \neq 0 \ \text{if } z \in A \ \text{for } z$ 
proof
  assume  $G \ z = 0$ 
  hence eventually  $(\lambda w. g \ w = 0) \ (\text{at } z) \ \text{using } G[OF \ \text{that}]$ 
  by  $(\text{auto } \text{simp: } \text{has\_fps\_expansion\_0\_iff } \text{at\_to\_0'} \ \text{eventually\_filtermap } \text{add\_ac} \ \text{eventually\_at\_filter } \text{nhds\_to\_0'} \ \text{elim: } \text{eventually\_mono})$ 
  hence eventually  $(\lambda_. \text{False}) \ (\text{at } z)$ 
  using ev_nonzero[OF that] unfolding eventually_at_filter by eventually_elim
auto
  thus False
  by simp
qed
have [simp]:  $F \ z \neq 0 \ \text{if } z \in A \ \text{for } z$ 
proof
  assume  $F \ z = 0$ 
  hence eventually  $(\lambda w. f \ w = 0) \ (\text{at } z) \ \text{using } F[\text{of } z] \ \text{that}$ 
  by  $(\text{auto } \text{simp: } \text{has\_fps\_expansion\_0\_iff } \text{at\_to\_0'} \ \text{eventually\_filtermap } \text{add\_ac} \ \text{eventually\_at\_filter } \text{nhds\_to\_0'} \ \text{elim: } \text{eventually\_mono})$ 
  hence eventually  $(\lambda_. \text{False}) \ (\text{at } z)$ 
  using ev_nonzero[OF that] unfolding eventually_at_filter
by eventually_elim  $(\text{use } \text{assms } \text{in } \text{auto})$ 

```

```

thus False
  by simp
qed
have [simp]:  $c z \neq 0$  if  $z \in A$  for  $z$ 
  using that by (simp add: c_def)

have h:  $h$  analytic_on  $\{z\} \wedge h z \neq 0$  if  $z \in A$  for  $z$ 
proof -
  show ?thesis
  proof (cases  $f z = 0$ )
    case False
    from False that have  $(\lambda z. g z / f z)$  analytic_on  $\{z\}$ 
      by (intro analytic_intros g_ana f_ana) auto
    also have ?this  $\longleftrightarrow$   $h$  analytic_on  $\{z\}$ 
    proof (rule analytic_at_cong)
      have eventually  $(\lambda w. f w \neq 0)$  (nhds  $z$ )
        using ev_nonzero[OF  $\langle z \in A \rangle$ ] unfolding eventually_at_filter
        by eventually_elim (use False  $\langle z \in A \rangle$  assms in auto)
      thus eventually  $(\lambda z. g z / f z = h z)$  (nhds  $z$ )
        by eventually_elim (auto simp: h_def)
    qed auto
    finally have  $h$  analytic_on  $\{z\}$  .
    moreover have  $h z \neq 0$ 
      using that assms by (simp add: h_def)
    ultimately show ?thesis by blast
  next
  case True
  with that have  $z: z \in A$   $f z = 0$ 
    by auto
  have zorder  $f z = \text{int}(\text{subdegree}(F z))$ 
    using F by (rule has_fps_expansion_zorder) (use  $z$  in auto)
  also have zorder  $f z = \text{zorder } g z$ 
    using  $z$  assms by auto
  also have zorder  $g z = \text{subdegree}(G z)$ 
    using G by (rule has_fps_expansion_zorder) (use  $z$  in auto)
  finally have subdegree_eq:  $\text{subdegree}(F z) = \text{subdegree}(G z)$ 
    by simp

  have  $(\lambda w. \text{if } w = 0 \text{ then } c z \text{ else } g(z + w) / f(z + w))$  has_fps_expansion
 $G z / F z$  (is ?P)
    using subdegree_eq  $z$  by (intro has_fps_expansion_divide F G) (auto simp:
c_def)
  also have ?this  $\longleftrightarrow$   $(\lambda w. h(z + w))$  has_fps_expansion  $G z / F z$ 
  proof (intro has_fps_expansion_cong)
    have eventually  $(\lambda w. w \neq z \longrightarrow f w \neq 0)$  (nhds  $z$ )
      using ev_nonzero[OF  $\langle z \in A \rangle$ ] unfolding eventually_at_filter
      by eventually_elim (use  $\langle z \in A \rangle$  assms in auto)
    hence eventually  $(\lambda w. w \neq 0 \longrightarrow f(z + w) \neq 0)$  (nhds  $0$ )
      by (simp add: nhds_to_0' eventually_filtermap)

```

```

      thus eventually ( $\lambda w. (if\ w = 0\ then\ c\ z\ else\ g\ (z + w) / f\ (z + w)) = h\ (z + w)$ ) (nhds 0)
    unfolding h_def by eventually_elim (use z in ‹auto simp: c_def h_def›)
  qed auto
  finally have h_analytic_on {z}
    using has_fps_expansion_imp_analytic by blast
  moreover have h_z ≠ 0
    using that z by (auto simp: h_def c_def)
  ultimately show ?thesis by blast
qed
qed

from h have h_ana: h_analytic_on A and h_nz:  $\forall z \in A. h\ z \neq 0$ 
  using analytic_on_analytic_at by blast+
moreover have g_z = h_z * f_z if z ∈ A for z
  using assms that by (auto simp: h_def)
ultimately show ?thesis
  using ‹open A› by (intro that[of h]) (auto simp: analytic_on_open)
qed

```

9.4 Constructing the sequence of zeros

The form of the Weierstraß Factorisation Theorem that we derived above requires an explicit sequence of the zeros that tends to infinity. We will now show that under mild conditions, such a sequence always exists.

More precisely: if A is an infinite closed set that is sparse in the sense that its intersection with any compact set is finite, then there exists an injective sequence f enumerating the values of A in ascending order by absolute value, and f tends to infinity for $n \rightarrow \infty$.

lemma *sequence_of_sparse_set_exists:*

fixes $A :: complex\ set$

assumes *infinite A closed A $\wedge r. r \geq 0 \implies finite (A \cap cball\ 0\ r)$*

obtains $f :: nat \Rightarrow complex$

where *mono (norm ∘ f) inj f range f = A filterlim f at_infinity at_top*

proof –

have $\exists f :: nat \Rightarrow complex. \forall n.$

$f\ n \in A \wedge$

$f\ n \notin f\ \{..<n\} \wedge$

$\{z \in A. norm\ z < norm\ (f\ n)\} \subseteq f\ \{..<n\} \wedge$

$(\forall k < n. norm\ (f\ k) \leq norm\ (f\ n))$

proof (*rule dependent_wf_choice[OF wf], goal_cases*)

case ($1\ f\ g\ n\ r$)

thus ?case by auto

next

case ($2\ n\ f$)

have $f: f\ k \in A \ \{z \in A. norm\ z < norm\ (f\ k)\} \subseteq f\ \{..<k\} \ \forall k' < k. cmod\ (f\ k') \leq cmod\ (f\ k)$

if $k < n$ for k

```

    using 2[of k] that by simp_all

  have infinite (A - f' {.. $n$ })
    using assms(1) by (intro Diff_infinite_finite) auto
  then obtain z0 where z0: z0 ∈ A - f' {.. $n$ }
    by (meson finite.intros(1) finite_subset subsetI)
  have finite (A ∩ cball 0 (norm z0))
    by (intro assms(3)) auto
  hence finite (A ∩ cball 0 (norm z0) - f' {.. $n$ })
    using finite_subset by blast
  moreover from z0 have A ∩ cball 0 (norm z0) - f' {.. $n$ } ≠ {}
    by auto
  ultimately obtain z where is_arg_min norm (λz. z ∈ A ∩ cball 0 (norm
z0) - f' {.. $n$ }) z
    using ex_is_arg_min_if_finite by blast
  hence z: z ∈ A norm z ≤ norm z0 z ∉ f' {.. $n$ }
    ∧ w. w ∈ A ⇒ norm w ≤ norm z0 ⇒ w ∉ f' {.. $n$ } ⇒ norm w ≥
norm z
    by (auto simp: is_arg_min_def)

  show ?case
  proof (rule exI[of _ z], safe)
    fix w assume w: w ∈ A norm w < norm z
    with z(4)[of w] z w show w ∈ f' {.. $n$ }
      by linarith
    next
    fix k assume k: k < n
    show norm (f k) ≤ norm z
      using f(2)[of k] z(1,3) k by auto
    qed (use z in auto)
  qed

  then obtain f :: nat ⇒ complex where f:
    ∧ n. f n ∈ A
    ∧ n. f n ∉ f' {.. $n$ }
    ∧ n. {z ∈ A. norm z < norm (f n)} ⊆ f' {.. $n$ }
    ∧ k n. k < n ⇒ norm (f k) ≤ norm (f n)
    by meson
  from f(2) have f_neq: f n ≠ f k if k < n for k n
    using that by blast

  have inj: inj f
  proof (rule injI)
    fix m n :: nat
    assume f m = f n
    thus m = n
      using f_neq[of m n] f_neq[of n m] by (cases m n rule: linorder_cases) auto
  qed

  have range: range f = A

```



```

proof safe
  fix  $z$  assume  $z: z \in A$ 
  show  $z \in \text{range } f$ 
  proof (rule ccontr)
    assume  $z \notin \text{range } f$ 
    hence  $\text{norm } (f\ n) \leq \text{norm } z$  for  $n$ 
      using  $f(3)[\text{of } n]$  by auto
    hence  $\text{range } f \subseteq A \cap \text{cball } 0 (\text{norm } z)$ 
      using  $f(1)$  by auto
    moreover have  $\text{finite } (A \cap \text{cball } 0 (\text{norm } z))$ 
      by (intro assms) auto
    ultimately have  $\text{finite } (\text{range } f)$ 
      using  $\text{finite\_subset}$  by blast
    moreover have  $\text{infinite } (\text{range } f)$ 
      using inj by (subst  $\text{finite\_image\_iff}$ ) auto
    ultimately show  $\text{False}$  by contradiction
  qed
qed (use  $f(1)$  in auto)

have  $\text{mono}: \text{mono } (\text{norm } \circ f)$ 
proof (rule  $\text{monoI}$ , unfold  $\text{o\_def}$ )
  fix  $m\ n :: \text{nat}$ 
  assume  $m \leq n$ 
  thus  $\text{norm } (f\ m) \leq \text{norm } (f\ n)$ 
    using  $f(4)[\text{of } m\ n]$  by (cases  $m < n$ ) auto
qed

have  $\neg \text{bounded } A$ 
proof
  assume  $\text{bounded } A$ 
  hence  $\text{bdd\_above } (\text{norm } \text{' } A)$ 
    by (meson  $\text{bdd\_above\_norm}$ )
  hence  $\text{norm } z \leq \text{Sup } (\text{norm } \text{' } A)$  if  $z \in A$  for  $z$ 
    using  $\text{that}$  by (meson  $\text{cSUP\_upper}$ )
  hence  $A \subseteq \text{cball } 0 (\text{Sup } (\text{norm } \text{' } A))$ 
    by auto
  also have  $\dots \subseteq \text{cball } 0 (\text{max } 1 (\text{Sup } (\text{norm } \text{' } A)))$ 
    by auto
  finally have  $A \subseteq A \cap \text{cball } 0 (\text{max } 1 (\text{Sup } (\text{norm } \text{' } A)))$ 
    by blast
  moreover have  $\text{finite } (A \cap \text{cball } 0 (\text{max } 1 (\text{Sup } (\text{norm } \text{' } A))))$ 
    by (intro assms) auto
  ultimately have  $\text{finite } A$ 
    using  $\text{finite\_subset}$  by blast
  hence  $\text{finite } (\text{range } f)$ 
    by (simp add:  $\text{range}$ )
  thus  $\text{False}$ 
    using inj by (subst (asm)  $\text{finite\_image\_iff}$ ) auto
qed

```

```

have lim: filterlim f at_infinity at_top
  unfolding filterlim_at_infinity_conv_norm_at_top filterlim_at_top
proof
  fix C :: real
  from ⟨¬bounded A⟩ obtain z where z ∈ A norm z > C
    unfolding bounded_iff by (auto simp: not_le)
  obtain n where [simp]: z = f n
    using range ⟨z ∈ A⟩ by auto
  show eventually (λn. norm (f n) ≥ C) at_top
    using eventually_ge_at_top[of n]
  proof eventually_elim
    case (elim k)
    have C ≤ norm (f n)
      using ⟨norm z > C⟩ by simp
    also have ... ≤ norm (f k)
      using monoD[OF ⟨mono (norm ∘ f)⟩, of n k] elim by auto
    finally show ?case .
  qed
qed

show ?thesis
  by (intro that[of f] inj range mono lim)
qed

lemma strict_mono_sequence_partition:
  assumes strict_mono (f :: nat ⇒ 'a :: {linorder, no_top})
  assumes x ≥ f 0
  assumes filterlim f at_top at_top
  shows ∃k. x ∈ {f k..<f (Suc k)}
proof -
  define k where k = (LEAST k. f (Suc k) > x)
  {
    obtain n where x ≤ f n
      using assms by (auto simp: filterlim_at_top eventually_at_top_linorder)
    also have f n < f (Suc n)
      using assms by (auto simp: strict_mono_Suc_iff)
    finally have ∃n. f (Suc n) > x by auto
  }
  from LeastI_ex[OF this] have x < f (Suc k)
    by (simp add: k_def)
  moreover have f k ≤ x
  proof (cases k)
    case (Suc k')
    have k ≤ k' if f (Suc k') > x
      using that unfolding k_def by (rule Least_le)
    with Suc show f k ≤ x by (cases f k ≤ x) (auto simp: not_le)
  qed (use assms in auto)

```

ultimately show *?thesis* by auto
qed

lemma *strict_mono_sequence_partition'*:
assumes *strict_mono* ($f :: \text{nat} \Rightarrow 'a :: \{\text{linorder}, \text{no_top}\}$)
assumes $x \geq f\ 0$
assumes *filterlim* f *at_top* *at_top*
shows $\exists!k. x \in \{f\ k..<f\ (Suc\ k)\}$
proof (*rule ex_ex1I*)
show $\exists k. x \in \{f\ k..<f\ (Suc\ k)\}$
using *strict_mono_sequence_partition*[*OF* *assms*].
fix $k1\ k2$ **assume** $x \in \{f\ k1..<f\ (Suc\ k1)\} \ x \in \{f\ k2..<f\ (Suc\ k2)\}$
thus $k1 = k2$
proof (*induction* $k1\ k2$ *rule: linorder_wlog*)
case (*le* $k1\ k2$)
hence $f\ k2 < f\ (Suc\ k1)$
by *auto*
hence $k2 < Suc\ k1$
using *assms*(1) *strict_mono_less* **by** *blast*
with *le* **show** $k1 = k2$
by *linarith*
qed *auto*
qed

lemma *sequence_of_sparse_set_exists'*:
fixes $A :: \text{complex set}$ **and** $c :: \text{complex} \Rightarrow \text{nat}$
assumes *infinite* A *closed* $A \ \bigwedge r. r \geq 0 \implies \text{finite}\ (A \cap \text{cball}\ 0\ r)$
assumes $c_pos: \bigwedge x. x \in A \implies c\ x > 0$
obtains $f :: \text{nat} \Rightarrow \text{complex}$ **where**
 $\text{mono}\ (\text{norm} \circ f)$ $\text{range}\ f = A$ *filterlim* f *at_infinity* *at_top*
 $\bigwedge z. z \in A \implies \text{finite}\ (f^{-1}\ \{z\}) \wedge \text{card}\ (f^{-1}\ \{z\}) = c\ z$
proof –
obtain $f :: \text{nat} \Rightarrow \text{complex}$ **where** f :
 $\text{mono}\ (\text{norm} \circ f)$ $\text{inj}\ f$ $\text{range}\ f = A$ *filterlim* f *at_infinity* *at_top*
using *assms* *sequence_of_sparse_set_exists* **by** *blast*
have f_eq_iff [*simp*]: $f\ m = f\ n \iff m = n$ **for** $m\ n$
using $\langle \text{inj}\ f \rangle$ **by** (*auto* *simp: inj_def*)

define $h :: \text{nat} \Rightarrow \text{nat}$ **where** $h = (\lambda n. \sum_{k < n}. c\ (f\ k))$

have [*simp*]: $h\ 0 = 0$
by (*simp* *add: h_def*)
have $h_ge: h\ n \geq n$ **for** n
proof –
have $h\ n \geq (\sum_{k < n}. 1)$
unfolding h_def **by** (*intro* *sum_mono*) (*use* $c_pos\ f$ **in** $\langle \text{auto}\ \text{simp:}\ Suc_le_eq \rangle$)
thus *?thesis* **by** *simp*
qed

```

have strict_mono h
  unfolding strict_mono_Suc_iff using f by (auto simp: h_def c_pos)
moreover from this have filterlim h at_top at_top
  using filterlim_subseq by blast
ultimately have Ex1:  $\exists!k. n \in \{h k..<h (Suc k)\}$  for n
  by (intro strict_mono_sequence_partition) auto

define g :: nat  $\Rightarrow$  nat where g = ( $\lambda n. THE k. n \in \{h k..<h (Suc k)\}$ )
have g:  $n \in \{h (g n)..<h (Suc (g n))\}$  for n
  using theI'[OF Ex1[of n]] by (simp add: g_def)
have g_eqI:  $g n = k$  if  $n \in \{h k..<h (Suc k)\}$  for n k
  using the1_equality[OF Ex1 that] by (simp add: g_def)
have g_h:  $g (h n) = n$  for n
  by (rule g_eqI) (auto intro: strict_monoD[OF  $\langle$ strict_mono h $\rangle$ ])

have mono g
  unfolding incseq_Suc_iff
proof safe
  fix n :: nat
  have h (g n) + 1  $\leq$  Suc n
    using g[of n] by auto
  also have Suc n < h (Suc (g (Suc n)))
    using g[of Suc n] by auto
  finally show g n  $\leq$  g (Suc n)
    by (metis  $\langle$ strict_mono h $\rangle$  add_lessD1 less_Suc_eq_le strict_mono_less)
qed

have filterlim g at_top at_top
  unfolding filterlim_at_top
proof
  fix n :: nat
  show eventually ( $\lambda m. g m \geq n$ ) at_top
    using eventually_ge_at_top[of h n]
  proof eventually_elim
    case (elim m)
    have n  $\leq$  g (h n)
      by (simp add: g_h)
    also have g (h n)  $\leq$  g m
      by (intro monoD[OF  $\langle$ mono g $\rangle$ ] elim)
    finally show ?case .
  qed
qed

have vimage:  $(f \circ g) - \{f n\} = \{h n..<h (Suc n)\}$  for n
  using g by (auto intro!: g_eqI)

show ?thesis
proof (rule that[of f  $\circ$  g])

```

```

have incseq ( $\lambda n. (norm \circ f) (g n)$ )
  by (intro monoI monoD[OF f(1)] monoD[OF <incseq g>])
thus incseq ( $norm \circ (f \circ g)$ )
  by (simp add: o_def)
next
have range ( $f \circ g$ )  $\subseteq A$ 
  using f(3) by auto
moreover have  $A \subseteq \text{range } (f \circ g)$ 
proof
  fix  $z$  assume  $z \in A$ 
  then obtain  $n$  where [simp]:  $z = f n$ 
    using f(3) by auto
  have  $f (g (h n)) \in \text{range } (f \circ g)$ 
    unfolding o_def by blast
  thus  $z \in \text{range } (f \circ g)$ 
    by (simp add: g_h)
qed
ultimately show  $\text{range } (f \circ g) = A$  by blast
next
fix  $z$  assume  $z \in A$ 
then obtain  $n$  where [simp]:  $z = f n$ 
  using f(3) by auto
have finite { $h n..<h (Suc n)$ }
  by auto
moreover have  $\text{card } \{h n..<h (Suc n)\} = c z$ 
  by (simp add: h_def)
ultimately show finite (( $f \circ g$ ) - '{ $z$ })  $\wedge$   $\text{card } ((f \circ g) - '{z}) = c z$ 
  using vimage[of n] by simp
next
show filterlim ( $f \circ g$ ) at_infinity at_top
  unfolding o_def by (rule filterlim_compose[OF f(4)] <filterlim g at_top
at_top>)]
qed
qed

```

9.5 The factorisation theorem for holomorphic functions

If g is a holomorphic function on an open connected domain whose zeros do not have an accumulation point on the frontier of A , then we can write g as a product of a function h holomorphic on A and an entire function f such that h is non-zero everywhere in A and the zeros of f are precisely the zeros of A with the same multiplicity.

In other words, we can get rid of all the zeros of g by dividing it with a suitable entire function f .

theorem weierstrass_factorization:

assumes g holomorphic_on A open A connected A

assumes $\bigwedge z. z \in \text{frontier } A \implies \neg z \text{ islimpt } \{w \in A. g w = 0\}$

obtains $h f$ **where**

```

  h holomorphic_on A f holomorphic_on UNIV
   $\forall z. f z = 0 \iff (\forall z \in A. g z = 0) \vee (z \in A \wedge g z = 0)$ 
   $\forall z \in A. \text{zorder } f z = \text{zorder } g z$ 
   $\forall z \in A. h z \neq 0$ 
   $\forall z \in A. g z = h z * f z$ 
proof (cases  $\forall z \in A. g z = 0$ )
case True
show ?thesis
proof (rule that[of  $\lambda_. 1 \lambda_. 0$ ]; (intro ballI allI impI)?)
  fix z assume z: z  $\in$  A
  have ev: eventually ( $\lambda w. w \in A$ ) (at z)
    using z assms by (intro eventually_at_in_open') auto
  show zorder ( $\lambda :: \text{complex}. 0 :: \text{complex}$ ) z = zorder g z
    by (intro zorder_cong eventually_mono[OF ev] refl) (use True in auto)
qed (use assms True in auto)
next
case not_identically_zero: False
define Z where Z = {z  $\in$  A. g z = 0}
have freq_nz: frequently ( $\lambda z. g z \neq 0$ ) (at z) if z  $\in$  A for z
proof -
  have  $\forall_F w$  in at z. g w  $\neq$  0  $\wedge$  w  $\in$  A
    using non_zero_neighbour_alt[OF assms(1,2,3) that(1)] not_identically_zero
by auto
  hence  $\forall_F w$  in at z. g w  $\neq$  0
    by eventually_elim auto
  thus ?thesis
    using eventually_frequently by force
qed

have zorder_pos_iff: zorder g z > 0  $\iff$  g z = 0 if z  $\in$  A for z
  by (subst zorder_pos_iff[OF assms(1,2) that]) (use freq_nz[of z] that in auto)

show ?thesis
proof (cases finite Z)
case True
  define f where f = ( $\lambda z. \prod_{w \in Z}. (z - w)^{\text{powi } (\text{zorder } g w)}$ )

  have eq_zero_iff: f z = 0  $\iff$  z  $\in$  A  $\wedge$  g z = 0 for z
    using True local.zorder_pos_iff
    unfolding f_def Z_def by fastforce
  have zorder_eq: zorder f z = zorder g z if z  $\in$  A for z
proof (cases g z = 0)
case False
  have g analytic_on {z}
    using that assms analytic_at by blast
  hence [simp]: zorder g z = 0
    using False by (intro zorder_eq_0I) auto
  moreover have f analytic_on {z}
    unfolding f_def by (auto intro!: analytic_intros)

```

```

hence zorder f z = 0
  using False by (intro zorder_eq_0I) (auto simp: eq_zero_iff)
ultimately show ?thesis
  by simp
next
case zero: True
have g analytic_on {z}
  using that assms(1,2) analytic_at by blast
hence zorder g z ≥ 0
  using that by (intro zorder_ge_0 freq_nz) auto
define f' where f' = (λz'. (∏ w∈Z- $\{z\}$ . (z' - w) powi (zorder g w)))
have zorder (λz'. (z' - z) powi (zorder g z) * f' z') z = zorder g z
proof (rule zorder_eqI)
  show open (UNIV :: complex set) f' holomorphic_on UNIV z ∈ UNIV
    using local.zorder_pos_iff
    by (fastforce intro!: holomorphic_intros simp: f'_def Z_def)+
  show f' z ≠ 0
    using True unfolding f'_def by (subst prod_zero_iff) auto
qed (use ⟨zorder g z ≥ 0⟩ in ⟨auto simp: powr_of_int⟩)
also have (λz'. (z' - z) powi (zorder g z) * f' z') = f
proof
  fix z' :: complex
  have Z = insert z (Z - {z})
    using that zero by (auto simp: Z_def)
  also have (∏ w∈... (z' - w) powi (zorder g w)) = (z' - z) powi (zorder
g z) * f' z'
    using True by (subst prod.insert) (auto simp: f'_def)
  finally show (z' - z) powi (zorder g z) * f' z' = f z'
    by (simp add: f'_def)
qed
finally show ?thesis .
qed

obtain h :: complex ⇒ complex where h:
  h holomorphic_on A ∧ z. z ∈ A ⇒ h z ≠ 0 ∧ z. z ∈ A ⇒ g z = h z * f z
proof (rule holomorphic_zorder_factorization[OF assms(1-3)])
  show f holomorphic_on A
    using local.zorder_pos_iff
    unfolding f_def Z_def by (fastforce intro: holomorphic_intros)
  show f z = 0 ⇔ g z = 0 if z ∈ A for z
    using that by (subst eq_zero_iff) auto
  show zorder f z = zorder g z if z ∈ A for z
    by (rule zorder_eq) fact
qedmetis

show ?thesis
proof (rule that[of h f]; (intro ballI)?)
  show h holomorphic_on A
    by fact

```

```

    show f holomorphic_on UNIV
      using local.zorder_pos_iff
      unfolding f_def Z_def by (fastforce intro: holomorphic_intros)
    qed (use True not_identically_zero in ⟨auto simp: eq_zero_iff zorder_eq h(2)
h(3)[symmetric]⟩)

next
case False
note infinite_zeroes = not_identically_zero False
define c where c = (λz. nat (zorder g z))

have ev_nz: eventually (λw. g w ≠ 0) (at z) if z ∈ A for z
proof -
  from infinite_zeroes(1) obtain z0 where z0: z0 ∈ A g z0 ≠ 0
  by auto
  have eventually (λw. g w ≠ 0 ∧ w ∈ A) (at z)
  by (rule non_zero_neighbour_alt[where ?β = z0]) (use assms z0 that in
auto)
  thus ?thesis
  by eventually_elim auto
qed

have no_limpt_Z: ¬z islimpt Z for z
proof
  assume z islimpt Z
  show False
  proof (cases z ∈ A)
  case False
  have z islimpt A
  by (rule islimpt_subset[OF ⟨z islimpt Z⟩]) (auto simp: Z_def)
  hence z ∈ closure A
  by (simp add: closure_def)
  with ⟨z ∉ A⟩ have z ∈ frontier A
  by (simp add: closure_Un_frontier)
  with assms and ⟨z islimpt Z⟩ show False
  by (auto simp: Z_def)
  next
  case True
  from True have eventually (λw. g w ≠ 0) (at z)
  using ev_nz by blast
  hence ¬z islimpt Z
  by (auto simp: islimpt_iff_eventually Z_def elim!: eventually_mono)
  with ⟨z islimpt Z⟩ show False by blast
  qed
qed
have closed Z
  using no_limpt_Z unfolding closed_limpt by blast

obtain a where a:

```



```

    incseq (norm ∘ a) range a = Z - {0}
    ∧z. z ∈ Z - {0} ⇒ finite (a - ' {z}) ∧ card (a - ' {z}) = c z
    filterlim a at_infinity at_top
  proof (rule sequence_of_sparse_set_exists')
    show infinite (Z - {0})
      using infinite_zeroes(2) by auto
  next
    show closed (Z - {0})
      unfolding closed_limpt using no_limpt_Z islimpt_subset by blast
  next
    show finite ((Z - {0}) ∩ cball 0 R) if R ≥ 0 for R
    proof (rule ccontr)
      assume *: infinite ((Z - {0}) ∩ cball 0 R)
      have ∃ z ∈ cball 0 R. z islimpt ((Z - {0}) ∩ cball 0 R)
        by (rule Heine_Borel_imp_Bolzano_Weierstrass) (use * in auto)
      then obtain z where z islimpt ((Z - {0}) ∩ cball 0 R)
        by blast
      hence z islimpt Z
        using islimpt_subset by blast
      thus False
        using no_limpt_Z by blast
    qed
  next
    show c z > 0 if z ∈ Z - {0} for z
      using zorder_pos_iff[of z] that by (auto simp: c_def Z_def)
    qed metis

  interpret f: weierstrass_product' a
  proof
    show a n ≠ 0 for n
      using a(2) by auto
    show finite (a - ' {z}) if z ∈ range a for z
      using a(3)[of z] a(2) that by simp
    qed fact+

  define m where m = (if 0 ∈ A then nat (zorder g 0) else 0)

  have zorder_eq: zorder (λz. z ^ m * f.f z) z = zorder g z if z ∈ A for z
  proof (cases g z = 0)
    case False
      have g analytic_on {z}
        using ⟨z ∈ A⟩ analytic_at assms by blast
      hence zorder g z = 0
        by (intro zorder_eq_0I False)
      have z ∉ range a
        using False Z_def a(2) by blast
      hence zorder (λz. z ^ m * f.f z) z = 0
        using False ⟨zorder g z = 0⟩
        by (intro zorder_eq_0I analytic_intros) (auto simp: f.zero m_def)
  
```

```

with ‹zorder g z = 0› show ?thesis
  by simp
next
case True
define F where F = fps_expansion f.f z
have frequently (λw. g w ≠ 0) (at z)
  using ev_nz[OF that] eventually_frequently by force
hence zorder g z ≥ 0
  by (intro zorder_ge_0) (use assms that in ‹auto simp: analytic_at›)

have ev: eventually (λz. z ∈ A) (nhds z)
  using that assms by (intro eventually_nhds_in_open) auto
have exp1: (λw. f.f (z + w)) has_fps_expansion F
  unfolding F_def by (intro analytic_at_imp_has_fps_expansion[OF
f.analytic])
have exp2: (λw. (z + w) ^ m * f.f (z + w)) has_fps_expansion (fps_const
z + fps_X) ^ m * F
  by (intro fps_expansion_intros exp1)
have [simp]: F ≠ 0
proof
  assume F = 0
  hence eventually (λz. f.f z = 0) (nhds z)
    using exp1 by (auto simp: has_fps_expansion_def nhds_to_0' eventu-
ally_filtermap)
  hence eventually (λz. g z = 0) (at z)
    by (auto simp: f.zero a_Z_def eventually_at_filter elim!: eventually_mono)
  hence eventually (λz::complex. False) (at z)
    using ev_nz[OF ‹z ∈ A›] by eventually_elim auto
  thus False by simp
qed

have zorder (λw. w ^ m * f.f w) z = int (subdegree ((fps_const z + fps_X)
^ m * F))
  using has_fps_expansion_zorder[OF exp2] by simp
also have ... = int (subdegree F) + (if z = 0 then m else 0)
  by auto
also have int (subdegree F) = zorder f.f z
  using has_fps_expansion_zorder[OF exp1] by simp
also have ... = int (card (a - ‹{z}›))
  by (rule f.zorder)
also have card (a - ‹{z}›) = (if z = 0 then 0 else c z)
proof (cases z = 0)
case True
  hence a - ‹{z}› = {}
    using a(2) by auto
  thus ?thesis using True by simp
next
case False
  thus ?thesis

```

```

    by (subst a(3)) (use True that in ⟨auto simp: Z_def⟩)
  qed
  also have int ... + (if z = 0 then m else 0) = zorder g z
    using ⟨zorder g z ≥ 0⟩ that by (auto simp: c_def m_def)
  finally show ?thesis .
qed

have eq_zero_iff: z ^ m * f.f z = 0 ⟷ g z = 0 if z ∈ A for z
  using that by (auto simp add: f.zero a m_def zorder_pos_iff Z_def)

obtain h :: complex ⇒ complex where h:
  h holomorphic_on A ∧ z. z ∈ A ⟹ h z ≠ 0 ∧ z. z ∈ A ⟹ g z = h z * (z
  ^ m * f.f z)
proof (rule holomorphic_zorder_factorization[OF assms(1-3)])
  show (λz. z ^ m * f.f z) holomorphic_on A
    by (intro holomorphic_intros)
  show z ^ m * f.f z = 0 ⟷ g z = 0 if z ∈ A for z
    by (rule eq_zero_iff) fact+
  show zorder (λz. z ^ m * f.f z) z = zorder g z if z ∈ A for z
    by (rule zorder_eq) fact+
qed metis

show ?thesis
proof (rule that[of h λz. z ^ m * f.f z]; (intro ballI allI impI)?)
  show h holomorphic_on A
    by fact
  show (λz. z ^ m * f.f z) holomorphic_on UNIV
    by (intro holomorphic_intros)
next
fix z :: complex
show (z ^ m * f.f z = 0) = ((∀ z ∈ A. g z = 0) ∨ z ∈ A ∧ g z = 0)
  using infinite_zeroes(1) a(2)
  by (auto simp: m_def zorder_eq eq_zero_iff zorder_pos_iff Z_def f.zero)
qed (use zorder_eq eq_zero_iff h in auto)
qed
qed

```

The following is a simpler version for entire functions.

```

theorem weierstrass_factorization_UNIV:
  assumes g holomorphic_on UNIV
  obtains h f where
    h holomorphic_on UNIV f holomorphic_on UNIV
    ∀ z. f z = 0 ⟷ g z = 0
    ∀ z. zorder f z = zorder g z
    ∀ z. h z ≠ 0
    ∀ z. g z = h z * f z
  using assms by (rule weierstrass_factorization, goal_cases) auto

```

9.6 The factorisation theorem for meromorphic functions

Let g be a meromorphic function on a connected open domain A . Assume that the poles and zeros of g have no accumulation point on the border of A . Then g can be written in the form $g(z) = h(z)f_1(z)/f_2(z)$ where h is holomorphic on A with no zeroes and f_1 and f_2 are entire.

Moreover, as direct consequences of that, the zeroes of f_1 are precisely the zeroes of g and the zeros of f_2 are precisely the poles of g (with the corresponding multiplicity).

theorem *weierstrass_factorization_meromorphic*:

assumes *mero*: g *nicely_meromorphic_on* A **and** A : *open* A *connected* A

assumes *no_limpt*: $\bigwedge z. z \in \text{frontier } A \implies \neg z \text{ islimpt } \{w \in A. g w = 0 \vee \text{is_pole } g w\}$

obtains $h f_1 f_2$ **where**

h *holomorphic_on* A f_1 *holomorphic_on* UNIV f_2 *holomorphic_on* UNIV

$\forall z \in A. f_1 z = 0 \iff \neg \text{is_pole } g z \wedge g z = 0$

$\forall z \in A. f_2 z = 0 \iff \text{is_pole } g z$

$\forall z \in A. \neg \text{is_pole } g z \implies \text{zorder } f_1 z = \text{zorder } g z$

$\forall z \in A. \text{is_pole } g z \implies \text{zorder } f_2 z = -\text{zorder } g z$

$\forall z \in A. h z \neq 0$

$\forall z \in A. g z = h z * f_1 z / f_2 z$

proof –

have *mero'*: g *meromorphic_on* A

using *mero* **unfolding** *nicely_meromorphic_on_def* **by** *auto*

define *pts* **where** $pts = \{z \in A. \text{is_pole } g z\}$

have $\{z. \text{is_pole } g z\}$ *sparse_in* A

using *meromorphic_on_imp_not_pole_cosparse*[*OF mero'*]

by (*auto simp: eventually_cosparse*)

hence *pts* *sparse_in* A

unfolding *pts_def* **by** (*rule sparse_in_subset2*) *auto*

have *open_diff_pts*: *open* ($A - pts'$) **if** $pts' \subseteq pts$ **for** pts'

proof (*rule open_diff_sparse_pts*)

show pts' *sparse_in* A

using $\langle pts \text{ sparse_in } A \rangle$ **by** (*rule sparse_in_subset2*) *fact*

qed (*use* $\langle \text{open } A \rangle$ **in** *auto*)

have *ev*: *eventually* ($\lambda w. w \in A - pts$) (*at* z) **if** $z \in A$ **for** z

proof (*cases* $z \in pts$)

case *False*

thus *?thesis*

using *that open_diff_pts*[*of pts*] **by** (*intro eventually_at_in_open'*) *auto*

next

case *True*

have *eventually* ($\lambda w. w \in (A - (pts - \{z\})) - \{z\}$) (*at* z)

using *that* **by** (*intro eventually_at_in_open open_diff_pts*) *auto*

also **have** $A - (pts - \{z\}) - \{z\} = A - pts$

using *True* **by** *auto*

finally **show** *?thesis* .

qed

show ?thesis

proof (cases $\forall z \in A - \text{pts}. g z = 0$)

case True

have no_poles: $\neg \text{is_pole } g z$ if $z \in A$ for z

proof -

have is_pole $g z \iff \text{is_pole } (\lambda _ :: \text{complex}. 0 :: \text{complex}) z$

by (intro is_pole_cong[OF eventually_mono[OF ev]]) (use that True in auto)

thus ?thesis

by simp

qed

hence [simp]: $\text{pts} = \{\}$

by (auto simp: pts_def)

have [simp]: $\text{zorder } g z = \text{zorder } (\lambda _ :: \text{complex}. 0 :: \text{complex}) z$ if $z \in A$ for z

by (intro zorder_cong[OF eventually_mono[OF ev]]) (use True that in auto)

show ?thesis

by (rule that[of $\lambda _. 1 \lambda _. 0 \lambda _. 1$]) (use True in $\langle \text{auto simp: no_poles} \rangle$)

next

case False

have is_pole_iff: $\text{is_pole } g z \iff z \in \text{pts}$ if $z \in A$ for z

using that by (auto simp: pts_def)

obtain h f1 where h_f1:

h holomorphic_on $A - \text{pts}$ f1 holomorphic_on UNIV

$\forall z. f1 z = 0 \iff (\forall z \in A - \text{pts}. g z = 0) \vee (z \in A - \text{pts} \wedge g z = 0)$

$\forall z \in A - \text{pts}. \text{zorder } f1 z = \text{zorder } g z$

$\forall z \in A - \text{pts}. h z \neq 0$

$\forall z \in A - \text{pts}. g z = h z * f1 z$

proof (rule weierstrass_factorization)

have g analytic_on $A - \text{pts}$

by (rule nicely_meromorphic_without_singularities)

(use mero in $\langle \text{auto simp: pts_def dest: nicely_meromorphic_on_subset} \rangle$)

thus holo: g holomorphic_on $A - \text{pts}$

by (rule analytic_imp_holomorphic)

show open $(A - \text{pts})$

by (rule open_diff_pts) auto

show connected $(A - \text{pts})$

by (rule sparse_imp_connected) (use A $\langle \text{pts sparse_in } A \rangle$ in auto)

show $\neg z \text{ islimpt } \{w \in A - \text{pts}. g w = 0\}$ if $z \in \text{frontier } (A - \text{pts})$ for z

proof -

from that have $z \in \text{frontier } A - \text{pts} \cup \text{pts}$

using $\langle \text{open } (A - \text{pts}) \rangle \langle \text{open } A \rangle \text{ closure_mono[of } A - \text{pts } A]$

by (auto simp: frontier_def interior_open)

thus ?thesis

proof

```

assume  $z \in pts$ 
hence  $is\_pole\ g\ z$ 
  by (auto simp: pts_def)
hence eventually ( $\lambda z. g\ z \neq 0$ ) (at  $z$ )
  using non_zero_neighbour_pole by blast
hence  $\neg z\ islimpt\ \{w. g\ w = 0\}$ 
  by (auto simp: islimpt_iff_eventually)
thus ?thesis
  using islimpt_subset[of  $z\ \{w \in A - pts. g\ w = 0\}\ \{w. g\ w = 0\}$ ] by blast
next
assume  $z: z \in frontier\ A - pts$ 
show  $\neg z\ islimpt\ \{w \in A - pts. g\ w = 0\}$ 
proof
  assume  $z\ islimpt\ \{w \in A - pts. g\ w = 0\}$ 
  hence  $z\ islimpt\ \{w \in A. g\ w = 0 \vee is\_pole\ g\ w\}$ 
    by (rule islimpt_subset) (auto simp: pts_def)
  thus False using no_limpt  $z$  by blast
qed
qed
qed
qed

have  $f1\_eq\_0\_iff: f1\ z = 0 \iff (z \in A - pts \wedge g\ z = 0)$  for  $z$ 
  using  $h\_f1(3)$  False by auto

define  $h'$  where  $h' = (\lambda z. if\ z \in pts\ then\ 0\ else\ inverse\ (h\ z))$ 

have  $isolated\_h: isolated\_singularity\_at\ h\ z$  if  $z \in pts$  for  $z$ 
proof -
  have open ( $A - (pts - \{z\})$ )
    by (rule open_diff_pts) auto
  moreover have  $z \in (A - (pts - \{z\}))$ 
    using that by (auto simp: pts_def)
  moreover have  $h$  holomorphic_on ( $A - (pts - \{z\}) - \{z\}$ )
    by (rule holomorphic_on_subset[OF  $h\_f1(1)$ ]) (use that in auto)
  ultimately show  $isolated\_singularity\_at\ h\ z$ 
    using isolated_singularity_at_holomorphic by blast
qed

have  $is\_pole\_h: is\_pole\ h\ z$  if  $z \in A$   $is\_pole\ g\ z$  for  $z$ 
proof -
  have  $f1: f1$  analytic_on  $\{z\}$ 
    by (meson analytic_on_holomorphic  $h\_f1(2)$  open_UNIV top_greatest)
  have eventually ( $\lambda w. g\ w \neq 0$ ) (at  $z$ )
    using  $\langle is\_pole\ g\ z \rangle$  non_zero_neighbour_pole by blast
  with  $ev[OF\ that(1)]$  have  $ev'$ : eventually ( $\lambda w. g\ w * f1\ w \neq 0$ ) (at  $z$ )
    by eventually_elim (use  $h\_f1(3)$  in auto)

  have  $is\_pole\ (\lambda w. g\ w / f1\ w)\ z$ 

```

```

proof (rule is_pole_divide_zorder)
  show isolated_singularity_at f1 z not_essential f1 z
using f1 by (simp_all add: isolated_singularity_at_analytic not_essential_analytic)
show isolated_singularity_at g z not_essential g z
  using mero' that unfolding meromorphic_on_altdef by blast+
show freq: frequently ( $\lambda w. g w * f1 w \neq 0$ ) (at z)
  using ev' by (rule eventually_frequently[rotated]) auto
from freq have freq': frequently ( $\lambda w. f1 w \neq 0$ ) (at z)
  using frequently_elim1 by fastforce
have zorder g z < 0
using ⟨is_pole g z⟩ ⟨isolated_singularity_at g z⟩ isolated_pole_imp_neg_zorder
by auto
  also have 0 ≤ zorder f1 z
    by (rule zorder_ge_0[OF f1 freq'])
  finally show zorder g z < zorder f1 z .
qed
also have ?this ↔ is_pole h z
  proof (intro is_pole_cong refl eventually_mono[OF eventually_conj[OF
ev[OF that(1)] ev']])
    fix w assume w ∈ A - pts ∧ g w * f1 w ≠ 0
    thus g w / f1 w = h w using h_f1(6)
      by (auto simp: divide_simps)
  qed
finally show is_pole h z .
qed

have h' analytic_on {z} if z ∈ A for z
proof (cases z ∈ pts)
  case False
    moreover have open (A - pts)
      by (rule open_diff_pts) auto
    ultimately have ( $\lambda z. \text{inverse } (h z)$ ) analytic_on {z}
      using that h_f1(1,2,5) ⟨open (A - pts)⟩ analytic_at False
      by (intro analytic_intros) (auto simp: f1_eq_0_iff)
    also have eventually ( $\lambda z. z \in A - \text{pts}$ ) (nhds z)
      using that False ⟨open (A - pts)⟩ by (intro eventually_nhds_in_open) auto
    hence ( $\lambda z. \text{inverse } (h z)$ ) analytic_on {z} ↔ h' analytic_on {z}
      by (intro analytic_at_cong) (auto elim!: eventually_mono simp: h'_def)
    finally show ?thesis .
  next
    case True
      have ( $\lambda w. \text{if } w = z \text{ then } 0 \text{ else } \text{inverse } (h w)$ ) holomorphic_on (A - (pts - {z}))
        proof (rule is_pole_inverse_holomorphic)
          from True have A - (pts - {z}) - {z} = A - pts
            by auto
          thus h holomorphic_on A - (pts - {z}) - {z}
            using h_f1(1) by simp
        next

```

```

    show open (A - (pts - {z}))
      by (rule open_diff_pts) auto
  next
    have is_pole g z
      using True that by (simp add: is_pole_iff)
    thus is_pole h z
      using is_pole_h that by auto
  next
    show  $\forall w \in A - (pts - \{z\}) - \{z\}. h w \neq 0$ 
      using h_f1(5) by auto
  qed
  also have ?this  $\longleftrightarrow$  h' holomorphic_on A - (pts - {z})
  proof (intro holomorphic_cong refl)
    fix w assume w:  $w \in A - (pts - \{z\})$ 
    show (if  $w = z$  then 0 else inverse (h w)) = h' w
      using True w by (cases w = z) (auto simp: h'_def)
  qed
  finally have h' holomorphic_on A - (pts - {z}) .
  moreover have  $z \in A - (pts - \{z\})$  open (A - (pts - {z}))
    using True that by (auto intro!: open_diff_pts)
  ultimately show h' analytic_on {z}
    using analytic_at by blast
  qed
  hence h': h' analytic_on A
    using analytic_on_analytic_at by blast

  have h'_eq_0_iff:  $h' w = 0 \longleftrightarrow$  is_pole g w if  $w \in A$  for w
    using that h_f1(5) is_pole_iff[of w] by (auto simp: h'_def)

  obtain h'' f2 where h''_f2:
    h'' holomorphic_on A f2 holomorphic_on UNIV
     $\forall z. f2 z = 0 \longleftrightarrow (\forall z \in A. h' z = 0) \vee (z \in A \wedge h' z = 0)$ 
     $\forall z \in A. h' z = 0 \longrightarrow zorder f2 z = zorder h' z$ 
     $\forall z \in A. h'' z \neq 0 \forall z \in A. h' z = h'' z * f2 z$ 
  proof (rule weierstrass_factorization[of h' A])
    show open A connected A
      by fact+
    show h' holomorphic_on A
      using h' <open A> by (simp add: analytic_on_open)
    show  $\neg z$  islimpt { $w \in A. h' w = 0$ } if  $z \in$  frontier A for z
  proof
    assume z islimpt { $w \in A. h' w = 0$ }
    also have { $w \in A. h' w = 0$ } = pts
      by (auto simp: h'_eq_0_iff pts_def)
    finally have z islimpt { $w \in A. g w = 0 \vee$  is_pole g w}
      by (rule islimpt_subset) (auto simp: pts_def)
    thus False using no_limpt[of z] that
      by blast
  qed
  qed

```



```

qed blast

show ?thesis
proof (rule that[of  $\lambda w. \text{inverse } (h'' w) f1 f2$ ]; (intro ballI allI impI)?)
  show  $(\lambda w. \text{inverse } (h'' w)) \text{ holomorphic\_on } A$ 
    using  $h''\_f2(1,2,5)$  by (intro holomorphic_intros) auto
next
  show  $f1 \text{ holomorphic\_on } UNIV f2 \text{ holomorphic\_on } UNIV$ 
    by fact+
next
  show  $f1 z = 0 \longleftrightarrow \neg \text{is\_pole } g z \wedge g z = 0$  if  $z \in A$  for  $z$ 
    using that by (subst  $f1\_eq\_0\_iff$ ) (auto simp: pts_def)
next
  show  $f2 z = 0 \longleftrightarrow \text{is\_pole } g z$  if  $z \in A$  for  $z$ 
  proof -
    have  $\neg(\forall z \in A. h' z = 0)$ 
      using  $False h''\_f2(6) h\_f1(6) h'\_eq\_0\_iff \text{is\_pole\_iff}$  by auto
    hence  $f2 z = 0 \longleftrightarrow h' z = 0$ 
      using  $h''\_f2(3)$  that by auto
    also have  $\dots \longleftrightarrow \text{is\_pole } g z$ 
      using that by (simp add:  $\text{is\_pole\_iff } h'\_eq\_0\_iff$ )
    finally show ?thesis .
  qed
next
  show  $\text{zorder } f1 z = \text{zorder } g z$  if  $z \in A \neg \text{is\_pole } g z$  for  $z$ 
    using  $h\_f1(4)$  that by (auto simp: pts_def)
next
  show  $\text{zorder } f2 z = -\text{zorder } g z$  if  $z \in A \text{is\_pole } g z$  for  $z$ 
  proof -
    have  $\text{zorder } f2 z = \text{zorder } h' z$ 
      using  $h''\_f2(4)$  that  $h'\_eq\_0\_iff[of z] \text{is\_pole\_iff}[of z]$  by auto
    also have  $\dots = \text{zorder } (\lambda w. \text{inverse } (h w)) z$ 
      using that by (intro  $\text{zorder\_cong eventually\_mono}[OF ev]$ ) (auto simp:
 $h'\_def$ )
    also have  $\dots = -\text{zorder } h z$ 
  proof (intro  $\text{zorder\_inverse}$ )
    have  $\text{is\_pole } h z$ 
      using that  $\text{is\_pole\_iff}[of z] \text{is\_pole\_h}[of z]$  by auto
    thus  $\text{not\_essential } h z$ 
      by force
    show frequently  $(\lambda w. h w \neq 0)$  (at  $z$ )
      using  $\text{non\_zero\_neighbour\_pole}[OF \langle \text{is\_pole } h z \rangle]$  eventually_frequently
  by force
  qed (use that in  $\langle \text{auto intro!; isolated\_h simp: pts\_def} \rangle$ )
  also have  $\text{zorder } f1 z = 0$ 
  proof (rule  $\text{zorder\_eq\_0I}$ )
    show  $f1 \text{ analytic\_on } \{z\}$ 
      using that  $h\_f1(2) \text{ holomorphic\_on\_imp\_analytic\_at}$  by blast
    show  $f1 z \neq 0$ 

```

```

    using that h_f1(3) False by (auto simp: pts_def)
  qed
  hence zorder h z = zorder f1 z + zorder h z
    by simp
  also have ... = zorder ( $\lambda w. f1 w * h w$ ) z
  proof (rule zorder_times [symmetric])
    have f1 analytic_on {z}
      using that h_f1(2) holomorphic_on_imp_analytic_at by blast
    thus isolated_singularity_at f1 z not_essential f1 z
      using isolated_singularity_at_analytic not_essential_analytic by blast+
    show isolated_singularity_at h z
      using that by (intro isolated_h) (auto simp: pts_def)
    have is_pole h z
      using is_pole_iff[of z] that by (intro is_pole_h) auto
    thus not_essential h z
      by force
    have z  $\in A$ 
      using that by auto
    have eventually ( $\lambda w. g w \neq 0$ ) (at z)
      using non_zero_neighbour_pole[of g z] that by auto
    hence eventually ( $\lambda w. f1 w * h w \neq 0$ ) (at z)
      using ev[OF  $\langle z \in A \rangle$ ] by eventually_elim (use h_f1(6) in auto)
    thus frequently ( $\lambda w. f1 w * h w \neq 0$ ) (at z)
      using eventually_frequently by force
  qed
  also have ... = zorder g z
  proof (rule zorder_cong)
    have eventually ( $\lambda w. w \in A - pts$ ) (at z)
      using ev[of z] that by auto
    thus eventually ( $\lambda w. f1 w * h w = g w$ ) (at z)
      by eventually_elim (use h_f1(6) in auto)
  qed auto
  finally show ?thesis .
next
qed
show inverse (h'' z)  $\neq 0$  if z  $\in A$  for z
  using h''_f2(5) that by auto
next
show g z = inverse (h'' z) * f1 z / f2 z if z: z  $\in A$  for z
  proof (cases is_pole g z)
  case False
  have *: g z = h z * f1 z h' z = h'' z * f2 z h'' z  $\neq 0$  h z  $\neq 0$ 
    using that h''_f2(5,6) h_f1(5,6) False unfolding pts_def by blast+
  have g z = h z * f1 z
    by fact
  also have ... = f1 z / h' z
    using * that False by (simp add: h'_def field_simps pts_def)
  also have h' z = h'' z * f2 z
    by fact

```

```

    also have  $f1\ z / (h''\ z * f2\ z) = inverse\ (h''\ z) * f1\ z / f2\ z$ 
      by (simp add: divide_inverse_commute)
    finally show ?thesis .
  next
  case True
  have  $\neg g\ -z \rightarrow g\ z$ 
  using True at_neq_bot is_pole_def not_tendsto_and_filterlim_at_infinity
by blast
  with mero and z and True have  $g\ z = 0$ 
    by (auto simp: nicely_meromorphic_on_def)
  moreover have  $f2\ z = 0$ 
    using True z by (simp add: h''_f2(3) h'_eq_0_iff)
  ultimately show ?thesis by simp
qed
qed
qed
qed

```

Again, we derive an easier version for functions meromorphic on the entire complex plane.

theorem *weierstrass_factorization_meromorphic_UNIV*:

assumes g nicely_meromorphic_on UNIV

obtains $h\ f1\ f2$ where

h holomorphic_on UNIV $f1$ holomorphic_on UNIV $f2$ holomorphic_on UNIV

$\forall z. f1\ z = 0 \longleftrightarrow \neg is_pole\ g\ z \wedge g\ z = 0$

$\forall z. f2\ z = 0 \longleftrightarrow is_pole\ g\ z$

$\forall z. \neg is_pole\ g\ z \longrightarrow zorder\ f1\ z = zorder\ g\ z$

$\forall z. is_pole\ g\ z \longrightarrow zorder\ f2\ z = -zorder\ g\ z$

$\forall z. h\ z \neq 0$

$\forall z. g\ z = h\ z * f1\ z / f2\ z$

proof (rule *weierstrass_factorization_meromorphic*)

show g nicely_meromorphic_on UNIV

by fact

show connected (UNIV :: complex set)

by (simp add: Convex.connected_UNIV)

show $\neg z$ islimpt $\{w \in UNIV. g\ w = 0 \vee is_pole\ g\ w\}$ if $z \in frontier\ UNIV$ for

z

using that by simp

show open (UNIV :: complex set)

by simp

qed auto

end

theory *Complex_Analysis*

imports

Riemann_Mapping

Residue_Theorem

Weierstrass_Factorization

begin

end

References

[1]