

δ -Complete Decision Procedure and dReal

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Based on the work of Sicun Gao and Soonho Kong

Outline

- Interval constraints propagation (ICP)
 - Branch and Prune Algorithm
 - Completeness
 - dReal Example
- Adding ODEs
 - dReach Example
 - SMT encoding
- dReal Tricks

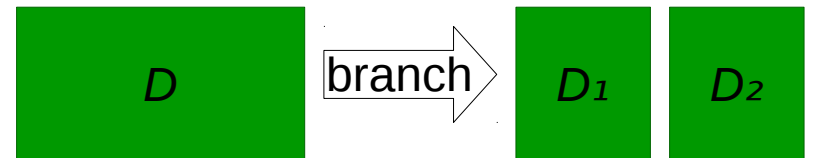
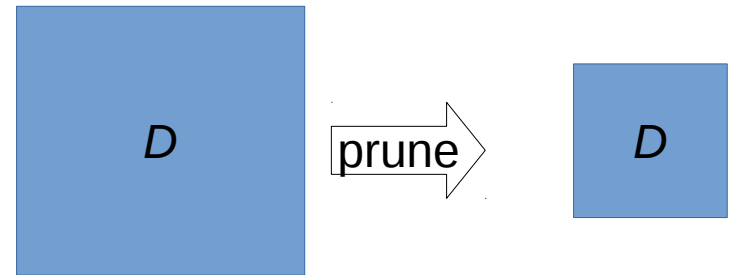
Interval Constraints Propagation

- Search for a solution using
 - Pruning: interval arithmetic to prune the search space.
 - Branching: when pruning is stuck, split the domain of a variable and continue recursively.
- Interval arithmetic on double precision numbers
 - Rounding errors taken into account
 - dReal uses IBEX and CAPD libraries
- Use $\delta > 0$ to guarantee the termination

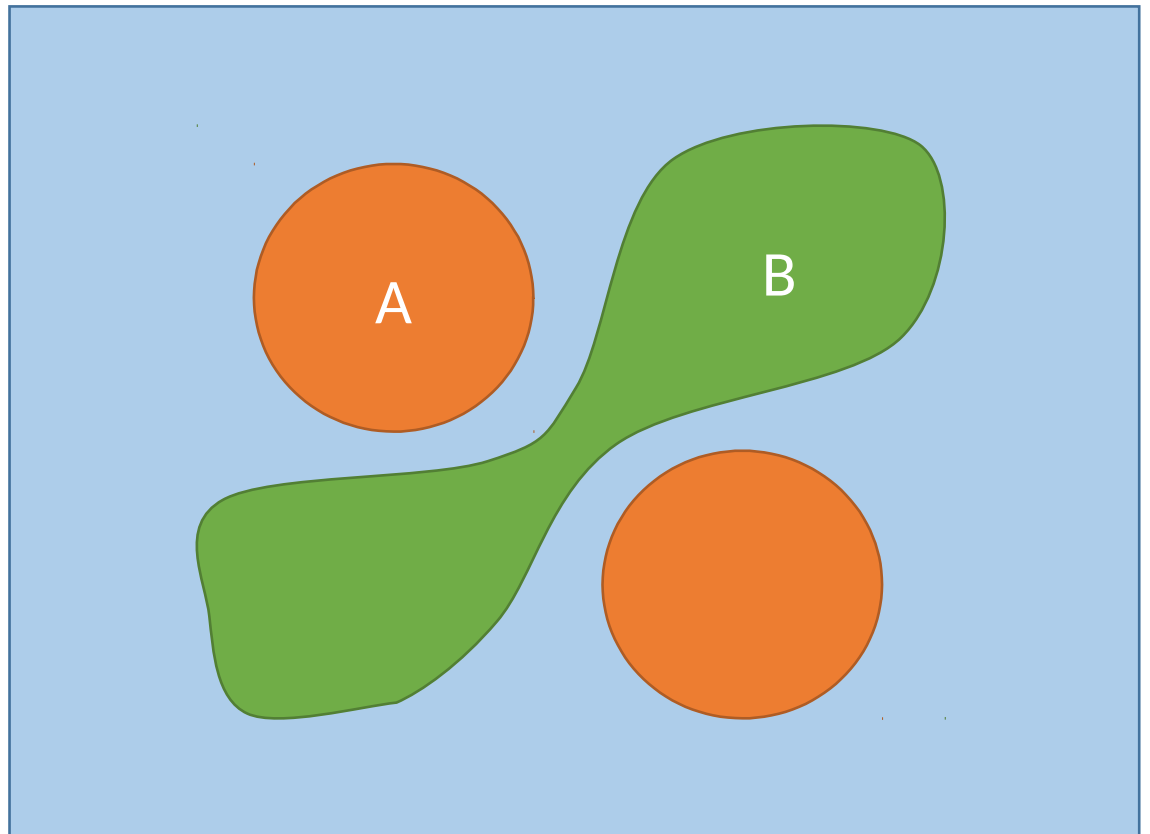
Branch and Prune ICP

Algorithm 1 $\text{ICP}(c_1, \dots, c_m, \vec{D} = D_1 \times \dots \times D_n, \delta)$

```
1:  $S.\text{push}(\vec{D})$ 
2: while  $S \neq \emptyset$  do
3:    $\vec{D} \leftarrow S.\text{pop}()$ 
4:   while  $\exists 1 \leq i \leq m, \vec{D} \neq_\delta \text{Prune}(\vec{D}, c_i)$  do
5:      $\vec{D} \leftarrow \text{Prune}(\vec{D}, c_i)$ 
6:   end while
7:   if  $\vec{D} \neq \emptyset$  then
8:     if  $\exists 1 \leq i \leq n, |D_i| \geq \varepsilon$  then  $\triangleright \varepsilon$  is some  
computable factor of  $\delta$ 
9:        $\{\vec{D}_1, \vec{D}_2\} \leftarrow \text{Branch}(\vec{D}, i)$ 
10:       $S.\text{push}(\vec{D}_1)$ 
11:       $S.\text{push}(\vec{D}_2)$ 
12:     else
13:       return sat
14:     end if
15:   end if
16: end while
17: return unsat
```

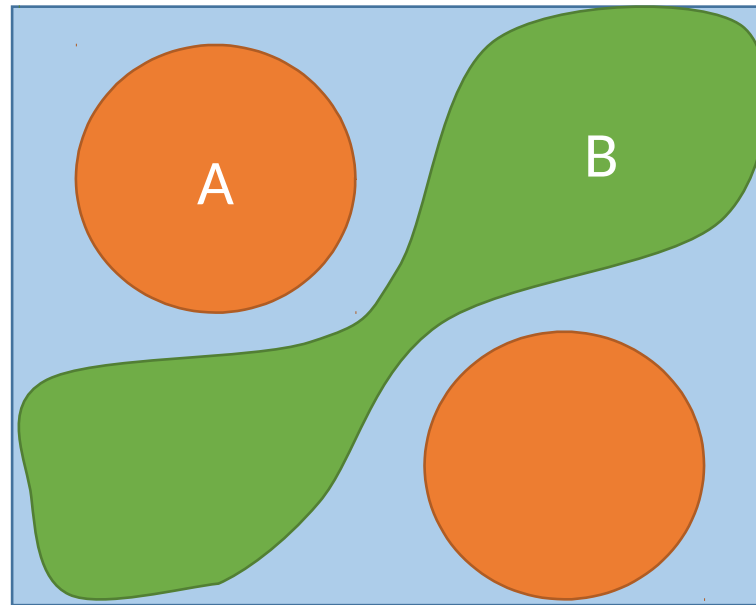


Branch-and-Prune Example



Branch-and-Prune Example

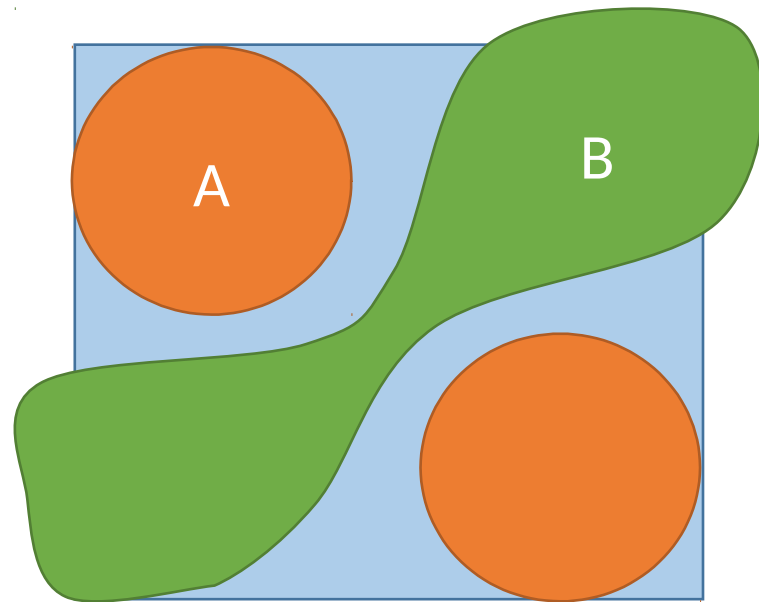
Prune by **B**



Branch-and-Prune Example

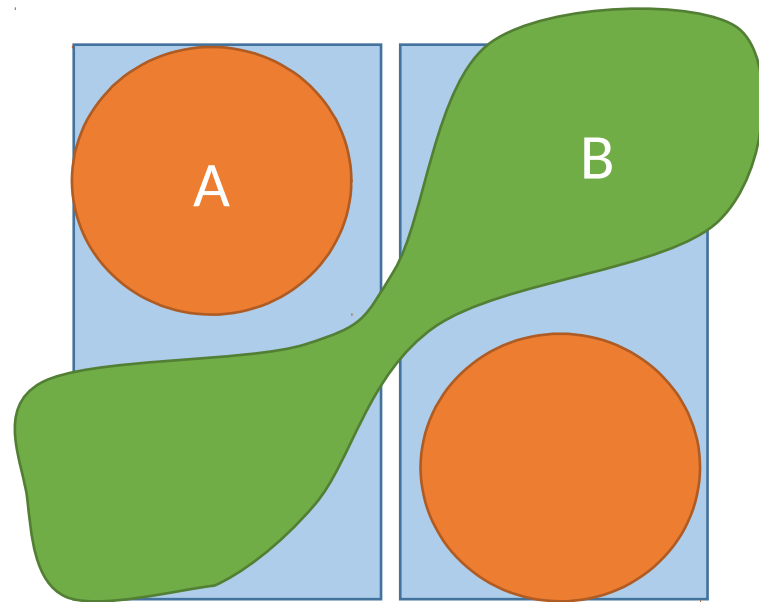
Prune by **B**

Prune by **A**



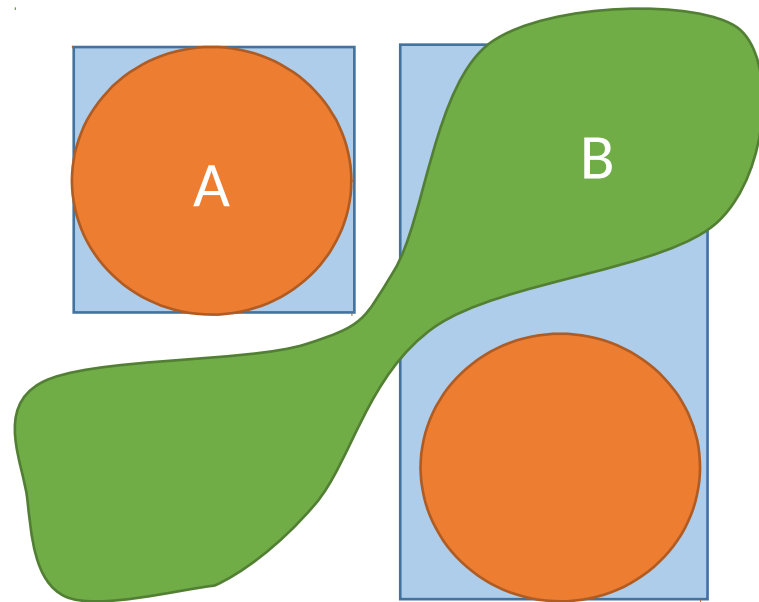
Branch-and-Prune Example

Prune by **B**
Prune by **A**
Branch



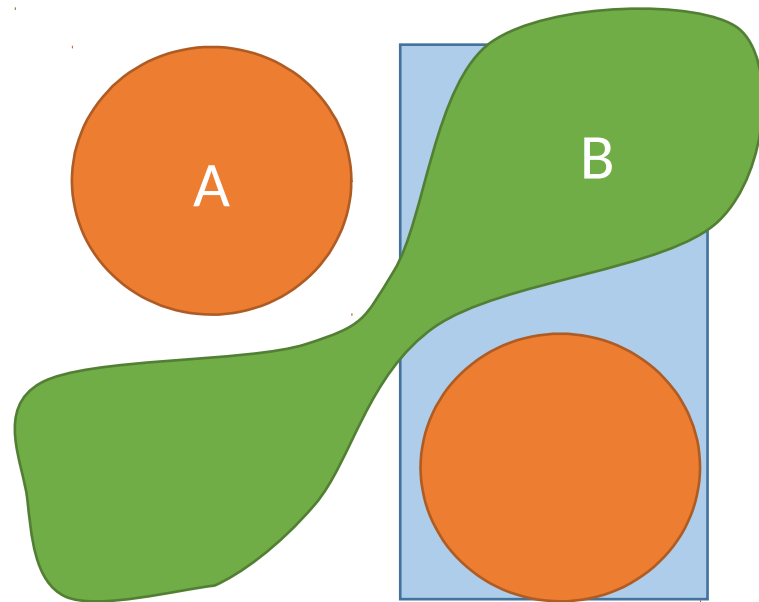
Branch-and-Prune Example

Prune by **B**
Prune by **A**
Branch
Prune by **A**



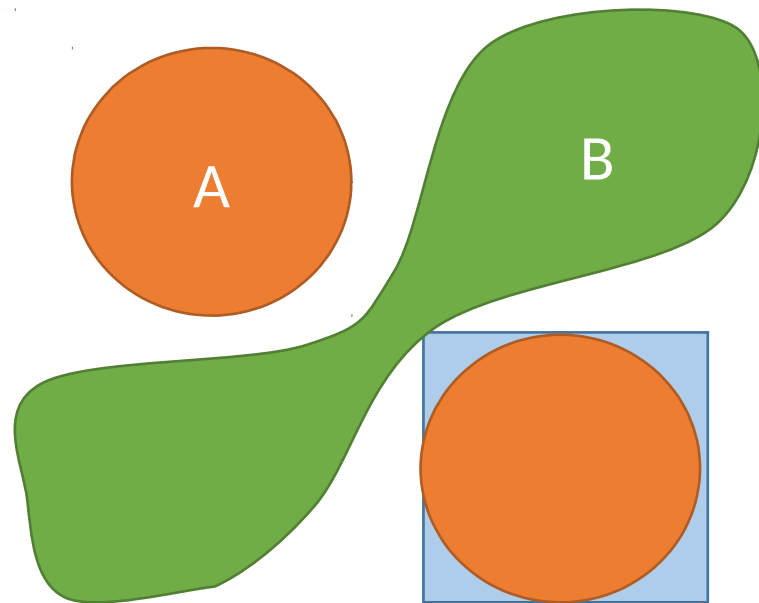
Branch-and-Prune Example

Prune by **B**
Prune by **A**
Branch
Prune by **A**
Prune by **B**



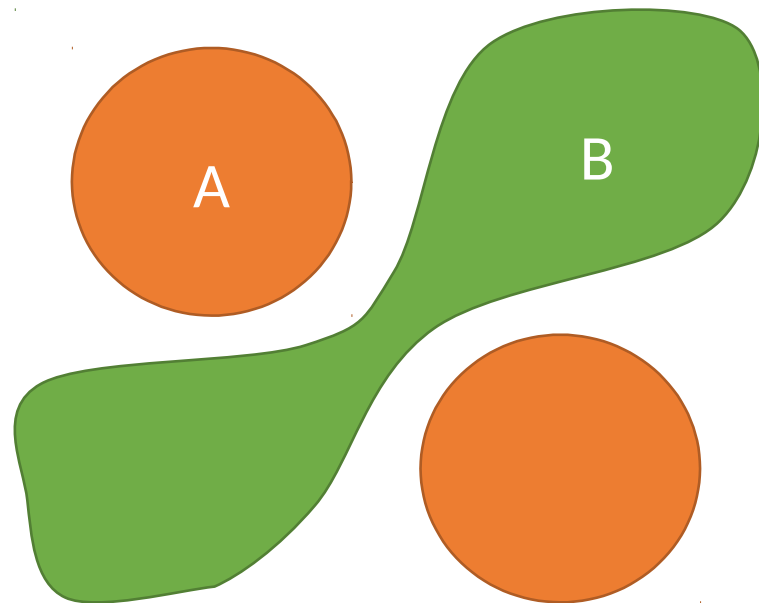
Branch-and-Prune Example

Prune by **B**
Prune by **A**
Branch
Prune by **A**
Prune by **B**
Prune by **A**



Branch-and-Prune Example

Prune by **B**
Prune by **A**
Branch
Prune by **A**
Prune by **B**
Prune by **A**
Prune by **B**



Completeness

- δ -satisfiability is NP (PSPACE with ODE).
- Idea:
 - If we can *guess a small enough box* containing the solution, we can check it in polynomial time using interval arithmetic.
 - If the problem is unsatisfiable, we need to explore a potentially exponential number of small boxes and show that all of them are empty.
- Takeaway message:

Nonlinear theories over the reals are *just* polynomially harder than SAT.

dReal

- Description: <http://dreal.github.io/>
- Getting the tool: <https://github.com/dreal/dreal3>
- GPL3 license
- Runs natively on Linux and Mac
- Runs on Windows via Docker

dReal Frontends

- SMT2

```
(set-logic QF_NRA)
(declare-fun x () Real)
(declare-fun y () Real)
(assert (< 2.4 x))
(assert (< x 2.6))
(assert (< -10.0 y))
(assert (< y 10.0))
(assert
  (and
    (= y (cos x))
  )
)
(check-sat)
(exit)
```

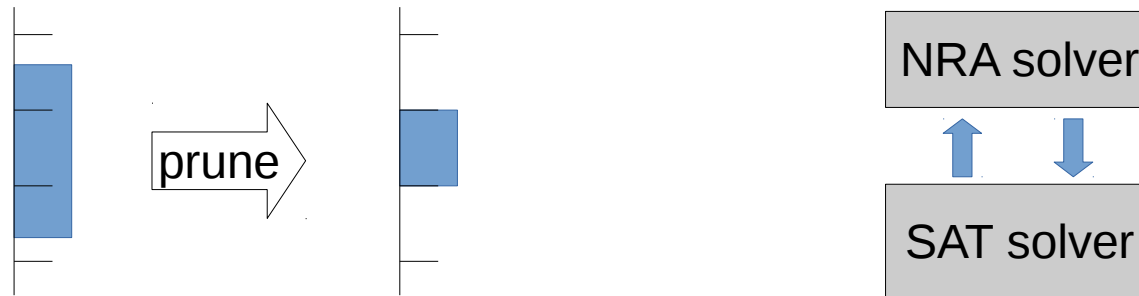
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```
var:
  [2.4, 2.6] x;
  [-10, 10] y;
ctr:
  y = cos(x);
```

dReal Example

What We Support

- Types: Real, Int, Bool
 - Int are handled in the ICP by a special contractor.
 - Bool are handled before the ICP by a SAT solver.



- Functions:
 - polynomials, trigonometric functions, logarithms, ...
 - (We will discuss very soon about the ODEs.)

ODEs and dReach

- dReal support ODEs directly in the SMT2 interface with a `QF_NRA_ODE` logic but the notation is non-standard.
- The dReach tool is much more user-friendly.
- dReach is a BMC that generates a dReal query from an hybrid automata

dReach Syntax

dReach Syntax

```
[0, 20] x;  
[-9.8] g;  
[-100, 100] v;  
[0, 10] time;
```

dReach Syntax

```
[0, 20] x;
[-9.8] g;
[-100, 100] v;
[0, 10] time;

{ mode 1;
  invt:
    (v <= 0);
    (x >= 0);
  flow:
    d/dt[x] = v;
    d/dt[v] = g;
  jump:
    (x = 0) ==>
    @2 (and (x' = x)
           (v' = (0 - v)));
}

{ mode 2;
  invt:
    (v >= 0);
    (x >= 0);
  flow:
    d/dt[x] = v;
    d/dt[v] = g;
  jump:
    (v = 0) ==>
    @1 (and (x' = x)
           (v' = v));
}
```

dReach Syntax

```
[0, 20] x;
[-9.8] g;
[-100, 100] v;
[0, 10] time;

{ mode 1;
  invt:
    (v <= 0);
    (x >= 0);
  flow:
    d/dt[x] = v;
    d/dt[v] = g;
  jump:
    (x = 0) ==>
    @2 (and (x' = x)
           (v' = (0 - v)));
}

{ mode 2;
  invt:
    (v >= 0);
    (x >= 0);
  flow:
    d/dt[x] = v;
    d/dt[v] = g;
  jump:
    (v = 0) ==>
    @1 (and (x' = x)
           (v' = v));
}

init:
  @1 (and (x = 10) (v = 0));
goal:
  @2 (and (x = 1) (v >= 1));
```

dReach Example

SMT Encoding (1)

- Variables

```
(declare-fun mode_1 () Real)
(declare-fun time_1 () Real)
(declare-fun x_1_0 () Real)
(declare-fun x_1_t () Real)
(declare-fun v_1_0 () Real)
(declare-fun v_1_t () Real)
```

- Mode invariants

```
(assert (and
  (forall_t 1 [0 time_1] (>= x_1_t 0) (<= v_1_t 0))
  (forall_t 2 [0 time_1] (>= x_1_t 0) (>= v_1_t 0))
))
```


SMT Encoding (2)

- Flow declaration

```
(declare-fun x () Real)
(declare-fun v () Real)
(define-ode flow_1 (
  (= d/dt[x] v)
  (= d/dt[v] g) ))
(define-ode flow_2 (
  (= d/dt[x] v)
  (= d/dt[v] g) ))
```

- Jump conditions

```
(assert (or (and (= mode_i 1) (= mode_j 2) (= x_i_t 0)
  (= x_j_0 x_i_t) (= v_j_0 (- v_i_t)))
  (and (= mode_i 2) (= mode_j 1) (= v_i_t 0)
  (= x_j_0 x_i_t) (= v_j_0 v_i_t))))
```

SMT Encoding (3)

- Connecting the flows

```
(assert (or
  (and (= mode_i 1)
    (= [x_i_t v_i_t] (integral 0. time_i [x_i_0 v_i_0] flow_1)))
  (and (= mode_i 2)
    (= [x_i_t v_i_t] (integral 0. time_i [x_i_0 v_i_0] flow_2))))
))
```

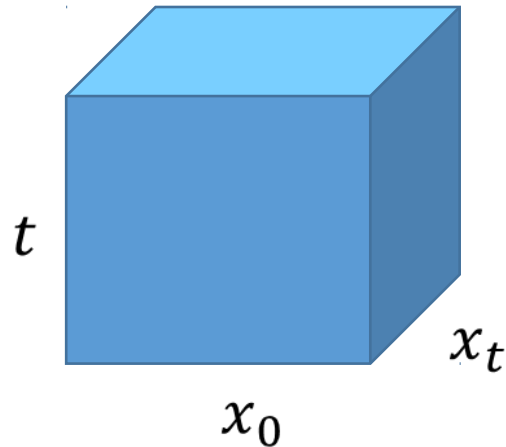
- Other elements

- Initial and final conditions
- Bounds for all the variables
- ...

ODEs, dReal, and Completeness

$$x_t = x_0 + \int_0^t f(x) dx \wedge 0 \leq t \leq 2$$

is just a pruning operator over the domain



dReal Tricks

- Julia bindings, C API, etc.
- Precision (δ)
 - Option: `--precision 0.1`
 - In SMT file: `(set-option :precision 0.1)`
- Model Generation
 - Option: `--model`
- Polytope contractor
 - Option: `--polytope`
- Branching heuristics
 - Options: `--gradbranch, --scoring-icp`

What Comes Next

- More efficient search heuristics (!!!)
- $\exists \forall$ formula
- More parallelism
- ...

Conclusion

- dReal is an SMT solver for nonlinear theories over the reals
- dReach is a bounded model checker for hybrid systems. dReach uses dReal as backend.
- If you have questions, contact us by email, open issues on github. Pull-requests on github are also welcome.