

# Continuous Relations and Richardson's Theorem

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# Background (Richardson's theorem)

Richardson, D.: Tassellations with local transformations. J. Computer and System Sciences 6, 373–388 (1972)



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- characterisation of transition **relations** determined by local rule of CA

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- characterisation of transition **relations** determined by local rule of CA
- using a notions of “topology” and “continuous relations” on the set of configurations.

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# Problems (Richardson's theorem)

“Topological” notions are introduced **indirectly**  
(via points of accumulations)

Richardson, D.: Tassellations with local transformations. J. Computer and System Sciences 6, 373–388 (1972)



# References for Solution

Continuity of relations has been studied in

Brattka, V. and Hertling, P.: Continuity and computability of relations. Informatik Berichte vol. 164. FernUniversität in Hagen (1994)

Ziegler, M.: Relative computability and uniform continuity of relations. Logic seminar 2011, Technische Universität Darmstadt (2011).



# What we have done

Reformulation of

- continuity of relations due to Brattke-Hertling
- nondeterministic CAs

by morphisms of *Rel* (relations)



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Proof of Richardson's theorem using topological and relational devices





# Notations

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$$\mathit{Rel}(X, Y) = \{\alpha \mid \alpha \subseteq X \times Y\}$$



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$$[\alpha] = \alpha\alpha^\sharp \sqcap \mathit{id}_X \text{ (domain of } \alpha: X \rightarrow Y)$$



# Basic Notions

$\alpha: X \rightarrow Y$

$\alpha: \text{univalent} \iff \alpha \# \alpha \sqsubseteq \text{id}_Y$   
(  $\forall x \in X \forall y, y' \in Y. (x, y) \in \alpha \wedge (x, y') \in \alpha \implies y = y'$  )



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$\alpha: \text{rectangular} \iff \alpha \nabla_Y X \alpha \sqsubseteq \alpha$   
(  $\exists W \sqsubseteq X \exists Z \sqsubseteq Y. \alpha = W \times Z$  )



# Elements

$X$ : set,  $I$ : singleton set

Elements of  $X$  is identified with tfn from  $I$  to  $X$



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$$y \in \{x, y, z\} \iff \begin{array}{c} I \longrightarrow X \\ \begin{array}{|c|} \hline \\ \hline \end{array} \begin{array}{c} x \\ y \\ z \end{array} \\ * \longmapsto y \end{array}$$





# Subsets

$X$ : set,  $I$ : singleton set

Subsets of  $X$  is identified with relations in  
 $Rel(I, X)$

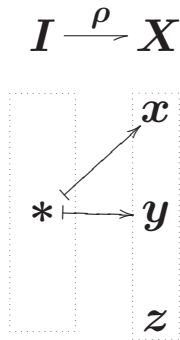


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$$X = \{x, y, z\}$$
$$\{x, y\} \subseteq X$$



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Subsets of  $X$  is identified with subidentities in  
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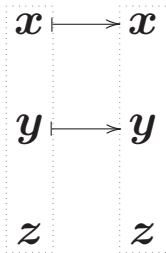
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$$X = \{x, y, z\} \\ \{x, y\} \subseteq X \quad \iff$$

$$X \xrightarrow{u} X$$



# Topological Space

$X$ : set,  $I$ : singleton set

$(X, \mathcal{O}(X))$



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$\mathcal{O}(X) \subseteq \underline{\text{Rel}}(I, X)$  closed under

- arbitrary union
- finite intersection



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Note:  $0_{IX}, \nabla_{IX} \in \mathcal{O}(X)$



# Continuous Relation

$X, Y$ : topological spaces

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inverse image of a open set is the intersection  
of  $\alpha$ 's domain and some open set



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basis element of the usual product topology of  $X \times Y$



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basis elements of the usual product topology of  $X \times Y$

$\alpha: X \rightarrow Y$  is closed  $\iff \alpha^-$  is open



# Nondeterministic CA over Group

$Y^X$ : the set of mappings from a set  $X$  to a set  $Y$

$Q$ : nonempty finite set,  $G$ : group

- $q \in Q$ : *state*
- $c \in Q^G$ : *configuration*
- $\delta: Q^G \rightarrow Q^G$ : *transition relation*



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- $\delta: Q^G \rightarrow Q^G$ : transition relation

$$\begin{array}{r}
 \{0, 1\}^{F(\{g\})} \ni \begin{array}{cccccccccc} \dots & g^{-3} & g^{-2} & g^{-1} & g^0 & g^1 & g^2 & g^3 & \dots \\ \dots & 1 & 0 & 1 & 0 & 1 & 0 & 1 & \dots \end{array} \\
 \downarrow \delta \\
 \{0, 1\}^{F(\{g\})} \ni \begin{array}{cccccccccc} \dots & g^{-3} & g^{-2} & g^{-1} & g^0 & g^1 & g^2 & g^3 & \dots \\ \dots & * & 1 & * & 1 & * & 1 & * & \dots \end{array} \\
 \hspace{15em} (* = 0 \text{ or } 1)
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$U \subseteq G, x \in G$

$p_U: Q^G \rightarrow Q^U$  denotes the projection



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$$\downarrow p_{\{g^{-1}, g^1\}}$$

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$$\downarrow p_{g^{-1}}$$

$$\{0, 1\} \ni$$

$$\frac{g^{-1}}{1}$$



# Shift Function

$x \in G$

shift function  $t_x: Q^G \rightarrow Q^G$  is defined by

$$\forall a \in G. t_x p_a = p_{x^{-1}a}$$





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# Local Rule

$N$ : finite subset of  $G$

$\lambda: Q^N \rightarrow Q$  is called a *local rule*



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 \downarrow \lambda & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\
 \{0, 1\} & * & & & 0 & & 1 \\
 & & & & & & (* = 0 \text{ or } 1)
 \end{array}$$



# Transition Determined by Local Rule

$\lambda: Q^N \rightarrow Q$ : local rule

$\tau_\lambda: Q^G \rightarrow Q^G$  defined by

$$\tau_\lambda = \prod_{x \in G} t_{x-1} p_N \lambda p_x^\#$$

is called the *transition relation determined by  $\lambda$*



# Transition Determined by Local Rule

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 & & * & & & & 0 & & 1 & & & & & (* = 0 \text{ or } 1)
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 \downarrow \lambda & \downarrow & \downarrow & \downarrow & \downarrow & & \\
 \{0, 1\} & * & & 0 & 1 & & (* = 0 \text{ or } 1)
 \end{array}$$

$$\tau\lambda = \prod_{x \in G} t_{x-1} p_N \lambda p_x^\#$$

$$\begin{array}{l}
 \{0, 1\}^{F(\{g\})} \\
 \downarrow t_{g^2} \\
 \{0, 1\}^{F(\{g\})} \\
 \downarrow p_{\{g^{-1}, g^1\}} \\
 \{0, 1\}^{\{g^{-1}, g^1\}} \\
 \downarrow \lambda \\
 \{0, 1\} \\
 \downarrow p_{g^{-2}}^\# \\
 \{0, 1\}^{F(\{g\})}
 \end{array}
 \begin{array}{c}
 \frac{\dots \quad g^{-3} \quad g^{-2} \quad g^{-1} \quad g^0 \quad g^1 \quad g^2 \quad g^3 \quad \dots}{\dots \quad 1 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \quad 1 \quad \dots} \\
 \downarrow \\
 \frac{\dots \quad -1 \quad g^0 \quad g^1 \quad g^2 \quad g^3 \quad 4 \quad 5 \quad \dots}{\dots \quad 1 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \quad 1 \quad \dots} \\
 \downarrow \\
 \frac{g^{-1} \quad g^1}{1 \quad 1} \\
 \downarrow \\
 1 \\
 \downarrow \\
 \frac{\dots \quad g^{-3} \quad g^{-2} \quad g^{-1} \quad g^0 \quad g^1 \quad g^2 \quad g^3 \quad \dots}{\dots \quad * \quad 1 \quad * \quad * \quad * \quad * \quad * \quad \dots}
 \end{array}$$



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$$\begin{array}{ccccccc}
 \{0, 1\}^{\{g^{-1}, g^1\}} & \frac{g^{-1} \quad g^1}{0 \quad 0} & \frac{g^{-1} \quad g^1}{0 \quad 1} & \frac{g^{-1} \quad g^1}{1 \quad 0} & \frac{g^{-1} \quad g^1}{1 \quad 1} & & \\
 \downarrow \lambda & \downarrow * & \downarrow & \downarrow 0 & \downarrow 1 & & \\
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# Topology on $Q^G$

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(( $\cong \wp(Q)$ ) discrete topology)



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- $\mathcal{O}(Q) = \text{Rel}(I, Q)$   
( $(\cong \wp(Q))$  discrete topology)
- $\mathcal{O}(Q^G)$  is the least topology s.t. for each  $x \in G$   $p_x : Q^G \rightarrow Q$  is continuous,  
i.e.  $\forall \rho \in \mathcal{O}(Q). \rho p_x^\# \in \mathcal{O}(Q^G)$   
(product topology)



# Richardson's Theorem (1/2)

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 $\tau_\lambda: Q^G \rightarrow Q^G$  satisfies



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(commutative with shifts)



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(commutative with shifts)

- $\prod_{x \in G} \tau_\lambda p_x p_x^\# = \tau_\lambda$

(determines states of each cell independently)



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$\tau_\lambda: Q^G \rightarrow Q^G$  satisfies

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(commutative with shifts)

- $\prod_{x \in G} \tau_\lambda p_x p_x^\# = \tau_\lambda$

(determines states of each cell independently)

- closed and continuous



# Richardson's Theorem (2/2)

For transition  $\delta: Q^G \rightarrow Q^G$  s.t.

- $\forall x \in G. t_x \delta = \delta t_x$   
(commutative with shifts)
- $\prod_{x \in G} \delta p_x p_x^\# = \delta$   
(determines states of each cell independently)
- closed and continuous



# Richardson's Theorem (2/2)

For transition  $\delta: Q^G \rightarrow Q^G$  s.t.

- $\forall x \in G. t_x \delta = \delta t_x$   
(commutative with shifts)
- $\prod_{x \in G} \delta p_x p_x^\# = \delta$   
(determines states of each cell independently)
- closed and continuous

there exists a local rule  $\lambda: Q^N \rightarrow Q$  s.t.

$$\delta = \lfloor \delta \rfloor \tau_\lambda.$$

( $\delta$ : restriction of  $\tau_\lambda$ ) 



# Another Result by Richardson

If local rule  $\lambda: Q^N \rightarrow Q$  is a **tfn**, then

$\forall$  finite  $S \subseteq G$ .  $\tau_S$  is surjective

$\implies \tau_\lambda$  is surjective



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where  $\tau_S: Q^{SN} \rightarrow Q^S$  is defined by

$$\tau_S = \prod_{x \in G} p_{SN}^\# t_{x^{-1}} p_N \lambda p_x^\# p_S$$



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# Slight Extension

If local rule  $\lambda: Q^N \rightarrow Q$  is **total**, then


$\forall$  finite  $S \subseteq G$ .  $\tau_S$  is surjective

$\implies \tau_\lambda$  is surjective



# Conclusion

What we have done:

- Investigation of the continuity of relations between topological spaces using relational notation
- Development fundamental properties of continuous relations
- Proof of Richardson's theorem using relational and topological devices
- Extension of a tiny part of Richardson's results 

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What we are going to do:

- More investigation of nondeterministic CA



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Relational and topological devices developed in this research should be useful for it!

