## Evaluation Strategy

Strict (or eager) evaluation.
Also known as call-by-value
Given an expression, which is a function application

$$
f\left(E_{1}, \ldots, E_{n}\right)
$$

evaluate $E_{1}, \ldots, E_{n}$ and then apply $f$ to the resulting values.

Call-by-name:
Substitute the expressions $E_{1}, \ldots, E_{n}$ into the definition of $f$ and then evaluate the resulting expression.

## Lazy Evaluation

Also known as call-by-need.
Like call-by-name, but sub-expressions that appear more than once are not copied. Pointers are used instead.

Potentially more efficient, but difficult to implement.

Standard ML uses strict evaluation.

## Lists

A list is an ordered collection (of any length) of elements of the same type

- [1,2,4];
> val it = [1, 2, 4] : int list
- ["a" , "", "abc", "a"];
> val it = . . . : string list
- [[1], [],[2,3]];
> val it = . . . : int list list
- [];
> val it = [] : 'a list
- 1:: [2,3];
> val it = [1, 2, 3] : int list


## Lists

There are two kinds of list:

## nil or [] is the empty list

$\mathrm{h}:: \mathrm{t}$ is the list with head h and tail t
$::$ is an infix operator of type
fn : 'a * 'a list -> 'a list
$\left[x_{1}, \ldots, x_{n}\right]$ is shorthand for

$$
x_{1}::\left(\cdots\left(x_{n}:: \text { nil }\right) \cdots\right)
$$

## Built-in Functions 1

null
fn : 'a list -> bool
determines if a list is empty
hd
fn : 'a list -> 'a
gives the first element of the list
tl
fn : 'a list -> 'a list
gives the tail of the list

## Built-in Functions 2

length
fn : 'a list -> int
gives the number of elements in a list
rev
fn : 'a list -> 'a list
gives the list in reverse order
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appends two lists NB: infix!

## List Functions

fun null $1=$

$$
\text { if } 1=[] \text { then true else false; }
$$

or, using pattern matching:
fun null [] = true
| null (_::_) = false;
fun hd (x::_) = x;
fun tl (_::1) = l;
NB: these functions are builtin and do not need to be defined

## Recursive definitions

$$
\begin{aligned}
\text { fun length }[] & =0 \\
\mid ~ r l e n g t h ~(h:: t) & =1+\operatorname{rlength}(t) ;
\end{aligned}
$$

fun append ([], 1) = 1
| append (h::t, l) = h::append(t,l);
fun reverse [] $=$ []
| reverse (h::t) = reverse (t)@[h];

Purely recursive definitions can be very inefficient

## Iterative Definitions

$$
\begin{aligned}
& \text { fun addlen }([], \mathrm{n})=\mathrm{n} \\
& \mathrm{l} \text { addlen }(\mathrm{h}:: \mathrm{t}, \mathrm{n})=\text { addlen }(\mathrm{t}, \mathrm{n}+1) ; \\
& \mathrm{fn}: \text { : a list } * \text { int }->\text { int }
\end{aligned}
$$

$$
\text { fun ilength } 1=\text { addlen }(1,0) \text {; }
$$

$$
\text { fun revto }([], 1)=1
$$

| revto (h::t, l) = revto (t, h::l);
fn : 'a list * 'a list -> 'a list

## Library List Functions

load "List";

We can then use List. take, List. drop

$$
\begin{gathered}
\text { fun take }(k,[])=[] \\
\mid \text { take }(k, h:: t)=
\end{gathered}
$$

$$
\text { if } k>0 \text { then } h:: \text { take }(k-1, t)
$$

else [];

$$
\text { fun } \operatorname{drop}(k,[])=[]
$$

| drop (k, h::t) =

$$
\text { if } k>0 \text { then } \operatorname{drop}(k-1,1)
$$

else h::t;
mn : int * 'a list -> 'a list

