

## Cpo's and domains

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A **chain complete poset**, or **cpo** for short, is a poset  $(D, \sqsubseteq)$  in which all countable increasing chains  $d_0 \sqsubseteq d_1 \sqsubseteq d_2 \sqsubseteq \dots$  have least upper bounds,  $\bigsqcup_{n \geq 0} d_n$ :

$$\forall m \geq 0 . d_m \sqsubseteq \bigsqcup_{n \geq 0} d_n \quad (\text{lub1})$$

$$\forall d \in D . (\forall m \geq 0 . d_m \sqsubseteq d) \Rightarrow \bigsqcup_{n \geq 0} d_n \sqsubseteq d. \quad (\text{lub2})$$

A **domain** is a cpo that possesses a least element,  $\perp$ :

$$\forall d \in D . \perp \sqsubseteq d.$$

"lub" = least upper bound

## Continuity and strictness

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- If  $D$  and  $E$  are cpo's, the function  $f$  is **continuous** iff
  1. it is monotone, and
  2. it preserves lubs of chains, *i.e.* for all chains  $d_0 \sqsubseteq d_1 \sqsubseteq \dots$  in  $D$ , it is the case that

$$f\left(\bigsqcup_{n \geq 0} d_n\right) = \bigsqcup_{n \geq 0} f(d_n) \quad \text{in } E.$$

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NB  
 $f(d_0) \subseteq f(d_1)$   
 $\subseteq f(d_2)$   
 $\subseteq \dots$   
 'cos  $f$  monotone

NB  $\forall i. d_i \sqsubseteq \bigsqcup_{n \geq 0} d_n \xrightarrow{\text{monotonicity}} \forall i. f(d_i) \subseteq f(\bigsqcup_{n \geq 0} d_n)$   
 $\implies \bigsqcup_{i \geq 0} f(d_i) \subseteq f(\bigsqcup_{n \geq 0} d_n)$

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So given 1, for 2 just need  $f(\bigsqcup_{n \geq 0} d_n) \sqsubseteq \bigsqcup_{n \geq 0} f(d_n)$

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- If  $D$  and  $E$  have least elements, then the function  $f$  is **strict** iff  $f(\perp) = \perp$ .

## Tarski's Fixed Point Theorem

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Let  $f : D \rightarrow D$  be a continuous function on a domain  $D$ . Then

- $f$  possesses a least pre-fixed point, given by

$$\text{fix}(f) = \bigsqcup_{n \geq 0} f^n(\perp).$$

- Moreover,  $\text{fix}(f)$  is a fixed point of  $f$ , *i.e.* satisfies  $f(\text{fix}(f)) = \text{fix}(f)$ , and hence is the **least fixed point** of  $f$ .

where

$$\begin{cases} f^0(\perp) \triangleq \perp \\ f^{n+1}(\perp) \triangleq f(f^n(\perp)) \end{cases}$$

## Pre-fixed points

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Let  $D$  be a poset and  $f : D \rightarrow D$  be a function.

An element  $d \in D$  is a **pre-fixed point of  $f$**  if it satisfies  $f(d) \sqsubseteq d$ .

The least pre-fixed point of  $f$ , if it exists, will be written

$$\boxed{\text{fix}(f)}$$

It is thus (uniquely) specified by the two properties:

$$f(\text{fix}(f)) \sqsubseteq \text{fix}(f) \quad (\text{lfp1})$$

$$\forall d \in D. f(d) \sqsubseteq d \Rightarrow \text{fix}(f) \sqsubseteq d. \quad (\text{lfp2})$$

# Proof of Tarski's Theorem

$\perp \subseteq f(\perp)$  because  $\perp$  is least elt of  $D$



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so  $f(\perp) \subseteq f(f(\perp)) \triangleq f^2(\perp)$  by monotonicity of  $f$

so  $f^2(\perp) \subseteq f(f^2(\perp)) = f^3(\perp)$

etc.

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etc.

We get a chain  $\perp \sqsubseteq f(\perp) \sqsubseteq f^2(\perp) \sqsubseteq f^3(\perp) \sqsubseteq \dots$

and can form its lub  $\bigcup_{n \geq 0} f^n(\perp)$

# Proof of Tarski's Theorem

Applying  $f$  to  $\perp \sqsubseteq f(\perp) \sqsubseteq f^2(\perp) \sqsubseteq \dots \sqsubseteq \bigcup_{n \geq 0} f^n(\perp)$

we get

$$f(\perp) \sqsubseteq f(f(\perp)) \sqsubseteq f(f^2(\perp)) \sqsubseteq \dots \sqsubseteq f\left(\bigcup_{n \geq 0} f^n(\perp)\right)$$

by monotonicity of  $f$

# Proof of Tarski's Theorem

Applying  $f$  to  $\perp \sqsubseteq f(\perp) \sqsubseteq f^2(\perp) \sqsubseteq \dots \sqsubseteq \bigcup_{n \geq 0} f^n(\perp)$

we get  $f(\perp) \sqsubseteq f(f(\perp)) \sqsubseteq f(f^2(\perp)) \sqsubseteq \dots \sqsubseteq f(\bigcup_{n \geq 0} f^n(\perp))$

by continuity of  $f$   $\longrightarrow$   $\parallel$   
 $\bigcup_{n \geq 0} f(f^n(\perp))$   
 $\parallel$   
 $\bigcup_{n \geq 0} f^{n+1}(\perp)$   
 $\parallel$   
 $\bigcup_{m \geq 1} f^m(\perp)$

# Proof of Tarski's Theorem

So  $\bigcup_{n \geq 0} f^n(\perp)$   
is a (pre-)fixed  
point for  $f$

$$\begin{aligned} & f\left(\bigcup_{n \geq 0} f^n(\perp)\right) \\ & \parallel \\ & \bigcup_{n \geq 0} f(f^n(\perp)) \\ & \parallel \\ & \bigcup_{n \geq 0} f^{n+1}(\perp) \\ & \parallel \\ & \bigcup_{m \geq 1} f^m(\perp) \end{aligned}$$

$$\bigcup_{m \geq 0} f^m(\perp) = \bigcup_{m \geq 1} f^m(\perp)$$

# Proof of Tarski's Theorem

For any pre-fixed point  $f(d) \sqsubseteq d$  we have

$\perp \sqsubseteq d$  because  $\perp$  is least elt of  $D$

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So  $f(\perp) \sqsubseteq f(d) \sqsubseteq d$       monotonicity + )

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So  $f(\perp) \sqsubseteq f(d) \sqsubseteq d$  *monotonicity +*

So  $f^2(\perp) = f(f(\perp)) \sqsubseteq f(d) \sqsubseteq d$

etc.



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So  $f^2(\perp) = f(f(\perp)) \sqsubseteq f(d) \sqsubseteq d$

etc.

We get  $f^n(\perp) \sqsubseteq d$  for all  $n \geq 0$

So  $\bigcup_{n \geq 0} f^n(\perp) \sqsubseteq d$

# Proof of Tarski's Theorem

For any pre-fixed point  $f(d) \subseteq d$  we have

So  $\bigcup_{n \geq 0} f^n(\perp)$  is  
a least pre-fixed point

We get

$$\bigcup_{n \geq 0} f^n(\perp) \subseteq d$$

QED

## Fixed point property of [[while B do C]]

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$$[[\text{while } B \text{ do } C]] = f_{[[B]], [[C]]}([[ \text{while } B \text{ do } C ]])$$

where, for each  $b : \text{State} \rightarrow \{\text{true}, \text{false}\}$  and  
 $c : \text{State} \rightarrow \text{State}$ , we define

as  $f_{b,c} : (\text{State} \rightarrow \text{State}) \rightarrow (\text{State} \rightarrow \text{State})$

$$f_{b,c} = \lambda w \in (\text{State} \rightarrow \text{State}). \lambda s \in \text{State}.$$

$$\text{if } (b(s), w(c(s)), s).$$

*we now know this  
is a domain*

- 
- Why does  $w = f_{[[B]], [[C]]}(w)$  have a solution?
  - What if it has several solutions—which one do we take to be [[while B do C]]?

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● Why does  $w = f_{[[B]], [[C]]}(w)$  have a solution?

● What if it has several solutions—which one do we take to be

[[while B do C]]?

← least (pre-)fixed point

← Tarski's Theorem  
(need to show  $f_{b,c}$   
is continuous)

# Continuity of $f_{b,c}$

Suppose  $C_0 \subseteq C_1 \subseteq C_2 \subseteq \dots$  in  $\text{State} \rightarrow \text{State}$

$$f_{b,c} \left( \bigcup_{n \geq 0} C_n \right) = \lambda s \in \text{State}. \text{ if } (b(s), (\bigcup_{n \geq 0} C_n)(c(s)), s)$$

that is

$$f_{b,c} \left( \bigcup_{n \geq 0} C_n \right) = \left\{ (s, s') \mid \begin{array}{l} b(s) = \text{true} \wedge \exists s''. (c(s) = s'' \wedge \\ (\bigcup_{n \geq 0} C_n)(s'') = s') \\ \vee \\ b(s) = \text{false} \wedge s = s' \end{array} \right\}$$

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$$= \bigcup_{n \geq 0} \{ (s, s') \mid \text{if } (b(s), C_n(c(s)), s) = s' \}$$

$$= \bigcup_{n \geq 0} f_{b,c}(C_n)$$

QED



# [[while B do C]]

---

[[while B do C]]

$$= \text{fix}(f_{[[B]], [[C]])}$$

Tarski Theorem

$$\equiv \bigsqcup_{n \geq 0} f_{[[B]], [[C]]}^n(\perp)$$

$$\equiv \lambda s \in \text{State}.$$

requires proof ...

$$\left\{ \begin{array}{ll} [[C]]^k(s) & \text{if } k \geq 0 \text{ is such that } [[B]]([[C]]^k(s)) = \text{false} \\ & \text{and } [[B]]([[C]]^i(s)) = \text{true for all } 0 \leq i < k \\ \text{undefined} & \text{if } [[B]]([[C]]^i(s)) = \text{true for all } i \geq 0 \end{array} \right.$$

# Example

Domain  $D = (P(\mathbb{N}), \subseteq)$  (same as  $\mathbb{N} \rightarrow 1$ )

Function  $f : D \rightarrow D$

$$f(S) \triangleq \{0\} \cup \{x+2 \mid x \in S\}$$

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$$f(S) \triangleq \{0\} \cup \{x+2 \mid x \in S\}$$

$S \in D$  is a pre-fixed point of  $f$  if

$$f(S) \subseteq S$$

ie.  $0 \in S$  &  $x+2 \in S$  for all  $x \in S$

ie.  $S$  is closed under the rules  $\frac{}{0 \in S}$  &  $\frac{x \in S}{x+2 \in S}$

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i.e.  $S$  is closed under the rules  $\frac{}{0 \in S}$  &  $\frac{x \in S}{x+2 \in S}$

So expect least pre-fixed point of  $f$   
to be Even =  $\{2x \mid x \in \mathbb{N}\}$

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$f$  is monotone :  $S \subseteq S' \Rightarrow f(S) \subseteq f(S')$  ✓

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$$f(S) \triangleq \{0\} \cup \{x+2 \mid x \in S\}$$

$f$  is monotone :  $S \subseteq S' \Rightarrow f(S) \subseteq f(S')$  ✓

$f$  is continuous :  $f(\bigcup_{n \geq 0} S_n) = \{0\} \cup \{x+2 \mid x \in \bigcup_{n \geq 0} S_n\}$   
 $= \{0\} \cup \bigcup_{n \geq 0} \{x+2 \mid x \in S_n\}$   
 $= \bigcup_{n \geq 0} f(S_n)$  ✓

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Function  $f : D \rightarrow D$

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Tarski's Theorem applies:

$$\text{fix}(f) = \bigcup_{n \geq 0} f^n(\emptyset)$$

$$f(\emptyset) = \{0\}$$

$$f^2(\emptyset) = \{0\} \cup \{0+2\}$$

$$f^3(\emptyset) = \{0, 2, 4\}$$

$$\vdots$$
$$f^n(\emptyset) = \{0, 2, 4, \dots, 2(n-1)\}$$

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Tarski's Theorem applies:

$$\text{fix}(f) = \bigcup_{n \geq 0} f^n(\emptyset) = \{0, 2, 4, 8, \dots\}$$

$$= \{2x \mid x \in \mathbb{N}\}$$

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(as expected).