

5.2 Fibonacci Heaps

Frank Stajano

Thomas Sauerwald

Lent 2016



Operation	Linked list	Binary heap	Binomial heap
Μаке-Неар	<i>O</i> (1)	<i>O</i> (1)	<i>O</i> (1)
INSERT	$\mathcal{O}(1)$	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$
Мілімим	$\mathcal{O}(n)$	$\mathcal{O}(1)$	$\mathcal{O}(\log n)$
Extract-Min	$\mathcal{O}(n)$	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$
Merge	$\mathcal{O}(n)$	$\mathcal{O}(n)$	$\mathcal{O}(\log n)$
DECREASE-KEY	$\mathcal{O}(1)$	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$
DELETE	$\mathcal{O}(1)$	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$



Operation	Linked list	Binary heap	Binomial heap	Fibon. heap
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T.S.

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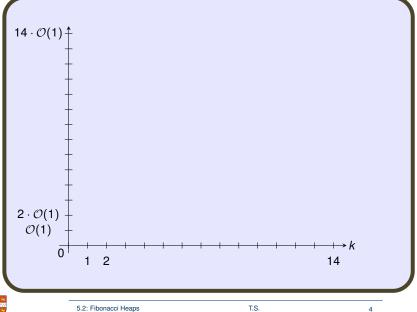
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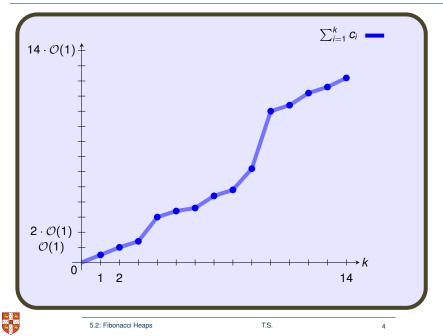
Fibonacci Heap: k/2DECREASE-KEY + k/2 INSERT

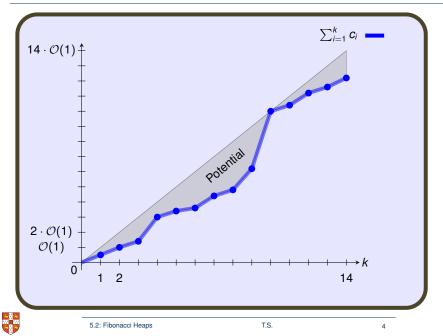
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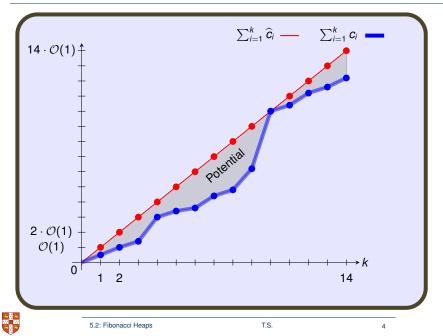
$$\Rightarrow \sum_{i=1}^{k} c_i \leq \sum_{i=1}^{k} \widehat{c}_i = \mathcal{O}(k)$$

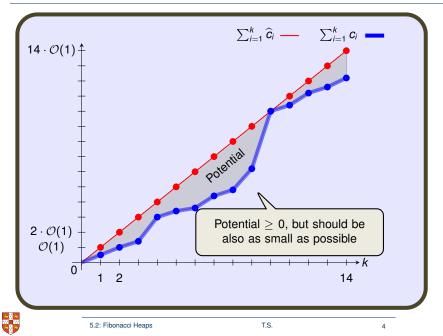












Structure

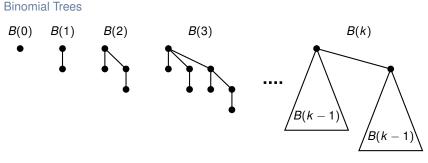
Operations

Glimpse at the Analysis

Amortized Analysis



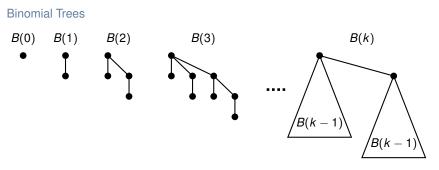
Reminder: Binomial Heaps



Binomial Heaps -

 Binomial Heap is a collection of binomial trees of different orders, each of which obeys the heap property

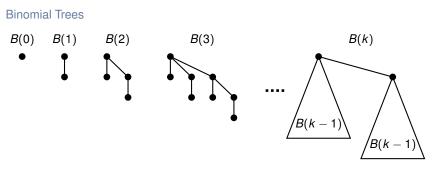




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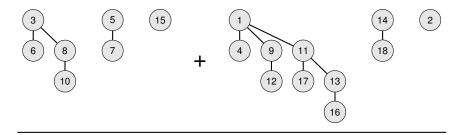




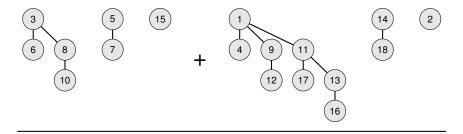
Binomial Heaps

- Binomial Heap is a collection of binomial trees of different orders, each of which obeys the heap property
- Operations:
 - MERGE: Merge two binomial heaps using Binary Addition Procedure
 - INSERT: Add B(0) and perform a MERGE
 - EXTRACT-MIN: Find tree with minimum key, cut it and perform a MERGE
 - DECREASE-KEY: The same as in a binary heap

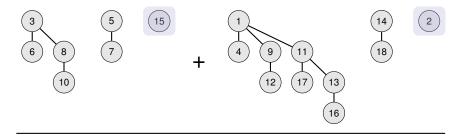




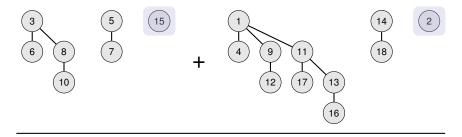






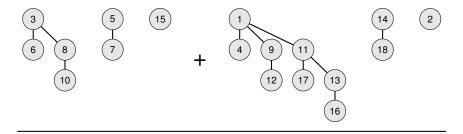






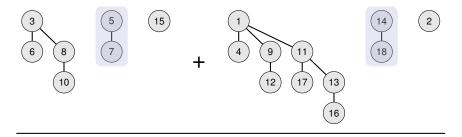




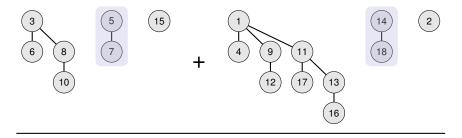








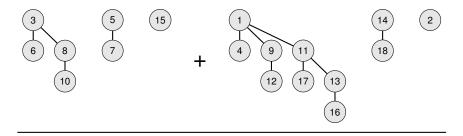


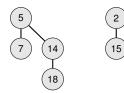




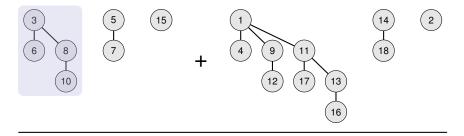


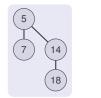








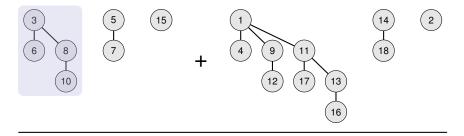


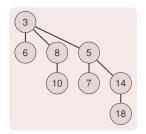




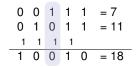
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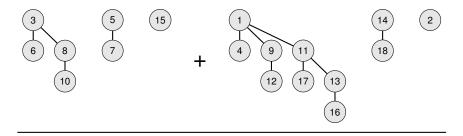


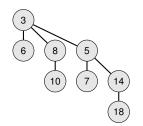


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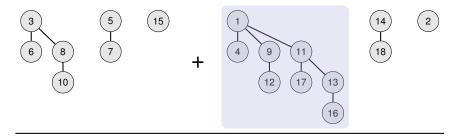


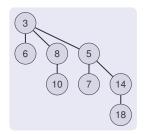




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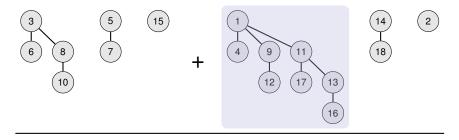
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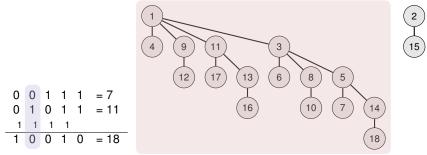




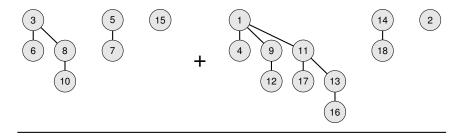
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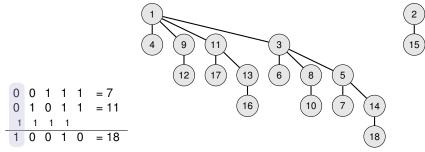




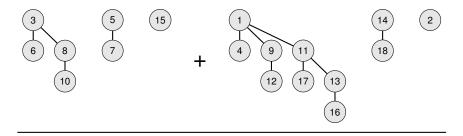


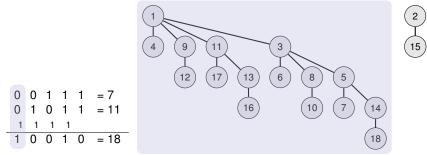






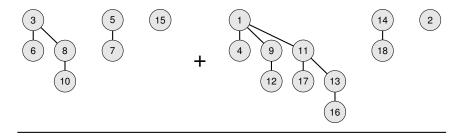


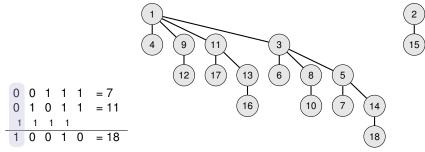






Merging two Binomial Heaps

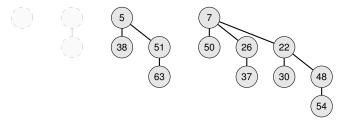






Binomial Heap:

- consists of binomial trees, and every order appears at most once
- immediately tidy up after INSERT or MERGE

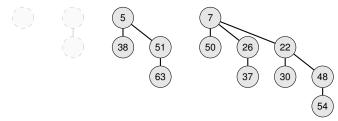




Binomial Heap vs. Fibonacci Heap: Structure

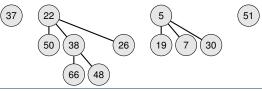
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Fibonacci Heap:

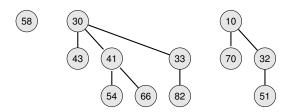
- forest of MIN-HEAPs
- Iazily defer tidying up; do it on-the-fly when search for the MIN







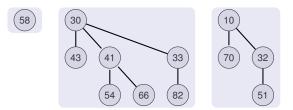
Forest of MIN-HEAPs





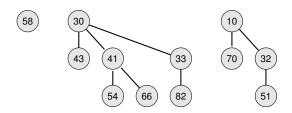
Fibonacci Heap _____

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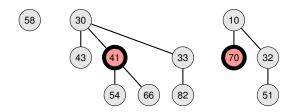


- Forest of MIN-HEAPs
- Nodes can be marked (roots are always unmarked)



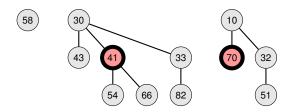


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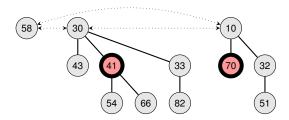


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- Tree roots are stored in a circular, doubly-linked list





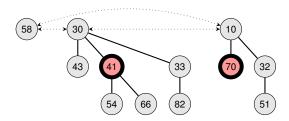
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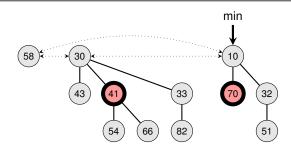
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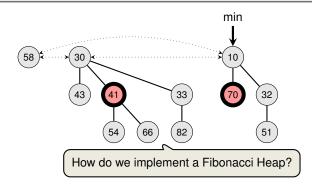
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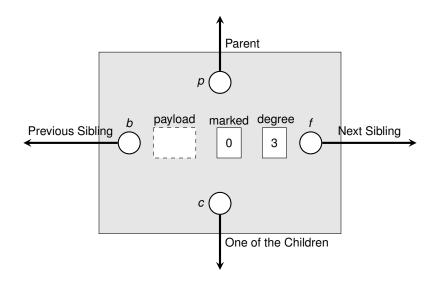


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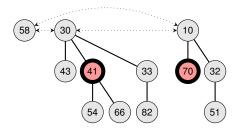






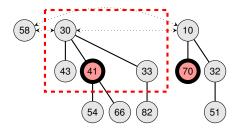


Magnifying a Four-Node Portion



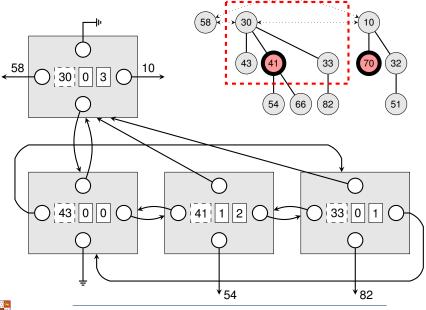


Magnifying a Four-Node Portion





Magnifying a Four-Node Portion





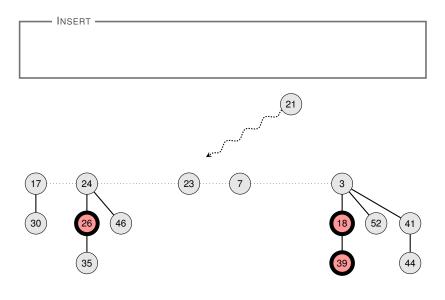
Structure

Operations

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Amortized Analysis

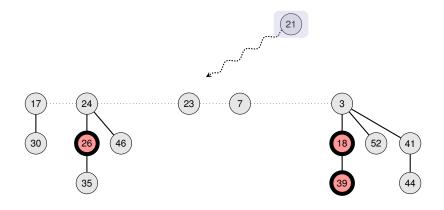






INSERT -

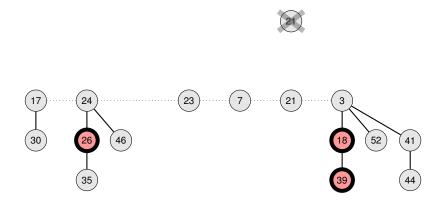
Create a singleton tree





INSERT -

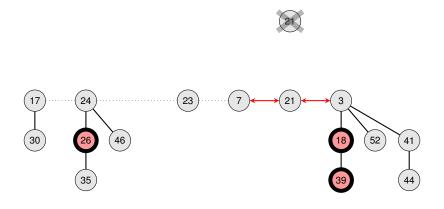
- Create a singleton tree
- Add to root list





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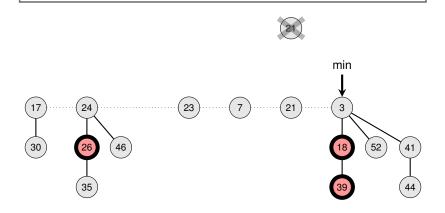
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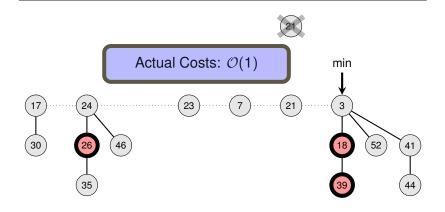
- Create a singleton tree
- Add to root list and update min-pointer (if necessary)





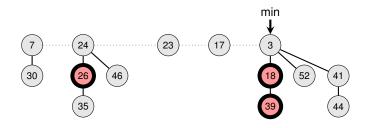
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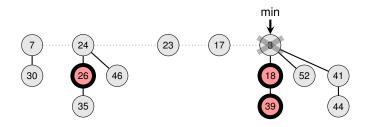






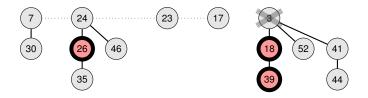






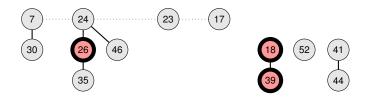






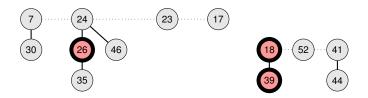


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them



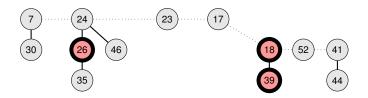


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them



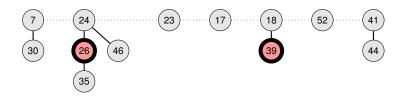


- Extract-Min ———
- Delete min \checkmark
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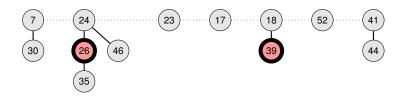


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark



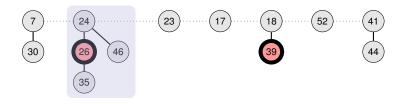


- Extract-Min –
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree



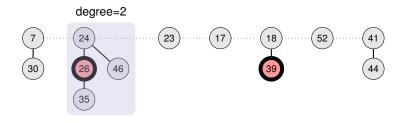


- Extract-Min -
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



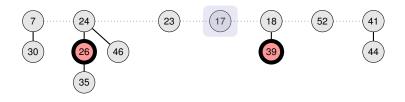


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



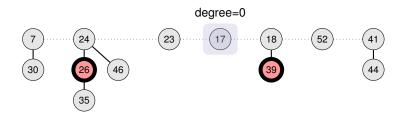


- Extract-Min -
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



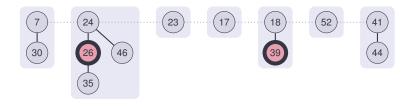


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



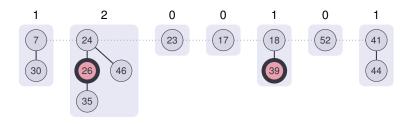


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



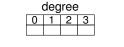


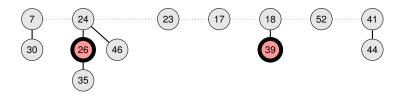
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- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)





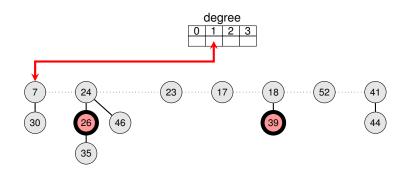
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- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)





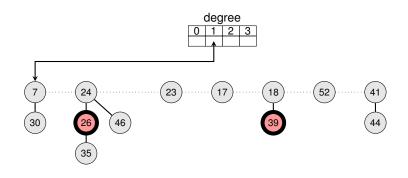


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



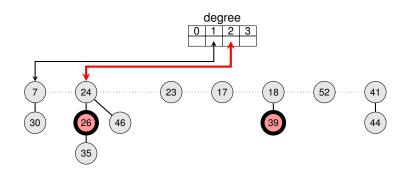


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



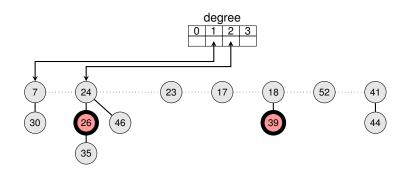


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



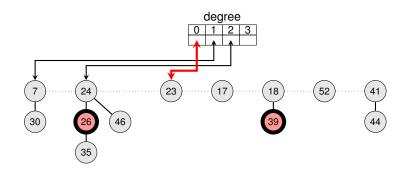


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



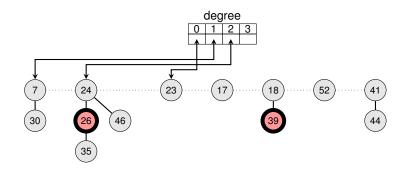


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



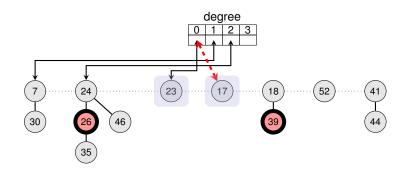


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
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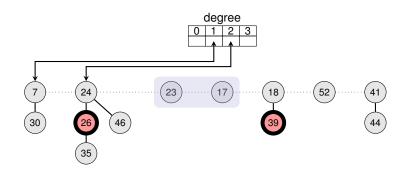


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



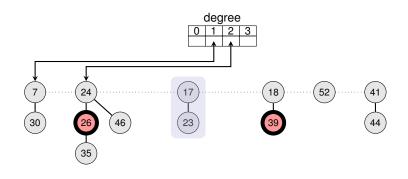


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- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



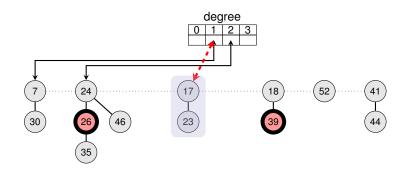


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



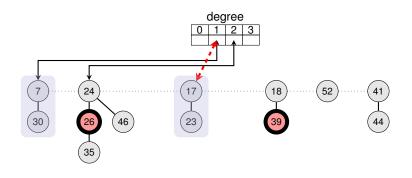


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



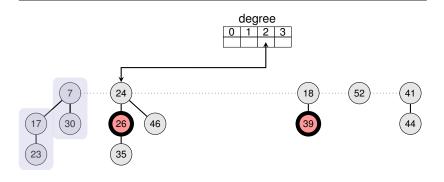


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



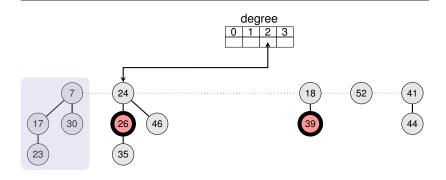


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



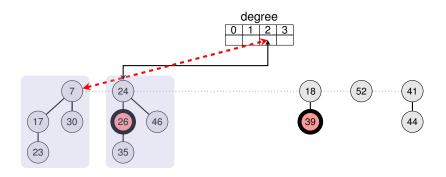


- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)





- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)





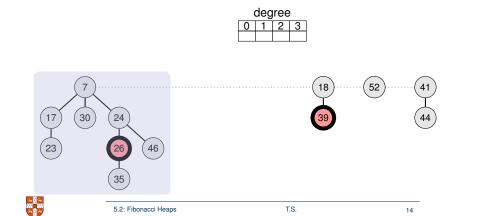
- Extract-Min ——
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



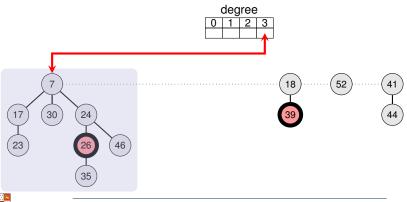




- Extract-Min ——
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)

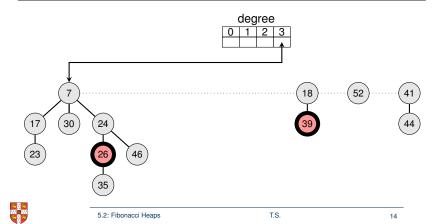


- Extract-Min ——
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)

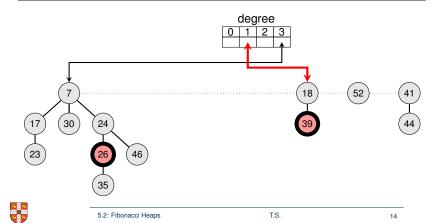




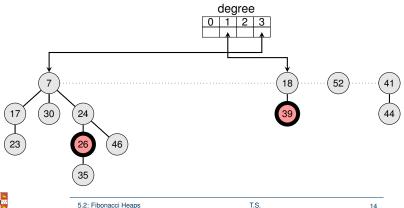
- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)

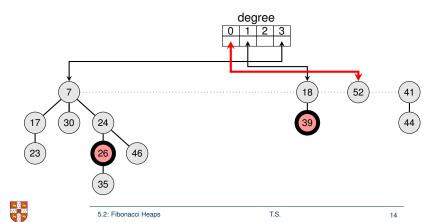


- Extract-Min ———
- Delete min √
- Meld childen into root list and unmark them √
- Consolidate so that no roots have the same degree (# children)

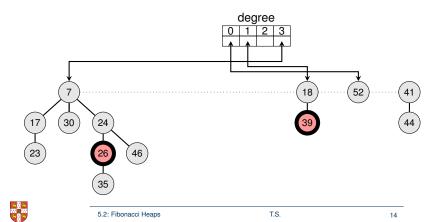




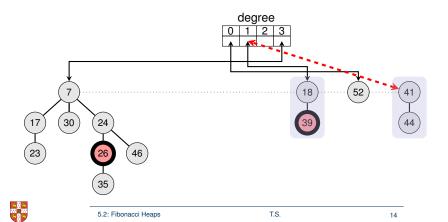
- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



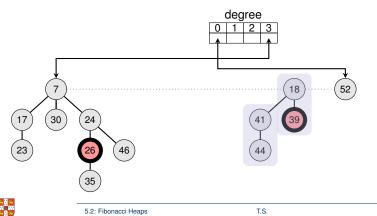
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- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



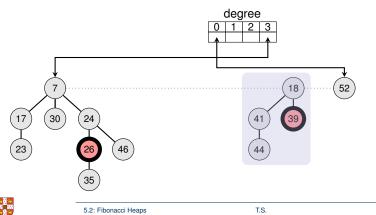
- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



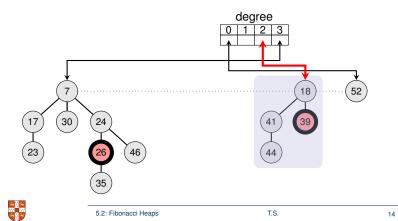
- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



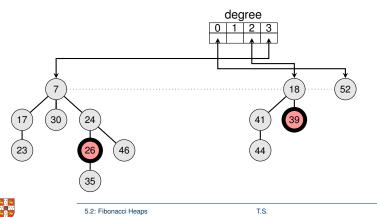
- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



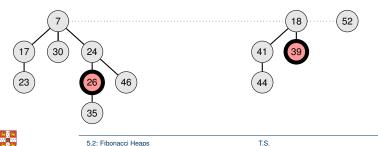
- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children)



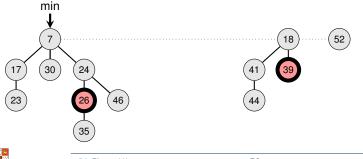
- Extract-Min ———
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children) \checkmark



- Extract-Min ——
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children) \checkmark
- Update minimum

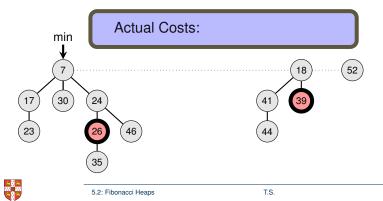


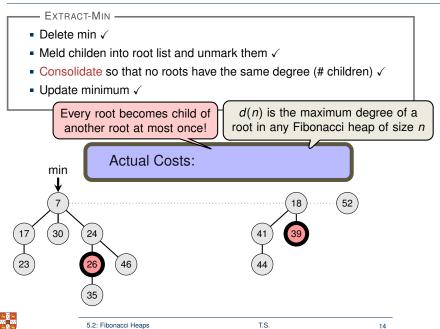
- Extract-Min ——
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children) \checkmark
- Update minimum \checkmark

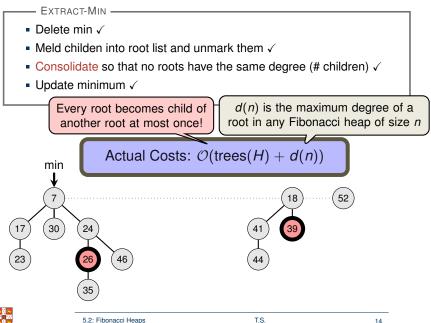




- Extract-Min –
- Delete min \checkmark
- Meld childen into root list and unmark them \checkmark
- Consolidate so that no roots have the same degree (# children) \checkmark
- Update minimum \checkmark



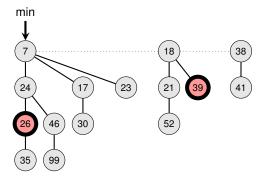




Fibonacci Heap: DECREASE-KEY (First Try)

DECREASE-KEY of node x —

Decrease the key of x (given by a pointer)

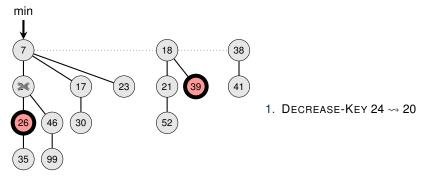




Fibonacci Heap: DECREASE-KEY (First Try)



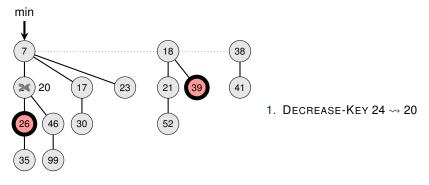
Decrease the key of x (given by a pointer)





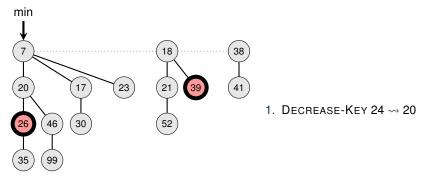


Decrease the key of x (given by a pointer)



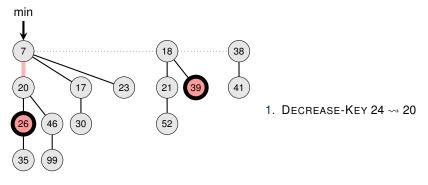


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated





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- Check if heap-order is violated

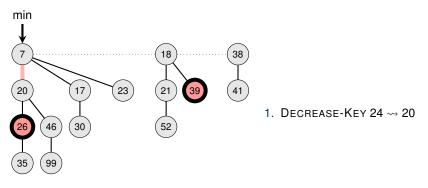




DECREASE-KEY of node x —

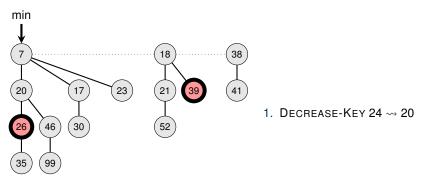
- Decrease the key of x (given by a pointer)
- Check if heap-order is violated

If not



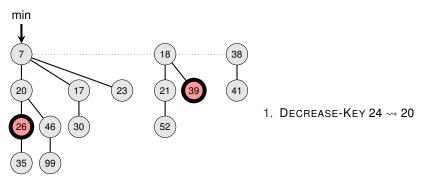


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
 - If not, then done.



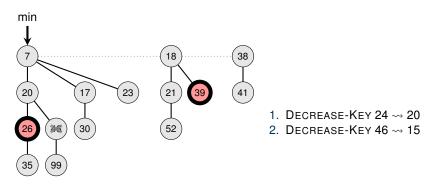


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
 - If not, then done.
 - Otherwise,



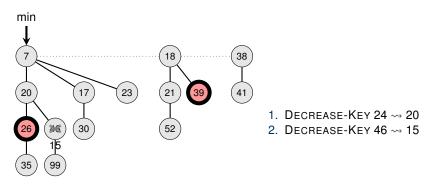


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
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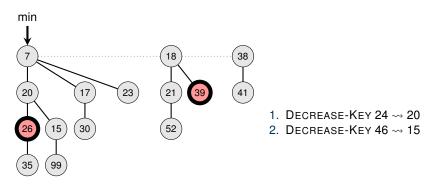


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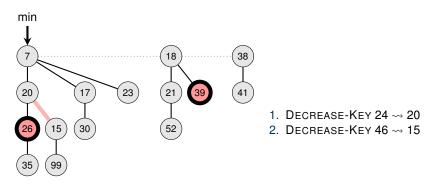


- Decrease the key of x (given by a pointer)
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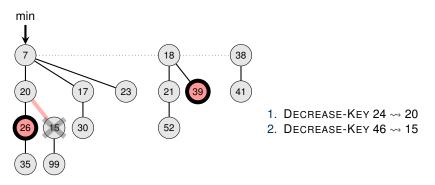


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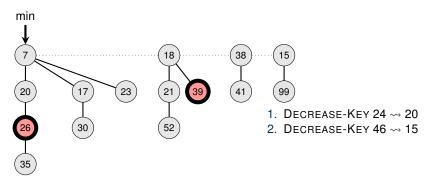


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
 - If not, then done.
 - Otherwise, cut tree rooted at *x* and meld into root list (update min).



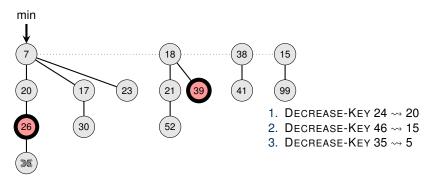


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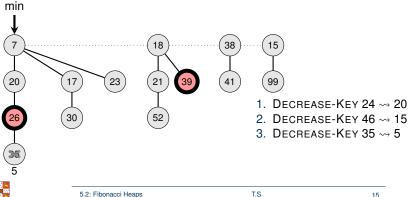


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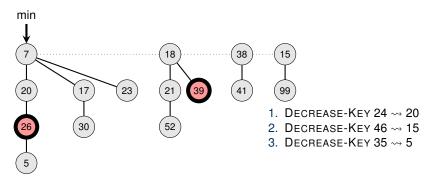




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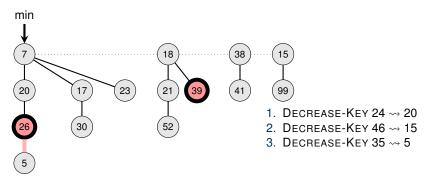


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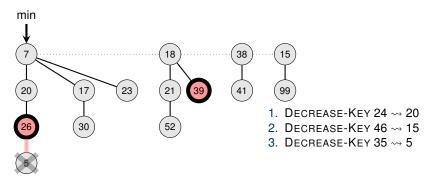


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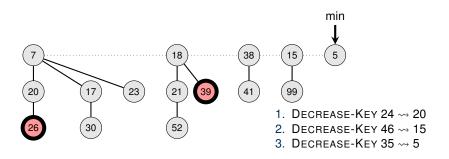


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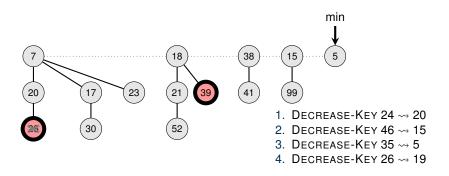


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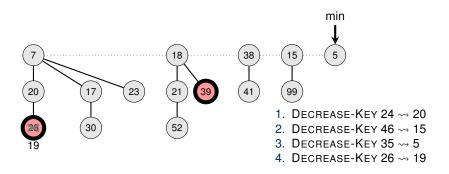


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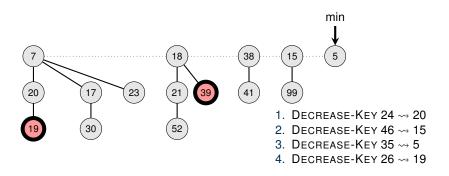


- Decrease the key of x (given by a pointer)
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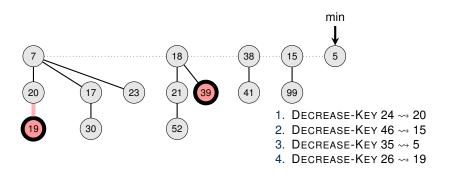


- Decrease the key of x (given by a pointer)
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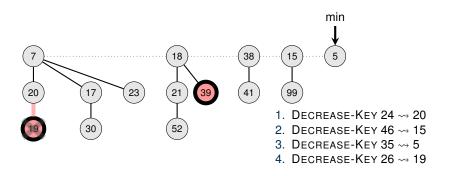


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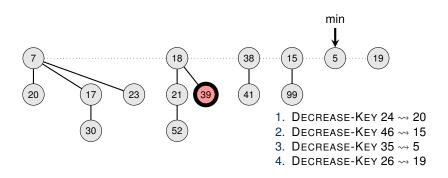


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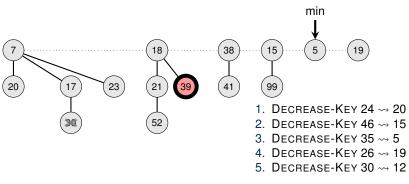


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- Check if heap-order is violated
 - If not, then done.
 - Otherwise, cut tree rooted at *x* and meld into root list (update min).



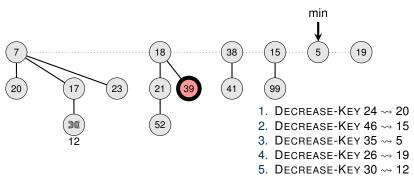


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
 - If not, then done.
 - Otherwise, cut tree rooted at *x* and meld into root list (update min).



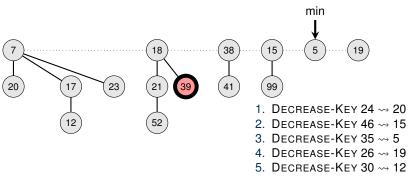


- Decrease the key of x (given by a pointer)
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 - If not, then done.
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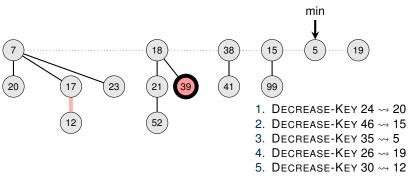


- Decrease the key of x (given by a pointer)
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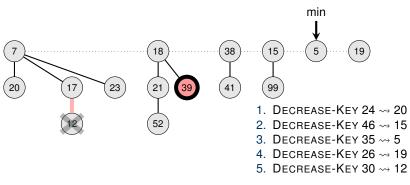


- Decrease the key of x (given by a pointer)
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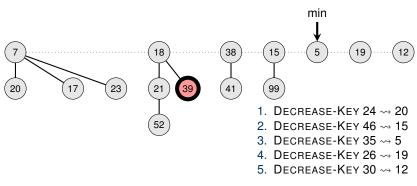


- Decrease the key of x (given by a pointer)
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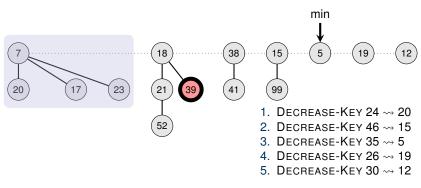


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
 - If not, then done.
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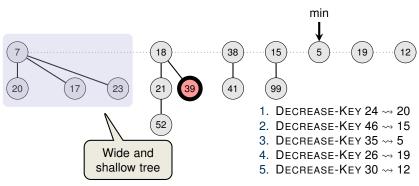


- Decrease the key of x (given by a pointer)
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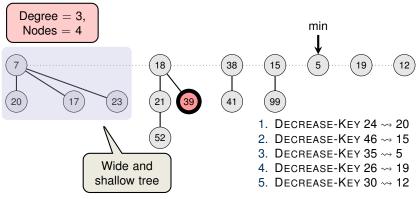


- Decrease the key of x (given by a pointer)
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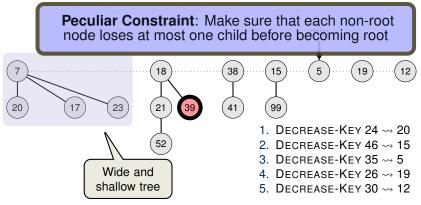


- Decrease the key of x (given by a pointer)
- Check if heap-order is violated
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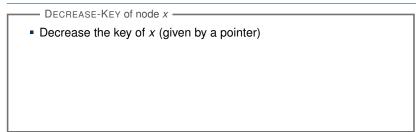


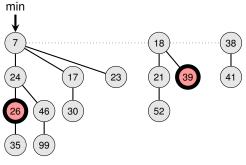
- Decrease the key of x (given by a pointer)
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 - If not, then done.
 - Otherwise, cut tree rooted at *x* and meld into root list (update min).





Fibonacci Heap: DECREASE-KEY

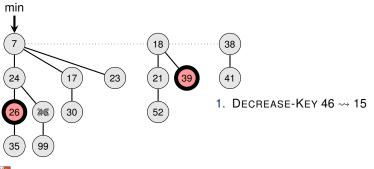






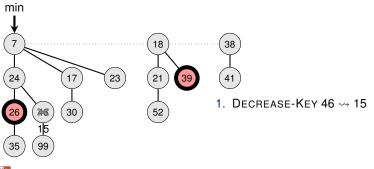
Fibonacci Heap: DECREASE-KEY

- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)



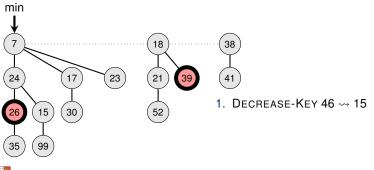


- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list



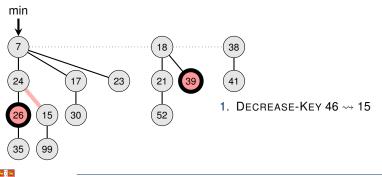


- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list

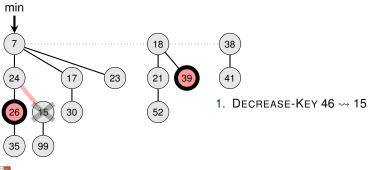




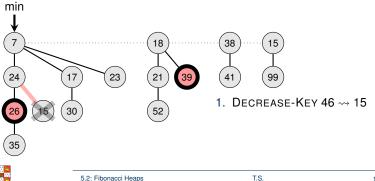
- Decrease the key of x (given by a pointer)
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- \Rightarrow Cut tree rooted at x, unmark x, meld into root list



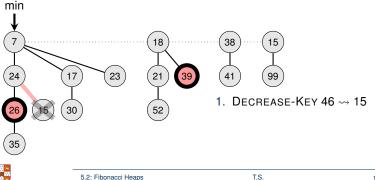
- Decrease the key of x (given by a pointer)
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- \Rightarrow Cut tree rooted at x, unmark x, meld into root list



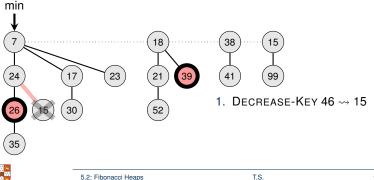
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:



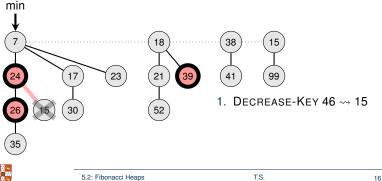
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked



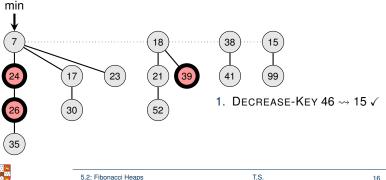
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)



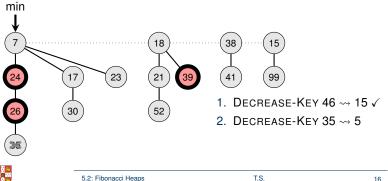
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)



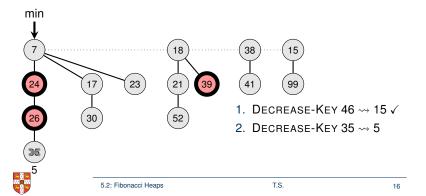
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)



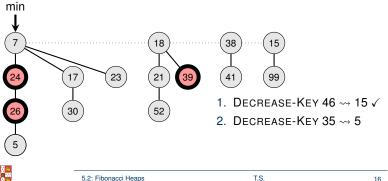
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)



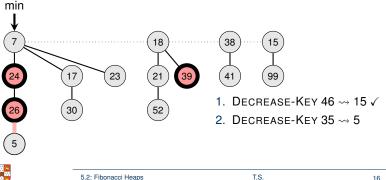
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)



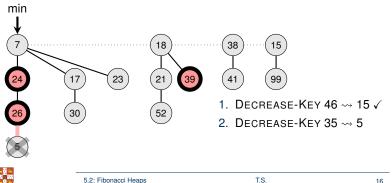
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)



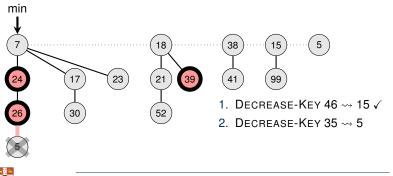
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)

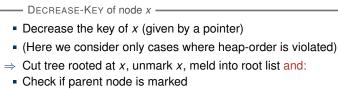


- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)

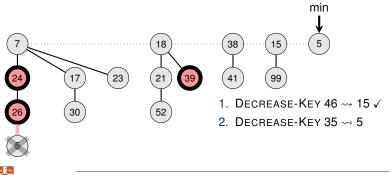


- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)

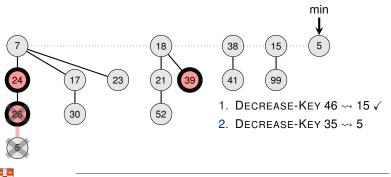




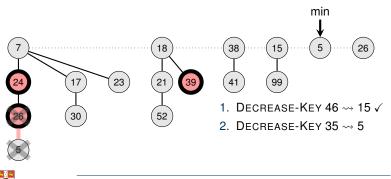
- If unmarked, mark it (unless it is a root)
- If marked,



- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)
 - If marked, unmark and meld it into root list and recurse (Cascading Cut)

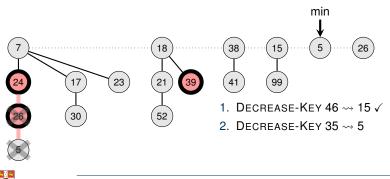


- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)
 - If marked, unmark and meld it into root list and recurse (Cascading Cut)

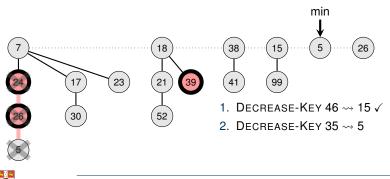


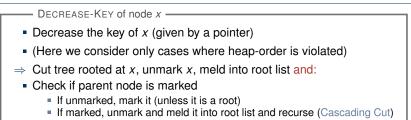
DECREASE-KEY of node x
Decrease the key of x (given by a pointer)

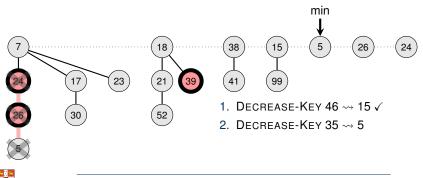
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)
 - If marked, unmark and meld it into root list and recurse (Cascading Cut)

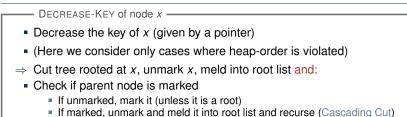


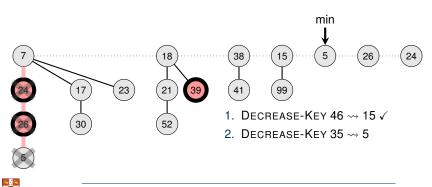
- Decrease the key of x (given by a pointer)
- (Here we consider only cases where heap-order is violated)
- \Rightarrow Cut tree rooted at x, unmark x, meld into root list and:
 - Check if parent node is marked
 - If unmarked, mark it (unless it is a root)
 - If marked, unmark and meld it into root list and recurse (Cascading Cut)

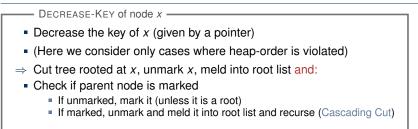


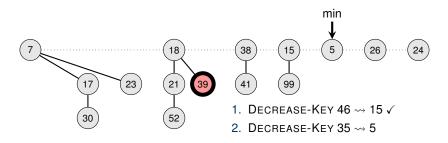




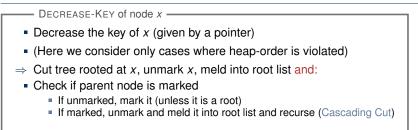


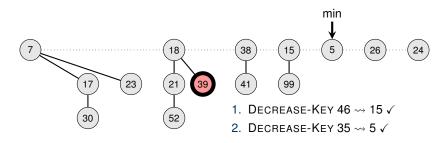




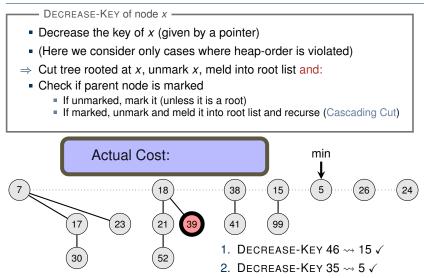




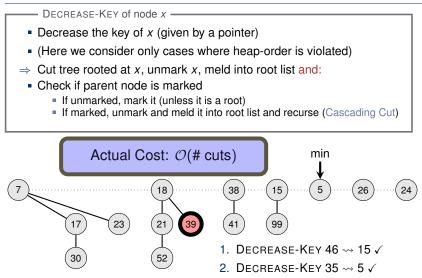




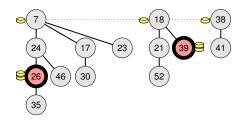












5.2 Fibonacci Heaps (Analysis)

Frank Stajano

Thomas Sauerwald

Lent 2016



Structure

Operations

Glimpse at the Analysis

Amortized Analysis



5.2: Fibonacci Heaps (Analysis)

- INSERT: actual $\mathcal{O}(1)$
- EXTRACT-MIN: actual O(trees(H) + d(n))
- DECREASE-KEY: actual O(# cuts) ≤ O(marks(H))



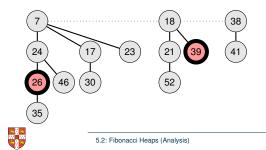
- INSERT: actual $\mathcal{O}(1)$
- EXTRACT-MIN: actual $\mathcal{O}(\text{trees}(H) + d(n))$
- DECREASE-KEY: actual O(# cuts) ≤ O(marks(H))

$$\Phi(H) = trees(H) + 2 \cdot marks(H)$$



- INSERT: actual $\mathcal{O}(1)$
- EXTRACT-MIN: actual $\mathcal{O}(\text{trees}(H) + d(n))$
- DECREASE-KEY: actual O(# cuts) ≤ O(marks(H))

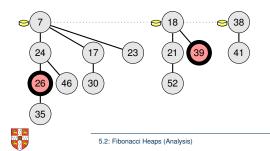
$$\Phi(H) = trees(H) + 2 \cdot marks(H)$$



T.S.

- INSERT: actual $\mathcal{O}(1)$
- EXTRACT-MIN: actual $\mathcal{O}(\text{trees}(H) + d(n))$
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$$\Phi(H) = trees(H) + 2 \cdot marks(H)$$

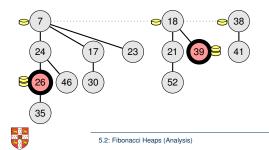


T.S.

з

- INSERT: actual $\mathcal{O}(1)$
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$$\Phi(H) = trees(H) + 2 \cdot marks(H)$$



T.S.

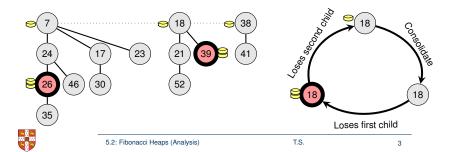
- INSERT: actual $\mathcal{O}(1)$
- EXTRACT-MIN: actual $\mathcal{O}(\text{trees}(H) + d(n))$
- DECREASE-KEY: actual $\mathcal{O}(\# \text{ cuts}) \leq \mathcal{O}(\text{marks}(H))$ amortized $\mathcal{O}(1)$

amortized $\mathcal{O}(1)$

Lifecycle of a node

amortized $\mathcal{O}(d(n))$

$$\Phi(H) = \text{trees}(H) + 2 \cdot \text{marks}(H)$$



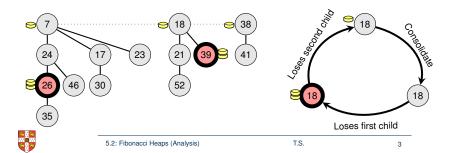
- INSERT: actual O(1)
- EXTRACT-MIN: actual $\mathcal{O}(\text{trees}(H) + d(n))$
- DECREASE-KEY: actual $\mathcal{O}(\# \text{ cuts}) \leq \mathcal{O}(\text{marks}(H))$ amortized $\mathcal{O}(1)$?

amortized $\mathcal{O}(1)$

Lifecycle of a node

amortized $\mathcal{O}(d(n))$?

$$\Phi(H) = \text{trees}(H) + 2 \cdot \text{marks}(H)$$



Structure

Operations

Glimpse at the Analysis

Amortized Analysis



Actual Cost —

• DECREASE-KEY: $\mathcal{O}(x+1)$, where *x* is the number of cuts.



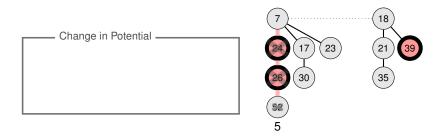
Actual Cost —

• DECREASE-KEY: $\mathcal{O}(x + 1)$, where *x* is the number of cuts.



Actual Cost —

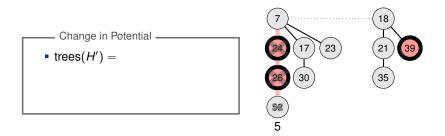
• DECREASE-KEY: $\mathcal{O}(x+1)$, where *x* is the number of cuts.





Actual Cost —

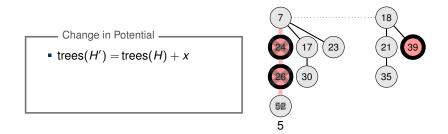
• DECREASE-KEY: $\mathcal{O}(x+1)$, where *x* is the number of cuts.





Actual Cost —

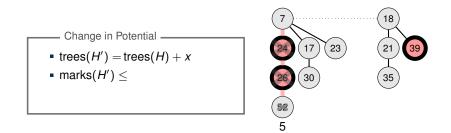
• DECREASE-KEY: $\mathcal{O}(x+1)$, where *x* is the number of cuts.





Actual Cost —

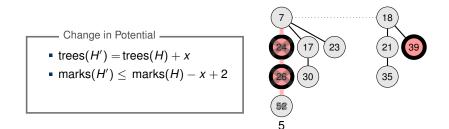
• DECREASE-KEY: O(x + 1), where x is the number of cuts.





Actual Cost –

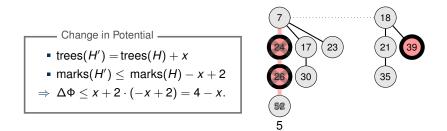
• DECREASE-KEY: $\mathcal{O}(x+1)$, where *x* is the number of cuts.





- Actual Cost -

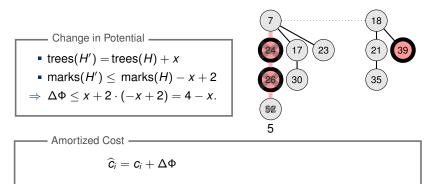
• DECREASE-KEY: $\mathcal{O}(x+1)$, where *x* is the number of cuts.





Actual Cost -

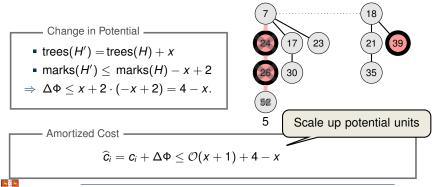
• DECREASE-KEY: O(x + 1), where x is the number of cuts.





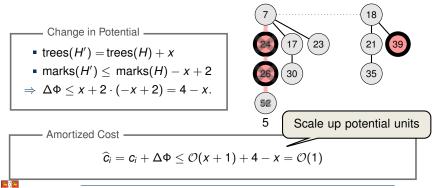
Actual Cost -

• DECREASE-KEY: $\mathcal{O}(x + 1)$, where *x* is the number of cuts.



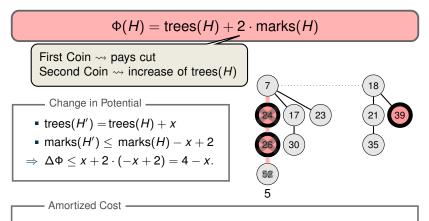
Actual Cost -

• DECREASE-KEY: $\mathcal{O}(x + 1)$, where *x* is the number of cuts.



Actual Cost -

• DECREASE-KEY: O(x + 1), where x is the number of cuts.



$$\widehat{c}_i = c_i + \Delta \Phi \leq \mathcal{O}(x+1) + 4 - x = \mathcal{O}(1)$$

