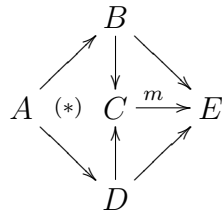


Category Theory for Computer Science Mid-Term Problems

November 2009

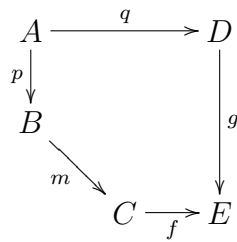
Problem 1.

(a) Consider the commutative diagram



where m is a mono. Prove that if $(*)$ is a pullback then so is the outer square.

(b) Consider the commutative diagram



where m is a mono.

Prove that if $\begin{array}{ccc} A & \xrightarrow{q} & D \\ m \circ p \downarrow & & \downarrow g \\ C & \xrightarrow{f} & E \end{array}$ is a pullback, then so is $\begin{array}{ccc} A & \xrightarrow{q} & D \\ p \downarrow & & \downarrow g \\ B & \xrightarrow{f \circ m} & E \end{array}$

Problem 2.

- A mono $m : X \rightarrow Y$ is *regular* if there exists a diagram $Y \begin{array}{c} \xrightarrow{f} \\ \rightrightarrows \\ \xrightarrow{g} \end{array} Z$ for which m is an equalizer.
- A mono $m : X \rightarrow Y$ is *strong*, if for every commutative diagram as on the left below where $e : U \rightarrow V$ is epi



there exists exactly one arrow $d : V \rightarrow X$ as on the right above such that both triangles commute.

- A mono $m : X \rightarrow Y$ is *extremal* if for every commutative diagram

$$\begin{array}{ccc} & & V \\ & e \nearrow & \downarrow v \\ X & \xrightarrow{m} & Y \end{array}$$

where $e : X \rightarrow V$ is epi, e is an isomorphism.

Prove that the following implications between properties of monomorphisms hold in any category:

$$\text{section} \implies \text{regular} \implies \text{strong} \implies \text{extremal} .$$

(None of these is in general an equivalence, but that is another story.)

Problem 3. Given an object A in a small category \mathbf{C} , consider the following diagram in the category \mathbf{Cat} of small categories and functors:

$$\begin{array}{ccc} \mathbf{C}^\rightarrow & & \mathbf{1} \\ & \searrow \text{Cod} & \swarrow K_A \\ & \mathbf{C} & \end{array}$$

where:

- \mathbf{C}^\rightarrow is the arrow category of \mathbf{C} and Cod is the *codomain functor*, defined by:
 - $\text{Cod}(f) = \text{cod}(f)$ on objects in \mathbf{C}^\rightarrow ,
 - $\text{Cod}(h, k) = k$ on arrows in \mathbf{C}^\rightarrow .
- $\mathbf{1}$ is the category with one object $*$ and one arrow 1_* , and K_A is the constant functor defined by:
 - $K_A(*) = A$,
 - $K_A(1_*) = 1_A$.

Prove that the slice category \mathbf{C}/A , together with suitable functors to \mathbf{C}^\rightarrow and $\mathbf{1}$, is a pullback of the above diagram.

Problem 4. Fix a set X with $* \notin X$.

For $n \in \mathbb{N}$, define

$$X_n \stackrel{\text{def}}{=} \{(x_1, \dots, x_n) \in (X \cup \{*\})^n \mid x_i = * \implies x_{i+1} = * \text{ for all } i = 1, \dots, n-1\} .$$

(a) Consider the following infinite diagram in the category **Sets** of sets and functions:

$$X_1 \xrightarrow{e_1} X_2 \xrightarrow{e_2} \cdots \longrightarrow X_n \xrightarrow{e_n} X_{n+1} \longrightarrow \cdots \quad (n \in \mathbb{N}) \quad (1)$$

where $e_n : X_n \rightarrow X_{n+1}$ is the function given by

$$e_n(x_1, \dots, x_n) = (x_1, \dots, x_n, *) .$$

Give a simple description of a colimit of (1), and prove its universal property.

(b) Consider the following infinite diagram in the category **Sets** of sets and functions:

$$X_1 \xleftarrow{p_1} X_2 \xleftarrow{p_2} \cdots \longleftarrow X_n \xleftarrow{p_n} X_{n+1} \longleftarrow \cdots \quad (n \in \mathbb{N}) \quad (2)$$

where $p_n : X_{n+1} \rightarrow X_n$ is the function given by

$$p_n(x_1, \dots, x_n, x_{n+1}) = (x_1, \dots, x_n) .$$

Give a simple description of a limit of (2), and prove its universal property.