The Fountain of Knowledge

Using Digital Fountains to provide reliable storage in the Data Centre Toby Moncaster, Andrew W. Moore, Jon Crowcroft

Data Centre Storage Requirements

Speed: Needs to keep up with application speed. Often read speed is more important than write. This leads to "move the application to the data" approaches.

Reliability: Data should not be corrupted, also need to ensure synchronisation between copies of the same data.

Resilience: Storage should be resilient against local and regional failures. Multiple copies of data may be required. Failures may come singly or in sequence.

Existing Approaches

Storage Area Networks (e.g. NetApp)

Provides dedicated specialist network to provide fast read/write access to centralised disk arrays. Storage appears as local disk.



Data Centre Knowledge Fountains

Data is split into chunks and encoded into a digital fountain. Encoded chunks are stored at locations across data centre. Resilience is achieved by creating multiple encodings of each tract. These are stored at separate physical locations.

To access the data simply download sufficient encoded chunks to decode the original data. Overhead is 5-10% depending on the encoding used.

Choosing Where to Store the Data

Use a hash-based indexing approach similar to Flat Datacenter Storage. Each storage location (or blob) has a tractserver. These work just as in FDS. Central metadata server stores simple index.

Central metadata (e.g. GFS)



Storage divided into chunks. A
single central metadata server
controls how data are
distributed and replicated
between these. Ideal for
oversubscribed networks.

Distributed (e.g. Flat Datacenter Storage)

Storage split into blobs. Data split into tracts. Each blob has a tract server which allocates space to each of the tracts.Central metadata server lists where the tract servers are.Uses hash function to split data among tracts. Only works in full bisection bandwidth networks.

Digital Fountains

Data is encoded using sparse erasure codes such as Tornado, Luby Transform or Raptor codes. Receiver needs to obtain slightly more encoded data than the original raw Index key is generated using simple hash of file descriptor and blob ID. Index returns list of tract servers where code blocks for that file exist.

Detailed Operation

Writing Data: Split the data into blocks. Decide on the degree of redundancy. Encode multiple versions of the data using a suitable encoding scheme. Ask the metadata server to choose which tract servers to store the code chunks.

Accessing data: Use the metadata server to identify all the tract servers with code chunks for the file. Select tract servers based on suitable heuristics (location, speed, etc). Download sufficient chunks to recover data.

Recovering from failures: If a storage blob fails the metadata server will choose alternative tract servers for the data on that blob. Because of redundant encoding many blobs need to fail before data is unrecoverable.

source. Can receive code chunks from multiple sources.



Receive C₁, C₂, C₇, C₃, C₄, C₅ Use C₁ to recover D₁ Use C₂ and D₁ to recover D₂ Then wait till you receive C₃ Use C₃ and D₁ to recover D₃ Use C₄ and D₂ to recover D₄ Use C₅ and D₃ to recover D₅ Use C₇ and D₃ to recover (D₅ + D₆) Use (D₅ + D₆) and D₅ to recover D₆

Data is divided into segments. Code chunks are created by XORing sets of data segments. To decode, repeatedly XOR the coded blocks until original data is recovered.

Is it Suitable for Data Centres?

Speed: System is fast since code chunks are downloaded in parallel. Digital Fountain encoding is efficient and fast.

Reliability: Data is encoded across many blocks. If a block is corrupted or lost there is plenty of redundancy to recover from this.

Resilience: The system has a high degree of resilience because the encoding only requires sufficient code blocks to be downloaded which can come from any storage blob.





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