# GraphChi: Large-Scale Graph Computation on Just a PC

Kyrola Et al.

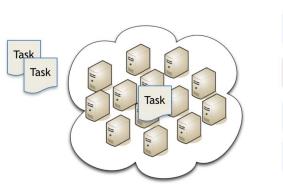
**James Trever** 

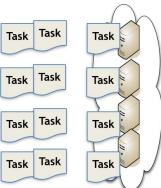
#### Could we compute Big Graphs on a single machine? Disk Based Computation



## Why would you want to?

- Distributed State is hard to program
  - Cluster crashes can occur
  - Cumbersome
- Efficient Scaling
  - Parallelise each task vs Parallelise across tasks
- Cost
  - Easier management and simpler hardware
- Energy Consumption
  - Full utilisation of a single computer
- Easier Debugging





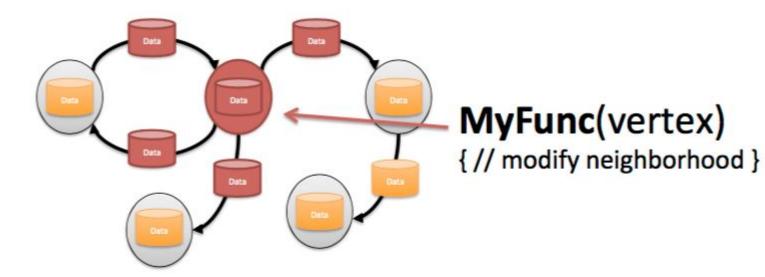


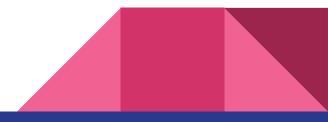
## Contents

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- Parallel Sliding Windows
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## **Computational Model**

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#### **Storage Model**

- Compressed Sparse Row (CSR) allows for fast loading of out-edges
- Compressed Sparse Column (CSC) allows for fast loading of in-edges



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- Compressed Sparse Column (CSC) allows for fast loading of in-edges

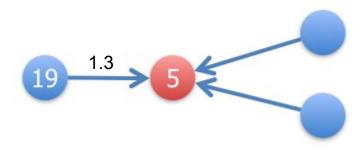
Why not both?



Challenges

#### **Random Access Problem**

- Symmetrised adjacency file with values

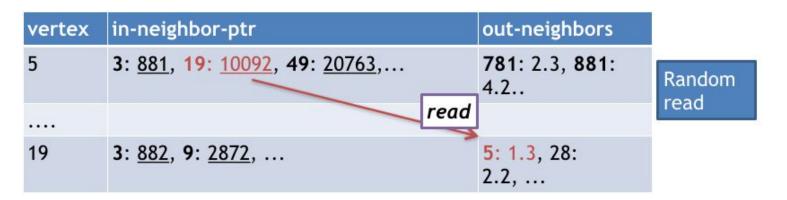


vertex	in-neighbors	out-neighbors	
5	<b>3</b> :2.3, <b>19</b> : <b>1</b> .3, <b>49</b> : 0.65,	<b>781</b> : 2.3, <b>881</b> : 4.2	Random
	sync	hronise	write
19	<b>3</b> : 1.4, <b>9</b> : 12.1,	<b>5:</b> 1.3, 28: 2.2,	



#### **Random Access Problem**

- File Index Pointers





1.3

19

#### **Possible Solutions**

- 1. Use SSD as memory extension
  - Too many small objects, need millions of reads and writes a second
- 2. Compress the graph structure to fit in RAM
  - Associated values do not compress well
- 3. Cachine the hot vertices
  - Unpredictable Performance



## Parallel Sliding Windows (PSW)

#### **PSW: Phases**

PSW processes the graph one sub-graph at a time

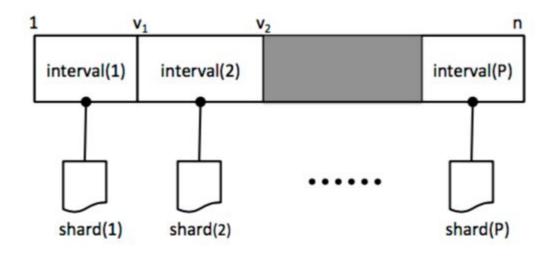
- 1. Load
- 2. Compute
- 3. Write

In one iteration the whole graph is processed

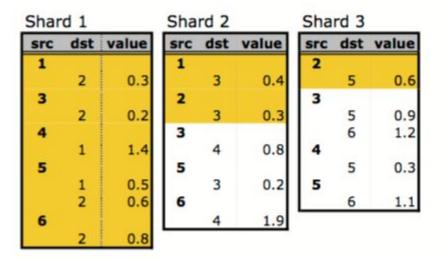


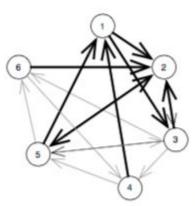
#### **PSW: Intervals and Shards - Load**

- Subgraph = Interval



#### **PSW: Example - Load**

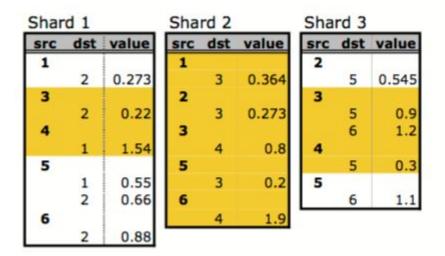


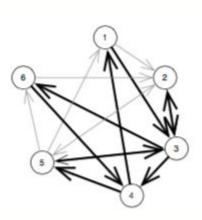


(a) Execution interval (vertices 1-2)

(b) Execution interval(vertices 1-2)

#### **PSW: Example - Load**

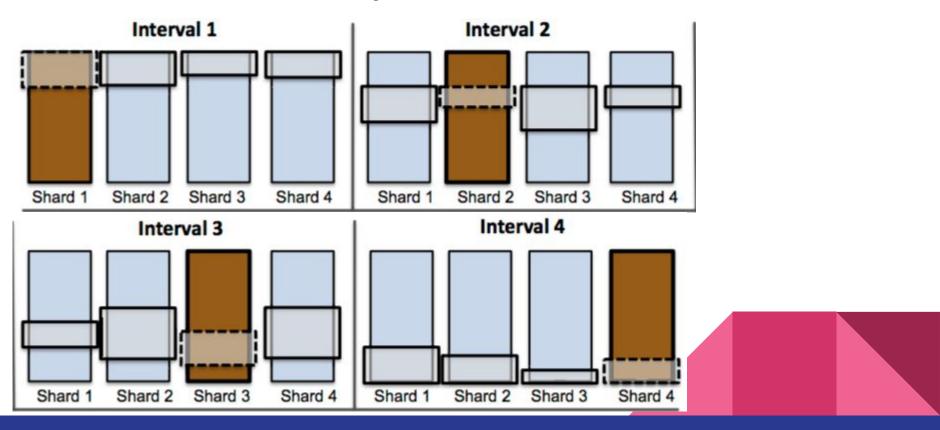




(c) Execution interval (vertices 3-4)

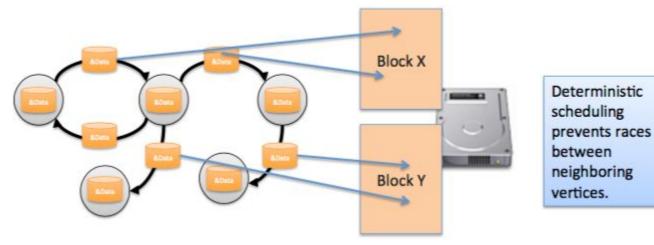
(d) Execution interval (vertices 3-4)

#### **PSW: General Example - Load**



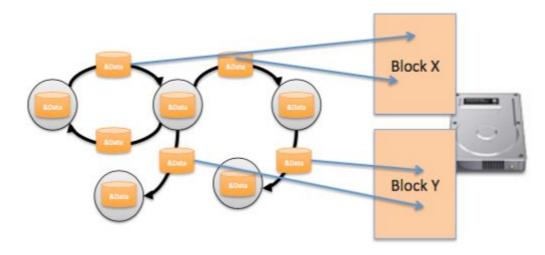
#### **PSW: Compute Phase**

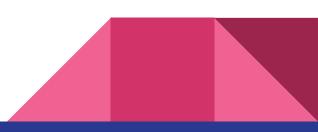
- UpdateFunction executes on intervals vertices in parallel
- Edges have pointers to the loaded data blocks



#### **PSW: Write Phase**

- Blocks are written back to disk asynchronously





## Implementation and Experiments

### **Preprocessing Step**

- Sharder program included with GraphChi

- 1. Counts the in-degree of each vertex and computes the prefix sum over the degree array so that each interval contains same number of in edges
- 2. Sharder writes each edge to temporary scratch file belonging to the shard
- 3. Sharder Processes each scratch file
- 4. Sharder computes binary degree file containing in and out degree for each vertex (used to calculate memory requirements)

#### **Preprocessing Experiment**

Graph name	Vertices	Edges	P	Preproc.
live-journal [3]	4.8M	69M	3	0.5 min
netflix [6]	0.5M	99M	20	1 min
domain [44]	26M	0.37B	20	2 min
twitter-2010 [26]	42M	1.5B	20	10 min
uk-2007-05 [11]	106M	3.7B	40	31 min
uk-union [11]	133M	5.4B	50	33 min
yahoo-web [44]	1.4B	6.6B	50	37 min

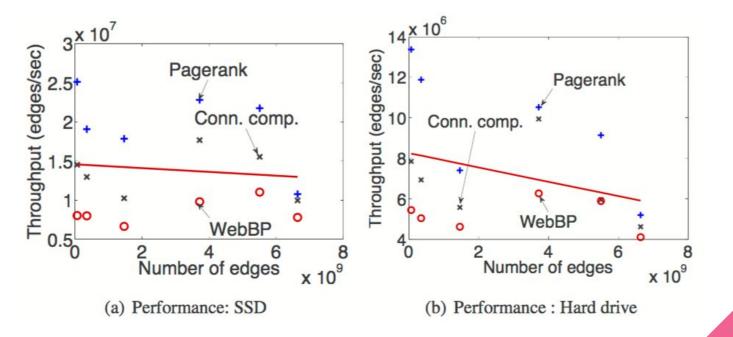
### **Comparison Experiment**

Application & Graph	Iter.	Comparative result	GraphChi (Mac Mini)	Ref
Pagerank & domain	3	GraphLab[30] on AMD server (8 CPUs) 87 s	132 s	-
Pagerank & twitter-2010	5	Spark [45] with 50 nodes (100 CPUs): 486.6 s	790 s	[38]
Pagerank & V=105M, E=3.7B	100	Stanford GPS, 30 EC2 nodes (60 virt. cores), 144 min	approx. 581 min	[37]
Pagerank & V=1.0B, E=18.5B	1	Piccolo, 100 EC2 instances (200 cores) 70 s	approx. 26 min	[36]
Webgraph-BP & yahoo-web	1	Pegasus (Hadoop) on 100 machines: 22 min	27 min	[22]
ALS & netflix-mm, D=20	10	GraphLab on AMD server: 4.7 min	9.8 min (in-mem)	
			40 min (edge-repl.)	[30]
Triangle-count & twitter-2010	-	Hadoop, 1636 nodes: 423 min	60 min	[39]
Pagerank & twitter-2010	1	PowerGraph, 64 x 8 cores: 3.6 s	158 s	[20]
Triange-count & twitter- 2010	-	PowerGraph, 64 x 8 cores: 1.5 min	60 min	[20]

Mac Mini Dual Core 2.5 GHz, 8GB Ram AMD Server 8 core server with 4 dual core CPU's



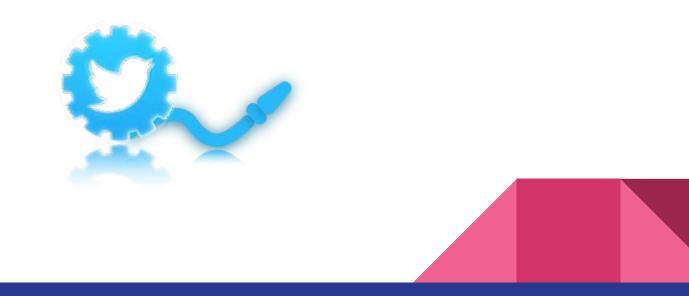
#### **Throughput Experiment**



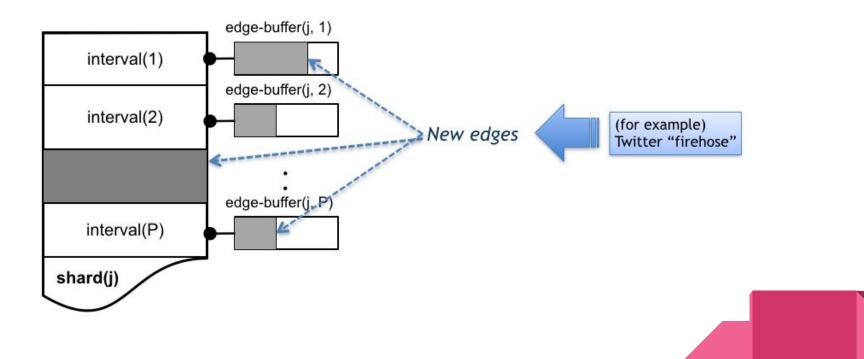
**Evolving Graphs** 

## **Evolving Graphs**

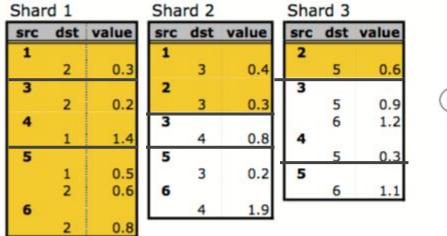
- Add and remove edges in streaming fashion whilst continuing computation
- Most interesting networks grow continuously

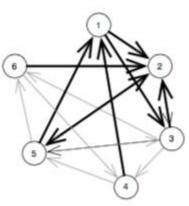


### **PSW and Evolving Graphs**



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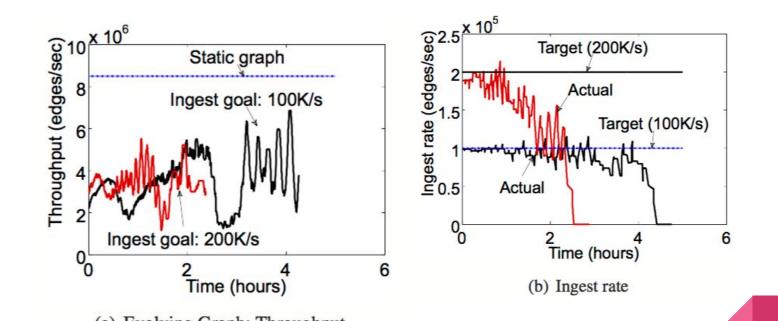




(a) Execution interval (vertices 1-2)

(b) Execution interval(vertices 1-2)

#### **Evolving Graphs - Experiment**



(a) Evolving Graph: Throughput

## Graphs Used

Graph name	Vertices	Edges	P	Preproc.
live-journal [3]	4.8M	69M	3	0.5 min
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## **Critical Evaluation**

- Few mistakes in the paper referencing incorrect tables or quoting wrong figures
- Cannot efficiently support dynamic ordering like priority ordering or efficiently support graph traversals or vertex queries
- Evolving graph experiments not very clear
- No monetary analysis

## Bibliography

A. Kyrola, G. Blelloch, and C. Guestrin, "Graphchi: Large-scale graph computation on just a pc," in Proceedings of the 10th USENIX Conference on Operating Systems Design and Implementation, OSDI'12, (Berkeley, CA, USA), pp. 31–46, USENIX Association, 2012.

And his original presentation found here:

https://www.usenix.org/sites/default/files/conference/protectedfiles/kyrola\_osdi12\_slides.pdf

